

BiiN

Object

Computing

Object Computing Benefits

- Productivity
 - Reusability
 - Maintainability
 - Extensibility
- Reliability

And

Object Computing Benefits

- To Store Knowledge

But

- OS File Structures and Relational DBMS are Inadequate
- Purely Software-based Persistent Object Support
 - is inefficient, and
 - does not address multi-user access

And

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Object Computing Benefits

- To Construct Secure Systems

But

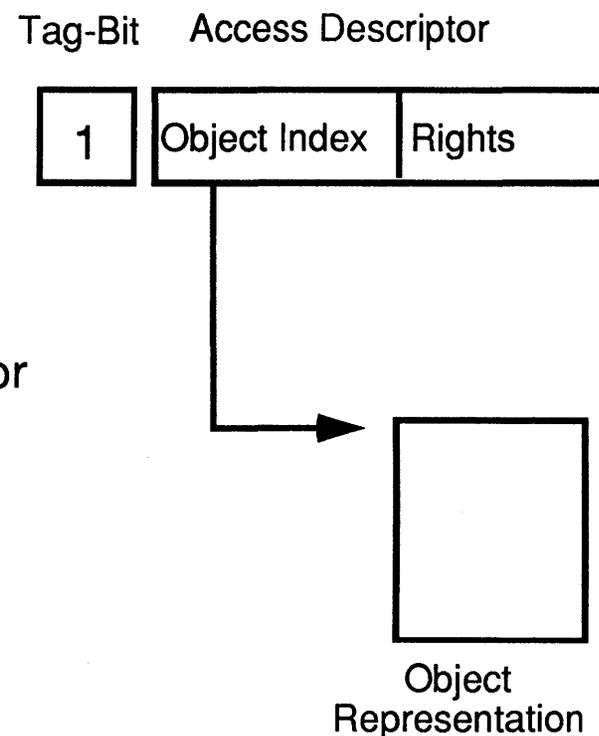
- Requires Run-time Enforcement
- Current Hardware Architectures do not Support Efficient Fine-Grained Protection

Topics

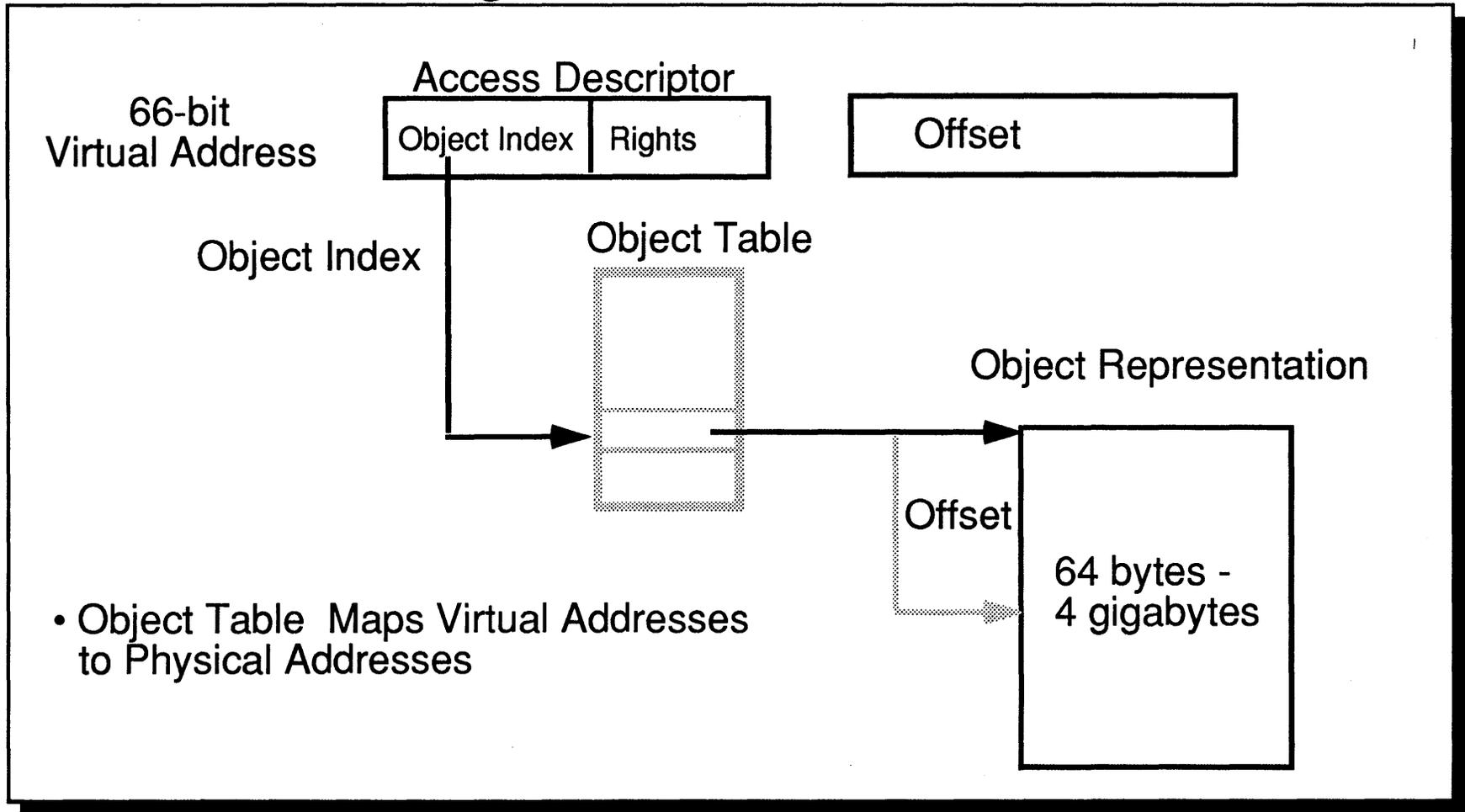
- Object Addressing and Protection
- Computational Model
- Type Manager Based Protection
- Inheritance
- Object Persistence

VLSI-Based Object Addressing and Protection

- Object is a Typed Memory Segment
 - From 64 to 4 gigabytes
 - If >4K bytes, then paged
 - up to 64M objects in virtual address space
- Object Accessed through Access Descriptor
 - AD is unforgeable (33rd bit)
 - AD is used just like a pointer – Access type in Ada
- Rights Determine Allowed Operations



Virtual Addressing



Object Structure

Tag-Bit

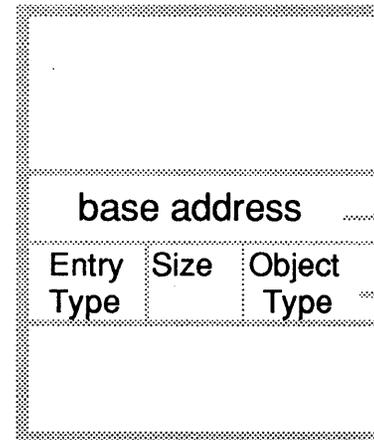
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Access Descriptor

