Burroughs

B 6700 / B 7700

NDL

LANGUAGE REFERENCE MANUAL

(RELATIVE TO MARK II.6 RELEASE)



\$7.00

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Page No.	Issue	Page No.	Issue
Title	Original	6-1 thru 6-19	Original
A	Original	6-20 blank	Original
i	Original	A-1 thru A-4	Original
ii blank	Original	B-1 thru B-5	Original
iii thru vii	Original	B-6 blank	Original
viii blank	Original	C-1 thru C-2	Original
1-1 thru 1-10	Original	D-1 thru D-7	Original
2-1 thru 2-3	Original	D-8 blank	Original
2-4 blank	Original	E-1 thru E-5	Original
3-1 thru 3-16	Original	E-6 blank	Original
4-1 thru 4-2	Original	Index-1 thru Index-11	Original
5-1 thru 5-187 .	Original	Index-12 blank	Original
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#### **PREFACE**

This document provides reference data for the experienced programmer who is familiar with the B 6700/B 7700 Network Definition Language (NDL) and the B 6700/B 7700 Data Communications System.

This reference manual is divided into the following six chapters and five appendixes.

- Chapter 1, INTRODUCTION, describes where the NDL Manual fits into the existing Data Communications System documentation, and defines the scope of NDL.
- Chapter 2, NDL SYNTAX CONVENTIONS, explains the syntactical notation used in defining the Network Definition Language.
- Chapter 3, LANGUAGE COMPONENTS, describes the elements that form the most primitive structures of the language.
- Chapter 4, SOURCE PROGRAM STRUCTURE, describes the basic structure of an NDL program.
- Chapter 5, DEFINITIONS, describes the various definitions that make up an NDL program.
- Chapter 6, VARIABLES, describes the program variables available to the NDL programmer.
- Appendix A, RESERVED WORDS, is a list of "words" that have been set aside for specific purposes within the Network Definition Language.
- Appendix B, TRANSMISSION CODES, provides useful data transmission code tables.
- Appendix C, SOURCE INPUT FORMAT AND CODING FORM, describes the input format and coding form to be used by the programmer.
- Appendix D, COMPILE-TIME OPTIONS, describes the compiler options available to the user.
- Appendix E, COMPILER SOURCE AND OBJECT FILES, describes how compiler communication is handled through various input and output files.

The information in the following documents pertains to and supplements the material presented in this reference manual:

<u>Title</u>	Form No.
B 6700/B 7700 Data Communications Functional Description	5000060
B 6700/B 7700 DCALGOL Reference Manual	5000052
B 6700 Input/Output Subsystem Information Manual	5000185

### TABLE OF CONTENTS

Chapter	이 보이 하는 이 이번 보는 사람들이 하는 사람들이 살아 하는 것은 것이 되었다.	Page
	PREFACE	i
1 .	INTRODUCTION	1-1
	General Scope of NDL Use of NDL	1-1 1-1 1-1
	Message Control System (MCS) Data Comm Controller (DCC) NDL Compiler DCP Interaction With the Main System	1-5 1-5 1-5 1-7
	DCP Tables	1-7
2	NDL SYNTAX CONVENTIONS	2-1
	Syntax Conventions Key Words Syntactic Variables Construct Terminator	2-1 2-2 2-2 2-2
3	LANGUAGE COMPONENTS	3-1
	Language Components Character Identifier	3-1 3-2 3-7
	Integer Label Remark Space	3-8 3-9 3-10 3-11
	String System Identifier TALLY Number Time	3-12 3-13 3-14 3-15 3-16
4	Toggle Number  SOURCE PROGRAM STRUCTURE	3-10 4-1
	NDL Program Unit	4-1
5	DEFINITIONS	5-1
	CONSTANT DEFINITION CONTROL DEFINITION Assignment Statement BREAK Statement CODE Statement	5-2 5-4 5-6 5-8 5-9
	Compound Statement CONTINUE Statement DELAY Statement ERROR Switch Statement	5-10 5-11 5-12 5-13

## TABLE OF CONTENTS (Cont)

Chapter		Page
5	<b>DEFINITIONS</b> (Cont)	
J		
	FINISH Statement	5-16
٠	FORK Statement	5-17
	GO TO Statement	5-18
		5-20
	IF Statement	5-21
•	INCREMENT Statement	5-23
	INITIALIZE Statement	5-24
	INITIATE Statement	5-25
	PAUSE Statement	5-28
	RECEIVE Statement	5-29
	SHIFT Statement	5-39
	SUM Statement	5-40
1	TRANSMIT Statement	5-42
	WAIT Statement	5-44
	DCP DEFINITION	5-45
	DCP EXCHANGE Statement	5-46
* *	DCP MEMORY Size Statement	5-51
	DCP TERMINAL Statement	5-52
	FILE DEFINITION	5-56
	FILE FAMILY Statement	557
	LINE DEFINITION	5 58
	LINE ADAPTER Class Statement	560
	LINE ADDRESS Statement	5-62
	LINE ANSWER Statement	5-63
	LINE DEFAULT Statement	5-64
	LINE ENDOFNUMBER Statement	5-65
	LINE MAXSTATIONS Statement	5-66
	LINE MODEM Statement	5-67
	LINE PHONE Statement	5-68
	LINE STATION Statement	5-69
	LINE TYPE Statement	5-70
	MCS DEFINITION	5-73
	MODEM DEFINITION	5-74
	MODEM ADAPTER Statement	5-75
*	MODEM LOSSOFCARRIER Statement	5-78
	MODEM NOISEDELAY Statement	5-79
	MODEM TRANSMITDELAY Statement	5-80
	REQUEST DEFINITION	5-81
141	Assignment Statement	5-83
	BACKSPACE Statement	5-85
	BREAK Statement	5-86
	CODE Statement	5-87
,	Compound Statement	5-88
	CONTINUE Statement	5-89
**	DELAY Statement	5-90
		E 01

## **TABLE OF CONTENTS (Cont)**

Chapter		Page
5	DEFINITIONS (Cont)	
	FETCH Statement	
	FINISH Statement	
	FORK Statement	
	GETSPACE Statement	
	GO TO Statement	
	IF Statement	
.*		5-102
		5-104
	INITIATE Statement	5-106
	PAUSE Statement	5-108
	RECEIVE Statement	5-109
	SHIFT Statement	5 - 121
	STORE Statement	5 - 122
	SUM Statement	5 - 124
	TERMINATE Statement	5 - 126
	TRANSMIT Statement	5-130
	WAIT Statement	5-133
	STATION DEFINITION	5-134
	STATION ADAPTER Statement	5 - 136
	STATION ADDRESS Statement	5 - 138
	STATION CONTROL Character Statement	5 - 139
	STATION DEFAULT Statement	5 - 140
	STATION ENABLEINPUT Statement	5 - 141
	STATION FREQUENCY Statement	5 - 142
	STATION INITIALIZE Statement	5 - 143
•	STATION LOGICALACK Statement	5-144
	STATION MCS Statement	5-145
	STATION MODEM Statement	5-146
	STATION MYUSE Statement	5 - 147
	STATION PHONE Statement	5 - 148
	STATION PHONE Statement	5 - 149
	STATION RETRY Statement	5 - 150
	STATION WIDTH Statement	5-151
	STATION WIDTH Statement	5 - 152
•	TERMINAL DEFINITION TERMINAL ADAPTER Statement	5-153
	TERMINAL ADDRESS Size Statement	5-156
	TERMINAL BACKSPACE Character Statement	5-157
•	TERMINAL BUFFER Size Statement	5-158
	TERMINAL CARRIAGE Character Statement	5-159
	TEDMINAL CI EAD Character Statement	5-160
	TERMINAL CLEAR Character Statement TERMINAL CODE Statement	5-161
	TERMINAL CONTROL Statement	5-162
	TERMINAL DEFAULT Statement	5-163
	TERMINAL DUPLEX Statement	5-164
	TERMINAL END Character Statement	5-166
	TERMINAL HOME Character Statement	5-167
4	I DIVINAL HOME Character Statement	5-168

## TABLE OF CONTENTS (Cont)

Chapter		Page
5	DEFINITIONS (Cont)	
	TERMINAL Illegal Character Statement	5-169
	TERMINAL INHIBITSYNC Statement	5-170
	TERMINAL Inter-Character Delay Statement	5-170
	TERMINAL LINEDELETE Character Statement	
	TERMINAL LINEFEED Character Statement	5-173
	TERMINAL MAXINPUT Statement	5-174
	TERMINAL PAGE Statement	5-175
	TERMINAL PARITY Statement	5-176
	TERMINAL REQUEST Statement	5-177
	TERMINAL SCREEN Statement	5-178
	TERMINAL TIMEOUT Statement	5-179
	TERMINAL TRANSMISSION Number Length Statement	5-180
	TERMINAL TURNAROUND Statement	5-181
	TERMINAL WIDTH Statement	5-182
	TERMINAL WRU Character Statement	5-183
	TRANSLATETABLE DEFINITION	5-184
6	VARIABLES	6-1
U		
	General	6-1
	Function of Variables	6-1
	Scope of Variables	6-2
	Scope of variables	
•	Description of Variables	6–2
A		D
Append		Page
, <b>A</b>		
A	Reserved Words	A-1
B	Transmission Codes	B-1
<b>C</b> .	Source Input Format and Coding Form	C-1
D	Compile-Time Options	D-1
E	Compiler Source and Object Files	E-1
	Index	Index-1
• •,		
	LIST OF ILLUSTRATIONS	
Figure	Title Title	Page
<u>-</u> .		
1-1	B 6700/B 7700 Data Communications System Documentation Hierarchy	1-2
1-2	Network Characteristics	1-3
1-3	NDL Sphere of Influence	1-4
1-4	Transfer of Control Within the DCP	1-7
5-1	Adapter Clusters Exchange	5-50
D-1	Option Control Card	D-4
E-1	NDL Compilation System	E-2

## LIST OF TABLES

Table	Title		Page
5-1	Relational Operators		 5-22
5-2	Allowable Combinations for \( \text{receive statement} \)		
5-3	Available Line Adapters		
5-4	Table of \(\langle communication \) type number\\\ s \qquad		
5-5	Relational Operators		
5-6	Allowable Combinations for \( \textit{receive statement} \)	)	 5-118
6-1	Table of Variables		
E-1	NDL Compiler Files		

#### 1. INTRODUCTION

#### GENERAL

This document is one of several documents concerning the B 6700/B 7700 Data Communications System. The hierarchy of these documents is illustrated in figure 1-1. Note that information contained in the B 6700/B 7700 Data Communications Functional Description, form no. 5000060, and B 6700 Input/Output Subsystem Information Manual, form no. 5000185, is prerequisite to this document.

#### SCOPE OF NDL

The Network Definition Language (NDL) source program describes a data communications network physically, logically, and functionally. Physical components of a data communications network include the hardware specifications and capabilities of the various elements which comprise the network. The logical characteristics of a network are the associations among the various components of a data communications subsystem (user programs, Message Control Systems, etc.), application-oriented characteristics (page size and width, special-purpose characters, etc.), and the symbolic names used to reference physical elements within the network. Figure 1-2 illustrates the physical and logical characteristics in their relation to the network elements. Figure 1-3 illustrates the sphere of influence of a NDL source program on the logical and physical components of a Data Communications System. The area enclosed by the broken line is the "global sphere of influence." The shaded area indicates the "local sphere of influence." The arrows indicate the flow of information.

NDL also specifies the functional behavior of the network or the way in which each data communications line is to be controlled. These specifications consist of individual routines, allowing the NDL programmer to implement the protocol required to meet the physical characteristics and application of the types of terminals that have been defined. The routines are compiled into a set of instructions which the Data Communications Processor (DCP) executes to perform the functions described by the NDL program.

The NDL source program is transformed into two files containing the information required to operate the defined network:

- a. The Network Information File (NIF), containing the logical and physical specifications of the network.
- b. The DCP Code File (DCPCODE), containing the Data Communication Processor (DCP) hardware instructions for operating the network.

#### **USE OF NDL**

Once the data communications hardware has been installed on a B 6700/B 7700 system, several software systems are required to generate and operate the data communications network. These packages, illustrated in figure 1-3, consist of one or more Message Control Systems (MCS), the Data Comm Controller (DCC), and the NDL compiler. The purpose, function, and use of each of these software items are described in the following paragraphs.

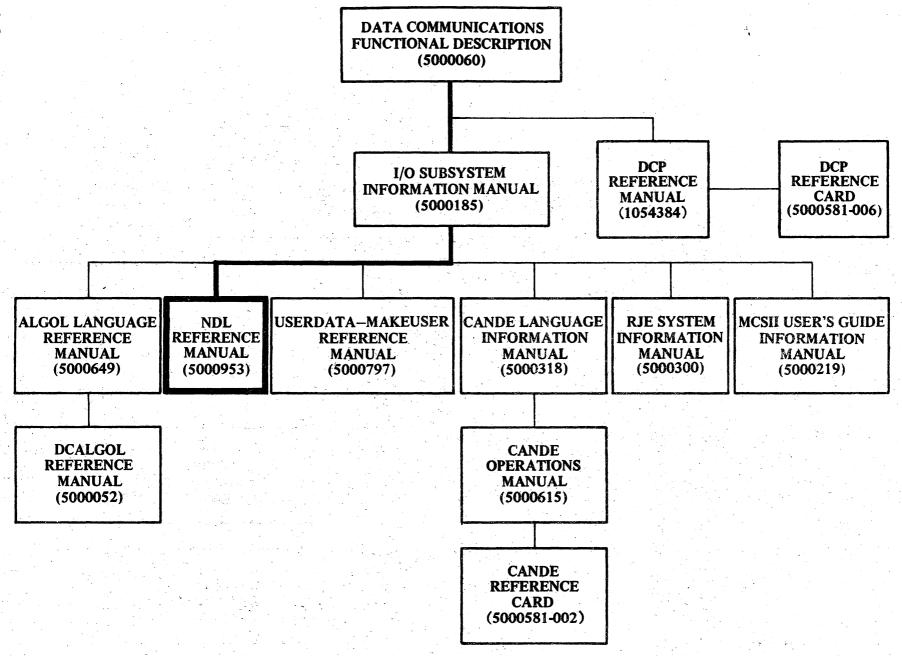


Figure 1-1. B 6700/B 7700 Data Communications System Documentation Hierarchy

Network Elements	Physical Characteristics	Logical Characteristics
DCPs	Memory size Reconfiguration capabilities	Set of terminals controlled by each DCP
LINEs	Physical location (address) Transmission speed Type of line and connection	Station line assignments Automatic answer capability
MODEMs	Physical delays Transmission speed and type Continuous vs. controlled carrier	Symbolic name
STATIONs	Terminal characteristics	Symbolic name Logical attributes Associated Message Control System
TERMINALs	Transmission code, speed and type Parity	Transmission numbers Special characters

Figure 1-2. Network Characteristics

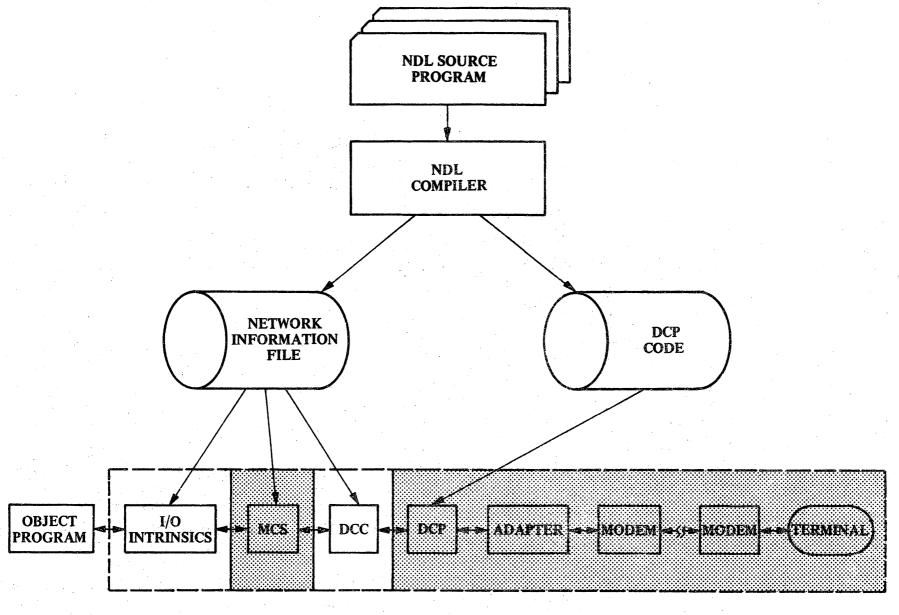


Figure 1-3. NDL Sphere of Influence

#### **Message Control System (MCS)**

An MCS is a special purpose DCALGOL program which may be a Burroughs-supplied MCS (SYSTEM/CANDE, SYSTEM/RJE, or SYSTEM/DIAGNOSTICMCS), or a user-written program. The primary function of an MCS is, as its name implies, to control the flow of data communications messages between the terminal and the main system. Information from the DCP, such as terminal input and status information, is forwarded to the MCS via the DCC. Messages from the MCS to the DCP, such as terminal output or network changes, are performed by the MCS invoking an intrinsic function called DCWRITE. This intrinsic, as well as the format of all MCS and DCP messages, is described in the B 6700/B 7700 DCALGOL Reference Manual, form no. 5000052. Each station which is defined by the NDL source program must have one, and only one, controlling MCS.

#### Data Comm Controller (DCC)

The DCC is the basic interface between the DCP and the main system. It exists as a subset of the basic B 6700/B 7700 Master Control Program (MCP) and operates as an independent task or stack, one such task for each active DCP.

Before a defined data communications network may be utilized, the DCPs which comprise the network must be initialized. As each DCP is initialized, the portion of the NDL-defined network which utilizes that DCP becomes active.

Once initialized, each DCC stack transfers messages between the associated DCP and the proper MCS.

#### **NDL** Compiler

Whereas the DCC and an MCS are required to operate a data communications network, the NDL compiler is used to generate the tables and DCP code which, to a large extent, control the way in which the network functions.

The NDL source program, then, must supply the NDL compiler with information which will allow the compiler to produce the proper NIF and DCPCODE files to operate all of the Data Comm Processors and their sub-components within the network. (In figure 1-3, the shaded areas indicate the areas of the data communications subsystem which are influenced by the NDL source program.) Although an NDL program may contain up to 11 discrete sections, it functionally consists of two interdependent pieces of information: the network description and the DCP programs.

#### **NETWORK DESCRIPTION**

The NDL programmer uses various sections of the NDL source program to describe the logical and physical characteristics of the network. The information supplied in those sections is used, in part, to supply the DCC with the proper tables and DCP code that are used to operate the network.

The NDL compiler performs consistency checks across the various definitions to ensure that the defined network is logically structured. For example, a line must not be associated with a particular modem if the defined speed range and transmission type of the modem do not permit a proper interface to the line. Similarly, a terminal defined to operate in an asynchronous mode must not be associated with a line which uses a synchronous adapter.

All of the information supplied by the NDL definitions is recorded within the NIF file. This enables an MCS or user program to gain access to many of the logical characteristics of the network, as well as permitting dynamic reconfiguration of the network by an MCS.

These definitions are also used to modify or include special areas of DCP code which are network dependent. For example, the DCP code for transmitting or receiving characters on a synchronous line is different from such code for transmitting or receiving on an asynchronous line. In addition, if any dial-out type lines are defined within the network, extra code must be generated for performing dial-out functions. Thus, the NDL compiler "tailors" the resultant DCPCODE file to fit all the requirements of the defined network.

#### DCP PROGRAMS

Once a data communications network is logically and physically defined, the functional operation of each line and station within the network must be described. These descriptions, called **CONTROL** definitions and **REQUEST** definitions, are individual programs which are executed by the DCP when required to perform the necessary line discipline. Each line must have one associated **CONTROL** definition, and each terminal may have one or more associated **REQUEST** definitions.

The CONTROL and REQUEST definitions consist of NDL statements which the compiler transforms into the DCP instructions to be executed when performing a particular network function. A RECEIVE REQUEST definition is invoked when input from a terminal is to be processed, and a TRANSMIT REQUEST definition is executed when output to a terminal has been requested by an MCS or user program. The line CONTROL definition is utilized to determine when and for which of the stations on the line a REQUEST definition is to be executed.

Since an NDL source program must handle many lines, the DCP must share its processing capabilities among the lines it services. Due to the fact that the data communications subsystem operates in a real-time environment, few network functions, if any, require the dedicated use of the DCP for an extended length of time. A RECEIVE REQUEST, for example, usually spends most of the time waiting for a character to be sent from a terminal. Likewise, a TRANSMIT REQUEST can only operate as fast as the line speed permits. Thus, while a REQUEST definition is waiting for an external event, or interrupt, from a line, the DCP is free to continue execution of a REQUEST or CONTROL definition for another line.

The allocation of the DCP for the servicing of its many lines is one of the duties of the basic DCP operating system and the CONTROL definitions. Figure 1-4 illustrates the means by which the control of the DCP is transferred between the operating system and the CONTROL and REQUEST definitions.

#### Line Control

Each CONTROL definition, or "Line Control Procedure," must perform two functions. First, it must select which station on the line is to receive attention next, and second, it must decide what particular function is to be performed for that station. If the function to be performed is an output request, control is transferred to the TRANSMIT REQUEST for the station. If the function is an input operation, the station's RECEIVE REQUEST is executed. Network functions which do not involve the reception or transmission of messages, such as status or network changes, are performed by invoking a common subprogram, or macro, within the DCP operating system itself.

#### Request Definitions

For each type of terminal which is capable of output, a **TRANSMIT REQUEST** must be named within the terminal definition. Likewise, if a terminal has input capabilities, a **RECEIVE REQUEST** must be supplied and named. Typically, many stations may share the same **REQUEST** definitions, just as many lines may utilize the same **CONTROL** definition. In some cases, more than one set of **REQUEST** definitions may be desired, and defined, for a station.

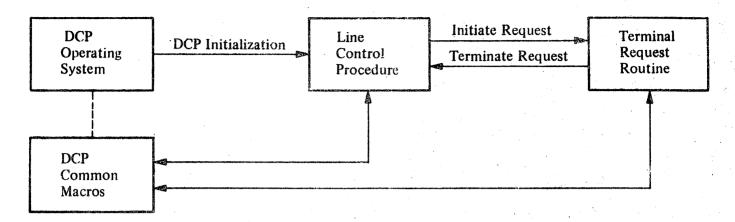


Figure 1-4. Transfer of Control Within the DCP

The functions of a REQUEST definition may be as simple or as complex as the application of the station dictates. One of the basic design goals of the DCP is to free the main system from the burden of performing basic terminal receptions and transmissions. However, by the proper application and coding of the CONTROL and REQUEST statements in NDL, a significant amount of intelligent message processing may be performed by the DCP, thereby allowing more of the main system's resources to be free to perform other work.

The NDL programmer must keep in mind, however, that the DCP runs at a finite rate, and that it is operating in real-time. Thus, if too much time is spent processing a message, other lines may fail to be serviced quickly enough to avoid transmission errors. Several NDL statements are provided to "break up" long strings of NDL code to ensure that the DCP may properly service all of its lines.

When a **REQUEST** definition has terminated the processing of an input or output function, it usually branches back to the beginning of the **CONTROL** definition. The **CONTROL** definition then selects the next station to be serviced and the process continues.

Thus, the functioning DCP can be visualized as a small multiprogramming system, where each line has its own program and operating environment and runs asynchronously and independently of the other lines. The **CONTROL** definition and its associated **REQUEST** definitions form the "main program" for each line, and the common DCP macros are "sub-programs."

#### DCP INTERACTION WITH THE MAIN SYSTEM

Although the DCP is a self-contained and asynchronous device with respect to the main system, it is not an autonomous unit, and requires the active participation of the main system and its resources to properly function. In particular, main memory storage space is required to contain tables and messages. In addition, the DCP requires the allocation of a pool of message areas in main memory for the gathering of input from terminals and reporting of error conditions.

#### **DCP Tables**

The NDL compiler constructs a series of tables which reflect the physical and logical characteristics of the network as defined by the NDL source program. The DCP uses these tables for the storage of status information, and for determining what types of functions are to be performed for each of the many different lines and stations which the DCP controls.

The compiler places a disk image of these tables within the DCPCODE file along with the DCP code itself. When the DCP is initialized, the tables are loaded into main memory by the DCC, which also provides the

DCP with a reference to the tables. If several DCPs exist which share hardware-exchanged adapter clusters, two DCPs may utilize the same set of tables if the network description indicates this mode of operation. In the case where a DCP is not "exchanged" in this manner, each DCP uses its own unique set of tables.

Each set of tables can be divided into two sets of information. Each line has a table, and each station has a table. The DCP uses a "line descriptor" to reference each line table. The descriptors for all lines controlled by the DCP are stored within a vector, which is then indexed by the physical line adapter address. Each line descriptor contains information concerning the status of the line (not ready, connected, busy, etc.), physical characteristics of the line (dial-out, switched, etc.), logical characteristics (automatic answer, etc.), and a reference to the CONTROL definition which is used for the line. In addition, the line descriptor contains the memory address of the actual line table.

Each line table contains additional information describing the logical and physical characteristics of the line. Much of the information in the line table can be referenced and/or modified directly by the NDL CONTROL and REQUEST definitions. Other information is reserved for use by the DCP operating system.

Immediately following the line information in the line table is a vector of station descriptors, one such descriptor for each station which can exist on the line. Similar to the line descriptor, each station descriptor addresses a table of information for a particular station. The DCP references the proper station table by indexing into the line table by the proper relative station address, or "station index," and using the addressed station descriptor to reference the proper station table. When a line CONTROL Procedure "selects" a station for the purposes of initiating a REQUEST definition, it is actually selecting the proper station table for use. Just as the line table contains a reference to the proper CONTROL definition to execute, the station table contains references to the appropriate TRANSMIT and RECEIVE REQUEST definitions. It should now be apparent that each station assigned to a particular line must utilize the same Line Control Procedure, since the stations on the line all share a common line descriptor and line table. However, each station may have a different set of REQUEST definitions, since these routines are station oriented.

Although each station table is of a fixed size, the line tables will vary in size directly proportional to the number of stations which can potentially exist on the line. The NDL source program specifically defines the maximum number of stations for each line, as well as which stations are assigned to what line. Not all station descriptors may be utilized for a given line, i.e., the number of real stations on a line at any given moment may be less than the true capacity of the line. A line may be declared with such "holes" when the NDL program is compiled, or a line may be reconfigured into such a state by an MCS. In some cases, a line may exist with no stations at all. At no time, however, may more stations be assigned to a line than the maximum number defined by the NDL program. Thus, it is the requirement of the DCP operating system and/or the CONTROL definitions to ensure that a selected station actually exists, or is valid, as defined by the current state of the network.

Just as a line may have no valid stations, it is possible and often desirable to define "spare" stations which have no line assignment. Such stations cannot be referenced or utilized by the DCP until they are logically assigned to a line by a reconfiguration request. Again, any such reconfiguration request will be disallowed by the DCC if the characteristics of the station conflict with those of the line to which it is being assigned, or with the stations which already exist on that line. Also, since the size of the line table cannot be altered, there must exist a "hole" or unused position on the line for the station.

Alternatively, an existing station may be subtracted from a line, thereby leaving a "hole," and either left in limbo with no line assignment, or moved to fill an existing "hole" on another line. Thus, stations which have special characteristics for a particular application may be logically moved about within the network while the data communications system is operating and without the further use of the NDL compiler.

#### DCP MESSAGE MAINTENANCE

With one exception, all functions performed by the DCP are the direct result of a DCP request message being sent to the DCP, usually by an MCS. In the case of terminal output, for example, a "write request" is sent to the DCP, which then invokes the action described within the request message itself by means of the appropriate station's TRANSMIT REQUEST definition. If spontaneous input from a terminal is to be received, there is normally no MCS request message associated with the input operation. When a station operates in this mode, the terminal is described as being "enabled for input," or simply, "enabled."

The process of gathering "enabled input messages," i.e., spontaneous input messages, is controlled by the **CONTROL** definition, and, of course, by the **RECEIVE REQUEST** defined for the terminal. In addition, the "enabled" state of a station is initially defined for each station, and may be dynamically changed by the controlling MCS.

When a station is enabled, and the RECEIVE REQUEST is invoked, the DCP must then acquire a message area in main memory in which to store the received message text. Such an area, which is called an "enable input space," is obtained by a DCP macro called GETSPACE. Since the DCP cannot directly participate in the main system memory management functions, a pool of such "enable input" spaces is maintained by the DCC. This pool of messages, sometimes referred to as the "available space pool," consists of a set of queues, each of which contains a linked list of available message areas of the same size. The NDL compiler computes the size of the enable input space required for each terminal based on the defined WIDTH, MAXINPUT, and BUFFERSIZE statements within the terminal definition. All terminals of a given size are assigned the use of the same queue. In order to reduce the number of different queues required, the NDL compiler rounds each terminal's input size up to a multiple of 16 words.

The available space areas are used for several purposes other than terminal enabled input. Error messages from the DCP and "switched line status" result messages are also spontaneous in nature and require an enable input space. In addition, it is possible for an NDL Request definition to invoke the GETSPACE macro and simply store the contents of variables in the obtained message space in order to communicate with the controlling MCS.

When the DCP GETSPACE macro is invoked, an area which is greater than or equal to the required message size is delinked from a message queue and assigned to the station. If no suitably sized areas exist within the space pool, a "no space" condition results, and the RECEIVE REQUEST must about reception of input.

The number of messages assigned to each queue is initially defined by the NDL compiler. By default, two areas are assigned to each size queue, although the NDL program may specify an alternate allotment on a terminal by terminal basis.

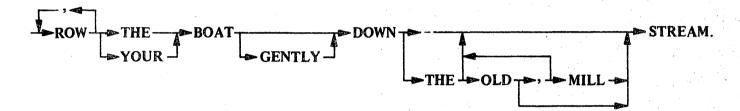
The DCC has the responsibility of maintaining the available space pool so that GETSPACE may always obtain a message area. As each available space area is returned to the DCC by the DCP in the normal course of completing an input operation, the queue from which the area was obtained is restored so that it contains the same number of areas as defined by the NDL compiler. Circumstances may arise, however, where all of the areas within a queue have been exhausted, but none of the areas has yet been returned to the DCC so that the queue can be replenished. In such an event, the DCP sets a global "space alarm" flag which is sensed by the DCC and causes it to immediately examine and replenish all of the available space queues. In addition, the DCC will then increase the target number of messages in each totally depleted queue, in order to reduce the possibility of future space alarms. During extended periods of DCP inactivity, the DCC will attempt to reduce the number of messages in each queue down toward the originally defined target value.

The DCC attempts to maintain the available space pool within the constraints specified by the NDL compiler. However, some networks may require more than the default number of message areas for some terminals if too many "no space" conditions occur. In such an event, the NDL program should specify a larger number of message areas for the affected terminals. Since the behavior of a network is difficult to predict under all circumstances, the NDL programmer will have to directly observe the effects of different message space specifications, and adjust the specifications so that the network operates efficiently without requiring excessive memory resources.

#### 2. NDL SYNTAX CONVENTIONS

#### SYNTAX CONVENTIONS

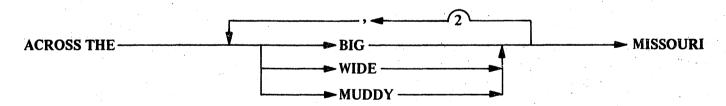
The syntax diagram is the method used to depict the Network Definition Language syntax. This method affords a very concise and lucid exposition of syntax, including defaults, alternatives, and iterations; it is rigorous without being cumbersome. There are few formal rules to remember: the basic rule is that any path traced along the forward directions of the arrows produces a syntactically valid command. The following examples illustrate the technique:



Valid constructs from this syntax diagram include:

ROW THE BOAT DOWN-STREAM.
ROW, ROW, ROW YOUR BOAT GENTLY DOWN THE STREAM.
ROW, ROW, ROW, ROW THE BOAT DOWN THE OLD STREAM.
ROW YOUR BOAT DOWN THE OLD, MILL STREAM.
ROW THE BOAT DOWN THE OLD, MILL STREAM.

The following convention is used to control iterations of options or constructs:



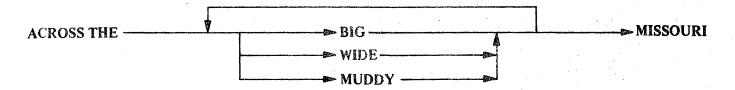
The "bridge" over the "2" can be crossed a maximum of two times, so a maximum of two commas (and three adjectives) can appear. Valid productions include:

ACROSS THE BIG MISSOURI ACROSS THE BIG, WIDE MISSOURI ACROSS THE BIG, WIDE, MUDDY MISSOURI

From the above example, it should be noted that the number of iterations is controlled by the "number" in the feed-back loop. When the "number" is not shown, there is no limit to the iterations. For example,

#### SYNTAX CONVENTIONS

Continued



would include the following valid combinations:

ACROSS THE BIG BIG WIDE WIDE MISSOURI
ACROSS THE BIG MISSOURI
ACROSS THE MUDDY MUDDY MUDDY MUDDY MISSOURI

If a comma were included in the above example, valid combinations would be as follows:

ACROSS THE BIG, BIG, MUDDY MISSOURI ACROSS THE BIG, WIDE, WIDE, MUDDY MISSOURI

#### **Key Words**

Boldface symbols and uppercase letters in syntax diagrams indicate symbols and words which appear literally in the instruction.

#### Syntactic Variables

In the syntax diagrams, left and right broken brackets ( $\langle \rangle$ ) are used to contain syntactic variables that represent information to be supplied by the programmer. A particular variable may represent a single character, a simple construct (such as an integer or text string), or a relatively complicated construct.

The following is an example of a syntactic variable that appears in a syntax diagram.

Braces ({ }) are used to enclose syntactic variable expressions defined by the meaning of the English language expression contained within the braces. For example, the following syntactic variable expression

ADAPTER  $\longrightarrow$  =  $\longrightarrow$  {\langle integer \rangle value of 1 thru 8}

would include the following valid constructs:

ADAPTER = 1.

ADAPTER = 6.

ADAPTER = 8.

#### **Construct Terminator**

Most constructs in the Network Definition Language must be terminated by a period (.). This is illustrated in the syntax diagrams as follows:

The period is part of the syntax and must appear following the construct.

# NDL Syntax Conventions SYNTAX CONVENTIONS Continued

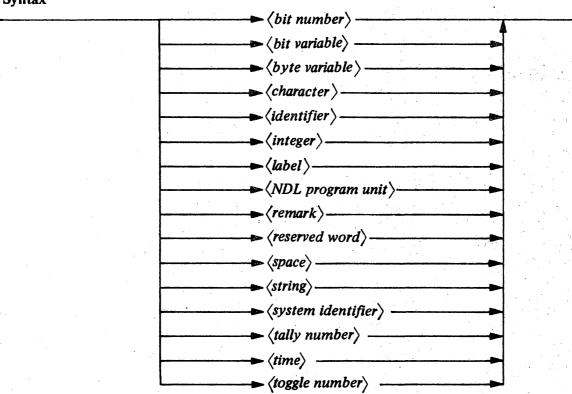
Some constructs, however, do not require a terminator, and can be followed by another construct. This is illustrated as follows:

The vertical bar ( | ) is not part of the syntax, but merely indicates the termination of the construct.

#### 3. LANGUAGE COMPONENTS

#### LANGUAGE COMPONENTS





#### **Examples**

A 450 6110 IF "B6700" SYSTEM/CANDE 30 MILLI

#### **Semantics**

(bit variable)s, (bit number)s, and (byte variable)s are all described in chapter 6.

A complete list of the \(\langle reserved \) word\(\rangle s\) is contained in appendix A.

(NDL program unit) is described in chapter 4.

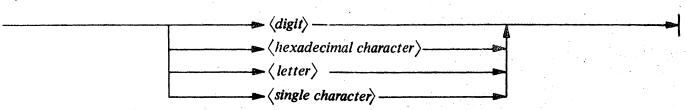
All other  $\langle language\ component \rangle s$  are described in this chapter.

#### Language Components

#### **CHARACTER**

#### **CHARACTER**

**Syntax** 



#### **Examples**

O Q A "Z"

#### **Semantics**

In all instances, a *(character)* is an entity whose exact form depends on the context of its usage. The normal inference is that of an 8-bit EBCDIC character.

 $\langle letter \rangle s$  and  $\langle digit \rangle s$  are usually used to create  $\langle identifier \rangle s$  and  $\langle string \rangle s$ .

Wherever \( \langle character \rangle \) occurs in the NDL syntax, an 8-bit character is needed. It is unique in that it may be formed using two \( \langle \) hexadecimal character \( \rangle s \).

# Language Components CHARACTER (algit)

n	1	п	r

#### **Syntax**

{one of the EBCDIC characters, 0 through 9, inclusive}

#### Examples

0

5

9

#### **Semantics**

Whenever the item of  $\langle digit \rangle$  appears in the NDL syntax, one of the 10 numeric EBCDIC characters from 0 through 9 is required.

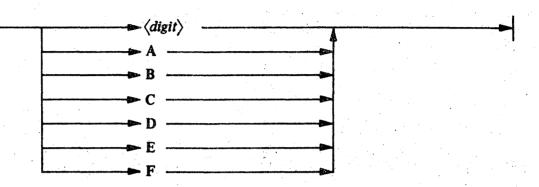
#### Language Components

#### **CHARACTER**

\(\lambda\) hexadecimal character\(\rangle\)

#### HEXADECIMAL CHARACTER

**Syntax** 



#### **Examples**

0

5 9

A

C F

#### **Semantics**

(hexadecimal character) s are defined as consisting of the characters in the decimal digit set plus the characters A, B, C, D, E, F. (hexadecimal character) s are generally used to define program values in terms of the hexadecimal (radix 16) number system, where A is equivalent to 10 in the decimal system, B is equivalent to 11 in the decimal system, etc.

Language Components				
	CHARACTER			
: '	(letter)			

				1
*	4	_	-	n
			м	~
	4 .			88

#### **Syntax**

one of the EBCDIC characters, A through Z, inclusive

#### **Examples**

A

Q Z

#### Semantics

Whenever the item of (letter) appears in the NDL syntax, one of the 26 alphabetic EBCDIC characters from A through Z is required.

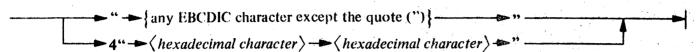
#### Language Components

#### **CHARACTER**

(single character)

#### SINGLE CHARACTER

#### **Syntax**



#### **Examples**

"A" 4"FF"

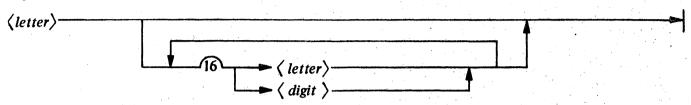
#### **Semantics**

The primary purpose of having a syntactical item of (single character) in NDL is for use in those places of syntax requiring an 8-bit character, which can be any combination of bits from "all off" through "all on."

For ease of programming and recognition of usage, the NDL programmer may use either normal EBCDIC graphic characters or \( \lambda exadecimal \) character\( \rangle s \) to create the needed bit pattern.

#### **IDENTIFIER**

#### **Syntax**



#### Examples

A

QV X3

B6700

**MINIMIZER** 

#### **Semantics**

 $\langle identifier \rangle$  s have no intrinsic meaning. They are used to give symbolic names to various definitions in NDL.

An  $\langle identifier \rangle$  must start with a  $\langle letter \rangle$ , which can be followed by any combination of  $\langle letter \rangle s$  and  $\langle digit \rangle s$ .

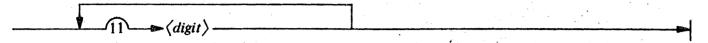
The maximum length of an (identifier) is 17 characters.

#### Language Components

#### INTEGER

#### **INTEGER**

Syntax



#### Examples

0 37

511

12345678901

#### **Semantics**

LABEL		
Syntax		
(integer) —		
Examples		

#### **Semantics**

22

9999999999

A  $\langle label \rangle$  is used to indicate where "execution can branch" within a given  $\langle control\ definition \rangle$  or  $\langle request\ definition \rangle$ .

 $\langle label \rangle$ s are "local" in scope; that is, each  $\langle label \rangle$  must be unique only within a given  $\langle control\ definition \rangle$  or  $\langle request\ definition \rangle$ . For example, the  $\langle label \rangle$  22 could appear more than once in an NDL program so long as it does not appear more than once in the same  $\langle control\ definition \rangle$  or  $\langle request\ definition \rangle$ .

Language Components	3
REMARK	

REMARK				
Syntax				
%	EBCDIC	characters}		-
Example				
% THIS IS A REMA	RK			•

#### **Semantics**

 $\langle remark \rangle$ s can appear anywhere in the source program. When the compiler encounters the percent sign (%), it skips immediately to the next source record before continuing the compilation process.

T EDCDIC graphic space, nexadecintal 40	-						4
1 EBCDIC graphic space, hexadecimal 40						en e	
Syntax		· ·					•
SPACE							

#### **Semantics**

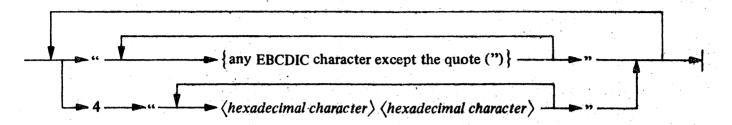
The NDL compiler looks at a contiguous sequence of  $\langle space \rangle s$  in a source program as a single  $\langle space \rangle$  (except when contained in a  $\langle string \rangle$ ). Therefore, wherever a single space is allowed, the programmer can use multiple  $\langle space \rangle s$  to improve readability of the program.

#### **Language Components**

#### **STRING**

#### **STRING**

#### **Syntax**



#### **Examples**

"THIS IS A STRING"

"AND" "SO" "IS" "THIS"

"AND SO" 4"C9E2" "THIS"

4"C2F6F7F0F0"

"%\*+?/@ (BIG B)"

#### **Semantics**

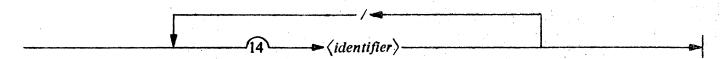
Only \(\langle hexadecimal character\rangle\) (4-bit) and EBCDIC character (8-bit) \(\langle string \rangle s\) can be constructed.

EBCDIC character (string) s are restricted in that they cannot contain the quote (") character.

The maximum allowed length of a (string) is 128 8-bit characters (1024 bits).

#### **SYSTEM IDENTIFIER**

#### **Syntax**



#### **Examples**

A B6700 SITE/MCS SYSTEM/RJE/DOWNTOWN X/Y/ZEBRA

#### **Semantics**

(system identifier)s have no intrinsic meaning. They are used to give symbolic names to various definitions in NDL. A (system identifier) is different from an (identifier) in that it is usually used to reference items belonging to the system and not simply to the NDL program.

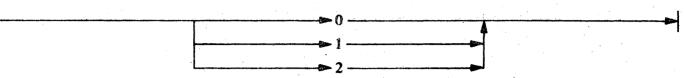
A maximum of 14 (identifier)s, each separated by a slash (/), is allowed.

#### Language Components

#### **TALLY NUMBER**

#### **TALLY NUMBER**

**Syntax** 



#### **Examples**

0 1 2

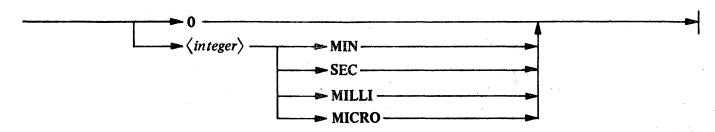
## Semantics

 $\langle tally\ number \rangle$  is required to designate one of the three  $\langle byte\ variable \rangle$  TALLYs; for example, TALLY[0].

TIME

# TIME

# **Syntax**



# **Examples**

0

5 MIN

**30 SEC** 

200 MILLI

9 MICRO

#### **Semantics**

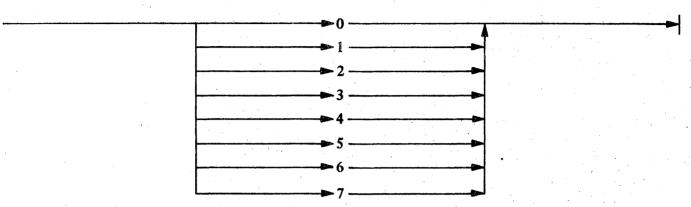
 $\langle time \rangle$  is used to express or define an increment of time. MIN denotes minutes, SEC denotes seconds, MILLI denotes milliseconds, and MICRO denotes microseconds.

The maximum amount of time that can be specified is 6 minutes 42 seconds.

# **TOGGLE NUMBER**

# **TOGGLE NUMBER**

**Syntax** 



# Examples

0 4 7

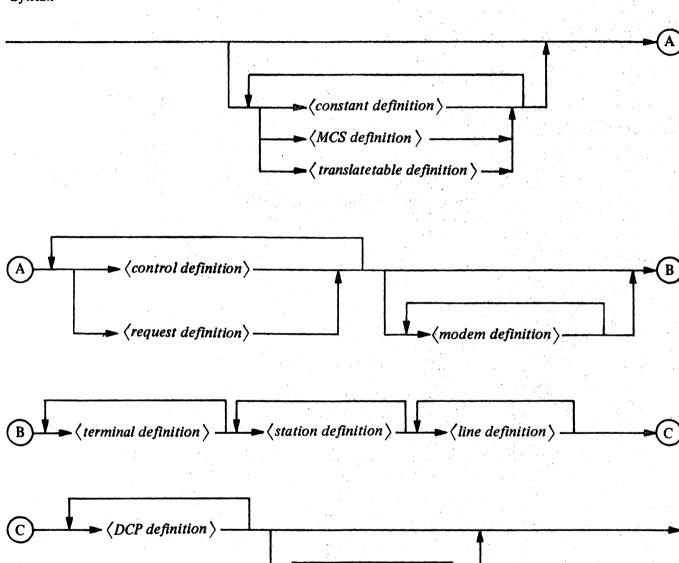
# Semantics

\(\langle toggle number \rangle\) is required to designate one of the eight \(\langle bit variable \rangle\) TOGs. For example, TOG[5].

# 4. SOURCE PROGRAM STRUCTURE

# NDL PROGRAM UNIT

# Syntax



 $\langle$  file definition $\rangle$ 

# Source Program Structure

#### **NDL PROGRAM UNIT**

Continued

# **Examples**

CONSTANT . . . MCS . . . TRANSLATETABLE . . . CONTROL . . . REQUEST . . . MODEM . . . TERMINAL . . . STATION . . . LINE . . . DCP . . . FILE . . .

#### **Semantics**

The NDL source program is divided into 11 program sections ordered as shown. An NDL program must include the control and request sections (in any order), and the terminal, station, line, and DCP sections. Each section is described in detail in chapter 5 of this manual.

# 5. DEFINITIONS

# **DEFINITIONS**

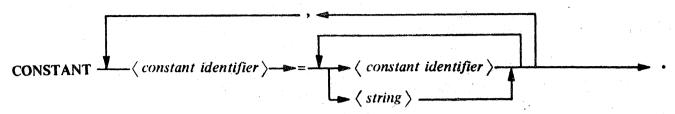
The NDL definitions, which comprise the 11 program sections of the source program structure shown in Section 4, are listed in alphabetical order and described in the same order:

- **CONSTANT** a.
- CONTROL b.
- **DCP** c.
- **FILE** d.
- LINE e.
- f. **MCS**
- **MODEM** g.
- REQUEST STATION h.
- i.
- **TERMINAL** j.
- k. **TRANSLATETABLE**

#### **CONSTANT**

#### **CONSTANT DEFINITION**

#### **Syntax**



#### **Examples**

```
CONSTANT NUL = 4"00".

CONSTANT SOH = 4"01", STX = 4"02".

CONSTANT C1 = SOH 4"00" STX, C2 = "123"4"F4".
```

#### **Semantics**

The  $\langle constant \ definition \rangle$  equates each of one or more  $\langle identifier \rangle s$  with a  $\langle string \rangle$ . Once that equation is made, any subsequent appearance of the  $\langle constant \ identifier \rangle$  is syntactically and semantically equivalent to the  $\langle string \rangle$ .

If a *constant identifier* appears after the equal sign, it must have been defined previously in the program.

# **Supplementary Examples**

Examples of Valid (constant definition)'s

Example 1

# CONSTANT GREETING = "WELCOME TO B 6700 TIME SHARING.".

This example equates the \( \string \) "WELCOME TO B 6700 TIME SHARING." to the \( \chiconstant identifier \) GREETING.

Example 2

# **CONSTANT**

```
CR = 4"0D", % A CARRIAGE RETURN
LF = 4"25", % A LINE FEED
CRANDLF = CR LF, % A CARRIAGE RETURN AND A LINE FEED
DELETELINE = "DELETED" CR LF. % THE STRING "DELETED",
% A CARRIAGE RETURN, AND
% A LINE FEED.
```

This example references other  $\langle constant \ identifiers \rangle$ s to define a  $\langle constant \ identifier \rangle$ .  $\langle String \rangle$ s and  $\langle constant \ identifier \rangle$ s may be interspersed to define a  $\langle constant \ identifier \rangle$ .

Examples of Invalid (constant definition)s

Example 1

#### **CONSTANT BADCNST = 4"123".**

The above  $\langle constant \ definition \rangle$  would cause a syntax error to be generated, because the  $\langle string \rangle$  is not properly formed. The length of the  $\langle string \rangle$  must be a multiple of eight bits.

# Example 2

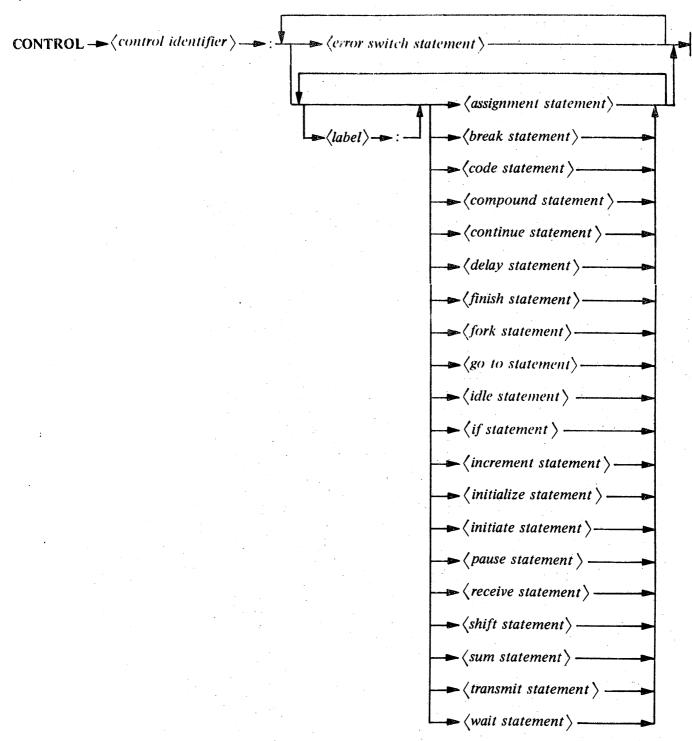
```
CRANDLF = CR LF, % LINE 1
CR = 4"0D", % LINE 2
LF = 4"25". % LINE 3
```

This example would cause a syntax error to be generated, because the *constant identifier* LF and CR are referred to in line 1, but not declared until line 2 and line 3. A *constant identifier* must be declared before any reference to that *constant identifier* can be made.

# **CONTROL**

# **CONTROL DEFINITION**

**Syntax** 



#### Examples

```
CONTROL CONTENTION:
    INITIATE REQUEST.
    INITIATE ENABLEINPUT.
    IDLE.
CONTROL POLL:
        IF STATION GTR 0 THEN
            BEGIN
            STATION = STATION -- 1.
    10:
            INITIATE REQUEST.
            INITIATE ENABLEINPUT.
            END.
        ELSE
            BEGIN
            STATION = MAXSTATIONS.
            GO TO 10.
            END.
```

GO TO 5.

#### **Semantics**

A \(\chicontrol\) definition\(\chi\) is an algorithm that describes the allocation of the use of a logical line to the stations assigned to that line. It is the \(\chicontrol\) definition\(\chi\) that decides if and when a station's Receive Request or Transmit Request should be initiated.

A single \(\langle\) control definition\) must control the logical line resource for all of the stations on a half-duplex line. In the case of full duplex, one \(\langle\) control definition\) must control the primary line, and one additional \(\langle\) control definition\) can be designated to control the auxiliary line. (One \(\langle\) control definition\) could be designated as the control for both the primary and auxiliary lines. If a \(\langle\) control definition\) is not designated as the control for an auxiliary line, then a default equivalent to an \(\langle\) idle statement\(\rangle\) is used.) The programmer, however, cannot directly define a particular \(\langle\) control definition\(\rangle\) for a logical line in its \(\langle\) line definition\(\rangle\). Instead, for each \(\langle\) terminal definition\(\rangle\), a single \(\langle\) control definition\(\rangle\) must be defined. (Two \(\langle\) control definition\(\rangle\) s can be named if the \(\langle\) terminal duplex statement\(\rangle\) specifies DUPLEX=TRUE.)

