

RT-11 Device Handlers Manual

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This manual describes the structure of device handlers, how to write your own device handler, and provides specific programming information about distributed RT-11 device handlers.

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Document Structure

This manual is divided into the following two chapters:

- Chapter 1, *Device Handlers*, describes the recommended structure of device handlers and provides detailed information on how to write a device handler.
- Chapter 2, *Programming for Specific Devices*, alphabetically presents programming information for specific distributed device handlers.

Audience

This manual is written for those users of the RT-11 operating system who want to understand distributed device handlers and write their own device handlers.

Conventions

The following conventions are used in this guide.

Convention	Meaning
Black print	In examples, black print indicates output lines or prompting characters that the system displays. For example: <pre>.BACKUP/INITIALIZE DL0:F*.FOR DU1:WRK Mount output volume in DU1;; continue? Y</pre>
Red print	In examples, red print indicates user input.
Braces ({ })	In command syntax examples, braces enclose options that are mutually exclusive. You can choose only one option from the group of options that appear in braces.
Brackets ([])	Square brackets in a format line represent optional parameters, qualifiers, or values, unless specified otherwise.
Lowercase characters	In command syntax examples, lowercase characters represent elements of a command for which you supply a value. For example: <pre>DELETE filespec</pre>

Convention	Meaning
UPPERCASE characters	In command syntax examples, uppercase characters represent elements of a command that should be entered exactly as given.
<code>[RET]</code>	<code>[RET]</code> in examples represents the RETURN key. Unless the manual indicates otherwise, terminate all commands or command strings by pressing <code>[RET]</code> .
<code>[CTRL/x]</code>	<code>[CTRL/x]</code> indicates a control key sequence. While pressing the key labeled Ctrl, press another key. For example: <code>[CTRL/C]</code>

Associated Documents

Basic Books

- *Introduction to RT-11*
- *Guide to RT-11 Documentation*
- *PDP-11 Keypad Editor User's Guide*
- *PDP-11 Keypad Editor Reference Card*
- *RT-11 Commands Manual*
- *RT-11 Quick Reference Manual*
- *RT-11 Master Index*
- *RT-11 System Message Manual*
- *RT-11 System Release Notes*

Installation Specific Books

- *RT-11 Automatic Installation Guide*
- *RT-11 Installation Guide*
- *RT-11 System Generation Guide*

Programmer Oriented Books

- *RT-11 IND Control Files Manual*
- *RT-11 System Utilities Manual*
- *RT-11 System Macro Library Manual*
- *RT-11 System Subroutine Library Manual*
- *RT-11 System Internals Manual*

- *RT-11 Volume and File Formats Manual*
- *DBG-11 Symbolic Debugger User's Guide*

Chapter 1

Device Handlers

The term *device handler* can mean three things, depending on the context in which it is used. A device handler can be:

- The source program

This is a .MAC file that is distributed with RT-11 or you write.

- The file image

This is a .SYS file that is distributed with RT-11 or the assembled and linked source program you write.

- The memory image

This is the part of the file image that resides in memory; the memory resident portion of the device handler. Not all of the file image is normally loaded in memory. The first block (block 0) of the file image, for example, is temporarily loaded when the monitor requires information that is stored in handler block 0. The memory resident portion of the device handler begins at block 1 of the file image. Therefore, block 1 of the file image is the beginning of the memory image.

To write a device handler, you first need to know what points to consider in the planning stage. These points are listed and cross-referenced in the first sections of this chapter. The points that have not been treated elsewhere in this manual are then described in detail. Device handler structure and a skeleton outline of a typical handler are covered here. After this, details are given on the optional features available to handlers and their implementation. Optional features include internal queuing, SET options, device I/O timeout support, special functions, error logging, and special services available in mapped systems.

To write a bootstrap for a system device, you first need to know the differences between a standard handler and a system device handler. These differences are discussed in several sections before the final sections of the chapter, where you will find explained the assembly, installation, testing, and debugging procedures for the new handler.

Be sure to also read Chapter 5 of the *RT-11 System Internals Manual*, as that chapter can help you decide whether you need to write an in-line interrupt service routine or a device handler.

1.1 How to Plan a Device Handler

The most important part of writing a device handler is taking the time to plan the whole process carefully. Follow these guidelines:

- Get to know your device
- Study the structure of a standard device handler
- Study the skeleton device handler
- Think about using the special features
- Study the sample handlers
- Prepare a flowchart of the device handler
- Write the code
- Install, test, and debug the handler

1.1.1 Get to Know Your Device

Learning about the characteristics of your device and the bus interface is crucial to writing a handler that works correctly. Review the appropriate material in Chapter 5 of the *RT-11 System Internals Manual* so that you can answer all the pertinent questions about your device before you attempt to write a handler for it.

1.1.2 Study the Structure of a Standard Device Handler

Section 1.2 describes the structure of a standard device handler. Read this section carefully; your handler should conform to this structure.

1.1.3 Study the Skeleton Device Handler

Section 1.2.10 contains a skeleton outline of a standard device handler. You can use this outline as a starting point when you begin to write your own handler.

1.1.4 Think About Using the Special Features

Sections 1.4 through 1.10 describe the special features available to device handlers. Read these sections carefully to determine whether any features are applicable to your handler.

1.1.5 Study the Sample Handlers

Appendix A contains assembly listings of three RT-11 device handlers (DL, DX, and XL) with extensive explanatory comments. Study these listings until you feel comfortable with the organization of the handlers, and you understand how they implement some of the special features. Obtain listings of handlers for other devices that resemble yours; you may be able to use some of the code that is already written.

1.1.6 Prepare a Flowchart of the Device Handler

Preparing a flowchart for your handler can help you plan the contents of the various sections. Flowcharting can also help you spot loose ends and errors in your programming logic. Unfortunately, flowcharts are not much help in pointing out potential race conditions. (A race condition is a situation in which two or more asynchronous processes attempt to modify the same data structure at the same time; as a result, the data structure is corrupted and the integrity of the processes is compromised.) Therefore, when you design the handler, examine every step carefully

and keep in mind what would happen if an interrupt occurred at each instruction. This kind of planning can help you avoid race conditions later.

1.1.7 Write the Code

If you have followed the recommended steps so far, writing the code for the device handler should be relatively simple. You must write Position-Independent Code (PIC) for the handler. Review the chapter on PIC code in the *PDP-11 MACRO-11 Language Reference Manual* if you are not already familiar with it. Copy as much code as possible from the commented device handlers in Appendix A or from other reliable sources. Start with a general outline that conforms to the structure presented in Section 1.2 and then add details to reflect the specifics of your particular device. When you have thoroughly checked the code for logic errors and it assembles properly, you are ready to test and debug it.

1.1.8 Install, Test, and Debug the Handler

Sections 1.14 and 1.15 show how to install a new device handler and how to begin testing and debugging it.

1.2 Structure of a Device Handler

For ease of explanation and understanding, the RT-11 handler source program is described as having the following six sections:

- Preamble

The preamble is the information section of the source. Much of the information you put in the preamble as arguments to macro parameters and as system conditionals is associated with symbols that are used by macros in other handler sections. The macros you use in the preamble section create many of the handler's data structures and further define the handler.

- Header

The header section is where you code the beginning of the memory resident portion of the handler.

- I/O initiation

The I/O initiation section contains the first executable instructions; the code to get the handler ready to perform data transfers. The I/O initiation section is able to use data structures and symbols that were defined in the previous sections and defines further handler characteristics.

- Interrupt service

The interrupt service section is the heart of the handler. It contains the code that processes interrupts as they are received from the device. It handles aborts and manages the handler queue.

- I/O completion

The I/O completion section contains code to inform the monitor of the success or failure of the interrupt processing and perform appropriate actions depending on success or failure.

- **Handler termination**

The handler termination section is the tail of the handler. It contains code to build tables and handler service routines. Being at the end of the handler, it defines a symbol that is used to determine the size of the handler.

