

FunCPU

*7 Bit Homebrew CPU
Designed For
Functional Programming

András Juhász

• „Árokparty”, 19.07.2014.

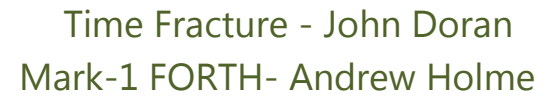


Motivation

- Software guy: to create something „real“
- IT engineer: to have a computer, which I fully understand
- Bored hobbyist: to find a challenge, a logical puzzle
- Myself: All of the above, plus: to create something „new“, unconventional, „exotic“.



A large pile of colorful, 3D-printed letters and numbers in various colors (green, yellow, blue, red, black) scattered on a white surface.



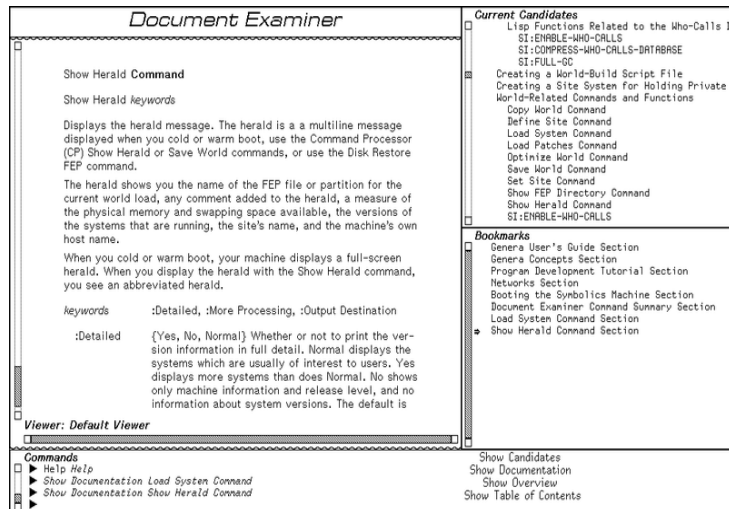
Goals

- Simple;
 - Unconventional;
 - Buildable/achievable;
 - „Useful“;
 - Native functional programming support.
-
- www.mycpu.blog.hu



„Lisp-machine“

- Graphical workstation
- General lisp op.system
- Symbolics, TI, Xerox



Dedicated CPUs

- VLSI-PLM, then PLUM for Prolog
- Nikolaus Wirth - CPU for Module 2
- INMOS Transputer for Occam
- AT&T Hobbit – to support CRISP C
- Java processors
- Ericsson ECOMP for Erlang
- Source: Wikipedia



Overview

- Concept
- Challenge
- Examples
- Architecture
- Implementation
- Computability
- Improvements



The Concept



FunCPU

- Natively supports functional programming;
- Special assembly instruction set;
- Simple programming paradigm;
- Typed (tagged) architecture;
- Main focus on numerical computations;
- Turing-complete.



Unconventional

- No PC.
- No SP, no stack.
- There are no flags.
- No jumping/branching instructions.
- Not even the accumulator exists.
- No I/O operations.
- So, what is what we have?



"Small Cooker"



"Small Cooker"



Architecture

- 7 bit literals;
- 3 built-in,
- 32 user-definable functions;
- ROM – to store functions (256 bytes);
- RAM – to evaluate expressions (256 bytes);
- 8 bit data bus;
- 8/9 bit address bus.



Functions

Built-in:

- **inc**(x) : $x+1=\text{inc}(x)$
- **dec**(x) : $x-1=\text{dec}(x)$
- **If** cond **then** exp1 **else** exp2

Plus user-definable functions.

Note: „0“ represents True, any other value is False.



Computational Model

- Library: function definitions

$$f(x,y) \quad := x+y$$

$$g(x,y,z) \quad := f(x,y)-f(y,z)$$

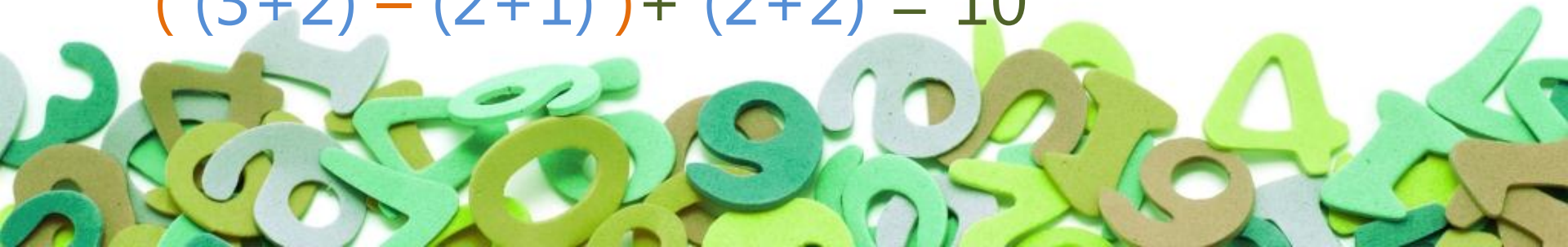
- Program: closed expression (const, func.)

$$g(3,2,1)+f(2,2)$$

- Execution: rewriting rules

$$(f(3,2) - f(2,1)) + (2+2) =$$

$$((3+2) - (2+1)) + (2+2) = 10$$



The Challenge



Challenges

- How to represent expressions;
- How to represent functions;
- Parenthesis, operation priorities;
- Argument passing;
- Evaluation strategy;
- How to model with integrated circuits;
- Physical implementation.