Next, in each \(\langle\) station definition\(\rangle\), a \(\langle\) terminal definition\(\rangle\) must be defined for the station (by means of the \(\langle\) station terminal statement\(\rangle\)). Finally, each \(\langle\) station definition\(\rangle\) is assigned to a logical line (in the \(\langle\) line station statement\(\rangle\) of the \(\langle\) line definition\(\rangle\)). This last association must be such that each station (i.e., \(\langle\) station definition\(\rangle\)) assigned to the logical line references (indirectly through its \(\langle\) terminal definition\(\rangle\)) the same \(\langle\) control definition\(\rangle\) as every other station assigned to the line.

A  $\langle control\ definition \rangle$  for a given line can deal only with one station at a time. All statements executed apply to and affect only one station assigned to the line. The  $\langle control\ definition \rangle$  chooses the station with which it wishes to deal by setting the value of the  $\langle byte\ variable \rangle$  STATION to the chosen station's station index.

(control identifier) has the same syntactic form as (identifier).

The statements in a  $\langle control\ definition \rangle$  are executed sequentially. In some cases, however, it is desirable to alter the order of execution of statements within the procedure. A  $\langle control\ statement \rangle$  preceded by a  $\langle label \rangle$  is one means of accomplishing this. The  $\langle go\ to\ statement \rangle$  is used to transfer control to a  $\langle label \rangle ed\ \langle control\ statement \rangle$ .

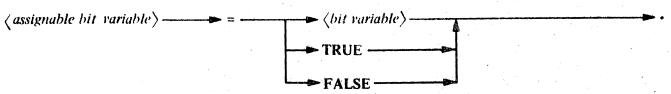
#### CONTROL

Assignment Statement

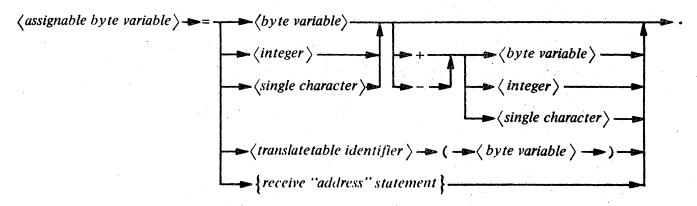
# **ASSIGNMENT STATEMENT**

# **Syntax**

#### FORM 1 - LOGICAL ASSIGNMENT



#### FORM 2 - VALUE ASSIGNMENT



#### Examples

TOG [0] = TRUE. TOG [1] = TOG [0]. LINE (BUSY) = FALSE. RETRY = STATION (TALLY). STATION = MAXSTATIONS.

TALLY [0] = STATION (FREQUENCY)—TALLY [1]. CHARACTER = TRANSTABLEID (CHARACTER). STATION = RECEIVE ADDRESS (TRANSMIT) [ADDERR:999].

#### **Semantics**

#### FORM 1

This form causes the value on the right side of the equal sign to replace the current value of  $\langle assignable bit variable \rangle$ .

#### FORM 2

Value assignment causes a calculated value on the right of the equal sign to be stored in the \( \assignable \) by te variable \( \alpha \). Arithmetic calculations are done in modulo 255 arithmetic.

 $\langle assignable\ byte\ variable \rangle = \langle translatetable\ identifier \rangle\ (\langle byte\ variable \rangle).$ 

This construct is the means to invoke user-defined character translation. User-defined translation is effected by three areas of the NDL source program.

a. In a \(\langle translate table \) definition\(\rangle\) the programmer must define the contents of a translation table and associate a \(\langle translate table \) identifier\(\rangle\) with it.

# Definitions CONTROL

#### Assignment Statement - Continued

b. In the \(\lambda\) terminal definition \(\rangle\) of a terminal type that requires special character translation the programmer should suppress automatic character translation by using either of the following forms of \(\lambda\) terminal code statement \(\rangle\):

CODE = BINARY.

or

#### CODE = EBCDIC.

c. In a  $\langle control\ definition \rangle$  or  $\langle request\ definition \rangle$ , the programmer invokes the translation by using this option of the value assignment. Any  $\langle byte\ variable \rangle$  can be designated as containing the character to be translated.

The \(\lambda\) translatetable identifier \(\rangle\) identifies the translation table to be used. An \(\lambda\) assignable by te variable \(\rangle\) is designated to the left of the equal sign, identifying where the resulting translated character is to be stored.

If N is the  $\langle source\ size \rangle$  (defined in the  $\langle translatetable\ definition \rangle$ ), then the N low-order bits of the  $\langle byte\ variable \rangle$  are used as an index into the translation table. The eight-bit character thus indexed is stored in the  $\langle assignable\ byte\ variable \rangle$ .

(assignable byte variable) = {receive "address" statement}.

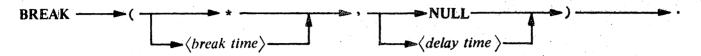
This construct attempts to RECEIVE the address characters of a terminal, ... d store in \( \alpha \) discrete in the semantics of the RECEIVE ADDRESS option of the \( \alpha \) receive statement \( \rangle \). The optional syntax in the \( \alpha \) receive "address" statement \( \alpha \) invokes the same actions as described in the \( \alpha \) receive statement \( \rangle \) semantics except for ADDERR. Any action specified for ADDERR is taken if no valid station assigned to the line is found with address characters equal to those received.

#### **CONTROL**

**Break Statement** 

# **BREAK STATEMENT**

**Syntax** 



# **Examples**

BREAK (\*, NULL).

BREAK (200 MILLI, 3 SEC).

BREAK (\*, 3 SEC).

BREAK (100 MILLI, NULL).

# **Semantics**

The  $\langle break \ statement \rangle$  causes binary zeros to be transmitted on the line, thus changing the state of the line to a "spacing" condition for a specified time.

 $\langle break \ time \rangle$  specifies the  $\langle time \rangle$  to break. An asterisk indicates that a standard break of 2 character times should be used.

 $\langle delay \ time \rangle$  specifies the  $\langle time \rangle$  to delay subsequent to the break and prior to when control continues.

#### **Semantics**

CODE=ASCII invokes the ASCII-to-EBCDIC translation for received data and the EBCDIC-to-ASCII translation for transmitted data.

CODE=BINARY inhibits any character translation on data transmitted or received.

#### **Pragmatics**

The \( \cdot \) code statement \( \rightarrow \) allows a programmer to either invoke or inhibit on a logical line the DCP ASCII-to-EBCDIC character code translation for input, and the EBCDIC-to-ASCII character code translation for output. Any \( \text{\centerminal definition} \right) \) that names, in its \( \text{\centerminal control statement} \rangle \), a \( \cdot \) control \( \definition \right) \) that utilizes the \( \langle \cdot \) code \( \statement \right) \), must define \( \text{ASCII(BINARY)} \) as its character code in the \( \left\) terminal \( \cdot \) code \( \statement \right) \). (Refer to the \( \left\) terminal \( \cdot \) code \( \statement \right) in this \( \character \) code \( \statement \right) in \)

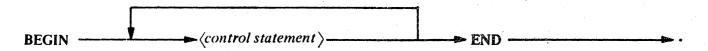
Once that translation has been invoked on a logical line, the translation continues until such time that it is inhibited. If translation is inhibited, translation will be inhibited on that line until invoked again by execution of CODE = ASCII, or if control is transferred to a \( \frac{request definition} \) which executes one of the following: CODE=ASCII, TERMINATE NORMAL, TERMINATE LOGICALACK, TERMINATE LOGICALACK (RETURN), TERMINATE ERROR, TERMINATE ENABLEINPUT, or (while executing a Receive Request) TERMINATE NOINPUT.

**CONTROL** 

**Compound Statement** 

#### **COMPOUND STATEMENT**

**Syntax** 



# **Examples**

BEGIN INITIATE TRANSMIT. TRANSMIT EOT. FINISH TRANSMIT. END.

IF STATION (VALID) THEN
IF STATION (READY) THEN
BEGIN % EXAMPLE 2
INITIATE TRANSMIT.
TRANSMIT "ON THE AIR".
FINISH TRANSMIT.
END.

% EXAMPLE 1

#### **Semantics**

The \(\chicompound statement\) groups several statements together to form a logical sequence. To execute more than one statement when the condition of an \(\langle if statement\rangle\) is satisfied, a \(\langle compound statement\rangle\) must be used.

# Definitions CONTROL Continue Statement

CONTINUES OF AT LANDAN			
Syntax		*	
		•	
CONTINUE -			

#### **Semantics**

CONTINUE STATEMENT

The \(\chicontinue statement\) can appear in only those \(\chicontrol definition\)\)s and \(\chicontent(request definition\)\)s written to communicate with full duplex terminal types. This statement causes the co-line to resume processing, if, and only if, it had been suspended by a \(\lambda wait statement\)\) or a \(\chicontent(receive statement\)\) with a CONTINUE option specified. If the co-line had not been suspended, this statement acts as a no-op. The \(\chicontinue statement\)\) has no effect upon the line on which it was executed.

# **Pragmatics**

Refer to the (fork statement) pragmatics.

CONTROL

**Delay Statement** 

**DELAY STATEMENT** 

**Syntax** 

DELAY (delay time)

Example

DELAY (10 MICRO).

# **Semantics**

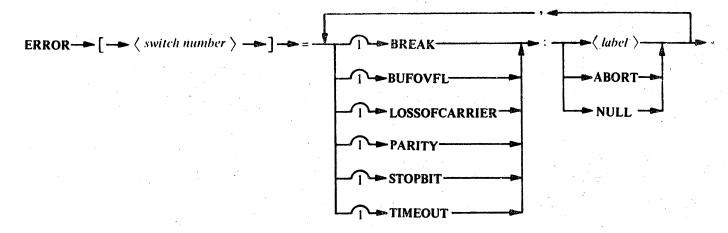
The  $\langle delay \ statement \rangle$  provides a means to delay a specified period of time before control proceeds to the next statement. The  $\langle control \ definition \rangle$  is suspended in a "sleep" state for the  $\langle delay \ time \rangle$  specified.

# **Pragmatics**

The "sleep" state induced by the \( \langle delay statement \rangle \) allows the DCP to service Cluster Attention Needed (CAN) interrupts for other logical lines.

#### ERROR SWITCH STATEMENT

#### Syntax



#### Examples

ERROR [0] = BREAK:0,BUFOVFL:NULL,LOSSOFCARRIER:ABORT PARITY:999, STOPBIT:999, TIMEOUT:NULL.

= BREAK:NULL.

ERROR[99] = BUFOVFL:NULL.

#### **Semantics**

The \(\langle error \) switch statement \(\rangle\) is a non-executable statement that allows the programmer to define a set of default actions that are to be taken in a \( \text{receive statement} \) if the specified errors occur.

 $\langle switch\ number \rangle$  has the syntactic form of  $\langle integer \rangle$ .

#### **BREAK**

The BREAK option variations cause actions as described if a break, that is, at least 2 character-times of a spacing line condition, is detected by the adapter cluster while receiving:

**BREAK:NULL** BREAK: (label)

causes no action. Execution proceeds as if the break did not occur. sets TRUE the \(\langle \) branches control

to \(\label\).

**BREAK: ABORT** 

sets TRUE the \(\langle bit variable \rangle \) BREAK[RECEIVE], and executes an

implicit TERMINATE ERROR.

#### **BUFOVFL**

The BUFOVFL option variations cause actions as described if the DCP is unable to service a cluster Attention Needed (CAN) interrupt before the Adapter Cluster receives another character (thus destroying the previous character):

**BUFOVFL:NULL** 

causes no action. Execution proceeds as if the error condition did not

occur.

**BUFOVFL**: (label) **BUFOVFL: ABORT**  sets TRUE the  $\langle bit \ variable \rangle$  BUFOVFL, and branches control to  $\langle labe_i \rangle$ . sets TRUE the \( \bit variable \) BUFOVFL, and executes an implicit

TERMINATE ERROR.

#### CONTROL

Error Switch Statement - Continued

#### LOSSOFCARRIER

The LOSSOFCARRIER option variations cause actions as described if a loss of carrier is detected while receiving.

LOSSOFCARRIER:NULL

causes no action. Execution proceeds as if the error

did not occur.

LOSSOFCARRIER: \(\langle \langle \lang

sets TRUE the \(\langle \text{bit variable} \rangle \text{LOSSOFCARRIER}, and

branches control to  $\langle label \rangle$ .

LOSSOFCARRIER: ABORT

sets TRUE the \(\langle bit variable \rangle LOSSOFCARRIER\), and

executes an implicit TERMINATE ERROR.

There is one exception to the actions described above. If a loss of carrier is detected while receiving, and if the terminal is modem-connect, and if the terminal's \( \station\) definition \( \rightarrow\) references a \( \sum\_{odem}\) definition \( \rightarrow\) that contains the statement LOSSOFCARRIER=DISCONNECT, then an implicit disconnect is done, regardless of the action associated with LOSSOFCARRIER in the \( \langle\) error switch statement \( \rightarrow\).

#### **PARITY**

The PARITY option variations cause actions as described if a parity bit error is detected by the adapter cluster:

**PARITY:NULL** 

causes no action. Execution proceeds as if the error did

not occur.

PARITY: (label)

sets TRUE the \( \frac{bit variable}{} \) PARITY, and branches control

to  $\langle label \rangle$ .

PARITY: ABORT

sets TRUE the \( \begin{aligned} bit variable \rangle PARITY, and executes a \end{aligned} \)

TERMINATE ERROR.

#### **STOPBIT**

The STOPBIT option variations cause the described actions if a stop bit error is detected by the adapter cluster:

STOPBIT: NULL

causes no action. Execution proceeds as if the error did

not occur.

**STOPBIT**: \(\langle \langle abel \rangle \)

sets TRUE the \( \frac{bit variable}{} \) STOPBIT, and branches control

to (label).

STOPBIT: ABORT

sets TRUE the \(\langle \) hit variable \(\rangle \) STOPBIT, and executes a

TERMINATE ERROR.

#### **TIMEOUT**

The TIMEOUT option variations cause the actions described if the time required to receive a character exceeds the \(\lambda timeout \time \rangle\). The \(\lambda timeout \time \rangle\) is defined in the \(\lambda terminal \timeout \timeout \time \rangle\), but can be overridden by including the \(\lambda timeout \time \rangle\)) or \((\text{NULL}\)) syntax options in the \(\lambda receive \text{statement}\).

TIMEOUT: NULL

causes no action. Execution proceeds as if the error did

not occur.

**TIMEOUT**: \(\label\)

sets TRUE the \( \frac{hit variable}{} \) TIMEOUT, and branches

control to  $\langle label \rangle$ .

TIMEOUT: ABORT

sets TRUE the \(\langle \text{hit variable} \rangle \text{TIMEOUT, and executes a}\)

TERMINATE ERROR.

# Definitions CONTROL

Error Switch Statement - Continued

# **Pragmatics**

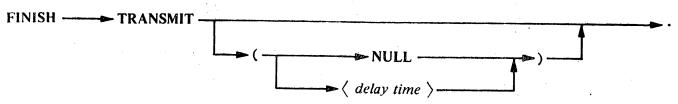
An \( \langle error switch statement \rangle \) must be associated with a \( \langle receive statement \rangle \) by means of a \( \langle switch number \rangle \) reference before any of the default actions will be invoked. The \( \langle error switch statement \rangle \) can appear in a \( \langle control definition \rangle \) as many times as the programmer deems convenient providing the following restriction is observed: Within a given \( \langle control definition \rangle \), \( \langle error switch statement \rangle s \) must have a unique \( \langle switch number \rangle \), and all \( \langle error switch statement \rangle s \) must precede all executable statements in the procedure.

#### CONTROL

Finish Statement

#### **FINISH STATEMENT**

**Syntax** 



# Examples

FINISH TRANSMIT. FINISH TRANSMIT (NULL). FINISH TRANSMIT (3 SEC).

#### **Semantics**

The purpose of the \( \frac{\text{finish statement}}{\text{character}} \) is to take a line out of the transmit ready state and prepare the line to receive information. The adapter cluster delays a period of time long enough for the last character TRANSMITted to be transmitted, plus 2 milliseconds, before the line is put in a receive ready state. Although the \( \frac{\text{finish statement}}{\text{puts puts the line in a receive ready state, the cluster hardware invokes a delay that inhibits any data from being received for 25 milliseconds. An INITIATE RECEIVE construct should precede any subsequent \( \text{receive statement} \), to override the 25 millisecond hardware delay.

The  $\langle delay \ time \rangle$  option allows the programmer to specify a software delay of  $\langle time \rangle$  before execution continues in the  $\langle control \ definition \rangle$ .

For example, the statement

FINISH TRANSMIT (3 SEC).

is equivalent to

FINISH TRANSMIT. DELAY (3 SEC).

The FINISH TRANSMIT (NULL) construct is equivalent to FINISH TRANSMIT.

FORK STATEMENT					
Syntax					
FORK —	→ \ label \>				
Example					

#### **Semantics**

**FORK 10.** 

The \( \fork\) statement\\ \) is allowed in only those \( \control\) definition\( \rangle \) and \( \control\) definition\( \rangle \) s that are written to communicate with full duplex terminal types. This statement can be executed in the \( \control\) definition\( \rangle \) or \( \control\) definition\( \rangle \) of the auxiliary line or the primary line. The execution of this statement causes the co-line control, if not busy, to branch to and begin executing code in the \( \control\) definition\( \rangle \) that executes the FORK at the \( \langle \langle \langle \langle \text{equest} \) specified, while control on the FORKing line executes an implicit PAUSE (i.e., a \( \chi \rangle \text{pause statement} \) and continues executing in parallel. The co-line is determined busy or not busy by testing the BUSY bit (i.e., LINE(BUSY)) or AUX(LINE(BUSY)), whichever is appropriate). If the co-line is busy, the \( \fork\) statement \( \rangle \) acts as a no-op.

#### **Pragmatics**

Synchronization problems can occur between the primary and auxiliary lines as a result of the \( fork \) statement \( \) executing the implicit PAUSE. The implicit PAUSE yields use of the DCP, to allow processing to proceed on other lines. Thus, processing on the co-line is actually started before the FORKing line exits the \( fork \) statement \( \). As a result, the programmer must, by some means (e.g., by setting and testing line TOGs), effect the synchronization of the lines. This is especially critical if the code contains \( \sqrt{wait} \) statement \( \) s and \( \sqrt{continue statement} \) s. The following example illustrates how full duplex lines could "hang" as a result of poor synchronization.

FORK 10. WAIT.

10: CONTINUE. WAIT.

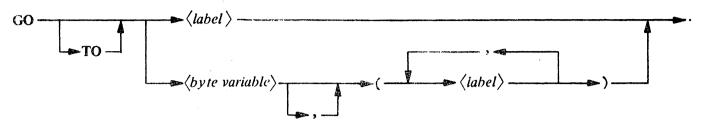
Assume that the primary line executes the FORK 10. At that point, the primary line temporarily yields use of the DCP to other lines. The auxiliary line starts up and executes the CONTINUE. Since primary control is still at the \( \langle fork statement \rangle \) and is not in a \( \langle wait statement \rangle \), the auxiliary line CONTINUE acts as a no-op. Next, the auxiliary line executes the WAIT. When the primary line is given use of the processor again, it executes its WAIT. At this point, the primary and auxiliary lines are "hung," each WAITing for a CONTINUE from its co-line.

#### **CONTROL**

Go To Statement

#### **GO TO STATEMENT**

#### **Syntax**



# **Examples**

GO 10.

GO TO 10.

GO TO TOGS, (0, 1, 2, 3).

GO TO STATION (5, 9, 12).

#### **Semantics**

The  $\langle go \ to \ statement \rangle$  alters the path of control, that is, the sequential flow of statement execution, within a  $\langle control \ definition \rangle$ .

GO TO (label)

This form of the  $\langle go \ to \ statement \rangle$  unconditionally transfers control to the  $\langle label \rangle$  specified.

**GO TO** \(\delta \) by te variable \(\rangle \)...

This form of the  $\langle go \ to \ statement \rangle$  provides a convenient means of dynamically selecting one or more  $\langle label \rangle s$  to which control could branch. The  $\langle label \rangle$  to branch to is selected by using the  $\langle bvte \ variable \rangle$  as an index value. If N represents the number of  $\langle label \rangle s$  in the  $\langle go \ to \ statement \rangle$ , then the  $\langle label \rangle s$  are numbered 0 to N-1. The  $\langle label \rangle$  corresponding to the index value is the  $\langle label \rangle$  to which control branches. If the index value is greater than N-1, then control continues at the statement following the  $\langle go \ to \ statement \rangle$ .

# Supplementary Example

GO TO STATION (5, 9, 12). % EXECUTION CONTINUES HERE IF STATION > 2.

- 5: TOG[0] = TRUE.
- 9: **TOG** [1] = **TRUE**.
- 12: TOG [2] = TRUE.

Go To Statement

Continued

This example illustrates the "GO TO  $\langle byte\ variable \rangle$ ..." option of the  $\langle go\ to\ statement \rangle$ . The value of STATION determines the next statement to be executed. If the value of STATION is 0, control branches to the  $\langle label \rangle$  5; if the value of STATION is 1, control branches to  $\langle label \rangle$  9; and if the value of STATION is 2, control branches to  $\langle label \rangle$  12. If the value of STATION is greater than 2, control continues at the next sequential statement.

Definitions
CONTROL

Idle Statement

B	D	8	ST	A"	ΓF	M	FN	JT
×	1,	ساة	 	Λ.		142		<b>,</b> ,

Syntax

IDLE-

#### **Semantics**

The execution of the *(idle statement)* causes a logical line to be suspended in an idle state. Specifically, IDLE causes the LINE(BUSY) *(bit variable)* to be set FALSE, the line to be suspended in a "sleeping" and "ready" status, and all subsequent inbound data to be discarded.

#### **Pragmatics**

The \(\langle idle statement\) suspends the execution of a \(\langle control definition\) for a logical line. Normally, this statement should be executed only when there are no outbound messages queued for any stations on the line and none of the stations on the line are ENABLED for input (or possibly, if the programmer wants any available inbound data discarded). Consider the following example of the contention-type \(\langle control \) definition\(\rangle\) taken from the Burroughs SYMBOL/SOURCENDL program.

#### **CONTROL CONTENTION:**

INITIATE REQUEST.
INITIATE ENABLEINPUT.
IDLE.

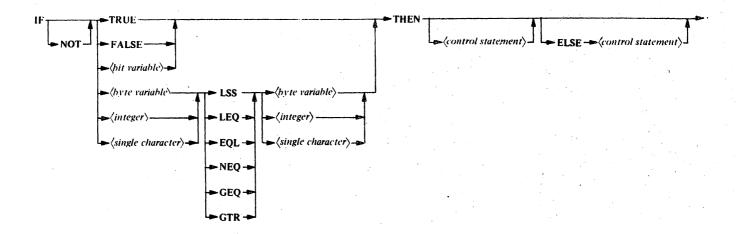
In this example, IDLE is executed only after it has been determined (by means of INITIATE REQUEST and INITIATE ENABLEINPUT) that the station is not QUEUED and not ENABLED for input.

Once a line is in an idle state, the line remains in an idle state until one of the following circumstances occurs:

- a. If the line TYPE is DIALIN and the line becomes connected (as a result of ANSWER = TRUE in the \(\langle \line \) definition\(\rangle\), a DIALOUT (TYPE = 98) DCWRITE from the MCS, or an ANSWER THE PHONE (TYPE = 100) DCWRITE from the MCS), the \(\langle \control \) definition\(\rangle\) is initiated for the line.
- b. If any of the following station-oriented DCWRITEs are executed for any station assigned to the line, then the \( \chion \) is initiated for the line.
  - 1. ENABLE INPUT (TYPE = 35)
  - 2. DISABLE INPUT (TYPE = 36)
  - 3. SET CHARACTERS (TYPE = 39)
  - 4. SET TRANSMISSION NUMBER (TYPE = 40)
  - 5. SET/RESET LOGICALACK (TYPE = 43)
  - 6. **NULL STATION REQUEST (TYPE = 48)**
  - 7. SET/RESET SEQUENCE MODE (TYPE = 49)
- c. If a WRITE request or a READ request is found in the DCP's Request Queue (placed there as a result of the MCS executing a WRITE (TYPE = 33) DCWRITE or a READ-ONCE ONLY (TYPE = 34) DCWRITE, or the I/O intrinsics) for a station on the line, then the appropriate \(\lambda request \) definition\(\rangle\) is initiated for the line.

#### IF STATEMENT

#### Syntax



#### **Examples**

IF TRUE THEN.

IF TOG[0] THEN TOG[0] = FALSE.

IF TALLY[0] LSS TALLY[1] THEN TALLY[0] = TALLY[1].

IF CHARACTER = 4"FF" THEN INITIATE BREAK.

**ELSE** 

BEGIN CHARACTER = 4"00". GO TO 0.

IF STATION(READY) THEN
IF STATION(QUEUED) THEN
LINE (TOG[0]) = TRUE.
ELSE
GO TO 10.

ELSE IDLE.

#### **Semantics**

The  $\langle if \, statement \rangle$  causes a condition (i.e., a Boolean expression) to be evaluated. The subsequent path of program control depends on whether the condition is evaluated as TRUE or FALSE.

If the condition is TRUE, the  $\langle control \ statement \rangle$  following the THEN, if present, is executed. Program control then resumes at a statement that follows the  $\langle if \ statement \rangle$ .

If the condition is FALSE, the  $\langle control statement \rangle$  following the ELSE is executed or, if the ELSE  $\langle control statement \rangle$  is omitted, program control resumes at the  $\langle control statement \rangle$  following the  $\langle if statement \rangle$ .

#### **CONTROL**

If Statement - Continued

The (control statement) can be any legal (control statement), including the (if statement) and (compound statement). The meanings of the relational operators are contained in table 5-1. The following diagrams illustrate the (if statement) semantics.

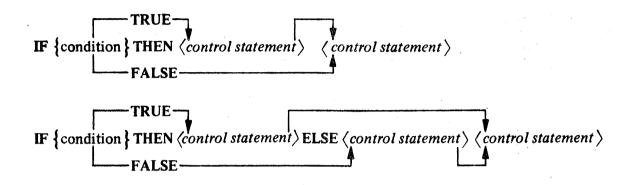
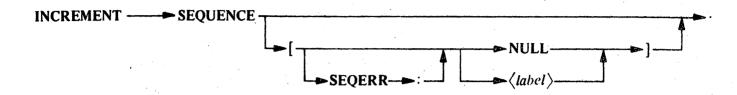


Table 5-1. Relational Operators

RELATIONAL OPERATOR	MEANING	SYNONYMS	
LSS	Less than	<and ls<="" td=""></and>	
LEQ	Less than or equal to	LE	
EQL	Equal to	= and EQ	
NEQ	Not equal to	NE	
GEQ	Greater than or equal to	GE	
GTR	Greater than	>and GT	

#### **INCREMENT STATEMENT**

#### Syntax



# **Examples**

INCREMENT SEQUENCE.
INCREMENT SEQUENCE (SEO)

INCREMENT SEQUENCE [SEQERR: NULL].

INCREMENT SEQUENCE [NULL].
INCREMENT SEQUENCE [999].

#### **Semantics**

The (increment statement) causes the sequence number stored in the DCP Station Table to be increased by the value of the increment (also stored in the DCP Station Table), providing that the station is in sequence mode; otherwise, this statement is a no-op.

When using the INCREMENT SEQUENCE construct, provision should be made for taking action if the increment caused the sequence number to exceed (overflow) the size of the sequence number field. The programmer can take such action by including the optional syntax. Failure to include overflow action results in an implicit TERMINATE ERROR if an overflow occurs.

The SEQERR:NULL and NULL options are semantically equivalent. These options set the SEQERR \( bit variable \) TRUE, and control continues at the next \( \) equential instruction.

The SEQERR:  $\langle label \rangle$  and  $\langle label \rangle$  options are semantically equivalent. They cause the SEQERR  $\langle bit \rangle$  variable to be set TRUE, and control to branch to  $\langle label \rangle$ .

Regardless of whether error action is specified or not, an overflow of the sequence number field destroys the contents of that field.

#### **Pragmatics**

#### **SEQUENCE MODE**

A station is considered to be in sequence mode whenever its SEQUENCE (bit variable) toggle is TRUE. SEQUENCE can be set TRUE only as a result of the controlling MCS executing the SET/RESET SEQUENCE MODE (TYPE = 49) DCWRITE. In addition, the TYPE 49 DCWRITE also stores the starting sequence number and increment in the appropriate fields of the LYP Station Table.

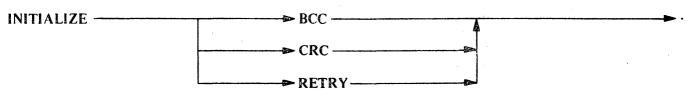
Sequence mode can be used for any application that the NDL programmer may see fit. Its use, however, requires common conventions between the NDL programmer and the MCS programmer. Burroughs has utilized sequence mode constructs in two \(\frac{request definition}{s}\) of S.\(\text{MBOL/SOURCENDL}:\)
READTELETYPE and WRITETELETYPE. Both require the coopera on of SYSTEM/CANDE to effect the execution of those statements. The reader is referred to those \(\frac{req}{est}\) definition\(\rangle\) s as an example of a particular application that Burroughs has implemented. Other statement \(\text{reasmit statement}\) (TRANSMIT SEQUENCE construct) and the \(\lambda\) statement\(\rangle\) (STORE SEQUENCE construct).

#### CONTROL

Initialize Statement

#### **INITIALIZE STATEMENT**

Syntax



# Examples

INITIALIZE BCC. INITIALIZE CRC. INITIALIZE RETRY.

#### **Semantics**

#### INITIALIZE BCC

The INITIALIZE BCC construct causes the \( \begin{align\*} byte variable \right) BCC to be initialized for purposes of accumulating a Block Check Character. The value to which BCC is initialized is dependent upon the horizontal parity defined for the station's associated \( \textit{terminal definition} \right) \) (in the \( \textit{terminal parity statement} \right) \). If horizontal parity is defined as HORIZONTAL:ODD, then BCC is initialized to all ones (i.e., 4"FF"). If defined as HORIZONTAL:EVEN, then INITIALIZE BCC initializes BCC to all zeroes (i.e., 4"00").

#### **INITIALIZE CRC**

The INITIALIZE CRC construct initializes CRC to the initial value required for calculating the Cyclic Redundancy Check. Any \(\lambda\text{terminal definition}\rangle\) referencing a \(\lambda\text{control definition}\rangle\) (in the \(\lambda\text{terminal parity statement}\rangle\)) that contains this instruction must define the horizontal parity (in the \(\lambda\text{terminal parity statement}\rangle\)) as HORIZONTAL:CRC(16); otherwise a syntax error is generated.

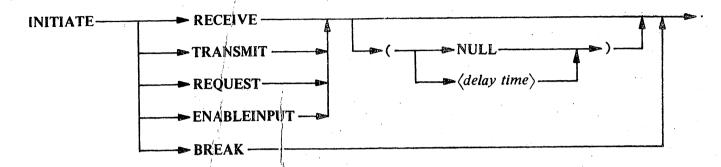
# **INITIALIZE RETRY**

The INITIALIZE RETRY construct causes the value stored in DCP INITIAL RETRY to be stored in DCP RETRY.

Initiate Statement

#### INITIATE STATEMENT

#### **Syntax**



# **Examples**

INITIATE RECEIVE, INITIATE TRANSMIT (3 SEC). INITIATE REQUEST (NULL). INITIATE ENABLEINPUT, INITIATE BREAK.

#### **Semantics**

#### INITIATE RECEIVE

The INITIATE RECEIVE const / ict causes the adapter cluster to initiate a receive delay calculated for the station. After the delay, the har laware is really to receive information.

The amount of time delayed, ferred to as the Initiate Receive delay, is unique to each station and is calculated at compile-time for each station. The algorithm that the compiler uses to calculate the Initiate Receive delay is described if the following three paragraphs.

- a. If the \( modem \) finition r ferenced in the \( station definition \) (in the \( station modem statement \) defines the m \( em \) NOISEDELAY as being greater than zero, then the Initiate Receive delay is 2 millisecond less than the combined \( \square time \) s defined in the \( modem noisedelay statement \) and the \( modem transmitdelay statement \).
- b. If the my lem NOISEDELAY is defined as zero and the modem TRANSMITDELAY is defined as being by than 7 milliseconds, then the Initiate Receive delay is zero.
- c. If the nodem NOISEDELAY is defined as zero and the modem TRANSMITDELAY is defined as being equal to or greater than 7 milliseconds, then the Initiate Receive delay is the lesser of 15 liliseconds or (1.5 milliseconds + modem TRANSMITDELAY).

The NULL option or the \( \langle \text{delay time} \rangle \) option can be used to override the calculated Initiate Receive delay.

NULL imediately readies the hardware so that it can receive information. \( \langle \text{delay time} \rangle \) specifies a \( \langle \text{time} \rangle \) ed in place of the Initiate Receive delay.

#### CONTROL

Initiate Statement - Continued

#### **Pragmatics**

An INITIATE RECEIVE instruction should precede the first \( \text{receive statement} \) following a transmission. If it does not, there is a possibility that execution of the \( \text{receive statement} \) will be delayed for a period of time of up to 25 milliseconds. The cause of the 25-millisecond delay is described under the semantics of the \( \text{finish statement} \).

#### **INITIATE TRANSMIT**

The INITIATE TRANSMIT construct causes the Adapter Cluster to be put in a transmit state after a calculated delay. The amount of time delayed is referred to as the Initiate Transmit Delay, and is unique to each station. It is derived by taking the greater of the NOISEDELAY \( \text{time} \) specified for the modem configured at the system end, or the TURNAROUND \( \text{time} \) specified by the station's associated \( \text{terminal definition} \). This construct must be executed prior to any attempt to transmit information.

The NULL option or the  $\langle delay \ time \rangle$  option can be used to override the calculated Initiate Transmit delay. NULL causes the adapter cluster to be put in a transmit state immediately.  $\langle delay \ time \rangle$  specifies a  $\langle time \rangle$  to be used in place of the Initiate Transmit delay.

#### INITIATE REQUEST

The INITIATE REQUEST construct conditionally initiates the next function as indicated by the message at the head of the Station Queue. The initiation of the function is conditional, subject to the following: the station must be valid, ready, and queued. Specifically, STATION(VALID), STATION(READY), and STATION(QUEUED) must be TRUE; otherwise, the instruction acts as a no-op.

The specific function invoked by this construct is dependent upon the type of message at the head of the Station Queue. Most commonly the message is a WRITE (TYPE=33) DCWRITE, thus causing the Transmit Request for the station to be entered. A READ-ONCE ONLY (TYPE=34) DCWRITE message at the head of the Station Queue would cause control to enter the Receive Request for the station. Other messages (unrelated to input or output) invoke their specific function and then transfer control to the beginning of the \( \chiontrol \text{definition} \rangle \). For example, a SET SEQUENCE MODE (TYPE=49) DCWRITE message would cause control to enter the subroutine of the DCP that handles setting sequence mode and, when finished, control would be transferred to the beginning of the \( \chiontrol \text{definition} \rangle \).

The  $\langle delay \ time \rangle$  option allows the programmer to specify that an implicit  $\langle delay \ statement \rangle$  for the  $\langle time \rangle$  specified, be executed before initiation of the next function from the Station Queue. For example, the statement

**INITIATE REQUEST (3 SEC).** 

is equivalent to

IF STATION(VALID) THEN

IF STATION(READY) THEN

IF STATION(QUEUED) THEN

BEGIN DELAY(3 SEC). INITIATE REQUEST. END.

The INITIATE REQUEST (NULL) construct is equivalent to INITIATE REQUEST.

Initiate Statement - Continued

# INITIATE ENABLEINPUT

The INITIATE ENABLEINPUT construct conditionally transfers control to the receive request appropriate for the station (that is, the station referenced by the \( \delta y te variable \) named STATION). The transfer of control is conditional, subject to the following: the station must be valid, ready, and enabled for input. More specifically, STATION(VALID), STATION(READY), AND STATION(ENABLED), must be TRUE; otherwise, the instruction acts as a no-op.

The NDL programmer can initially enable a station for input by means of the *(station enableinput statement)*. Additionally, after DCP initialization, the station's MCS can enable or disable the station for input by means of the **TYPES** 35 and 36 **DCWRITES**.

(NULL) and ( $\langle delay\ time \rangle$ ) allow the programmer to specify that an implicit  $\langle delay\ statement \rangle$ , for time specified, be executed before the transfer of control.  $\langle delay\ time \rangle$  has the syntactic form of time  $\rangle$ . For example, the statement

INITIATE ENABLEINPUT (3 SEC).

is equivalent to:

IF STATION(VALID) THEN

IF STATION(READY) THEN

IF STATION(ENABLED) THEN

BEGIN DELAY (3 SEC). INITIATE ENABLEINPUT. END.

The (NULL) option specifies zero delay.

#### **INITIATE BREAK**

The INITIATE BREAK construct causes binary zeroes to be transmitted on the line, thus changing the state of the line to a "spacing" condition. The line remains in the spacing condition until some subsequent construct causes the adapter cluster to change the state of the line. Constructs that would change the line's state are INITIATE TRANSMIT, INITIATE RECEIVE, FINISH TRANSMIT, BREAK, and IDLE.

Definitions
CONTROL
Pause Statement

PA	I	ISF	CT	Δ	rfn	/FN	JT

**Syntax** 

PAUSE:

#### **Semantics**

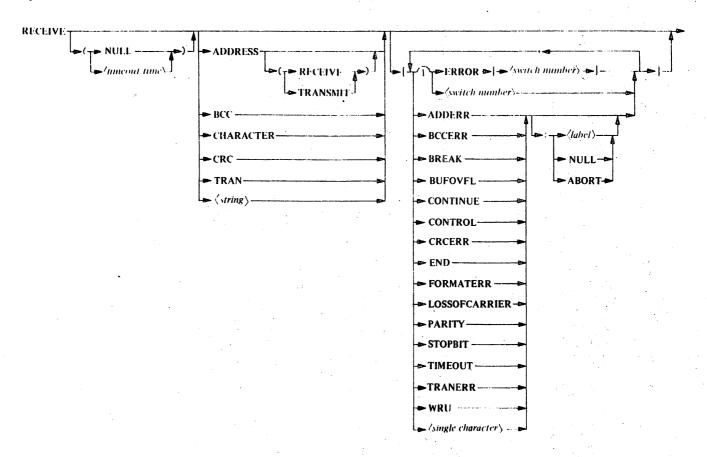
The \(\langle pause statement \rangle \) suspends a \(\langle control definition \rangle \) in a "sleep" state for a minimum period of time (200 microseconds for the B 6358 Model II DCP, and 6 microseconds for the B 6350 Model I DCP) to allow the DCP to service other lines. It is recommended that a \(\langle pause statement \rangle \) be used in any kind of loop that would tie up processor time and thereby prevent the servicing of other lines. The failure to do so results in a high number of timeout faults.

#### **Pragmatics**

Instances may occur in which the DCP requires an even greater period of "sleep" to service other lines. Repeated timeout faults, despite utilization of the  $\langle pause\ statement \rangle$ , are indications of such conditions. A greater period of "sleep" time can be effected by means of a  $\langle delay\ statement \rangle$ , with the  $\langle time \rangle$  specified greater than "sleep" time effected by the  $\langle pause\ statement \rangle$ .

#### **RECEIVE STATEMENT**

# **Syntax**



# **Examples**

RECEIVE.

RECEIVE CHARACTER.

RECEIVE (3 SEC) ADDRESS (RECEIVE) [0, ADDERR:10].

RECEIVE (NULL) [

PARITY:999, LOSS OF CARRIER:999, END, WRU:NULL

RECEIVE CRC [ERROR [1], CRCERR:10].

#### CONTROL

Receive Statement - Continued

#### **Semantics**

The \(\frac{receive statement}\) causes the adapter cluster to attempt to receive information from the appropriate logical line.

The following two syntax items define a maximum amount of time that the adapter cluster should wait for receipt of the first character, and then each subsequent character, if applicable, before assuming that the terminal has "timed out." If neither of these options is included, the \(\lambda timeout \time\rangle\) defined (in the \(\lambda terminal \timeout \time\rangle\)) for the station's \(\lambda terminal \text{definition}\rangle\) is implicitly used as the \(\lambda timeout \text{timeout}\) in this statement.

#### (NULL)

This option specifies that the adapter cluster should wait an infinite amount of time.

# (\langle time \rangle)

The \(\lambda time\rightarrow\) defines a \(\lambda time\rightarrow\) that the adapter cluster should wait for a character. If this \(\lambda time\rightarrow\) is exceeded before receipt of a character, and the TIMEOUT syntax appears, then the action specified for TIMEOUT is taken (refer to TIMEOUT). If the \(\lambda timeout \time\rightarrow\) is exceeded and TIMEOUT syntax does not appear, an implicit TERMINATE ERROR is executed.

The following syntax options define the nature of the information to be received, the amount of information to be received, and how the information is to be handled. If these options are omitted, it is semantically equivalent to specifying CHARACTER (i.e., "RECEIVE." is semantically equivalent to "RECEIVE CHARACTER.")

#### **ADDRESS**

The proper number of address characters (as defined by the station's \(\lambda terminal \) definition\(\rangle\) in the \(\lambda terminal \) address \(size \) statement\(\rangle\)) are received and checked for agreement against the actual address characters defined in the \(\lambda \) station address \(station \) an implicit \(text{TERMINATE ERROR}\) is \(station \) address \(station \) address \(station \) an implicit \(text{TERMINATE ERROR}\) is \(station \) address \(station \) and \(station \) address \

#### ADDRESS (RECEIVE)

This option is equivalent to ADDRESS, except that ADDRESS (RECEIVE) must be used when an address pair is defined in the \( \station address statement \) and the programmer needs to check for the proper receive address.

#### ADDRESS (TRANSMIT)

This option is equivalent to ADDRESS, except that ADDRESS (TRANSMIT) must be used when an address pair is defined in the \( \station address statement \) and the programmer needs to check for the proper transmit address.

#### **BCC**

One character is received and checked against the \( \frac{byte variable}{\rm} \) BCC. If the character received and BCC are not equal, a Block Check Character error condition results. If the BCCERR syntax appears, then specified action is taken; otherwise an implicit TERMINATE ERROR is executed.

Presumably, if the RECEIVE BCC instruction appears, the programmer has defined horizontal parity in the \(\lambda terminal parity statement\rangle\), and the accumulated Block Check Character is contained in BCC.

Receive Statement - Continued

#### **CHARACTER**

One character is received and stored in CHARACTER.

#### **CRC**

Two characters are received. The first character is checked against CRC[0], and the second compared against CRC[1]. If the characters received and CRC are not equal, a Cyclic Redundancy Check error condition results. If the CRCERR syntax appears, then specified action is taken; otherwise an implicit TERMINATE ERROR is executed.

Presumably, if the RECEIVE CRC construct appears, the programmer has defined horizontal parity HORIZONTAL:CRC(16) in the \(\lambda terminal parity statement \rangle\), and the Cyclic Redundancy Check is contained in CRC[0] and CRC[1].

#### TRAN

The proper number of transmission number characters (as defined by the station's associated \(\lambda terminal \) definition\(\rangle\) in the \(\lambda terminal \) transmission number length statement\(\rangle\)) are received and checked for agreement with the Receive Transmission Number maintained in the DCP Station Table. If the characters received and the Receive Transmission Number are not equal, a transmission number error results. If the TRANERR syntax option appears, then specified action is taken; otherwise, an implicit TERMINATE ERROR is executed.

(string)

The number of characters as indicated by the length of the  $\langle string \rangle$  are received and checked against those characters in the  $\langle string \rangle$ . If the  $\langle string \rangle$  and the characters received are not equal, then a format error condition results. If the **FORMATERR** syntax option appears, then that action is taken; otherwise an implicit **TERMINATE ERROR** is executed.

The following syntax options specify actions to be taken upon either the receipt of defined characters or occurrences of specific error conditions:

# **ERROR**[\(\langle\) switch number\(\rangle\)]

Associates a previously defined Error Switch with the \( \text{receive statement} \)\). This allows the programmer to associate a set of previously defined error actions with the \( \text{receive statement} \)\), thus reducing the amount of coding required for each \( \text{receive statement} \)\). BREAK, BUFOVFL, LOSSOFCARRIER, PARITY, STOPBIT, and TIMEOUT syntax options are not allowed if the ERROR[\( \sqrt{switch number} \)\)] syntax appears in the \( \text{receive statement} \)\). Refer to the \( \text{error switch statement} \)\) for more information.

(switch number)

Semantically equivalent to **ERROR**[ (switch number)].

#### CONTROL

#### Receive Statement - Continued

#### ADDERR

The **ADDERR** option variations cause the following actions if an address error is detected when attempting to receive the address characters of a terminal:

**ADDERR** 

sets TRUE the ADDERR (bit variable) and branches control to the next

sequential statement.

ADDERR:NULL

causes no action. Execution proceeds as if the error condition did not

occur.

ADDERR: \(\label\rangle\)

sets TRUE the ADDERR (bit variable) and branches control to (label).

ADDERR: ABORT

Not allowed.

#### **BCCERR**

The BCCERR option variations cause the following actions if the character received is not equal to the data stored in BCC.

**BCCERR** 

sets TRUE the (bit variable) BCCERR, and branches control to the next

sequential statement.

BCCERR:NULL

causes no action. Execution proceeds as if the error condition did not

occur.

BCCERR: \(\langle \langle \langle \rangle \)

sets TRUE the \( \langle bit variable \rangle \) BCCERR and branches control to \( \langle label \rangle \).

BCCERR: ABORT

Not allowed.

#### BREAK

The BREAK option variations cause the following actions if a break, that is, at least two character-times of a spacing line condition, is detected by the adapter cluster while receiving:

**BREAK** 

sets TRUE the \( \frac{bit variable}{} \) BREAK[RECEIVE], and branches control

to the next sequential statement.

**BREAK:NULL** 

causes no action. Execution proceeds as if the break did not occur.

BREAK: \(\label\)

sets TRUE the \( \begin{array}{c} bit variable \rangle \) BREAK [RECEIVE], and branches control

to  $\langle label \rangle$ .

**BREAK: ABORT** 

sets TRUE the (bit variable) BREAK[RECEIVE], and executes an

implicit TERMINATE ERROR.

Receive Statemen. - Continued

#### **BUFOVFL**

The BUFOVFL option variations cause the following actions if the DCP is unable to service a Cluster Attention Needed (CAN) interrupt before the adapter cluster receives another character (thus destroying the previous character):

BUFOVFL sets TRUE the \( \begin{array}{ll} bit variable \rangle BUFOVFL, and branches control to the next \end{array}

sequential statement.

BUFOVFL:NULL causes no action. Execution proceeds as if the error condition did not

occur.

BUFOVFL: (label) sets TRUE the (bit variable) BUFOVFL, and branches control to

 $\langle label \rangle$ .

BUFOVFL: ABORT sets TRUE the \(\langle \text{hit variable} \rangle \text{BUFOVFL}\), and executes an implicit

TERMINATE ERROR.

# CONTINUE

This option is allowed only in \( \text{receive statement} \) s of \( \text{control definition} \) s and \( \text{request definition} \) s that are written to communicate with full duplex terminal types. CONTINUE syntax causes action as described below if the co-line executes a \( \text{continue statement} \) before all information specified by the \( \text{receive statement} \) is received.

**CONTINUE** branches control to the next sequential statement.

CONTINUE: NULL causes no action. Execution proceeds as if the (continue statement) had

not been executed.

**CONTINUE**:  $\langle label \rangle$  branches control to  $\langle label \rangle$ .

CONTINUE: ABORT Not allowed.

#### **CONTROL**

The **CONTROL** option variations cause the following actions if the control character of the station (as defined in the (station control character statement) is received:

control sets TRUE the (bit variable) control sets TRUE the (bit variable) control

to the next sequential statement.

CONTROL: NULL sets TRUE the \(\begin{array}{l} bit variable \) CONTROLFLAG, and execution continues

as if the character was not the station's control character.

control: (label) sets TRUE the (bit variable) control to

 $\langle label \rangle$ .

CONTROL: ABORT Not allowed.

CONTROL

Receive Statement - Continued

#### **CRCERR**

The following CRCERR option variations cause the following actions if the first character received does not equal CRC[0], or the second character received does not equal CRC[1]. (This item is appropriate only for the RECEIVE CRC form of the (receive statement); refer to the CRC option.)

**CRCERR** 

sets TRUE the \( \frac{bit variable}{} \) CRCERR, and branches control to the next

sequential statement.

**CRCERR:NULL** 

causes no action. Execution proceeds as if the error did not occur.

CRCERR: \(\langle \langle \lan

sets TRUE the \( bit variable \) CRCERR, and branches control to \( label \).

**CRCERR: ABORT** 

Not allowed.

# **END**

The END option variations cause the following actions if the "end" character of the station (as defined by the \(\lambda terminal end character statement\rangle\) in the \(\lambda terminal definition\rangle\) associated with the station) is received:

**END** 

causes control to branch to the next sequential statement.

**END:NULL** 

causes no action. Execution proceeds as if the character was not the

" end" character.

END: \(\label\)

branches control to (label).

**END:ABORT** 

Not allowed.

# **FORMATERR**

The following variations of the **FORMATERR** option cause the following actions if the characters received are not equal to those in the  $\langle string \rangle$  (this item is appropriate only for the **RECEIVE**  $\langle string \rangle$  construct of the  $\langle receive\ statement \rangle$ ):

**FORMATERR** 

sets TRUE the \( \begin{aligned} bit variable \rangle FORMATERR, and branches control to the next sequential statement. \end{aligned}

FORMATERR: NULL

causes no action. Execution proceeds as if the

error did not occur.

FORMATERR: \(\label\rangle\)

sets TRUE the (bit variable) FORMATERR, and

branches control to  $\langle label \rangle$ .

FORMATERR: ABORT

not allowed.

# LOSSOFCARRIER

The LOSSOFCARRIER option variations cause the following actions if a loss of carrier is detected while receiving.

LOSSOFCARRIER sets TRUE the (bit variable) LOSSOFCARRIER, and branches

control to the next sequential statement.

LOSSOFCARRIER: NULL causes no action. Execution proceeds as if the error did not

occur.

LOSSOFCARRIER: \(\langle label \rangle \) sets TRUE the \(\langle bit variable \rangle \) LOSSOFCARRIER, and branches

control to  $\langle lahel \rangle$ .

LOSSOFCARRIER: ABORT sets TRUE the (bit variable) LOSSOFCARRIER, and executes

an implicit TERMINATE ERROR.

There is one exception to the actions described above. If a loss of carrier is detected while receiving, and if the terminal is modem-connect, and if the terminal's \( \station definition \) references a \( \square modem definition \) that contains the statement LOSSOFCARRIER=DISCONNECT, then an implicit disconnect is done, regardless of the action specified.

# **PARITY**

The **PARITY** option variations cause the following actions if a parity bit error is detected by the adapter cluster:

PARITY sets TRUE the (bit variable) PARITY, and branches control to the

next sequential statement.

PARITY: NULL causes no action. Execution proceeds as if the error did not occur.

PARITY: \(\lambda \) label \(\rangle \) sets TRUE the \(\lambda \) bit variable \(\rangle \) PARITY, and branches control to

 $\langle label \rangle$ .

PARITY: ABORT sets TRUE the (bit variable) PARITY, and executes a

TERMINATE ERROR.

#### **STOPBIT**

The STOPBIT option variations cause the described actions if a stop bit error is detected by the adapter cluster:

sets TRUE the (bit variable), and branches control to the next

sequential statement.

STOPBIT:NULL causes no action. Execution proceeds as if the error did not occur.

sets TRUE the \( \begin{array}{ll} bit variable \right) STOPBIT, and branches control to

 $\langle label \rangle$ .

STOPBIT: ABORT sets TRUE the \( \langle \text{bit variable} \rangle \) STOPBIT, and executes a

TERMINATE ERROR.

#### CONTROL

Receive Statement - Continued

#### TIMEOUT

The **TIMEOUT** option variations cause the actions described if the time required to receive a character exceeds the  $\langle timeout \ time \rangle$ . The  $\langle timeout \ time \rangle$  is defined in the  $\langle terminal \ timeout \ statement \rangle$ , but can be overridden by including the ( $\langle timeout \ time \rangle$ ) or (NULL) syntax options in the  $\langle receive \ statement \rangle$ .