The complexity of the coding you must write is reduced because the RT-11 system macro library (SYSMAC.SML) provides device handler macros to generate much of the required code.

You should read and think about the following points before working through this section:

- Although the various macro parameters are listed and briefly described in this chapter, you should consult the *RT-11 System Macro Library Manual* for complete parameter argument descriptions. Refer to that manual as you read this chapter.

Some of the macros that you use to write a device handler are interdependent. For example, the device status word is created from symbols that SYSMAC.SML equates based on arguments you supply to .DRDEF parameters. Those symbols are then used by .DRBEG to create the device status word and store it into the handler's block 0.

- RT-11 distributes a library of the system data structures (SYSTEM.MLB), described in the *RT-11 System Internals Manual*. In this section, the symbols that identify handler data structures and the elements in those structures are as defined in SYSTEM.MLB. If your device handler is assembled with SYSTEM.MLB, you can use those symbols and need not define them explicitly in your handler.
- As you work through the parts of this section, you should look at the skeletal device handler in Section 1.2.10. The skeletal handler illustrates the overall structure.

For examples of specific handler structure, look at the sample device handlers in Appendix A.

Also refer to Table 1-11, which illustrates the layout of a device handler .SYS file image.

1.2.1 Preamble Section

Begin the device handler source file with the preamble section. Include a .MCALL directive for the .DRDEF macro and any other macros you use that this chapter does not explicitly mention. Also in the preamble, you should define system conditionals that you will use later.

As shown in the skeletal handler, Figure 1-1, you include macros in the preamble section that build various data structures and define symbols. The following macros can be used in the preamble section:

- **.DRDEF**
Provides the primary definition of the device handler and is the only mandatory device macro. Many of the values you supply as arguments to **.DRDEF**'s parameters are equated during assembly to symbols that are then used by other handler macros.
- **.DREST**
Provides information about the handler, which is stored in block 0 of the handler's file image.
- **.DRPTR**
Points to handler service routines that can be run when the handler is loaded, unloaded, fetched, and released. Those routines do not reside in memory (keeping the memory resident portion of the handler smaller), but are read into and executed from the USR buffer.
- **.DRSPF**
Defines which special functions the handler supports.
- **.DRINS**
Points to any installation checking code and defines how the handler CSRs are to be displayed.
- **.DRSET**
Defines the handler SET commands.

As you work through this section, look at Table 1–3 to see which offsets in block 0 are written by those macros.

1.2.1.1 **.DRDEF Macro**

Use the **.DRDEF** macro near the beginning of your device handler. In the following list of functions performed by **.DRDEF**, *dd* represents the device name you specify in the macro's *name* parameter. The **.DRDEF** macro's functions are to:

- Issue **.MCALL** directives for all handler-related macros
- Provide default values for the key system conditionals
- Invoke the **.QELDF** macro to define queue element offsets
- Define bit patterns for device characteristics
- Define *dd*DSIZ as the device size in blocks
- Define *dd*\$COD as the device identification
- Set up the device status word from information in *dd*DSIZ and *dd*\$COD
- Provide default values for the device CSR in *dd*\$CSR and vector in *dd*\$VEC
- Make the symbols *dd*\$CSR and *dd*\$VEC global

- Indicate whether the handler supports extended device units
- Indicate whether the handler supports DMA (direct memory access)
- Define the required number of permanent UNIBUS mapping registers if this handler supports DMA on UNIBUS processors
- Indicate whether the handler requires serialized I/O request satisfaction

The format of the `.DRDEF` macro call is as follows:

Macro Call: `.DRDEF name,code,stat,size,csr,vec
[,UNIT64=str][,DMA=str][,PERMUMR=n][,SERIAL=str]`

<i>name</i>	is the two-letter handler name, stored in H.HAN (offset 0 of handler block 0) by <code>.DREST</code> .
<i>code</i>	is the device identifier byte, stored in H.DSTS (offset 56 of handler block 0) by <code>.DRBEG</code> .
<i>stat</i>	is the device status bit pattern, stored in H.DSTS (offset 56 of handler block 0) by <code>.DRBEG</code> .
<i>size</i>	is the device size, stored in H.DSIZ (offset 54 of handler block 0) by <code>.DRBEG</code> .
<i>csr</i>	is the default value for the device's control and status register, stored in H.ICSR (offset 176 of block 0) by <code>.DRBEG</code> . To suppress storing a value in 176, specify <code>*NO*</code> as the argument to <i>csr</i> .
<i>vec</i>	is the default value for the device's interrupt vector, stored in H1.VEC (offset 0 of block 1) by <code>.DRBEG</code> .
<i>UNIT64=str</i>	is the number of device units to be supported by this handler, stored in H.UNIT (offset 76 of handler block 0) by <code>.DRDEF</code> .
<i>DMA=str</i>	indicates whether this handler supports direct memory access, stored in symbol DV2.DM of H1.FLG (offset 10 of block 1) by <code>.DRBEG</code> .
<i>PERMUMR=n</i>	indicates this handler should be assigned <i>n</i> permanent UNIBUS mapping registers, stored in H.64UM (offset 100 of handler block 0) by <code>.DRDEF</code> .
<i>SERIAL=str</i>	indicates handler requires serialized I/O completion, stored in symbol HF2.SR of H1.FG2, (offset 16 of block 1) by <code>.DRBEG</code> .

The `.DRDEF` macro also issues the `.MCALL` directive for the following macros:

```
.DRAST  .DRBEG  .DREST  .DRFIN
.DRBOT  .DREND  .DRINS  .DRSPF
.DRSET  .DRVTB  .FORK   .QELDF
.DRTAB  .DRUSE
```

In addition, if you assemble your handler with the conditional `TIM$IT` set to 1, `.DRDEF` issues a `.MCALL` directive for the `.TIMIO` and `.CTIMIO` macros.

1.2.1.1.1 System Conditionals

RT-11 source files make extensive use of conditional assembly directives. Sections of source code are included or omitted at assembly time, based on the value of conditional symbols. For example, RT-11 uses the conditional ERL\$G to indicate whether routines for error logging should be assembled.

If you use conditional symbols in your handler, they should conform to RT-11 standard usage by setting the conditional equal to 0 to indicate that the feature it represents is not to be included and by setting the conditional to 1 to include the feature. (Note that RT-11 uses only the values 0 and 1 to indicate absence or presence of a feature.) See the *PDP-11 MACRO-11 Language Reference Manual* for information on the conditional assembly directives (.IF EQ, .IF NE, and so on).

The .DRDEF macro sets to 0 the system generation conditionals TIM\$IT (for device timeout), MMG\$T (for extended memory support), and ERL\$G (for error logging), if you do not define them in a prefix file at assembly time. In addition, if the symbols have values other than 0, .DRDEF sets them to 1.

1.2.1.1.2 Queue Element Offsets

The .DRDEF macro invokes .QELDF to define queue element offsets and define symbols for those offsets.

As shown in Table 1-1, the size of a queue element is determined by whether or not a monitor supports mapping.

Unmapped Monitors

For unmapped monitors, each queue element contains 16₈ bytes.

Mapped Monitors

Device handlers in a mapped environment require two more words of information to locate the actual user buffer in physical memory. The offsets, Q.PAR and Q.MEM, are values for PAR1 that, when combined with the user virtual buffer address (Q.BLKN), provide the physical address of the buffer.

Q.PAR and Q.MEM initially contain the same PAR1 value. The value in Q.PAR varies from Q.MEM only with UNIBUS Mapping Register (UMR) support; if the UMR handler UB is loaded, Q.PAR becomes a relocation constant to load UMRs. Q.MEM remains the PAR1 displacement bias for CPU memory management (MMU) address values. If there is no UMR support, Q.PAR and Q.MEM continue to contain the same PAR1 value. Therefore, you should use Q.MEM as the PAR1 displacement bias because it is not affected by the presence of UMR support.