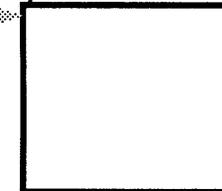
Object Index

Rights

Object Table



Object Representation

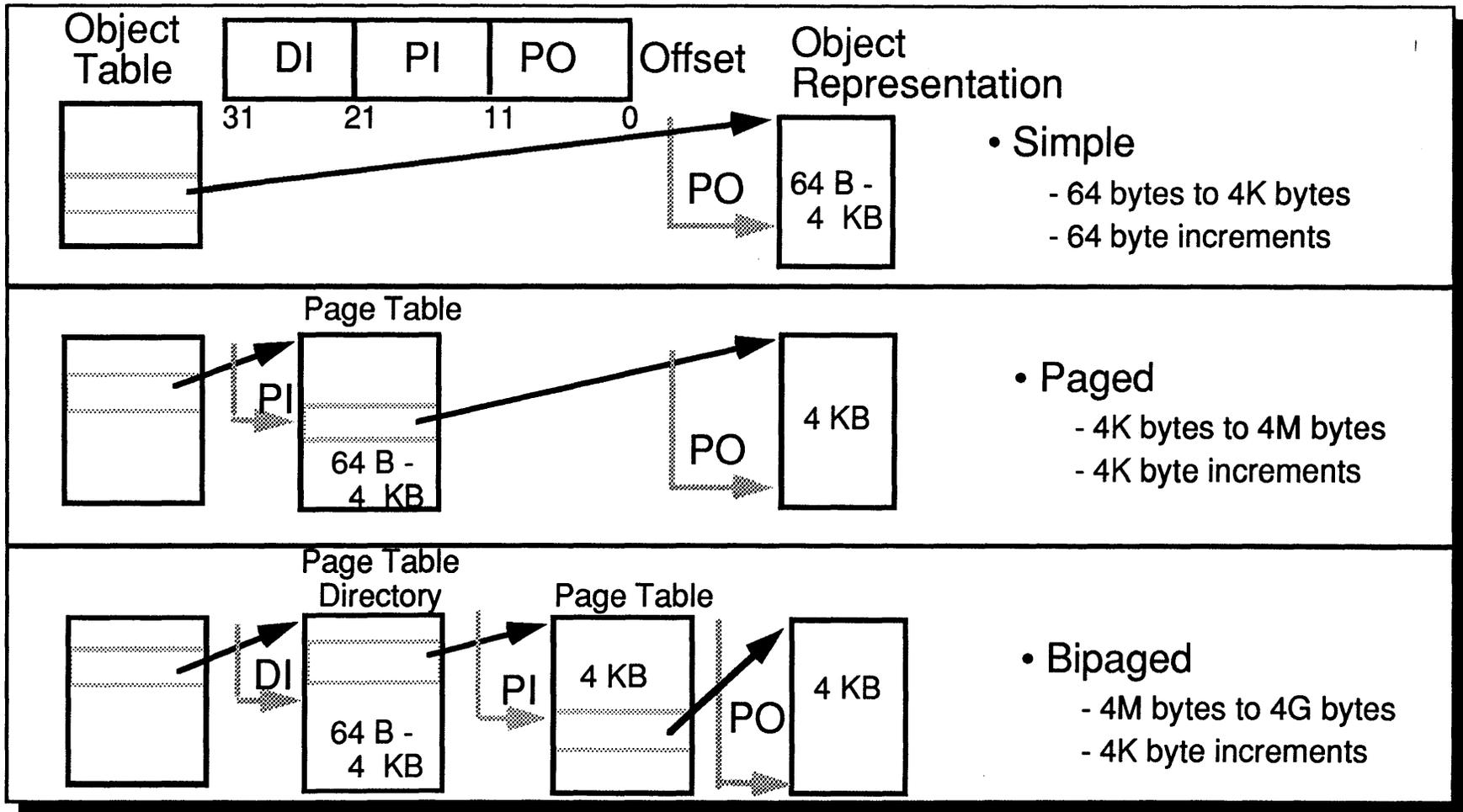


Type Definition



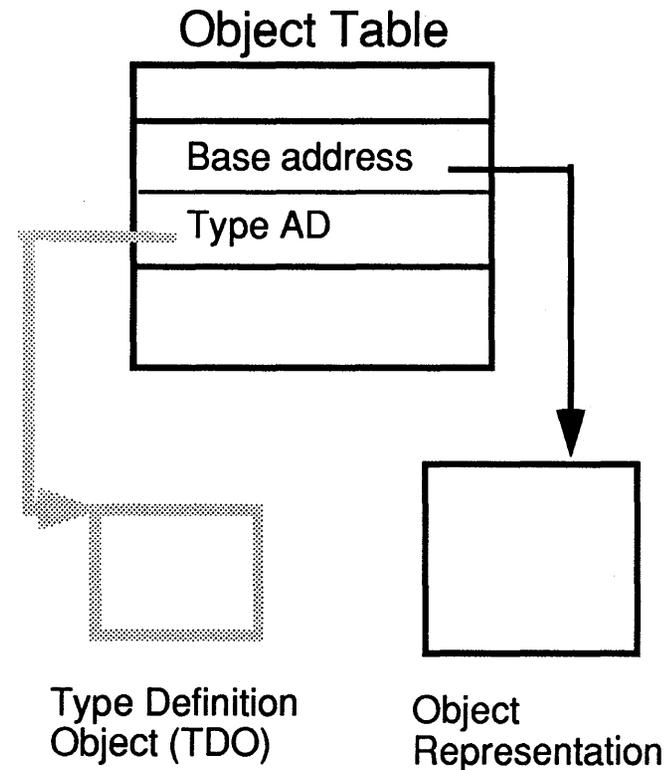
- Object Defined by Object Table Entry
 - Used for virtual to physical address translation
 - Visible to just OS memory manager
 - On-chip TLB (MMU) caches mapping information
- Translation is Independent of Execution Context
 - TLB not flushed on process switch

Address Translation

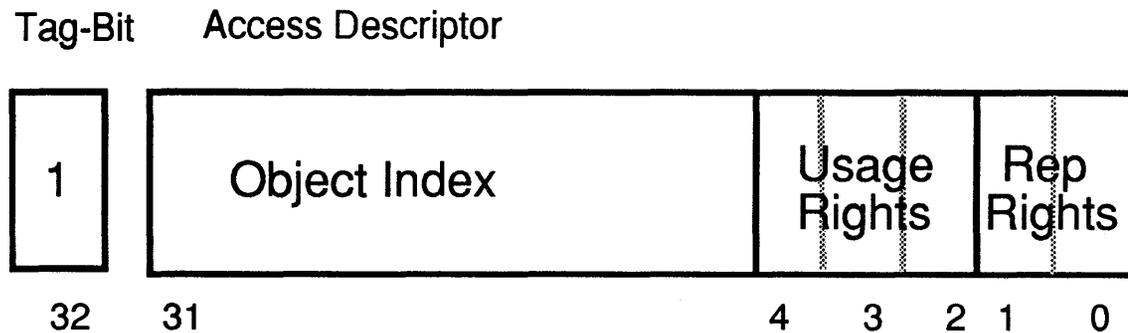


Object Typing

- Object Types are Defined by Type Definition Objects (TDO)
- Each Object Descriptor Contains an AD Pointing to the TDO of its Type
- H/W recognized types:
 - Semaphores
 - Ports
 - Processes
 - Domains
 - TDOs



Representation Rights

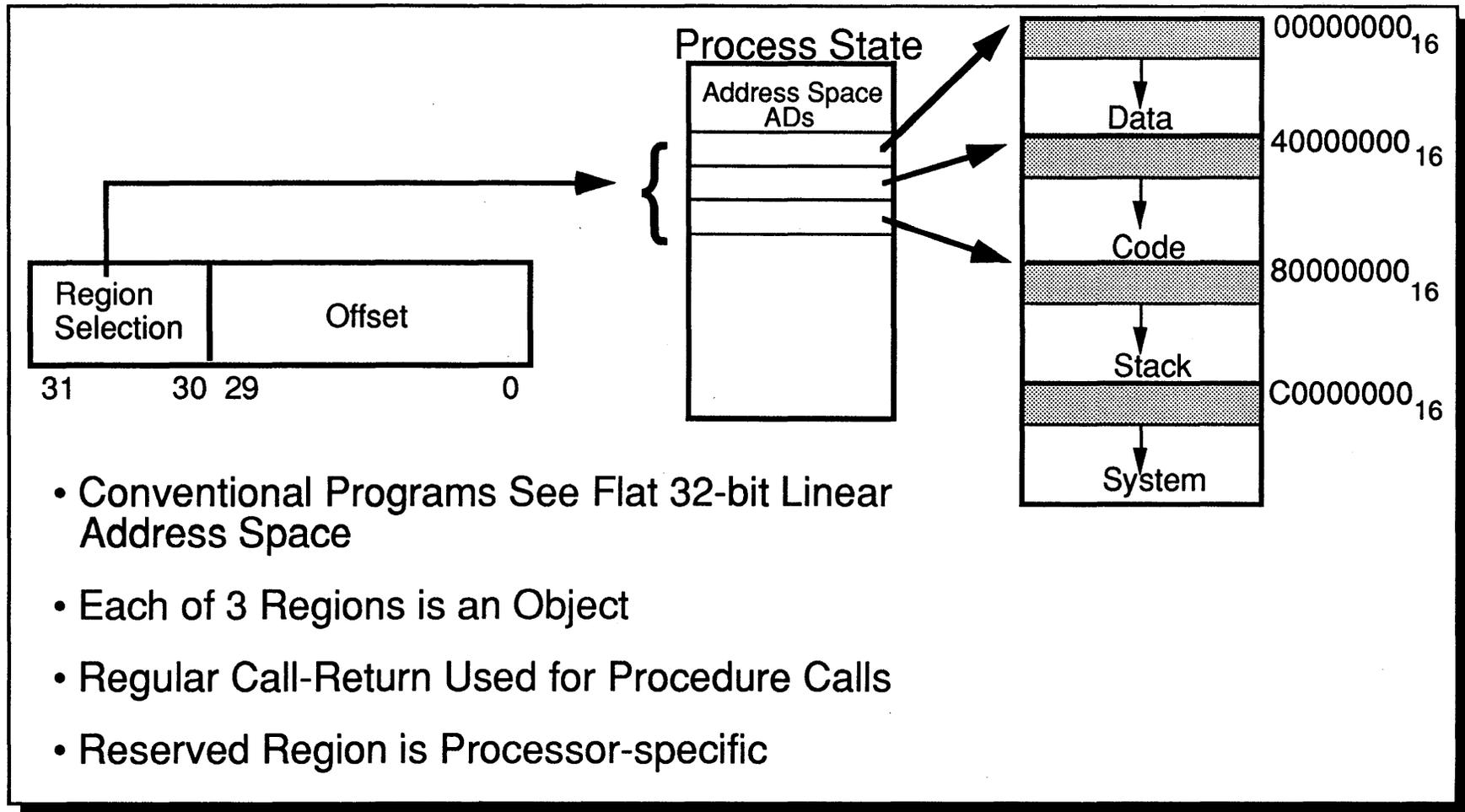


- Representation Rights Control Direct Access to Object
 - Read rights checked on Load instructions
 - Write rights checked on Store Instructions
- Rights Checking Occurs in Parallel with Address Translation

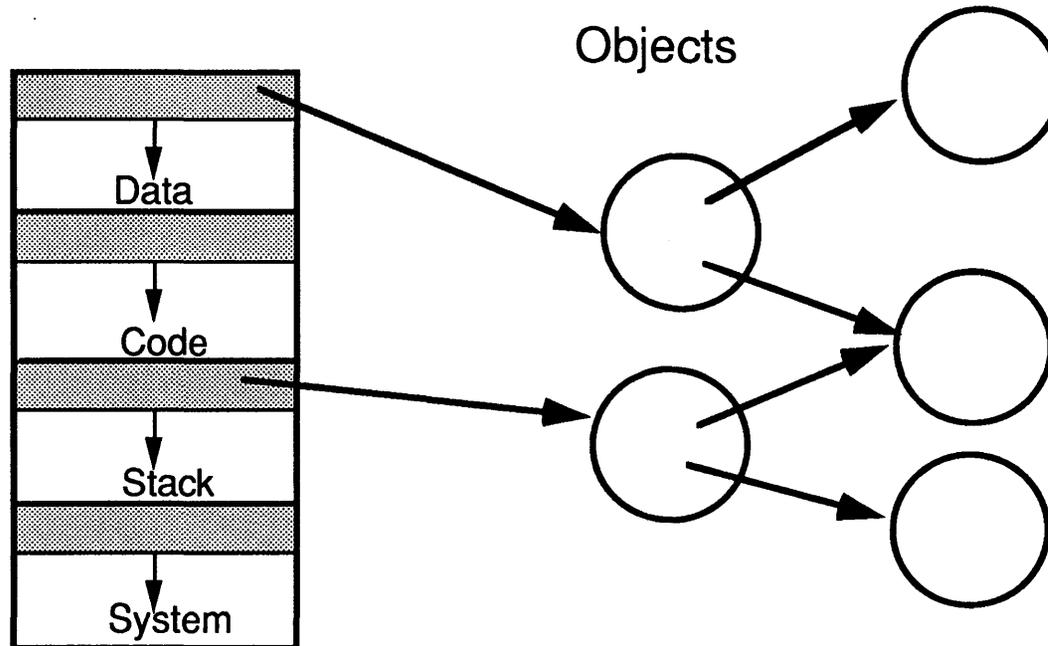
Topics

- Object Addressing and Protection
- **Computational Model**
- Type Manager Based Protection
- Inheritance
- Object Persistence