TIMEOUT

sets the (bit variable) TIMEOUT, and branches control to the

next sequential statement.

TIMEOUT:NULL

causes no action. Execution proceeds as if the error did not occur.

TIMEOUT: \(\label\rangle\)

sets TRUE the \(\langle \text{bit variable} \rangle \text{TIMEOUT, and branches control to}\)

(label).

TIMEOUT: ABORT

sets TRUE the \( \frac{bit variable}{} \) TIMEOUT, and executes a

TERMINATE ERROR.

#### **TRANERR**

The TRANERR option variations cause the described actions if the characters received and the Receive Transmission Number stored in the Station Table are not equal (this item is allowed only in the RECEIVE TRAN construct of the \( \text{receive statement} \):

**TRANERR** 

sets TRUE the \( bit variable \) TRANERR, and branches control

to the next sequential statement.

TRANERR: NULL

causes no action. Execution proceeds as if the error did not occur.

TRANERR: (label)

sets TRUE the \(\langle\) bit variable TRANERR, and branches control to

(label).

TRANERR: ABORT

not allowed.

# WRU

The WRU option causes the following actions if the WRU character of the station is received (the \( \station WRU \) character statement \( \) defines the WRU character):

WRU

sets TRUE the WRU (hit variable), and branches control to the

next sequential statement.

WRU:NULL

sets TRUE the WRU \( bit variable \), and execution proceeds as if

the character received was not the WRU character.

WRU: \(\label\)

sets TRUE the (bit variable) WRU, and branches control to

 $\langle label \rangle$ .

WRU: ABORT

not allowed.

#### CONTROL

Receive Statement — Continued

*{single character}* 

The  $\langle single\ character \rangle$  syntax causes the following actions if a character received is equal to the  $\langle single\ character \rangle$ :

\( \single \character \rangle \)

branches control to the next sequential statement.

\(\langle\) single character \(\rangle\): NULL

causes no action. Execution proceeds as if the character received

was not equal to the (single character).

\(\langle character\rangle : \langle label\rangle

branches control to \(\lambda \lambda bel \rangle.

\(\langle\) character\(\rangle\): ABORT

not allowed.

The allowable combinations of the  $\langle receive\ statement \rangle$  syntax options are defined in table 5-2. The (NULL) and ( $\langle timeout\ time \rangle$ ) options are allowed in any form of the  $\langle receive\ statement \rangle$ . Allowed combinations of the other syntax options are denoted by a "X" in the appropriate columns and rows.

# Supplementary Examples

#### Statement

RECEIVE (3 SEC) [TIMEOUT:10].

RECEIVE ADDRESS [ADDERR:99].

RECEIVE CHARACTER [CONTINUE: 10, CONTROL: 20, TIMEOUT: 30, "\*": 40].

# Explanation

Causes the adapter cluster to attempt to receive a character. If the character is not received within 3 seconds, the \( \delta it variable \right) TIMEOUT is set TRUE and control branches to 10.

If the character(s) received do not equal those defined in the \( station address statement \), the \( \lambda bit variable \) ADDERR is set TRUE, and control branches to 99.

This statement would only be allowed in a \(\langle control \, definition \rangle \) or \(\langle request \, definition \rangle \) that is written to communicate with full duplex terminal types because it contains the CONTINUE option.

**CONTINUE: 10** would cause a branch to 10 if the co-line  $\langle control\ definition \rangle$  executes a  $\langle continue\ statement \rangle$  before a character is received.

CONTROL: 20 would set CONTROLFLAG TRUE and branch to 20 if the character received is the station's control character.

TIMEOUT:30 would set TIMEOUT TRUE and branch to 30 if a character is not received within the \(\lambda timeout \) time\(\rangle\) defined in the \(\lambda terminal \) timeout statement\(\rangle\).

"\*":40 would cause a branch to 40 if the character received is the asterisk character.

# Definitions **CONTROL**

# Receive Statement - Continued

# Statement

# Explanation

RECEIVE[ERROR[0]].

An attempt is made to receive one character and store it in CHARACTER. If any errors described in the associated *(error switch statement)* occur while receiving, then the action defined in that *(error switch statement)* is taken.

RECEIVE[0]

Same as above.

Table 5-2. Allowable Combinations for (receive statement)

	ADDERR	BCCERR	BREAK	BUFOVFL	CONTINUE	CONTROL	CRCERR	END	FORMATERR	LOSSOFCARRIER	PARITY	STOPBIT	TIMEOUT	TRANERR	WRU	⟨single character⟩
ADDRESS	X		X	X	X			•		X	X	X	X			
ADDRESS(RECEIVE)	X		X	X	X					X	X	X	X			
ADDRESS(TRANSMIT)	X		X	$\mathbf{X}$	X					X	X	X	$\mathbf{X}$			•
BCC		X	X	X	X					X	X	X	X			
CHARACTER			X	$\mathbf{X}$	$^{-}$ X	X		X		X	X	X	X		X	X
CRC			X	X	X		X			$\mathbf{X}^{-}$	$\mathbf{X}^{-}$	X	X			
(string)			X	X	X			,	X	X	- X	X	X			
TRAN			X	$\mathbf{X}_{i}$	<b>X</b>					X	X	X	X	X		

SHIFT STATEMENT			
Syntax			
SHIFT —	<b>&gt;</b> UP	<u> </u>	
	DOWN —	J	

#### **Semantics**

The  $\langle shift\ statement \rangle$  is to be used in a  $\langle control\ definition \rangle$  that communicates with stations using the Baudot (5-bit) character code set. (The character code set is defined in the  $\langle terminal\ code\ statement \rangle$  of the associated  $\langle terminal\ definition \rangle$ ).

SHIFT UP indicates that received characters are to be translated to their respective uppercase graphics (usually referred to as FIGS).

SHIFT DOWN indicates that receive characters are to be translated to their respective lowercase graphics (usually referred to as LTRS).

If the station does not use Baudot code, the  $\langle shift \ statement \rangle$  acts as a no-op.

# **Pragmatics**

In the Baudot character code set, most bit patterns have two graphic representations: one is referred to as FIGS (the uppercase graphic), and the other as LTRS (the lowercase graphic).

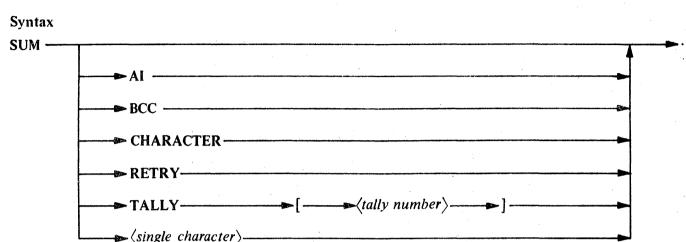
When transmitting to a terminal that uses Baudot code, the terminal prints LTRS until it receives a specially designated character indicating that it should shift to printing FIGS. The terminal continues printing the FIGS until it receives a specially designated character indicating that it should resume printing the LTRS.

When the information is received from a terminal that uses Baudot, the same conventions hold true; that is, the terminal communicates whether FIGS or LTRS follow by the transmission of a designated character. The terminal initially transmits LTRS.

#### CONTROL

**Sum Statement** 

#### **SUM STATEMENT**



# **Examples**

SUM AI. SUM CHARACTER. SUM "A". SUM TALLY [1].

#### **Semantics**

The purpose of the \( \sum \statement \) is to affect the calculation of the horizontal parity check (whether that be a Block Check Character or a Cyclic Redundancy Check). The specific effect of the \( \sum \statement \) is dependent upon two factors: the SUMmed item, and whether the station's \( \statement \) for which \( \sum \control \) definition \( \sigma \) is running, defines horizontal parity as CRC(16).

Following is a description of the effect that each form of the \( \sum statement \) has on the calculation of the horizontal parity check. Any reference to CRC means CRC[0] and CRC[1] collectively.

# **SUM**

Semantically equivalent to SUM CHARACTER.

# **SUM AI**

If the horizontal parity check is a Block Check Character or is undefined, the contents of AI are exclusively OR-ed with the contents of BCC, and the result is stored in BCC.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of AI and CRC, and the result is stored in CRC.

#### **SUM BCC**

If the horizontal parity check is a Block Check Character or is undefined, then the contents of BCC are exclusively OR-ed with itself, and the result is stored in BCC. (The result in BCC would be zero in this case.)

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of CRC[0] and CRC, and the result is stored in CRC.

Sum Statement - Continued

# **SUM CHARACTER**

If the horizontal parity check is a Block Check Character or is undefined, the contents of CHARACTER are exclusively OR-ed with the contents of BCC, and the result is stored in BCC.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of CHARACTER and CRC, and the result is stored in CRC.

#### **SUM RETRY**

If the horizontal parity check is a Block Check Character or is undefined, the contents of RETRY are exclusively OR-ed with the contents of BCC, and the result stored in BCC.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of **RETRY** and **CRC**, and the result is stored in **CRC**.

# SUM TALLY [\langle tally number \rangle]

If the horizontal parity check is a Block Check Character or is undefined, the contents of TALLY [ $\langle tally number \rangle$ ] are exclusively OR-ed with the contents of BCC, and the result is stored in BCC.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of TALLY [\langle tally number \rangle] and the result is stored in CRC.

# **SUM** \(\single \character\)

If the horizontal parity check is a Block Check Character or is undefined, the  $\langle single\ character \rangle$  is exclusively OR-ed with the contents of BCC, and the result is stored in BCC.

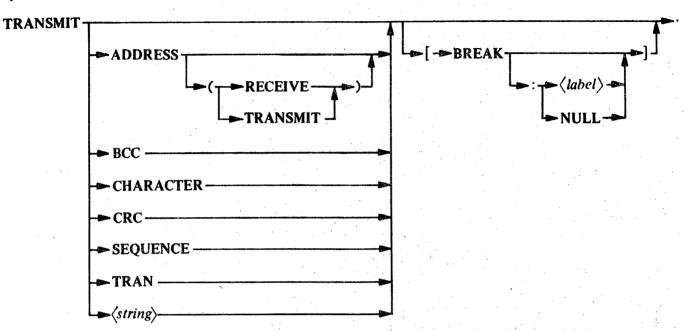
If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the (single character) and CRC, and the result is stored in CRC.

CONTROL

**Transmit Statement** 

#### TRANSMIT STATEMENT

**Syntax** 



# Examples

TRANSMIT,
TRANSMIT CHARACTER [BREAK:NULL],
TRANSMIT SOH STX 4"00"[BREAK:10],
TRANSMIT TRAN,
TRANSMIT ADDRESS (TRANSMIT) [BREAK].

#### **Semantics**

The \(\langle transmit statement\rangle\) causes the adapter cluster to transmit information to a terminal. The following group of syntax options specifies the information to be transmitted. All options except CHARACTER use the \(\langle byte variable\rangle\) CHARACTER as a temporary storage area; thus, any information contained in CHARACTER before execution of the \(\langle transmit statement\rangle\) shall be destroyed by the \(\langle transmit statement\rangle\). If none of the first group of options is chosen, it is semantically equivalent to specifying CHARACTER (i.e., "TRANSMIT." is equivalent to "TRANSMIT CHARACTER.").

#### **ADDRESS**

The proper number of characters (as specified by the station's associated  $\langle terminal \ definition \rangle$  in the  $\langle terminal \ address \ size \ statement \rangle$ ) are taken from the address field in the Station Table and transmitted.

# **ADDRESS (RECEIVE)**

This option is equivalent to ADDRESS, except that ADDRESS (RECEIVE) must be used when an address pair is defined in the *(station address statement)* and the programmer wants to transmit the receive address.

#### ADDRESS (TRANSMIT)

This option is equivalent to ADDRESS, except that ADDRESS (TRANSMIT) must be used when an address pair is defined in the (station address statement) and the programmer wants the transmit address transmitted.

#### **BCC**

The BCC option causes the contents of the (byte variable) BCC to be transmitted.

#### **CHARACTER**

The CHARACTER causes the contents of the (byte variable) CHARACTER to be transmitted.

#### **CRC**

This option causes two bytes to be transmitted; the contents of CRC [0] are transmitted first, followed by CRC [1]. If the station's associated \(\lambda terminal definition\rangle\) does not define horizontal parity as CRC ([16]), the use of this \(\lambda option\rangle\) causes a syntax error to be generated at compile time.

#### **SEOUENCE**

The SEQUENCE option causes the character representation of the value stored in the Sequence field of the Station Table to be transmitted if the station is in sequence mode (i.e., this \( \delta \text{it variable} \) SEQUENCE is TRUE); otherwise, the \( \lambda \text{transmit statement} \rangle \) is a no-op.

#### TRAN

The proper number of transmission number characters (as defined by the station's associated \(\lambda\) terminal definition\(\rangle\) in the \(\lambda\) terminal transmission number length statement\(\rangle\)) are extracted from the Transmit Transmission Number field in the Station Table and then transmitted.

(suring)

Each character of the *(string)* is transmitted.

The BREAK syntax allows the programmer to specify action if a "break" is received from the terminal while the adapter cluster is still transmitting. If this option is omitted and a break occurs, an implicit TERMINATE ERROR is executed. The following describes the actions of the three syntactical forms:

BREAK sets TRUE the \( bit variable \) BREAK [TRANSMIT], and causes

a branch of control to the next statement.

BREAK: (label) sets TRUE the (bit variable) BREAK [TRANSMIT], and causes

a branch of control to  $\langle label \rangle$ .

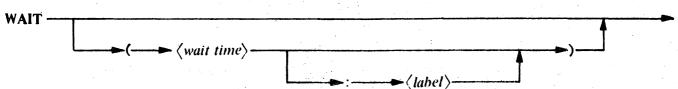
**BREAK:NULL** causes no action. Execution proceeds as if the break did not occur.

# CONTROL

#### **Wait Statement**

#### WAIT STATEMENT





# Examples

WAIT. WAIT (3 SEC). WAIT (5 MILLI:6).

#### **Semantics**

The \(\sqrt{wait statement}\) is only allowed in \(\scrt{control definition}\)\)s that are written to communicate with full duplex terminal types. Execution of this statement causes the \(\scrt{control definition}\)\) to be suspended until the co-line executes a \(\scrt{continue statement}\)\). The optional syntax effects the statement as described below:

(wait time)

defines the maximum amount of \( \text{time} \) that the \( \control \) definition \( \text{should} \) should be suspending waiting for the \( \continue \) statement \( \text{continue} \) is exceeded and the co-line has not executed a \( \continue \) statement \( \text{continue} \), execution resumes at the next sequential statement.

\langle wait time\rangle: \langle label \rangle

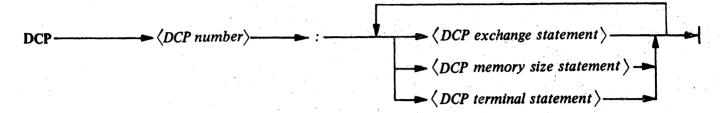
same as above except execution resumes at  $\langle label \rangle$  if a  $\langle continue\ statement \rangle$  is not executed within  $\langle wait\ time \rangle$ .

# **Pragmatics**

Refer to the \( \fork \) statement \( \rightarrow \) pragmatics.

#### **DCP DEFINITION**

**Syntax** 



# Example

DCP 1:
 MEMORY = 8196.
 EXCHANGE = 2.
 TERMINAL = SOMETERMINALNAME.

# **Semantics**

The  $\langle DCP \ definition \rangle$  is the means by which the programmer defines attributes of each Data Communications Processor (DCP) in the Data Communications System.

The  $\langle DCP \ number \rangle$  identifies the DCP and must correspond to an address (ranging from 0 through 7) wired into each DCP by the field engineer. The attributes of the DCP are defined subsequently by means of  $\langle DCP \ statement \rangle s$ . A maximum of eight DCP definitions may appear in the NDL source program.

Each  $\langle DCP statement \rangle$  is described subsequently.

**DCP** 

DCP Exchange Statement

#### **DCP EXCHANGE STATEMENT**

**Syntax** 

 $= \longrightarrow \langle DCP number \rangle$ 

#### Example

EXCHANGE = 4.

#### **Semantics**

The  $\langle DCP \ exchange \ statement \rangle$  specifies that the DCP shares hardware-exchanged adapter clusters with another DCP.  $\langle DCP \ number \rangle$  defines the other DCP.

This statement is required in any  $\langle DCP \ definition \rangle$  referenced by a  $\langle DCP \ exchange \ statement \rangle$  in another  $\langle DCP \ definition \rangle$ , or in any  $\langle DCP \ definition \rangle$  that does not have lines defined for it in the  $\langle line \ definition \rangle$  section of the source program.

# **Pragmatics**

The maximum number of DCPs that can share a set of adapter clusters is 2. The definitions of both DCPs that share adapter clusters must contain a  $\langle DCP \ exchange \ statement \rangle$  naming the  $\langle DCP \ number \rangle$  of the DCP with which it shares the adapter clusters. For example, if DCP 1 and DCP 2 share adapter clusters, then the definition of DCP 1 must contain the statement

# EXCHANGE = 2.

and the definition of DCP 2 must contain the statement

# EXCHANGE = 1.

If a DCP shares adapter clusters with another DCP, then any adapter cluster connected to either of the DCPs must be shared by both. A DCP is not allowed to share only a portion of its adapter clusters.

# LINE SECTION REQUIREMENTS

If two DCPs share adapter clusters, it is required that the  $\langle line\ definition \rangle s$  for each DCP be given addresses (by means of a  $\langle line\ address\ statement \rangle$ ) such that both DCPs do not have lines defined on the same cluster.

The following program segment would cause the compiler to generate a syntax error because both DCPs have lines defined on adapter cluster 0.

#### LINE L100:

ADDRESS = 1:0:0. % ADDRESS =  $\langle DCP \rangle$ : $\langle ADAPTER CLUSTER \rangle$ : $\langle LINE \rangle$ .

LINE L201:

ADDRESS = 2:0:1.

DCP 1:

MEMORY = 8192 EXCHANGE = 2.

DCP 2:

MEMORY = 8192. EXCHANGE = 1.

# MCS RECONFIGURATION

The EXCHANGE CLUSTERS (TYPE = 129) DCWRITE function allows a Message Control System to transfer control of any or all adapter clusters, that are exchanged by two DCPs, from the DCP that currently controls the designated adapter clusters to the DCP with which it is exchanged. This aspect of the reconfiguration feature may be invoked in order to provide an installation with the ability to effect "load-leveling" between two DCPs that share hardware-exchanged adapter clusters or to transfer all of the work load of a DCP to its partner if the DCP malfunctions. For more information regarding reconfiguration, refer to the B 6700/B 7700 DCALGOL Reference Manual, form no. 5000052.

# Supplementary Example

The following is a program segment describing the data communications system illustrated in figure 5-1. This example illustrates how the \( \left( \text{line definition} \right) \) and \( \left( \text{DCP definition} \right) \) sections can be written to describe a data communications system in which two DCPs share hardware-exchanged adapter clusters.

**DCP** 

DCP Exchange Statement - Continued

% %

%STATION DEFINITION SECTION.

%

STATION DEFAULT ALLSTATIONS:

**STATION STA1:** 

DEFAULT = ALLSTATIONS. TERMINAL = TTY.

**STATION STA2:** 

DEFAULT = ALLSTATIONS. TERMINAL = TTY.

**STATION STA3:** 

DEFAULT = ALLSTATIONS. TERMINAL = TTY.

**STATION STA4:** 

DEFAULT = ALLSTATIONS. TERMINAL = TTY.

**STATION STA5:** 

DEFAULT = ALLSTATIONS. TERMINAL = TTY.

**STATION STA6:** 

**DEFAULT = ALLSTATIONS. TERMINAL = TTY.** 

% %

%LINE DEFINITION SECTION.

> ADDRESS = 0:0:0. ADAPTER = 1(DIRECT). STATION = STA1.

LINE LOO1:

ADDRESS = 0:0:1. ADAPTER = 1(DIRECT). STATION = STA2.

# Definitions **DCP**

# DCP Exchange Statement - Continued

# **LINE L020:**

ADDRESS = 0:2:0. ADAPTER = 1(DIRECT).STATION = STA5.

# LINE LO21:

ADDRESS = 0:2:1. ADAPTER = 1(DIRECT).STATION = STA6.

#### 

# **LINE L110:**

ADDRESS = 1:1:0. ADAPTER = 1(DIRECT).STATION = STA3.

# LINE L111:

ADDRESS = 1:1:1. ADAPTER = 1(DIRECT).STATION = STA4.

# %DCP DEFINITION SECTION.

# DCP 0:

MEMORY = 8192.**EXCHANGE** = 1.

# DCP 1:

MEMORY = 8192. EXCHANGE = 0.

# DCP Exchange Statement - Continued

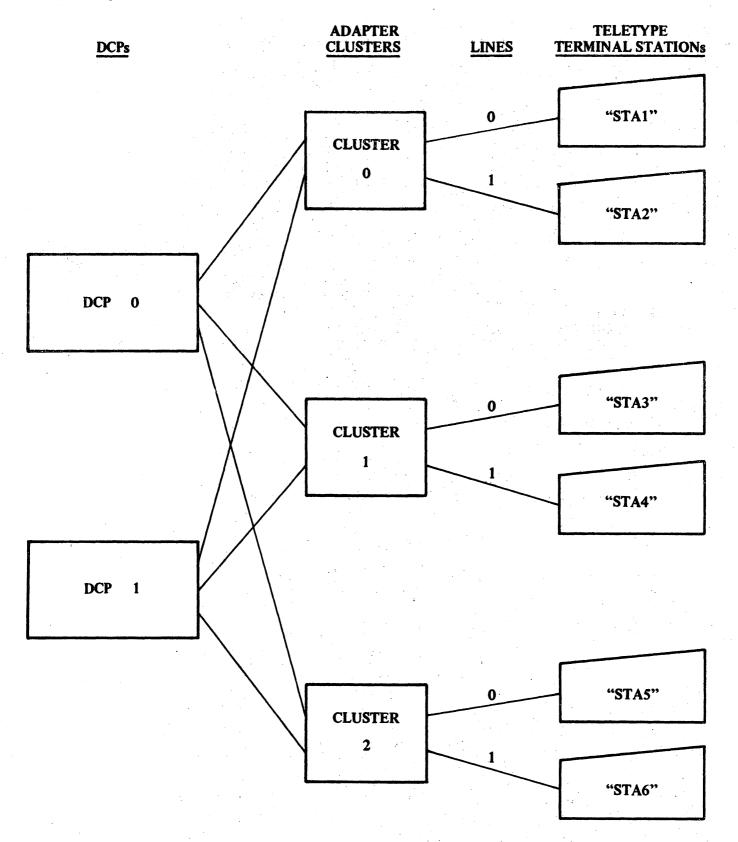


Figure 5-1. Adapter Clusters Exchange

# DCP MEMORY SIZE STATEMENT

**Syntax** 

 $MEMORY \longrightarrow = \longrightarrow \langle integer \rangle$ 

# **Examples**

MEMORY = 4096.MEMORY = 0.

# **Semantics**

The  $\langle DCP \ memory \ size \ statement \rangle$  defines the number of words of local memory in the DCP being defined.

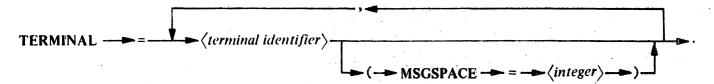
A zero value for (integer) indicates that the DCP has no local memory and that all code generated for the DCP shall reside in main system memory. A non-zero value for (integer) that is less than the amount of local memory required, as determined by the compiler, results in a compile-time error.

**DCP** 

**DCP Terminal Statement** 

# DCP TERMINAL STATEMENT

Syntax



# Examples

TERMINAL = TELETYPE.
TERMINAL = M33, TD800 (MSGSPACE = 5), TELETYPE (MSGSPACE = 2).

#### Semantics

The purpose of the  $\langle DCP \ terminal \ statement \rangle$  is twofold. Each aspect of this statement is discussed in the subsequent two paragraphs.

The primary purpose of the  $\langle DCP \ terminal \ statement \rangle$  is to provide the means of specifying which terminal types in the data communications network that the DCP must be able to control. Only those terminal types specified in this statement will have the object code required to control them included in the object code generated for the DCP. If this statement is omitted from a DCP definition, the compiler includes the object code required to control all terminal types in the data communications network.

The second purpose of the  $\langle DCP \ terminal \ statement \rangle$  is to provide a means of specifying the initial number of message spaces allotted for each terminal type controlled by the DCP.

The \(\lambda terminal identifier\rangle\) must name a terminal type defined by a \(\lambda terminal definition\rangle\), and specifies a terminal type for which the DCP must have access to the controlling code.

The (MSGSPACE = \(\langle integer \rangle \)) option specifies the number of message spaces initially allotted for the terminal type. If this option is omitted, two message spaces are allotted by default.

# **Pragmatics**

Note that if any terminal type is not named in the  $\langle DCP \ terminal \ statement \rangle$ , the data communications network may not be reconfigured (by means of a reconfiguration DCWRITE in an MCS) such that it adds that terminal type to those terminal types controlled by the DCP. Refer to the supplementary example that follows.

# Supplementary Example

The program segment below illustrates the pragmatics. A station whose terminal type is SCREENDEVICE cannot be added on the spare line L003 of DCP 1, because DCP 1 does not have the code available to control SCREENDEVICE.

```
%CONTROL & REQUEST DEFINITION SECTION.
    REQUEST READTTY:
                                               The object code generated from these state-
                                               ments is required to control TTY terminal
    REQUEST WRITETTY:
                                               types.
    REQUEST READSCREENDEVICE:
                                               The object code generated from these state-
                                               ments is required to control SCREENDEVICE
    REQUEST WRITESCREENDEVICE:
                                               terminal types.
%TERMINAL DEFINITION SECTION.
%
    TERMINAL DEFAULT DEFAULTLIST:
        BLOCK
                             FALSE.
        SCREEN
                             FALSE.
        TURNAROUND
                             0.
        ICTDELAY
                             0.
        TRANSMISSION
                             0.
                             FALSE.
        DUPLEX
        TIMEOUT
                             3 SEC.
        ADDRESS
                             0.
        PAGE
                             0.
                             ASC67.
        CODE
        INHIBITSYNC
                             FALSE.
                             NULL.
        BUFFER
        MAXINPUT
                             80.
        WIDTH
                             72.
        PARITY
                             NULL.
        ADAPTER
                             4.
        WRU
                             ENO.
```

ETX(DYNAMIC).

BS(DYNAMIC).

CONTENTION.

**END** 

**BACKSPACE** 

CONTROL

**DCP** 

#### DCP Terminal Statement - Continued

#### TERMINAL TTY:

**DEFAULT** 

DEFAULTLIST.

REQUEST

WRITETTY:TRANSMIT,READTTY:RECEIVE,≺

#### TERMINAL SCREENDEVICE:

**DEFAULT** 

DEFAULTLIST.

**SCREEN** 

= TRUE.

REQUEST

= WRITESCREENDEVICE:TRANSMIT,READSCREENDEVICE: ◄

RECEIVE.

These statements specify the \( \text{request} \)
definitions \( \text{required} \) required to control the defined terminal type. The object code generated by the procedures named here must be accessible by a DCP that has the terminal type attached to any of its lines.

%
%STATION DEFINITION SECTION.
%

# STATION DEFAULT ALLSTATIONS:

ENABLEINPUT = TRUE.

LOGICALACK = FALSE.

MCS = SYSTEM/CANDE.

CONTROL = QM. RETRY = 15.

MYUSE = OUTPUT.INPUT.

**STATION STA1:** 

**DEFAULT** = ALLSTATIONS.

TERMINAL = TTY.

**STATION STA2:** 

**DEFAULT** = ALLSTATIONS.

TERMINAL = TTY.

**STATION STA3:** 

**DEFAULT** = ALLSTATIONS.

TERMINAL = TTY.

**STATION STA4:** 

**DEFAULT** = ALLSTATIONS.

TERMINAL = SCREENDEVICE.

**STATION STA5:** 

**DEFAULT** = ALLSTATIONS.

TERMINAL = SCREENDEVICE.

**STATION STA6:** 

**DEFAULT** = **ALLSTATIONS**.

TERMINAL = SCREENDEVICE.

A \( \langle \text{line definition} \) naming any of these stations must be a \( \langle \text{line definition} \rangle \) for DCP 0. DCP 1 does not have access to code required to control terminals associated with these stations.

A \( \langle \text{line definition} \) naming any of these stations must be a \( \langle \text{line definition} \rangle \) for DCP 1. DCP 0 does not have access to code required to control terminals associated with these stations.

```
%
%LINE DEFINITION SECTION.
                       LINES FOR DCP 0
LINE LOOO:
        ADDRESS
                              0:0:0.
        ADAPTER
                              1(DIRECT).
                              STA1.
        STATION
    LINE LOO1:
                              0:0:1.
        ADDRESS
                          =
                                                  The \langle line\ station\ statement \rangle of any \langle line\ 
        ADAPTER
                          =
                              1(DIRECT).
                                                 definition for DCP 0 must name a station
        STATION
                              STA2.
                                                  that has a TTY terminal type associated with
                                                 it. DCP 0 does not have access to code
    LINE LOO2:
                                                  required to control SCREENDEVICE ter-
        ADDRESS
                              0:0:2.
                                                  minal types.
        ADAPTER
                              1(DIRECT).
        STATION
                              STA3.
                 % THIS IS A SPARE LINE
    LINE LOO3:
        ADDRESS
                              0:0:3.
        MAXSTATIONS
                              1.
                          LINES FOR DCP 1
LINE L100:
        ADDRESS
                              1:0:0.
                              1(DIRECT).
        ADAPTER
        STATION
                              STA4.
    LINE L101:
                                                  The \(\langle \) line station statement \(\rangle \) of any \(\langle \) line
                                                  definition for DCP 1 must name a station
        ADDRESS
                              1:0:1.
                                                  that has a SCREENDEVICE terminal type
                          =
                              1(DIRECT).
        ADAPTER
                                                  associated with it. DCP I does not have
        STATION
                              STA5.
                                                  access to code required to control TTY
    LINE L102:
                                                  terminal types.
        ADDRESS
                              1:0:2.
        ADAPTER
                              1(DIRECT).
        STATION
                              STA6.
%DCP DEFINITION SECTION.
    DCP 0:
                              8192.
        MEMORY
                                                  The effect of this statement is that this DCP
        TERMINAL
                              TTY.
                                                  has access to control code for TTY terminal
                                                  types only.
    DCP 1:
                              8192.
        MEMORY
                                                  The effect of this statement is that this DCP
```

SCREENDEVICE.

**TERMINAL** 

5 - 55

has access to control code for SCREEN-

**DEVICE** terminal types only.

FILE

#### **FILE DEFINITION**

**Syntax** 

FILE \( \sqrt{file identifier} \) \( \sqrt{file family statement} \)

# Example

FILE NETWORK: FAMILY = STATIONID1, STATION ID2, FILEID1.

#### **Semantics**

The  $\langle file\ definition \rangle$  provides the means to define a data communications file and specify the stations associated with that file. The  $\langle file\ identifier \rangle$  is the external name (TITLE) of the file, and has the syntactical form of a  $\langle system\ identifier \rangle$ .

A single-station file is a file that has one station associated with it. A single-station file can, but need not, be formally defined in a  $\langle file\ definition \rangle$ . The reason that a single-station file does not need to be defined is that each station is itself a file. The external name (TITLE) of such a file would be the  $\langle station\ identifier \rangle$  of the station.

A multi-station file is, as the name implies, a file that has more than one station associated with it. Multi-station files must be defined in  $\langle file\ definition \rangle s$ .

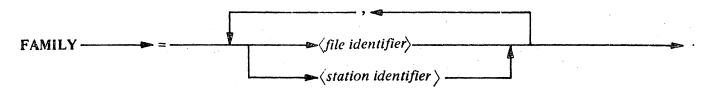
# **Pragmatics**

A general discussion of data communication files and their peculiarities can be found in chapter 2 of the B 6700 Input/Output Subsystem Information Manual, form no. 5000185, under the heading "DATA COMM FILES." The information contained in that discussion is a prerequisite to understanding the significance of \( \frac{file \text{ definition}}{s} \). Chapter 3 of the same manual contains a table that lists all file attributes and provides an explanation of each attribute. Attributes relative to data communication files are found by examining the "KIND" column of the table for the key word "Datacom." The information found in the explanation of each data communications-relative attribute is also a prerequisite.

A detailed discussion of data communications object job I/O can be found in appendix B of the <u>B 6700/B 7700 DCALCOL</u> Reference Manual, form no. 5000052, under the semantics of STATION ASSIGNMENT TO FILE (TYPE = 64). The information found there is not considered a prerequisite; however, it does contribute toward a deeper understanding of data communications files and data communications object job I/O.

#### FILE FAMILY STATEMENT

**Syntax** 



# Example

FAMILY = STATIONID1, STATIONID2, FILEID1.

#### **Semantics**

The  $\langle file\ family\ statement \rangle$  defines the stations associated with a data communications file. If a  $\langle file\ identifier \rangle$  is named, all of the stations associated with the file named will also be associated with the file being defined. Any duplication of an  $\langle identifier \rangle$  in a  $\langle file\ family\ statement \rangle$  is ignored.

# Supplementary Example

The following example is the (file definition) section of a hypothetical NDL program. Assume that the stations STATION1, STATION2, STATION3, STATION4, STATION5, STATION6, STATION7, and STATION8 have been defined in the (station definition) section.

#### **FILE TTYS:**

FAMILY = STATION1, STATION2, STATION3.

# **FILE CRTS**:

FAMILY = STATION4, STATION5, STATION6.

# **FILE EXECUTIVES:**

FAMILY = STATION1, CRTS, STATION6, STATION 7.

# FILE THE/ENTIRE/NETWORK:

FAMILY = STATION, STATIONS, TTYS, CRTS.

TTYS is the (identifier) of this file. The FAMILYSIZE is 3. STATION1, STATION2, and STATION3 are the stations associated with this file.

CRTS also has a FAMILYSIZE of 3. The stations associated with the file are STATION4, STATION5, and STATION6.

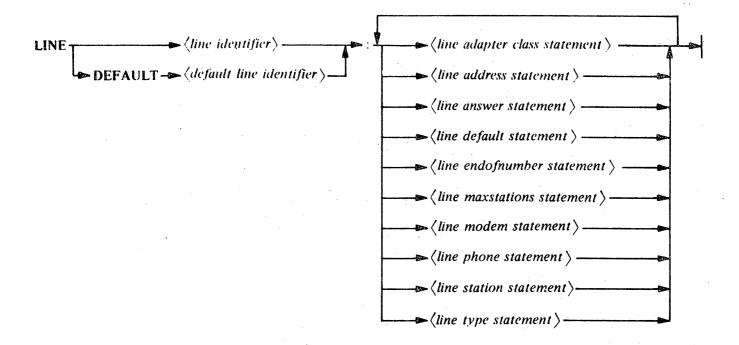
EXECUTIVES has a FAMILY-SIZE of 5. The stations associated with the file are STATION1, STATION4, STATION5, STATION6, and STATION7.

THE/ENTIRE/NETWORK has a FAMILYSIZE of 8. The stations associated with THE/ENTIRE/NETWORK are STATION1, STATION2, STATION3, STATION4, STATION5, STATION6, STATION7, and STATION8.

LINE

# LINE DEFINITION

# **Syntax**



# Examples

# LINE TTYDIALIN:

TYPE = DIALIN.
ADAPTER = 1 (MODEM).
MODEM = TTYMODEM.
ANSWER = TRUE.

PHONE = 2139686521.

ADDRESS = 0:0:0.

STATION = TTYSTATION.

MAXSTATIONS = 1.

# LINE DEFAULT LINEDEFAULTLIST1:

ADAPTER = 1 (MODEM).

ANSWER = TRUE. ENDOFNUMBER = FALSE.

MAXSTATIONS = 1.

TYPE = DIALIN.

MODEM = TTYMODEM.

# **Semantics**

 $\langle line\ identifier \rangle$  and  $\langle default\ line\ identifier \rangle$  both have the same syntactical form as  $\langle identifier \rangle$ . Each form of the  $\langle line\ definition \rangle$  syntax is described subsequently.

# **LINE** \(\langle \text{line identifier} \rangle : \tau. \tag{...}

This form of the (line definition) defines the attributes of a logical line in the data communications network. Line attributes are defined in one of the following ways:

- a. Each attribute is defined explicitly by means of a  $\langle line \ statement \rangle$  in the  $\langle line \ definition \rangle$ .
- b. Each attribute is defined implicitly by an explicit reference to a set of default attribute values.
- c. Some of the line attributes are defined implicitly as in b, and the remainder are defined explictly as in a.

Some \(\langle \text{line statement} \rangle \sin \text{must be defined for each \(\langle \text{line definition} \rangle \); others do not. Some of the statements may or may not require definition, depending upon the appearance of other statements. The semantics portion of each \(\langle \text{line statement} \rangle \) states, among other things, whether the attribute must be defined and its effect upon the requirements of other attribute definitions.

To define the attributes of a line as described in item a above, this syntax form must be used.

To define the attributes of a line as described in items b and c above, this syntax form, the following syntax form, and the \( \line \text{default statement} \) must be used in conjunction (this is described under the following syntax form).

# LINE DEFAULT (default line identifier): . . .

This syntax form is referred to as a Default  $\langle line \ definition \rangle$ . Its purpose is to decrease the number of source statements required to define all of the logical lines in the data communications system. This is accomplished in the following manner. Attributes common to several logical lines are defined by means of a Default  $\langle line \ definition \rangle$ . Associated with each Default  $\langle line \ definition \rangle$  is a  $\langle default \ line \ identifier \rangle$ . Subsequent to the Default  $\langle line \ definition \rangle$ , any  $\langle line \ definition \rangle$  that has those attributes in common can reference the  $\langle default \ line \ identifier \rangle$ , instead of repeating the list. (A  $\langle default \ line \ identifier \rangle$  is referenced by means of a  $\langle line \ default \ statement \rangle$ .) The NDL compiler uses the last definition of a line attribute, and therefore the programmer can reference a Default  $\langle line \ definition \rangle$  and change any attributes by redefining them in the  $\langle line \ definition \rangle$ .

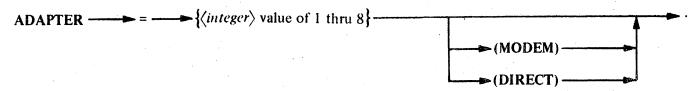
In appearance, the Default (line definition) is similiar to the (line definition). The differences are that the reserved word DEFAULT follows the reserved word LINE, and that there are no statements that are required to be defined in a Default (line definition).

#### LINE

Line Adapter Class Statement

#### LINE ADAPTER CLASS STATEMENT

Syntax



# Examples

ADAPTER = 5.

ADAPTER = 4 (MODEM).

#### **Semantics**

The \(\langle \text{line adapter class statement}\rangle\) identifies the Adapter Class of the line adapter for the logical line and, optionally, names the connection type (i.e., modem connect or direct connect).

The Adapter Class must be compatible with the \( \chiconnumication type number \) specified in the \( \station \) adapter statement \( \rightarrow \) of any station assigned to the line. (Note that all stations assigned to a line must have the same \( \chiconnumication type number \rangle \) defined.) Table 5-3 lists the compatible Adapter Classes for each \( \chiconnumication type number \rangle \). For example, a line having stations assigned to it that define a \( \chiconnumication type number \rangle \) of 4 can name as an Adapter Class either 1, 2, 3, 4, or 5 (refer to table 5-3). On the other hand, a line having stations assigned to it that defines 15 as the \( \chiconnumication type number \rangle \) can name only 5 as an Adapter Class.

If the connection type is named in the statement, it is considered by the compiler as documentation only. The compiler determines whether the line adapter or a modem-connect line adapter, by the presence or absence of a  $\langle line\ modem\ statement \rangle$  for the  $\langle line\ definition \rangle$ . A syntax error is generated, however, if **DIRECT** is named and a  $\langle line\ modem\ statement \rangle$  is present.

# **Pragmatics**

#### LINE ADAPTERS AND ADAPTER CLASSES

There are 13 available line adapters. Three of the 13 are special-purpose line adapters; they are used for Touch-Tone<sup>(R)</sup> telephone input, Audio-Response lines, and Automatic Calling Units (ACU). The remaining 10 are general-purpose line adapters.

The 13 available line adapters are divided into eight "Adapter Classes." The Touch-Tone Addio-Response, and ACU line adapters comprise Adapter Classes 6, 7, and 8, respectively. The 10 general-purpose line adapters comprise Adapter Classes 1 through 5. Adapter Classes 1 through 5 differ primarily in the maximum transmission speed at which the line adapters may be operated. Adapter Classes 1 through 5 each consist of two line adapters, one being a "direct connect," and the other being a "modem connect." The direct connect has a terminal attached to it by means of a two-wire or four-wire direct connection. The modem connect has a terminal attached to it through modems using an RS232† defined interface. Refer to table 5-3 for the Adapter Class and the use of each line adapter.

<sup>†</sup>A technical specification published by the Electronic Industries Association establishing the interface requirements between modems and terminals or computers.

Table 5-3. Available Line Adapters

MARKETING NUMBER*	CONNECTION	USE	CLASS
В 6650-1	Direct	Two-wire direct connect, asynchronous bit-serial transmission up to a maximum line speed of 600 BPS, simplex or half-duplex.	1
В 6650-1	Modem	Modem-connected with 100-Series type modem using RS232-defined interface, asynchronous bit-serial transmission up to a maximum line speed of 600 BPS, simplex or half-duplex. (Two required for full duplex.)	1
В 6650-2	Direct	Same as B 6650-1D, except maximum line speed is 1800 BPS.	2
В 6650-2	Modem	Same as B 6650-1M, except with 202- Series type modem and up to a maxi- mum line speed of 1800 BPS.	2
В 6650-3	Direct	Same as B 6650-1D, except maximum line speed is 2400 BPS.	3
В 6650-3	Modem	Modem-connected with 202-Series (asynchronous) or 201-Series (synchronous) modem using RS232-defined interface, bit-serial transmission up to a maximum line speed of 2400 BPS, simplex or half-duplex. (Two required for full duplex.)	3
В 6650-4	Direct	Same as B 6650-1D, except maximum line speed is 4800 BPS.	4
B 6650-4	Modem	Same as B 6650-3M, except maximum line speed is 4800 BPS.	4
В 6650-5	Direct	Same as B 6650-1D, except maximum line speed is 9600 BPS.	5
В 6650-5	Modem	Same as B 6650-3M, except maximum line speed is 9600 BPS.	5
В 6650-6		For Touch-Tone <sup>®</sup> telephone input.	. 6
В 6650-7		For Audio-Response line.	7
В 6650-8	1	For Automatic Calling Unit (ACU).	8 .

<sup>\*</sup>The above marketing numbers refer to B 6700 line adapters. B 7700 line adapters have similar numbers, the difference being a leading 7 instead of a 6; e.g., B 6650-1 for B 6700, and B 7650-1 for B 7700. In this table only, a distinction is made between modem-connected line adapters and direct-connected line adapters by affixing either a D (for direct-connected) or an M (for modem-connected) to the field marketing numbers (under the "Use" column) which require that distinction.

LINE

Line Address Statement

#### LINE ADDRESS STATEMENT

Syntax

 $ADDRESS \longrightarrow = \longrightarrow \langle DCP \ number \rangle \longrightarrow : \longrightarrow \langle adapter \ cluster \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ adapter \ number \rangle \longrightarrow : \longrightarrow \langle line \ n$ 

# Example

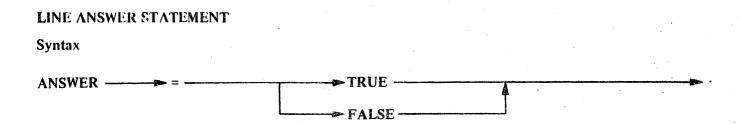
ADDRESS = 2:0:15.

The above example would appear in the  $\langle line \ definition \rangle$  of the line at the 15th line adapter position in adapter cluster number 0 of DCP number 2.

#### **Semantics**

The  $\langle line \ address \ statement \rangle$  identifies the DCP number, the adapter cluster number, and the line adapter number of the defined logical line. If two DCPs share hardware-exchanged adapter clusters (as defined by the  $\langle DCP \ exchange \ statement \rangle$  in a  $\langle DCP \ definition \rangle$ ), then the  $\langle DCP \ number \rangle$  defined in this statement is the DCP initially expected to service the adapter cluster of which the line is a part. This statement, which is required, must be defined explicitly in each  $\langle line \ definition \rangle$ .

Line Answer Statement



#### **Semantics**

The  $\langle line\ answer\ statement \rangle$  defines whether or not (TRUE or FALSE, respectively) the DCP is to automatically answer an incoming call. This statement is required if the  $\langle line\ type\ statement \rangle$  in the  $\langle line\ definition \rangle$  defines the line configuration as DIALIN only, or DIALIN and DIALOUT.

If ANSWER = FALSE, an incoming call causes the following actions to be taken by the DCP. A SWITCHED STATUS RESULT (CLASS = 7) message is sent to the MCS of the station that is the first entry in the Line Table for that line. (Unless an MCS has reconfigured the line so that it changes the first entry, the first entry in the Line Table will be the entry for the first station listed in the \( \line \text{station statement} \rangle \) of the \( \line \text{dine definition} \rangle \).) The message has a bit set in it that indicates the line is in a "ringing" status. Presumably, upon notification of a line in a ringing status, the MCS programmer instructs the DCP to answer the phone, or it takes appropriate action to clear the line.

If ANSWER = TRUE, an incoming call causes the DCP to take the following actions. A SWITCHED STATUS RESULT (CLASS = 7) message is sent to the controlling MCS of the station that is the first entry in the Line Table for that line. In this case the message has a bit set indicating that there has been an incoming call, and that the DCP is in the process of answering the call.

An MCS may change the value of ANSWER after DCP initialization, by means of a SET/RESET AUTO-ANSWER (TYPE = 102) DCWRITE.

#### LINE

Line Default Statement

#### LINE DEFAULT STATEMENT

Syntax

# Example

**DEFAULT = DFLTLIST1.** 

#### **Semantics**

The  $\langle line\ default\ statement \rangle$  allows the programmer to specify the  $\langle default\ line\ identifier \rangle$  of a set of default line attributes to be used for a  $\langle line\ definition \rangle$  whose description is incomplete. It is advantageous to group attributes that several lines have in common under a Default  $\langle line\ definition \rangle$  and list the remaining attributes under each individual  $\langle line\ definition \rangle$ . The compiler will then refer to the Default  $\langle line\ definition \rangle$  to complete the  $\langle line\ definition \rangle$ . The  $\langle line\ default\ statement \rangle$  is not required to appear in a  $\langle line\ definition \rangle$ ; however, a  $\langle line\ definition \rangle$  must define all required attributes if a  $\langle line\ default\ statement \rangle$  does not appear.

The  $\langle line\ default\ statement \rangle$  can appear in a  $\langle line\ definition \rangle$  or a Default  $\langle line\ definition \rangle$ . Thus,  $\langle line\ default\ statement \rangle s$  can be "nested" to combine the attributes of one or more Default  $\langle line\ definitions \rangle s$ .

Line Endofnumber Statement

LINE	<b>ENDOF</b>	NUMBER	STATEMENT

**Syntax** 

ENDOFNUMBER =	→TRUE	
LI (DOI I (ONIDE)	•	
i	<b>───</b> FALSE ───	

# **Semantics**

The  $\langle line\ endofnumber\ statement \rangle$  applies only to  $\langle line\ definition \rangle s$  that specify the Automatic Calling Unit (ACU) Adapter Class in its  $\langle line\ adapter\ class\ statement \rangle$  (e.g., ADAPTER = 8). This statement is required for those  $\langle line\ definition \rangle s$ , and specifies whether or not (TRUE or FALSE, respectively) the ACU has an "end of number" option.

LINE

Line Maxstations Statement

LINE MAXSTATIONS STATEMENT

Syntax

MAXSTATIONS — = \( \sqrt{integer} \)

Example

MAXSTATIONS = 25.

#### **Semantics**

The  $\langle line\ maxstations\ statement \rangle$  specifies the number of stations that may be assigned to the defined line. If this statement does not appear for a line having assigned stations (the  $\langle line\ station\ statement \rangle$  lists all stations initially assigned to a line), it is assumed that MAXSTATIONS is the number of stations explicitly specified as assigned to the line in the  $\langle line\ station\ statement \rangle$  The  $\langle integer \rangle$  specified must not equal 0, exceed 255, or be less than the number of stations listed in the  $\langle line\ station\ statement \rangle$  (if the  $\langle line\ station\ statement \rangle$  is defined).

# **Pragmatics**

This statement informs the compiler of the maximum number of station descriptors required in the Line Table of the DCP's table structure. By defining MAXSTATIONS to be greater than the number of stations listed in the \( \langle \text{line station statement} \), an MCS may reconfigure more stations onto the line at some point in time after DCP initialization. For information regarding reconfiguration, refer to the \( \frac{B 6700/B 7700}{B 6700/B 7700} \)
DCALGOL Reference Manual, form number 5000052.

LINE MODEM STATEMENT			
Syntax			
MODEM	(modem id	lentifier\	

# Example

MODEM = BELL103A.

# **Semantics**

The \langle line modem statement \rangle specifies the modem type that exists on the system end of the physical line. (The \langle station modem statement \rangle in a \langle station definition \rangle specifies the modem type connected to the line on the terminal end.)

# **Pragmatics**

The compiler references other portions of the program with this statement, checking for consistency. If, for example, the \( \)modem definition\( \) of the \( \)modem identifier\( \) specified in this statement lists any \( \)communication type number\( \)s in its \( \)modem adapter statement\( \) that are not compatible with the Adapter Class specified in the \( \)line adapter class statement\( \) of the \( \)line definition\( \), then a syntax error is generated. Another situation that causes a syntax error to be generated is if the compiler discovers that the modem type specified in this statement is not compatible, in respect to the \( \)communication type number\( \), with the modem type specified in the \( \)station definition\( \)\) of a station assigned to the line.

Definitions LINE

Line Phone Statement

LINE PHONE STATEMENT

**Syntax** 

Example

PHONE = 12136572385.

**Semantics** 

The  $\langle line\ phone\ statement \rangle$ , implemented for documentation purposes only, documents the telephone number of a **DIALIN** type line. This statement is optional in a  $\langle line\ definition \rangle$ .

#### LINE STATION STATEMENT

**Syntax** 



#### **Examples**

STATION = RJE1. STATION = DAKOTA/KID, BIDS.

#### **Semantics**

The \(\langle \text{line station statement}\rangle\) is the means by which the NDL programmer associates one or more stations with a line. A station that is associated with a particular line is said to be "assigned" to that line.

This statement is required in those  $\langle line\ definition \rangle$  s that specify **DUPLEX** in the  $\langle line\ type\ statement \rangle$ . In all other variations of  $\langle line\ type\ statement \rangle$ , this statement is optional.

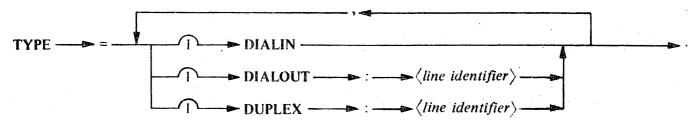
If more than one station is named, each station must have the same (communication type number) defined in its respective (station adapter statement).

LINE

Line Type Statement

#### LINE TYPE STATEMENT

**Syntax** 



#### **Examples**

TYPE = DIALIN.

**TYPE = DIALOUT: ACULINE.** 

TYPE = DUPLEX: AUXLINE.

TYPE = DIALIN, DIALOUT: AUTOCALL, DUPLEX: SUPERVISORY.

#### **Semantics**

The  $\langle line\ type\ statement \rangle$  provides the compiler with specific information concerning special logical line configurations. This statement is required for  $\langle line\ definition \rangle s$  whose line utilize either dial-in, dial-out, or full duplex hardware facilities.

#### DIALIN

This form identifies the line as a dial-in line. A line that may be dialed from a remote site is a dial-in line. The appropriate  $\langle line\ type\ statement \rangle$  for this configuration would be:

#### TYPE=DIALIN.

A logical line defined in this manner must include the  $\langle line \ answer \ statement \rangle$  and the  $\langle line \ modem \ statement \rangle$ . The  $\langle line \ definition \rangle$  for such a line could appear as follows:

#### LINE DIALUPLINE:

TYPE = DIALIN.

ADDRESS = 0:0:0.

MODEM = TTY103A.

**STATION** = DIALUPSTATION.

**ANSWER** = TRUE.

ADAPTER = 1(MODEM).

#### DIALOUT

This form identifies the line as a dial-out line. A dial-out line is defined as a line that can become connected to a remote site as a result of a Message Control System issuing a DIALOUT (TYPE = 98) DCWRITE to the line (thereby causing an Automatic Calling Unit (ACU) to dial the phone number of the remote site). The TYPE=DIALOUT: \langle line identifier \rangle syntax of the statement specifies such a configuration. The \langle line identifier \rangle names the \langle line definition \rangle that defines the associated ACU. The following example illustrates how the \langle line definition \rangle s could appear for a dial-out configuration.