Table 1-1: Queue Element Offsets

Name	Offset	Meaning
------	--------	---------

With All Monitors:

Q.LINK	0	Link to next queue element
--------	---	----------------------------

Table 1–1 (Cont.): Queue Element Offsets

Name	Offset	Meaning
Q.CSW	2	Pointer to channel status word
Q.BLKN	4	Physical block number
Q.FUNC	6	Special function code
Q.JNUM	7	Job number
Q.UNIT	7	Device unit number
Q.BUFF	^O10	User virtual buffer address
Q.WCNT	^O12	Word count
Q.COMP	^O14	Completion routine code
With Unmapped Monitors:		
Q.ELGH	^O16	Length of queue elements
	^O20– ^O24	Reserved
With Mapped Monitors:		
Q.PAR	^O16	Is initially PAR1 value. See text above
Q.MEM	^O20	Is always PAR1 value. See text above
	^O22	Reserved
Q.ELGH	^O24	Length of queue elements

Since the handler usually deals with queue element offsets relative to offset Q.BLKN, the .QELDF macro also defines the following symbolic offsets:

Symbolic Offset	From Q.BLKN
Q\$LINK	–4
Q\$CSW	–2
Q\$BLKN	0
Q\$FUNC	2
Q\$JNUM	3
Q\$UNIT	3
Q\$BUFF	4
Q\$WCNT	6
Q\$COMP	^O10

Symbolic Offset	From Q.BLKN
Q\$PAR	^O12
Q\$MEM	^O14

1.2.1.1.3 Symbol Definitions

Use direct assignment statements to define symbols that you will use later in the handler. Typically, the definitions include the device registers and other useful internal symbols. Some examples from the DY handler for mapped monitors follow:

```

; FIXED OFFSETS EQUATES (.FIXDF)
    $PNPTR  =:      000404  ;RMON OFFSET OF PNAME TABLE
    P1$EXT  =:      000432  ;RMON OFFSET OF $P1EXT ADDRESS
    $H2UB   =:      000460  ;RMON OFFSET OF UB ENTRY VECTOR PTR
    MMG$T   = 1

; EXTENDED MEMORY SUBROUTINE OFFSETS FROM $P1EXT (.P1XDF)
    $MPMEM  =:      -22.    ;OFFSET TO MAP KT-11 VIRTUAL TO PHYSICAL
    NOUMRS  = 1          ; NUMBER OF PERMANENT UMRS REQUIRED

; DY CHARACTERISTICS
    DDNBLK  = DYDSIZ*2    ;DOUBLE DENSITY SINGLE-SIDED
    DYNREG  = 3           ;# OF REGISTERS TO READ FOR ERROR LOG.
    RETRY   = 8.         ;RETRY COUNT
    SPFUNC  = 100000     ;SPECIAL FUNCTIONS FLAG
                                ; (IN COMMAND WORD)

; SPECIAL FUNCTION CODES
    SIZ$FN  = 373        ;373 - GET DEVICE SIZE
                                ;374 - UNUSED
    WDD$FN  = 375        ;375 - WRITE WITH DELETED DATA
    WRT$FN  = 376        ;376 - WRITE ABSOLUTE SECTOR
    RED$FN  = 377        ;377 - READ ABSOLUTE SECTOR
;NOTE: if you add a SPFUN code here also add it to .DRSPF

```

The .DRDEF macro also defines the following symbols for you:

```

HDERR$ = 1          ;HARD ERROR BIT IN THE CSW
EOF$   = 20000     ;END OF FILE BIT IN THE CSW

```

1.2.1.1.4 Device-Identifier Byte

The low byte of the device status word, the device-identifier byte, identifies each device in the system. You specify the correct device identifier as the *code* argument to .DRDEF. The values are defined in octal and listed under .DRDEF in the *RT-11 System Macro Library Manual*.

To create device-identifier codes for devices that are not already supported by RT-11, start by using code 377₈ for the first device, 376 for the second, and so on. This procedure should avoid conflicts with codes that RT-11 will use in the future for new hardware devices.

1.2.1.1.5 Device Status Word

The device status word identifies each unique physical device in an RT-11 system and provides other information about it, such as whether it is random or sequential access. The `.DRDEF` macro sets up symbols based on the parameter arguments for `code` and `stat`. The `.DRBEG` macro takes those symbols, builds the device status word, and stores it in block 0 of the handler file at the offset `H.DSTS` and in the `$STAT` table when the device is installed. The `.DSTATUS` programmed request can return this value to a running program.

Table 1-2 shows the meaning of the bits in the device status word. Except for `ABTIO$` and `HNDLR$`, all bits have an individual meaning. The meaning of `ABTIO$` and `HNDLR$` is determined by their combination; they should be thought of as a pair. More information on the `ABTIO$/HNDLR$` pair is found in Sections 1.3.1 and 1.3.2.

Table 1-2: Device Status Word

Bit	Symbol	Meaning
0-7	—	Device-identifier byte (see Section 1.2.1.1.4)
8	<code>VARSZ\$</code>	0 = <code>SF.SIZ</code> (special function code 373) requests are invalid for this handler 1 = <code>SF.SIZ</code> (code 373) requests (return volume size) are valid for this handler
9	<code>ABTIO\$†</code>	0 = Handler is not entered at abort entry point on normal program exits 1 = Handler is entered at abort entry point whenever a program terminates
10	<code>SPFUN\$</code>	0 = <code>.SPFUN</code> requests are invalid 1 = Handler accepts <code>.SPFUN</code> requests
11	<code>HNDLR\$‡</code>	0 = Enter handler at abort entry point only if there is an active queue element belonging to the aborted job 1 = Enter handler at abort entry point on all aborts
12	<code>SPECL\$</code>	1 = Special directory-structured device (examples are <code>MS</code> and <code>MU</code>)
13	<code>WONLY\$</code>	1 = This is a write-only device
14	<code>RONLY\$</code>	1 = This is a read-only device
15	<code>FILST\$</code>	0 = This is a sequential-access device (examples are <code>LP</code> , <code>LS</code> , <code>MS</code>) 1 = This is a random-access device (examples are <code>DU</code> and <code>DY</code>)

†`ABTIO$` works in combination with `HNDLR$`. See Section 1.3.1.

‡`HNDLR$` works in combination with `ABTIO$`. See Section 1.3.1.

The bit combinations for handlers that internally queue I/O requests are described in Section 1.4. See Section 1.9 for details on special devices (such as magtape).

All device handlers that have bit 15 set are assumed to be RT-11 file-structured devices by most of the system utility programs.

An easy way to define the device status word is to use the symbols for the bit patterns that `.DRDEF` defines for you. Thus, you can create the `stat` argument by ORing together the appropriate symbols from the list below.

```
FILST$ == 100000 ;File-structured random access
RONLY$ == 40000 ;Read-only
WONLY$ == 20000 ;Write-only
SPECL$ == 10000 ;Special directory structured device
HNDLR$ == 4000 ;Enter handler on abort
SPFUN$ == 2000 ;Accepts special functions
ABTIO$ == 1000 ;Always take abort entry
VARSZ$ == 400 ;Handler supports variable-size volumes
```

For example, form the `stat` argument for the DY, MS, and LS handlers as follows:

- For DY: `FILST$!SPFUN$!VARSZ$`
- For MS: `SPECL$!SPFUN$`
- For LS: `WONLY$!SPECL$`

1.2.1.1.6 Device Size Word

The `size` argument for the `.DRDEF` macro defines `ddDSIZ` to be the size of the device in 256-word blocks. The `.DRDEF` macro stores the value of `ddDSIZ` in `H.DSIZ`, offset 54 in the handler's block 0.

The `.DSTAT` programmed request returns the value of the device size word to a running program. For examples of the `.DRDEF` macro, see the device handler listings in Appendix A.