Encoding Scheme

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Hex	Value
0	0	0	0	0	0	0	0	00	zero / true
0	c	c	c	c	c	c	c	00-	constant
	1	1	1	1	0	1	1	7B	last const.
0	1	1	1	1	1	a	a	7C	argument
								7F	1..4
1	F	f	f	f	f	a	a	80-FE	function
1	1	1	1	1	1	0	0	FC	dec
1	1	1	1	1	1	0	1	FD	If
1	1	1	1	1	1	1	0	FE	inc
1	1	1	1	1	1	1	1	FF	EOX

Function arity: 1..4

map 8 bit value to 3 bit class

Note: Constant functions must have at least one argument.



Function Encoding

„%1fff ffaa“

Function address in binary: %ffff f000.
(\$00, \$08, \$10, ... , \$E8, \$F0)

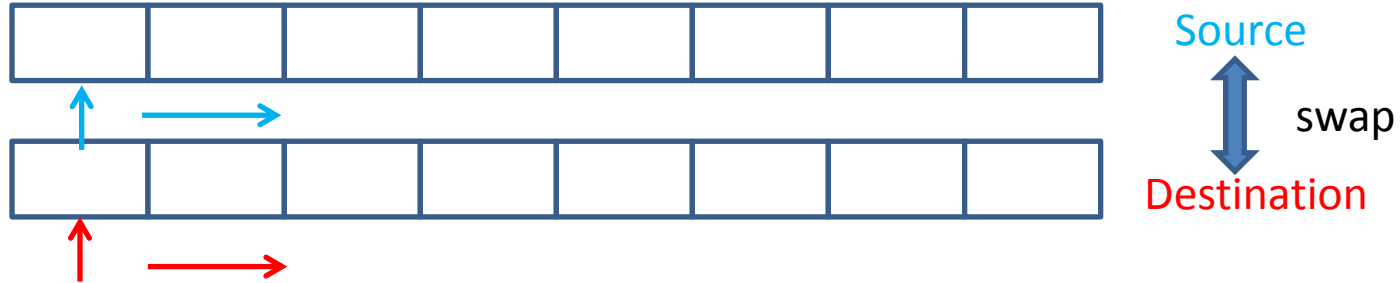
Number of arguments: aa=00, 01, 10, 11
denote 4, 3, 2, 1 arguments respectively.

- Arity is encoded in function „id“ -> efficient
- No real function „id“, id gives instantly the function physical address.

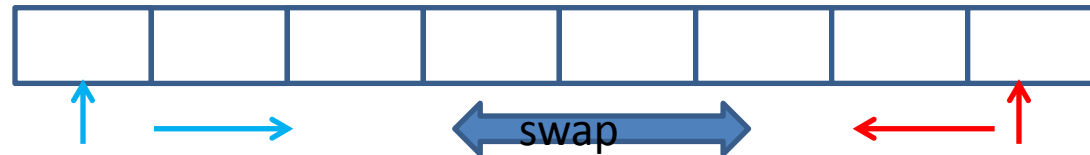


Memory Models

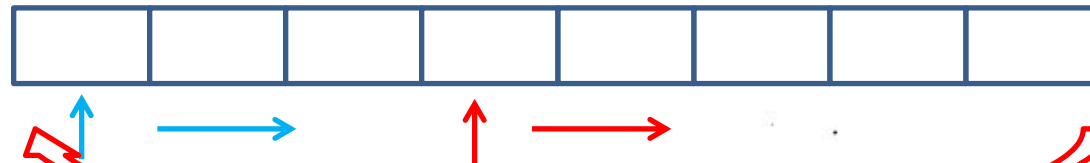
1.) Separate Source and Destination Memory



2.) Shared Memory, Source increasing/Destination decreasing



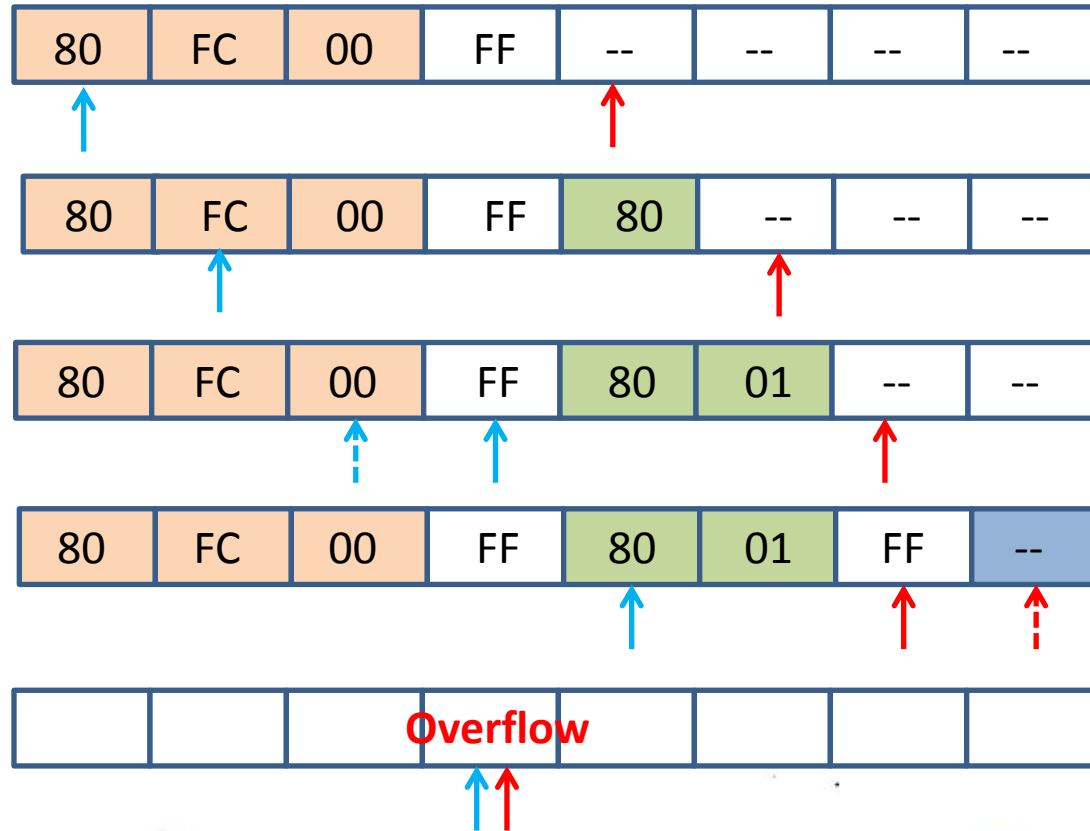
3.) Shared Memory, Source is „following” Destination