#### LINE DIALOUTLINE:

**ADAPTER** 

TYPE = DIALOUT: ACULINE.

= 1(MODEM).

**ADDRESS** = 0:0:1. **MODEM** = TTY103A.

#### LINE

Line Type Statement - Continued

#### LINE ACULINE:

ENDOFNUMBER = FALSE. ADDRESS = 0:0:5. ADAPTER = 8.

The \(\langle \text{line definition}\rangle\) for the dial-out line must include a \(\langle \text{line modem statement}\rangle\) and cannot include a \(\langle \text{line modem statement}\rangle\). The \(\langle \text{line definition}\rangle\) for the ACU must include a \(\langle \text{line endofnumber statement}\rangle\), and it must define an address (in the \(\langle \text{line address statement}\rangle\)) that is on the same adapter cluster as the associated dial-out line.

#### **DUPLEX**

This form identifies the line as the primary of a line pair, for purposes of simultaneous transmission and receptions. The  $\langle line \ identifier \rangle$  names the auxiliary line's  $\langle line \ definition \rangle$ . The line referenced as the auxiliary cannot contain a  $\langle line \ type \ statement \rangle$  nor a  $\langle line \ statement \rangle$ .

The following is an example of how full duplex primary and auxiliary lines could be defined.

#### LINE DUPLEXPRIMARY:

TYPE = DUPLEX:DUPLEXAUXILIARY.
ADDRESS = 0:0:5.
MODEM = SUPERMODEM.
STATION = MODEL37.
ADAPTER = 1.

#### LINE DUPLEXAUXILIARY:

**ADDRESS** = 0:0:6. **ADAPTER** = 1.

#### **Pragmatics**

#### COMBINED CONFIGURATIONS

A dial-in/dial-out line is characterized by both the ability to be dialed from a remote site, and the ability to become connected to a remote site as a result of a Message Control System issuing a DIALOUT (TYPE = 98) DCWRITE. This type of configuration requires the DIALIN and DIALOUT: \(\langle \text{line identifier} \rangle\) options to appear in the \(\langle \text{line type statement} \rangle\). The following example illustrates how a dial-in and dial-out \(\langle \text{line definition} \rangle\) could appear:

#### LINE IOLINE:

TYPE = DIALIN, DIALOUT: AUTOCALLUNIT.

= 1(MODEM).

ADDRESS = 0:1:0.

MODEM = TTY103A.

STATION = REMOTETTY.

ANSWER = TRUE.

#### LINE AUTOCALLUNIT:

**ADAPTER** 

ENDOFNUMBER = FALSE. ADDRESS = 0:1:1. ADAPTER = 8.

#### LINE

#### Line Type Statement - Continued

The full duplex syntax could be combined with the dial-in and dial-out syntax as follows:

#### LINE IODUPLEX:

TYPE = DIALIN, DIALOUT: AUTOCALLUNIT, DUPLEX: AUXLINE.

ADDRESS = 0:2:0.

MODEM = SUPERMODEM.

**STATION** = **REMOTEDUPLEXDEVICE**.

ANSWER = TRUE. ADAPTER = 1(MODEM).

LINE AUXLINE:

ADDRESS = 0:2:1.

ADAPTER = 1(MODEM).

LINE AUTOCALLUNIT:

ENDOFNUMBER = FALSE. ADDRESS = 0:2:2.

ADDRESS = 0:2ADAPTER = 8.

#### MCS DEFINITION

#### **Syntax**

MCS — 
$$\langle$$
 MCS identifier  $\rangle$  —  $\rangle$  CONTROL —  $\rangle$  FALSE

#### **Examples**

MCS SYSTEM/CANDE:CONTROL = FALSE.
MCS SYSTEM/DIAGNOSTICMCS:CONTROL = TRUE.

#### Semantics

The purpose of the  $\langle MCS \ definition \rangle$  is twofold: First, the  $\langle MCS \ definition \rangle$  adds the  $\langle MCS \ identifier \rangle$  to the list (contained in the Network Information File) of valid Message Control System (MCS) programs; and second, the  $\langle MCS \ definition \rangle$  specifies whether or not (CONTROL = TRUE, or CONTROL = FALSE, respectively) the named MCS is allowed to execute a limited set of DCWRITE functions that perform DCP diagnostic functions in addition to the standard DCWRITEs.  $\langle MCS \ identifier \rangle$  has the syntactic form of a  $\langle system \ identifier \rangle$ .

#### **Pragmatics**

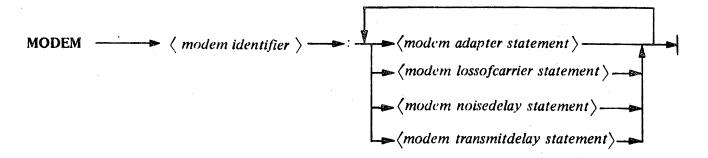
A list of valid MCSs is maintained in the Network Information File in order to restrict unauthorized DCALGOL programs from becoming an MCS. (A DCALGOL program becomes an MCS when it successfully executes an INITIALIZE PRIMARY QUEUE (TYPE = 0) DCWRITE.) The MCS declaration is one means of adding a name of an MCS to that list. (One other means is the \(\sigma \text{station MCS statement}\) in a \(\sigma \text{station definition}\).)

The diagnostic DCWRITE functions allow an MCS to perform on-line tests of components in the Data Communications System. Those DCWRITEs that may be utilized in an MCS when CONTROL = TRUE have DCWRITE TYPE numbers greater than 159.

**MODEM** 

#### **MODEM DEFINITION**

**Syntax** 



#### Example

**MODEM MABELL103A:** 

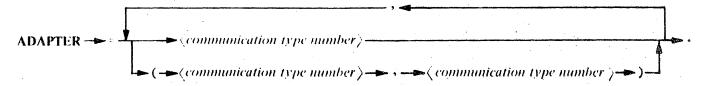
ADAPTER = 4.
LOSSOFCARRIER = DISCONNECT.
NOISEDELAY = 0.
TRANSMITDELAY = 0.

#### **Semantics**

The \(\text{modem definition}\) defines the attributes of a modem type in the data communications system. The \(\text{modem identifier}\) names the \(\text{modem definition}\), and has the syntactic form of \(\text{identifier}\). The \(\text{modem statement}\)'s are described subsequently.

#### MODEM ADAPTER STATEMENT

#### **Syntax**



#### **Examples**

ADAPTER = 1,2,3,4. ADAPTER = 10. ADAPTER = (2,3), (4,5), 6,7.

#### **Semantics**

The *(modem adapter statement)* defines one or more combinations of character format, synchronous/ asynchronous communication, and line speed (in the case of asynchronous communications) with which the modem is compatible. This is done by supplying one or more *(communication type number)* s (or number pairs).

Table 5-4 lists the allowed  $\langle communication \ type \ number \rangle s$  and the characteristics associated with each. For example, the statement

ADAPTER 
$$= 4$$
.

defines an 11-bit character format, asynchronous communication, at a line speed of 110 bits per second.

If the modem is to be used in a full duplex mode, and the primary and auxiliary lines have different characteristics, then one or more (communication type number) pairs must be supplied. For example, the statement

$$ADAPTER = (11,6).$$

defines for the primary line a 10-bit character format, synchronous communication, at a speed of 1800 bits per second. The characteristics associated with the auxiliary line are the same as for the primary line, except that the auxiliary line runs at a line speed of 150 bits per second.

#### **Pragmatics**

#### **COMMUNICATION TYPE NUMBERS**

A (communication type number) is an integer that has associated with it a set of attributes that define three line characteristics. Those characteristics are the format of the characters transmitted (start information, data information, parity information, and stop information), whether the line is to be driven synchronously or asynchronously, and the speed of the transmissions (in the case of asynchronous communications). Table 5-4 lists the allowed (communication type number)s and the line characteristics associated with each.

Most of the electronics that directly control a line are located in the adapter cluster that contains the line adapter for that line (rather than being located in the line adapter itself). The adapter cluster is somewhat general purpose in its design in that it can run at various line speeds and handle various character formats. The DCP can cause the adapter cluster to function in a suitably special-purpose way (with respect to a single line) by supplying it a number derived from the (communication type number).

#### **MODEM**

#### Modem Adapter Statement - Continued

There are three areas in an NDL program that require the programmer to supply one or more (communication type number)s:

- a. In the (modern adapter statement) of each (modern definition),
- b. In the \(\lambda terminal adapter statement\rangle\) for each \(\lambda terminal definition\rangle\), and
- c. In the \(\station\) adapter statement\(\rangle\) for each \(\station\) definition\(\rangle\).

As it encounters each area, the NDL compiler cross-checks to determine if the areas are compatible in their description. If inconsistencies in component compatibility arise, syntax errors are generated. Restrictions are described in the \(\lambda terminal adapter statement\rangle\) and \(\lambda station adapter statement\rangle\) semantics.

#### **EXPLANATION OF TABLE 5-4**

Table 5-4 lists the allowed (communication type number)s in the column labeled "COMM. TYPE NUM." To the right of each (communication type number) are the three line characteristics associated with it, under the columns labeled "SPEED (BPS)," "CHARACTER FORMAT," and "SYNCHRONOUS OR ASYNCHRONOUS." The rightmost column, labeled "COMPATIBLE ADAPTER CLASSES," is referenced and described in the (line definition) section of this chapter.

Table 5-4. Table of  $\langle communication \ type \ number \rangle s$ 

			CHARACTER FORMAT							
COMM. TYPE NUM.	SPEED (BPS)	CHAR. SIZE (BITS)	START INFO.	DATA INFO.	PARITY INFO.	STOP INFO.	SYNCHRONOUS OR ASYNCHRONOUS	COMPATIBLE ADAPTER CLASSES 1 2 3 4 5 6 7 8		
1	45.5	7.5	1	5.0		1.5	ASYNC.	xxxxx		
2	56.9	7.5	1	5.0		1.5	ASYNC.	XXXXX		
3	75.0	7.5	1	5.0		1.5	ASYNC.	XXXXX		
4	110.0	11.0	1	7.0	1	2	ASYNC.	XXXXX		
5	134.5	9.0	1	6.0	1	1	ASYNC.	XXXXX		
6	150.0	10.0	1	7.0	1	1	ASYNC.	xxxxx		
7	300.0	10.0	1	7.0	1	1	ASYNC.	XXXXX		
8	600.0	10.0	1	7.0	1	1	ASYNC.	xxxxx		
9	1200.0	10.0	1	7.0	1	1	ASYNC.	XXXX		
10	1200.0	6.0	1	4.0		1	ASYNC.	XXXX		
11	1800.0	10.0	1	7.0	1	1	ASYNC.	xxxx		
12	2400.0	10.0	1	7.0	1	1	ASYNC.	XXX		
13	3600.0	10.0	1	7.0	1	1	ASYNC.	X X		
14	4800.0	10.0	1	7.0	1	1	ASYNC.	ХX		
15	9600.0	10.0	1	7.0	1	1	ASYNC.	X		
16	2000.0	7.0		6.0	1		SYNC.	xxx		
17	2000.0	8.0		7.0	1		SYNC.	XXX		
18	2000.0	9.0		8.0	1		SYNC.	XXX		
19	2400.0	7.0		6.0	1		SYNC.	XXX		
20	2400.0	8.0		7.0	1		SYNC.	XXX		
21	2400.0	9.0		8.0	1	<u>-</u> _	SYNC.	XXX		
22	4800.0	7.0		6.0	1	<del></del>	SYNC.	хх		
23	4800.0	8.0	<u></u> -	7.0	1 .		SYNC.	хх		
24	4800.0	9.0		8.0	1		SYNC.	X X		
25	9600.0	7.0		6.0	1		SYNC.	X		
26	9600.0	8.0		7.0	1		SYNC.	X		
27	9600.0	9.0		8.0	1		SYNC.	X		
28	40.0	4.0	<u></u> -	4.0			SYNC.	X		
29	16.0	8.0		7.0	1		SYNC.	X		
30	40.0	4.0		4.0			SYNC.	X		

#### **MODEM**

Modem Lossofcarrier Statement

#### MODEM LOSSOFCARRIER STATEMENT

**Syntax** 

#### **Pragmatics**

Certain modems (Western Electric (Bell System) 103 series modems, and possibly others) maintain continuous carrier in both directions while the line is properly connected. As such, CF (Carrier Detected) and CB (Clear to Send) are maintained TRUE while connected. Additionally, if each modem is equipped with both the Initiate Disconnect and the Respond to Disconnect options, each modem employs the "long space disconnect" convention. This convention allows one modem to determine if the other is disconnecting, and itself go "on-hook" and drop CC (Data Set Ready).

Two problems arise, however, when only one such modem is configured at the system end, and the terminal is interfaced with an acoustic coupler at the terminal end. At the time of making a connection, establishment of carrier is difficult. In fact, the system modem may detect carrier from the coupler while the telephone receiver is near the coupler and before the receiver is properly seated. In this case, CF and CB are raised prematurely, and if the system takes this as a cue to begin transmission of a greeting, the two signals (the data transmitted from the system, and carrier from the acoustic coupler) interact with each other, and the system modem detects loss of carrier. At the time of terminating a call, if the terminal initiates the disconnect and has no "long space disconnect" facility, or if the terminal operator does not use it, the system modem detects only loss of carrier. In this case, the system modem drops CF and CB, the modem remains "off-hook" and maintains CC (Data Set Ready). Thereafter, any incoming calls would receive a "busy" signal.

The \(\langle modem \) loss of carrier statement\(\rangle\) is implemented for such a configuration. If this statement is included in the definition of a modem, special logic is invoked, in addition to the normal logic, when dealing with that modem type.

In the case of the system calling out, normal logic waits for CC to be raised by the modem. If CC is raised within 25 seconds, the line is immediately released as connected. A timeout of 25 seconds causes CD (Data Terminal Ready) to be dropped, the modem goes "on-hook," and the line reverts to a disconnected state. The special logic is invoked after CC is found TRUE. With the 25-second timeout in effect, the special logic then waits until CF and CB are both raised by the modem. After CF and CB are detected, the logic then delays approximately 5 seconds before notifying the system that the line is connected. This gives the terminal operator sufficient time to place the receiver in the acoustic coupler.

In the case of a terminal-initiated disconnect, that condition is detected in the normal logic by either the "long space disconnect" adapter cluster interrupt or by a CC (Data Set Ready) FALSE condition. In addition to the normal logic, the special logic also interprets CF FALSE or CB FALSE as a terminal-initiated disconnect.

### Definitions **MODEM**

Modem Noisedelay Statement

#### MODEM NOISEDELAY STATEMENT

**Syntax** 

NOISEDELAY → = → ⟨ delay time ⟩ -----

#### **Examples**

NOISEDELAY = 0. NOISEDELAY = 200 MILLI.

#### **Semantics**

The \langle modem noisedelay statement \rangle defines the amount of time that should be delayed when the modem enters a Clear to Send (CB) status to avoid receiving "noise" on the line. \langle delay time \rangle must be expressed as \langle time \rangle, and affects the amount of time delayed after an INITIATE RECEIVE or INITIATE TRANSMIT construct is executed and before the next statement is executed in a \langle control definition \rangle or \langle request \definition \rangle. The \langle delay time \rangle defined in this statement is used in a compiler algorithm that calculates the delay. The compiler algorithm is discussed in the semantics of the INITIATE RECEIVE and INITIATE TRANSMIT constructs under the \langle initiate statement \rangle. This statement must appear in each \langle modem \definition \rangle.

#### **MODEM**

Modem Transmitdelay Statement

#### MODEM TRANSMITDELAY STATEMENT

**Syntax** 

TRANSMITDELAY = \( \sqrt{delay time} \)

#### **Examples**

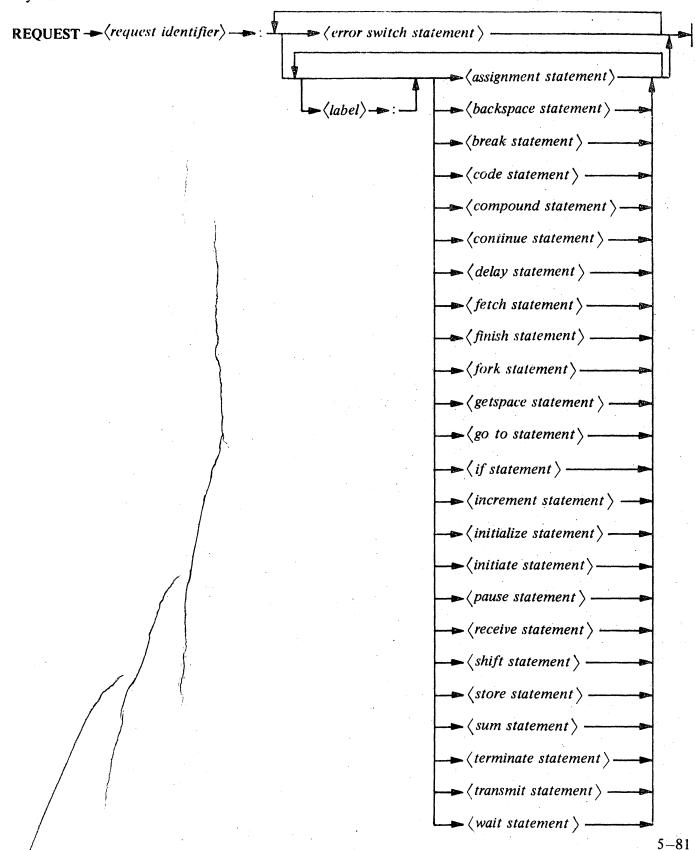
TRANSMITDELAY = 0. TRANSMITDELAY = 150 MICRO.

#### Semantics

The \( \langle modern transmitdelay statement \rangle \) defines the amount of time required for the modern to switch to a Clear to Send (CB) state after receiving a Request to Send (CA). \( \langle delay \) time \( \rangle \) must be expressed as \( \langle time \rangle \) and affects the amount of time delayed after an INITIATE RECEIVE or INITIATE TRANSMIT construct is executed and before the next statement is executed in a \( \langle control \) definition \( \rangle \) or \( \langle request \) definition \( \rangle \). The \( \langle \) defined in this statement is used in a compiler algorithm that calculates the delay. The compiler algorithm is discussed in the semantics of the \( \text{INITIATE RECEIVE} \) and \( \text{INITIATE} \) TRANSMIT constructs of the \( \langle \) initiate statement \( \rangle \). This statement must appear in each \( \langle modern \) definition \( \rangle \).

#### **REQUEST DEFINITION**

**Syntax** 



**REQUEST** 

Continued

#### Example

#### **REQUEST READTTY:**

INITIATE RECEIVE. RECEIVE TEXT [END]. TERMINATE NORMAL.

#### **Semantics**

\(\langle request \) definition\(\rangle s\), sometimes referred to as Requests, are coded line disciplines (protocols) that are used in communicating with the various terminal types in the data communications network. A \(\langle request\) definition\(\rangle\) must be coded for each capability of a terminal type; if it is possible for a terminal type to send input to the system and receive output from the system, then two \(\langle\) request definition\(\rangle\) s must be specified for that terminal type in its \(\langle\) terminal definition\(\rangle\). The input \(\langle\) request definition\(\rangle\) is generally referred to as the "Receive Request," and the output \(\langle\) request definition\(\rangle\) the "Transmit Request." (The specific \(\langle\) request definition\(\rangle\) to be used for each of these capabilities is specified by the \(\langle\) terminal request statement\(\rangle\).

When there is a message to be sent to a particular station on a line, the  $\langle control\ definition\rangle$  initiates the Transmit Request specified for the  $\langle terminal\ definition\rangle$  associated with the station. The Transmit Request procedure handles the transmission of the message. If the transmission of the message is successful, the Transmit Request is terminated, and a branch of control is made back to the  $\langle control\ definition\rangle$  for the initiation of the next Request.

If the terminal associated with a station is allowed to input data, the  $\langle control\ definition\rangle$  designated for that line normally initiates the Receive Request specified for the terminal type. If the terminal has information to transmit, the Receive Request procedure obtains a message space in which to store the received text, receives and stores the text, and then terminates in a manner that forwards the message to the MCS. If the terminal has nothing to transmit, the Receive Request procedure usually notes that there was no input, and terminates. In either case, upon termination, control returns to the  $\langle control\ definition\rangle$  for the initiation of the next Request.

 $\langle request \ identifier \rangle$  has the syntactic form of  $\langle identifier \rangle$ .

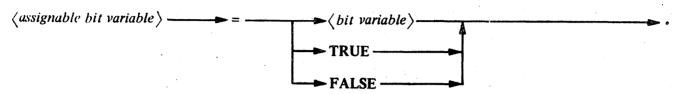
Statements in  $\langle request \ definition \rangle s$  are executed sequentially. In some cases, however, it is desirable to alter the order of execution of statements within the procedure. A  $\langle request \ statement \rangle$  preceded by a  $\langle label \rangle$  is one means of accomplishing this. The  $\langle go \ to \ statement \rangle$  is used to transfer control to a  $\langle label \rangle ed \langle request \ statement \rangle$ .

A  $\langle request\ statement \rangle$  must be appropriate for the type of Request in which it appears. That is, some  $\langle request\ statement \rangle s$  are allowed only in Receive Requests, some are allowed only in Transmit Requests, and some are allowed in either type. Subsequently, the semantics portion of each statement defines, among other things, in which type of  $\langle request\ definition \rangle$  the statement can appear.

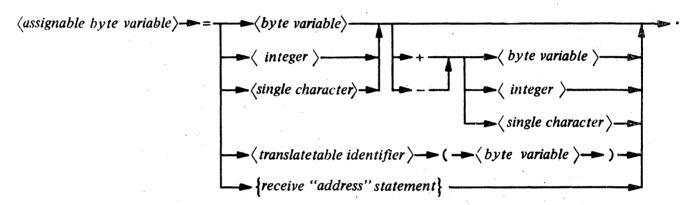
#### ASSIGNMENT STATEMENT

#### Syntax

#### FORM 1 LOGICAL ASSIGNMENT



#### FORM 2 - VALUE ASSIGNMENT



#### **Examples**

```
TOG [0] = TRUE.

TOG [1] = TOG [0].

LINE (BUSY) = FALSE.

RETRY = STATION (TALLY).

TALLY [0] = STATION (FREQUENCY) - TALLY [1].

CHARACTER = TRANSTABLEID (CHARACTER).

STATION (TALLY) = RECEIVE ADDRESS (TRANSMIT) [ADDERR:999].
```

#### **Semantics**

#### FORM 1

This form causes the value on the right side of the equal sign to replace the current value of  $\langle assignable bit variable \rangle$ .

#### FORM 2

Value assignment causes a calculated value on the right of the equal sign to be stored in the (assignable byte variable). Arithmetic calculations are done in modulo 255 arithmetic.

 $\langle assignable\ byte\ variable \rangle = \langle translatetable\ identifier \rangle\ (\langle byte\ variable \rangle).$ 

This construct is the means to invoke user-defined character translation. User-defined translation is effected by three areas of the NDL source program.

a. In a \(\langle translatetable definition \rangle\), the programmer must define the contents of a translation table and associate a \(\langle translatetable identifier \rangle\) with it.

#### **REQUEST**

#### Assignment Statement - Continued

b. In the \(\lambda terminal definition\rangle\) of a terminal type that requires special character translation, the programmer should suppress automatic character translation by using either of the following forms of the \(\lambda terminal code statement\rangle\):

CODE = BINARY.

or

#### CODE = EBCDIC.

c. In a  $\langle control\ definition \rangle$  or  $\langle request\ definition \rangle$ , the programmer invokes the translation by using this option of the value assignment. Any  $\langle byte\ variable \rangle$  can be designated as containing the character to be translated.

The \(\lambda\) translatetable identifier\(\rangle\) identifies the translation table to be used. An \(\lambda\) assignable by te variable\(\rangle\) is designated to the left of the equal sign, identifying where the resulting translated character is to be stored.

If N is the  $\langle source \ size \rangle$  (defined in the  $\langle translatetable \ definition \rangle$ ), then the N low-order bits of the  $\langle byte \ variable \rangle$  are used as an index into the translation table. The eight-bit character thus indexed is stored in the  $\langle assignable \ byte \ variable \rangle$ .

 $\langle assignable \ byte \ variable \rangle = \{receive "address" \ statement \}.$ 

This construct attempts to RECEIVE the address characters of a terminal, and store in \( \assignable \) byte \( \text{variable} \) the station index of a station whose address characters are equal to those received. The \( \text{receive} \) "address" statement \( \text{ is the same as described in the semantics of the RECEIVE ADDRESS construct of the \( \text{receive statement} \) The optional syntax in the \( \text{receive "address" statement} \) invokes the same actions as described in the \( \text{receive statement} \) semantics except for \( \text{ADDERR} \). Any action specified for \( \text{ADDERR} \) is taken if no valid station assigned to the line is found with address characters equal to those received.

Definitions
REQUEST
Backspace Statement

R	A	CK	SP	4	CE	ST	A٦	<b>FEN</b>	1	FN	ľ	Γ

**Syntax** 

BACKSPACE-

#### **Semantics**

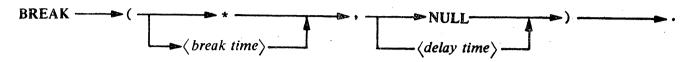
The \(\langle backspace statement \rangle \) causes the message text pointer to be moved backwards one character. This statement can only appear in a Receive Request. The \(\langle backspace statement \rangle \) may be executed repeatedly; however, the message text pointer will never be stepped back so far that it points into the message header.

#### **REQUEST**

**Break Statement** 

#### **BREAK STATEMENT**

**Syntax** 



#### Examples

BREAK (\*, NULL). BREAK (200 MILLI, 3 SEC). BREAK (\*, 3 SEC). BREAK (100 MILLI, NULL).

#### **Semantics**

The \(\langle break statement \rangle \) causes binary zeroes to be transmitted on the line, thus changing the state of the line to a "spacing" condition for a specified time.

The  $\langle break \ time \rangle$  specifies the  $\langle time \rangle$  to break. An asterisk indicates that a standard break of 2 character times should be used.

The  $\langle delay \ time \rangle$  specifies the  $\langle time \rangle$  to delay subsequent to the break and prior to when control continues.

#### **Semantics**

**CODE=ASCII** invokes the ASCII-to-EBCDIC translation for received data and the EBCDIC-to-ASCII translation for transmitted data.

**CODE=BINARY** inhibits any character translation on data transmitted or received.

#### **Pragmatics**

The \( \cdot \) code statement \( \rightarrow \) allows a programmer to either invoke or inhibit on a logical line the DCP ASCII-to-EBCDIC character code translation for input, and the EBCDIC-to-ASCII character code translation for output. Any \( \text{terminal definition} \right) \) that names, in its \( \text{terminal control statement} \right), a \( \cdot \) control \( \definition \right) \) that utilizes the \( \langle \cdot \) code statement \( \right), \text{ must define ASCII (BINARY)} \) as its character code in the \( \text{terminal code statement} \right). (Refer to the \( \langle \) terminal code statement \( \right) \) in this chapter.)

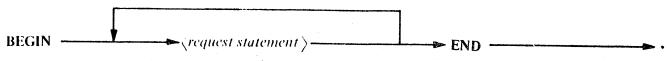
Once that translation has been invoked on a line, the translation continues until such time that it is inhibited. If translation is inhibited, translation will be inhibited on that line until invoked again by any of the following constructs: CODE=ASCII, TERMINATE NORMAL, TERMINATE LOGICALACK, TERMINATE LOGICALACK(RETURN), TERMINATE ERROR, TERMINATE ENABLEINPUT, or (while executing a Receive Request) TERMINATE NOINPUT.

#### **REQUEST**

**Compound Statement** 

#### **COMPOUND STATEMENT**

**Syntax** 



#### Example

BEGIN
INITIATE TRANSMIT.
TRANSMIT TEXT.
FINISH TRANSMIT.
END.

#### **Semantics**

The \(\langle compound statement \rangle \) groups several statements together to form a logical sequence. To execute more than one statement when the condition of an \(\langle if statement \rangle \) is satisfied, a \(\langle compound statement \rangle \) must be used.

# Definitions REQUEST Continue Statement

CONTINUE STATEMENT		•		
Syntax				
CONTINUE -		 	<del></del>	 _

#### **Semantics**

The \(\chicontinue statement\) can appear in only those \(\lambda request definition\rangle s\) and \(\lambda control definition\rangle s\) written to communicate with full duplex terminal types. This statement causes the co-line to resume processing, if, and only if, it had been suspended by a \(\lambda wait statement\rangle\), or a \(\lambda receive statement\rangle\) with a CONTINUE option specified. If the co-line had not been suspended, this statement acts as a no-op. The \(\lambda continue statement\rangle\) has no effect upon the line on which it was executed.

#### **Pragmatics**

Refer to the \( \frac{fork statement}{} \right) pragmatics.

#### **REQUEST**

**Delay Statement** 

#### **DELAY STATEMENT**

**Syntax** 

DELAY — ( delay time ) — .

#### Examples

DELAY (3 SEC). DELAY (0).

#### **Semantics**

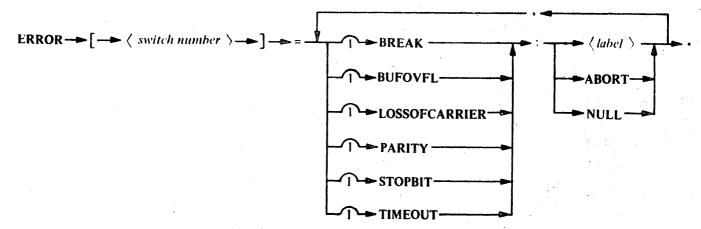
The \( \langle delay \) statement \( \rangle \) provides a means to delay a specified period of time before control proceeds to the next statement. The \( \langle \) request definition \( \rangle \) is suspended in a "sleep" state for the \( \langle \) delay time \( \rangle \) specified.

#### **Pragmatics**

The "sleep" state induced by the (delay statement) allows the DCP to service other logical lines.

#### **ERROR SWITCH STATEMENT**

#### Syntax



#### **Examples**

ERROR [0] = BREAK: 0, BUFOVFL: NULL, LOSSOFCARRIER: ABORT, PARITY: 999,

STOPBIT: 999. TIMEOUT: NULL.

ERROR [1] = BREAK: NULL.

ERROR [99] = BUFOVFL: NULL.

#### **Semantics**

The  $\langle error switch statement \rangle$  is a non-executable statement that allows the programmer to define a set of default actions that are to be taken in a  $\langle receive statement \rangle$  if the specified errors occur.  $\langle switch number \rangle$  has the syntactic form of  $\langle integer \rangle$ . The semantics of each option is described subsequently.

#### **BREAK**

The BREAK option variations cause the following actions if a break, that is, at least two character-times of a spacing line condition, is detected by the adapter cluster while receiving:

**BREAK: NULL** causes no action. Execution proceeds as if the break did not occur.

BREAK: (label) sets TRUE the (bit variable) BREAK [RECEIVE], and branches

control to  $\langle label \rangle$ .

BREAK: ABORT sets TRUE the (bit variable) BREAK [RECEIVE], and executes an

implicit TERMINATE ERROR.

#### **BUFOVFL**

The BUFOVFL option variations cause the following actions if the DCP is unable to service a Cluster Attention Needed (CAN) interrupt before the adapter cluster receives another character (thus destroying the previous character):

BUFOVFL: NULL causes no action. Execution proceeds as if the error conditions did not

occur.

BUFOVFL: \(\lambda \text{label}\rangle\) sets TRUE the \(\lambda \text{bit variable}\rangle\) BUFOVFL, and branches control to

(label).

#### REQUEST

Error Switch Statement - Continued

**BUFOVFL: ABORT** 

sets TRUE the \( \frac{bit variable}{} \) BUFOVFL, and executes an implicit

TERMINATE ERROR.

#### LOSSOFCARRIER

The LOSSOFCARRIER option variations cause the following actions if a loss of carrier is detected while receiving.

LOSSOFCARRIER: NULL causes no action. Execution proceeds as if the error did not occur.

LOSSOFCARRIER: \(\langle label \rangle \) sets TRUE the \(\langle bit variable \rangle \) LOSSOFCARRIER, and branches

control to  $\langle label \rangle$ .

LOSSOFCARRIER: ABORT sets TRUE the \( \begin{array}{l} bit variable \rangle \text{LOSSOFCARRIER}, and executes an \)

implicit TERMINATE ERROR.

There is one exception to the actions described in the above. If a loss of carrier is detected while receiving, and if the terminal is modem-connect, and if the terminal's (station definition) references a (modem definition) that contains the statement LOSSOFCARRIER=DISCONNECT, then an implicit disconnect is done, regardless of the action associated with LOSSOFCARRIER in the (error action statement).

#### **PARITY**

The PARITY option variations cause the following actions if a parity bit error is detected by the adapter cluster:

PARITY: NULL causes no action. Execution proceeds as if the error did not occur.

PARITY: (label) sets TRUE the (bit variable) PARITY, and branches control to

 $\langle label \rangle$ .

PARITY: ABORT sets TRUE the (bit variable) PARITY, and executes a TERMINATE

ERROR.

#### **STOPBIT**

The **STOPBIT** option variations cause the following actions if a stop bit error is detected by the adapter cluster:

STOPBIT: NULL causes no action. Execution proceeds as if the error did not occur.

STOPBIT: \(\lambda \lambda \text{label}\rangle \) sets TRUE the \(\lambda \text{bit variable}\rangle \) STOPBIT, and branches control to

 $\langle label \rangle$ .

STOPBIT: ABORT sets TRUE the \( \begin{aligned} \text{bit variable} \right) STOPBIT, and executes a TERMINATE

ERROR.

#### **TIMEOUT**

The TIMEOUT option variations of the TIMEOUT syntax shown below cause the actions described if the time required to receive a character exceeds the  $\langle timeout \ time \rangle$ . The  $\langle timeout \ time \rangle$  is defined in the  $\langle terminal \ timeout \ statement \rangle$ , but can be overridden by including the ( $\langle timeout \ time \rangle$ ) or (NULL) syntax options in the  $\langle receive \ statement \rangle$ .

**TIMEOUT: NULL** causes no action. Execution proceeds as if the error did not occur.

TIMEOUT: (label) sets TRUE the (bit variable) TIMEOUT, and branches control to

 $\langle label \rangle$ .

### Definitions **REQUEST**

Error Switch Statement - Continued

**TIMEOUT: ABORT** 

sets TRUE the (bit variable) TIMEOUT, and executes a TERMINATE ERROR.

#### **Pragmatics**

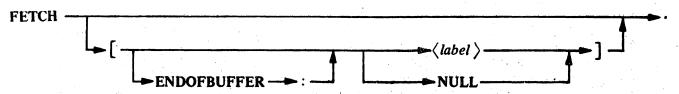
An \( \left(\text{error switch statement} \right)\) must be associated with a \( \left(\text{receive statement} \right)\) by means of a \( \left(\text{switch number} \right)\) reference before any of the default actions are invoked. The \( \left(\text{error switch statement} \right)\) can appear in a \( \left(\text{request definition} \right)\) as many times as the programmer deems convenient, providing the following restriction is adhered to: within a given \( \left(\text{request definition} \right)\), \( \left(\text{error switch statement} \right)\) s must have a unique \( \left(\text{switch number} \right)\), and all \( \left(\text{error switch statement} \right)\) s must precede all executable statements in the procedure.

REQUEST

Fetch Statement

#### **FETCH STATEMENT**

#### **Syntax**



#### Examples

FETCH.
FETCH [10].
FETCH [ENDOFBUFFER:NULL].

#### **Semantics**

The execution of the \( \) fetch statement \( \) loads into CHARACTER, the character pointed to by the message text pointer and updates the pointer to point forward one character position.

When using the  $\langle fetch \ statement \rangle$ , provision should be made for taking action if the end-of-the-text buffer is encountered. The programmer can specify this action by including the optional syntax shown in the syntax diagram.

NULL specifies that no action should be taken.

(label) specifies that control should branch to (label) if the end of buffer is encountered.

If the end of buffer is encountered and no action is specified, an implicit TERMINATE ERROR is executed.

For program documentation, the **ENDOFBUFFER** syntax can be added to the error action specification.

Supplementary Example

INITIATE TRANSMIT.

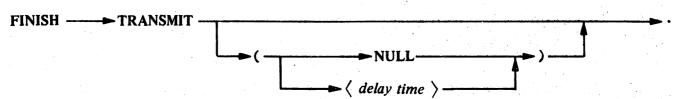
3: FETCH [ENDOFBUFFER:5].
TRANSMIT CHAR.

GO TO 3.

5: FINISH TRANSMIT.

#### **FINISH STATEMENT**

**Syntax** 



#### **Examples**

FINISH TRANSMIT.
FINISH TRANSMIT (NULL).
FINISH TRANSMIT (3 SEC).

#### **Semantics**

The purpose of the \( \frac{\( finish \) statement \\ \) is to take a line out of the transmit ready state and prepare the line to receive information. The adapter cluster delays a period of time long enough for the last character TRANSMITted to be transmitted, plus 2 milliseconds, before the line is put in a receive ready state. Although the \( \langle \( finish \) statement \( \rangle \) puts the line in a receive ready state, the cluster hardware invokes a delay that inhibits any data from being received for 25 milliseconds. An INITIATE RECEIVE construct should precede any subsequent \( \langle \) receive statement \( \rangle \) to override the 25-millisecond hardware delay.

The  $\langle delay \ time \rangle$  option allows the programmer to specify a software delay of  $\langle time \rangle$  before execution proceeds in the  $\langle control \ definition \rangle$ .

For example, the statement

FINISH TRANSMIT (3 SEC).

is equivalent to

FINISH TRANSMIT. DELAY (3 SEC).

The FINISH TRANSMIT (NULL) form is equivalent to FINISH TRANSMIT.

Definitions **REOUEST** 

Fork Statement

FORK STATEMENT

Syntax:

FORK ───── | label |

Example

**FORK 10.** 

#### **Semantics**

The \langle tork statement \rangle is allowed in only those \langle control definition \rangles and \langle request definition \rangles s that are written to communicate with full duplex terminal types. This statement can be executed in the \langle control definition \rangle or \langle request definition \rangle of the auxiliary line or the primary line. The execution of this statement causes the co-line control, if not busy, to branch to and begin executing code in the \langle request definition \rangle that executes the FORK at the \langle label \rangle specified, while control on the FORKing line executes an implicit PAUSE (i.e., a \langle pause statement \rangle) and continues executing in parallel. The co-line is determined busy or not busy by testing the BUSY bit (i.e., LINE(BUSY)) or AUX(LINE(BUSY)), whichever is appropriate). If the co-line is busy, the \langle fork statement \rangle acts as a no-op.

#### **Pragmatics**

Synchronization problems can occur between the primary and auxiliary lines as a result of the \( \fork \) statement\( \rightarrow \) executing the implicit PAUSE. The implicit PAUSE yields use of the DCP, to allow processing to to proceed on other lines. Thus, processing on the co-line is actually started before the FORKing line exits the \( \fork \) statement\( \rightarrow \). As a result, the programmer must, by some means (e.g., by setting and testing line TOGs), effect the synchronization of the lines. This is especially critical if the code contains \( \lambda \) wait statement\( \rightarrow \) s and \( \lambda \) continue statement\( \rightarrow \) s. The following example illustrates how full duplex lines could "hang" as a result of poor synchronization.

FORK 10. WAIT.

10: CONTINUE. WAIT.

Assume that the primary line executes the FORK 10. At that point, the primary line temporarily yields use of the DCP to other lines. The auxiliary line starts up and executes the CONTINUE. Since primary control is still at the (fork statement) and is not in a (wait statement), the auxiliary line CONTINUE acts as a no-op. Next. the auxiliary line executes the WAIT. When the primary line gets use of the processor again, it executes its WAIT. At this point, the primary and auxiliary lines are "hung", each WAITing for a CONTINUE from its co-line.

<b>GETSPACE STATEME</b>	NT		· · · · · · · · · · · · · · · · · · ·		
Syntax					
GETSPACE ———	<b></b> [	→ ⟨ label ⟩	-	<b>&gt;</b> ]	

#### Example

**GETSPACE** [10]

#### **Semantics**

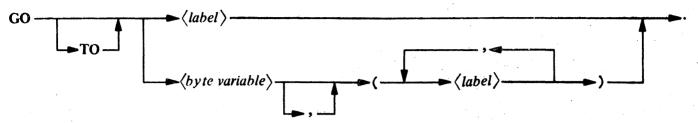
The  $\langle getspace\ statement \rangle$  provides the means for a Receive Request to explicitly acquire a message space for input. The message space (if obtained) is linked into the head of the Station Queue, thereby setting STATION (QUEUED) to TRUE. If there is no message space available at the time the  $\langle getspace\ statement \rangle$  is executed, control branches to the  $\langle label \rangle$ . If a message space has already been acquired, this instruction acts as a no-op. This statement is also treated as a no-op if it appears in a Transmit Request.

#### REQUEST

Go To Statement

#### **GO TO STATEMENT**

#### **Syntax**



#### Examples

GO 10.

GO TO 10.

GO TO TOGS, (0,1,2,3).

GO TO STATION (5,9,12).

#### **Semantics**

The  $\langle go \ to \ statement \rangle$  alters the path of control, that is, the sequential flow of statement execution, within a  $\langle request \ definition \rangle$ .

GO TO \(\langle \langle abel \rangle \)

This form of the \( \go \) to statement \( \) unconditionally transfers control to the \( \lambda \) specified.

GO TO \(\langle byte variable \rangle \). . .

This form of the  $\langle go \ to \ statement \rangle$  provides a convenient means of dynamically selecting one or more  $\langle label \rangle s$  to which control could branch. The  $\langle label \rangle$  to branch to is selected by using the  $\langle byte \ variable \rangle$  as an index value. If N represents the number of  $\langle label \rangle s$  in the  $\langle go \ to \ statement \rangle$ , then the  $\langle label \rangle s$  are numbered 0 to N-1. The  $\langle label \rangle$  corresponding to the index value is the  $\langle label \rangle$  to which control branches. If the index value is greater than N-1, then control continues at the statement following the  $\langle go \ to \ statement \rangle$ .

#### **Supplementary Example**

**GO TO STATION (5,9,12).** 

% EXECUTION CONTINUES HERE IF STATION > 2.

5: TOG [0] = TRUE.

9: TOG[1] = TRUE.

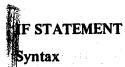
12: TOG[2] = TRUE.

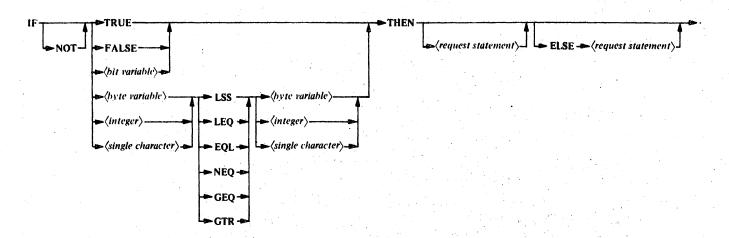
## Definitions REQUEST

Go To Statement - Continued

This example illustrates the GO TO  $\langle byte\ variable \rangle$  construct of the  $\langle go\ to\ statement \rangle$ . The value of STATION determines the next statement to be executed. If the value of STATION is 0, control branches to the  $\langle label \rangle$  5; if the value of STATION is 1, control branches to  $\langle label \rangle$  9; and if the value of STATION is 2, control branches to  $\langle label \rangle$  12. If the value of STATION is greater than 2, control continues at the next sequential statement.

Definitions
REQUEST
If Statement





#### **Examples**

```
IF TRUE THEN.

IF TOG [0] THEN TOG [0] = FALSE.

IF TALLY [0] LSS TALLY [1] THEN TALLY [0] = TALLY [1].

IF CHARACTER = 4"FF" THEN

INITIATE BREAK.

ELSE

BEGIN

CHAR = 4"00".

GO TO 0.

END.
```

#### **Semantics**

The  $\langle if statement \rangle$  causes a condition (i.e., a Boolean expression) to be evaluated. The subsequent path of program control depends on whether the condition is evaluated as TRUE or FALSE.

If the condition is TRUE, the  $\langle request \ statement \rangle$  following the THEN, if present, is executed. Program control then resumes at the statement that follows the  $\langle if \ statement \rangle$ .

If the condition is FALSE, the (request statement) following the ELSE is executed or, if the ELSE (request statement) is omitted, program control resumes at the (request statement) following the (if statement).

The  $\langle request \ statement \rangle$  can be any legal  $\langle request \ statement \rangle$ , including the  $\langle if \ statement \rangle$  and  $\langle compound \ statement \rangle$ .

The meanings of the relational operators are contained in table 5-5.

The following diagrams illustrate the above semantics.

The following diagrams illustrate the above semantics.

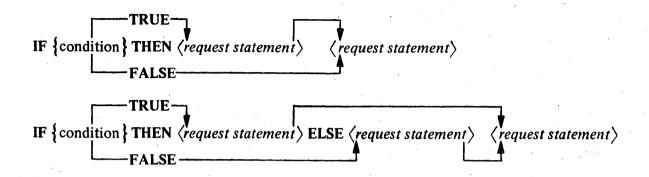


Table 5-5. Relational Operators

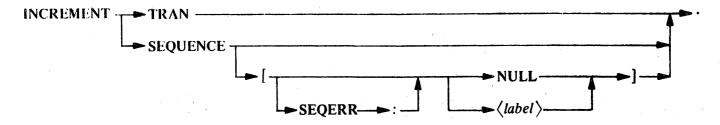
RELATIONAL OPERATOR	MEANING	SYNONYMS
LSS	Less than	< and LS
LEQ	Less than or equal to	LE
EQL	Equal to	= and EQ
NEQ	Not equal to	NE
GEQ	Greater than or equal to	GE
GTR	Greater than	> and GT

**REQUEST** 

Increment Statement

#### **INCREMENT STATEMENT**

**Syntax** 



#### **Examples**

INCREMENT TRAN.
INCREMENT SEQUENCE [SEQERR:10].
INCREMENT SEQUENCE [NULL].

#### **Semantics**

#### **INCREMENT TRAN**

This construct of the  $\langle increment \ statement \rangle$  is only allowed in those  $\langle request \ definition \rangle s$  in the  $\langle terminal \ request \ statement \rangle s$  of  $\langle terminal \ definition \rangle s$  that contain a  $\langle terminal \ transmission \ number \ length \ statement \rangle$  defining the transmission number length as nonzero and non-NULL.

INCREMENT TRAN causes 1 to be added to the receive transmission number stored in the Station Table when it is executed in a Receive Request, and causes 1 to be added to the transmit transmission number stored in the Station Table when it is executed in a Transmit Request.

The transmission numbers are stored and incremented in EBCDIC.

If INCREMENT TRAN causes the transmission number to exceed (overflow) the size of the transmission number field, the carry is truncated and the result will be zeros (i.e., EBCDIC zeros) in that field.

#### **INCREMENT SEQUENCE**

This construct causes the sequence number stored in the DCP Station Table to be increased by the value of the increment (also stored in the DCP Station Table), providing that the station is in "sequence mode"; otherwise, this statement is a no-op.

When using the INCREMENT SEQUENCE construct, provision should be made for taking action if the increment caused the sequence number to exceed (overflow) the size of the sequence number field. The programmer can take such action by including the optional syntax. Failure to include overflow action results in an implicit TERMINATE ERROR if an overflow occurs.

SEQERR: NULL and NULL are semantically equivalent. These options set the SEQERR \( \) bit variable \( \) TRUE, and control continues at the next sequential instruction.

SEQERR:  $\langle label \rangle$  and  $\langle label \rangle$  are semantically equivalent. They cause the SEQERR  $\langle bit \ variable \rangle$  to be set TRUE, and control to branch to  $\langle label \rangle$ .

Regardless of whether error action is specified or not, an overflow of the sequence number field destroys the contents of that field.

### Definitions REQUEST

Increment Statement - Continued

#### **Pragmatics**

A station is considered to be in sequence mode whenever its **SEQUENCE** (bit variable) is **TRUE**. **SEQUENCE** can be set **TRUE** only as a result of the Message Control System (MCS) executing the SET/RESET SEQUENCE MODE (TYPE = 49) DCWRITE. In addition, the TYPE 49 DCWRITE also stores the starting sequence number and increment in the appropriate fields of the DCP Station Table.

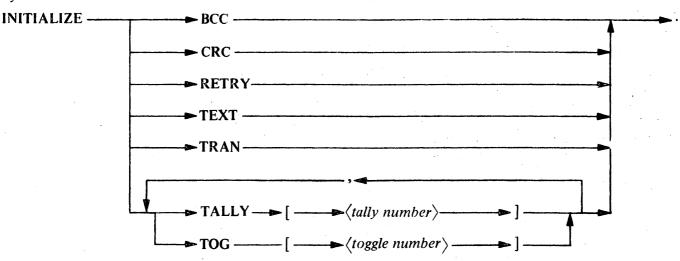
Sequence mode can be used for any application that the NDL programmer may see fit. Its use, however, requires common conventions between the NDL programmer and the MCS programmer. Burroughs has utilized sequence mode constructs in two \( \frac{request definition}{s} \) of SYMBOL/SOURCENDL: READTELETYPE and WRITETELETYPE. Both require the cooperation of SYSTEM/CANDE to effect the execution of those statements. The reader is referred to those \( \frac{request definition}{s} \) as an example of a particular application that Burroughs has implemented.

#### **REQUEST**

Initialize Statement

#### INITIALIZE STATEMENT

**Syntax** 



#### Examples

INITIALIZE BCC.
INITIALIZE CRC.
INITIALIZE RETRY.

#### Semantics

#### **INITIALIZE BCC**

This construct causes the \( \begin{align\*} \lambda \text{trevariable} \rangle \text{ BCC} \) to be initialized for purposes of accumulating a Block Check Character. The value to which BCC is initialized is dependent upon the horizontal parity defined for the station's associated \( \lambda \text{terminal definition} \rangle \) (in the \( \lambda \text{terminal definition parity statement} \rangle \)). If horizontal parity is defined as HORIZONTAL:ODD, then BCC is initialized to all ones (i.e., 4"FF"). If defined as HORIZONTAL:EVEN, then INITIALIZE BCC initializes BCC to all zeros (i.e., 4"00").

#### **INITIALIZE CRC**

This instruction initializes CRC to the initial value required for calculating the Cyclic Redundancy Check. Any \(\lambda terminal definition\rangle\) referencing a \(\lambda request definition\rangle\) (in the \(\lambda terminal request statement\rangle\)) as that contains this instruction must define the horizontal parity (in the \(\lambda terminal parity statement\rangle\)) as HORIZONTAL:CRC(16); otherwise a syntax error is generated.

#### INITIALIZE RETRY

This instruction causes the value stored in DCP INITIALRETRY to be stored DCP RETRY.

#### **INITIALIZE TEXT**

The function of this form is to initialize the message text pointer to zero. When initialized to zero, the message text pointer points to the first text character of the message.

# Definitions **REQUEST**

Initialize Statement - Continued

# **INITIALIZE TRAN**

This form causes zeroes (i.e., EBCDIC zeroes, 4"F0F0F0") to be stored in the appropriate Transmission Number fields of the Station Table. In a Receive Request, zeroes are stored in the Receive Transmission Number field; in a Transmit Request, zeroes are stored in the Transmit Transmission Number field.

# **INITIALIZE TALLY** [\(\langle tally number \rangle \)]

This form causes the specified station TALLY to be initialized from the appropriate message header field if a message is present; otherwise the specified TALLY is initialized to zero.

# INITIALIZE TOG [ \( \text{toggle number} \) ]

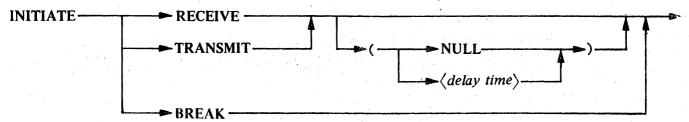
This form causes the specified station TOGGLE to be initialized from the appropriate message field if a message is present; otherwise the specified TOGGLE is initialized FALSE.

# REQUEST

Initiate Statement

# INITIATE STATEMENT

# Syntax



# Examples

INITIATE RECEIVE.
INITIATE TRANSMIT (3 SEC).
INITIATE BREAK.

#### **Semantics**

#### **INITIATE RECEIVE**

The INITIATE RECEIVE construct causes the adapter cluster to initiate a receive delay calculated for the station. After the delay, the hardware is ready to receive information.

The amount of time delayed, referred to as the Initiate Receive delay, is unique to each station and is calculated at compile-time for each station. The algorithm that the compiler uses to calculate the Initiate Receive delay is described in the following three paragraphs.