1.2.1.2 .DREST Macro

The `.DREST` macro places device specific information about the handler into handler block 0:

- The device class and any variation
- The presence of bad-block replacement information
- How the handler can be installed, loaded, and mounted

The format of the `.DREST` macro call is as follows:

Macro Call: `.DREST [CLASS=str][,MOD=str][,DATA=dptr]
[,TYPE=str][,REPLACE=rptr][,STAT2=symb]`

`CLASS=str` stores the class symbol in `H.CLAS`, offset 20 in handler block 0.

`MOD=str` stores the classification modifier in `H.MOD`, offset 21 in handler block 0.

`DATA=dptr` stores an internal table file address in `H.DATA`, offset 72 in handler block 0.

<i>TYPE=str</i>	stores an internal table device classification in H.TYPE, offset 70 in handler block 0.
<i>REPLACE=rptr</i>	stores a pointer to a bad-block replacement table in H.REPL, offset 32 in handler block 0.
<i>STAT2=symb</i>	stores a second status word in H.STS2, offset 36 in handler block 0.

See Section 1.2.1.9 for more information on the contents of handler block 0, including those offsets written by .DREST. For information on using the .DREST macro, see the *RT-11 System Macro Library Manual*.

1.2.1.3 .DRINS Macro

The .DRINS macro sets up the installation code area in the handler's block 0:

- Defines the display CSR addresses (displayed by RESORC)
- Defines the installation CSR addresses (used by INSTALL command) and monitor bootstrap
- Defines system device (INSSYS) and data device (INSDAT) installation entry points

INSSYS is located at symbol H.ISY, offset 202. INSDAT is located at symbol H.IDK, offset 200.

The format of the .DRINS macro call is as follows:

Macro Call: `.DRINS name,<csr,csr,...>`

<i>name</i>	is the device handler name. If <i>name</i> is preceded by a minus sign (-), it indicates that the specified CSR is for display purposes only; there is no installation CSR for this invocation of .DRINS.
<i>csr</i>	creates a symbolic reference to a CSR for this device. The first (or only) specified is both the installation CSR and the first display CSR. The .DRBEG macro stores the installation CSR in H.ICSR, offset 176 in block 0. The .DRINS macro stores the first display CSR in H.DCSR, offset 174 in handler block 0. (You must also specify <i>csr</i> = *NO* in .DRDER for this to take effect.) If more than one CSR is specified, the second and any subsequent in the list are the secondary (and subsequent) display CSRs. Those are written to offset 172, 170, and so forth. The list is terminated with a word containing a zero value. (There remains a single installation CSR.)

See Section 1.2.1.9 for more information on the contents of handler block 0, including those offsets written by .DRINS. For information on using the .DRINS macro, see the *RT-11 System Macro Library Manual*.

1.2.1.4 .DRPTR Macro

The .DRPTR macro sets up pointers to handler service routines that can assist the handler when it is fetched, loaded, released, or unloaded.

The pointers are located in handler block 0. The service routines are not normally located in handler block 0 and are not located in the handler memory image. When called, any service routine is read from the handler file image into the shared area of the USR and used by the handler.

The format of the .DRPTR macro call is as follows:

Macro Call: .DRPTR [FETCH=*n*][,RELEASE=*n*][,LOAD=*n*][,UNLOAD=*n*]

FETCH=*n* stores a pointer to a fetch service routine in H.FETC, offset 2 in handler block 0.

RELEASE=*n* stores a pointer to a release service routine in H.RELE, offset 4 in handler block 0.

LOAD=*n* stores a pointer to a load service routine in H.LOAD, offset 6 in handler block 0.

UNLOAD=*n* stores a pointer to an unload service routine in H.UNLO, offset 10 in handler block 0.

See Section 1.2.1.9 for more information on the contents of handler block 0, including those offsets written by .DRPTR. For information on using the .DRPTR macro, see the *RT-11 System Macro Library Manual*.

1.2.1.5 .DRSPF Macro

The .DRSPF macro defines a handler's support for special functions. As explained in the *RT-11 System Macro Library Manual*, two methods can be used to create that support.

The format of the .DRSPF macro call is as follows:

Macro Call: .DRSPF *arg*[,*arg2*][,TYPE=*n*]

Up to three groups of special functions can be described in symbols H.SPF1, H.SPF2, and H.SPF3, beginning at offset 22 in handler block 0. Any further groups require the *extension table method*, which are stores in the pointer symbol H.SPFX at offset 30 in handler block 0. The offset H.SPFX points to that extension table of other supported special functions.

See Section 1.2.1.9 for more information on the contents of handler block 0, including those offsets written by .DRSPF. For information on using the .DRSPF macro, see the *RT-11 System Macro Library Manual*.

1.2.1.6 .DRTAB Macro

The .DRTAB macro is normally reserved for use by Digital. Although .DRTAB is described in the *RT-11 System Macro Library Manual*, you should use .DRUSE in your handler.

1.2.1.7 .DRUSE Macro

The .DRUSE macro defines a list of data tables for the device handler. There are three levels of definition.

1. You write a data table (or tables) at some file address (or addresses) in your device handler. You invoke .DRUSE enough times to define each data table. To invoke .DRUSE, see the *RT-11 System Macro Library Manual*.
2. At a file address, the .DRUSE macro creates a descriptor table of those data tables. The descriptor table is described in Section 1.2.8.7.
3. The .DRUSE macro places a pointer to the descriptor table file address in H.USER, at offset 106 in the handler's block 0.

The format of the .DRUSE macro call is as follows:

Macro Call: .DRUSE *type*,*addr*,*size*

type stores the value of *type* at symbol DT.ID in the descriptor table
addr stores the value of *addr* as symbol DT.PTR in the descriptor table
size stores the value of *size* as symbol DT.SIZ in the descriptor table

1.2.1.8 .DRSET Macro

The .DRSET macro must be invoked in the preamble section of the device handler. Invoking .DRSET and the structure of the SET tables it creates are described in Section 1.5.2.

1.2.1.9 Information in File Image Block 0

Table 1-3 describes the contents of block 0 of the assembled handler file image. This is the informational block and is not normally loaded into memory.

The symbol names in the table are those used in the distributed system definition library file, SYSTEM.MLB. The macros are those that actually write the offset; they are not necessarily the originating macro. Where appropriate, the description indicates where you can find more information about the offset, its contents, or the structure pointed to by an address in the offset.

Table 1-3: Contents of .SYS Image Block 0

Offset	Symbol	Macro	Description
000000	H.HAN	.DREST	Handler identifier in RAD50
	H.HANV		Value for H.HAN (RAD50 HAN)
000002	H.FETC	.DRPTR	Pointer to a FETCH service routine; See Section 1.2.8.1
000004	H.RELE	.DRPTR	Pointer to a RELEASE service routine; See Section 1.2.8.1
000006	H.LOAD	.DRPTR	Pointer to a LOAD service routine; See Section 1.2.8.1

Table 1–3 (Cont.): Contents of .SYS Image Block 0

Offset	Symbol	Macro	Description
000010	H.UNLO	.DRPTR	Pointer to an UNLOAD service routine; See Section 1.2.8.1
000012– 000016			Reserved
000020	H.CLAS	.DREST	Device classification; See Section 1.2.1.2
000021	H.MOD	.DREST	Device classification modifier; See Section 1.2.1.2
000022	H.SPF1	.DRSPF	First special function (index method) list; See Section 1.2.8.2
000024	H.SPF2	.DRSPF	Second special function (index method) list; See Section 1.2.8.2
000026	H.SPF3	.DRSPF	Third special function (index method) list; See Section 1.2.8.2
000030	H.SPFX	.DRSPF	Pointer to further special functions (extension table method); See Section 1.2.8.2
000032	H.REPL	.DREST	Pointer to bad-block replacement table; See Section 1.2.8.3
000034			Reserved
000036	H.STS2	.DREST	Second status word; See Section 1.2.8.5
000040– 000050			SYSCOM area for runnable handlers.
000052	H.SIZ	.DRBEG	Handler size (<i>ddEND–ddSTRT</i>)
000054	H.DSIZ	.DRBEG	Device size (<i>ddDSIZ</i>); See Section 1.2.1.1.6
000056	H.DSTS	.DRBEG	Device status word (<i>ddSTS</i>); See Section 1.2.1.1.5
000060	H.GEN	.DREND	Result of standard SYSGEN conditionals OR'd with the value of the FORCE= parameter; See Section 1.2.8.6
000061			Reserved
000062	H.BPTR	.DRBOT	Pointer to the primary bootstrap; See Section 1.11.2.2
000064	H.BLEN	.DRBOT	Bootstrap size in bytes; See Section 1.11.2.1
000066	H.READ	.DRBOT	Pointer to the bootstrap read routine; See Section 1.11.2.5
000070	H.TYPE	.DRTAB .DREST	If contains value –1, indicates written by .DRTAB (only Digital distributed handlers)—otherwise: If contains a RAD50 value, indicates invoked by .DREST and is the device type classification for an internal table