Optimized storage

„Step over”

Index Chase



Parenthesis

- $2 * \text{fac}(\text{fac}(2) + \text{fac}(1))?$
- Everything is right-associative (PPN):
- $* \ 2 \ \text{fac} \ + \ \text{fac} \ 2 \ \text{fac} \ 1$
- Function Scope? (simple, hw-based!):
- Store arity in AC in the beginning.
- $\text{AC} := \text{AC} - 1 + \text{arity of the next symbol}.$
- End of Scope, if $\text{AC} = 0.$



Example: Parenthesis

- Scope of the first „fac“:
- if 2 **fac** + fac 2 fac 1 if ... $ac=1$
- if 2 **fac** + fac 2 fac 1 if ... $ac=2=1-1+2$
- if 2 **fac** + **fac** 2 fac 1 if ... $ac=2=2-1+1$
- if 2 **fac** + **fac** 2 fac 1 if ... $ac=1=2-1$
- if 2 **fac** + **fac** 2 **fac** 1 if ... $ac=1=2-1+1$
- if 2 **fac** + fac 2 fac 1 if ... $ac=\mathbf{0}=1-1$



Evaluation Issue

- $\text{fac}(n) := \text{if } n=0 \text{ then } 1 \text{ else } n * \text{fac}(n-1)$
- $\text{fac}(1) = ?$
- $\text{if } 1=0 \text{ then } 0 \text{ else } 1 * \text{fac}(0)$
- $\text{If } 1=0 \text{ then } 0 \text{ else } 1 * (\text{if } 0=0 \text{ then } 1 \text{ else } 0 * \text{fac}(-1))$
- $\text{If } 1=0 \text{ then } 0 \text{ else } 1 * (\text{if } 0=0 \text{ then } 1 \text{ else } 0 * (\text{if } -1=0 \text{ then } 0 \text{ else } -1 * \text{fac}(-1))) \dots\dots$



Evaluation Goals

- Must terminate (applicative results in infinite loop)
- Must be efficient.
- Must be simple (suitable to be represented in hardware directly).



Passing Arguments

- $\text{fac}(n) := n * \text{fac}(n-1)$
- $\text{fac}(\text{add}(5,4))?$
- $\text{ack}(\text{add}(\text{fac}(\text{fac}(2)) + \text{fac}(1)), \text{fac}(1)), 1)?$
- No stack;
- No dedicated memory for parameters;
- Still „arbitrary“ depth of function calls.



Evaluation - If

- `if cond exp1 exp2` \Rightarrow
- `exp1`, if `cond=00`
- `exp2`, if `cond<>00`, but it is a literal
- Otherwise „if“ is not evaluated.



Evaluating – inc, dec...

- `inc exp` \Rightarrow
- `exp+1`, if „`exp`“ is constant.
- `dec` is reduced analogously.
- Other functions are evaluated/"called", if all of their arguments are constants.
- Corollary: all parameters passed to functions are constants.



Evaluation Strategy

- Many strategies co-exist.
- Scan input expression symbol by symbol.
- 1.) Copy symbol from source to output.
- 2.) Evaluate functions (if, inc, dec, and user-defined) if possible.
- Next cycle: output becomes the input.
- Stop, if the first symbol is constant.



Termination

- inc, dec terminates immediately, thus reduces the length of expression.
- if cond exp1 exp2 – if cond is constant, then the length of expression is decreased.
- Eventually everything is based on/can be reduced to the three built-in functions.
- Sooner or later, all „if“, „inc“, „dec“ are evaluated/reduced (if possible).



Examples



Elementary Functions

- $I(x) := x$

7F FF

- $C(x) := c$

„c“ FF, where $c = 00..7B$

constant 7C can be encoded as: `inc(7B)`

FC 7B FF



Few Predicates

- $\text{not}(x) := \text{if } x \text{ then } 1 \text{ else } 0$

FE 7F 01 00 FF

- $\text{dec}(x,y) := \text{if } y=0 \text{ then } x \text{ else } \text{dec}(\text{dec}(x), \text{dec}(y))$

FE 7E 7F 81 F8 7F F8 7E FF

- $>(x,y) :=$ if $x=0$ then not(y)
else $>(\text{dec}(x), \text{dec}(y))$



Some Functions

- $\text{add}(x,y) :=$ if $y=0$ then x
 else $\text{inc}(\text{add}(x,\text{dec}(y)))$
- $\text{mul}(x,y) :=$ if $y=0$ then 0
 else $\text{add}(x,\text{mul}(x,\text{dec}(y)))$
- $\text{fac}(n) :=$ if $n=0$ then 1
 else $\text{mul}(n,\text{fac}(\text{dec}(n)))$



Case...

case x

when c0 then b0

when c1 then b1

...

when cn then bn

otherwise b :=

if=(x,c0) then b0

else if =(x,c1) then b1

else ...