- a. If the \(\langle modem definition \rangle\) referenced in the \(\langle station definition \rangle\) (in the \(\langle station modem statement \rangle\) defines the modem NOISEDELAY as being greater than zero, then the Initiate Receive delay is 2 milliseconds less than the combined \(\langle time \rangle s\) defined in the \(\langle modem noisedelay statement \rangle\) and the \(\langle modem transmitdelay statement \rangle\).
- b. If the modem NOISEDELAY is defined as zero and the modem TRANSMITDELAY is defined as being less than 7 milliseconds, then the Initiate Receive delay is zero.
- c. If the modem NOISEDELAY is defined as zero and the modem TRANSMITDELAY is defined as being equal to or greater than 7 milliseconds, then the Initiate Receive delay is the lesser of 15 milliseconds or
  - $(1.5 \text{ milliseconds} + \frac{\text{modem TRANSMITDELAY}}{2}).$

The NULL option or the \( \langle delay \) time \( \rangle \) option can be used to override the calculated Initiate Receive delay.

NULL immediately readies the hardware so that it can receive information. \( \langle delay \) time \( \rangle \) specifies a \( \langle \) time \( \rangle \) to be used in place of the Initiate Receive delay.

# **Pragmatics**

An INITIATE RECEIVE instruction should precede the first \( \text{receive statement} \) following a transmission. If it does not, there is a possibility that execution of the \( \text{receive statement} \) will be delayed for a period of time of up to 25 milliseconds. The cause of the 25-millisecond delay is described under the semantics of the \( \text{finish statement} \).

# **INITIATE TRANSMIT**

The INITIATE TRANSMITconstruct causes the adapter cluster to be put in a transmit state after a calculated delay. The amount of time delayed is referred to as the Initiate Transmit Delay, and is unique to each station. It is derived by taking the greater of the NOISEDELAY \(\lambda time \rangle\) specified for the modem configured at the system end, or the TURNAROUND \(\lambda time \rangle\) specified by the station's \(\lambda terminal \) definition\(\rangle\).

This construct must be executed prior to any attempt to TRANSMIT.

The NULL option or the  $\langle delay\ time \rangle$  option can be used to override the calculated Initiate Transmit delay. NULL causes the adapter cluster to be put in a transmit state immediately.  $\langle delay\ time \rangle$  specifies a  $\langle time \rangle$  to be used in place of the Initiate Transmit delay.

# INITIATE BREAK

The INITIATE BREAK construct causes binary zeroes to be transmitted on the line, thus changing the state of the line to a "spacing" condition. The line remains in the spacing condition until some subsequent instruction causes the adapter cluster to change the state of the line. Constructs that would change the line's state are INITIATE TRANSMIT, INITIATE RECEIVE, FINISH TRANSMIT, BREAK and IDLE.

Definitions **REQUEST** 

Pause Statement

PAUSE STATEMENT.

**Syntax** 

PAUSE-

# **Semantics**

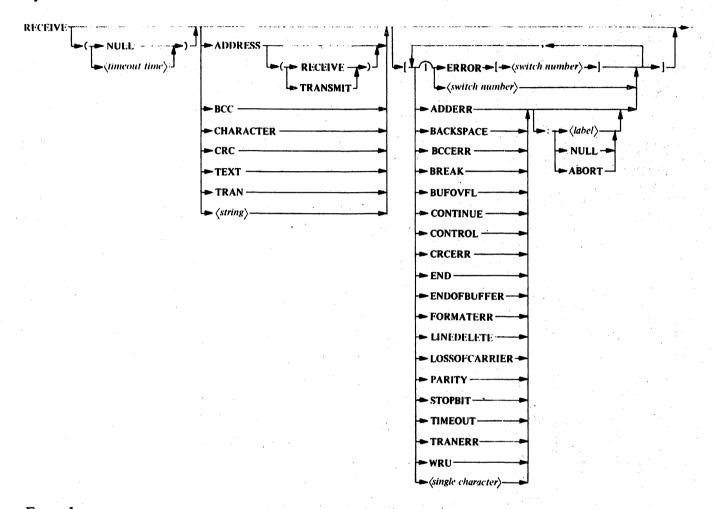
The \( \text{pause statement} \) suspends the \( \text{request definition} \) in a "sleep" state for a minimum period of time (200 microseconds for the B 6358 Model II DCP, and 6 microseconds for the B 6350 Model I DCP) to allow the DCP to service other lines. It is recommended that a \( \text{pause statement} \) be used in any kind of loop that would tie up processor time and thereby prevent the servicing of other lines. The failure to do so results in a high number of timeout faults.

# **Pragmatics**

Instances may occur in which the DCP requires an even greater period of "sleep" to service other lines. Repeated timeout faults, despite utilization of the  $\langle pause\ statement \rangle$ , are indications of such conditions. A greater period of "sleep" time can be effected by means of a  $\langle delay\ statement \rangle$ , with the  $\langle delay\ time \rangle$  specified greater than "sleep" time effected by the  $\langle pause\ statement \rangle$ .

# RECEIVE STATEMENT

# **Syntax**



# **Examples**

```
RECEIVE CHARACTER.
RECEIVE (3 SEC) ADDRESS (RECEIVE) [0, ADDERR:10].
RECEIVE (NULL) [
PARITY:999,
LOSSOFCARRIER:999,
BACKSPACE:NULL,
END,
WRU:NULL
].
RECEIVE CRC [ERROR [1], CRCERR:10].
RECEIVE "LITERAL STRING" [FORMATERR:NULL].
RECEIVE EOT SOH.
RECEIVE TEXT [END:10].
```

# **REQUEST**

Receive Statement - Continued

# **Semantics**

The  $\langle receive\ statement \rangle$  causes the adapter cluster to attempt to receive information from the appropriate logical line.

The following two syntax items define a maximum amount of time that the adapter cluster should wait for receipt of the first character, and then each subsequent character, if applicable, before assuming that the terminal has "timed out." If neither of these options is included, the \(\lambda timeout \time\rangle\) defined (in the \(\lambda terminal \timeout \time\rangle\)) for the station's associated terminal type is implicitly used as the \(\lambda timeout \time\rangle\) in this statement.

# (NULL)

This option specifies that the adapter cluster should wait an infinite amount of time.

(\langle time \rangle)

The \(\langle time \rangle \) defines a \(\langle time \rangle \) that the adapter cluster should wait for a character. If this \(\langle time \rangle \) is exceeded before receipt of a character, and the TIMEOUT syntax appears, then the action specified for TIMEOUT is taken (refer to TIMEOUT). If the \(\langle time \rangle \) is exceeded and TIMEOUT syntax does not appear, an implicit TERMINATE ERROR is executed.

The following syntax options define the nature of the information to be received, the amount of information to be received, and how the information is to be handled. If none of the options are used, it is semantically equivalent to specifying CHARACTER (e.g., "RECEIVE." is semantically equivalent to "RECEIVE CHARACTER.").

# **ADDRESS**

The proper number of address characters (as defined by the station's associated \(\lambda terminal definition\rangle\) in the \(\lambda terminal address size statement \rangle\)) are received and checked for agreement against the actual address characters defined in the \(\lambda station address statement \rangle\). If the address characters do not correspond, an address error condition results; if the ADDERR syntax appears, then the specified action is taken. Otherwise an implicit TERMINATE ERROR is executed. (Refer to the ADDERR semantics.)

# ADDRESS (RECEIVE)

This option is equivalent to ADDRESS, except that ADDRESS (RECEIVE) must be used when an address pair is defined in the *(station address statement)* and the programmer needs to check for the proper receive address.

# ADDRESS (TRANSMIT)

This option is equivalent to ADDRESS, except that ADDRESS (TRANSMIT) must be used when an address pair is defined in the \( \station address statement \) and the programmer needs to check for the proper transmit address.

# **BCC**

One character is received and checked against the  $\langle byte\ variable \rangle$  BCC. If the character received and BCC are not equal, a Block Check Character error condition results; if the BCCERR syntax appears, then the specified action is taken. Otherwise an implicit TERMINATE ERROR is executed.

Presumably, if the **RECEIVE BCC** construct appears, the programmer has defined horizontal parity in the \(\lambda terminal parity statement\rangle\), and the accumulated Block Check Character is contained in **BCC**.

Receive Statement - Continued

#### **CHARACTER**

One character is received and stored in CHARACTER.

# **CRC**

Two characters are received. The first character is checked against CRC [0], and the second compared against CRC [1]. If the characters received and CRC are not equal, a Cyclic Redundancy Check error condition results; if the CRCERR syntax appears, then specified action is taken. Otherwise an implicit TERMINATE ERROR is executed.

Presumably, if the RECEIVE CRC instruction appears, the programmer has defined horizontal parity as HORIZONTAL:CRC(16) in the \(\lambda\text{terminal parity statement}\rangle\), and the Cyclic Redundancy Check is contained in CRC [0] and CRC [1].

# **TEXT**

Characters are received into **CHARACTER** and stored in the text portion of the message space obtained until either a syntax option results in a branch from the \( \textit{receive statement} \), or a non-recoverable error, such as a disconnect, occurs. If the occurrence of a particular character results in a branch outside of the \( \text{receive statement} \) (as specified by a syntax option), then that character is not stored but remains in **CHARACTER**.

The RECEIVE TEXT construct is, in effect, the same as the following loop:

1: RECEIVE CHARACTER. STORE CHARACTER. GO TO 1.

In nearly all cases, the *(receive statement)* should contain optional syntax to avoid the "endless" loop and an eventual implicit **TERMINATE ERROR** as a result of a timeout, end-of-buffer condition, etc.

# **TRAN**

The proper number of transmission number characters (as defined by the station's associated \(\text{terminal definition}\) in the \(\text{terminal transmission number length statement}\)) are received and checked for agreement with the Receive Transmission Number maintained in the DCP Station Table. If the characters received and the Receive Transmission Number are not equal, a transmission number error results. If the TRANERR syntax appears, then specified action is taken; otherwise an implicit TERMINATE ERROR is executed.

 $\langle$  string $\rangle$ 

The number of characters as indicated by the length of the \(\string\)\) are received and checked against those characters in the \(\string\)\. If the \(\string\)\ and the characters received are not equal, then a format error condition results. If the FORMATERR syntax option appears, then that action is taken; otherwise an implicit TERMINATE ERROR is executed.

The following syntax options specify actions to be taken upon either the receipt of defined characters or occurrences of specific error conditions.

# REQUEST

Receive Statement - Continued

# **ERROR** [ \( \switch number \) ]

This syntax option associates a previously defined Error Switch with the \( \text{receive statement} \). This allows the programmer to associate a set of previously defined error actions with the \( \text{receive statement} \), thus reducing the amount of coding required for each \( \text{receive statement} \). BREAK, BUFOVFL, LOSSOFCARRIER, PARITY, STOPBIT, and TIMEOUT syntax options are not allowed if the ERROR \( \left\) switch number \( \right\) syntax appears in the \( \text{receive statement} \right\). Refer to the \( \text{error switch statement} \right\) for more information.

(switch number)

Semantically equivalent to ERROR [\( \switch number \)].

# **ADDERR**

The ADDERR option variations cause the following actions if an address error is detected when an attempt is made to receive a terminal's address characters:

**ADDERR** 

sets TRUE the ADDERR (bit variable) and branches control to the next sequential statement.

ADDERR: NULL

causes no action. Execution proceeds as if the error condition did not occur.

ADDERR: (lahel)

sets TRUE the ADDERR (bit variable) and branches control to (label).

ADDERR: ABORT

not allowed.

# BACKSPACE

The following BACKSPACE option variations cause the following actions if the terminal's backspace character (as defined by the \(\lambda terminal\) backspace character statement\(\rangle\)) is received:

**BACKSPACE** 

moves the message text pointer backwards one character position, and branches control to the next sequential statement.

**BACKSPACE: NULL** 

moves the message text pointer backwards one character. Control remains within the \( \chi receive \) statement \( \chi \) if of the form RECEIVE TEXT.

BACKSPACE: \(\langle \langle abel \rangle \)

moves the message text pointer backwards one character, and branches control to  $\langle label \rangle$ .

**BACKSPACE: ABORT** 

not allowed.

# **BCCERR**

The following BCCERR option variations cause the following actions if the character received is not equal to the data stored in BCC.

**BCCERR** 

sets TRUE the *(bit variable)* BCCERR, and branches control to the next sequential statement.

**BCCERR:NULL** 

causes no action. Execution proceeds as if the error condition did not occur.

# Definitions **REQUEST**

# Receive Statement - Continued

BCCERR: \(\langle \langle abel \rangle \)

sets TRUE the \( \frac{bit variable}{} \rightarrow BCCERR \) and

branches control to (label).

**BCCERR:ABORT** 

not allowed.

# BREAK

The BREAK option variations cause the following actions as if a break, that is, at least two character-times of a spacing line condition, is detected by the adapter cluster while receiving:

**BREAK** 

sets TRUE the \( \begin{aligned} bit variable \rangle \text{BREAK} \end{aligned} \)

[RECEIVE], and branches control to the next

sequential statement.

**BREAK:NULL** 

causes no action. Execution proceeds as if the

break did not occur.

BREAK: \(\label\)

sets TRUE the \( \begin{aligned} bit variable \rightarrow BREAK \end{aligned} \)

[RECEIVE], and branches control to \( \lambda \lambda \text{label} \).

**BREAK: ABORT** 

sets TRUE the *(bit variable)* BREAK [RECEIVE], and executes an implicit

TERMINATE ERROR.

#### BUFOVFL

The following variations of the BUFOVFL option cause the following actions if the DCP is unable to service a Cluster Attention Needed (CAN) interrupt before the Adapter Cluster receives another character (thus destroying the previous character):

**BUFOVFL** 

sets TRUE the \(\langle bit variable \rangle \) BUFOVFL, and

branches control to the next sequential

statement.

**BUFOVFL:NULL** 

causes no action. Execution proceeds as if the

error condition did not occur.

**BUFOVFL**: \(\langle label \rangle \)

sets TRUE the \(\langle bit variable \rangle \) BUFOVFL, and

branches control to \(\lambda \lambda abel \rangle.

**BUFOVFL: ABORT** 

sets TRUE the *(bit variable)* BUFOVFL, and executes an implicit TERMINATE ERROR.

# **CONTINUE**

This item is allowed only in \( \text{receive statement} \) s of \( \text{control definition} \) s and \( \text{request definition} \) s that are written to communicate with full duplex terminal types. CONTINUE syntax causes action as described below if the co-line executes a \( \text{continue statement} \) before all information specified by the \( \text{receive statement} \) is received.

CONTINUE

branches control to the next sequential

statement.

CONTINUE: NULL

causes no action. Execution proceeds as if the *(continue statement)* had not been executed.

CONTINUE: (label)

branches control to  $\langle label \rangle$ .

**CONTINUE: ABORT** 

not allowed.

# REQUEST

Receive Statement - Continued

# CONTROL

The following variations of the **CONTROL** option cause the following actions if the control character of the station (as defined in the *(station control character statement)*) is received:

CONTROL

sets TRUE the *(bit variable)* CONTROLFLAG, and branches control to the next sequential statement.

**CONTROL:NULL** 

sets TRUE the \( \begin{array}{l} bit variable \rightarrow CONTROL FLAG, \) and execution continues if the character was not the station's control character.

CONTROL: \label

sets TRUE the \( \lambda \) it variable \( \cong \) CONTROLFLAG, and branches control to \( \lambda \) label \( \lambda \).

not allowed.

CONTROL: ABORT

# **CRCERR**

The following variations of the CRCERR option cause the following actions if the first character received does not equal CRC [0], or the second character received does not equal CRC [1]. (This item is appropriate only for the RECEIVE CRC construct of the \( \text{receive statement} \); refer to the CRC option.)

**CRCERR** 

sets TRUE the \( \begin{array}{l} \begin

CRCERR:NULL

cause no action. Execution proceeds as if the

error did not occurr.

CRCERR: \(\label\)

sets TRUE the \(\langle bit variable \rangle CRCERR\), and

branches control to  $\langle label \rangle$ .

**CRCERR: ABORT** 

not allowed.

#### **END**

The following variations of the **END** option cause the following actions if the "end" character of the station (as defined by the \(\lambda terminal end character statement\) in the \(\lambda terminal definition\) associated with the station) is received:

**END** 

causes control to branch to the next sequential statement.

**END:NULL** 

causes no action. Execution proceeds as if the character was not the "end" character.

END: \(\label\)

branches control to  $\langle label \rangle$ .

**END:ABORT** 

not allowed.

# REQUEST

Receive Statement - Continued

# **ENDOFBUFFER**

This syntax item is allowed in the **RECEIVE TEXT** contruct of the *(receive statement)*. The variations of the **ENDOFBUFFER** option shown below cause the following actions if either of the following conditions arises:

- a. There is no message space and an attempt is made to store information into a message space (the store function is an implicit action of the **RECEIVE TEXT** construct), or
- b. The number of characters stored in the message exceeds the maximum allowed (the maximum is defined by either the \(\lambda terminal \) maximum statement \(\rangle\) or the \(\lambda terminal \) buffer size statement \(\rangle\).

**ENDOFBUFFER** 

sets TRUE the \(\langle bit variable \rangle ENDOFBUFFER\), and branches control to the next sequential statement.

ENDOFBUFFER:NULL

causes no action. Execution proceeds as if the error did not occur.

**ENDOFBUFFER**: \( \langle \langle abel \)

sets TRUE the (bit variable) ENDOFBUFFER,

and branches control to \( \lambda \lambda bel \rangle \).

**ENDOFBUFFER: ABORT** 

not allowed.

#### **FORMATERR**

The following variations of the **FORMATERR** option cause the following actions if the characters received are not equal to those in the  $\langle string \rangle$  (this item is appropriate only for the **RECEIVE**  $\langle string \rangle$  construct of the  $\langle receive statement \rangle$ ):

**FORMATERR** 

sets **TRUE** the *\langle bit variable \rangle* **FORMATERR**, and branches control to the next sequential statement.

FORMATERR:NULL

causes no action. Execution proceeds as if the error did not occur.

**FORMATERR**: *(label)* 

sets **TRUE** the *(bit variable)* **FORMATERR**, and branches control to *(label)*.

FORMATERR: ABORT

not allowed.

#### LINEDELETE

The following variations of the LINE DELETE option cause the following actions if the station's linedelete character is received (the LINEDELETE character is defined by the \(\lambda terminal linedelete \) character statement\(\rangle\):

LINEDELETE

alters the value of the message text pointer to point to the first character position in the message text, and branches control to the next sequential statement.

LINEDELETE:NULL

alters the value of the message text pointer to point to the first character position in the message text, and execution proceeds as if the character was not the linedelete character.

# REQUEST

Receive Statement - Continued

LINEDELETE: \(\langle \langle \langle

alters the value of the message text pointer to point to the first character position in the message

text, and branches control to (label).

LINEDELETE: ABORT

not allowed.

# LOSSOFCARRIER

The following variations of the LOSSOFCARRIER syntax cause the following actions if a loss of carrier is detected while receiving.

LOSSOFCARRIER

sets TRUE the \(\langle bit variable \rangle LOSSOFCARRIER\),

and branches control to the next sequential

statement.

LOSSOFCARRIER:NULL

causes no action. Execution proceeds as if the

error did not occur.

LOSSOFCARRIER: \(\langle \langle abel \rangle \)

sets TRUE the \(\langle bit variable \rangle LOSSOFCARRIER\),

and branches control to \( \langle label \rangle \).

LOSSOFCARRIER: ABORT

sets TRUE the \( \begin{array}{l} bit variable \rangle LOSSOFCARRIER, \) and executes an implicit TERMINATE ERROR.

There is one exception to the actions described above. If a loss of carrier is detected while receiving, and if the terminal is modem-connect, and if the terminal's \( \station \) definition \( \rightarrow \) references a \( \lambda \) modem definition \( \rightarrow \) that contains the construct LOSSOFCARRIER=DISCONNECT, then an implicit disconnect is done, regardless of the action specified.

# **PARITY**

The following variations of the **PARITY** option cause the following actions if a parity bit error is detected by the adapter cluster:

**PARITY** 

sets TRUE the *(bit variable)* PARITY, and branches control to the next sequential state-

ment.

PARITY:NULL

causes no action. Execution proceeds as if the

error did not occur.

PARITY: \(\label\rangle\)

sets TRUE the \( \frac{bit variable}{} \) PARITY, and

branches control to \(\lambda \lambda \text{label}\rangle.

**PARITY: ABORT** 

sets TRUE the (bit variable) PARITY, and

executes a TERMINATE ERROR.

# **STOPBIT**

The following variations of the **PARITY** option cause the described actions if a stop bit error is detected by the adapter cluster:

**STOPBIT** 

sets TRUE the (bit variable), and branches control

to the next sequential statement.

STOPBIT:NULL

causes no action. Execution proceeds as if the

error did not occur.

# Definitions REQUEST

# Receive Statement - Continued

STOPBIT: \(\label\)

sets TRUE the \(\langle bit variable \rangle STOPBIT\), and

branches control to \(\lambda \lambda \text{label}\rangle.

STOPBIT: ABORT

sets TRUE the \(\langle \text{iit variable} \rangle STOPBIT\), and

executes a TERMINATE ERROR.

#### TIMEOUT

The variations of the **TIMEOUT** syntax shown below cause the actions described if the time required to receive a character exceeds the \(\lambda timeout \time \rangle\). The \(\lambda timeout \time \rangle\) is defined in the \(\lambda terminal \timeout \time \rangle\) or (NULL) syntax options in the \(\lambda receive \text{statement} \rangle\).

**TIMEOUT** 

sets the \( \frac{bit variable}{\) TIMEOUT, and branches control to the next sequential statement.

TIMEOUT:NULL

causes no action. Execution proceeds as if the

error did not occur.

TIMEOUT: \(\label\rangle\)

sets TRUE the (bit variable) TIMEOUT, and

branches control to (label).

TIMEOUT: ABORT

sets TRUE the (bit variable) TIMEOUT, and

executes a TERMINATE ERROR.

# TRANERR

The following variations of the **TRANERR** option cause the described actions if the characters received and the Receive Transmission Number stored in the Station Table are not equal (this item is allowed only in the **RECEIVE TRAN** construct of the *(receive statement)*):

TRANERR

sets TRUE the (bit variable) TRANERR, and branches control to the next sequential statement.

TRANERR:NULL

causes no action. Execution proceeds as if the

error did not occur.

TRANERR: (label)

sets TRUE the (bit variable) TRANERR, and

branches control to (label).

TRANERR: ABORT

not allowed.

# WRU

The following variations of the WRU syntax cause the following actions if the WRU character of the station is received (the (station WRU character statement) defines the WRU character):

WRU

sets TRUE the WRU (bit variable), and branches control to the next sequential statement.

WRU:NULL

sets TRUE the WRU (bit variable), and execution proceeds as if the character received was not the WRU character.

WRU: \(\label\rangle\)

sets TRUE the \(\langle bit variable \rangle \) WRU, and branches control to \(\langle label \rangle \).

WRU: ABORT

not allowed.

# Definitions **REOUEST**

# Receive Statement - Continued

(single character)

The following variations of the  $\langle single\ character \rangle$  syntax cause the following actions if a character received is equal to the  $single\ character$ :

(single character)

branches control to the next sequential state-

ment.

(single character): NULL

causes no action. Execution proceeds as if the character received was not equal to the \( \single \)

character).

(single character): (label)

branches control to (label).

(single character): ABORT

not allowed.

The allowable combinations of the  $\langle receive\ statement \rangle$  syntax options are defined in table 5-6 below. The (NULL) and ( $\langle timeout\ time \rangle$ ) options are allowed in any construct of the  $\langle receive\ statement \rangle$ . Allowed combinations of the other syntax options are denoted by a "X" in the appropriate columns and rows.

Table 5-6. Allowable Combinations for \( \text{receive statement} \)

ADDRESS(RECEIVE)         X	STOPBIT	LINEDELETE LOSSOFCARRIER	LOSSOFCARRIER PARITY STOPBIT	TIMEOUT TRANERR WRU  single character
ADDRESS(TRANSMIT)         X	X X	X	x x x	X
ADDRESS(TRANSMIT)         X	X X	X	XXX	X
CHARACTER         X	X X	$\mathbf{X}$	$\mathbf{X} \mathbf{X} \mathbf{X}$	Χ .
CRC         X X X X         X X X X           TEXT         X X X X X X X X X X X X X X X X X X X	XX	X	$\mathbf{x} \mathbf{x} \mathbf{x}$	X
TEXT         X	X X	$\mathbf{X} \cdot \mathbf{X}$	$\mathbf{X} \mathbf{X} \mathbf{X}$	<b>x x x</b>
TRAN X X X	$\mathbf{X}^{\top}\mathbf{X}^{\top}$	X	$\mathbf{X} \mathbf{X} \mathbf{X}$	X
	XX	ΧX	$\mathbf{X} \mathbf{X} \mathbf{X}$	$\mathbf{x} = \mathbf{x} \cdot \mathbf{x}$
\( \string \rangle \) X X X X X X X X	X X	$\mathbf{X}$	X X X	X X
	XX	X	x x x	X

# Definitions **REQUEST**

Receive Statement - Continued

**Supplementary Examples** 

Statement

RECEIVE (3 SEC) [TIMEOUT:10].

RECEIVE ADDRESS [ADDERR:99].

RECEIVE CHARACTER [CONTINUE:10, CONTROL:20, TIMEOUT:30, "\*":401.

RECEIVE [ERROR[0]].

RECEIVE [0].

RECEIVE (1 SEC) TEXT [LINEDELETE:NULL, CONTROL:NULL].

# Explanation

Causes the adapter cluster to attempt to receive a character. If the character is not received within 3 seconds, the *(bit variable)* TIMEOUT is set TRUE and control branches to 10.

If the character(s) received do not equal those defined in the *(station address statement)*, the *(bit variable)* ADDERR is set TRUE, and control branches to 99.

This statement would only be allowed in a \(\langle control definition \rangle \) or \(\langle request definition \rangle \) that is written to communicate with full duplex terminal types, because it contains the **CONTINUE** item.

**CONTINUE:10** would cause a branch to 10 if the co-line  $\langle control \ definition \rangle$  executes a  $\langle continue \ statement \rangle$  before a character is received.

**CONTROL:20** would set **CONTROLFLAG TRUE** and branch to **20** if the character received is the station's control character.

TIMEOUT:30 would set TIMEOUT TRUE and branch to 30 if a character is not received within the \(\lambda timeout \time\rangle\) defined in the \(\lambda terminal \) timeout statement\(\rangle\).

"\*":40 would cause a branch to 40 if the character received is the asterisk character.

An attempt is made to receive one character and store it in **CHARACTER**, If any errors described in the associated *(error switch statement)* occur while receiving, then the action defined in that *(error switch statement)* is taken.

An attempt is made to receive one character and store it in CHARACTER. If any errors described in the associated *(error switch statement)* occur while receiving, then the action defined in that *(error switch statement)* is taken.

LINEDELETE:NULL causes the message text pointer to be set to the first character position if the linedelete character (as defined in the \(\lambda terminal linedelete character statement\rangle\) is received, and characters continue to be received and stored in the message text beginning at the first character position.

# **REQUEST**

# Receive Statement - Continued

# Statement

# Explanation

CONTROL:NULL causes the \( \begin{align\*} bit variable \\ CONTROLFLAG \to be set TRUE if the control character of the station (defined in the \( \station \) control character statement \( \)) is received, and characters continue to be received.

# SHIFT STATEMENT Syntax SHIFT UP

# **Semantics**

The  $\langle shift\ statement \rangle$  is to be used in a  $\langle control\ definition \rangle$  that communicates with stations using the Baudot (5-bit) character code set. (The character code set is defined in the  $\langle terminal\ code\ statement \rangle$  of the associated  $\langle terminal\ definition \rangle$ .)

SHIFT UP indicates that received characters are to be translated to their respective uppercase graphics (usually referred to as FIGS).

SHIFT DOWN indicates that received characters are to be translated to their respective lowercase graphics (usually referred to as LTRS).

If the station does not use Baudot code, the (shift statement) acts as a no-op.

# **Pragmatics**

In the Baudot character code set, most bit patterns have two graphic representations; one is referred to as FIGS (the uppercase graphic), and the other as LTRS (the lowercase graphic).

When transmitting to a terminal that uses Baudot code, the terminal prints LTRS until it receives a specially designated character indicating that it should shift to printing FIGS. The terminal continues printing the FIGS until it receives a specially designated character indicating that it should resume printing the LTRS.

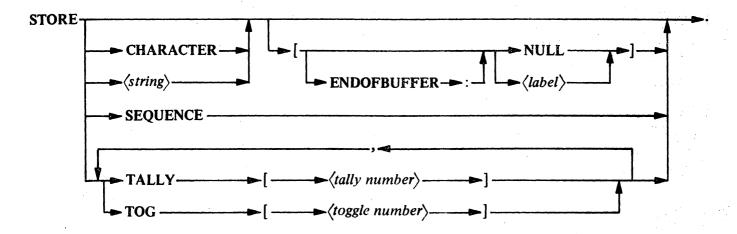
When information is received from a terminal that uses Baudot, the same conventions hold true; that is, the terminal communicates whether FIGS or LTRS follow, by the transmission of a designated character. The terminal initially transmits LTRS.

# REQUEST

Store Statement

#### STORE STATEMENT

Syntax



# **Examples**

STORE.

STORE CHARACTER [ENDOFBUFFER:20].

STORE "ABC" [NULL].

STORE SEQUENCE.

STORE TALLY [0].

**STORE TOG [0], TOG [1], TALLY [0].** 

#### **Semantics**

#### STORE

This form is semantically equivalent to the STORE CHARACTER construct.

# STORE CHARACTER

This form causes the data contained in **CHARACTER** to be stored in the message space. If no message space is associated with the  $\langle request\ definition \rangle$ , then an implicit  $\langle getspace\ statement \rangle$  is executed. The data is stored in the character position pointed to by the message text pointer, and the text pointer is updated after the **STORE** to point to the next forward character position.

It is possible to encounter the end-of-the-text buffer when using this instruction. It is recommended that the optional syntax be included whenever using this statement. The optional syntax specifies action to be taken if the end of buffer is encountered. The NULL option specifies that the only action that should be taken is to set ENDOFBUFFER to TRUE. The \( \lambda \lambda bel \rangle \) option specifies that the only action that should be tered, control should branch to \( \lambda \lambda bel \rangle \) and also set ENDOFBUFFER TRUE. The ENDOFBUFFER: part can be included for documentation. An implicit TERMINATE ERROR is executed if no end-of-buffer action is specified.

# STORE (string)

This form causes  $\langle string \rangle$  to be stored in the message space. If no message space is currently associated with the  $\langle request \ definition \rangle$ , an implicit  $\langle getspace \ instruction \rangle$  is executed. The  $\langle string \rangle$  is stored in the message space beginning at the character position pointed to by the message text pointer, and the text pointer is updated after the **STORE** to point to the first character position following the  $\langle string \rangle$ .

Store Statement - Continued

This instruction uses CHARACTER as a temporary storage area to store each character of (string). Thus, any data in CHARACTER prior to a STORE (string) instruction will be destroyed.

It is possible to encounter the end-of-the-text buffer when using this instruction. Therefore, it is recommended that this instruction include the optional syntax. Refer to the STORE CHARACTER construct for the semantics of this syntax.

# **STORE SEQUENCE**

Providing the station is in sequence mode (i.e., SEQUENCE is TRUE), the STORE SEQUENCE construct causes the current value of the sequence number to be stored in message word [5].[26:27] is a binary integer, and message word [5].[27.1] is set TRUE to indicate its presence. If the station is not in sequence mode (i.e., SEQUENCE is FALSE), then the instruction is a no-op. If no message space is present at the time of the STORE, then an implicit  $\langle getspace\ instruction \rangle$  is executed first.

# **STORE TALLY** [ \( \tally number \) ]

This form causes the TALLY specified to be stored in the message space header. If no message space is present, an implicit (getspace statement) is executed just prior to the store operation.

# **STORE TOG** [ \langle toggle number \rangle ]

This form causes the **TOGGLE** specified to be stored in the message space header. If no message space is present, an implicit (getspace statement) is executed just prior to the store operation.

# **Pragmatics**

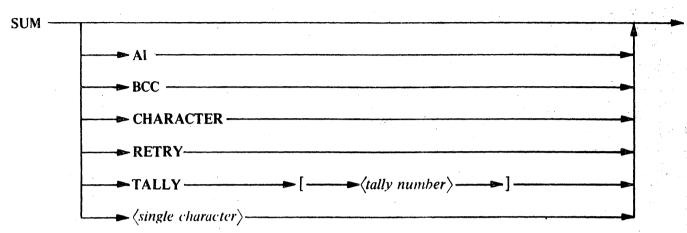
The application of the STORE TALLY and STORE TOG constructs rests solely with the programmer. Since the message space is usually returned to a Message Control System (MCS), some mutual convention could be established between the NDL programmer and the MCS programmer as to the meaning of the contents of the TALLYs and TOGGLEs.

# REQUEST

Sum Statement

#### SUM STATEMENT

Syntax



# Examples

SUM AI. SUM CHARACTER. SUM "A". SUM TALLY [1].

# **Semantics**

The purpose of the \( \sum \) statement \( \) is to affect the calculation of the horizontal parity check (whether that be a Block Check Character or a Cyclic Redundancy Check). The specific effect of the \( \sum \) statement \( \) is dependent upon two factors: The SUMmed item, and whether the station's \( \sum \) terminal \( \definition \) , for which \( \lambda \) request \( \definition \rangle \) is running, defines horizontal parity as \( \mathbb{CRC(16)} \). Following is a description of the \( \sum \) statement \( \rangle \) has on the calculation of the horizontal parity check.

#### **SUM**

Semantically equivalent to SUM CHARACTER.

# SUM AI.

If the horizontal parity check is a Block Check Character or is undefined, the contents of **AI** are exclusively OR-ed with the contents of **BCC**, and the result is stored in **BCC**.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of AI and CRC, and the result is stored in CRC.

# **SUM BCC**

If the horizontal parity check is a Block Check Character or is undefined, then the contents of **BCC** are exclusively OR-ed with itself, and the result is stored in **BCC**. (The result in BCC would be zero in this case.)

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of CRC[0] and CRC, and the result is stored in CRC.

# **SUM CHARACTER**

If the horizontal parity check is a Block Check Character or is undefined, the contents of **CHARACTER** are exclusively OR-ed with the contents of **BCC**, and the result is stored in **BCC**.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of **CHARACTER** and **CRC**, and the result is stored in **CRC**.

#### **SUM RETRY**

If the horizontal parity check is a Block Check Character or is undefined, the contents of RETRY are exclusively OR-ed with the contents of BCC, and the result stored in BCC.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of RETRY and CRC, and the result is stored in CRC.

# SUM TALLY [\langle tally number \rangle]

If the horizontal parity check is a Block Check Character or is undefined, the contents of **TALLY** [\(\lambda tally number\)] are exclusively OR-ed with the contents of **BCC**, and the result is stored in **BCC**.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the contents of TALLY [\langle tally number \rangle \right] and CRC, and the result is stored in CRC.

# **SUM** (single character)

If the horizontal parity check is a Block Check Character or is undefined, the (single character) is exclusively OR-ed with the contents of BCC, and the result is stored in BCC.

If the horizontal parity check is a Cyclic Redundancy Check, the Cyclic Redundancy Check algorithm is computed with the (single character) and CRC, and the result is stored in CRC.

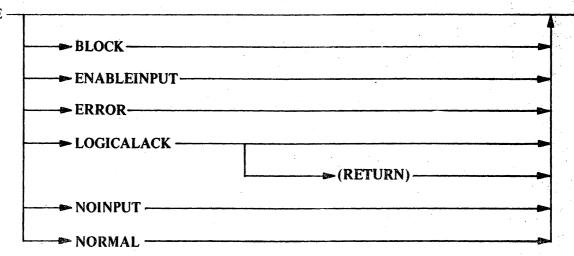
# REQUEST

Terminate Statement

# TERMINATE STATEMENT

**Syntax** 

# TERMINATE -



# Examples

TERMINATE NORMAL.
TERMINATE LOGICALACK.
TERMINATE LOGICALACK(RETURN).
TERMINATE.

#### **Semantics**

Each form of the \(\langle terminate statement \rangle \) is described in the following paragraphs.

# **TERMINATE**

This construct causes control to branch from a  $\langle request \ definition \rangle$  and to begin executing the appropriate  $\langle control \ definition \rangle$ . Any message that may be queued is left in the Station Queue (regardless of whether the message is incoming or outgoing) and STATION(QUEUED) is untouched.

# TERMINATE BLOCK

In a Receive Request, this construct causes the following actions:

- a. an implicit  $\langle getspace\ instruction \rangle$  is executed (in case the  $\langle request\ definition \rangle$  may have been terminated without ever having acquired a message space);
- b. the Error Flag field, Last Flag Set field, and DCP RETRY are stored into the appropriate message fields;
- c. the "More-Blocks-to-Follow" bit (programmatically referenced by BLOCK) in the message (message word [0].[29:1]) is set TRUE;
- d. the message is delinked from the Station Queue and linked into the DCP Result Queue; and
- e. control continues at the next sequential statement.

In a Transmit Request, the TERMINATE BLOCK construct causes the following:

a. the Error Flag field, Last Flag Set field, and DCP RETRY are stored into the appropriate message fields;

- b. the message is linked into the DCP Result Queue; and
- c. the \(\text{request definition}\) is continued at the next sequential statement if STATION(QUEUED) is TRUE; otherwise, the \(\text{request definition}\) is suspended and put in a "sleep" state until STATION(QUEUED) becomes TRUE.

# TERMINATE ENABLEINPUT

This construct is allowed in Transmit Requests only.

This instruction causes the following actions:

- a. the STATION(ENABLED) bit is tested; if STATION(ENABLED) is FALSE, then this instruction acts as a no-op; otherwise, steps b through d are executed;
- b. the Error Flag field, Last Flag Set field, and DCP RETRY are stored into the appropriate message fields;
- c. the message is linked into the DCP Result Queue; and
- d. control leaves the Transmit Request and the station's Receive Request is entered.

# TERMINATE ERROR

This construct serves to inform the station's MCS of an unsuccessful attempt to complete a Receive or Transmit Request. This instruction inhibits the initiation of any new functions for the station.

The result of the **TERMINATE ERROR** construct is as follows:

- a. STATION(READY) bit variable is set FALSE;
- b. a minimum-size message space is obtained, filled with error information for the MCS, and linked into the DCP Result Queue (its destination being the MCS); and
- c. the line is idle until the MCS takes some action.

Additionally, if the TERMINATEing (request definition) was a Receive Request, any message space that may have been acquired to store a RECEIVEd message is discarded.

The error message sent to the MCS contains the following information:

```
MSG[0].[47:8] = 99.

[39:8] = AC Register contents.

[31:8] = AI Register contents.

[23:24] = Logical Station Number.

MSG[1].[47:8] = Result Byte Index

[39:6] = Line status prior to TERMINATE ERROR.

[33:1] = LINE(TOG[1]).

[32:1] = LINE(TOG[0]).

[31:8] = Last Flag Set in MSG[1].[23:24]

[23:24] = Error Flag field.
```

# **REQUEST**

Terminate Statement - Continued

MSG[2].[47:8] = CHARACTER.

[39:16] = Last DCP "Sleep" address.

MSG[4].[23:24] = Original DCWRITE TYPE. (Contains the original contents of MSG[0],[47:24] prior to presentation of the message to the DCP.)

Refer to appendix E, "The Error Result Message," in the B 6700/B 7700 DCALGOL Reference Manual, form number 5000052, for more information regarding this message.

# TERMINATE LOGICALACK

This construct is allowed in Receive Requests only. This instruction tests the LOGICALACK bit in the Station Table. (The semantics of the \( \)station logicalack statement \( \) describe how the LOGICALACK bit is set.) If LOGICALACK is FALSE, the instruction acts as a no-op and control continues at the next sequential statement. If LOGICALACK is TRUE, the following occurs:

- a. an implicit \( \frac{getspace statement}{\} \) is executed (in case the \( \frac{request definition}{\} \) is terminating without ever having acquired any message space);
- b. the ACKNOWLEDGEREADY bit in the Line Table is set (the consequences of this action are described subsequently);
- c. the "Message to be ACKNOWLEDGEd" bit is set in the Error Flag Field;
- d. the Error Flag Field, Last Flag Set Field, and DCP RETRY are stored into the appropriate message fields;
- e. the message is delinked from the Station Queue and linked into the DCP Result Queue;
- f. the line is put in a "sleep" state until the station's MCS responds to the message with an ACKNOWLEDGE (TYPE = 44) DCWRITE; and
- g. upon receipt of the ACKNOWLEDGE, the Receive Request is allowed to continue at the next sequential statement.

The ACKNOWLEDGEREADY bit is inaccessible to the NDL programmer, and it exists for each logical line in its Line Table. The only time that this bit will be TRUE is when a station's LOGICALACK bit is TRUE and its Receive Request has executed the TERMINATE LOGICALACK construct or the TERMINATE LOGICALACK(RETURN) construct. Once TRUE, the ACKNOWLEDGEREADY bit will not be set FALSE, and the \( \frac{request definition}{request} \) will not be allowed to continue until the MCS executes the ACKNOWLEDGE (TYPE = 44) DCWRITE.

# TERMINATE LOGICALACK(RETURN)

This construct is allowed in Receive Requests only. This instruction tests the LOGICALACK bit in the Station Table. (The semantics of the \( \station \logicalack \) statement \( \station \) describe how this bit gets set.) If this bit is found TRUE, this statement functions exactly as does TERMINATE LOGICALACK; refer to that form for semantics. If LOGICALACK is FALSE, the following occurs:

- a. an implicit \( \langle getspace statement \rangle \) is executed (in case the \( \langle request definition \rangle \) is terminating without ever having acquired a message space);
- b. the Error Flag field, Last Flag Set field, and DCP RETRY are stored into the appropriate message fields;
- c. the message is delinked from the Station Queue and linked into the DCP Result Queue; and
- d. control continues at the next sequential instruction in the (request definition).

Terminate Statement - Continued

#### TERMINATE NOINPUT

If executed in a Transmit Request, this form is semantically equivalent to the TERMINATE construct (refer to that construct for semantics). When executed in a Receive Request, the following occurs:

- a. any message space that may have been acquired is discarded;
- b. LINE(BUSY) is set FALSE; and
- c. control branches to the appropriate (control definition).

# TERMINATE NORMAL

The purpose of this construct is to signal the satisfactory completion of a  $\langle request \ definition \rangle$ . If executed in a Receive Request, the following occurs:

- a. an implicit  $\langle getspacestatement \rangle$  is executed (in case the  $\langle request definition \rangle$  is terminating without having ever acquired a message space);
- b. the Error Flag field, Last Flag Set field, and DCP ENTRY are stored into the appropriate message fields;
- c. the message space is delinked from the Station Queue and linked to the DCP Result Queue;
- d. LINE(BUSY) is set FALSE; and
- e. control branches from the \( \frac{request definition} \) and (providing the DCP does not take advantage of LINE(BUSY) set FALSE to initiate a \( \frac{request definition} \)) the appropriate \( \frac{control}{definition} \) is entered.

If TERMINATE NORMAL is executed in a Transmit Request, the following occurs:

- a. the Error Flag field, Last Flag Set field, and DCP RETRY are stored into the appropriate message fields;
- b. the message is linked into the DCP Result Oueue;
- c. LINE(BUSY) is set FALSE; and
- d. control branches from the \( \text{request definition} \) and (providing the DCP does not take advantage of LINE BUSY) set FALSE to initiate a \( \text{request definition} \) the appropriate \( \text{control definition} \) is entered.

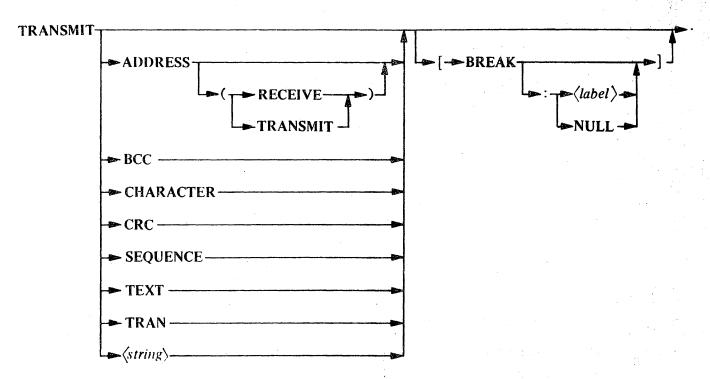
In the Transmit Request case, the message linked to the DCP Result Queue is a result message (specifically, a GOOD RESULTS (CLASS = 5) Message). The intended destination is the MCS; however, the MCS has the option of whether to accept GOOD RESULTS Messages or to have the DCC discard them.

# **REQUEST**

Transmit Statement

#### TRANSMIT STATEMENT

Syntax



# Examples

TRANSMIT.

TRANSMIT CHARACTER [BREAK:NULL].

TRANSMIT SOH STX 4"00" [BREAK:10].

TRANSMIT TRAN.

TRANSMIT ADDRESS(TRANSMIT)[BREAK].

TRANSMIT TEXT[BREAK].

TRANSMIT "LITERAL STRING".

#### **Semantics**

The \(\langle transmit statement\)\ causes the adapter cluster to transmit information to a terminal. The following group of syntax options specifies the information to be transmitted. All options, except CHARACTER, use the \(\langle b\) te variable \(\rangle\) CHARACTER as a temporary storage area; thus, any information contained in CHARACTER before execution of the \(\langle transmit statement\)\) shall be destroyed by the \(\langle transmit statement\). If none of the first group of options are chosen, it is semantically equivalent to specifying CHARACTER (i.e., TRANSMIT is equivalent to TRANSMIT CHARACTER).

# **ADDRESS**

The proper number of characters (as specified by the station's \(\lambda\text{terminal definition}\rangle\) in the \(\lambda\text{terminal address size statement}\rangle\)) are taken from the Address field in the Station Table and transmitted.

# ADDRESS(RECEIVE)

This option is equivalent to ADDRESS, except that ADDRESS(RECEIVE) must be used when an address pair is defined in the *(station address statement)* and the programmer wants to transmit the receive address.

# ADDRESS(TRANSMIT)

This option is equivalent to ADDRESS, except that ADDRESS(TRANSMIT) must be used when an address pair is defined in the \( \station address statement \) and the programmer wants the transmit address transmitted.

# **BCC**

The BCC option causes the content of the \( \langle byte variable \rangle \) BCC to be transmitted.

#### **CHARACTER**

The CHARACTER option causes the content of the (byte variable) CHARACTER to be transmitted.

# **CRC**

This option causes two bytes to be transmitted; CRC[0] is transmitted first, followed by CRC[1]. If the station's \(\lambda terminal definition\rangle\) does not define horizontal parity as CRC(16), the use of this option causes a syntax error to be generated at compile time.

# **SEQUENCE**

This option causes the character representation of the value stored in the Sequence field of the Station Table to be transmitted if the station is in sequence mode (i.e., the \( \frac{bit variable}{\} \) SEQUENCE is TRUE); otherwise, the \( \frac{transmit statement}{\} \) is a no-op.

# **TEXT**

This option extracts characters, one at a time, from the associated message, using CHARACTER as a temporary storage area, and transmits the characters until the end of the text buffer is encountered. At that point, control branches to the next statement. The TRANSMIT TEXT construct is, in effect, the same as the following loop:

1: FETCH [ENDOFBUFFER:2]. TRANSMIT CHARACTER. GO TO 1.

2:

This option can only be used with the \(\lambda transmit statement \rangle \) in Transmit Requests.

#### TRAN

The proper number of transmission number characters (as defined by the station's \(\lambda terminal definition\rangle\) in the \(\lambda terminal transmission number length statement\rangle\)) are extracted from the Transmit Transmission Number field in the Station Table and then transmitted.

*(string)* 

Fach character of (string), using CHARACTER as a temporary storage area, is transmitted.

# REQUEST

# Transmit Statement - Continued

The BREAK option allows the programmer to specify action if a "break" is received from the terminal while the adapter cluster is still transmitting. If this option is omitted and a break occurs, an implicit TERMINATE ERROR instruction is executed. The following describes the actions of the three syntactical forms:

BREAK sets TRUE the (bit variable) BREAK [TRANSMIT] and causes a branch of

control to the next statement.

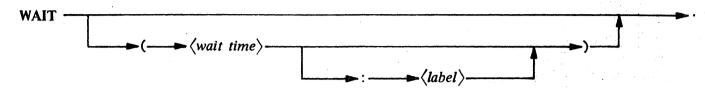
BREAK: \(\langle \langle \text{label}\rangle \) sets TRUE the \(\langle \text{bit variable} \rangle \text{BREAK[TRANSMIT]}\) and causes a branch of

control to \(\langle label \rangle.

BREAK: NULL causes no action. Execution proceeds as if the break did not occur.

# WAIT STATEMENT

# **Syntax**



# Examples

WAIT. WAIT (3 SEC). WAIT (5 MILLI:6).

# **Semantics**

The  $\langle wait \ statement \rangle$  is only allowed in  $\langle request \ definition \rangle s$  that are written to communicate with full duplex terminal types. Execution of this statement causes the  $\langle request \ definition \rangle$  to be suspended until the co-line executes a  $\langle continue \ statement \rangle$ . The optional syntax effects the statement as described below.

(wait time)

defines the maximum amount of  $\langle time \rangle$  that the  $\langle request\ definition \rangle$  should be suspending waiting for the  $\langle continue\ statement \rangle$ . If  $\langle wait\ time \rangle$  is exceeded and the co-line has not executed a  $\langle continue\ statement \rangle$ , execution resumes at the next sequential statement.

\langle wait time \rangle : \langle label \rangle

same as above except execution resumes at \( \lambda label \rangle \) if a \( \lambda continue \) statement \( \rangle \) is not executed within \( \lambda wait \) time \( \rangle \).

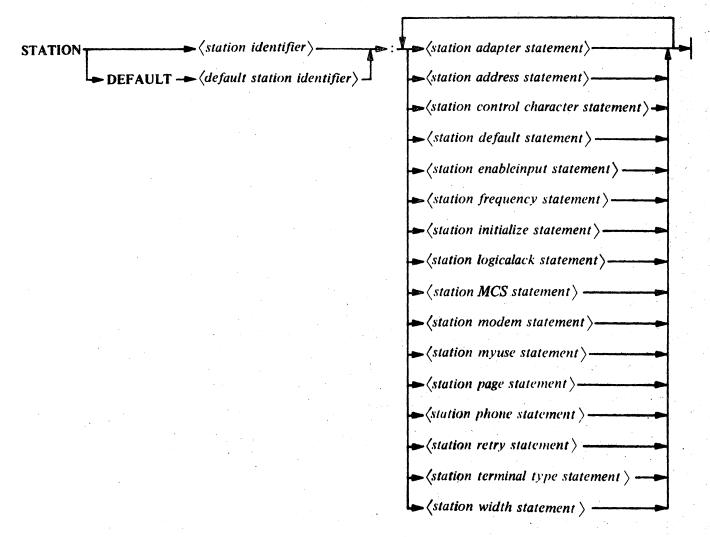
# **Pragmatics**

Refer to the  $\langle fork \ statement \rangle$  pragmatics.

# **STATION**

# STATION DEFINITION

**Syntax** 



# **Examples**

# STATION KMET:

ENABLEINPUT = FALSE.

MCS = SYSTEM/CANDE.

CONTROL = 4"6F". RETRY = 15. LOGICALACK = FALSE.

MYUSE = INPUT, OUTPUT. TERMINAL = TELETYPE.

# STATION DEFAULT STADFLT2:

CONTROL = "?".

MCS = SYSTEM/CANDE.

ADAPTER = 4.

DEFAULT = STADFLT1.

# **Semantics**

(station identifier) and (default station identifier) have the syntactical form of a (system identifier), Each syntactical form of the (station definition) is described subsequently.

# **STATION** \(\station identifier\): . . .

This form of the \( \station \) defines the attributes of a station. The attributes must be defined in one of the following ways:

- a. Each attribute is explicitly defined by means of a (station statement).
- b. Each attribute is defined implicitly by means of an explicit reference to a set of previously defined default attribute values.
- c. Some of the attributes are defined implicitly as in b, and the remainder are defined explicitly as in a.

Some of the station attributes must be defined for each station; others do not. Some of the statements may or may not be required, depending upon the appearance of other statements. The semantics portion of each (station statement) states, among other things, whether the attribute must be defined and its effect upon the requirements of other (station statement)s.

To define the attributes of a station as described in item a above, only this syntax form is used.

To define the attributes of a station as described in items b and c above, this syntax form, the following syntax form, and the (station default statement) must be used in conjunction (this is described under the following syntax form).