Simple Program Model

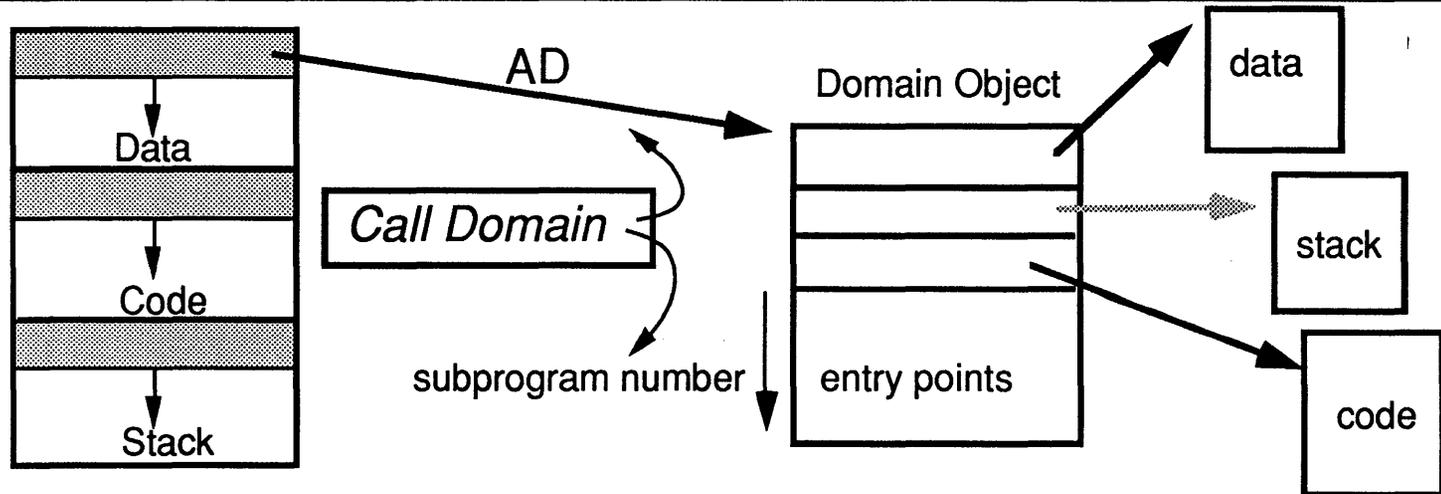


Extended Program Model



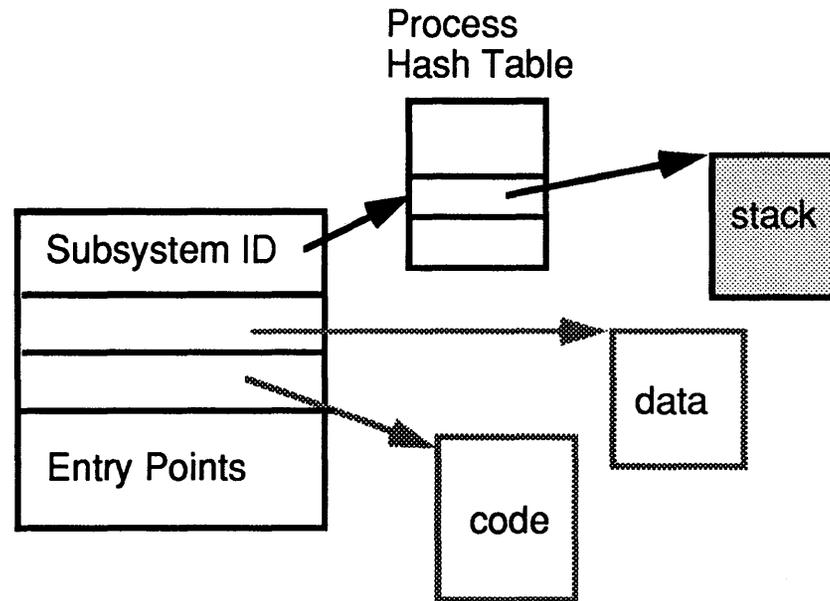
- Linear Address Space Can be Root of an Object Network
 - Interconnected by ADs
 - ADs used just like pointers

Multiple Address Spaces Per Process



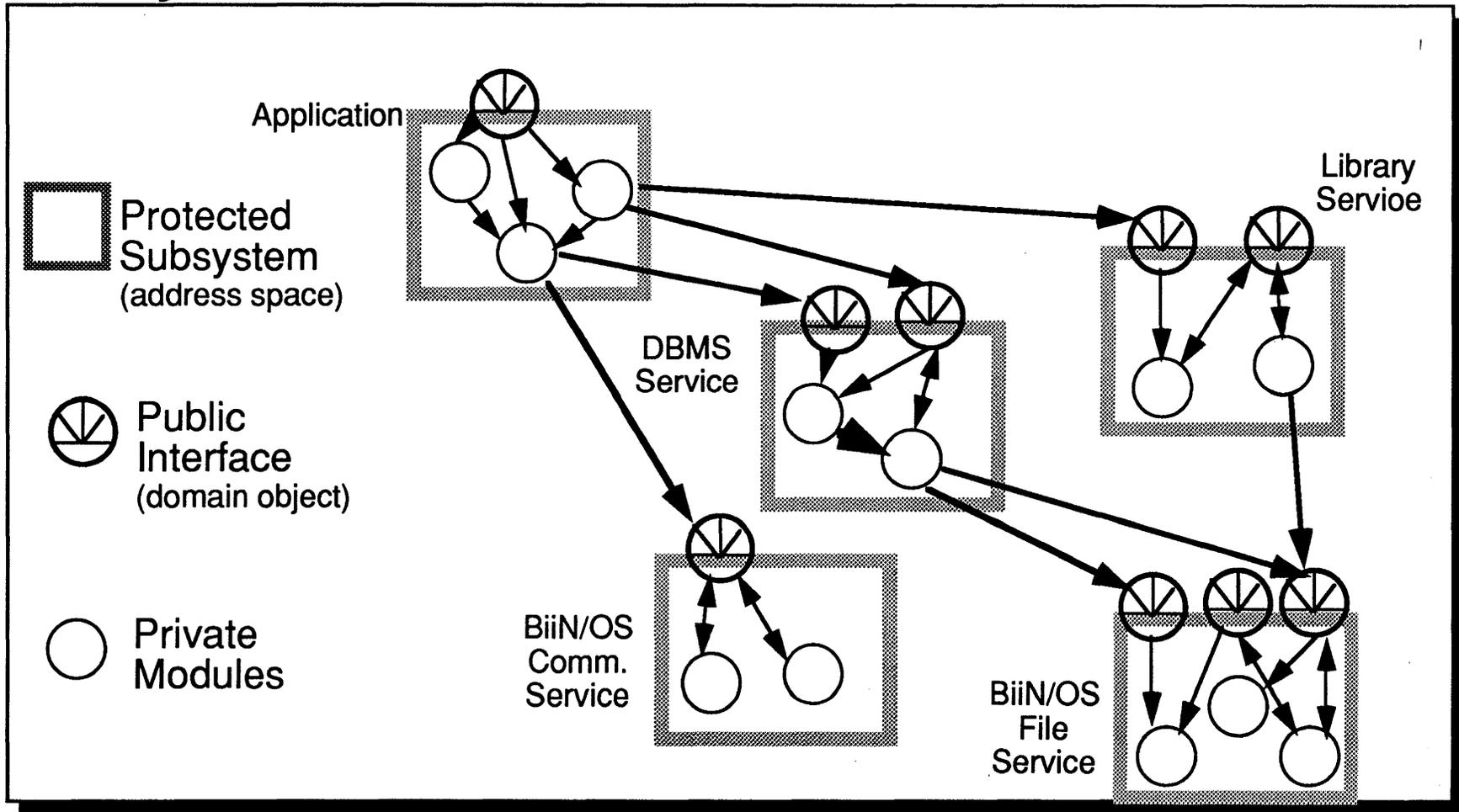
- Destination Address Space Defined by Domain Object
- Interdomain Call/Return Instructions
 - Linkage kept in process object
 - No implicit access between address spaces
 - Parameters may include ADs
 - Performance comparable to other architectures' supervisor calls
- No Limit on Number of Address Spaces per Process

Closer Look At Call Domain



- Subsystem_ID Selects Stack from Process Object
 - Null: use caller's stack
 - World: use the default program stack
 - ID: A unique (named) stack is selected per process

Subsystem Based Protection



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Subsystem Based Protection Benefits

- Reliability, Maintainability, and Extensibility Without Compromising Performance
- Increased Productivity in Integration and Test
 - Decompose application into protected subsystems
 - Since each subsystem is linked independently, turnaround time (recompile/relink) is faster
 - Since errors confined to subsystem, they are easier to find
- Increased Performance Without Compromising Security
 - Services can safely execute in user's process
 - Other architectures require separate process, which results in:
 - higher invocation overhead
 - potential bottlenecks in symmetric multiprocessors

Topics

- Object Addressing and Protection
- Computational Model
- **Type Manager Based Protection**
- Inheritance
- Object Persistence

Object-Oriented Design

- Define Abstract Data Type
- Define Set of Operations on Type
- Set of Operations form a Module
- Module Hides Implementation
 - Representation of data type
 - Operations (Algorithms) on data types