Table 1–3 (Cont.): Contents of .SYS Image Block 0

Offset	Symbol	Macro	Description
000072	H.DATA	.DRTAB .DREST	If H.TYPE written by .DRTAB, then H.DATA is a pointer to the list of handler data table descriptors. If H.TYPE written by .DREST, then H.DATA is the file address of the internal data tables. See Section 1.2.8.7
000074	H.DLEN	.DRTAB .DREST	Size in bytes of total list of handler data table descriptors; See Section 1.2.8.7
000076	H.UNIT	.DRDEF	Pointer to extended device-unit ownership table
000100	H.64UM	.DRDEF	Letter name of extended device-unit handler and device characteristics for UMR support; See Section 1.2.8.8
000102– 000104			Reserved
000106	H.USER	.DRUSE	Pointer to the file address of the handler data descriptor table; See Section 1.2.8.7
000110– 000173		.AUDIT .MODULE	Information written by those two macros. Terminated by –1. This list and the display CSR list cannot overlap.
000164– 000174	H.DCSR	.DRINS	Display CSRs read by RESORC. If more than one, each written into previous offset; See Section 1.2.1.3
000176	H.ICSR	.DRDEF .DRINS	Installation CSR; See Sections 1.2.1.1 and 1.2.1.3.
000200	H.IDK	.DRINS	Data device installation entry point (INSDAT); See Section 1.2.1.3
000202	H.ISY	.DRINS	System device installation entry point (INSSYS); See Section 1.2.1.3
000204– 000377			Installation code; See Section 1.14.3.5
000400– 000777	H.SET	.DRSET	SET code; See Section 1.5.2

1.2.2 Header Section

The second part of an RT–11 device handler is the header section. The header section is the beginning of the memory resident portion of the handler and starts at the base of file image block 1. In the header section, you invoke the .DRBEG macro to build a data structure of variable size at the beginning of the handler’s memory image. This macro also stores information in the handler file at offsets 52 through 60 of block 0, and creates some global symbols.

The data you set up in the header section is used when the handler is brought into memory with the .FETCH programmed request or LOAD monitor command. The

contents of location 176, described below, are used by the bootstrap when it checks for the presence of device hardware at handler installation time.

As shown in the skeletal handler, Figure 1–1, you include macros in the preamble section that build various data structures and define symbols. The following macros can be used in the header section:

- **.DRBEG**
Defines the handler queue entry point and provides other information about the handler. Writes locations in the handler file image blocks 0 and 1.
- **.DRVTB**
Defines multiple vectors if the handler supports more than one interrupt vector.

1.2.2.1 .DRBEG Macro

The **.DRBEG** macro sets up offsets in block 0 and the header information in block 1. This macro also generates the appropriate global symbols for your handler. Before you invoke **.DRBEG**, invoke **.DRDEF** to define various symbols that **.DRBEG** uses internally. The format for **.DRBEG** is as follows:

.DRBEG name[,SPFUN=spsym][,NSPFUN=nspsym]

name is the two-character device name.

spsym is the label on the list of DMA standard special functions. Sets HF2.SD in offset H1.FG2 of handler block 1.

nspsym is the label on list of DMA nonstandard special functions. Sets HF2.ND in offset H1.FG2 of handler block 1.

For examples of **.DRBEG**, see the DL handler listing in Appendix A and the UB example in Chapter 2.

1.2.2.2 Multivector Handlers: .DRVTB Macro

An RT–11 device handler can service multiple controllers where each controller has an interrupt vector. The handler can also service a device that has more than one vector.

Device handlers support a single vector through the **.DRDEF** macro's *vec* parameter. A device handler that supports multiple vectors must contain the **.DRVTB** macro. Invoke the **.DRVTB** macro once for each vector. Each invocation creates a table with three entries. The table for each vector consists of the vector location, the interrupt entry point, and the Processor Status, or PS, value.

You can invoke **.DRVTB** anywhere between the **.DRBEG** macro and the **.DREND** (or **.DRBOT**) macro, as long as it does not interfere with the flow of control within the handler. You must invoke this macro once for each vector, and the macro calls must appear one after the other in the handler.

The format of the **.DRVTB** macro is as follows:

.DRVTB name,vec,int[,ps]

name is the two-character device name. Specify it on the first .DRVTB call; leave this argument blank on all subsequent calls.

vec is the location of the vector; it must be between 0 and 474. The first vector is usually *dd\$VEC*. The value must be a multiple of 4. The .DRBEG stores the value for *dd\$VEC* in H1.VEC, offset 0 of block 1.

int is the symbolic name of the interrupt handling routine; it must appear elsewhere in the handler. It generally takes the form *ddINT*, where *dd* represents the two-character device name. The .DRBEG stores the value for *ddINT* in H1.ABT, offset 2 of block 1.

ps is an optional value you can use to specify the low-order four bits of the new Processor Status word in the interrupt vector. If you omit this argument, it defaults to 0.

An example of a handler that can use two vectors is the DY handler, when that handler is built to support a second controller. The following example shows the source lines and the code the macros generate:

```
.IF NE DYT$O ; If we support two controllers:
.DRVTB DY,DY$VEC,DYINT ; DY$VEC symbol for first vector table
.DRVTB ,DY$VC2,DYINT ; DY$VC2 symbol for second vector table
.ASSUME . LE DYSTRT+1000
.ENDC ;NE DYT$O
```

Generates:

```
.IF NE DYT$O
.DRVTB DY,DY$VEC,DYINT
.WORD DY$VEC&^C3.,DYINT-.,^o340!0,^o100000
.DRVTB ,DY$VC2,DYINT
.WORD DY$VC2&^C3.,DYINT-.,^o340!0,^o100000
.ASSUME . LE DYSTRT+1000
.ENDC ;NE DYT$O
```

In the example above, the priority bits of the PS are always set to PR7, even if you omit the *ps* argument.

PS Condition Codes

In the .DRVTB macro, only the condition code bits of the *ps* argument are significant. These can be useful if you have a common interrupt service entry point for two or more vectors and you need to determine through which vector the interrupt occurred. For example, the skeletal handler (Figure 1–1) has a single interrupt entry point for its two vectors. For the handler to determine the source of the interrupt, one is serviced with the carry bit clear and the other (*INT2*), when the carry bit is set.

1.2.2.3 Information in File Image Block 1

The following table describes the contents of block 1 of the assembled handler file image that are written by the `.DRBEG` macro. This is the first block that is normally loaded into memory and is therefore block 0 of the handler memory image.

The symbol names used in the table are from the distributed system definition library file, `SYSTEM.MLB`. All defined offsets are written by `.DRBEG` but `.DRBEG` is not the originating macro for all locations. As appropriate, the description indicates where you can find more information about each offset, its contents, or the structure pointed to by an address in the offset.