# STATION DEFAULT (default station identifier): ...

This form is referred to as a Default  $\langle station \ identifier \rangle$ . Its purpose is to decrease the number of source statements required to define all of the stations. This is accomplished in the following manner. Attributes common to several stations are defined by means of a Default  $\langle station \ definition \rangle$ . Associated with each Default  $\langle station \ definition \rangle$  is a  $\langle default \ station \ identifier \rangle$ . Subsequent to the Default  $\langle station \ definition \rangle$ , any  $\langle station \ definition \rangle$  can reference the  $\langle default \ station \ identifier \rangle$ , instead of repeating the list. A  $\langle default \ station \ identifier \rangle$  is referenced by means of a  $\langle station \ default \ statement \rangle$ . The NDL compiler uses the last definition of a station attribute, and therefore the programmer can reference a Default  $\langle station \ definition \rangle$  and change any attributes by redefining them in the  $\langle station \ definition \rangle$ .

In appearance, the Default (station definition) is similar to the (station definition). The differences are that the reserved word DEFAULT follows the reserved word STATION, and that there are no statements that are required to appear in a Default (station definition).

# **STATION**

Station Adapter Statement

# STATION ADAPTER STATEMENT

Syntax

Examples

ADAPTER = 
$$4$$
.  
ADAPTER =  $(11.6)$ .

#### Semantics

The \(\station\) adapter statement\) defines a combination of character format, synchronous/asynchronous communication, and line speed (in the case of synchronous communications) that the DCP must use to communicate with the terminal associated with the station. This is done by supplying a \(\chicommunication\) type number\(\rangle\) (or number pair). Table 5-4 lists the allowed \(\chicommunication\) type number\(\rangle\)s and the characteristics associated with each.

For example,

$$ADAPTER = 4.$$

This statement defines an 11-bit character format, asynchronous communication, at a line speed of 110 bits per second.

If the station's associated terminal type utilizes full duplex (i.e., the \(\lambda\) terminal duplex statement\(\rangle\) specifies DUPLEX=TRUE), and the primary and the auxiliary lines have different characteristics, then a \(\lambda\) communication type number\(\rangle\) pair must be supplied.

For example,

ADAPTER = 
$$(11,6)$$
.

This statement defines for the primary line a 10-bit character format, asynchronous communication, at a speed of 1800 bits per second. The characteristics associated with the auxiliary line are the same except that it runs at a line speed of 150 bits per second.

The statement:

$$ADAPTER = (6.6).$$

is semantically equivalent to:

$$ADAPTER = 6.$$

The \(\communication type number\)\(\) (or number pair) defined in this statement must be one of those listed in the \(\lambda terminal adapter statement\rangle\) of the station's associated \(\lambda terminal definition\rangle\). The \(\lambda station adapter statement\rangle\) is required unless the \(\lambda terminal adapter statement\rangle\) lists only one \(\lambda communication type number\rangle\) (or number pair), in which case, the \(\lambda station adapter statement\rangle\) may be omitted and the \(\lambda terminal adapter statement\rangle\) specification is used.

# Definitions STATION

Station Adapter Statement - Continued

# **Supplementary Examples**

The following program fragments illustrate valid adapter statement specifications.

# Example 1

# **MODEM AMODEM:**

ADAPTER=1,2,3,4,5,6,7,8,9,10.

# **TERMINAL ATERMINAL:**

ADAPTER=6,7,8,9,11,12,13,14,15.

# **STATION ASTATION:**

ADAPTER=7.
MODEM=AMODEM.
TERMINAL=ATERMINAL.

# Example 2

# **MODEM DUPLEXMODEM:**

ADAPTER=6, (11,6), (12,6), 12.

# TERMINAL DUPLEXTERMINAL:

ADAPTER=6, (11,6).

# STATION DUPLEXSTATION:

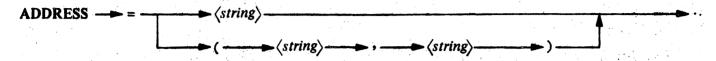
ADAPTER=(11,6).
MODEM=DUPLEXMODEM.
TERMINAL=DUPLEXTERMINAL.

# STATION

Station Address Statement

# STATION ADDRESS STATEMENT

# **Syntax**



# **Examples**

#### **Semantics**

The  $\langle station \ address \ statement \rangle$  defines the actual address characters of the station's terminal that are required for operations such as polling and selecting. The number of characters in the  $\langle string \rangle (s)$  must be equal to the number defined in the  $\langle terminal \ address \ size \ statement \rangle$  of the associated terminal. This statement is not allowed in Default  $\langle station \ definitions \rangle$ .

$$ADDRESS = \langle string \rangle$$
.

This form of the statement is used when the receive address and the transmit address are the same.

ADDRESS = 
$$(\langle string \rangle, \langle string \rangle)$$
.

This form of the statement must be used if the receive address and the transmit address differ. The first \( \string \ranger\) defines the receive address characters, and the second \( \string \ranger\) defines the transmit address characters.

# **Pragmatics**

The address characters of a station can be changed as a result of the Message Control System (MCS) executing a SET CHARACTERS (TYPE = 39) DCWRITE.

# Definitions **STATION**

Station Control Character Statement

# STATION CONTROL CHARACTER STATEMENT

**Syntax** 

CONTROL → = → ⟨single character⟩

Example

CONTROL = "?".

# **Semantics**

The \(\station\) control character statement\(\righta\) defines the control character of the station. The control character can be recognized by the DCP when RECEIVEd in a message text from the station, and any action to be taken can be specified by the programmer using the CONTROL syntax in the \(\scalegar{receive}\) statement\(\righta\).

#### STATION

Station Default Statement

# STATION DEFAULT STATEMENT

**Syntax** 

# Example

DEFAULT = STADFLT1.

# **Semantics**

The \( \station default statement \) allows the programmer to specify the \( \default station identifier \) of a set of previously defined default station attributes to be used for a \( \station definition \) whose description is incomplete. It is advantageous to group common attributes under a Default \( \station definition \) and list the remaining attributes under each individual \( \station definition \). The compiler will then refer to the Default \( \station definition \) to complete the \( \station definition \). The \( \station default statement \) is not required to appear in each \( \station definition \); however, a \( \station definition \) must define all required attributes locally if a \( \tation default statement \) does not appear.

The (station default statement) can appear in a (station definition) or a Default (station definition).

# Supplementary Example

The following is an example of how a Default  $\langle station \ definition \rangle$  can be used in conjunction with a  $\langle station \ definition \rangle$ .

# STATION DEFAULT STADFLT:

MCS = SYSTEM/CANDE.
CONTROL = "?".

RETRY = 15.

LOGICALACK = FALSE.

MYUSE = INPUT, OUTPUT.

TERMINAL = TELETYPE.

LISE. {Default (station definition)}

**ENABLEINPUT = TRUE.** 

# STATION TESTSTATION:

DEFAULT = STADFLT. MODEM = MABELL103A.

MCS = SYSTEM/DIAGNOTICMCS.

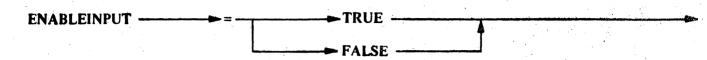
ADAPTER = 4

(station default statement) references Default (station definition) above to complete the (station definition).

Station Enableinput Statemen

#### STATION ENABLEINPUT STATEMENT

**Syntax** 



#### Semantics

The \(\station\) enableinput statement\) defines the initial state of the station's "enabled" bit (programmatically referred to as STATION(ENABLED)). This statement must be defined in each \(\station\) definition\(\station\).

# **ENABLEINPUT = TRUE.**

This construct causes the "enabled" bit to be initially TRUE after DCP initialization, and the station is said to be "enabled for input," or simply "enabled."

# **ENABLEINPUT = FALSE.**

This construct causes the "enabled" bit to be initially FALSE after DCP initialization, and the station is said to be "disabled for input," or "disabled."

# **Pragmatics**

Whether a station is enabled or disabled for input can directly affect the execution sequence of instructions in the (control definition) and (request definition) (s) designated for that station. Specifically, if the station is disabled for input, control will never branch to the Receive Request for that station as a result of either an INITIATE ENABLEINPUT or a TERMINATE ENABLEINPUT construct. Refer to the INITIATE ENABLEINPUT and TERMINATE ENABLEINPUT constructs for more detailed information.

The MCS of the station may change the state of the "enabled" bit, after DCP initialization, by means of the ENABLE INPUT (TYPE = 35) DCWRITE or the DISABLE INPUT (TYPE = 36) DCWRITE.

# **STATION**

**Station Frequency Statement** 

STATION FREQUENCY STATEMENT

**Syntax** 

Example

FREOUENCY = 10.

#### Semantics

The \(\station\) frequency statement\(\righta\) defines the initial value of the \(\lambda\) byte variable\(\righta\) programmatically referred to as STATION (FREQUENCY). The \(\lambda\) control definition\(\righta\) specified for the station can reference the \(\lambda\) byte variable\(\righta\) and use the value stored there in any way that the programmer sees fit; however, the intended use of the variable is to influence in some way the rate at which a polled station is polled. In the polling \(\lambda\) control definition\(\righta\) provided by Burroughs Corporation in SYMBOL/SOURCENDL, FREQUENCY specifies a relative polling rate: 0 means poll at the highest rate, 1 means to poll at a slower fate, 2 means to poll at a still lower rate, etc.

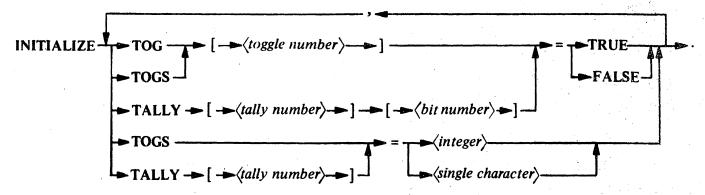
The (integer) must not exceed a value of 255.

The MCS of the defined station can change the value of **STATION**(**FREQUENCY**) by means of an ENABLE INPUT (TYPE = 35) DCWRITE.

# Station Initialize Statement

# STATION INITIALIZE STATEMENT

# **Syntax**



# **Examples**

```
INITIALIZE TOGS = 4"FF".
INITIALIZE TOG|0| = TRUE,
TALLY|1| = 25,
TALLY|0| |7| = TRUE.
INITIALIZE TALLY|0| = "?".
```

# **Semantics**

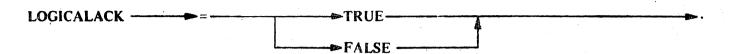
The \(\langle station initialize statement \rangle \) provides the means to define initial values for the station TOGGLEs and TALLYs. Any initial values defined for station TOGGLEs and TALLYs are stored in the TOGGLEs and TALLYs at DCP initialization time only.

# Definitions STATION

Station Logicalack Statement

STATION LOGICALACK STATEMENT

**Syntax** 



# **Semantics**

The \(\station\) logiculack statement\) defines the initial state of a bit, referred to as the Logicalack bit, in the Station Table. TRUE or FALSE can be specified, indicating the initial state as on or off, respectively. If the LOGICALACK bit is on, special action is taken if the Receive Request executes either the TERMINATE LOGICALACK or TERMINATE LOGICALACK(RETURN) constructs of the \(\lambda\) terminate statement\(\rangle\). This statement is required in \(\station\) definition\(\rangle\)s.

The MCS of the station can change the value of the Logicalack bit after DCP initialization by means of the SET/RESET LOGICALACK (TYPE = 43) DCWRITE.

# STATION MCS STATEMENT

**Syntax** 

# **Examples**

MCS = SYSTEM/RJE. MCS = SYSTEM/APL. MCS = UTILITY/MCS.

#### **Semantics**

The  $\langle station \ MCS \ statement \rangle$  defines the Message Control System (MCS) that is responsible for handling messages to and from the station. If the MCS named is not an MCS defined in a  $\langle MCS \ definition \rangle$ , it is added to the list of valid MCS programs to be contained in the Network Information File, and the MCS will not be allowed to execute diagnostic DCWRITEs. Refer to the semantics of the  $\langle MCS \ definition \rangle$  for information regarding the diagnostic DCWRITEs. This statement is required in  $\langle station \ definition \rangle s$ .

#### **STATION**

Station Modem Statement

#### STATION MODEM STATEMENT

**Syntax** 

Example

MODEM = BELL201.

#### **Semantics**

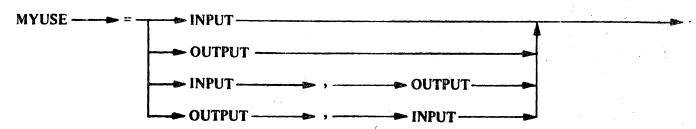
The \( \station \) modem statement \( \) applies to a station that has associated with it a terminal type that must communicate with the Data Communications System through the use of a modem. This statement associates the modem type (i.e., a \( \) modem definition \( \)) used for that purpose with the station. If this statement is omitted from the \( \) station definition \( \), and the \( \) line definition \( \) for the line to which the station is assigned (if, in fact, the station is assigned to a line) does not contain a \( \) line modem statement \( \), then the compiler assumes a direct connection between the terminal and the line adapter.

The \(\)modern identifier\) must name a \(\)modern definition\) that is compatible with the defined station attributes. To be more specific, the \(\)communication type number\) specified in the \(\)station adapter statement\) (or in the \(\)terminal adapter statement\) of the station's \(\)terminal definition\) if no \(\)station adapter statement\) appears) must be one of the \(\)communication type number\) s listed in the \(\)modern adapter statement\) of the modern named.

After DCP initialization, the MCS of the station may change the *(modem definition)* associated with the station, by means of the MOVE/ADD/SUBTRACT STATION (TYPE = 130) DCWRITE.

#### STATION MYUSE STATEMENT

#### **Syntax**



# **Semantics**

The \(\station\) myuse statement\(\righta\) defines to what extent an object job can use the station as in input or output device.

MYUSE=INPUT specifies that an object job can use the station as an input file only.

MYUSE=OUTPUT specifies that an object can use the station as an output file only.

MYUSE=INPUT,OUTPUT or MYUSE=OUTPUT,INPUT specifies that an object job can use the station as an input and/or output file.

A \(\langle terminal request statement \rangle\) must be defined by the station's \(\langle terminal definition \rangle\) for handling input and/or output capabilities as specified in the \(\langle station myuse statement \rangle\). Thus, if the station is to send input to, and receive output from, an object job, the station's \(\langle terminal definition \rangle\) must specify a Transmit Request and a Receive Request.

Note that the *(station myuse statement)* restricts the use of the station by object jobs only. The MCS can communicate with the station to the extent specified in the *(terminal request statement)* of the station's *(terminal definition)*. That is, regardless of what is specified in the *(station myuse statement)*, the MCS can receive information from, or send information to, a station, provided that the station's *(terminal definition)* specified a Receive and Transmit Request.

The station MYUSE attribute can be interrogated by an object job through reference to the MYUSE file attribute. For further information, refer to the <u>B 6700 Input/Output Subsystem Information Manual</u>, form number 5000185.

**STATION** 

**Station Page Statement** 

#### STATION PAGE STATEMENT

**Syntax** 

### **Examples**

**PAGE = 12. PAGE = 0.** 

## **Semantics**

The \(\station\) page statement\(\rangle\) defines the number of logical lines per logical page. The \(\langle\)integer\(\rangle\) specified must be less than or equal to the number of lines specified in the \(\langle\)terminal page statement\(\rangle\) of the station's \(\langle\)terminal definition\(\rangle\) (unless that number is zero, indicating pagination is arbitrary). If a \(\langle\)station page statement\(\rangle\) is not included in the \(\langle\)station definition\(\rangle\), the station's \(\langle\)terminal definition\(\rangle\) specifications for pagination are used.

An object job may obtain the PAGE value of a station, if the station is attached to a file, and that file is open, by interrogating the PAGESIZE file attribute and supplying the File Relative Station Number (FRSN). Refer to the PAGESIZE attribute in the <u>B 6700 Input/Output Subsystem Information Manual</u>, form number 5000185, for more information.

ST	A'	ri	റ	N	Pŀ	10	N	E	ST	A	TEN	/	FN	IT

**Syntax** 

# Example

**PHONE = 12136572385.** 

# **Semantics**

The *(station phone statement)* is implemented for documentation purposes only. This statement documents the telephone number that the system would have to dial to reach the station's terminal.

STATION	v*					
Station Retry Statement						
STATION RETRY STATEMENT	·					
Syntax						
RETRY ==		· (integer)	)		<del></del>	
Example						
RETRY = 3.					•	

Semantics

**Definitions** 

The (station retry statement) defines a default value for DCP INITIAL RETRY. Refer to the RETRY (hvte variable) for more information.

# Definitions STATION Station Terminal Type Statement

# STATION TERMINAL TYPE STATEMENT

**Syntax** 

# **Examples**

TERMINAL = APLTERM. TERMINAL = TTY.

# **Semantics**

The (station terminal type statement) associates a terminal type with the station. This statement is required in a (station definition).

After DCP initialization, the MCS of the station can change the terminal type associated with the station by means of the MOVE/ADD/SUBTRACT STATION (TYPE = 130) DCWRITE.

<b>Definitions</b>	
STATION	·:••

Station Width Statement

# STATION WIDTH STATEMENT

**Syntax** 

to the second of						
WIDTH -	<b>&gt;</b> = -	<del></del>	 integer)	) ———		·

# **Examples**

WIDTH = 72. WIDTH = 132.

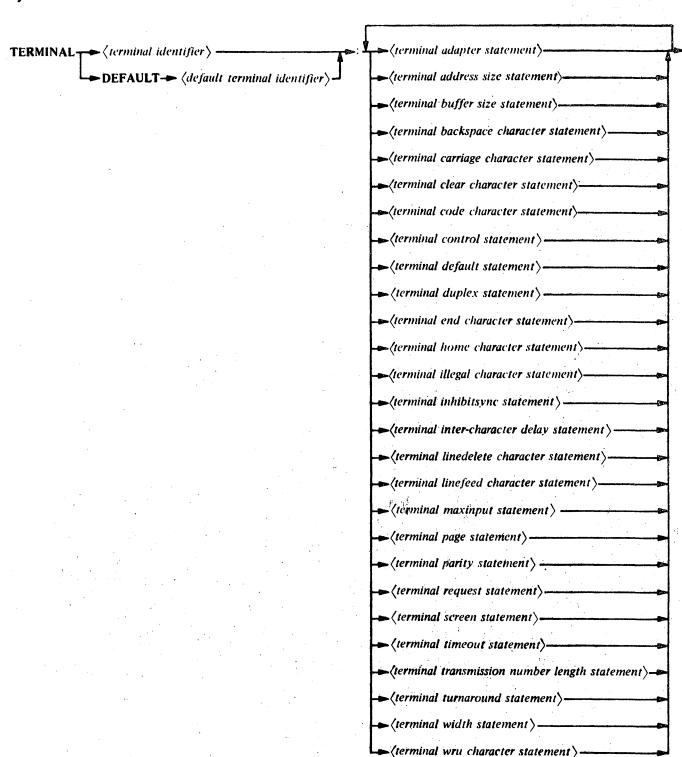
# **Semantics**

The (station width statement) defines the number of characters in a logical display line of output on the station's terminal. If this statement is not included in a (station definition), then the WIDTH defined for the station's (terminal definition) is the default station WIDTH.

An object job can interrogate a station's WIDTH by testing the value of the WIDTH file attribute. Refer to the <u>B 6700 Input/Output Subsystem Information Manual</u>, form number 5000185, for further information.

# **TERMINAL DEFINITION**

**Syntax** 



# Definitions TERMINAL Continued

# Examples

# **TERMINAL TTY:**

CODE ASC67. **PARITY** NULL. FALSE. **SCREEN** BUFFER NULL. **DUPLEX** FALSE. **ADDRESS** NULL. 72: WIDTH **MAXINPUT** 72. TIMEOUT 300 SEC. REQUEST RECEIVE: READTTY.

TRANSMIT: WRITETTY.

**CONTROL** = **CONTENTIONDEVICE**.

# **TERMINAL DEFAULT DEFAULTLIST1:**

CODE = ASC67.
PARITY = NULL.
SCREEN = FALSE.
BUFFER = NULL.

#### Semantics

 $\langle terminal\ identifier \rangle$  and  $\langle default\ terminal\ identifier \rangle$  each have the syntactic form of  $\langle identifier \rangle$ .

Each construct of the  $\langle terminal \ definition \rangle$  is described subsequently.

# TERMINAL (terminal identifier):...

This form of the \(\lambda terminal definition\rangle\) syntax defines the attributes of a terminal type in the data communications network. Most terminal attributes are hardware-dependent. The attributes of the terminal type are defined in one of the following ways:

- a. Each attribute is defined explicitly by means of a \(\lambda\terminal\) attribute statement \(\rangle\) in the \(\lambda\terminal\) definition \(\rangle\).
- b. Each attribute is defined implicitly by an explicit reference to a set of previously defined default attribute values.
- c. Some of the attributes are defined implicitly as in b, and the remainder are defined explicitly as in a

Some of the \(\lambda terminal statement \rangle s\) must be defined for each \(\lambda terminal definition \rangle\); others do not. Some of the statements may or may not be required, depending upon the appearance of other statements. The semantics portion of each \(\lambda terminal attribute statement \rangle\) states, among other things, whether the attribute must be defined and its effect upon the requirement of other attribute definitions.

To define the attributes of a **TERMINAL** as described in item a above, this syntax form must be used.

To define the attributes of a terminal type as described in items b and c above, this syntax form, the following syntax form, and the \(\lambda terminal default statement\rangle\) must be used in conjunction (this is described under the following syntax form).

# **TERMINAL DEFAULT** \(\langle default terminal identifier \rangle : \ldots \).

This form is referred to as a Default \(\lambda terminal definition \rangle.\)

Its purpose is to decrease the number of source statements required to define all of the terminal types in the data communications system. This is accomplished in the following manner. Attributes common to several terminal types are defined by means of a Default \(\lambda terminal definition\rangle\). Associated with each Default \(\lambda terminal definition\rangle\) is a \(\lambda default terminal identifier\rangle\). Subsequent to the Default \(\lambda terminal definition\rangle\) any \(\lambda terminal definition\rangle\) that has those attributes in common may reference the \(\lambda default terminal identifier\rangle\) is referenced by means of a \(\lambda terminal default statement\rangle\).) The NDL compiler uses the last definition of a terminal attribute, and therefore the programmer can reference a Default \(\lambda terminal definition\rangle\) and change any attributes by redefining them in the \(\lambda terminal definition\rangle\).

In appearance, the Default \(\lambda terminal definition\rangle\) is similar to the \(\lambda terminal definition\rangle\). The differences are that the reserved word DEFAULT follows the reserved word TERMINAL, and that no statements are required to appear in a Default \(\lambda terminal definition\rangle\).

# Supplementary Example

Below is an example of how a Default \(\lambda terminal definition \rangle \) can be used in conjunction with a \(\lambda terminal definition \rangle \).

# TERMINAL DEFAULT DEFAULTLIST1:

CODE = ASC67.
PARITY = NULL.
SCREEN = FALSE.
BUFFER = NULL.

**DEFAULTLIST1** is the \( \default \)
terminal identifier \( \) of this Default \( \text{terminal definition} \). The set of default attributes that follows is referenced by this name.

\(\langle \text{terminal statement} \rangle s \text{ define the default attributes associated with DEFAULTLIST!.}

Above, DEFAULTLIST1 has associated with it four attributes. Any subsequent \(\lambda terminal definition \rangle\) in a source program can reference these default attributes by the appearance of a \(\lambda terminal default \) statement \(\rangle\) in the \(\lambda terminal definition \rangle\). The \(\lambda terminal default \) statement \(\rangle\) has the form:

where the \(\langle default\) terminal identifier \rangle\) must name a previously defined Default \(\langle terminal\) definition \(\rangle\). More information regarding the use of Default \(\langle terminal\) definition \(\rangle\) s in conjunction with \(\langle terminal\) default statement \(\rangle\) s can be found in the \(\langle\) terminal default statement \(\rangle\) semantics.

Below, TTY uses the \(\lambda terminal default statement\rangle\) to reference DEFAULTLIST1. DEFAULTLIST1 contains the attribute information required to complete the \(\lambda terminal definition\rangle\) TTY.

# **TERMINAL TTY:**

DEFAULT = DEFAULTLIST1. ← This \( \text{terminal default statement} \)

DUPLEX = FALSE. ADDRESS = NULL. 

This \( \text{terminal default statement} \)

definition \( \text{definition} \)

defined previously.

WIDTH = 72. MAXINPUT = 72.

TIMEOUT = 300 SEC.

**CONTROL** = **CONTENTIONDEVICE**.

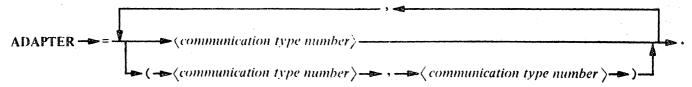
REQUEST = RECEIVE: READTTY, TRANSMIT: WRITETTY.

#### TERMINAL

Terminal Adapter Statement

#### TERMINAL ADAPTER STATEMENT

# Syntax



# Examples

ADAPTER = 4. ADAPTER = (6,10), (10,6). ADAPTER = 5, (5,6), 6.

#### Semantics

The  $\langle terminal\ adapter\ statement \rangle$  defines one or more combinations of character format, synchronous/asynchronous communication, and line speed (in the case of asynchronous communications), with which the terminal type is compatible. This is done by supplying one or more  $\langle communication\ type\ number \rangle_s$  (or number pairs). Table 5-4 lists the allowed  $\langle communication\ type\ number \rangle_s$  and the characteristics associated with each.

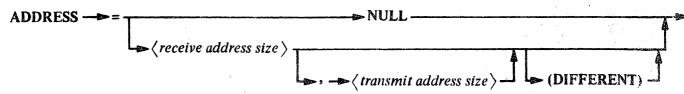
If the terminal type is to be operated in a full duplex mode, and the primary and the auxiliary lines have different characteristics, then a *(communication type number)* pair must be supplied.

If the terminal is to be modem-connected (i.e., connected to the system through the use of modems), then at least one of the  $\langle communication\ type\ number \rangle s$  (or number pairs) must be compatible with those numbers listed for the connecting modem in the  $\langle modem\ adapter\ statement \rangle$ .

Terminal Address Size Statement

#### TERMINAL ADDRESS SIZE STATEMENT

**Syntax** 



## Examples

ADDRESS = 2.

ADDRESS = 2 (DIFFERENT).

ADDRESS = 3, 2 (DIFFERENT).

ADDRESS = 2.3.

#### Semantics

The \(\text{terminal address size statement}\) defines the number of address characters that the terminal type transmits and receives. The number of address characters must not be confused with the actual address \(\text{characters}\) used in polling and selecting; the \(\station\) address statement \(\rightarrow\) defines the actual address characters. This attribute must be defined when actual address characters are defined in the \(\station\) address statement \(\rightarrow\) of a \(\station\) definition \(\rightarrow\) that references the \(\text{terminal definition}\).

\(\text{receive address size}\) and \(\text{transmit address size}\) must be integers greater than zero and less than 4. The \(\text{receive address size}\) defines the number of address characters the terminal expects to receive, and the \(\text{transmit address size}\) defines the number of address characters that the terminal transmits. If the \(\text{transmit address size}\) is not defined, it is assumed the \(\text{transmit address size}\) is equal in length to the \(\text{receive address size}\). The \(\text{receive address size}\) and the \(\text{transmit address size}\) for a given terminal must concur with the length of the character \(\lambda \text{string}\)(s) defined as the actual address characters in the \(\text{station address statement}\) of any \(\lambda \text{station definition}\) which references the \(\text{terminal definition}\); otherwise, a syntax error results.

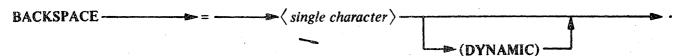
The (DIFFERENT) option must be used if the *(receive address)* and the *(transmit address)*, as defined in the *(station address statement)*, are not identical.

#### TERMINAL

Terminal Backspace Character Statement

#### TERMINAL BACKSPACE CHARACTER STATEMENT

Syntax



# Examples

BACKSPACE = 4"16".
BACKSPACE = "-" (DYNAMIC).

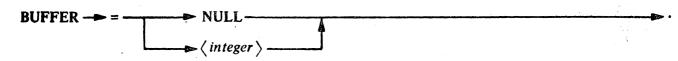
#### Semantics

The \(\langle terminal backspace character statement\rangle\) defines the backspace character of the terminal type (i.e., the character that the terminal type would transmit to indicate that the previous character should be deleted). If defined, the backspace character can be recognized by the DCP when RECEIVEd (in a \(\langle receive statement \rangle\)), and any action to be taken can be specified by the programmer (using the BACKSPACE syntax).

(DYNAMIC) indicates that the controlling MCS of a station referencing the \(\lambda\) terminal definition\(\rangle\) is allowed to change the backspace character for the station by means of a SET CHARACTERS (TYPE=39) DCWRITE.

# **TERMINAL BUFFER SIZE STATEMENT**

**Syntax** 



# **Examples**

BUFFER = NULL. BUFFER = 960.

# **Semantics**

The \(\lambda\terminal\) buffer size statement \(\rangle\) applies to buffered devices and defines the size, in characters, of the terminal type buffer. If the terminal type is an unbuffered device, the form:

# **BUFFER = NULL.**

can be used, or the statement may be omitted; additionally, if the device is unbuffered, the \(\lambda terminal \) must be defined for the \(\lambda terminal \) definition\(\rangle\).

# **TERMINAL**

Terminal Carriage Character Statement

# TERMINAL CARRIAGE CHARACTER STATEMENT

**Syntax** 

# Example

CARRIAGE = 4"0D".

# **Semantics**

This statement is implemented for program documentation purposes only. This statement provides a means of documenting the carriage return character of a terminal type. The documentation of this character is optional in a \(\lambda terminal definition\rangle\).

Terminal Clear Character Statement

# TERMINAL CLEAR CHARACTER STATEMENT

**Syntax** 

# Example

**CLEAR = 4"11".** 

# **Semantics**

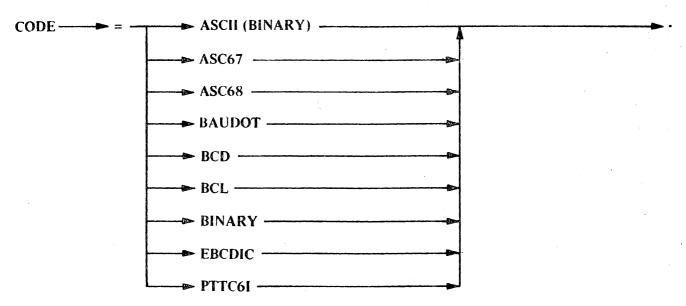
This statement is implemented for program documentation purposes only. It provides a means to document the clear character of a terminal type. The documentation of this character is optional in a \(\lambda terminal definition\rangle\).

#### **TERMINAL**

**Terminal Code Statement** 

#### TERMINAL CODE STATEMENT

Syntax



#### **Semantics**

The \(\lambda terminal code statement\rangle\) specifies the character code translation required for the DCP to communicate with the terminal type. The internal code of the DCP is EBCDIC, and the DCP translates from EBCDIC to the code specified for transmissions, and from the code specified to EBCDIC for receptions.

**BINARY** and **EBCDIC** specify that translation is not required.

ASC67 and ASC68 specify the standard software translation tables for the ASCII character code.

ASCII (BINARY) allows a *control definition* or *request definition* to switch back and forth between ASCII code translation and no translation. The *code statement* in a *request definition* or *control definition* or *control definition* to enter a "transparent" mode in Binary Synchronous communications procedures.

BAUDOT, BCD, BCL, and PTTC61 specifications are all indicative of the translation they invoke. For example, BAUDOT invokes the Baudot character code set, PTTC61 invokes PTTC/6, etc.

#### **Pragmatics**

For special applications a programmer can define and invoke non-standard character codes by:

- a. defining a translation table in a \(\langle translatetable definition \rangle \);
- b. specifying **BINARY** or **EBCDIC** in the \(\langle terminal code statement \rangle\); and
- c. invoking the translation in a \(\chicontrol\) definition\(\chi\) or \(\chicontrol\) definition\(\chi\) by means of the appropriate option of the \(\lambda assignment \) statement\(\chi\).

Refer to the \(\langle translatetable definition\rangle\) in this chapter for more information.

#### TERMINAL CONTROL STATEMENT

**Syntax** 



# Examples

CONTROL = CONTENTION.
CONTROL = PRIMARYCONTROL,
AUXILIARYCONTROL.

#### Semantics

The \(\lambda terminal control statement\rangle\) specifies the \(\lambda control definition\rangle(s)\) responsible for allocation of the logical line(s) to which a terminal type is associated. This attribute must be defined for all \(\lambda terminal definitions\rangle\).

Terminal types that do not utilize full duplex, reverse channel, or voice response features require that only one  $\langle control\ identifier \rangle$  be named.

Terminal types that utilize full duplex, reverse channel, or voice response features (i.e., DUPLEX = TRUE) may optionally specify a second  $\langle control\ identifier \rangle$ . The first  $\langle control\ identifier \rangle$  names the  $\langle control\ definition \rangle$  for the primary line, and the second  $\langle control\ identifier \rangle$  names the  $\langle control\ definition \rangle$  for the auxiliary line. If only one  $\langle control\ identifier \rangle$  is specified, it is assumed to be the  $\langle control\ definition \rangle$  for the primary line, and the default equivalent of an  $\langle idle\ statement \rangle$  is used for auxiliary line control.

# **TERMINAL**

**Terminal Default Statement** 

#### TERMINAL DEFAULT STATEMENT

Syntax

DEFAULT	$\cdot \langle$ default terminal identifie $r  angle$	
---------	---	--

Example

DEFAULT = TTYDFLT.

#### **Semantics**

The \(\lambda terminal default statement\rangle\) allows the programmer to specify the \(\lambda default terminal identifier\rangle\) of a set of default terminal attributes previously defined to be used for a \(\lambda terminal definition\rangle\) whose description is incomplete. It is advantageous to group common attributes under a Default \(\lambda terminal definition\rangle\) and list the remaining attributes under each individual \(\lambda terminal definition\rangle\). The compiler will then refer to the Default \(\lambda terminal definition\rangle\) to complete the \(\lambda terminal definition\rangle\); however, a \(\lambda terminal definition\rangle\) must define all required attributes if a \(\lambda terminal default statement\rangle\) does not appear.

The  $\langle terminal \ default \ statement \rangle$  can appear in a  $\langle terminal \ definition \rangle$  or a Default  $\langle terminal \ definition \rangle$ .

# **Supplementary Example**

The following example illustrates how  $\langle terminal \ default \ statement \rangle s$  may be "nested" to combine the attributes of one or more Default  $\langle terminal \ definition \rangle s$ .

The effect of referencing GENERALDEFAULT within the Default \(\lambda\) terminal definition \(\rangle\) TTYDEFAULT is that the attributes associated with TTYDEFAULT are equivalent to all attributes as defined by GENERALDEFAULT plus the attributes explicitly defined in TTYDEFAULT.

If a \(\text{terminal definition}\) or Default \(\text{terminal definition}\), references a Default \(\text{terminal definition}\), the compiler does not compare the two definitions for contradictory statements. If contradictory statements exist within the two definitions, the last value defined for the attribute takes precedence. In the example, TTY2 defines the value of the PAGE attribute as 66, and the Default \(\text{terminal definition}\) that TTY2 references defines the value of the PAGE attribute as 0. The compiler uses 66 as the value of the PAGE attribute for TTY2.

#### TERMINAL

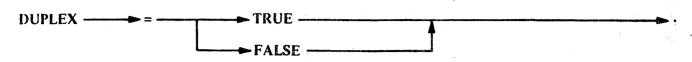
# Terminal Default Statement - Continued

```
TERMINAL DEFAULT GENERALDEFAULT:
    TURNAROUND = 0.
    ICTDELAY
                   = 0.
                                                  Default (terminal definition)
    TRANSMISSION = 0.
    ADDRESS
                   = 0.
    PAGE
                   = 0.
    BUFFER
                   = 0.
TERMINAL DEFAULT TTYDEFAULT:
                                                  (terminal default statement) references
    DEFAULT
                   = GENERALDEFAULT.
                                                 above Default (terminal definition)s.
    BLOCK
                   = FALSE.
                   = FALSE.
    SCREEN
                   = NULL.
    PARITY
    SYNCS
                   = FALSE.
                                                  Default (terminal definition)
                   = 3 SEC.
    TIMEOUT
                   = 72.
    MAXINPUT
    WIDTH
                   = 72.
    ADAPTER
                   = 4.
    CODE
                   = ASC67.
TERMINAL TTY1:
                                                 \(\langle\) terminal default statement \(\rangle\) references
    DEFAULT
                   = TTYDEFAULT.
                                                 above Default (terminal definition)s.
    DUPLEX
                   = FALSE.
    WRU
                   = ENQ.
    END
                   = ETX (DYNAMIC).
    BACKSPACE
                   = BS (DYNAMIC).
    CONTROL
                   = CONTEND.
    REQUEST
                   = WRITETTY: TRANSMIT, READTTY: RECEIVE.
TERMINAL TTY2:
                                                  (terminal default statement) references
    DEFAULT
                   = TTYDEFAULT.
                                                 above Default (terminal definition)s.
    DUPLEX
                   = FALSE.
    WRU
                   = 4"98".
                   = 4"97"
    BACKSPACE
    CONTROL
                   = SPECIALCNTRL.
                   = READER: RECEIVE, WRITER: TRANSMIT.
    REQUEST [1]
    REQUEST [2]
                   = READPPT: RECEIVE, WRITEPPT: TRANSMIT.
    PAGE
```

**Terminal Duplex Statement** 

TERMINAL DUPLEX STATEMENT

**Syntax** 



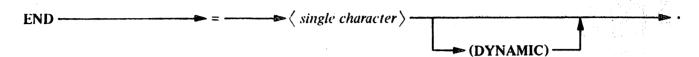
# **Semantics**

The \(\lambda terminal duplex statement\rangle\) defines whether or not (TRUE or FALSE, respectively) the terminal type utilizes full duplex, reverse channel, or voice response features. If DUPLEX = TRUE, then the \(\lambda line \) definition\rangle for any line that has this terminal type assigned must contain the \(\lambda line type statement \rangle\) constructs that specify full duplex. This attribute must be defined for each \(\lambda terminal definition \rangle\).

# Terminal End Character Statement

# TERMINAL END CHARACTER STATEMENT

Syntax



# **Examples**

END = 4"0D". END = "&" (DYNAMIC).

#### **Semantics**

The \(\lambda terminal end character statement\rangle\) defines the "end" character of the terminal type (i.e., the character that the terminal type would transmit to indicate an end-of-text). If defined, the "end" character can be recognized by the DCP when RECEIVEd (in a \(\lambda receive statement\rangle\)), and any action to be taken can be specified by the programmer (using the END syntax).

(DYNAMIC) indicates that the Message Control System of a station referencing the \(\lambda terminal definition\rangle\) is allowed to change the character for the station by means of a SET CHARACTERS (TYPE=39) DCWRITE.

# **TERMINAL**

**Terminal Home Character Statement** 

# TERMINAL HOME CHARACTER STATEMENT

Syntax

 $HOME \longrightarrow = \longrightarrow \langle single \ character \rangle \longrightarrow$ 

Example

HOME = 4"0C".

# **Semantics**

This statement is implemented for program documentation purposes only. It provides a means of documenting the home character of the terminal type. The documentation of this character in a \( \lambda terminal \) definition\( \rangle \) is optional.

Terminal Illegal Character Statement

# TERMINAL ILLEGAL CHARACTER STATEMENT

**Syntax** 

Example

ILLEGALCHR = 4"FF".

# **Semantics**

The  $\langle terminal \ illegal \ character \ statement \rangle$  is implemented for documentation purposes only. The documentation of this character is not required in a  $\langle terminal \ definition \rangle$ .

#### **TERMINAL**

Terminal Inhibitsync Statement

#### TERMINAL INHIBITSYNC STATEMENT

Syntax



#### **Semantics**

The \(\lambda terminal inhibits ync statement\rangle\) affects only terminal types that specify any of the \(\lambda communication type number\rangles 17\) through 27 in its \(\lambda terminal adapter statement\rangle\). This statement has no affect upon, and need not be defined for, terminal types that do not specify any of those \(\lambda communication type number\rangles\).

If INHIBITSYNC = FALSE, then the following occurs during a synchronous transmission. The transmission begins with the transmission of four sync characters by the adapter cluster. As the fourth sync character is being transmitted, the first character of the message is requested from the DCP. The DCP should respond to this request by supplying the first character of the transmission. As each supplied character is transmitted, the adapter cluster requests another character. If the DCP is unable to respond in time to the request, the adapter cluster transmits a sync character; this process is called "sync filling." Sync filling is repeated as necessary until the DCP responds with another character or the DCP directs the adapter cluster to "finish transmit" for the line.

When INHIBITSYNC = FALSE during a synchronous reception, the following occurs. At the beginning of the reception, bit patterns from the line are examined by the adapter cluster and the bits discarded until a sync character is recognized. The recognition of a sync character establishes that the next bit to be received by the adapter cluster is the first bit of the next character. The sync character is discarded, instead of being made available to the DCP. All characters in the transmission that are not sync characters are made available to the DCP. The DCP may then fetch these characters. Any sync characters received in the transmission are discarded.

If INHIBITSYNC = TRUE, then the following occurs during a synchronous transmission. All actions occur that would occur if INHIBITSYNC = FALSE. In addition, if a sync fill is required, a "sync fill interrupt" occurs so that the DCP can determine when one or more undesired sync characters have been inserted into the transmission. System software responds to the interrupt by executing a TERMINATE ERROR. The controlling MCS is notified of all such situations so that corrective action (MAKE LINE READY (TYPE = 96) DCWRITE, for example) can be taken.

When INHIBITSYNC = TRUE during a synchronous reception, the following occurs. At the beginning of the reception, bit patterns from the line are examined by the adapter cluster and the bits discarded until a sync character is recognized. The recognized sync character is discarded, as is the next character if it is also a sync character. Thereafter, all subsequent characters (sync characters or otherwise) are made available to the DCP as data.

The reserved word SYNCS is a synonym for INHIBITSYNC.

Terminal Inter-Character Delay Statement

### TERMINAL INTER-CHARACTER DELAY STATEMENT

**Syntax** 

**Examples** 

ICTDELAY = 0. ICTDELAY = 200 MILLI.

#### **Semantics**

The \(\lambda terminal inter-character delay statement\)\) provides the user a means to insert a timed delay between each character transmitted to the terminal type. The \(\lambda delay time\rangle\) specified defines the interval of \(\lambda time\rangle\) between the transmission of the start of one character to the start of the next character. If the time specified is less than the time required to transmit a character, this statement has no effect. This attribute must be defined for all \(\lambda terminal definition\rangle s\).

# Supplementary Example

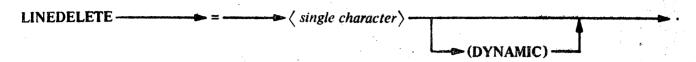
A Model 33 TELETYPE can receive characters at a maximum rate of one character every 100 milliseconds. If, for some reason, the programmer needs to insert a 100-millisecond delay between each character transmitted to the terminal, this can be done by specifying:

ICTDELAY = 200 MILLI.

Terminal Linedelete Character Statement

# TERMINAL LINEDELETE CHARACTER STATEMENT

**Syntax** 



# **Examples**

LINEDELETE = 4"07". LINEDELETE = 4"A0".

# **Semantics**

The (terminal linedelete character statement) defines the linedelete character of the terminal type. If defined, the linedelete character can be recognized by the DCP when RECEIVEd (in a (receive statement)), and any action to be taken can be specified by the programmer (using the LINEDELETE syntax).

(DYNAMIC) indicates that the Message Control System of a station referencing the \(\lambda\)terminal definition\(\rangle\) is allowed to change the character for the station by means of a SET CHARACTERS (TYPE=39) DCWRITE.

Terminal Linefeed Character Statement

# TERMINAL LINEFEED CHARACTER STATEMENT

**Syntax** 

LINEFEED - < single character >.

Example

**LINEFEED = 4"25".** 

# **Semantics**

This statement is provided for program documentation only. It documents the linefeed character of the terminal type. The documentation of this character in a \( \text{terminal definition} \) is optional.

**Terminal Maxinput Statement** 

TERMINAL MAXINPUT STATE	FA	/FNT
-------------------------	----	------

**Syntax** 

Example

MAXINPUT = 72.

### **Semantics**

The \(\lambda terminal maximput statement\rangle\) applies to unbuffered terminals and defines the maximum size text, in characters, that a terminal is allowed to transmit in one message. This attribute must be defined in all \(\lambda terminal definition\rangle\)s in which the \(\lambda terminal buffer size statement\rangle\) is not defined or is defined as \(\text{BUFFER} = \text{NULL}\). This statement applies only to unbuffered devices; it is meaningless to define maximput if the \(\lambda terminal buffer size statement\rangle\) is defined as non-\text{NULL}\.

Terminal Page Statement

# TERMINAL PAGE STATEMENT

**Syntax** 

 $PAGE \longrightarrow = \longrightarrow \langle integer \rangle$ 

# **Examples**

**PAGE = 0. PAGE = 12.** 

#### **Semantics**

The \(\lambda terminal\) page statement\(\rangle\) defines the maximum number of output lines per page as restricted by the hardware of the terminal. There are, for example, devices that can only print/display a defined number of lines before some type of carriage/cursor control information must be supplied. If the terminal type being defined has no such restrictions, then

$$PAGE = 0.$$

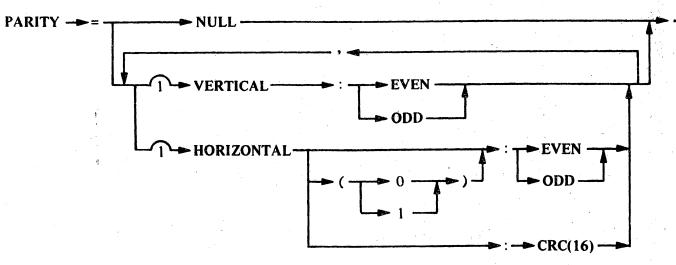
should be specified, thus indicating that pagination is arbitrary. This attribute must be defined for all (terminal definition)s.

#### **TERMINAL**

**Terminal Parity Statement** 

#### TERMINAL PARITY STATEMENT

Syntax



# Examples

PARITY = NULL.

PARITY = VERTICAL:ODD.

PARITY = HORIZONTAL:CRC(16).

PARITY = VERTICAL:ODD, HORIZONTAL(0):EVEN.

# **Semantics**

The \(\lambda terminal parity statement\rangle\) defines the type of parity checking and generation to be performed by the DCP when communicating with the terminal type. If the form:

# PARITY=NULL.

is used, parity is not checked or generated.

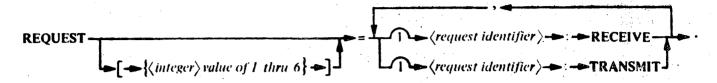
The VERTICAL option refers to the vertical parity bit of a character, and can be defined as ODD or EVEN.

The HORIZONTAL option specifies the type of horizontal parity. If horizontal parity is a Block Check Character, then ODD or EVEN must be specified. If horizontal parity is a Cyclic Redundancy Check, then CRC(16) must be specified.

The 0 or 1 option defines the function of the vertical parity bit of the Block Check Character. If this bit is a parity bit for the Block Check Character, then this option must be omitted or defined as 0 (zero). If undefined, the option is assumed to be 0 (zero). If the bit is to be considered as a horizontal parity bit of all high-order bits in the message, then this option must be defined as 1.

#### **TERMINAL REQUEST STATEMENT**

Syntax



#### **Examples**

REQUEST = READTTY:RECEIVE.
REQUEST = WRITETTY:TRANSMIT, READTTY:RECEIVE.
REQUEST[2] = TTYTAPEIN:RECEIVE, TTYTAPEOUT:TRANSMIT.

#### **Semantics**

The \(\lambda terminal request statement\rangle\) specifies a \(\lambda request identifier\rangle\), or a pair of \(\lambda request identifier\rangle\)s, that designates the \(\lambda request definition\rangle\) to handle input from (the RECEIVE option) and/or output to (the TRANSMIT option) the terminal type. The \(\lambda request definition\rangle\) that handles input is commonly referred to as the Receive Request, and the \(\lambda request definition\rangle\) that handles output is commonly referred to as the Transmit Request. This statement must appear in each \(\lambda terminal definition\rangle\), and cannot appear in a Default \(\lambda terminal definition\rangle\).

The  $\{\langle integer \rangle \}$  value of 1 through 6 allows the specification of up to six pairs of Transmit and Receive Requests for the same device. Normally, these Request pairs differ for some application-dependent reasons. Only one pair of  $\langle request \ definition \rangle s$  can be the controlling  $\langle request \ definition \rangle s$  at any instant of time. The  $\langle request \ definition \rangle s$  in control of the terminal type immediately after DCP initialization has an  $\{\langle integer \rangle \ value \ of \ 1 \ through \ 6 \}$  of 1; they retain control until the Message Control System (MCS) of a station associated with the terminal type executes a SET APPLICATION NUMBER (TYPE = 38) DCWRITE.

## Definitions TERMINAL

Terminal Screen Statement

TERMINAL	SCREEN STATEMEN			
Syntax				
SCREEN —		TRUE —	1	
		FALSE —		

#### **Semantics**

The \(\lambda terminal screen statement\rangle\) defines whether or not (TRUE or FALSE, respectively) the terminal type is a screen (i.e., CRT) device. This attribute must be defined in each \(\lambda terminal definition\rangle\).

# Definitions TERMINAL Terminal Timeout Statement

TERMINAL	TIMEOUT STATE	MENT			
Syntax					
TIMEOUT -	=		timeout time $ angle$	-	•
Example					
TIMEOU	UT = 3 SEC.				

#### **Semantics**

The \(\lambda terminal timeout statement\rangle\) defines the interval of \(\lambda time\rangle\) that the adapter cluster should wait from the receipt of one character to the start of the next (in a \(\lambda receive statement\rangle\)) before assuming that the terminal has "timed out." The action taken upon a timeout condition can be specified in a \(\lambda receive statement\rangle\) by means of the TIMEOUT syntax.

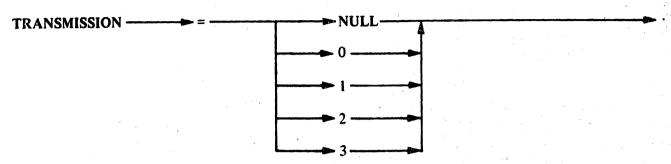
#### **Definitions**

#### **TERMINAL**

Terminal Transmission Number Length Statement

#### TERMINAL TRANSMISSION NUMBER LENGTH STATEMENT

**Syntax** 



#### **Semantics**

The \(\lambda terminal transmission number length statement\)\) defines the number of characters that the terminal transmits and receives as the message transmission number. The 0 and NULL options are semantically equivalent and specify that no transmission number is used. A non-NULL transmission number length must be specified if a \(\lambda control definition\rangle\) or \(\lambda request definition\rangle\) that references the item TRAN is defined for the terminal type. This statement may be omitted from a \(\lambda terminal definition\rangle\) if the terminal does not transmit or receive transmission numbers.

## Definitions TERMINAL

#### Terminal Turnaround Statement

#### TERMINAL TURNAROUND STATEMENT

**Syntax** 

**Examples** 

TURNAROUND = 0. TURNAROUND = 200 MILLI.

#### **Semantics**

The \(\lambda terminal turnaround statement\rangle\) defines the time required for the terminal to shift from transmitting data to receiving data. The \(\lambda time \rangle\) defined is a parameter of a compiler algorithm for calculating the initiate transmit delay. Refer to the semantics of the \(\lambda control definition \rangle\) or \(\lambda request definition \rangle\) \(\lambda initiate statement \rangle\) in this chapter for more information. This attribute must be defined for each \(\lambda terminal definition \rangle\).

## Definitions TERMINAL

Terminal Width Statement

						SI				

**Syntax** 

 $\mathsf{WIDTH} \longrightarrow = \longrightarrow \langle integer \rangle$ 

Example

WIDTH = 80.

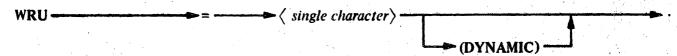
#### **Semantics**

The \(\lambda terminal \) width statement\(\rangle\) defines the width, in characters, of a display line of output on the terminal type. The \(\lambda integer\rangle\) must be greater than 0 and less than 256; additionally, the value of the \(\lambda integer\rangle\) must be less than or equal to the size defined in the \(\lambda terminal \) buffer size statement\(\rangle\), if present. It is not required that the \(\lambda terminal \) width statement\(\rangle\) appear in a \(\lambda terminal \) definition\(\rangle\). If the \(\lambda terminal \) width statement\(\rangle\) is not defined in the \(\lambda terminal \) definition\(\rangle\), then the buffer size value is substituted for this value, if present; otherwise, the value of MAXINPUT is substituted by default.