Ackermann function

ack(m,n):=

if m=0 then inc(n)

elseif n=0 then ack(dec(m),1)

else ack(dec(m),ack(m,dec(n)))

FE 7E FC 7F FE 7F 81 F8 7E 01 81 F8 7E 81 7E
F8 7F FF



Functions Encoded

add:= 00/81: FE 7E 7F FC 81 7F F8 7E FF

mul:= 10/89: FE 7E 00 81 7F 89 7F F8 7E FF

fac:= 20/90: FE 7F 01 89 90 F8 7F 7F FF

FE

if

FC

inc

F8

dec

00, 01

constant

FF

end of exp.

7E, 7F

arguments

y, x



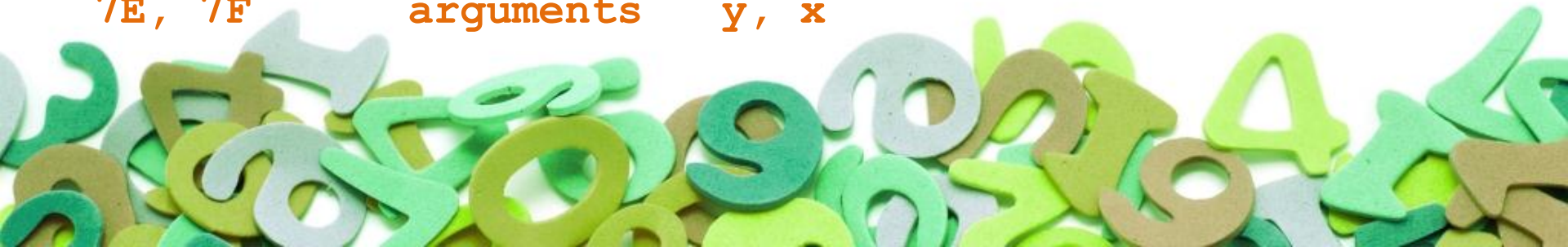
1 + 1 = 2

add(x,y) := FE 7E 7F FC 81 7F F8 7E FF

[illegible]

↳ $1+1=$ „02“

FE	if	81	Add
FC	inc	F8	dec
00, 01, 02	constant	FF	end of exp.
7E, 7F	arguments	y, x	



fac(5) = 120

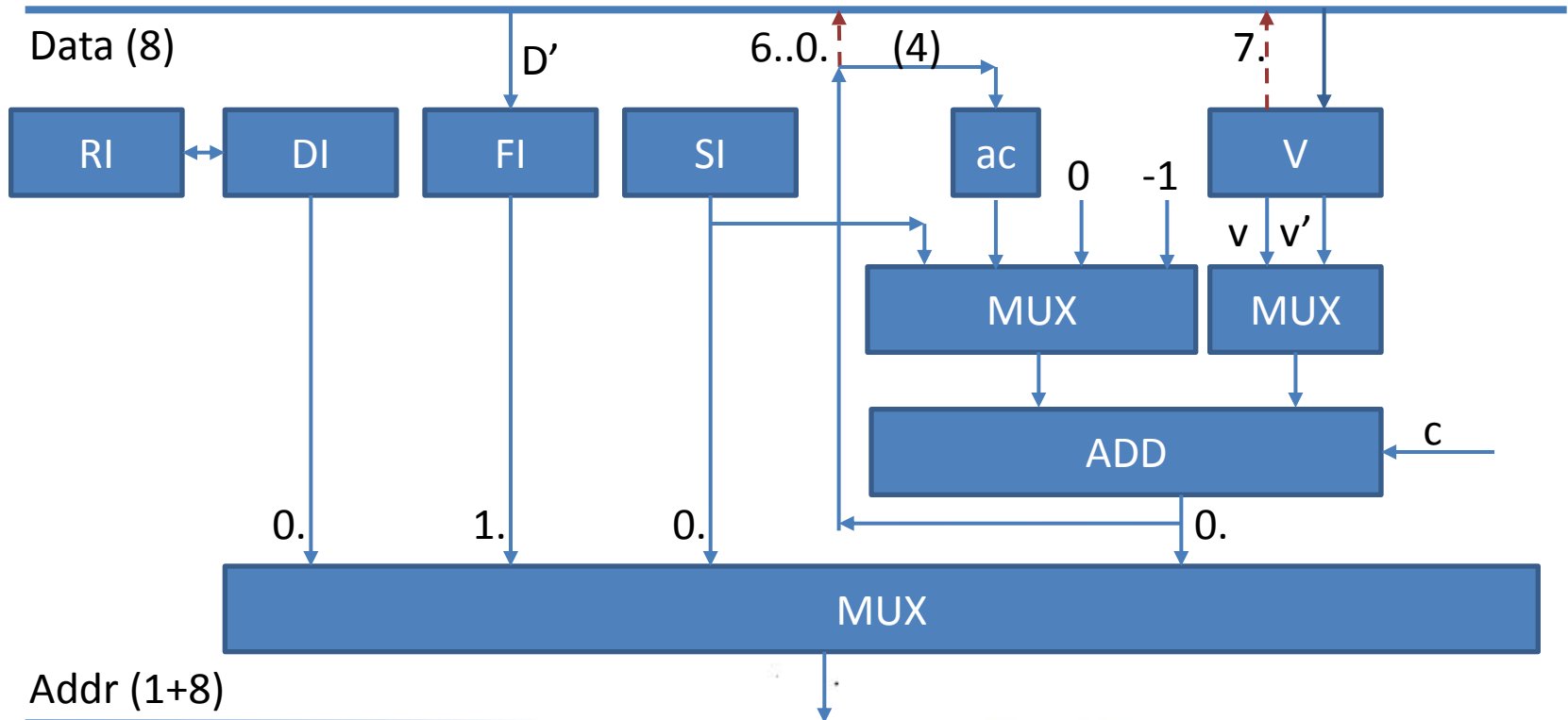
--:	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F
00:	FF	FC	FC	FC	75	FF	FC	FC	76	FF	FC	77	FF	78	FF	FC
10:	FC	FE	03	71	FC	81	71	F8	03	FF	FC	FC	FC	FC	FC	81
20:	71	02	FF	FC	FC	FC	FC	FC	FE	71	02	FC	81	02	F8	71
30:	FF	FC	FC	FC	FC	FC	FC	81	02	70	FF	FC	FC	FC	FC	FC
40:	FC	FE	02	70	FC	81	70	F8	02	FF	FC	FC	FC	FC	FC	FC
50:	FC	81	70	01	FF	FC	FC	FC	FC	FC	FC	FC	FE	70	01	FC
60:	81	01	F8	70	FF	FC	FC	FC	FC	FC	FC	FC	FC	81	01	6F
70:	FF	FC	FC	FC	FC	FC	FC	FC	FC	FE	01	6F	FC	81	6F	F8
80:	01	FF	FC	FC	FC	FC	FC	FC	FC	FC	FC	81	6F	00	FF	FC
90:	FC	FC	FC	FC	FC	FC	FC	FC	FE	6F	00	FC	81	00	F8	6F
A0:	FF	FC	FC	FC	FC	FC	FC	FC	FC	FC	FC	81	00	6E	FF	FC
B0:	FC	FC	FC	FC	FC	FC	FC	FC	FC	FE	00	6E	FC	81	6E	F8
C0:	00	FF	FC	FC	FC	FC	FC	FC	FC	FC	FC	FC	6E	FF	FC	FC
D0:	FC	FC	FC	FC	FC	FC	FC	6F	FF	FC	FC	FC	FC	FC	FC	FC
E0:	FC	70	FF	FC	FC	FC	FC	FC	FC	FC	71	FF	FC	FC	FC	FC
F0:	FC	FC	72	FF	FC	FC	FC	FC	FC	73	FF	FC	FC	FC	FC	74