Mapping to Ada

Object-oriented Design

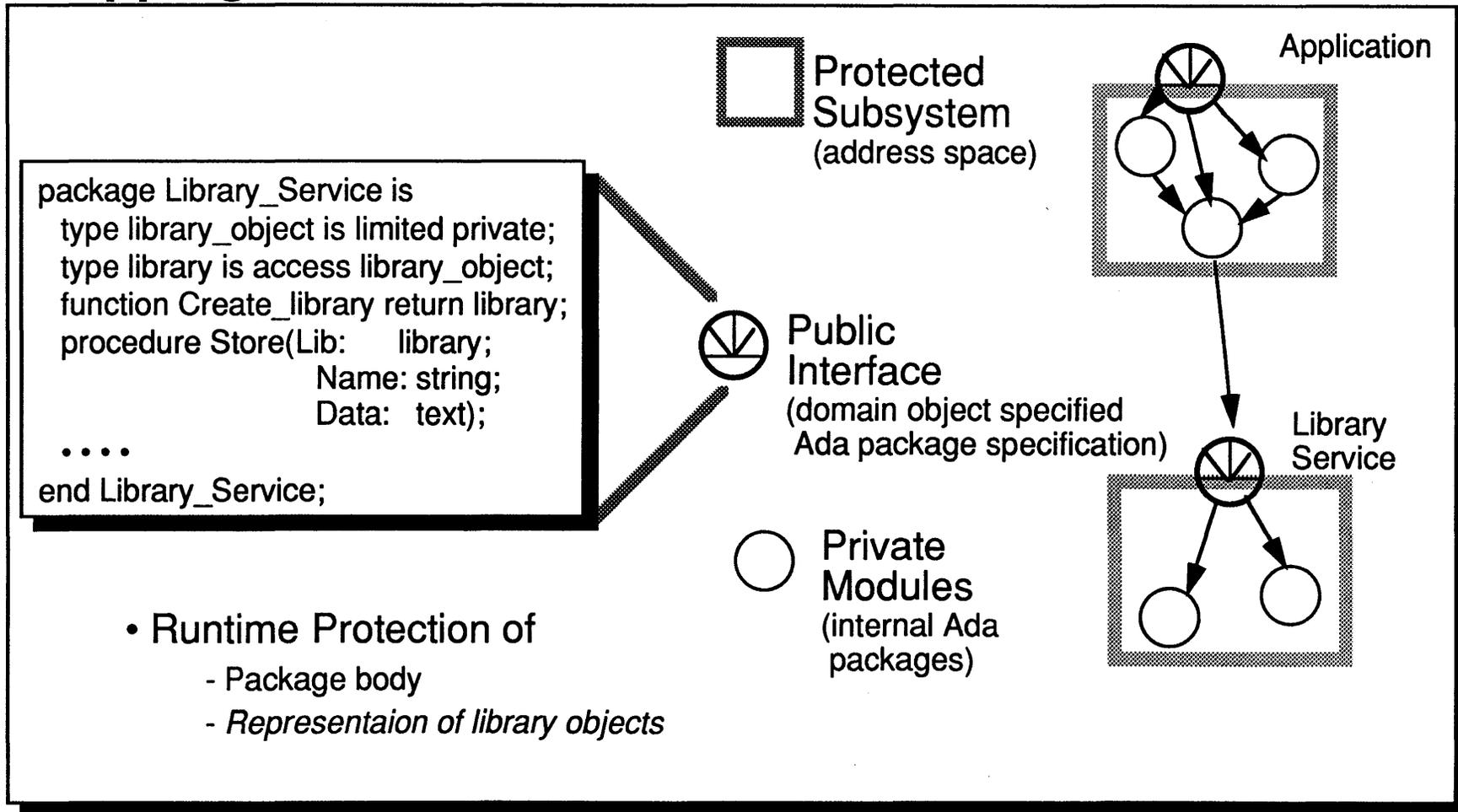
- Define Abstract Data Type
- Define Set of Operations on Type
- Set of Operations form a Module
- Module Hides Implementation
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Maps to Ada Package

```
package Library_Service is
    type library_object is limited private;
    type library is access library_object;
    function Create_library return library;
    procedure Store(Lib: library;
                   Name: string;
                   Data: text);
    . . . .
end Library_Service;

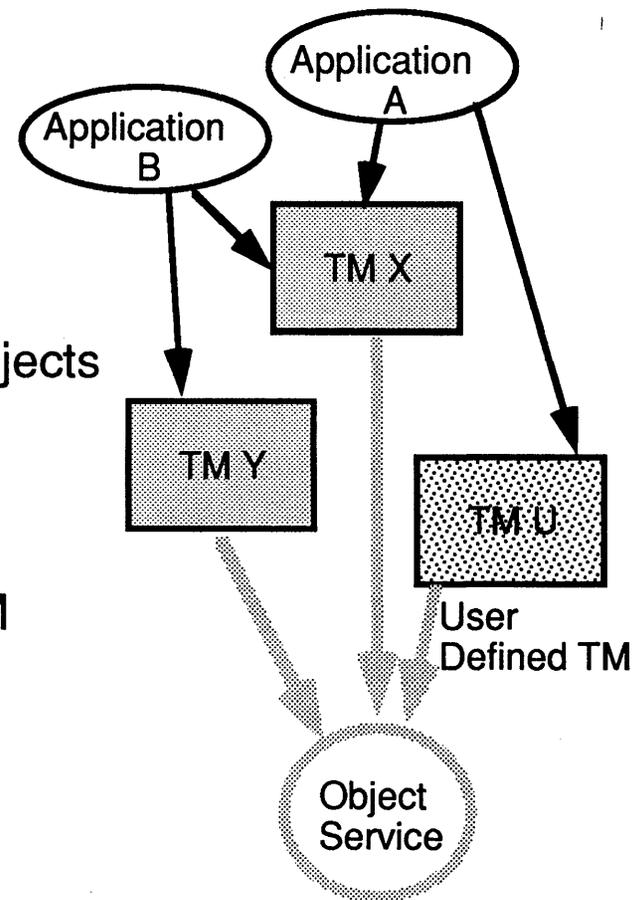
package body Library_Service is
    -- Contains Implementation
    -- Hidden from users of package
```

Mapping To Architecture

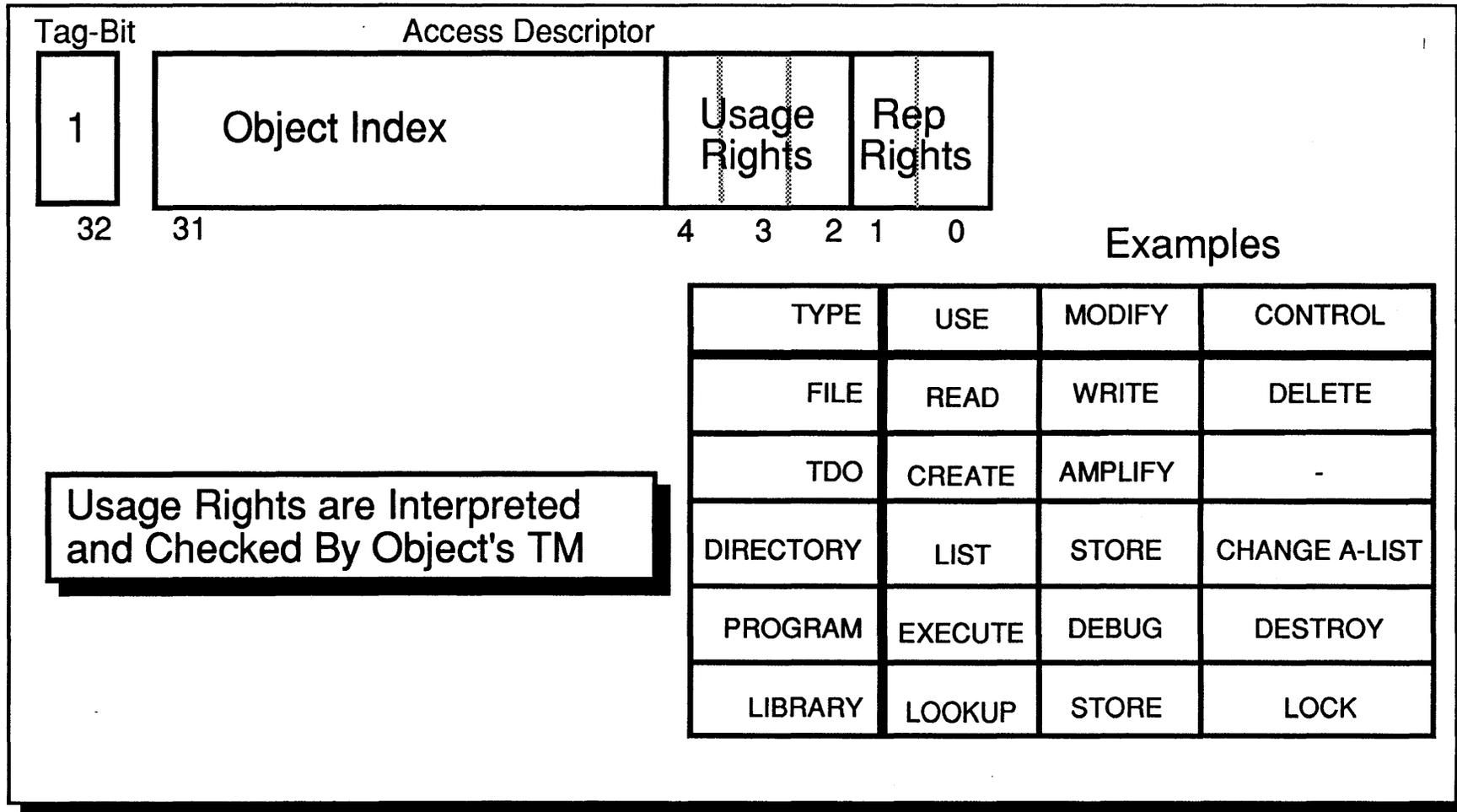


Principles Of Type Manager Based Protection

- Objects are Typed
- The *X Service* is the Type Manager (TM) for Objects of Type X
- Only TM X Can Access Representation of X Objects
- Applications Can Pass Around ADs (without Representation Rights) for X objects
- Anyone Can Create a New Object Type and TM
- BiiN/OS Provides Object Management Service



Usage Rights



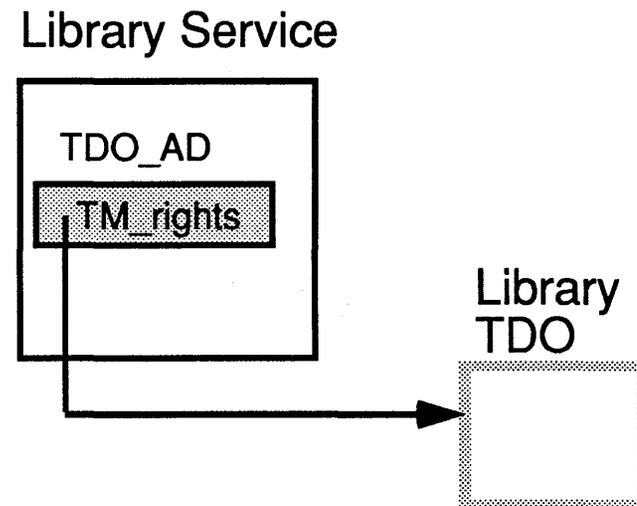
A TM Example Using Libraries

Outline

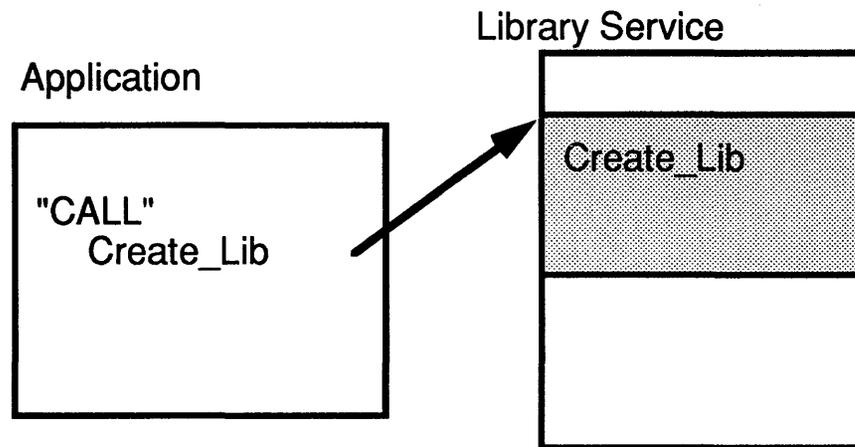
- Creation of the Library Type
- Creating an Object of Type Library
- User-Level Protection on Library Objects
- Invoking the Lookup Procedure on a Library Object