Table 1–4: Contents of `.SYS` Image Block 1

Offset	Symbol	Macro	Description
001000	H1.VEC	<code>.DRBEG</code>	Either the device vector if a single vector device or an offset to the table of vectors for multivector devices (<i>ddSTRT</i>)
001002	H1.ABT	<code>.DRBEG</code>	Offset to the interrupt service entry point
001004	H1.HLD	<code>.DRBEG</code>	Priority (340)
001006	H1.LQE	<code>.DRBEG</code>	Pointer to the last queue element (<i>ddLQE</i>)
001010	H1.CQE	<code>.DRBEG</code>	Pointer to the current queue element (<i>ddCQE</i>) in handler memory image
001010	H1.FLG	<code>.DRBEG</code>	Flag word (in handler file image); See Section 1.2.9.1
001012	H1.NOP	<code>.DRBEG</code>	NOP instruction OR'd with flags; See Section 1.2.9.2
001014	H1.BR	<code>.DRBEG</code>	Branch instruction (optional)
001016	H1.FG2	<code>.DRBEG</code>	Second flag word (optional); See Section 1.2.9.3
001020	H1.SCK	<code>.DRBEG</code>	Pointer to SPFUN address check routine (optional)
001022	H1.SDF	<code>.DRBEG</code>	Pointer to standard DMA SPFUN table (optional)
001024	H1.LDT	<code>.DRBEG</code>	Pointer to LD translation table (optional)
001026	H1.NDF	<code>.DRBEG</code>	Pointer to nonstandard DMA SPFUN table (optional)

1.2.3 I/O Initiation Section

The I/O initiation section contains the first executable instructions of the handler and must follow the call to `.DRBEG`. The purpose of the code in this section is to start a data transfer. Remember that you must write Position-Independent Code (PIC) for the handler.

When a program issues a programmed request that requires device I/O, such as `.READ` or `.WRITE`, control first passes to the Resident Monitor, which then calls the device handler for the peripheral device with the `CALL` instruction. The monitor calls the handler at the handler's sixth word—that is, the first word immediately after the five-word data header. The monitor makes the call whenever a new queue element becomes the first element in a handler's queue. This situation occurs when

an element is added to an empty queue, or when an element becomes first in a queue because a prior element was released. If any parameters in the I/O request are invalid for the device (for example, the block number is too large, the unit number is too high, and so on), the handler should proceed immediately to the I/O completion section and signal a hard (fatal) error.

The I/O initiation code executes at processor priority 0 in system state, which means that no context switch can occur, no completion routines can run, and any traps to 4 and 10 cause a system fatal halt. All registers are available for you to use in this section. The fifth word of the handler header, *ddCQE*, contains a pointer to the current queue element at its third word, *Q.BLKN*.

The queued I/O system guarantees that requests for data transfers are serialized so that RT-11 device handlers need not be re-entrant. Therefore, you can minimize the size of a handler by mixing, rather than separating, the pure code and the data segments.

1.2.3.1 Guidelines for Starting the Data Transfer

Since the purpose of the I/O initiation section is to start up the data transfer, you must now supply the instructions to do this. The following steps (from the RK handler) represent guidelines for a generalized I/O initiation section:

1. You should have already decided how many times the handler will retry a transfer should an error occur. Initialize a retry counter by moving the maximum number of retries to it. The following two lines of code illustrate this step:

```

MOV      #RKCNT, (PC)+          ;RKCNT = MAXIMUM # OF RETRIES
RETRY:   .WORD    0             ;THE RETRY COUNTER

```

2. Put the pointer to the current queue element into a register, and get the device unit number and the block number for the transfer from the queue element. The following lines of code illustrate this.

```

MOV      RKCQE, R5              ;GET CURRENT QUEUE ELEMENT POINTER
MOV      @R5, R2                ;PICK UP BLOCK NUMBER
MOV      Q$UNIT-1(R5), R4       ;GET REQUESTED UNIT NUMBER
ASR      R4                     ;SHIFT UNIT NUMBER
ASR      R4                     ; TO HIGH 3 BITS
ASR      R4                     ; OF LOW BYTE
SWAB     R4                     ;PUT UNIT NUMBER IN HIGH 3 BITS
BIC      #^C<DAUNIT>, R4       ;ISOLATE UNIT IN DRIVE SELECT BITS

```

3. Next, perform the steps to calculate the address on the device for the data transfer to begin. The instructions you use depend on the device's structure, of course. Once you have calculated the correct address, save it in a memory location. If you need to retry this transfer, you will not have to recalculate the address.

```

.
.
.
MOV      R3, (PC)+              ;SAVE ADDRESS IN DISKAD
DISKAD:  .WORD    0             ;SAVE CALCULATED ADDRESS HERE

```

4. Steps 1 through 3 outlined above are executed only once for each data I/O request from a running program. However, in case of a soft error, you may need to restart a transfer as part of the retry operation. So, by placing a label here to use as the retry entry point, you avoid repeating steps 1 through 3.

The following steps can be performed more than once. They are executed once for the first I/O startup, and they can be executed again if an I/O error causes a retry.

At this point, the handler should determine whether the I/O request is a read, a write, or a seek. It should then generate the appropriate op code for the operation and move it to the device control and status register. This step actually initiates the I/O transfer.

```

        CSIE =          100          ;INTERRUPT ENABLE
        FNWRITE =       12          ;WRITE
        CSGO =          1          ;GO BIT
        .
        .
        .
AGAIN:  MOV      RKCQE,R5          ;POINT TO QUEUE ELEMENT
        MOV      #CSIE!FNWRITE!CSGO,R3 ;ASSUME A WRITE
        MOV      #RKDA,R4          ;POINT TO DISK
        .                          ;ADDRESS REGISTER
        .
        .

```

5. Finally, return to the interrupted program by going through the monitor first. Then when the I/O transfer finishes, the device will interrupt, and control will pass to the handler at the interrupt entry point in the interrupt service section of the handler.

```

        RTS      PC              ;AWAIT INTERRUPT

```

1.2.3.2 Transferring the Data

Data can be transferred between a device and the user buffer as individual bytes, words, or by direct memory access (DMA). How the data is transferred is largely determined by whether or you are using a mapped or unmapped monitor. This section describes transferring the three types of data into both unmapped and mapped memory.

1.2.3.2.1 Byte Transfer from the User Buffer to the Device

The following examples are from the XL handler and illustrate transferring a byte from the user buffer.

Unmapped Monitor

```
GNXTCH: MOV      XOCQE,R4          ;R4->current output queue element
          BEQ      10$             ;None available...
          ADD      #Q$WCNT,R4      ;R4->word count
          TST      @R4             ;Any characters left to output?
          BEQ      20$             ;Nope, this request is complete
          INC      @R4             ;Yes, now there is one less to do
          MOVB     @-(R4),R5        ;Get the byte to output
          INC      @R4             ;bump pointer to next byte
```

Mapped Monitor

RT-11 provides the \$GTBYT routine to perform the address translation between a user buffer in mapped memory and the device. The \$GTBYT routine is described in more detail in Section 1.10.4.1.

Before the call:

R4 must point to *Q.BLKN*, the third word in the queue element.

After the call:

(*SP*), the first word on the stack, contains the next byte from the user buffer in the low byte. The contents of the high byte are not defined.

R4 is unchanged.

```
GNXTCH: MOV      XOCQE,R4          ;R4->current output queue element
          BEQ      10$             ;None available...
          TST      Q$WCNT(R4)      ;Any characters left to output?
          BEQ      20$             ;Nope, this request is complete
          INC      Q$WCNT(R4)      ;Yes, now there is one less to do
          CALL     @$GTBYT         ;Get the byte to output
          MOV      (SP)+,R5
```

The buffer address (*Q.BUFF*) in the queue element is updated by 1. If *Q.BUFF* is greater than 20077, a 1 is added to *Q.PAR* and *Q.MEM* and *Q.BUFF* is reduced by 100.

1.2.3.2.2 Byte Transfer from the Device to the User Buffer

The following examples are from the XL handler and illustrate transferring a byte into the user buffer.