Terminal WRU Character Statement

#### TERMINAL WRU CHARACTER STATEMENT

Syntax



#### **Examples**

#### **Semantics**

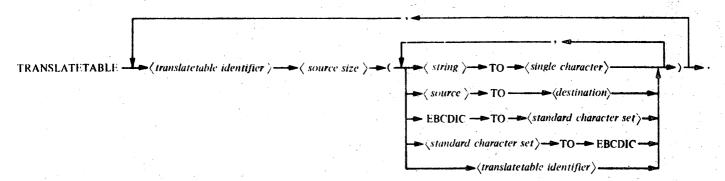
The \(\langle terminal WRU \) character statement\(\rangle\) defines the WRU character for the terminal type (i.e., the character the terminal type would transmit to request a response from the DCP). If defined, the WRU character can be recognized by the DCP when RECEIVEd (in a \(\langle\) receive statement\(\rangle\)), and any action to be taken can be specified by the programmer (using the WRU syntax).

(DYNAMIC) indicates that the Message Control System of a station referencing the \(\lambda\) terminal definition\(\rangle\) is allowed to change the character for the station by means of the SET CHARACTERS (TYPE=39) DCWRITE.

#### TRANSLATETABLE

#### TRANSLATETABLE DEFINITION

#### **Syntax**



#### **Examples**

TRANSLATETABLE	ATABLE	8("STRING" TO "X").
TRANSLATETABLE	<b>BTABLE</b>	8(4"000102" TO 4"AABBCC").
TRANSLATETABLE	<b>CTABLE</b>	7(4"02820B" TO 4"AACCDD").
TRANSLATETABLE	DTABLE	7(4"00" TO 4"AA"),
	ETABLE	8(DTABLE, 4"01" TO 4"BB").
TRANSLATETABLE	<b>FTABLE</b>	8(4"00" TO 4"AA"),
	<b>GTABLE</b>	8(EBCDIC TO BCL,FTABLE,
		4"01" TO 4"BB").
TRANSLATETABLE	TRANID	7("1" TO 4"01",
· ·		"2" TO 4"02".
	· .	"34" TO 4"0304",
		4"F5F6" TO 4"0607")

#### **Semantics**

The  $\langle translatetable definition \rangle$  allows the definition of tables that may be used in  $\langle control definition \rangle$  s or  $\langle request definition \rangle$  s to translate characters of one character set to those of another character set.

Translation tables need to be defined in an NDL program only if non-standard character sets must be dealt with in the Data Communications System. Terminals that transmit and receive a standard character set do not require a translation table definition; instead, the character set is merely named in the \(\lambda terminal\) code statement\(\rangle\) of the \(\lambda terminal\) definition\(\rangle\). The character sets that do not require a \(\lambda translatetable\) definition\(\rangle\) are ASCII, BAUDOT, BCD, BCL, EBCDIC, and PTTC/6.

The \(\langle translatetable identifier \rangle \) that follows the key word TRANSLATETABLE names the translation table, and must be in the syntactic form of an \(\langle identifier \rangle \).

(source size) defines the character size, in bits, of characters to be translated. (source size) must be an /integer greater than 0 and less than 9.

#### TRANSLATION TABLE STRUCTURE

Each element of the translation table consists of eight bits. If N represents the (source size), then the size of the table is 2 raised to the Nth power. The elements of the table are selected by an index that ranges from 0 through 2 to the Nth power minus 1.

At execution time, translation is done in the following manner. The binary weight of the low-order N bits of the character to be translated is used as an index into the specified translation table. The element of the table thus indexed is the translated result.

Continued

#### INSERTING DATA INTO THE TRANSLATION TABLE

Every translation table has a default base in which each element in the table is 0 (all bits off). Data can be placed into the translation table by various specifications within the parentheses. If more than one specification appears for a given translation table, each succeeding specification overrides, within its scope, previous specifications.

\(\langle\) to \(\langle\) single character\(\rangle\)

This form inserts data into the translation table in the following manner. Each eight-bit character in the  $\langle string \rangle$  is examined from left to right. If a character in the  $\langle string \rangle$  is numerically greater than the size of the table, no entry is placed in the translation table; otherwise, the  $\langle single\ character \rangle$  is stored in the element of the table whose index is the binary weight of the N low-order bits of the  $\langle string \rangle$  character (where N is the  $\langle source\ size \rangle$  specified).

(source) TO (destination)

(source) and (destination) must be (string)s of equal length. This form of specification inserts data into the translation table in the following manner. Translation is based upon corresponding characters in (source) and (destination), starting from left and proceeding to right. The first character of (source) corresponds to the first character of (destination), the second character of (source) corresponds to the second character of (destination), etc. If a character in (source) is numerically greater than the size of the table, then no entry is placed in the translation table; otherwise, the corresponding character in (destination) is stored in the element of the table whose index is the binary weight of the N low-order bits of the corresponding character in (source).

(standard character set) TO EBCDIC and EBCDIC TO (standard character set)

This form specifies a standard system software translation table from the NDL compiler that is to be copied into the translation table. The (standard character set)s that may be specified are EBCDIC, ASC67, and BCL. These forms provide a way of obtaining a legitimate base upon which additional specifications can be made.

(translatetable identifier)

This form of specification indicates that the contents of a previously defined translation table is to be copied into the translation table. The \(\lambda translate table \) identifier \(\rangle\) must be the \(\lambda identifier \rangle\) of a previously defined translation table. This form provides a means of obtaining a legitimate base upon which additional specifications can be made.

#### **Pragmatics**

Those tables, and only those tables, that are used by a DCP reside in the local memory of that DCP (unless a DCP does not have local memory, in which case they reside in main system memory). Memory for translation tables is allocated in blocks of 256 words, regardless of the space required for those tables. Tables are densely packed and all elements are used before another block of 256 words is allocated. Unless consideration is given to the translation requirements of devices in the data communications system while in the planning and programming stages, translation tables can be very costly in terms of local memory. Although it is beyond the scope of this manual to describe the planning of a data communications system, this fact should not escape the NDL programmer.

#### **Definitions**

#### TRANSLATETABLE

Continued

#### Supplementary Examples

#### Example 1.

## TRANSLATETABLE ATABLE 8("STRING" TO "X").

Character to be Trans	Character to be Translated									
"S"			"X"							
"T"			"X"							
"R"			"X"							
··I"			"X"							
"N"			"X"							
"G"			"X"							

ATABLE is a translation table containing 256 elements. The (source size), 8 in this example, determines the table size. All characters from (source) are translated to the (single character).

#### Example 2

#### TRANSLATETABLE BTABLE 8(4"000102" TO 4"AABBCC").

Charact	ter to be T	ranslate	<u>d</u>		Result
	4"00" 4"01" 4"02" 4"03"				4"AA" 4"BB" 4"CC" 4"00"

BTABLE contains 256 elements. Characters from (source), 4"000102", are translated to the corresponding characters in the (destination), 4"AABBCC". The character 4"03" is translated to 4"00" because there is no specification in (source) for 4"03".

#### Example 3

#### TRANSLATETABLE CTABLE 7(4"02820B" TO 4"AACCDD").

Character to be Translated	Result
4"01"	4"00"
4"02"	4"AA"
4"82"	4"AA"
4"0B"	4"DD"

In this example, the translation table CTABLE contains 128 elements. The character 4"01" is translated to a 4"00" character, because 4"01" is unspecified in the \( \source \). The character 4"82" is translated to the character 4"AA" because only the low-order seven bits of 4"82" are used to index the translation table.

#### Example 4

## TRANSLATETABLE DTABLE 8(4"00" TO 4"AA"). ETABLE 8(DTABLE, 4"01" TO 4"BB").

The above \(\langle\tanslate\t

Continued

#### Example 5

TRANSLATETABLE FTABLE 8(4"00" TO 4"AA").
GTABLE 8(EBCDIC TO BCL, FTABLE, 4"01" TO 4"BB").

GTABLE is defined to contain 256 elements, and specifies the standard EBCDIC-to-BCL translation table upon which subsequent specifications modify. FTABLE also contains 256 elements and appears as a specification in GTABLE. Since each succeeding specification overrides within its scope any previous specification, FTABLE in effect overlays all elements. The result is the same as if only the following had appeared:

TRANSLATETABLE FTABLE 8(4"00" TO 4"AA"), GTABLE 8(FTABLE, 4"01" TO 4"BB").

The above example points out that any table appearing as a specification indicates all elements of that table, not just those elements explicitly defined. The example is not intended to illustrate an acceptable programming practice.

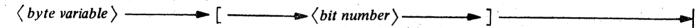
#### 6. VARIABLES

#### **GENERAL**

The NDL compiler does not allow a programmer to declare and use program variables, as do other language compilers such as ALGOL, PL/I, and COBOL. Instead, the NDL programmer can use only predefined program variables.

The \( \begin{align\*} bit variable \rangle s \) and \( \begin{align\*} byte variable \rangle s \) are the two types of variables the programmer can use. The \( \begin{align\*} bit variable \rangle s \) are one-bit variables that can only assume logical values (i.e., TRUE or FALSE). The \( \begin{align\*} byte variable \rangle s \) are all eight-bit variables, and can assume integer values from 0 through 255, except for the IR variable, which is a 10-bit variable. The IR variable is included as a \( \begin{align\*} byte variable \rangle \) as a matter of convenience.

Individual bits of a \( \langle \text{byte variable} \rangle \) can be referenced and used like a \( \langle \text{bit variable} \rangle \), if referenced in the form illustrated below.



where \( \lambda \text{it number} \rangle \text{ is an \( \lambda \text{integer} \rangle \text{ not greater than the number of bits contained in the variable minus 1.} \)

For example, bit 5 of IR is referenced as IR[5].

#### **FUNCTION OF VARIABLES**

Functionally, variables fall into one of three general categories:

- a. Variables that are available to the programmer for general information storage.
- b. Variables that can be used for system/station communication.
- c. Variables that contain control information.

General information variables can be used within their scope by the programmer for data storage, calculations, etc. Additionally, some variables in this category could (by convention) be used as communication paths between (request definition)s executing on a given line. (The use of a given variable for this application is restricted by the scope of that variable.)

Variables whose intended function is communication to and from the main system and stations are generally contained in the message header of a message sent to the main system from a station, or sent to the station from the main system. Messages from the main system to a station are originated either by the MCS or by an application program (via the I/O Intrinsics).

The format of message variables within a message header is described in detail in the <u>B 6700/B 7700</u>

<u>DCALGOL Reference Manual</u>. Generally, message variables are contained in five fields of the message header:

- a. Message Toggles (word [1].[39:8])
- b. Message Tallys (word [3].[23:24])

#### Continued

- c. Message Error Flags (word [1], [23:24])
- d. Variant "Carriage Control" (word [0] . [39:16])
- e. Message Retry Count (word [2], [47:8])

Message Toggles and Message Tallys provide storage area in the header for some of the station general information variables. The meaning of values stored in these fields must be established by mutual convention between the MCS writer and the NDL programmer.

Message Error Flags are used for the station to communicate to an MCS that some exceptional event has occurred in a \( \text{request definition} \) or \( \text{control definition} \). These variables reference bits in the message header of "result" messages returned to the MCS as a result of execution of a \( \text{terminate statement} \).

Carriage Control is valid for Transmit Requests, and provides information regarding the kind of carriage control to be performed by a Transmit Request. These variables reference bits or bytes in the message header of WRITE (TYPE=33) DCWRITE messages.

The Message Retry Count is described under RETRY in this chapter.

Variables whose function is to contain control information are used by both the DCP operating system and the programmer. Generally, these variables provide information to control the logic paths of *(control definition)s*, *(request definition)s*, and the DCP operating system.

#### **SCOPE OF VARIABLES**

The scope of the variables in NDL is described as being:

- a. Station-oriented.
- b. Line-oriented.
- c. Global.

Station-oriented variables exist for each station in the network. TALLY [0] is an example of a station-oriented (byte variable); thus, each station has its own TALLY [0]. The variables of a given station are visible to a line only while STATION is set to that station's "station index."

Line-oriented variables exist for each line on a DCP. The variables of a given line are visible to every station assigned to that line. MAXSTATIONS is an example of a line-oriented variable. Each line on a DCP has its own MAXSTATIONS, and every station assigned to a given line can access the MAXSTATIONS variable of that line.

A global variable is a variable that is visible to all stations on a DCP.

#### **DESCRIPTION OF VARIABLES**

The remainder of this chapter contains descriptions of each  $\langle bit \ variable \rangle$  and  $\langle byte \ variable \rangle$ . The variables (listed in table 6-1) are described in alphabetical order. The name of the variable precedes a summary of the variable characteristics, followed by a detailed description of the variable.

The summary of the variable characteristics includes the places in the source program that the variable can be interrogated or altered, and the size, in bits, of the variable. In the summary, the word "Interrogate" indicates that the programmer can interrogate the variable. The word "Alter" indicates that the programmer can use the \( \delta \text{it variable} \rangle / \left \delta yte variable \rangle \) as an \( \left \left \text{assignable bit variable} \rangle / \left \left \text{assignable byte variable} \rangle \). The corresponding letters "C", "T", and "R" in the summary refer to \( \left \left \text{control definition} \right \), Transmit Request, and Receive Request, respectively. The last item to appear in the summary is the size, in bits, of the variable. If no size is defined, then the size of the variable is one bit.

For example, the summary:

**EXAMPLE1** Interrogate, CTR, 8

can be expanded as follows:

**EXAMPLE1** is an 8-bit variable. It can be interrogated in a *control definition*, Transmit Request, or Receive Request.

The summary:

EXAMPLE2
Interrogate/Alter, CTR/TR

can be expanded as follows:

EXAMPLE2 is a \( \begin{align\*} bit variable \right) . It can be interrogated in a \( \control \) definition \( \right) , \)
Transmit Request, or Receive Request.
Additionally, EXAMPLE2 can be altered (i.e., appear as an \( \langle assignable \) bit variable \( \right) in a Transmit Request or Receive Request, but not in a \( \langle control \) definition \( \right).

Table 6-1 contains the summaries of each variable for quick reference.

Table 6-1. Table of Variables

NAME	SIZE (in bits)	INTERROGATE	ALTER
1 14 2172 22			
ADDERR		CTR	CTR
Ai	8	CTR	CTR
AUX (LINE (BUSY))		CTR	CTR
AUX (LINE (QUEUED))		CTR	CTR
AUX (LINE (TALLY [{0 or 1}]	8	CTR	CTR
AUX (LINE (TOG [{0 or 1}])		CTR	CTR
BCC	8	CTR	CTR
BCCERR		TR	TR
BLOCK		Т	
BLOCKED		Т	
BREAK [RECEIVE]		CTR	CTR
BREAK [TRANSMIT]		CTR	CTR
BUFOVFL		CTR	CTR
CARRIAGE		T	
CHARACTER	8	CTR	CTR
CONTROLFLAG		TR	TR
CRC		CTR	CTR
CRC [{0 or 1}]	8	CTR	CTR
CRCERR		TR	TR
DISCONNECT		TR	TR
ENDOFBUFFER		TR	TR
FORMATERR		TR	TR
INHIBITSYNC		CTR	CTR
IR	10	CTR	
LINE (BUSY)		CTR	CTR
LINE (QUEUED)		CTR	CTR
LINE (TALLY [{0 or 1}]	8	CTR	CTR
LINE (TOG [{0 or 1}]		CTR	CTR
LINEFEED		T	
LOSSOFCARRIER		CTR	CTR
MAXSTATIONS	Ω	CTR	
NAKFLAG	0	TR	TR

Table 6-1. Table of Variables (Cont)

	SIZE			
NAME	(in bits)	INTERROGATE	ALTER	
NAKONSELECT		TR	TR	
NOSPACE		CTR		
PAGE		T		
PAPERMOTION		T		
PARITY		CTR	CTR	
RETRY	8	CTR	CTR	
SEQERR		TR	TR	
SEQUENCE		CTR	CTR	
SKIP		T	-	
SKIPCONTROL	8	T		
SPACE		T		
STATION	8	C	С	
STATION (ENABLED)		CTR		
STATION (FREQUENCY)	8	CTR		
STATION (QUEUED)		CTR	4	
STATION (READY)		CTR		
STATION (TALLY)	8	CTR	CTR	
STATION (VALID)		CTR		
STOPBIT		CTR	CTR	
SYNCS		CTR	CTR	
TAB		T		
TALLY [ \(\langle tally number \rangle \)]	8	CTR	CTR	
TIMEOUT		CTR	CTR	
TOG [\langle toggle number\rangle]		CTR	CTR	
TOGS	8	CTR	CTR	
TRANERR		TR	TR	
WRUFLAG		TR	TR	
			4	

Continued

#### **ADDERR**

Interrogate/Alter, CTR/CTR

ADDERR references bit 8 in the Error Flag Field of a message header, and normally indicates that an address character error has occurred while executing a *(receive statement)*. Refer to the ADDRESS option of the *(receive statement)*.

#### AI

Interrogate/Alter, CTR/CTR, 8

This variable addresses a volatile register and should not be used for data storage. Its main purpose is to allow access to the untranslated byte just received rather than to the translated byte in **CHARACTER**, particularly when executing the (sum statement).

#### AUX(LINE(BUSY))

Interrogate/Alter, CTR/CTR

AUX(LINE(BUSY)) is used to allow or inhibit the interruption of the execution of a *(control definition)* or *(request definition)* on the auxiliary line of a full duplex line pair. If this bit is TRUE, it indicates to the DCP operating system that the line is engaged in functions that must not be interrupted. If FALSE, it indicates to the DCP operating system that the line can be interrupted to initiate another function. AUX(LINE(BUSY)) is line-oriented, but may be altered only by the auxiliary line. Both the auxiliary and primary line may interrogate this bit.

A (control definition) or (request definition) will be interrupted when AUX(LINE(BUSY)) is FALSE if the primary line executes a (fork statement). (Note that an interruption causes control to leave a (control definition) or (request definition), and that control is not returned to the point where the interruption occurred.) AUX(LINE(BUSY)) is set TRUE by system software when:

- a. The primary line executes a \( \) fork statement \( \) and AUX(LINE(BUSY)) is FALSE or
- b. The auxiliary line (control definition) is entered, or
- c. The auxiliary line enters a Receive or Transmit Request.

If AUX(LINE(BUSY)) is TRUE when the primary line executes a \( \fork \) statement \\ \), the \( \fork \) statement \\ \) will act as a no-op.

## AUX(LINE(QUEUED)) Interrogate/Alter, CTR/CTR

This is a line-oriented (bit variable) that refers to the queued status of the auxiliary line of a full duplex line pair. The bit is set by the DCP operating system if and when an input message space is explicitly acquired by executing a (getspace statement) on the auxiliary line.

## AUX(LINE(TALLY[{0 or 1}])) Interrogate/Alter, CTR/CTR, 8

These are line-oriented (byte variable)s for the auxiliary line of a full duplex line pair, and can be used for any purpose by the NDL programmer. They can be accessed by either the primary or auxiliary line at any time.

### AUX(LINE(TOG [{0 or 1}])) Interrogate/Alter, CTR/CTR

These are line-oriented (bit variable)s for the auxiliary line of a full duplex pair, and may be used for any purpose by the NDL programmer. They may be accessed by either the primary or auxiliary line at any time.

#### **BCC**

Interrogate/Alter, CTR/CTR, 8

BCC is used by system software for the purpose of accumulating a Block Check Character when a station \(\lambda terminal definition\rangle\) defines horizontal parity as ODD or EVEN in the \(\lambda terminal parity statement\rangle\).

Block Check Character accumulation is an automatic function, if appropriate, of the \( \textit{receive statement} \) and \( \text{transmit statement} \). Block Check Character accumulation is based upon execusive-OR logic, that is, as characters are received or transmitted, they are exclusively OR-ed with the contents of BCC. It is the responsibility of the programmer to initialize BCC when appropriate. (Refer to the \( \text{initialize statement} \) under \( \text{request definition} \) or \( \text{control definition} \).)

If a station \(\lambda terminal definition\rangle\) does not define horizontal parity, BCC can be used as a temporary data storage area. It should be pointed out, however, that the value in BCC is destroyed by most constructs of the \(\lambda terminate statement\rangle\). Furthermore, since the intended purpose of BCC is to contain parity information, BCC and CRC[0] address the same data space. BCC cannot be used if a terminal uses Cyclic Redundancy Check.

When accumulating a Block Check Character, a convenient means to eliminate a specific character from the value accumulated in BCC is the \( \sum \) statement \( \).

Continued

#### **BCCERR**

Interrogate/Alter, CTR/CTR

BCCERR refers to bit 7 in the Error Flag Field of a result message, and conventionally indicates that a horizontal parity (BCC) error occurred while executing a *(receive statement)*. Refer to the semantics of the BCC option of the *(receive statement)*.

#### **BLOCK**

Interrogate, T

This bit references bit 29 in word zero of a message header. If TRUE, this bit indicates that more blocks (or messages) of a blocked transmission are to follow. Use of this bit implies a convention between the MCS and the NDL programmer for the purposes of providing blocked transmissions.

**BLOCK** is set **TRUE** implicitly as a result of execution of a **TERMINATE BLOCK** construct in a Receive Request.

#### **BLOCKED**

Interrogate, T

A synonym for BLOCK. Refer to BLOCK.

#### BREAK [RECEIVE]

Interrogate/Alter, CTR/CTR

This \( \lambda \text{bit variable} \) refers to bit 3 in the Error Flag Field of a message, and normally indicates that a break condition was sensed in a \( \lambda \text{receive statement} \rangle \). Refer to the semantics of the BREAK option of the \( \lambda \text{receive statement} \rangle \).

Note that if this bit is TRUE in a message to be returned to the MCS, the message is returned as a STATION EVENT (CLASS=1) message. Refer to the <u>B 6700/B 7700 DCALGOL Language Reference Manual</u> for more information regarding this message.

#### BREAK[TRANSMIT]

Interrogate/Alter, CTR/CTR

This \( \frac{bit variable}{\} \) refers to bit 5 in the Error Flag Field of a message, and normally indicates that a break condition was sensed while executing a \( \langle \text{transmit statement} \rangle \). Refer to the semantics of the \( \text{BREAK} \) option of the \( \langle \text{transmit statement} \rangle \).

Note that if this bit is TRUE in a message to be returned to the MCS, the message is returned as a STATION EVENT (CLASS=1) message. Refer to the B 6700/B 7700 DCALGOL Language Reference Manual for more information regarding this message.

#### **BUFOVFL**

Interrogate/Alter, CTR/CTR

This (bit variable) refers to bit 2 in the Error Flag Field of a message, and normally indicates that a cluster buffer overflow condition occurred while executing a (receive statement). Refer to the semantics under the BUFOVFL option of the (receive statement).

#### **CARRIAGE**

Interrogate, T

CARRIAGE is a carriage control variable, and is used to indicate if a carriage return is desired at the completion of the text transmission.

CARRIAGE is TRUE if message word [0]. [25:1] is zero.

This bit can be set by the I/O Intrinsics for a data communications file, or by the MCS.

#### CHARACTER

Interrogate/Alter, CTR/CTR, 8

CHARACTER is a line-oriented (byte variable).

CHARACTER contains the last character TRANSMITted or RECEIVEd on the line, unless otherwise altered by a \( \frac{fetch statement}{\rangle} \) or an \( \langle assignment statement \rangle \).

#### CONTROLFLAG

Interrogate/Alter, CTR/CTR

This \( \delta it variable \) refers to bit 12 in the Error Flag Field of a message, and normally indicates that the station defined control character was received. Refer to the CONTROL option of the \( \lambda receive statement \rangle \).

Note that if this bit is on in a message to be returned to an MCS, and the first character of the message is the control character of the station, the message is returned as a STATION EVENT (CLASS=1). Refer to the B 6700/B 7700 DCALGOL Language Reference Manual for more information regarding this message.

#### **CRC**

Interrogate/Alter, CTR/CTR

In (request definition)s and (control definition)s that use the Cyclic Redundancy Check, system software tests the status of the (bit variable) CRC before the execution of any (receive statement) or (transmit statement). If CRC is TRUE, the byte (or bytes) transmitted or received are calculated into the Cyclic Redundancy Check stored in the (byte variable)s CRC[0] and CRC[1]. If CRC is FALSE, bytes transmitted or received do not affect the Cyclic Redundancy Check.

Continued

CRC[{0 or 1}]
Interrogate/Alter, CTR/CTR, 8

System software uses the (byte variable)s CRC[0] and CRC[1] as a concatenated 16-bit information field to contain Cyclic Redundancy Check information for those stations whose (terminal definition)s define horizontal parity as CRC(16). If the (bit variable) CRC is TRUE, Cyclic Redundancy Check calculation is done using CRC[0] and CRC[1] as a 16-bit field, and the characters TRANSMITted or RECEIVEd. If CRC is FALSE, Cyclic Redundancy Check calculation is inhibited.

If a station \(\lambda terminal definition\rangle\) does not define horizontal parity, then CRC[0] and CRC[1] can be used as a temporary storage area. It should be pointed out, however, that the values in CRC[0] and CRC[1] are destroyed by most constructs of the \(\lambda terminate statement\rangle\). Additionally, since the intended purpose of these variables is storage of parity information, CRC[0] and BCC address the same byte. CRC[0] cannot be used for temporary data storage if the \(\lambda control definition\rangle\) or \(\lambda request definition\rangle\) uses BCC for Block Check Character accumulation.

#### CRCERR

Interrogate/Alter, CTR/CTR

CRCERR references bit 7 in the Error Flag Field of a result message, and conventionally indicates that an error in the Cyclic Redundancy Check occurred while executing a *(receive statement)*. Refer to the semantics of the CRC option of the *(receive statement)*.

## DISCONNECT

Interrogate/Alter, TR/TR

DISCONNECT references bit 12 in the Error Flag Field of a message, and indicates that a disconnect occurred on the line while executing a (request definition).

#### **ENDOFBUFFER**

Interrogate/Alter, TR/TR

ENDOFBUFFER references bit 17 in the Error Flag Field of a result message, and is conventionally used by a (request definition) to indicate when an overflow of the text buffer has occurred. Refer to the semantics of the ENDOFBUFFER option of the (receive statement).

#### **FORMATERR**

Interrogate/Alter, TR/TR

This bit references bit 10 in the Error Flag Field of a result message, and is conventionally used to indicate that a format error occurred while executing a *(receive statement)*. Refer to the RECEIVE *(string)* construct of the *(receive statement)*.

#### **INHIBITSYNC**

Interrogate/Alter, CTR/CTR

INHIBITSYNC is a line-oriented variable that causes actions as described under the \( \text{terminal inhibitsync} \) statement \( \text{.} \)

#### IR

Interrogate, CTR, 10-bit

IR addresses the 10-bit Input Register of the adapter cluster. This register contains hardware related control and data information for a line adapter.

IR can be interrogated using a *bit number* specification. *bit number* for IR range from zero through 9. For example, IR[0] addresses bit number zero of the Input Register.

Refer to the Burroughs Data Communications Processor Reference Manual or the DCP Reference Card for the meaning of the bits in IR.

### LINE(BUSY)

Interrogate/Alter, CTR/CTR

LINE(BUSY) is a line-oriented control information bit, and is used to allow or inhibit the interruption of the execution of a \( \chicontrol definition \) or \( \chicontrol definition \) on a single line. In the case of a full duplex line pair, LINE(BUSY) refers to the primary line. If this bit is TRUE, it indicates to the DCP operating system that the line is engaged in functions that must not be interrupted. If FALSE, it indicates to the DCP operating system that the line can be interrupted to initiate another function. LINE(BUSY) can be altered only by the primary line of a full duplex line pair.

A \(\chicontrol\) definition\(\rangle\) or \(\chicontrol\) definition\(\rangle\) is interrupted when LINE(BUSY) is FALSE if the DCP receives in its Request Queue a station-oriented DCWRITE message, and STATION(QUEUED) is FALSE for that station. If the message is a READ - ONCE ONLY (TYPE=34), STATION is set to that station index, and control is transferred to the Receive Request for that station. If the message is a WRITE (TYPE=33) DCWRITE message, STATION is set to that station index, and control is transferred to the Transmit Request for that station. If the message TYPE is neither of the above, the function associated with the message is executed and control resumes at the beginning of the line \(\chicontrol\) definition\(\rangle\), with the value of STATION equal to the index of the station for which the function was initiated. (Note that an interruption causes control to leave a \(\chicontrol\) definition\(\rangle\) or \(\rangle\) request definition\(\rangle\), and that control is not returned to the point where the interruption occurred.)

Continued

### LINE(BUSY) is set TRUE by system software when:

- a. The (control definition) is entered,
- b. A (request definition) is entered, or
- c. The line is the primary of a full duplex line pair, LINE(BUSY) is FALSE, and the auxiliary line executes a (fork statement).

Note that if LINE(BUSY) is TRUE when the auxiliary line of a full duplex line pair executes a \( \) fork statement \( \) acts as a no-op.

#### LINE(QUEUED)

Interrogate/Alter, CTR/CTR

LINE(QUEUED) is a line-oriented variable used to indicate whether or not (TRUE or FALSE, respectively) a message has been queued for any station on the line. It is set TRUE by system software when a message is inserted into an empty Station Queue of a station assigned to the line. It is the programmer's responsibility to set it FALSE when appropriate.

LINE(TALLY[{0 or 1}])
Interrogate/Alter, CTR/CTR, 8

LINE(TALLY[ $\{0 \text{ or } 1\}$ ]) are line-oriented variables for data storage, etc., available to the programmer.

An MCS may dynamically alter the line tallies by performing a SET/RESET LINE TOG/TALLY (TYPE=103) DCWRITE request.

LINE(TOG[{0 or 1}])
Interrogate/Alter, CTR/CTR

LINE(TOG[ $\{0 \text{ or } 1\}$ ]) are line-oriented variables for general information storage, etc., available to the programmer.

An MCS may dynamically alter the line toggles by performing a SET/RESET LINE TOG/TALLY (TYPE=103) DCWRITE request.

#### **LINEFEED**

Interrogate, T

LINEFEED is a carriage control variable and is TRUE when message word [0]. [24:1] is zero. If TRUE, LINEFEED indicates that a new line is required at the completion of the text transmission. This bit can be set by the I/O Intrinsics for a data communications file, or by the MCS.

#### LOSSOFCARRIER

Interrogate/Alter, CTR/CTR

**LOSSOFCARRIER** references bit 18 in the Error Flag Field of a result message, and is conventionally used to indicate that a loss of carrier occurred while executing a *(receive statement)*. Refer to the **LOSSOFCARRIER** option of the *(receive statement)*.

#### **MAXSTATIONS**

Interrogate, CTR, 8

**MAXSTATIONS** is a line-oriented  $\langle byte\ variable \rangle$  whose value is the maximum number of stations that can be assigned to the line.

MAXSTATIONS is initialized to the value defined in the \( \lambda \) line maxstations statement \( \rangle \) of the \( \lambda \) line maxstations statement \( \rangle \) does not appear in a \( \lambda \) line definition \( \rangle \), then MAXSTATIONS is initialized to the number of stations listed in the \( \lambda \) line station statement \( \rangle \). If neither statement appears, MAXSTATIONS is zero.

Within a \(\langle\control\) definition\(\rangle\), the valid range of values which may be assigned to the \(\langle\) by te variable\(\rangle\) STATION is between zero and MAXSTATIONS -1, inclusive.

#### **NAKFLAG**

Interrogate/Alter, TR/TR

NAKFLAG references bit 11 in the Error Flag Field of a result message, and conventionally indicates that a transmission was NAKed by the terminal. This bit is not set by system software, and its use is at the option of the programmer.

#### **NAKONSELECT**

Interrogate/Alter, TR/TR

NAKONSELECT references bit 16 of the Error Flag Field of a result message, and is conventionally used to indicate that a Transmit Request was NAKed when it attempted to select the terminal. This bit is not set by system software, and its use is at the option of the programmer.

#### **NOSPACE**

Interrogate, CTR

NOSPACE is a global variable that, when TRUE, indicates that a "no space" condition exists in the available space pool. NOSPACE is set by system software when the condition exists, and reset when the condition no longer exists.

Continued

#### **PAGE**

Interrogate, T

PAGE is a carriage control variable, and conventionally indicates whether a new page is required for the output device. For example, on a screen device, PAGE = TRUE could indicate to the Transmit Request that a home/clear sequence should be transmitted before or after the text is transmitted to the terminal. Refer to PAPERMOTION.

PAGE is set TRUE is message word [0].[26:1] = 1.

#### **PAPERMOTION**

Interrogate, T

**PAPERMOTION** is a carriage control variable that is conventionally used to indicate whether carriage control is desired before or after the message text is transmitted. If message word [0].[30:1] = 1, **PAPERMOTION** is set **TRUE**, and carriage control should be done before the text is transmitted; otherwise, carriage control after the text is transmitted.

#### **PARITY**

Interrogate/Alter, CTR/CTR

PARITY references bit 6 of the Error Flag Field in a result message, and indicates that a vertical parity error was detected when executing a *(receive statement)*. Refer to the PARITY option of the *(receive statement)*.

#### RETRY

Interrogate/Alter, CTR/CTR, 8

**RETRY** is a station-oriented variable, and is referred to as DCP RETRY.

The purpose of DCP RETRY is to record the number of attempts a \( \textit{request definition} \) has made to communicate with a terminal but failed as the result of some abnormal condition. Conventionally, the NDL programmer decrements RETRY (i.e., DCP RETRY) by one for each unsuccessful attempt at an operation until RETRY equals zero, then executes a TERMINATE ERROR.

When a \( \text{request definition} \) is initiated by the DCP, DCP RETRY is implicitly set to an initial value called DCP INITIAL RETRY. The default value of DCP INITIAL RETRY is specified by the NDL program in the \( \text{station retry statement} \).

By using the Message Retry Field in the message header (message word [2].[47:8]), the MCS can control the value assigned to DCP INITIAL RETRY, and therefore, is the initial value of DCP RETRY. If the Message Retry Field is 255, the value specified in the *(station retry statement)* assigned to DCP INITIAL RETRY, otherwise the value of the Message Retry Field is assigned to DCP INITIAL RETRY. The NDL program can restore the value of DCP RETRY to the value of DCP INITIAL RETRY at any time by executing the INITIALIZE RETRY construct.

All forms of the \(\lambda terminate statement \rangle \) which result in a message being returned to an MCS cause the current value of DCP RETRY to be stored in the Message Retry Field of the result message.

#### **SEOERR**

Interrogate/Alter, TR/TR

**SEQERR** references bit 14 in the Error Flag Field of a result message, and conventionally indicates that a sequence number overflow occurred as the result of the execution of an **INCREMENT SEQUENCE** construct. Refer to the *(increment statement)*.

#### **SEQUENCE**

Interrogate/Alter, CTR/CTR

SEQUENCE is a station-oriented bit that indicates whether or not (TRUE or FALSE, respectively) a \(\langle request \ definition \rangle \) is to perform automatic sequencing. SEQUENCE is controlled by a SET/RESET SEQUENCE MODE (TYPE=49) DCWRITE from the MCS. SEQUENCE can be set FALSE by the NDL program but can be set TRUE only by the controlling MCS. SEQUENCE can be set TRUE only if the \(\langle request \ definition \rangle \) for a terminal employs sequence number constructs such as TRANSMIT SEQUENCE, INCREMENT SEQUENCE, and STORE SEQUENCE. Use of automatic sequencing is the option and responsibility of the NDL programmer.

#### **SKIP**

Interrogate/T

SKIP is a carriage control variable. SKIP is used in conjunction with SKIPCONTROL to indicate a "skip to channel N" on an output device. If message word [0].[27:1] = 1, SKIP is set TRUE and SKIPCONTROL contains the channel number to skip to. Both SKIP and SKIPCONTROL can be set by the I/O Intrinsics for a data communications file, or by the MCS.

#### SKIPCONTROL

Interrogate, T, 8

SKIPCONTROL is used in conjunction with the \( \frac{bit variable}{s} \) SKIP and SPACE. If SKIP is TRUE, then SKIPCONTROL applies to SKIP. If SPACE is TRUE, SKIPCONTROL applies to SPACE. If neither are TRUE, SKIPCONTROL is undefined. Both SKIP and SPACE should not be TRUE concurrently. For a description of the function of this byte, refer to SPACE and SKIP. SKIPCONTROL is transferred to the Transmit Request in the message header of a WRITE (TYPE=33) DCWRITE in message word [0].[39:8], and can be set by the I/O Intrinsics for a data communications file, or by the MCS.

#### **SPACE**

Interrogate, T

SPACE is a carriage control variable. SPACE is used in conjunction with SKIPCONTROL to indicate the number of vertical lines to skip. If message word [0] [28:1] = 1, SPACE is set TRUE and SKIPCONTROL indicates the number of lines to skip. SPACE and SKIPCONTROL can be set by the I/O Intrinsics for a data communications file, or by the MCS.

Continued

#### **STATION**

Interrogate/Alter, C/C, 8

STATION is a line-oriented (byte variable) used in a (control definition) of a multi-station line to select a particular station with which the (control definition) wishes to interact. That is, to access the variables of a particular station, or INITIATE the Receive Request or Transmit Request of a station, the station index value associated with the station must be stored in STATION.

A station index value is associated with each station that is assigned to a logical line. At DCP initialization time, station index values are assigned sequentially, beginning at zero, to each station on a given line in the order that the stations were named in the \( \line \) line station statement \( \rangle \) of the \( \line \) line definition \( \rangle \).

After DCP initialization, an MCS can cause a station to be logically added to a line. When this occurs, a station index value becomes associated with the station. An MCS can also cause the logical removal of a station from a line. After such action, the station index value that was associated with the station no longer references a valid station. Thus, after DCP initialization, "holes" can exist in the sequence of valid station index values for a given line. A station index value can be "tested" to determine if it references a valid station by interrogating the STATION(VALID) (bit variable).

There is a maximum valid station index value associated with each line. That value is determined either by the \(\langle \line \text{maxstations statement}\rangle\) or by the \(\langle \line \text{station statement}\rangle\). (Refer to the \(\langle \line \text{maxstations statement}\rangle\) for more information.) This value can be obtained in a \(\langle \control \text{definition}\rangle\) by interrogating MAXSTATIONS.

#### STATION(ENABLED)

Interrogate, CTR

This is a station-oriented (bit variable) which refers to the "enabled" state of a station. When this variable is TRUE, the station is enabled for input, and the station Receive Request can be invoked. If STATION(ENABLED) is FALSE, the station is disabled for input, and attempts to invoke the Receive Request will be disallowed.

The setting of STATION(ENABLED) is initially defined by the (station enableinput statement) in the station definition, and may be altered by an MCS via the ENABLE INPUT (TYPE=35) and DISABLE INPUT (TYPE=36) DCWRITE.

#### STATION(FREQUENCY)

Interrogate, CTR, 8

STATION(FREQUENCY) is a station-oriented (byte variable), and is conventionally used to contain a relative polling frequency for polled stations. The initial value for STATION(FREQUENCY) is supplied by the (station frequency statement) for a station. It can be altered by an MCS via the ENABLE INPUT (TYPE=35) DCWRITE. Refer to the (station frequency statement).

#### **STATION(QUEUED)**

Interrogate, CTR, 8

STATION(QUEUED) is a station-oriented variable that indicates whether or not (TRUE or FALSE, respectively) there are any messages (output or enableinput) in the station queue. Note that if this variable is FALSE, the execution of an INITIATE REQUEST construct acts as a no-op.

#### STATION(READY)

Interrogate, CTR

If STATION(READY) is TRUE, the station associated with the station index stored in STATION is logically ready. No function (e.g., a Transmit Request or Receive Request) can be INITIATEd for the station if it is not ready. Stations can become not-ready as the result of the execution of a TERMINATE ERROR in one of its (request definition)s or as the result of the MCS executing a MAKE STATION NOT-READY (TYPE=37) DCWRITE.

#### STATION(TALLY)

Interrogate/Alter, CTR/CTR, 8

STATION(TALLY) is a station-oriented (byte variable) and is a general purpose variable which may be used by the NDL program for data storage. The initial value of STATION(TALLY) is zero. Note that STATION(TALLY) differs from TALLY [(tally number)] in that it cannot be directly STOREd in a message header.

#### STATION(VALID)

Interrogate, C

The STATION(VALID) bit indicates whether or not (TRUE or FALSE, respectively) there is a valid station associated with the station index value stored in STATION. Refer to the (byte variable) STATION.

#### **STOPBIT**

Interrogate/Alter, CTR/CTR

**STOPBIT** references bit 1 in the Error Flag Field of a result message, and conventionally indicates that a stop bit error was detected while executing a *(receive statement)*. Refer to the **STOPBIT** option of the *(receive statement)*.

#### **SYNCS**

Interrogate/Alter, CTR/CTR

SYNCS is a synonym for INHIBITSYNC. Refer to the INHIBITSYNC description.

Continued

#### TAB

Interrogate, T

TAB is a carriage control variable, and is conventionally used to indicate tabulation for the terminal. This bit is not set by I/O Intrinsics, and its use implies some established convention between the MCS and the NDL programmer. TAB is set TRUE if message word [0].[30:1] = 1.

TALLY [\(\langle tally number \rangle\)]
Interrogate/Alter, CTR/CTR, 8

TALLY [0], TALLY [1], and TALLY [2] are general purpose station-oriented (byte variable)s. They can be used for storage of 8-bit quantities such as counters, characters, etc. When the DCP is initialized, the TALLYs are initially zero unless a value is specified in a (station initialize statement). TALLYs may be initialized directly from a message header (message word [3].[23:24]) by utilizing the INITIALIZE TALLY [(tally number)] construct, thereby enabling an MCS to supply additional information to the DCP. The DCP can likewise transfer the value of a TALLY back to an MCS in a result message by utilizing the (store statement). Once a TALLY has been assigned a value, that TALLY retains that value until explicitly altered by the NDL program.

#### **TIMEOUT**

Interrogate/Alter, CTR/CTR

TIMEOUT references bit 0 (zero) of the Error Flag Field in a result message, and conventionally indicates that a timeout occurred while executing a *(receive statement)*. Refer to the TIMEOUT option of the *(receive statement)*.

TOG [\(\langle\) toggle number\(\rangle\)]
Interrogate/Alter, CTR/CTR

TOG [0] through TOG [7] are general purpose station-oriented (bit variable) s, often referred to as toggles. They can be used for storage of logical values (TRUE and FALSE). When the DCP is initialized, the value of the toggles is set to the value specified in the (station initialize statement), or, if such initialization is not specified, the initial value will be FALSE. Toggles can be assigned a value directly from a message header (message word [1].[39:8]) by utilizing the INITIALIZE TOG [(toggle number)] construct. Toggles can be stored into a result message by utilizing the (store statement). Once a toggle has been assigned a value, that toggle retains that value until explicitly altered by the NDL program.

#### **TOGS**

Interrogate/Alter, CTR/CTR, 8

TOGS addresses the eight  $\langle bit \ variable \rangle$ s TOG[0] through TOG [7]. For example, TOGS = 4"FF" sets TOG[0] through TOG [7] TRUE. TOG [0] is considered the low-order bit, and TOG [7] the high-order bit.

#### TRANERR

Interrogate/Alter, CTR/CTR

TRANERR references bit 9 in the Error Flag Field of a result message, and is conventionally used to indicate that a transmission number error occurred. Refer to the TRAN option of the (receive statement).

#### **WRUFLAG**

Interrogate/Alter, CTR/CTR

WRUFLAG references bit 13 in the Error Flag Field of a result message. If this bit is TRUE upon termination of a (request definition), the result message is returned to the MCS as a STATION EVENT (CLASS = 1) message. Refer to the <u>B 6700/B 7700 DCALGOL Language Reference Manual</u> for more information regarding this message.

#### APPENDIX A. RESERVED WORDS

The following is a complete list of reserved words used in the Network Definition Language. These words have special meaning to the compiler and cannot be used as *(identifier)s* or in any manner other than their defined meaning. Any synonym of a reserved word is shown adjacent to the word, in parentheses.

ABORT		BLOCKED	(BLOCK)
ADAPTER	(ADAPTOR)	BREAK	
ADAPTOR	(ADAPTER)	BUFFER	
ADDERR		BUFOVFL	
ADDRESS		BUSY	
AI		CARRIAGE	
ALTERNATE		CHAR	(CHARACTER)
ANSWER		CHARACTER	(CHAR)
ASCII		CLEAR	
ASC63		CLUSTERS	
ASC67		CODE	
ASC68		CONNECTION	
AUX	(AUXILIARY)	CONSTANT	
AUXILIARY	(AUX)	CONTINUE	
BACKSPACE	(BKSP)	CONTROL	
BAUDOT		CRC	
BCC		CRCERR	(BCCERR)
BCCERR	(CRCERR)	DCP	
BCD		DEFAULT	
BCL		DEFINE	
BEGIN		DELAY	
BINARY		DIALIN	
BKSP	(BACKSPACE)	DIALOUT	
BLKN		DIFFERENT	
BLKNERR		DIRECT	
BLOCK	(BLOCKED)	DISCONNECT	

#### **RESERVED WORDS (Cont)**

**HORIZONTAL DOWN DUPLEX ICTDELAY IDLE** DYNAMIC **EBCDIC** IF **ELSE ILLEGALCHR INCREMENT ENABLED ENABLEINPUT INHIBITSYNC** (SYNCS) **END INITIALIZE ENDOFBUFFER** INITIATE INPUT **ENDOFNUMBER** IR EQ (EQL) (LINEDELETE) (EQ) LD **EQL ERROR** LE (LEQ) (LE) **EVEN LEQ** LINE **EXCHANGE FALSE** LINEDELETE (LD) **FAMILY** LINEFEED **FETCH** LOGICALACK FILE LOGIN **FINISH** LOSSOFCARRIER (LSS) **FOR** LS LSS **FORK** (LS) **FORMAT MAXINPUT FORMATERR MAXSTATIONS FREQUENCY** MCS **GE** (GEQ) **MEMORY GEQ** (GE) **MICRO GETSPACE MILLI** GO MIN GT (GTR) MODE **GTR** (GT) **MODEM** 

**MSGSPACE** 

**HOME** 

#### **RESERVED WORDS (Cont)**

MYUSE SHIFT

NAKFLAG SKIP

NAKONSELECT SKIPCONTROL

NE (NEQ) SPACE NEQ (NE) SPO

NOINPUT STANDARD

NOISEDELAY STATION
NORMAL STORRIT

NORMAL STOPBIT
NOSPACE STORE

NOT SUM

NULL SYNCS (INHIBITSYNC)

ODD TAB

OUTPUT TALLY

PAGE TASK

PAPERMOTION TERMINAL

PARITY TERMINATE

PASSIVE TEXT
PAUSE THEN

PHONE TIMELIMIT

PTTC6I TIMEOUT

QUEUED TO READY TOG

RECEIVE TOGS

REMOTE TRAN

REQUEST TRANERR

RETRY TRANSLATETABLE

RETURN TRANSLATOR

SCREEN TRANSMISSION

SEC TRANSMIT

SECURITY TRANSMITDELAY

SEQERR TRUE

SEQUENCE TURNAROUND

## RESERVED WORDS (Cont)

TYPE WIDTH

UP WRAPAROUND

USER WRU

VALID WRUFLAG

WAIT

## APPENDIX B. TRANSMISSION CODES

## **BAUDOT CODE**

·									
B 6 -						0	0	1	1
i t b5-					<b>.</b>	0	1	0	1
5	b4 b3 b2 b1 \		Co l umn	0	1	2	3		
	0.	0	0	0	0	BLK	<b>T</b>	BLK	5
	0	0	0	1	Ì	Ε	Z	3	11
	0	0	1	0	2	LF	L	LF	3/4
	0	0	1	1	3	Α	W		2
	0	1	0	0	4	SPACE	Н	SPACE	DIAMOND
	0	1	0	1	5	S	Y	BELL	6
·	0	1	1	0	6	I	Р	8	, 0
	0	1	1	1	7	U	Q	7	1
	1_	0	0	0	8	CR	0	CR	9
	1	0	0	1	9	D	В	\$	5/8 ?
	1	0	1	0	10 (A)	R	G	4	8
	1	0	1	1	11(B)	J	FIGS		FIGS
	1	1	0	0	12(C)	N	М	7/8	•
	1	1	0	1	13(D)	F	X	1/4	/
	1	1	1	0	14 (E)	С	V	1/8	3/8;
	1	1	1	1	15(F)	K	LTRS	1/2	LTRS

## **DATA REPRESENTATION**

EBCDIC GRAPHIC	BCL GRAPHIC	HEX.	EBCDIC INTERNAL	DECIMAL VALUE	EB CARI	CDIO		OCTAL	BCL INTERNAL	BCL EXTERNAL	CAR	BCL D C	ODE	USASCII X3.4-1967
Blank		40	0100 0000	64	No P	unch	es	60	11 0000	01 0000	No F	unc	hes	010 0000
[	*	4A	0100 1010	74	12	8	2	33	01 1011	11 1100	12	8	4	101 1011
		4B	0100 1011	75	12	8	3	32	01 1010	11 1011	12	8	3	010 1110
<		4C	0100 1100	76	12	8	4	36	01 1110	11 1110	12	8	6	011 1100
(		4D	0100 1101	77	12	8	5	35	01 1101	11 1101	12	8	5	010 1000
+		4E	0100 1110	. 78	12	8	6		•	11 1010				010 1011
1	<b>←</b>	4F	0100 1111	79	12	8	7	37	01 1111	11 1111	12	8	7	111 1100
&		- 50	0101 0000	80	12			34	01 1100	11 0000	12			010 0110
]		5 <b>A</b>	0101 1010	90	11	8	2	76	11 1110	01 1110	0	8	6	101 1101
\$		5 <b>B</b>	0101 1011	91	11	8	3	52	10 1010	10 1011	11	8	3	010 0100
*		5C	0101 1100	92	11	8	4	53	10 1011	10 1100	11	8	4	010 1010
)		5 <b>D</b>	0101 1101	93	11	8	5	55	10 1101	10 1101	11	8	5	010 1001
;		5E	0101 1110	94	11	8	6	56	10 1110	10 1110	11	8	6	011 1011
-	$\leq$	5F	0101 1111	95	11	8	7	57	10 1111	10 1111	11	8	7	
- ·		60	0110 0000	96	11			54	10 1100	10 0000	11			101 1111
1		61	0110 0001	97	0	1		61	11 0001	01 0001	0	1		010 1111
	•	6B	0110 1011	107	0	. 8	3	72	11 1010	01 1011	0	8	3	010 1100
%		6C	0110 1100	108	0	8	4	73	11 1011	01 1100	0	8	4	010 0101
	<b>≠</b>	6D	0110 1101	109	0	8	5	74	11 1100	01 1010	0	8	2	010 1101
>		6E	0110 1110	110	0	8	6	16	00 1110	00 1110	8	6		011 1110
?		6F	0110 1111	111	0	8	7	14	00 1100	00 0000	*			011 1111
:		. 7A	0111 1010	122	8	2		15	00 1101	00 1101	8	- 5		011 1010
#	***	7B	0111 1011	123	8	3		12	00 1010	00 1011	8	3		010 0011
@		7C	0111 1100	124	8	4		13	00 1011	00 1100	8	4		100 0000
	≥	. 7D	0111 1101	125	8	5		17	00 1111	00 1111	8	.7		010 0111
=		7E	0111 1110	126	8	6		75	11 1101	01 1101	0	8	5	011 1101
"		7 <b>F</b>	0111 1111	127	8	7		77	11 1111	01 1111	0	8	7	010 0010

# **DATA REPRESENTATION (Cont)**

EBCDIC GRAPHIC	BCL GRAPHIC	HEX.	EBCDIC INTERNAL	DECIMAL VALUE	EBCDIC CARD CODE	OCTAL	BCL INTERNAL	BCL EXTERNAL	BCL CARD CODE	USASCII X3.4-1967
(+)PZ	+	C0	1100 0000	192	12 0	20	01 0000	11 1010	12 0	
A		Cl	1100 0001	193	12 - 1	21	01 0001	11 0001	12 1	100 0001
В		C2	1100 0010	194	12 2	22	01 0010	11 0010	12 2	100 0010
$\bar{\mathbf{C}}$		C3	1100 0011	195	12 3	23	01 0011	11 0011	12 3	100 0011
D		C4	1100 0100	196	12 4	24	01 0100	11 0100	12 4	100 0100
E		C5	1100 0101	197	12 5	25	01 0101	11 0101	12 5	100 0101
F		<b>C6</b>	1100 0110	198	12 6	26	01 0110	11 0110	12 6	100 0110
G		C7	1100 0111	199	12 7	27	01 0111	11 0111	12 7	100 0111
Н		C8	-1100 1000	200	12 8	30	01 1000	11 1000	12 8	100 1000
1		<b>C</b> 9	1100 1001	201	12 9	31	01 1001	11 1001	12 9	100 1001
(!)MZ	x	D0	1101 0000	208	11 0	40	10 0000	10 1010	11 0	010 0001
J		D1	1101 0001	209	11 1	41	10 0001	10 0001	11 1	100 1010
K		D2	1101 0010	210	-11 2	42	10 0010	10 0010	11 2	100 1011
L		<b>D3</b>	1101 0011	211	11 3	43	10 0011	10 0011	11 3	100 1100
M		D4	1101 0100	212	11 4	44	10 0100	10 0100	11 4	100 1101
N		D5	1101 0101	213	11 5	45	10 0101	10 0101	11 5	100 1110
0		D6	1101 0110	214	11 6	46	10 0110	10 0110	11 6	100 1111
P	•		· 1101 0111	215	11 7	47	10 0111	10 0111	11 7	101 0000
Q		D8	1101 1000	216	- 11 8	50	10 1000	10 1000	11 8	101 0001
R		<b>D9</b>	1101 1001	217	11 9	51	10 1001	10 1001	11 9	101 0010
¢		E0	1110 0000	224	0 8 2		•	00 0000		•
S		<b>E2</b>	1110 0010	226	0 2	62	11 0010	01 0010	0 2	101 0011
T		E3	1110 0011	227	0 3	63	11 0011	01 0011	0 3	101 0100
U		<b>E4</b>	1110 0100	228	0 4	64	11 0100	01 0100	0 4	101 0101
V		E5	1110 0101	229	0 5	65	11 0101	01 0101	0 5	101 0110
$\mathbf{W}$	e filosofie de la compansión de la compa	<b>E6</b>	1110 0110	230	0 6	66	11 0110	01 0110	0 6	101 0111
X		<b>E</b> 7	1110 0111	231	0 7	67	11 0111	01 0111	0 7	101 1000
Y		<b>E8</b>	1110 1000	232	0 8	70	11 1000	01 1000	0 8	101 1001
Z		E9	1110 1001	233	0 9	71	11 1001	01 1001	0 9	101 1010
0		F0	1111 0000	240	0	00	00 0000	00 1010	0	011 0000
1		Fl	1111 0001	241	1	01	00 0001	00 0001	1	011 0001
2	$\mathcal{L}_{\mathcal{L}_{\mathbf{a}}}$	F2	1111 0010	242	2	02	00 0010	00 0010	2	011 0010
3		F3	1111 0011	243	3	03	00 0011	00 0100	3	011 0100
4	∓ <del>C</del> r	F4	1111 0100	244	4	04	00 0100	00 0100	4	011 0100

## **DATA REPRESENTATION (Cont)**

EBCDIC GRAPHIC	BCL GRAPHIC	HEX.	EBCDIC INTERNAL	DECIMAL VALUE	EBCDIC CARD CODE	OCTAL	BCL INTERNAL	BCL EXTERNAL	BCL CARD CODE	USASCII X3.4-1967
5		F5	1111 0101	245	5	05	00 0101	00 0101	5	011 0101
6		F6	1111 0110	246	6	06	00 0110	00 0110	6	011 <b>0110</b>
7		F7	1111 0111	247	7	07	00 0111	00 0111	7	011 0111
8		F8	1111 1000	248	8	10	00 1000	00 1000	8	011 1000
9		F9	1111 1001	249	9	11	00 1001	00 1001	9	011 1001

## **NOTES**

- 1. EBCDIC 0100 1110 also translates to BCL 11 1010.
- 2. EBCDIC 1100 1111 is translated to BCL 00 0000 with an additional flag bit on the most-significant bit line (8th bit). This function is used by the unbuffered printer to stop scanning.
- 3. EBCDIC 1110 0000 is translated to BCL 00 0000 with an additional flag bit on the next to most significant bit line (8th bit). As the print drums have 64 graphics and spaces, this signal can be used to print the 64th graphic. The 64th graphic is a "CR" for BCL drums and a "t" for EBCDIC drums.