Type Creation

- To Define a New Data Type, TM Creates a Type Definition Object (TDO)
 - AD to TDO is never given out
 - Usage Rights of that AD have TM_rights semantics
 - Create (use) right
 - Amplify (modify) right



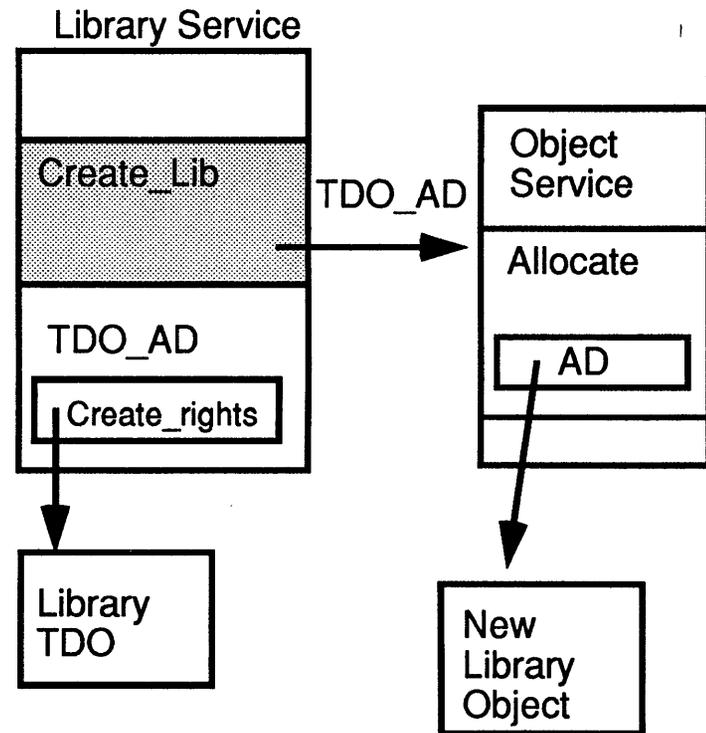
Object Creation



- Application Cannot Create Library Object Directly
 - It cannot get AD with create rights for Library TDO
- Application Must Call Create_Library Subprogram
- Call Requires AD for Library Service Domain Object

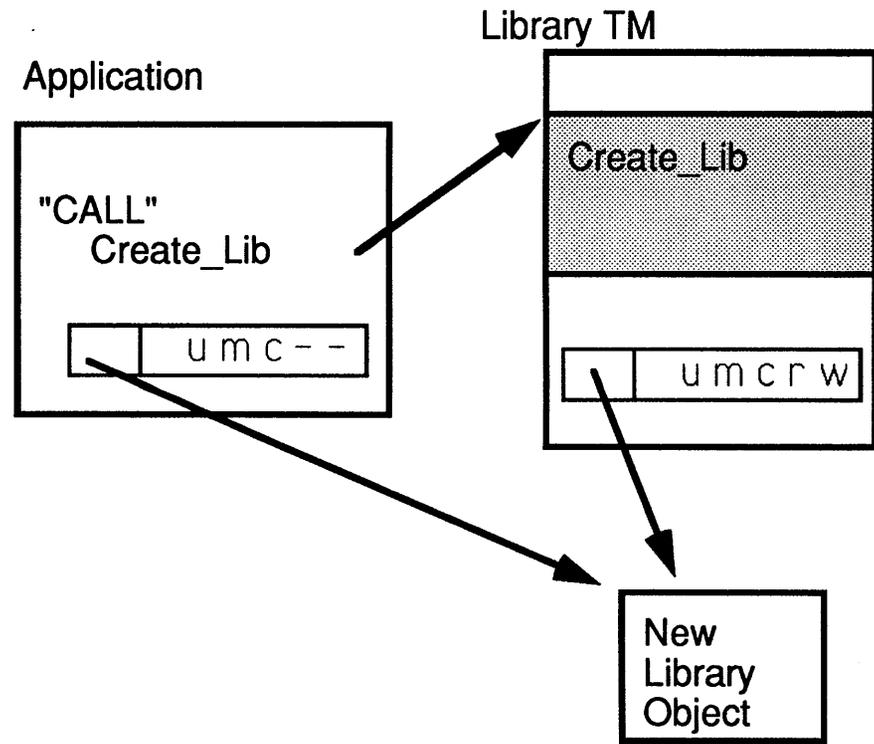
Object Creation (cont'd)

- Library Service Calls BiIN/OS Object Service
 - AD for Library TDO is passed
 - Object service checks for create_rights in AD
- Object Service Allocates Library Object
 - Returns AD with all rights to Library TM

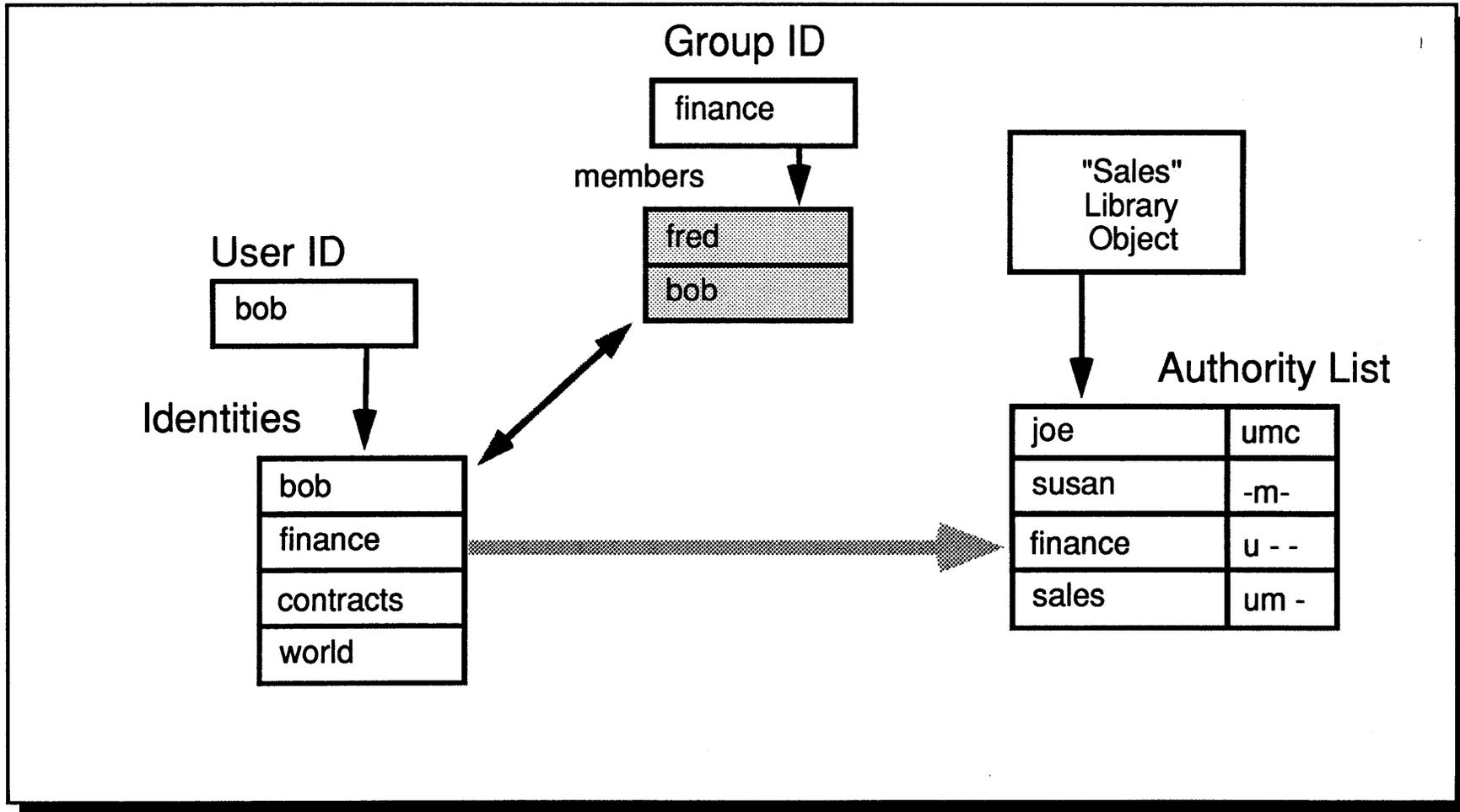


Object Creation (cont'd)

- Library Service Then
 - Initializes library object
 - Removes representation rights in AD (using restrict rights instruction)
 - Returns AD to application
 - Does NOT keep AD to library object
- Application Cannot Access Representation
- Application Controls Usage Rights to Just its Library Object