Architecture



Architecture (registers)

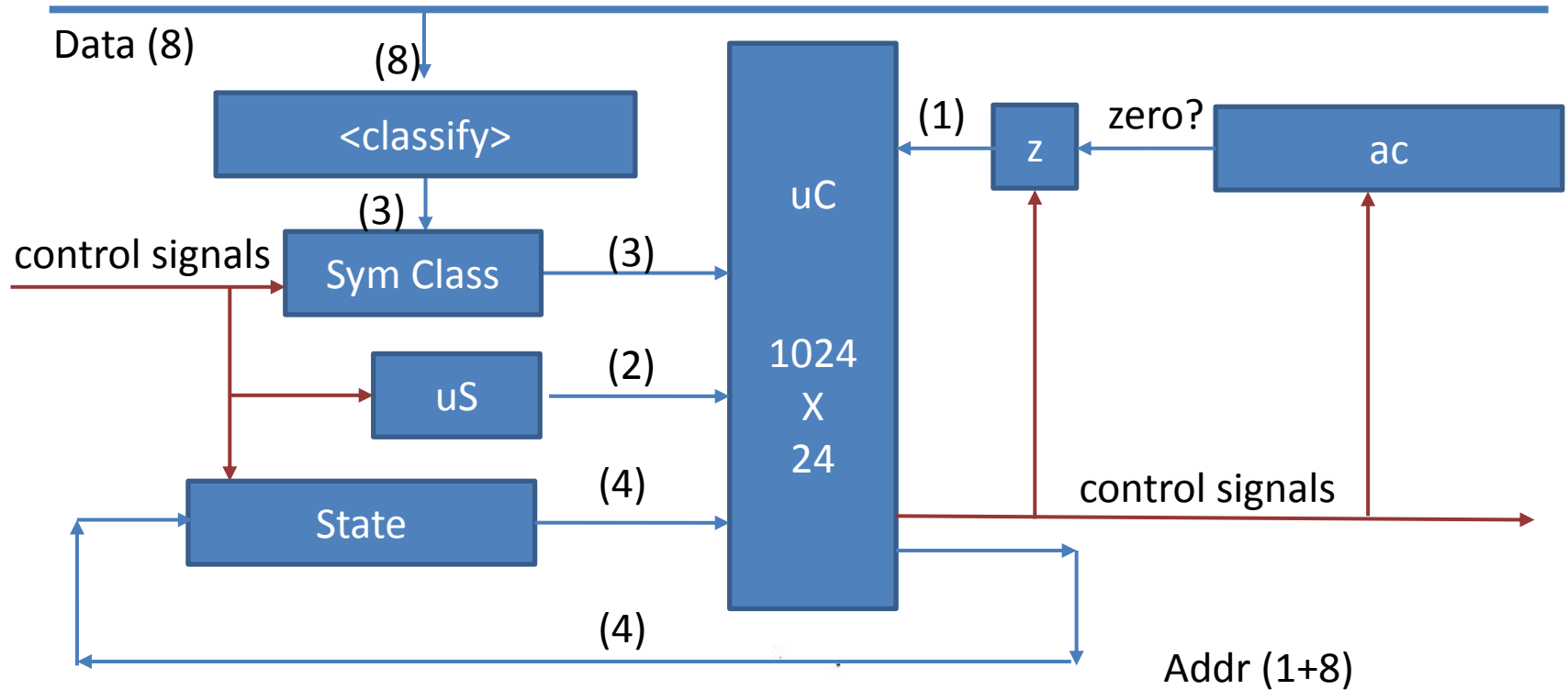


Registers

- SI – source index (8 bit)
- DI – destination index (8 bit)
- FI – function index (8 bit)
- RI – redex index (8 bit)
- AC – argument count (4 bit)
- V – value (8 bit)
- S – state (4), Z – zero (1), C – class (3)

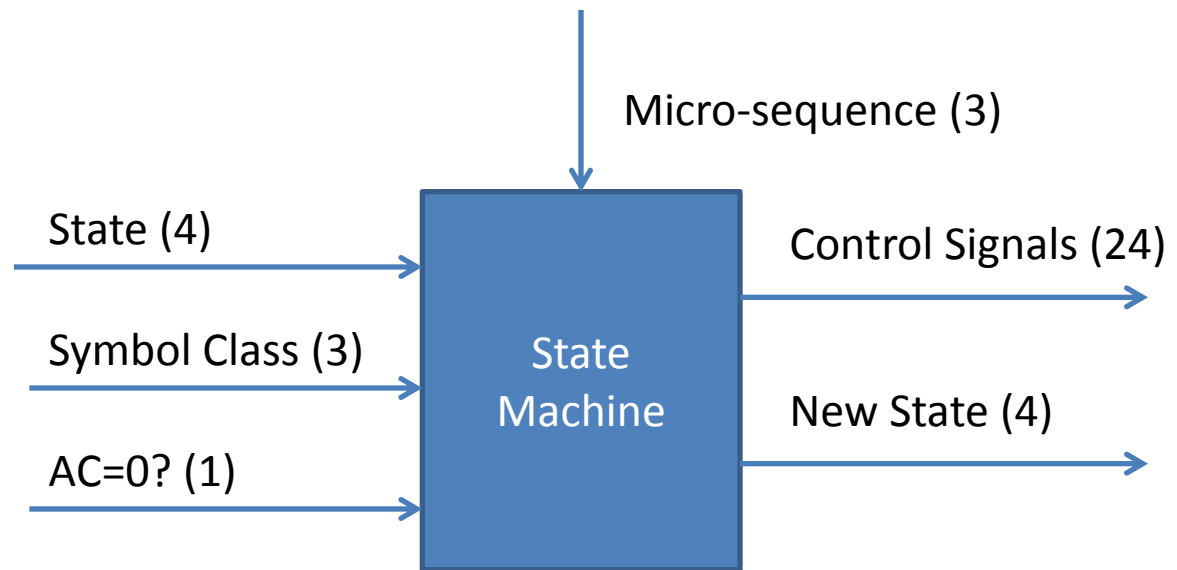


Architecture (ctrl unit)