Unmapped Monitor

```
30$:      ADD      #Q$WCNT,R4      ;R4->Word count
          MOVB     R5,@-(R4)        ;Return the character
          INC      (R4)+           ;Bump the buffer pointer
          DEC      (R4)            ;Is transfer complete?
          ; (z-bit=1 if so)
```

Mapped Monitor

RT-11 provides the \$PTBYT routine to perform the address translation between a user buffer in mapped memory and the device. The \$PTBYT routine is described in Section 1.10.4.2.

Before the call:

R4 must point to *Q.BLKN*, the third word in the queue element.

The byte to transfer to the user buffer must be on the top of the stack. The character must be in the low byte of the stack's first word. The high byte is unpredictable.

After the call:

The word containing the character to transfer is removed from the stack and transferred to the user buffer.

R4 is unchanged.

```
30$:      MOVB      R5,-(SP)           ;Put character here for PUTBYT
          CALL     @$PTBYT         ;Call the routine
          DEC      Q$WCNT(R4)      ;Is transfer complete?
                                   ; (z-bit=1 if so)
```

The buffer address (*Q.BUFF*) in the queue element is updated by 1. If *Q.BUFF* is greater than 20077, a 1 is added to *Q.PAR* and *Q.MEM* and *Q.BUFF* is reduced by 100.

1.2.3.2.3 Word Transfer from the Device to the User Buffer

The handler may have to change a word in user memory. The following examples are taken from the *DY* handler and return a word of size information.

Unmapped Monitor

```
; DRIVER IS DUAL DENSITY ONLY

          BIS      #CSDN,R4        ;ALWAYS USE DOUBLE DENSITY
          CMPB     R1,#SIZ$FN     ;SPECIAL SIZE FUNCTION?
          BNE      3$             ;NO, CONTINUE
          MOV      #DDNBLK,@(R5)+ ;RETURN DOUBLE DENSITY SIZE
          JMP      DYDONE         ;DONE WITH SIZE OPERATION
```

Mapped Monitor

RT-11 provides the routine *\$PTWRD* to perform the address translation between the device and a user buffer. The *\$PTWRD* routine is described in Section 1.10.5.

```
; DRIVER IS DUAL DENSITY ONLY

          BIS      #CSDN,R4        ;ALWAYS USE DOUBLE DENSITY
          CMPB     R1,#SIZ$FN     ;SPECIAL SIZE FUNCTION?
          BNE      3$             ;NO, CONTINUE
          MOV      #DDNBLK,-(SP)  ;RETURN DOUBLE DENSITY SIZE
          MOV      DYCQE,R4       ;CURRENT QUEUE ELEMENT
          CALL     @$PTWRD        ;STORE SIZE IN BUFFER
          JMP      DYDONE         ;DONE WITH SIZE OPERATION
```

The buffer address (*Q.BUFF*) in the queue element is updated by 2. If *Q.BUFF* is greater than 20077, a 1 is added to *Q.PAR* and *Q.MEM* and *Q.BUFF* is reduced by 100.

1.2.3.2.4 Non-DMA Transfers

The following examples are from the DY handler and illustrate getting a pointer to the user buffer for use in DMA transfer initialization.

Unmapped Monitor

```
3$:
      MOV      (R5)+,R0      ;GET THE USER'S BUFFER ADDRESS
      MOV      @R5,WRDCNT   ;GET WORD COUNT
      BPL     4$            ;POSITIVE MEANS READ, SO ALL SET UP
```

Mapped Monitor

RT-11 provides the \$MPMEM routine to perform address translation for non-DMA transfers between the device and a user buffer. Non-DMA transfers are typically done with the MOV instruction. The \$MPMEM routine is described in Section 1.10.3.1.

```
3$:
      CALL    @$MPPTR      ;CONVERT MAPPED ADDRESS TO PHYSICAL ADDRESS
      MOV     (SP)+,R0     ;GET PHYSICAL BUFFER ADDRESS LOW ORDER BITS
      MOV     R4,(PC)+    ;SAVE CURRENT COMMAND WORD
35$:
      .BLKW
      MOV     (SP)+,R4     ;GET HIGH-ORDER ADDRESS BITS <21:18>
      BIT    #1700,R4     ;22-BIT ADDRESS SPECIFIED?
      BNE    DYERR        ;YES, NOT VALID FOR THIS CONTROLLER
      SWAB   R4           ;MOVE TO CORRESPONDING POSITIONS IN HIGH BYTE
      BIS    35$,R4       ;NOW MERGE COMMAND WORD WITH EXTENSION BITS

      MOV     @R5,WRDCNT  ;GET WORD COUNT
      BPL    4$          ;POSITIVE MEANS READ, SO ALL SET UP
```

1.2.3.2.5 DMA Transfers

The address translation for DMA transfers is performed by the \$MPPHY routine, described in Section 1.10.3.2. A complete description of doing DMA transfers using UNIBUS mapping registers (UMRs) is in Section 2.13.

1.2.4 Interrupt Service Section

Control passes to the interrupt service section of the handler when a device interrupts, when the program requesting the I/O transfer aborts, or a .ABORT is issued for the channel. The code in this section must first determine if the data transfer had an error, if it was incomplete, or if it was complete, and then take the appropriate action. The same register usage restrictions that apply to the interrupt entry point also apply to the abort entry point. See Chapter 5 in the *RT-11 System Internals Manual* for information on interrupt service routines.

Your first step in coding the interrupt service section is to set up the interrupt entry point and the abort entry point by using the .DRAST macro. (These entry points are sometimes referred to as the asynchronous trap entry points.) The default name for the interrupt entry point is *ddINT*, where *dd* is the device name. Under normal conditions, the handler is called at the interrupt entry point when an interrupt occurs. However, under some circumstances, the handler is called at the abort entry point located at *ddINT-2*. The various situations are discussed in the following sections.

1.2.4.1 .DRAST Macro

Use the .DRAST macro to set up the interrupt entry point and the abort entry point, and to lower the processor priority. The .DRDEF and .DRVTB macros fill in the structure at bootstrap (for the system device) or at .FETCH time (for a data device).

The format of the .DRAST macro is as follows:

.DRAST name,pri[,abo]

name is the two-character device name.

pri is the priority of the device, and the priority at which the interrupt service code is to execute.

abo is an optional argument that represents the label of the abort entry point. If you omit this argument, the macro expansion generates a RETURN instruction at the abort entry point. Either the branch to the specified label or the RETURN instruction is the word immediately preceding the interrupt entry point *ddINT*.

The following example from the DY handler shows the .DRAST macro call. In the example, DYABRT is the label for the abort routine which would generate the instruction BR DYABRT in the word preceding the interrupt entry point DYINT.

```
.SBTTL  INTERRUPT ENTRY POINT
.DRAST  DY,5,DYABRT           ; AST entry point
BR      DYABRT               ; Jump to abort entry point
DYINT:: JSR    R5,@$INPTR     ; Jump to monitor INTEN code
        .WORD  ^C<5*^o40>&^o340 ; New priority
        .FORK  DYFBLK        ; Request fork level immediately
        JSR    R5,$FKPTR     ; Jump to monitor fork code
        .WORD  DYFBLK-      ; Offset to fork queue element
        CALL  SETDY         ; Setup registers
        BMI   DYERR2        ; Check out the error and retry
INTDSP: JMP    @(PC)+        ; No error, return to called
INTRTN: .WORD  0             ; : Address of waiting routine
```

The next example, from the RK handler, does not have an abort routine. Notice the instruction, RETURN, in the word immediately preceding the interrupt entry point RKINT.

```
.DRAST  RK,5
.GLOBL  $INPTR               ;MAKE THIS SYMBOL GLOBAL
RETURN  ;JUST RETURN ON ABORT
RKINT:: JSR    R5,$INPTR     ;JUMP TO MONITOR INTEN CODE
        .WORD  ^C<5*^o40>&^o340 ;NEW PRIORITY
```

1.2.4.2 Abort Entry Point

As described in Section 1.3, there are a number of situations that cause an abort in the queued I/O system. The response to the abort situation by the handler and RMON depends on the ABTIO\$ and HNDLR\$ bits in the device status word.