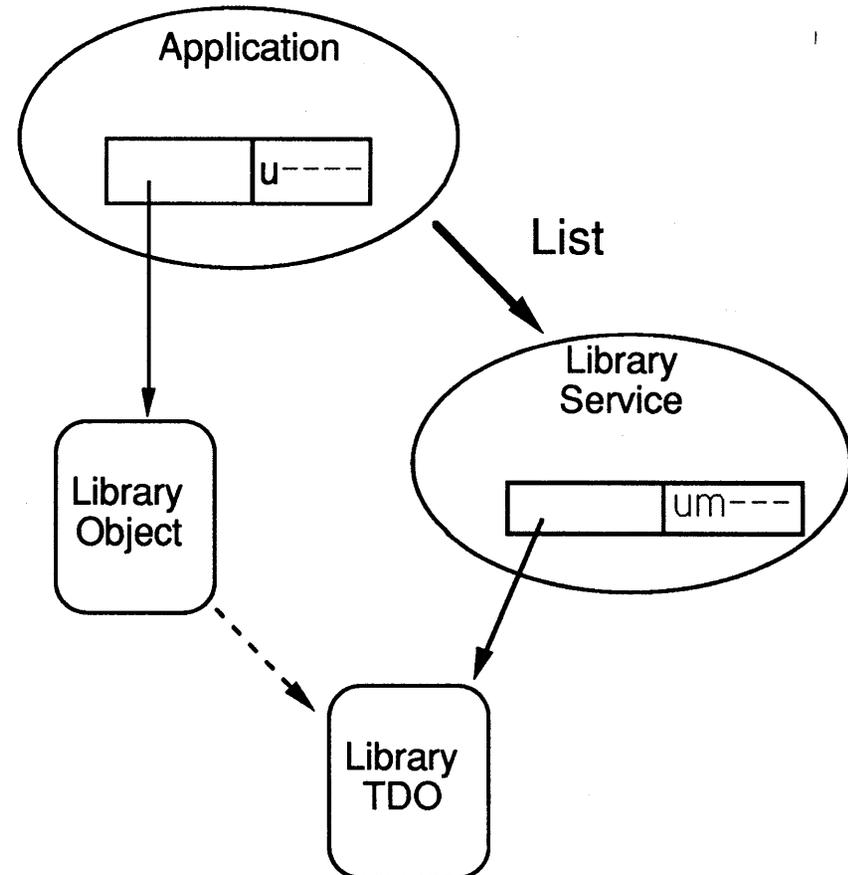


Authority-List Determines Usage Rights



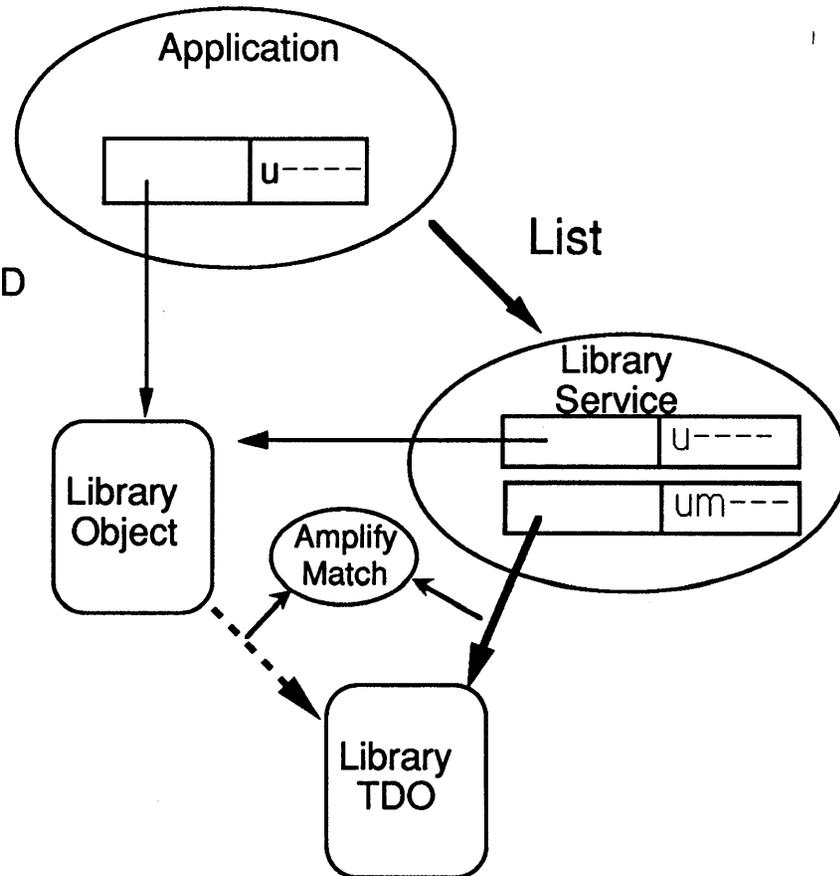
Type Specific Operations

- Application has AD for Library Object with Just Use (Lookup) Rights
- Application Invokes List Subprogram Passing AD



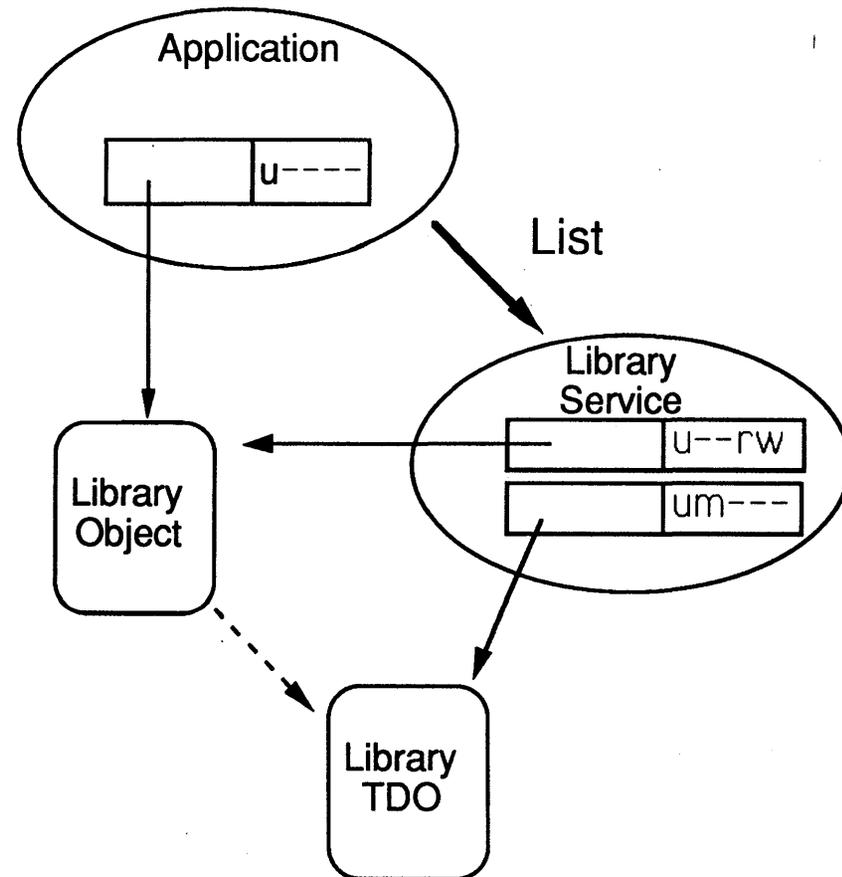
Type Specific Operations (cont'd)

- Application has AD for Library Object with Just use (lookup) Rights
- Application Invokes List Subprogram Passing AD
- Library Service Executes Amplify Instruction
 - Takes AD to-be-amplified and AD with Amplify (modify) rights for a TDO
 - Verifies type match
 - Adds representation rights



Type Specific Operations (cont'd)

- Application has AD for Library Object with Just use (lookup) Rights
- Application Invokes List Subprogram Passing AD
- Library Service Executes Amplify Instruction
 - Takes AD to-be-amplified and AD with Amplify (modify) rights for a TDO
 - Verifies type match
 - Adds representation rights
- Library Service Can Now Access Representation of Library Object



Relationship To Security

"The TCB shall be designed and structured to use a complete, conceptually simple protection mechanism with precisely defined semantics. This mechanism shall play a central role in enforcing the internal structuring of the TCB and the system. The TCB shall incorporate significant use of layering, abstraction and data hiding. "

From the Orange Book (DOD 5200.28-STD) section 3.3.3.1.1, System Architecture for B3 level security

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Relationship to Security (cont'd)

Corresponds to SAT Type Enforcement

- Secure Ada Target
- Extends Bell and LaPadula Model Beyond A1
- NCSC/Honeywell Research

"Domains are essentially a mechanism for encapsulating managers for different data types and transformations between data types. This provides a way to decompose the proof of security for the system into manageable pieces and to tailor the security policy for a system in an application dependent fashion. ... Thus, type enforcement is more than a mere convenience. It provides a way to unify the treatment of trusted subjects with that of generic untrusted subjects."

From "Extending the Noninterference Version of MLS for SAT", IEEE Transactions on Soft. Eng. Feb. 1987.



Topics

- Object Addressing and Protection
- Computational Model
- Type Manager Based Protection
- **Inheritance**
- Object Persistence