Control Unit

- Turing-machine
- 16 states

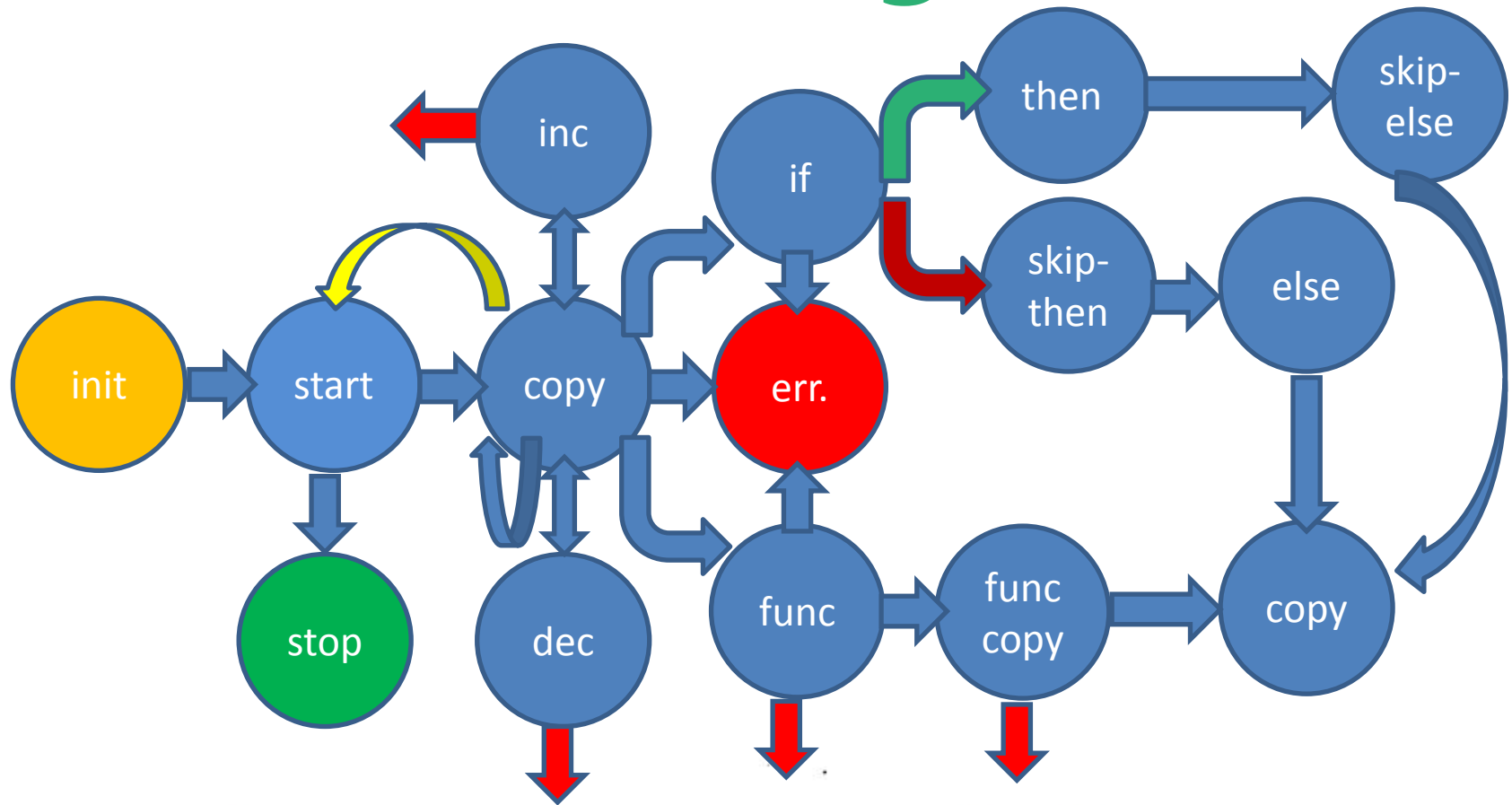


States

- Init - initialization
- Copy loop – copying
- If – evaluation of „if“
- inc/dec – evaluation of inc/dec
- Func/copy – evaluation of user functions
- Error – syntactical or internal error
- Stop – final state