- 4. The remaining 189 EBCDIC codes are translated to BCL 00 0000 (?code).
- 5. The **EBCDIC** graphics and **BCL** graphics are the same except as follows:

	<u>BCL</u>		<b>EBCDIC</b>
≥		•	(single quote)
<b>x</b>	(multiply)	!	
≤	· · · · · · · · · · · · · · · · · · ·	· · · · · · · · · · · · · · · · · · ·	(not)
≠ ≠			(underscore)
←		·	(or)

# PTTC/6 CODE

	\b <sub>7</sub>						0	0	0	0	1	1	1	1
\	B b6						0	0	1	, 1	0	0	1	1
	t b5						0	1	0	1	0	1	0	1
	5	ь4 ↓	b3 ↓	b2 ▼	b1 ↓	Co l umn Row	0	1	2	3	4	5	6	7
		0	0	0	0	0	SPACE	9	-	+ &	SPACE	DELTA	BACK/	• <
		0	0	0	1	1	1	/	j	a	>	QUES	J	Α
	,	0	0	-	0	2	2	S	k	b	)	S	К	В
		0	0	1	1	3	3	t	<u>l</u>	С	;	• 1	L	С
	•	0	1	0	0	4	4	سرن	m	- d	SBLANK	U	М	D
• .		0	1	0	1	5	_5	٧	n	e	(	V	N	Ε
		0	1	ī	0	5	6	W	0	f	•	W	0	F
		0	1	سنا		7	7	×	Р	9	11	X	Р	G
		[]	0	0	0	8	8	У	q	k	*	Υ	Q	Н
-		1	0	0	1	9	9	Z	i	i	[	Z	R	1
		1	0	1	0	10(A)	0	#***	M7	PZ	]	GRPMRK	GAMMA	SQ. RT
		1	0	1	1	11(B)	+	,	\$	•	SEGMARK	•	V.BAR	. •
		1	1	0	0	12(C)	PN	ВҮ	RFS	PF	PN	ВҮ	RES	PF
		1	1	0	1	13(D)	Rs	LF	NL	нт	RS	LF	NL	нт
		1	1	1	0	14(E)	UC	EOR	BS	LC	UC	EOB	BS	LC
		1	1	1	1	15 (F)	EOT	PRE	IL	DEL	EOT	PRF	11	DEL

#### APPENDIX C. SOURCE INPUT FORMAT AND CODING FORM

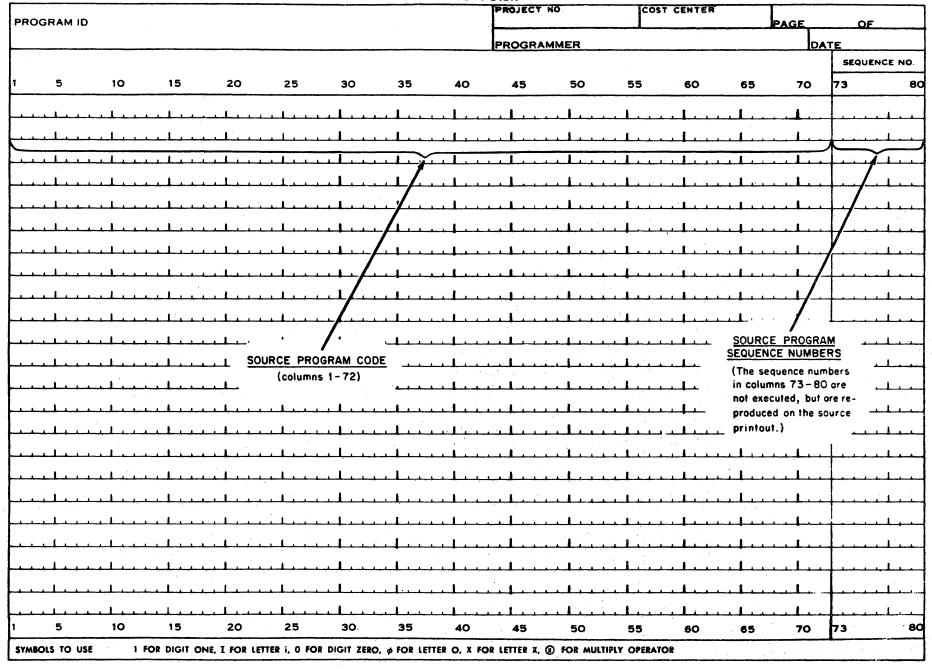
### **SOURCE INPUT FORMAT**

An NDL source program is represented by a set of ordered external records. The external records could come from cards, tape, disk, remote device, or a combination of these. The source information on any given record must be divided into two areas. Character positions 1-72 are assumed to contain elements of the Network Definition Language (described in sections 2 through 6 of this manual) for compilation by the compiler. Character positions 73 through 80 are assumed to contain information regarding the sequence of the input record; specifically, this area is for sequence numbers. Sequence numbers are optional.

There is no fixed format for source information in character positions 1 through 72. This information can appear in a free format form, with the following exceptions: elements contained on the card must comply with any syntactical restrictions, and syntactical items cannot be continued from record to the next. For example, the reserved word TERMINAL cannot begin on one source record and continue on the next.

#### **CODING FORM**

To facilitate keypunching, as well as to provide the programmer with a suggested format to follow in writing his source program, printed programming forms are often used. An example of such a form appears on the following page.



#### APPENDIX D. COMPILE-TIME OPTIONS

#### **COMPILER CONTROL STATEMENTS**

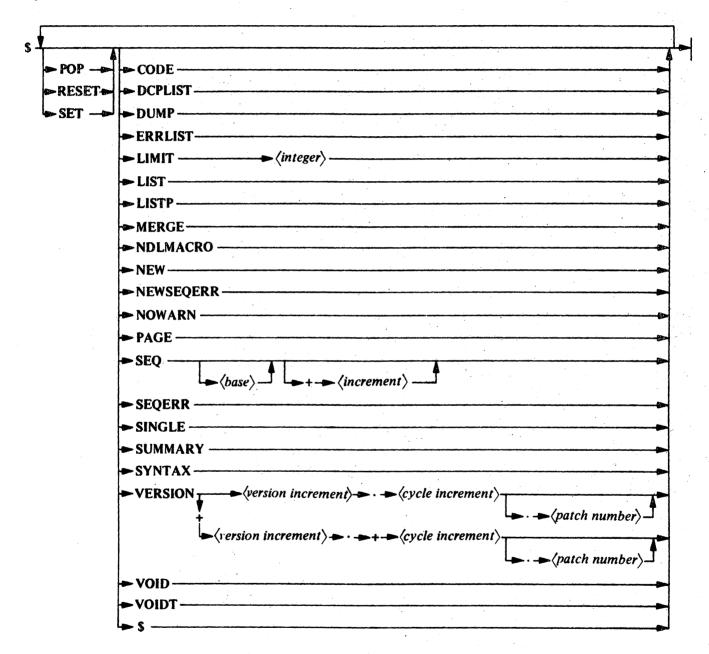
The user is provided with the compile-time ability to control the manner in which the compiler processes the source input that it accepts. The user can specify the manner in which the compiler is to receive the source input, the consequences of certain syntax errors, and the form of the generated compiler output. The compiler control statement is the medium by which these constraints are communicated to the compiler. Such statements are entered into the compiler by cards in the same manner as source language statements. Compiler control statements, entered as input to the compiler via option control cards, can occur at any point in the compiler input files and must contain only compiler control information.

An option control card is identified by the appearance of a dollar sign (\$) in the first or second column of the card. If the \$ is placed in card column 2, the option control card image is placed in the updated symbolic file (NEWTAPE) if such a file is generated. Compilation control information is punched in the succeeding columns through column 72, with an eight-digit sequence number in columns 73 through 80. All blanks in columns 73 through 80 represent the lowest-value sequence number. An option control card with no other compiler information causes the card image in the secondary input file that has the same sequence number to be ignored.

The basic element of compiler control information is the compiler option, which can be invoked by the appearance of its name on an option control card. Two mutually exclusive states are associated with the majority of these options: SET and RESET; various compiler functions are dependent upon the states of such options. Default states are assigned to these compiler options, and the desired state of such an option can be specified on an option control card. Such option control cards can also contain arguments associated with the option. The balance of compiler options are parameter options with which no states are associated. The functions performed by these latter options are initiated by the appearance on an option control card of the appropriate option name and any related arguments.

## **OPTION CONTROL CARDS**

## **Syntax**



#### **Semantics**

The purpose of a compiler control statement is the assignment of a desired value or state (SET or RESET) to an indicated compiler option(s). Such a control statement must begin with either an explicit or an implicit option action. An explicit option action is defined as one of the following mnemonics: SET, RESET, or POP.

An implicit option action is indicated when a compiler control statement contains only the names of options and no explicit option action. In the latter case, all options named in the compiler control statement are assigned the state SET, and all other options are assigned the state RESET.

If a compiler control statement begins with the option action SET, the options following the option action are assigned the state SET; the states of all other options are unchanged. If the compiler control statement begins with the option action RESET, the options following the option action are assigned the state RESET; the states of all other options are unchanged. If the specified option action is POP, then the options have not been changed previously from their default states. The states of all other options are unchanged. The following statements are examples of compiler control statements employing the SET, RESET, and POP option actions.

- **\$ SET LIST SINGLE**
- **\$ RESET VOID**
- **S POP NEW NEWSEOERR**
- \$ SET SEQ 0+100

An option that has a default state of RESET is initially assigned a 48-bit stack word filled with zeros; an option that has a default state of SET is initially assigned a 48-bit stack word with a 1 on top and zeros in the remaining positions. The top stack position denotes the state of the option at any time. Each SET option action causes the stacks allocated to the designated standard options to be pushed down one bit and a 1 to be placed at the top of each of these stacks. Each RESET causes the appropriate option stacks to be pushed down one bit and a 0 to be placed at the tops of these stacks. POP causes the stacks correponding to the designated options to be POPped up one bit, causing the associated options to revert to their immediate previous states. Since the size of these option stacks is 48 bits, a maximum history of 48 states can be recorded. When an option control card appears that has a standard option name and an implicit option action, the resultant action is identical to that which would have resculted had all 48 bits of each standard option stack been RESET and followed by an explicit SET performed on each indicated option. For example, after the appearance of an option control card containing:

#### \$ SINGLE

the history stack for the SINGLE option contains a 1 in the top stack position and all zeros in the following positions. The history stack for each of the other compiler options would then contain all zeros. A compiler control statement that applies to compiler options begins with an explicit or implicit option action and contains a list of options to which the option action is to apply. This statement ends when the next implicit option action is encountered on the compiler control card or when a percent sign is encountered on the compiler control card or when a percent sign is encountered or column 72 of the card is reached. The compiler options affected by the compiler control card retain the indicated states for all input cards with sequence numbers greater than the sequence number on the compiler control card that has the control statement, or the physically succeeding input cards for a deck in which all sequence numbers are blank, until another compiler control card is encountered that alters the option states. The following illustration (figure D-1) is an example of a card that has compiler control statements employing option actions:

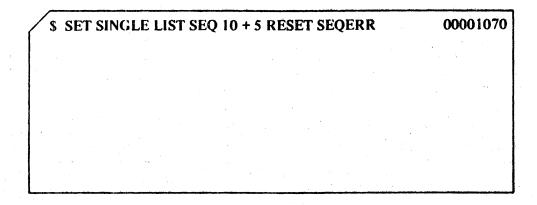


Figure D-1. Option Control Card

The option control card assigns the state SET to the options SINGLE, LIST, and SEQ, with the sequencing arguments of 10 and +5. It also assigns the state RESET to the option SEQERR. The card has the sequence number 00001070 in columns 73 through 80.

#### **OPTIONS**

The compiler recognizes the following identifiers as valid compiler option names:

CODE	NOWARN	
DCPLIST	PAGE	
DUMP	SEQ	
ERRLIST	SEQERR	
LIMIT	SINGLE	
LIST	SUMMAR	Y
LISTP	SYNTAX	
MERGE	VERSION	
NDLMACRO	VOID	
NEW	YOIDT	. •
NEWSEQERR	\$	

The compiler options are discussed alphabetically in the following paragraphs. The default state of each option is indicated in parentheses following the option name; the function performed by the option is discussed in the paragraph accompanying the same.

If an option control card is empty, it has no effect on other options; however, if there is a card image on the symbolic file with the same sequence number as the empty option control card, the image on the symbolic file is deleted.

The compiler options are as follows:

#### CODE (RESET)

The code option causes the printout to contain the compiler-generated object code.

## DCPLIST (RESET)

If SET, lists code addresses of each source statement on the LINE file. (There will be two separate lists addresses if used for two DCPs, three for three DCPs, etc.)

## DUMP (RESET)

If SET, causes a "raw dump" listing of NIF on the LINE file.

#### **ERRLIST (RESET)**

The ERRLIST option causes syntax error information for CANDE to be written on the ERRORFILE file. When a compilation error is detected in the source input, an error message is written in the ERRORFILE file. This option is provided primarily for use when the compiler is called from a remote terminal by the CANDE language, but it can be used regardless of the manner in which the compiler is called. When the compiler is called from CANDE, the default state of the ERRLIST option is SET and ERRORFILE is automatically equated to the remote device involved.

#### LIMIT (cannot be SET or RESET)

The integer parameter allows the user to control compiler error terminations. The proper format for the **LIMIT** option is as follows:

## LIMIT (integer)

Compilation is terminated if the number of errors detected by the compiler equals or exceeds the *(integer)* If no LIMIT statement appears, a default error limit of 150 is assigned unless the compilation is initiated through CANDE, in which case the default error is 10.

### LIST (SET; RESET for CANDE)

The LIST option causes a printout to be generated on the compiler output LINE file. The contents of such printouts are specified in the preceding paragraphs describing compiler features. If the LIST option is RESET, only syntax error messages and compilation information are listed.

#### LISTP (RESET)

When SET, the LISTP option causes patches and input records from the compiler CARD file to be included on the printout while records from the compiler TAPE file are excluded. This option is effective only if the LIST option is RESET. If the LIST option is SET, the state of LISTP is ignored. Therefore, the LISTP or the LIST option causes a printout to be generated when SET.

#### **MERGE (RESET)**

When SET, the MERGE compiler option causes primary input, CARD file, to be merged with secondary input, TAPE file, to form the total input to the compiler. If matching sequence numbers occur, the primary input overrides. If the MERGE option is RESET, only primary input is used and secondary input is totally ignored. Therefore, the total input to the compiler when the MERGE option is SET consists of all card images from the CARD file, and all card images from the TAPE file that do not have sequence numbers that can be found on cards in the CARD file.

#### **NDLMACRO (RESET)**

If **SET**, the **NDL MACRO** interface code will be printed following each statement within a *(request definition)* or *(control definition)*.

#### **NEW (RESET)**

When the state of the NEW option is SET, the merged input from the CARD and TAPE files is placed on the updated symbolic output file NEWTAPE. This file is coded in EBCDIC and is structured in 15-word records and 450-word blocks. Therefore, it can later be used as input to the compiler through the TAPE file. All option control cards in the merged CARD and TAPE file input are placed on the NEWTAPE file when NEW is SET and only if the initial \$ sign on these cards is in card column 2.

The NEW option can be SET and RESET as necessary by option control cards appearing at any point in the input file. Such option control cards can also be placed on the NEWTAPE file if the \$ signs on these cards are in column 2.

The NEWTAPE file is created despite the occurrence of syntax errors in the source input. This file can be used as a secondary input for a later compilation.

The NEWTAPE file can be label-equated so that, for example, the output goes to magnetic tape.

#### **NEWSEOERR (RESET)**

The NEWSEQERR option causes sequence errors on the NEWTAPE file to be flagged. If sequence errors occur and the NEWSEQERR option is SET, the NEWTAPE file is not locked, and the message NEWTAPE NOT LOCKED {number of errors} NEWTAPE SEQUENCE ERRORS is printed on the printout. NEWTAPE, NIF, and DCPCODE files are not locked.

#### **NOWARN (RESET)**

When SET, suppresses any compiler warnings from appearing on the LINE file.

PAGE (cannot be SET or RESET)

The PAGE compiler option must appear on a option card without an option action preceding it. When a PAGE option card appears, the printout is spaced to the top of the next page, but only if the LIST option is SET.

## SEQ (RESET)

The proper format of the SEQ option is as follows:

If the SEQ option is SET, the printout and the new secondary source language file, NEWTAPE, contain new sequence numbers as defined by the \langle base \rangle and \langle increment \rangle are unspecified, a base of 0 and increment of 10 are assumed.

This option has effect only when the LIST and/or NEW options are also SET. The sequence numbers that appear on the card images in these files when the SEQ option is RESET are identical to the sequence numbers on the corresponding cards in the input file.

#### Example

This compiler control card specifies that, when the state of the SEQ option is SET, sequencing begins with the sequence number 00000100 and proceeds in increments of 100.

#### SEQERR (RESET)

The SEQERR option causes sequence errors on the TAPE file to be flagged. If sequence errors occur and the SEQERR option is SET, DCPCODE and NIF files are not locked, and the message CODE FILE NOT LOCKED { number of errors } TAPE SEQUENCE ERRORS is printed on the printout.

#### SINGLE (RESET)

The SINGLE option causes the printout to be single-spaced. When the SINGLE option is RESET, the printout is double-spaced. (Note that double-spacing is default.)

### **SUMMARY (RESET)**

If SET, lists on the LINE file the memory space allocations for user translation tables and terminal message space allocations for each DCP.

### SYNTAX (RESET)

When SET, the source program is checked for syntax errors only. DCPCODE and NIF files are not generated.

VERSION (SET, RESET, and POP are ignored by the compiler)

The VERSION compiler option allows the user to specify an initial version number for a source program, to replace an existing version number, or to append an existing version number.

## Examples

- \$ VERSION 25.010.010
- \$ VERSION +01,+001.010

When compiling with the NEW compiler option SET and a VERSION compiler card appears in the symbolic, and if the patch deck contains a VERSION compiler option, the new symbolic is updated to the version, cycle, and patch number on the last VERSION compiler card in the patch deck. The sequence number must be less than the one in the symbolic.

#### **VOID (RESET)**

If the VOID option is SET, all input, other than \$ cards, from the TAPE and the CARD files is ignored by the compiler until the VOID option is RESET or POPped into a RESET state. The ignored input is neither listed nor included in the updated symbolic file regardless of the states of the LIST and NEW options. The VOID option can be RESET, once it is SET, only by a option control card in the CARD file.

#### **VOIDT (RESET)**

If the VOIDT option is SET, all secondary input, other than \$ cards, from the TAPE file is ignored by the compiler until the VOIDT option is RESET or POPped into a RESET state. Therefore, while the VOIDT option is SET, only primary input is compiled. The ignored input is neither listed nor included in the updated symbolic file regardless of the states of the LIST and NEW options. The VOID option can be RESET, once it is SET, only by an option control card in the CARD file.

#### \$ (RESET)

When SET, the dollar sign (\$) option causes the printout of all subsequent  $\langle option\ control\ card \rangle$  images when the LIST option is SET. This option appears as \$SETS or \$ \$.

#### APPENDIX E. COMPILER SOURCE AND OBJECT FILES

#### **COMPILER FILES**

Compiler communication is handled through various input and output files (figure E-1). Cards, disk, or magnetic tape can be specified as source language input media. Input must be in the input format defined in the preceding sections. The compiler has the capability of merging, on the basis of sequence numbers, input from cards, tape, or disk. When inputs are being merged, indications of text insertions or replacements can be made to appear on the printout. In addition to the printout, the compiler can also generate updated symbolic files. These files can be created in addition to the compiler-generated output code file.

## Input Files

The primary compiler input file is a card file with the internal name CARD; the secondary input file is a serial disk file with the internal name TAPE. The presence of the primary file (CARD) is required for each compilation; the presence of the secondary file (TAPE) is optional for each compilation. When two card images, one from the CARD file and the other the TAPE file have the same sequence number, the former is primary and is compiled, and the latter is ignored. This is the standard mode of handling source language input. File CARD can be either BCL-coded with 10-word records or EBCDIC-coded with 14-word records and can be either blocked or unblocked. File TAPE can be BCL-coded with 10-word records and 150-word blocks, or EBCDIC-coded with a 14- or 15-word record and 420- or 450-word blocks. Both the CARD file and the TAPE file can be label-equated (via the FILE system control card) to change the TITLE and KIND of the file. The TAPE file is used as input only when the MERGE compiler option is SET.

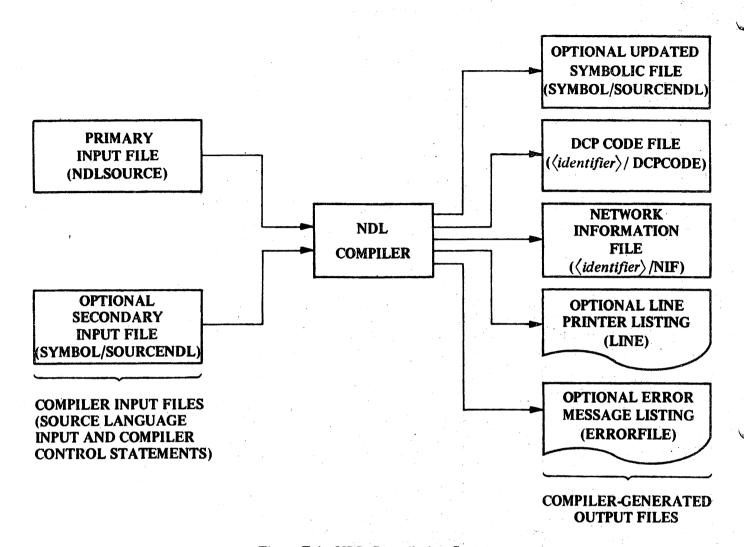


Figure E-1. NDL Compilation System

#### **Output Files**

Output files produced by the compiler consist of the DCP code file, the Network Information File, an updated symbolic file, a line printer printout, and an error message file. The DCP code file has the internal name DCPCODE and is saved on disk after the compilation unless the COMPILE system control card specifies compilation for syntax only, or unless syntax errors are detected in the source language input by the compiler. If compilation for library is specified, then the DCPCODE and NIF files are saved on disk. The title of the saved DCP code file is identical to the program name \( \langle identifier \rangle \) appearing on the COMPILE system control card with the suffix of \( \DCPCODE. \)

The title of the saved Network Information File is identical to the program name (identifier) appearing on the COMPILE system control card with the suffix of /NIF.

The updated symbolic file is, by default, a disk file generated only if the compiler option NEW is SET. This file contains the compilation source input or a selected portion of this input as specified by the state of the NEW compiler option. It can be used as the TAPE file for a succeeding compilation. This output file has the internal file name NEWTAPE and contains EBCDIC-coded 15-word records in 450-word blocks.

The line printer printout is an optional print file that is created unless the compiler option LIST is RESET. (The LIST option is SET by default unless the compilation is initiated through CANDE.) The file has the internal name LINE, consists of 22-word EBCDIC-coded records, and contains the following information:

- a. Source and compiler control statements used as input to the compiler.
- b. Error messages and error count.
- c. Number of input card images scanned.
- d. Elapsed compilation time.
- e. Processing time required for compilation.
- f. Total number of words of DCP code generated.
- g. Number of disk segments required for the DCP code file.
- h. Title of the generated code file.

Depending upon the specified setting of the LIST and CODE compiler options, the line printer printout can contain more (or less) information than the basic items listed above. Card images from the CARD file are denoted on the printout by a C after the card contents. Card images from the TAPE file are denoted by a T in this location. A P denotes a patch of a TAPE card image.

The output error-message file with the internal file name and assigned title of ERRORFILE is an optional line printer file that is created when the ERRLIST compiler option is SET. This file is normally employed for compilations initiated through CANDE, in which case ERRLIST is SET by default and the ERRORFILE file is assigned to the remote device involved. The ERRORFILE file can also be used for compilations initiated through the card reader. This file is assigned EBCDIC-coded 12-word records that result in a line width of 72 characters, allowing the file to be used as output to a remote terminal or card punch without truncation of text. When a syntax error is detected, an error message is written following the line of text. The error message consists of an explanatory message and indicates the probable cause of the error.

#### Compiler File Table

Table E-1, NDL Compiler Files, lists the external name of the file (i.e., the name one would label-equate to), the internal name of the file (i.e., the name used when the file is declared within the compiler), the purpose served by the file, the default KIND of the file, the code used to store file data, the default record size (MAXRECSIZE) and block size (BLOCKSIZE) of the file, and a brief commentary on the specific file. The attributes of any of these files can be changed by the use of FILE system control cards directed to the compiler.

Table E-1. NDL Compiler Files

	T		<del>, , , , , , , , , , , , , , , , , , , </del>		·	<del>,</del>	
EXTERNAL NAME	INTERNAL NAME	PURPOSE	KIND	CODE	RECORD SIZE	BLOCK SIZE	COMMENTS
NDLSOURCE	CARD	Input Card File	CARD READER	EBCDIC	14 Words	Blocked or Unblocked	Required for each compilation. Primary compiler input file; may be label-equated to change file attributes.
				BCL	10 Words		CANDE file is equated to this file automatically for compilations initiated through CANDE. Default title is NDLSOURCE.  BUFFERS = 2.  FILETYPE = 8.
SYMBOL/SOURCENDL	ТАРЕ	Input Disk File	DISK	EBCDIC	14 or 15 Words	420 or 450 Words	Optional file; need not be present for each compilation. Secondary compiler input file; selected as input by SETting
				BCL	10 Words	150 Words	MERGE compiler option. Can be label- equated to change file attributes as desired. The default title is  SYMBOL/SOURCENDL.  FILETYPE = 8.
⟨identifier⟩/DCPCODE	DCPCODE	DCP Code File	DISK	Hexadecimal	30 Words	420 Words	Generated DCP code file. Saved or discarded and assigned a title as indicated by compilation method. For CANDE compilations, the title becomes:  OBJECT/(identifier)/DCPCODE.
SYMBOL/SOURCENDL	NEWTAPE	Updated Symbolic Output File	DISK	EBCDIC	15 Words	450 Words	Optional output file produced when NEW compiler option is SET. This file contains portions of the source input and is label-equatable. It is suitable for use as a TAPE file for a later
							compilation.  BUFFERS = 2.  AREASIZE = 1000.  AREAS = 20.

Table E-1. NDL Compiler Files (Cont)

EXTERNAL NAME	INTERNAL NAME	PURPOSE	KIND	CODE	RECORD SIZE	BLOCK SIZE	COMMENTS
⟨identifier⟩/NIF	NIF	Network Information File	DISK		30 Words	420 Words	Generated Network Information File (NIF). Saved or discarded and assigned a title as indicated by compila- tion method. For CANDE compilations the title becomes: OBJECT/\(\( identifier \right) \)/NIF.
LINE	LINE	Line Printer Printout	LINE PRINTER or REMOTE	EBCDIC	22 Words	22 Words	Optional and label-equatable file. Produced when the compiler option LIST is SET.
ERRORFILE	ERROR- FILE	Error Listing Output File	LINE PRINTER	EBCDIC	12 Words	12 Words	Optional error listing file produced when ERRLIST compiler is SET. Contains card images and error messages. Automatically provided for CANDE input.

## **INDEX**

Item	Page
ADDERR	5-60, 5-75, 5-136, 5-156 5-32, 5-112, 6-6
ADDRESS	5-30, 5-42, 5-62, 5-110, 5-130, 5-138, 5-157 5-157
AI	5-40, 5-124, 6-6
	5-63
	5-6, 5-83, 6-3
	5-6, 5-83, 6-3
(assignment statement)	5-6, 5-83
AUX(LINE(BUSY))	
AUX(LINE(QUEUED))	
AUX(LINE(TALLY[{U or I}]))	6–7 6–7
AUX(LINE(IUG[{U O7 I}]))	
available line adapters	5–61
RACKSDACE	5-85, 5-112, 5-158
	5-85.
	5–39
	5-30, 5-40, 5-43, 6-7
	5-32, 5-112, 6-8
	6–1
\(\delta bit variable \rangle \cdots	
	5-126, 6-8
	5-126, 6-8
	5-8, 5-13, 5-32
	5-8, 5-86
	5–8, 5–86
	5-13, 5-33, 5-91, 5-113, 6-9 6-1, 6-3
,	·
CARRIAGE	5–160, 6–9
CHARACTER 5-31.5-41.5	<b>-42</b> , 5 <b>-43</b> , 5 <b>-111</b> , 5 <b>-122</b> , 5 <b>-125</b> , 5 <b>-130</b> , 5 <b>-131</b> , 6 <b>-9</b>
	3–2
\(\lambda\) \(\lam	3–3
(hexadecimal character)	3-4
\(\lefter\rightarrow \cdot\)	
(single character)	
character translation	5-6, 5-9, 5-83, 5-87, 5-162, 5-184
CLEAR	5-161 5-7, 5-9, 5-87, 5-162
CODE	5-7, 5-9, 5-87, 5-162
(code statement)	5–9, 5–87
coding form	C-2
	5-75, 5-77, 5-156, 5-170 E-2
	D-1
	E-3

Item			Page
compiler files			<b>177</b> 4
compiler files			
compound statements	• • • • • • • • • • • • • • • • • • • •	• • • • • • • • •	5-10, 5-88
\(\langle if statement \rangle \cdots \)			E 01 E 100
"GO TO byte variable" construct		• • • • • • • •	5-21, 5-100
CONSTANT			
(constant definition)			
$\langle constant$ identifier $\langle \ldots \rangle$			
construct terminator			2-2
CONTINUE	• • • • • • • • • • • • • • • • • • • •	5-11, 5-33	, 5–89, 5–113
$\langle continue\ statement  angle \ \dots $			5-11, 5-89
CONTROL			
(control definition)			
$\langle control\ identifier angle \qquad \ldots \qquad \ldots$			5-163
control statement's			
(assignment)			
BREAK			
CODE			5–9
compound			5–10
CONTINUE			5-11
DELAY			5-12
ERROR switch			
FINISH			
FORK			
GO TO			
IDLE			
IF			
INCREMENT			
INITIALIZE			
INITIATE			
PAUSE			
RECEIVE			
SHIFT			
SUM			
TRANSMIT			
WAIT			
CONTROLFLAG		•••••	3 <del>-44</del>
CDC		F 111 F 1	21 6 0 6 10
CRC	5-31, 5-43, 5-104,	3-111, 3-1	31, 0-9, 0-10
CRCERR	• • • • • • • • • • • • • • • • • • • •	3-3	4, 5-114, 6-9
D 4 C C 4 !!			٠, .
Data Comm Controller			
data communication files			
FAMILY			
data representation			
(DCP definition)	• • • • • • • • • • • • • • • • • • • •		5–45
$\langle \mathit{DCP}$ exchange statement $ angle \ldots \ldots$		• • • • • • • • •	5–46
$\langle DCP$ memory size statement $ angle \ldots \ldots$			5-51
〈DCP number〉			5–60

Item						Page
DCP programs						1-6
(DCP statement)s		•		. •		
EXCHANGE						5-46
MEMORY						5-51
TERMINAL						
DCP Tables						
<b>(DCP terminal stateme</b>	ent)	• • • • • • • • • • •	• • • • • • • • • •			5-52
DEFAULT	·····		• • • • • • • • • •	5-64, 5-135	, 5–140, 5–	155, 5-164
(default line identifier)	· <b>&gt;</b>					5-59, 5-64
default station identif	fier)					5–135
default terminal idens	tifier >			• • • • • • • • •		5-155
Definitions	• • • • • • • • • • • •					5-1
CONSTANT						
CONTROL						
DCP						
FILE						
LINE						
MCS						
MODEM						
REQUEST						5-81
STATION						5-134
TERMINAL						
DELAY						
(delay statement)						
delay time						
DIALIN						
DIALOUT	• • • • • • • • • • •		• • • • • • • • • •	• • • • • • • •		5-70
(digit)						3–3
DISCONNECT						6-9
DUPLEX					5-70, 5-	-71, 5-166
	and the second second					
ENABLEINPUT					5—	127.5-141
END					5_33 5_	114 5-167
ENDOFBUFFER						
ENDOFNUMBER						
ERROR		• • • • • • • • • •	* * * * * * * * * * * *	5-13, 5-3	1, 3-91, 3-	112, 3-12/
(error switch statemen						
EXCHANGE	• • • • • • • • • • •	• • • • • • • • • •				5-46
FAMILY			• • • • • • • • • •			5-57
FETCH						5-94
(fetch statement)						594
file definition						
(file family statement)						
(file identifier)				• • • • • • • • • •		,-30, 3+3/
(file statement)						
FAMILY	• • • • • • • • • • • • • • • • • • • •					5–57
files	• • • • • • • • • •					5-56
the state of the s			<ul> <li>A second of the s</li></ul>		·	

Item			•	Page
〈finish statement〉				5-16, 5-95
FINISH TRANSMIT	• • • • • • • • • • • • • • • • • • • •			5-16, 5-95
FORK				
(fork statement)	• • • • • • • • • • • •	· • • • • • • • • • • •		5-17, 5-96
FORMATERR				
FREQUENCY	• • • • • • • • • • • • • • • • • • • •		,	5-142
full duplex constructs, executable				
(continue statement)				
\(\fork statement\rangle \cdot\tag{fork statement}\)				
\(\rightarrow\ait\) \(\ldots\)	• • • • • • • • • • • • • • • • • • • •	• • • • • • • • • •	• • • • • • • • • • • •	. 5–44, 5–132
GETSPACE				5-97
(getspace statement)				
GO TO				
(go to statement)				
(hexadecimal character)				
HOME				
HORIZONTAL				
(horizontal parity variant)	• • • • • • • • • • • • •		• • • • • • • • • • •	5–176
WARRING TO R. A. W. P.				P 481
ICTDELAY	• • • • • • • • • • • • • • • • • • • •		• • • • • • • • • • • •	5-1/1
(identifier)				
IDLE				
(idle statement)				
〈if statement〉	• • • • • • • • • • • • •	• • • • • • • • • •	• • • • • • • • • • • • •	5 21 5 100
ILLEGALCHR	• • • • • • • • • • • • • • • • • • • •	• • • • • • • • • • • • • • • • • • •	• • • • • • • • • • •	3-21, 3-100
(increment statement)	• • • • • • • • • • • • • • • • • • • •	• • • • • • • • • • •		5 5-109
INHIBITSYNC	• • • • • • • • • • • • • •	• • • • • • • • • • •		5_170 6 11
INITIALIZE	• • • • • • • • • • • • • • • • • • • •	• • • • • • • • • •		5_24 5_104
(initialize statement)				
INITIATE				
initiate receive delay				
(initiate statement)				
initiate transmit delay				
input files, compiler				
input format, source	• • • • • • • • • • • •		· • • • • • • • • • • • • • • • • • • •	C-1
(integer)	• • • • • • • • • • •			3-8
`IR				
keywords			•	2-2
Roy Words			• • • • • • • • • • • •	2-2
⟨label⟩	• • • • • • • • • • • • •			3–9
language components				
(letter)				
LINE				5-58
line adapters and adapter classes	• • • • • • • • • • • •			5-60

Item		Page
〈line adapter class statement〉		
(line address statement)		
line control		
(line default statement)		
$\langle$ line definition $\rangle$		
(line endofnumber statement)		
cline identifier		
cline maxstations statement		
(line modem statement)		
(line phone statement)		
line section requirements	• • • • • •	5-46
(line statement)s		
ADAPTER		
ADDRESS		
ANSWER		
DEFAULT		
ENDOFNUMBER		
MAXSTATIONS		
MODEM		
PHONE		
STATION		
TYPE		
(line station statement)		
(line type statement)		
LINEDELETE		
LINE(BUSY)		
LINE(QUEUED)		
LINE(TALLY[{0 or 1}])		
LINE(TOG[{0 òr 1}])		
LINEFEED		
logical assignment		
LOGICALACK		
LOSSOFCARRIER 5–14, 5–35, 5–78, 5-	-92, 5-	116, 6–13
MAXINPUT		5-174
MAXSTATIONS		
MEMORY		
(MCS definition)      MCS reconfiguration		5 A7
Message Control System.		
MODEM 5-		
(modem adapter statement)	-07, 5-	74, 3-140 5-75
(modern daupter statement)	•••••	5-73 5-74
(modem definition)		
(modern lossofcarrier statement)		
(modern noisedelay statement)		
(modern noisedeury statement)		
\modelin biai cinterit/	• • • • • •	5-15

Item	4					Page
/modern state		•			· ·	i de la filosofia de la composición de La composición de la
(modern statement)s						
ADAPTER						
LOSSOFCARRIER						
TRANSMITDELAY .						
(modem transmitdelay state	$ment$ $\rangle$ $\ldots$					5-80
MYUSE	· • • • • • • • • • • • • • • • • • • •					5-147
			• •	r		
NAKFLAG						6-13
NAKONSELECT					The second secon	
NDL program unit						
NDL syntax convention						
NOINPUT						
NOISEDELAY						
NOSPACE					• • • • • •	6–13
object files						E-1
options, compiler						D-4
output files						
PAGE					5-148	5-175 6-14
PAPERMOTION						
PARITY	• • • • • • • • • • • • •		5 1.4	5 35 5 0	5 116	5 176 6 1A
PAUSE		• • • • • • • • • • • • • • • • • • • •	• • • • • • • •	• • • • • • • • •	• • • • • • •	5 20 5 100
\(\rho_{\text{pause statement}}\) \(\cdots\)						
PHONE	• • • • • • • • • • • •	• • • • • • • • •	• • • • • • • • •		• • • • • •	5-68, 5-149
RECEIVE	• • • • • • • • • • •					5-29, 5-109
⟨receive address size⟩						
Receive Request						5–82
(receive statement)				5-29, 5-3	7, 5–38, 5	<del>-109,</del> 5-118
⟨receive statement⟩, allowab	le combinations			: • • • • • • • • • •		5-118
relational operators					• • • • • •	5-22
synonyms						5-22
(remark)						
REQUEST						5-177
\(\request\) definition\(\rangle\) \(\ldots\)						
\request identifier \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \						
request statement						
, .		• • • • • • • • •	• • • • • • • •	• • • • • • • • •	• • • • • •	3-62, 3-1//
<pre> ⟨request statement⟩s assignment</pre>				and the second		
assignment		• • • • • • • •			• • • • • •	5-83
BACKSPACE						
BREAK					• • • • • •	5-86
compound						
CONTINUE						
DELAY			• • • • • • • •			

Item				Page
(reaues	st statement\s (Cont)	1 S		
F	ERROR			5–91
	FETCH			
F	FINISH		 	5-95
	ORK			
	GETSPACE			
	GO TO			
	F			
	NCREMENT			
	NITIALIZE			
	NITIATE			
	AUSE			
	RECEIVE			
	TORE			
				5–124
_	ERMINATE			
	RANSMIT		41	
	VAIT			
Reque				
F	Receive Request		 	5-82
	ransmit Request			
reserve	ed words		 	
RETR	Y	• • • • • • • • • • •	 $\dots 5-41, 5-10$	<b>14,</b> 5–125, 5–150, 6–14
	Υ		5-41, 5-10	<b>94,</b> 5–125, 5–150, 6–14
scope				
scope	of NDL			
scope o	of NDL			1–1 6–1
scope o o SCRE	of NDL			1–1 6–1 5–178
scope o o SCRE SEOE	of NDL			
scope o SCRE SEQE SEQU	of NDL of variables EN RR ENCE		5–23, 5–4	1-1 
scope O SCRE SEQE SEQUE Seque	of NDL of variables EN ER ENCE nce mode		5–23, 5–4	1-1 5-178 6-15 ,5-103,5-123,5-131 5-23,5-103
scope SCRE SEQE SEQU sequer SHIFT	of NDL of variables EN RR ENCE nce mode		5-23, 5-4	1-1 
scope SCRE SEQE SEQU sequer SHIFT	of NDL of variables EN RR ENCE nce mode f tatement		5–23, 5–4	1-1 
scope  SCRE SEQE SEQU sequer SHIFT (shift s.	of NDL of variables EN RR ENCE nce mode f tatement		5–23, 5–4	1-1 
scope SCRE SEQE SEQU sequer SHIFT Shift s. SKIP SKIP	of NDL of variables EN RR ENCE nce mode f tatement		5–23, 5–4	1-1 
scope SCRE SEQE SEQU sequer SHIFT Shift s SKIP SKIP Source source	of NDL of variables EN RR ENCE nce mode f tatement control e files input format		5-23, 5-4	1-1 6-1 5-178 6-15 , 5-103, 5-123, 5-131 5-23, 5-103 5-39, 5-121 6-15 6-15 E-1 C-1
scope  SCRE SEQE SEQU sequer SHIFT Shift s SKIP SKIPC source source source	of NDL of variables EN RR ENCE nce mode f tatement CONTROL e files e input format e program structure		5-23,5-4	1-1
scope  SCRE SEQU sequer SHIFT Shift s SKIP SKIPC source source source source	of NDL of variables EN RR ENCE nce mode f tatement CONTROL e files e input format e program structure e size		5-23,5-4	1-1
scope  SCRE SEQE SEQU sequer SHIFT Shift s SKIP SKIPC source source source SPAC	of NDL of variables EN RR ENCE nce mode f tatement control e files e input format e program structure e size E		5-23,5-4	1-16-15-1786-15 , 5-103, 5-123, 5-1315-23, 5-1035-395-39, 5-1216-156-156-156-156-15
scope  SCRE SEQE SEQU sequer SHIFT  Shift s SKIP SKIPC source source source SPACI  (space)	of NDL of variables EN RR ENCE nce mode of tatement control e files e input format e program structure e size E		5-23,5-4	1-16-15-1786-15 , 5-103, 5-123, 5-1315-23, 5-1035-395-39, 5-1216-156-156-156-156-15
scope  SCRE SEQU sequer SHIFT  Shift s SKIP SKIPC source source source source space  space statem	of NDL of variables EN RR ENCE nce mode ftatement CONTROL e files e input format e program structure e size E		5-23, 5-4	1-1
scope  SCRE SEQU sequer SHIFT (shift s SKIP SKIPC source source source (space) statem	of NDL of variables EN RR ENCE nce mode f tatement  CONTROL e files e input format e program structure e size E		5-23, 5-4	1-1
scope  SCRE SEQU sequer SHIFT Shift s SKIP SKIPC source source source source source source source statem	of NDL of variables EN RR ENCE nce mode f tatement  CONTROL e files e input format e program structure e size  Ence mode CONTROL		5-23,5-4	1-1
scope  SCRE SEQE SEQU sequer SHIFT Shift s SKIP SKIPC Source source source source space statem	of NDL of variables EN RR ENCE nce mode f tatement  control statement control statement correct control statement correct corr		5-23,5-4	1-1 6-1 5-178 6-15 , 5-103, 5-123, 5-131 5-23, 5-103 5-39 5-39, 5-121 6-15 6-15 E-1 C-1 4-1 5-184 5-184 5-45
scope  SCRE SEQU sequer SHIFT Shift s SKIP SKIPC Source source source SPACI (space) statem	of NDL of variables EN RR ENCE nce mode f tatement  control statement control statement correct control statement correct corr		5-23,5-4	1-1 6-1 5-178 6-15 , 5-103, 5-123, 5-131 5-23, 5-103 5-39, 5-121 6-15 6-15 E-1 C-1 4-1 5-184 5-45 5-57 5-58

ltem	Page
statements (Cont)	
(request statement)	5–83
station statement	
(terminal statement)	
STATION	
\(\station\) adapter statement \(\station\).	
station address statement	
station control character statement	
station default statement	
station definition \	
station enableinput statement	
station frequency statement	
station identifier	
station initialize statement	
\(\rangle\)station logicalack statement\(\rangle\) \(\rangle\)\(\rangle\)	
station MCS statement	
\(\rangle\)station modem statement\(\rangle\)\(\rangle\)	
\(\station\) myuse statement\(\rangle\)	
\(\station\) my use statement \(\station\) \(\station\) page statement \(\station\)	
\(\rangle\) station place statement \(\rangle\) \(\rangle\) station phone statement \(\rangle\) \(\rangle\)	
station retry statement	
\(\rangle\) station retry statement \(\rangle\) s	3-130
ADAPTER	5-136
ADDRESS.	
CONTROL	
DEFAULT.	
ENABLEINPUT. FREQUENCY.	
INITIALIZE	
LOGICALACK	
MCS	,
MODEM	
MYUSE	
PAGE	
PHONE	
RETRY	
TERMINAL	
WIDTH	
(station terminal type statement)	
(station width statement)	5–152
STATION(ENABLED)	6-16
STATION(FREQUENCY)	6-16
STATION(QUEUED).	
STATION(READY)	
STATION(TALLY)	
STATION(VALID)	
STOPBIT 5–14, 5–35, 5–92	
STORE	5-122

Item						•			Page
(store statement)	• • • • • • • • • • • • •	• • • • • • •						• • • • • •	5-122
(string)									
SUM			· • • • • •				5-40,	5-124,	5 - 125
(sum statement)								. 5-40,	5 - 124
\(\rangle switch number \rangle \)									5-31
SYNCS								. 5-170	6-17
syntactic variable	s				· • • • • •			• • • • • • •	. 2-2
syntax convention	ns								. 2-1
key words.									. 2-2
	riables								
construct ter	rminator		. <b></b> .						. 2-2
(system identifier)	>								3-13
				. •					
TAB							• • • • • • •	· · · · · · ·	6-18
TALLY					5-41,5	-105, 5-1	23, 5-12	5, 5-143	6-18
TERMINAL							5-151,	5-154.	5-155
	statement >								
	size statement) .								
terminal backspace	ce character statem	$ ent\rangle$ .							5-158
	ze statement)								
terminal carriage	character statemen	$ t\rangle$							5-160
	racter statement								
	tement \								
	statement \								
	AULT								
(terminal default s	statement \							5-154.	5-164
	$\langle n \rangle$								
	tatement \								
	racter statement								
terminal home ch	aracter statement								5-168
terminal identifie									
	naracter statement								
	$\gamma$ nc statement $\rangle$ .								
	racter delay staten								
`	te character stateme	. /							5-172
	character statemen	. /							5-173
terminal maxinpu									5-174
terminal page stat									5-175
terminal parity st									5-176
terminal request s						· • • • • • • • • •			5-177
terminal screen st	• •	. t		*					5-178
terminal statemen									5-153
terminal statemen									5-153
ADAPTER.									5-156
									5-157
	E								5-158
BUFFER									5-159

Item						Page
(terminal statement)s	(Cont)	• • • • • • • • • •				5-153
LINEFEED				• • • • • • • • • •	• • • • • • • • • •	5-173
REQUEST		• • • • • • • • • •		• • • • • • • • •	• • • • • • • • • •	5-170
TDANGMICCIO		• • • • • • • • • •	• • • • • • • • • • •			5 100
	N					
terminal timeout state						
terminal transmission						
terminal turnaround s	ctatomont	iemeni/	• • • • • • • • • • •	•••••	• • • • • • • • • •	5-181
rerminal width staten						
terminal wru characte						
TERMINATE						
(terminate statement)	,					5-126
TEXT						
\(\time\) \(\time\)						,
TIMEOUT		5–14.	5-30, 5-36, 5	5-92, 5-110,	5-117.5-17	9. 6–18
(timeout time)						5-110
TOG				5-105.	5-123, 5-14	3.6-18
\(\lambda toggle number \rangle \dots \) TRAN \(\)			• • • • • • • • •		3–16	5,5-123
TRAN				5-31, 5-43, 5	-104, 5-111	,5-131
TRANERR					.5-36.5-11	7.6 - 19
(translatetable definiti	$on \rangle \dots \dots$				• • • • • • • • •	5-184
translation, character			5–6	6, 5–9, 5–83,	5-87, 5-162	, 5–184
translation table struc	cture					5-184
data insertion						5-185
TRANSLATETABLE						
(translatetable identifi						
transmission codes						
TRANSMIT					5–42	. 5-130

Item				* ************************************			Page
(transmit address size)			 				 5-157
Transmit Request	• • • • • • • • •		 				 5-82
(transmit statement)			 : • • • • •		· • • • •		 . 5-42, 5-130
TURNAROUND			 				 5-181
TYPE	· · · · · · · · · · ·		 · · · · · ·	· • • • • • •	. <b></b> .		 5-70
use of NDL			 				 1-1
				· .	: '		
Value assignment			 <b></b>				 5-6, 5-83
variables							
⟨byte variable⟩							
(bit variable)							
description of							
function of							
scope of							
VERTICAL		• • • • • •	 	• • • • • •	· • • • •		 5–176
, WAIT		• • • • • •	 	• • • • • •	. <b></b> .	• • • • • •	 . 5-44, 5-133
(wait statement)		•••••	 • • • • •	• • • • • •		• • • • • •	 . 5-44, 5-133
(wait time)							
WIDTH							
WRU							
WRUFLAG			 				 6-19

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