Behavior Inheritance

- Multiple Implementations for Same Behavior

- Example

Behavior

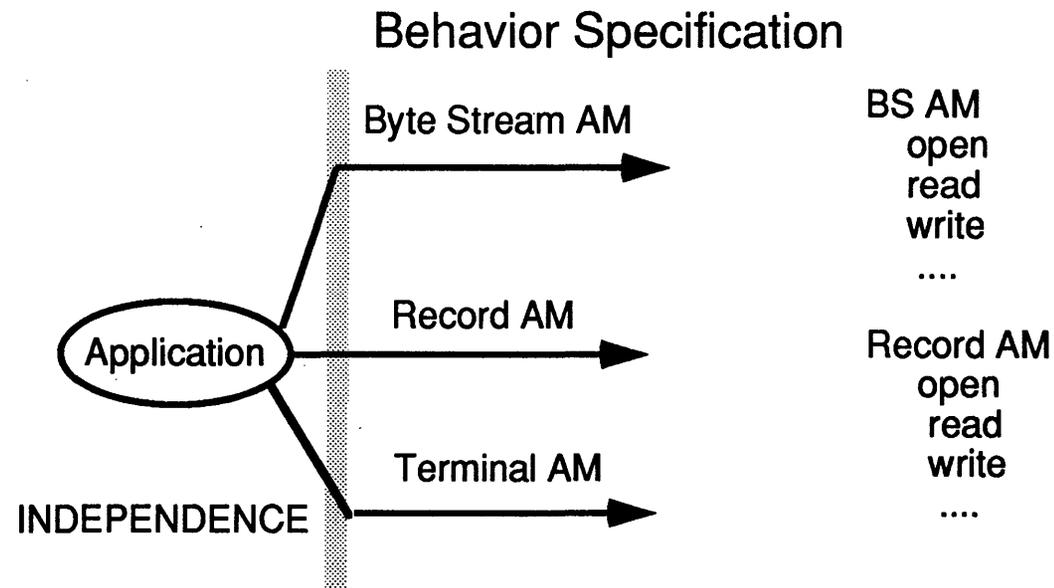
Byte-Stream Access Method

Implementations

File, Pipe, Magtape, Terminal

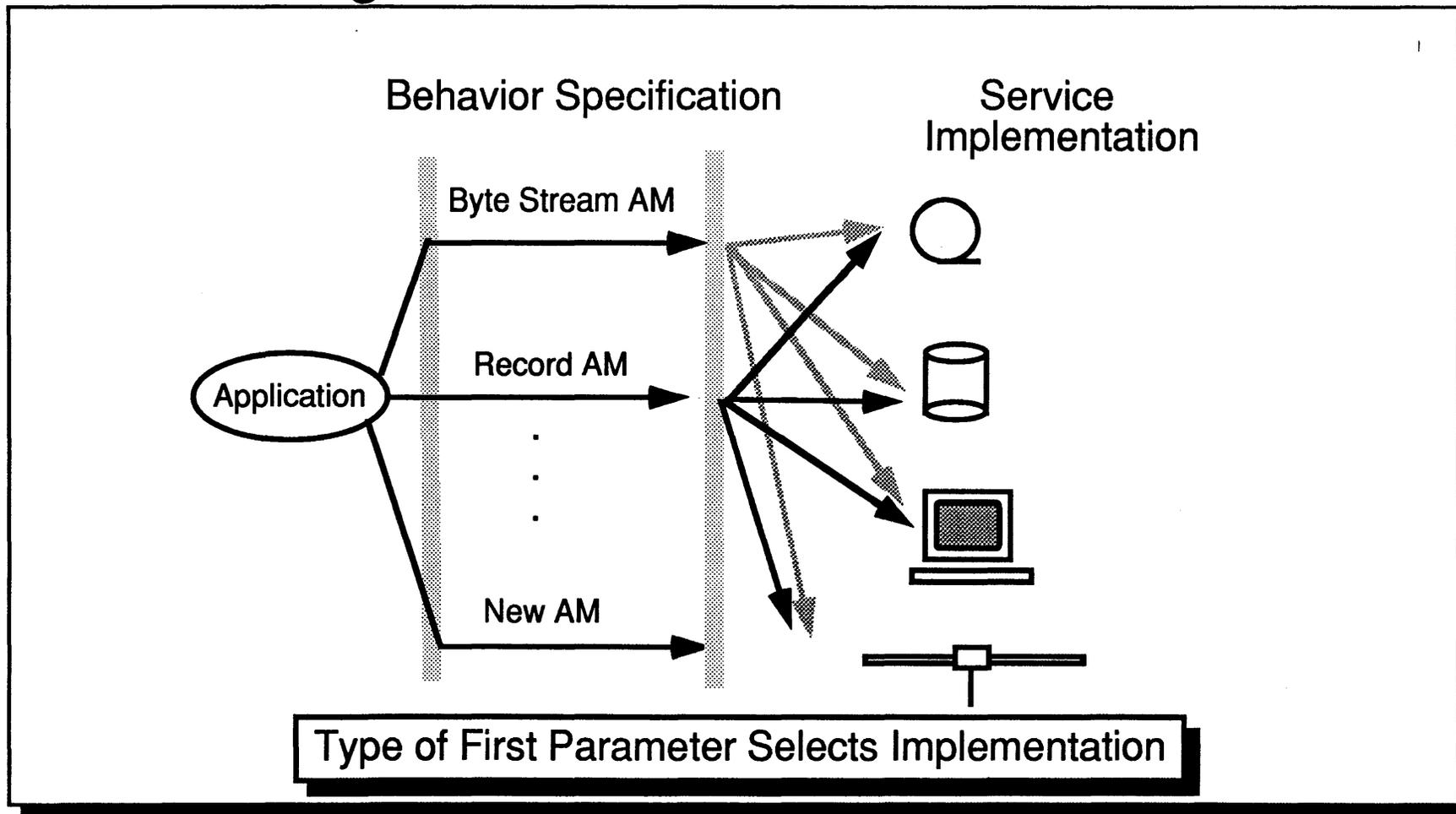
- Implementation is Selected in Call Instruction Based on Type of Object

Application Independent of Implementation

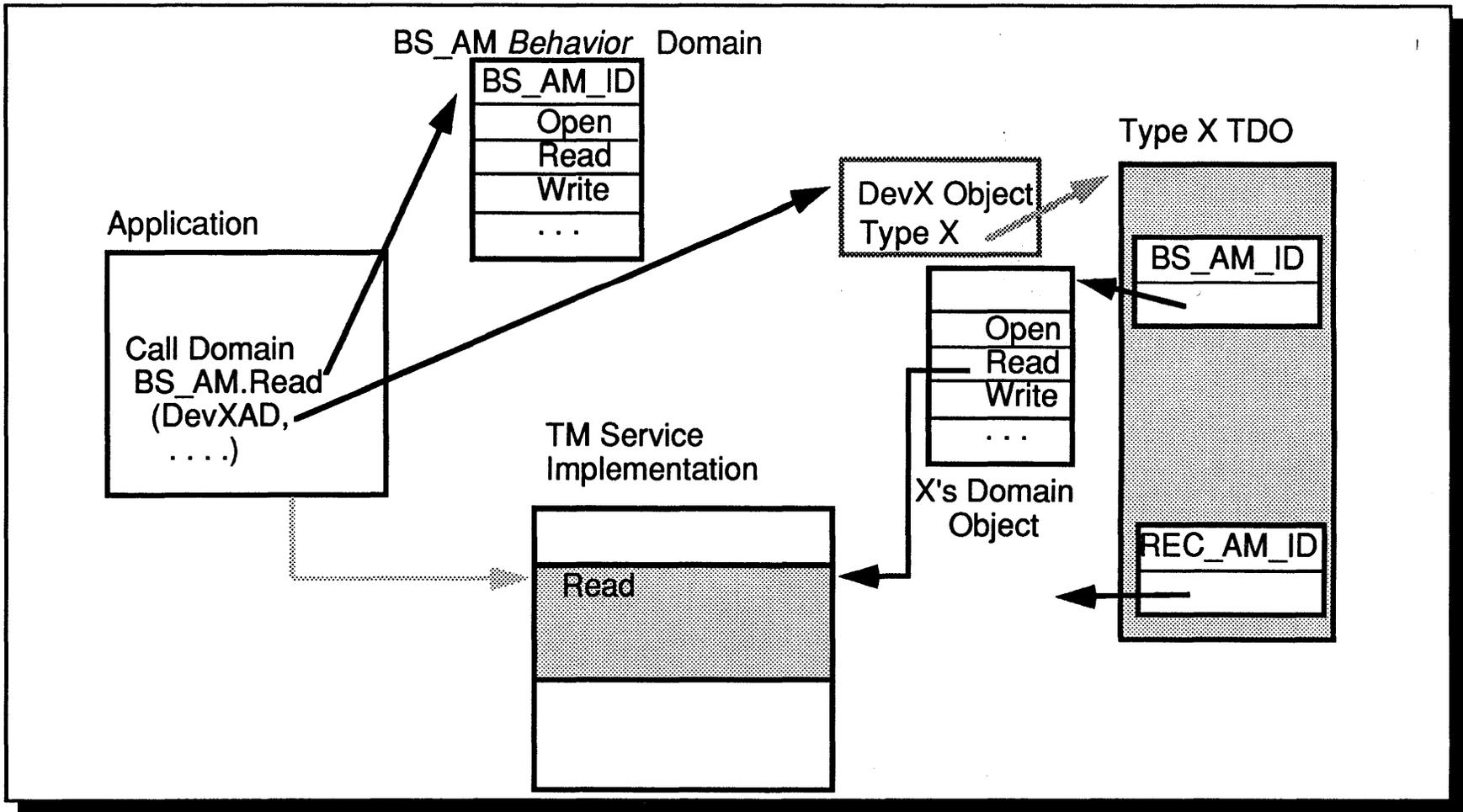


- Ada Package Specification Specifies Behavior
- Provides Semantic Definition to Interface
- Services Provide Implementation (package body)
- New Behaviors can be Added

Call Vectoring



Implementation Selected by Object's Type



Behavior Inheritance Summary

- Dynamic Binding on Every Call Based on Type
- Old Program Binaries Work with New Implementations Without Even Relinking
- *Call Behavior* Instruction Same as Call Domain
 - Different effect due to difference in domain objects
 - Thus, invocable from any language
- Service Can Dynamically Add Implementations for New Behaviors

Topics

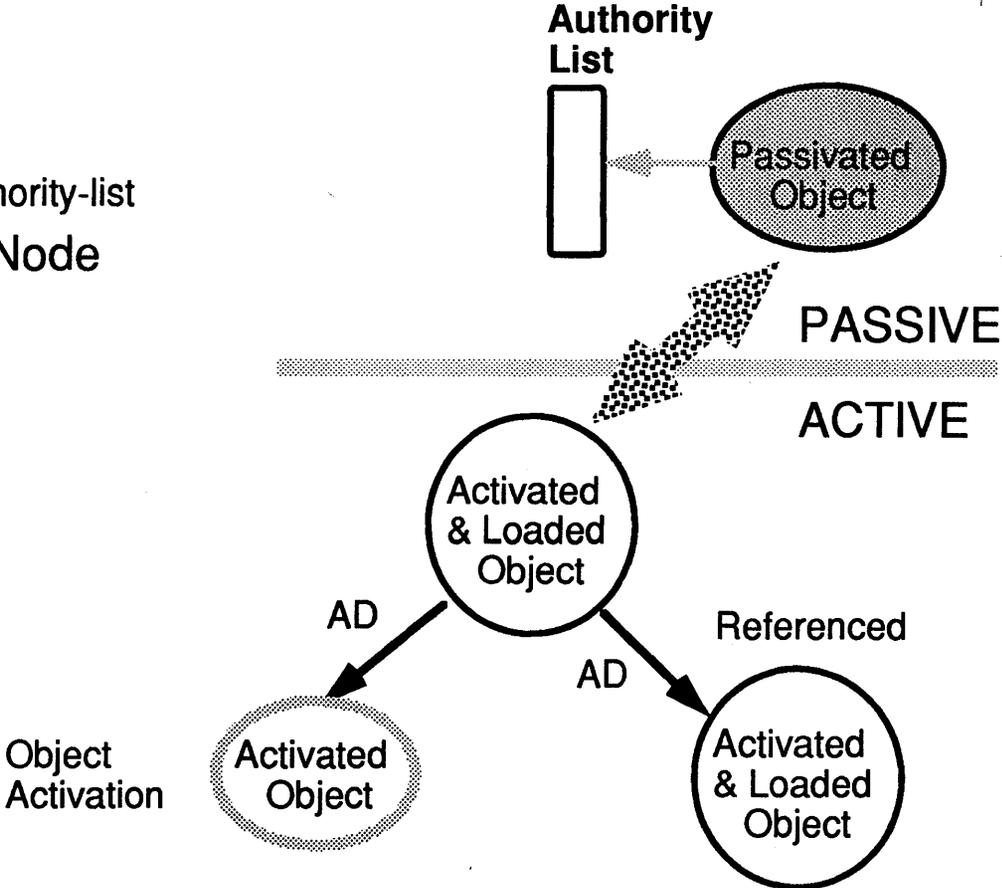
- Object Addressing and Protection
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- Inheritance
- **Object Persistence**

Object States

- Two object states: *Passive* and *Active*
- A Passivated Object is
 - Stored in permanent storage (disk)
 - Managed by Passive Store Management
 - Protected by Authority List
- An Activated Object is
 - Stored in virtual memory
 - Managed by object service
 - Protected by VLSI-based object addressing