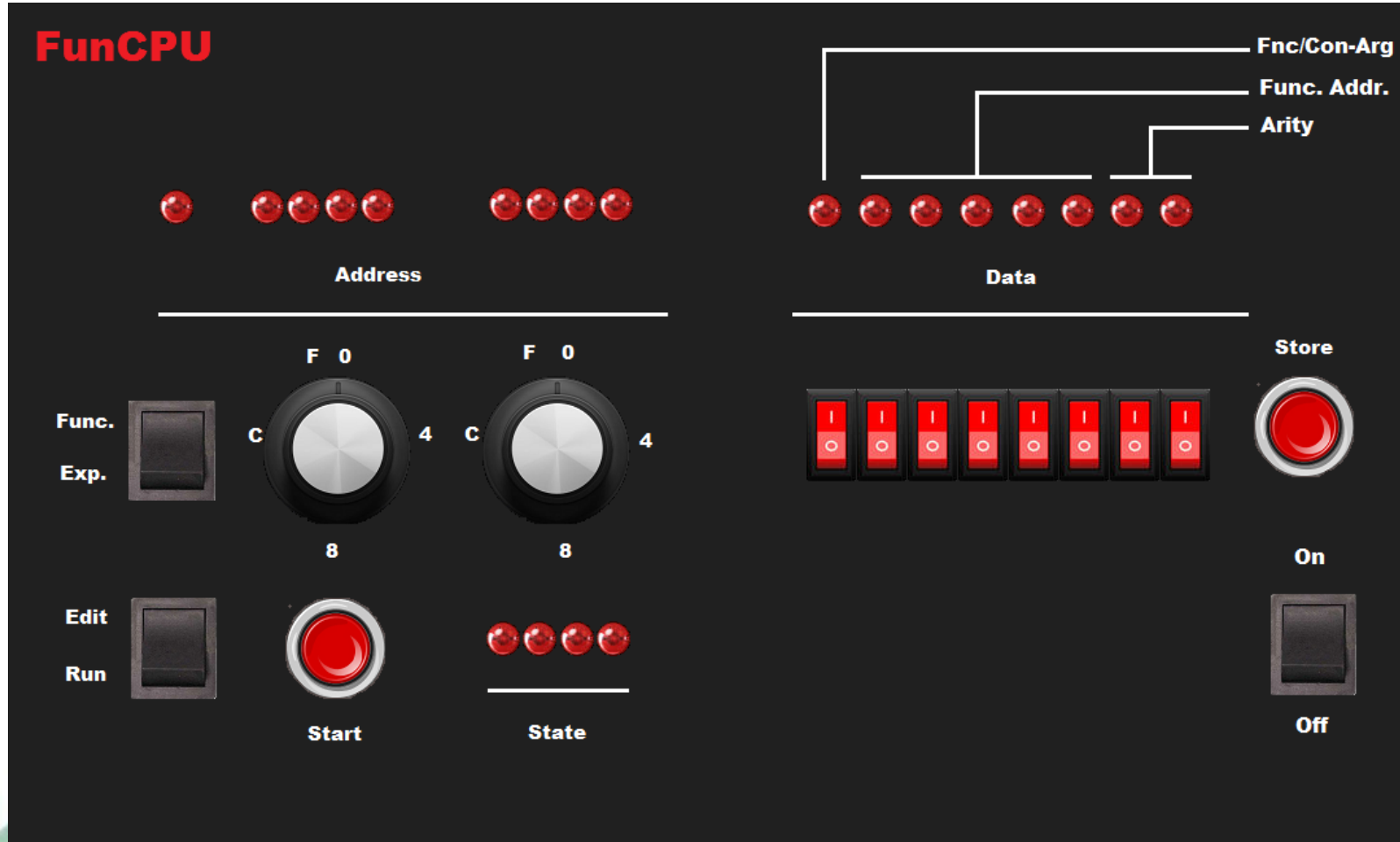
A horizontal strip of colorful, 3D letter cutouts in various shades of green, yellow, and brown, arranged in a dense, overlapping row. The letters are of different sizes and are scattered across the width of the image, creating a vibrant, textured border.



Implementation



Control Board



Physical Implementation

- 74HC163/161 – FI, AC / uC, SI, DI
- 74HC273 - RI, V
- 74HC273 - S, Z, C
- 74HC153 - 4 to 1 MUX (Addr, ALU src 1)
- 74HC157 - 2 to 1 MUX (ALU src 2)
- 74HC241 – octal tristate buffer
- MCM2708 – 1Kb EPROM (uCode, Classifier)



Computability



Overcoming Limitations

- Increase argument count: pairing functions, e.g. Cantor:
- $\pi(x, y) := \frac{1}{2} * (x + y) * (x + y + 1) + y$
- List: Gödel-encoding
- $list(x_1, x_2, \dots, x_n) := 2^{x_1} * 3^{x_2} * \dots * p_n^{x_n}$
- Only theoretical approaches due to the low precision of number representation.



Atomic Functions

- $c(x)$ CC FF
- $s(x)$ FC
- $\pi_i(x_1, x_2, x_3, x_4)$ 7F/7E/7D/7C FF
- FE 7F 7E FE F8 7F 7D 7C FF



Composition

- $h(x_1, x_2, x_3, x_4) = f(g_1(x_1, x_2, x_3, x_4), \dots, g_k(x_1, x_2, x_3, x_4))$
- f 83
- g_1 93, g_2 A3, g_3 B3, g_4 C3

83 93 7F 7E 7D 7C A3 7F 7E 7D 7C
B3 7F 7E 7D 7C C3 7F 7E 7D 7C FF



Primitive Recursion

- $h(0, x_1, x_2) = f(x_1, x_2)$
- $h(s(y), x_1, x_2) =$
 $g(y, h(y, x_1, x_2), x_1, x_2)$
- h 82, f 91, g A3

FE 7F 91 7E 7D A3 F8 7F

82 F8 7F 7E 7D 7E 7D FF



μ -operator

- $\mu(f) (x_1, x_2, x_3) = z \equiv \text{def}$
 $f(z, x_1, x_2, x_3) = 0$
 $f(i, x_1, x_2, x_3) > 0 \quad \forall i < z$
- f 82

FE 82 7F 7E 7D 7C 7F

82 FC 7F 7E 7D 7C FF



Improvements



Limitations

- Only integer types with limited precision.
- No advanced and/or user-defined types (e.g. list, record, tuples, etc.).
- Supports only 4 arguments.
- Small memory.
- No support for higher order functions.



Limitations – cont.

- Expression evaluation is slow and inefficient, requires more storage due to low-level basic functions.
- Functions cannot be embedded in arbitrary depth:
 $f(f(f(f(\dots$ is fine, but $f(f(f(f(f(f(\dots$ is not.
- Note: $f(f(f(f(\text{const}, f(\dots$ is ok again.
- Where f is a function with 4 arguments.

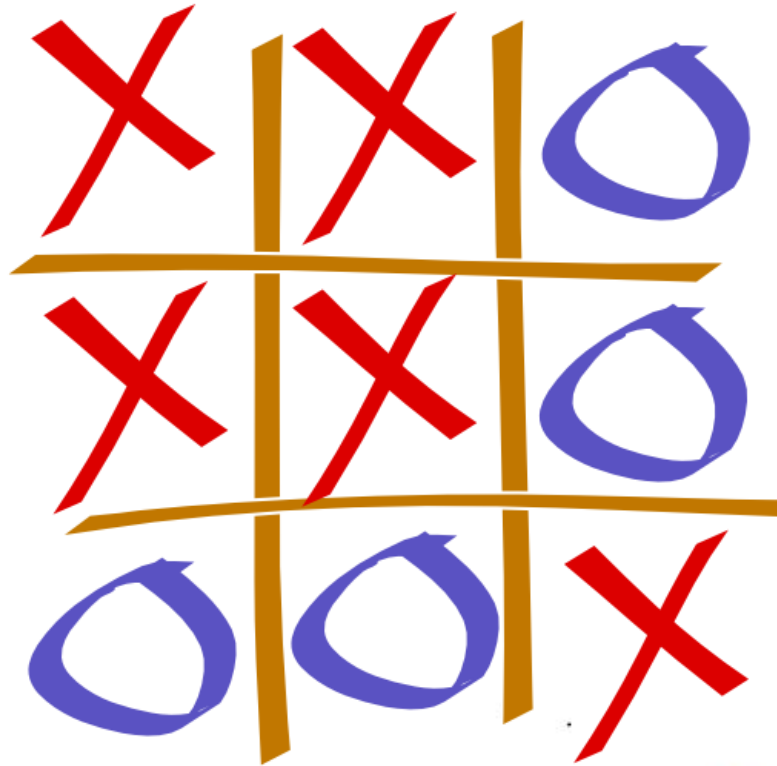


Enhancements

- More built-in functions (e.g. $+$, $-$, $*$, etc.)
- Increased memory
- Advanced timing: clock period reduced
- Increased register size (higher precision, deeper function embedding)
- More efficient expression evaluation
- Exploit parallelism



More Challenges...



Questions?



References / Sources

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<http://web.cecs.pdx.edu/~harry/Relay/>