Activation

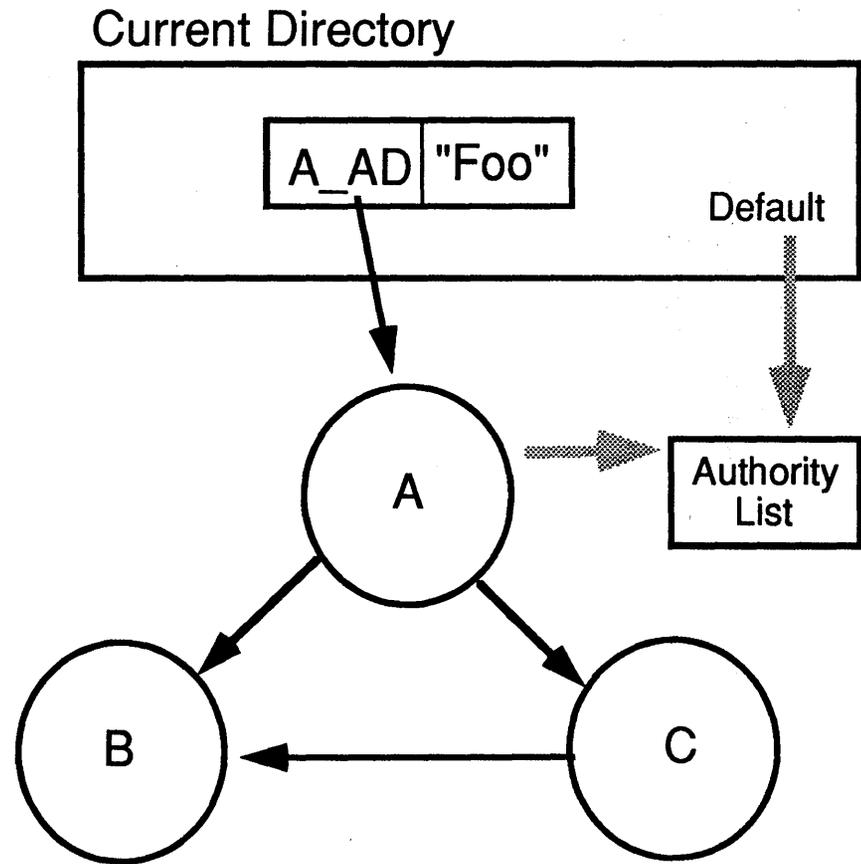
- Implicit
 - Similar to VM fault
 - Rights in AD determined by authority-list
- Object Networks Can Cross Node Boundaries



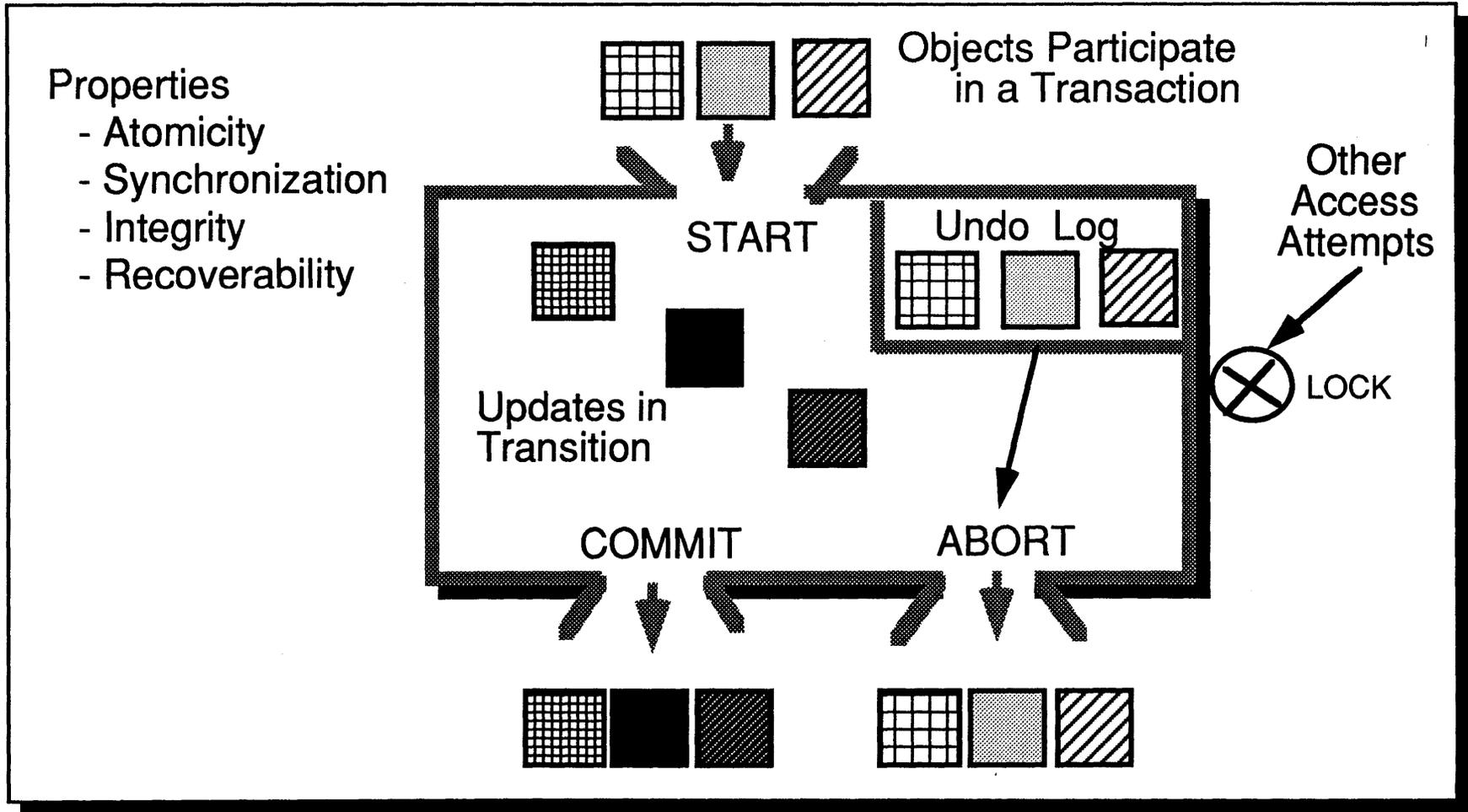
Passivation

- Explicit
 - Transaction-based for synchronization and recovery
 - Controlled by type manager to assure consistency

```
Start_Transaction;  
  Store("Foo", A_AD);  
  Update(A_AD );  
  Update(B_AD );  
  Update(C_AD );  
Commit_Transaction;
```



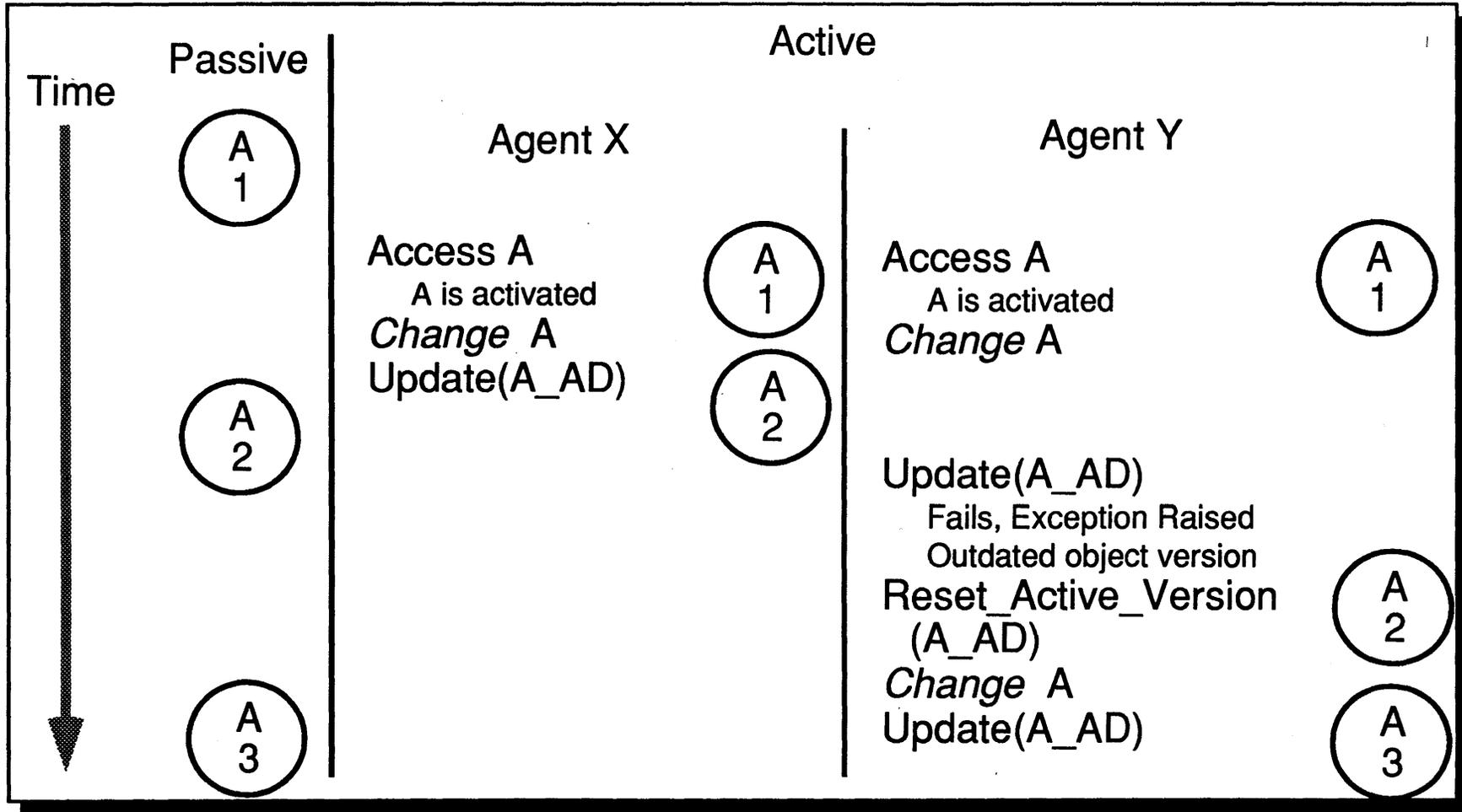
Transaction Concepts



Transaction Service Embedded in BiiN/OS

- Transaction Service Acts as Coordinator
- Multiple Services Can Participate in Same Transaction
 - Files
 - Directories
 - Libraries
- Extendable to New Services
- Distributed
 - using 2-phase commit protocol

Distributed Optimistic Concurrency



Distributed Pessimistic Concurrency

Start_Transaction;

Reserve(A_AD);

Synchronizes. Does Reset, if necessary.

Change A

Update(A_AD);

Commit_Transaction;

Persistent Object Summary

- Supports a Permanent Network of Distributed Typed Objects
 - Network can cross disk and node boundaries
 - Accessible independent of location
- Supports Concurrent Distributed Access
 - Based on transactions for synchronization and data integrity
 - Both optimistic and pessimistic synchronization are supported
- Activation is Implicit for Ease-of-use
- Passivation is Explicit for Data Integrity

Summary