When an abort occurs, it is important to stop I/O on some devices. Character-oriented devices, such as the communications handler XL, fall into this category.

So, character-oriented devices generally contain an abort routine; the abort entry point is simply a branch instruction to that routine. The following lines are from the XL handler:

```
XLDONE:
.
.
.
    BIC      #RC.IE,@XIS          ;Turn off input interrupts
.
.
.
    RTS     PC                    ;Return to monitor
```

Other devices, such as disks, should be allowed to complete an I/O transfer attempt, even if an abort occurs. In fact, trying to abort in the middle of an operation can corrupt data or formatting information on a disk. So, instead of having a separate abort routine, most handlers for disks ignore an abort. Thus, a RETURN instruction is located at the abort entry point, which simply returns control to the monitor.

The abort entry point is always located at the word previous to the interrupt entry point (*ddINT-2*). If the optional *.DRAST abo* parameter is specified, the abort entry point is a branch instruction to the label specified as the *abo* parameter argument. If *abo* is not specified, the *.DRAST* macro expansion places a RETURN instruction at the abort entry point (*ddINT-2*).

If you use *.FORK* in your handler, there is a special procedure you must follow if an abort occurs. You must move 0 to *F.BADR* (the fork routine address, at offset 2) in the fork block. This prevents the monitor from attempting to execute a meaningless fork routine after the abort.

1.2.4.3 Lowering the Priority to Device Priority

When the interrupt occurs, the handler is entered at priority 7. As with interrupt service routines, the handler's first task is to lower the processor priority to the priority of the device, thus permitting more important devices to interrupt this service routine. Instead of using the *.INTEN* call, as in an interrupt service routine, use the *.DRAST* macro to lower the priority.

1.2.4.4 Guidelines for Coding the Interrupt Service Section

Since the purpose of this section is to evaluate the results of the last device activity, you must now supply the instructions to do this. Essentially, the code must determine if the transfer was in error, if it was incomplete, or if it was complete.

1. If an Error Occurred

If an error occurred during the transfer, the handler must distinguish between a hard error and a soft error that might vanish if the operation is retried.

If the error is hard, the handler should immediately exit through the I/O completion section after setting *HDERR\$* in the CSW.

If the error is soft, the handler should prepare to retry the transfer. It should decrement the count of available retries. Then, possibly at fork level, it should branch back to the I/O initiation section to restart the transfer. If the transfer has already been retried enough times (the retry count is 0), treat the failure as though it were a hard error. In that case, the handler should proceed to the I/O completion section after setting HDERR\$ in the CSW.

Note that dropping to fork level is not strictly required to process an error. Whether or not to use .FORK depends on the length of time required for setting up the retry. The .FORK call is especially useful because it gives you use of R0 through R3, thus permitting you to use common routines for the retry. If you do not use .FORK, only R4 and R5 are available.

2. Perform Retries at Fork Level

As also described in the *RT-11 System Internals Manual*, the .FORK macro causes a return to the Resident Monitor, which dismisses the current interrupt. The code that follows .FORK executes at priority 0, rather than at device priority, after all other interrupts have been serviced, but before any jobs or their completion routines can execute. The code following .FORK executes, as does the main body of the interrupt service section of the handler, in system state. (This is the same state the I/O initiation section runs in.) Thus, context switching is prevented while the fork level code is executing, and any traps to 4 and 10 cause a system fatal halt.

The following example from the RK handler illustrates how the handler drops priority to fork level to retry data transfers after a soft error occurred. Fork level is ideal for performing the retries, since this may be a lengthy process. The .FORK call and its expansion are as follows:

```

.FORK    RKFBLK                ;THE FORK CALL
JSR     R5,@$FKPTR            ;(JUMP TO MONITOR FORK CODE)
.WORD   RKFBLK - .           ;(OFFSET TO FORK QUEUE ELEMENT)
RKRETR: CLRB    RETRY+1       ;RESET A FLAG
        BR      AGAIN        ;BRANCH INTO I/O INIT SECTION

```

3. If the Transfer Was Incomplete

In general, a transfer is considered to be incomplete when there are more characters or more blocks of data left to transfer. The handler should restart the device and exit with a RETURN instruction to wait for the next interrupt.

4. If the Transfer Was Complete

When the transfer is complete, the handler can simply exit through the I/O completion section.

1.2.5 I/O Completion Section

The I/O completion section provides a common exit path to inform the monitor that the handler is done with the current request, so that the monitor can release the current queue element.

The I/O completion section is an extension of the interrupt service section. Control passes from the interrupt service section to the I/O completion section when a data transfer completes, when a hard error is detected, or when a soft error condition exhausts the number of allowed retries.

(Note that you can branch directly to this section from the I/O initiation section if you immediately detect a hard error.)

1. If an Error Occurred

There are two kinds of errors that cause control to pass to the I/O completion section: hard errors, which should cause a branch to this section immediately, and soft errors that have exhausted their allotted number of retries, which cause a branch to this section after the last retry fails. Treat both cases alike in handling the exit to the monitor.

First, set the hard error bit (HDERR\$), bit 0, in the Channel Status Word for the channel. The second word of the I/O queue element, Q.CSW, points to the Channel Status Word. Then jump to the I/O completion routine in the Resident Monitor. Use the .DRFIN macro, described below, to generate the code for this jump.

The following lines of code are from the DY handler. They illustrate how the handler sets the hard error bit and jumps back to the monitor.

```
10$:    BIC      #<CSINIT!CSINT>,@DYCSA ;DISABLE FLOPPY INTERRUPTS
                                     ;AND INHIBIT DRIVE RESET
11$:    .DRFIN  DY                      ;GO TO I/O COMPLETION
      .
      .
DYERR:  MOV     DYCQE,R4                ;R4 -> CURRENT QUEUE ELEMENT
        BIS     #HDERR$,@-(R4)        ;SET HARD ERROR IN CSW
        BR     10$                    ;EXIT ON HARD ERROR
```

2. If the Transfer Was Complete

For a block-oriented device, such as a disk or diskette, the handler simply disables interrupts and performs the jump to the monitor. The .DRFIN macro generates the code to perform the jump.

For a character- or word-oriented device, the procedure is slightly more complicated because the handler may have to report end-of-file to the job that requested the I/O transfer. When the handler actually detects the EOF condition on a READ operation, it should set an internal EOF flag, put the last character in the user's buffer, and then zero-fill the rest of the buffer. Then the handler should jump back to the monitor, as it would if EOF were not detected but the

buffer had simply filled up. The handler waits until it is called again to signal EOF to the user.

This convention for indicating end-of-file makes character-oriented devices appear to programs as random-access devices, which is in keeping with the RT-11 philosophy of device independence.

.DRFIN Macro

Use the .DRFIN macro to generate the instructions for the jump back to the monitor at the end of the handler I/O completion section. The macro makes the pointer to the current queue element a global symbol, and it generates Position-Independent Code for the jump to the monitor. When control passes to the monitor after the jump, the monitor releases the current queue element.

The format of the .DRFIN macro is as follows:

.DRFIN name

name is the two-character device name.

For examples of the .DRFIN macro, see the handler listings in Appendix A.

1.2.6 Handler Termination Section

The purpose of the handler termination section is to declare some global symbols and to establish a table of pointers to locations in the Resident Monitor. The pointers are filled in by the bootstrap, if the handler is for the system device. Otherwise, they are filled in when the handler is made resident with .FETCH or LOAD. The termination section also provides a symbol to determine the size of the handler. Use the .DREND macro to generate the handler termination code.

1.2.6.1 .DREND Macro

The format of the .DREND macro is as follows:

.DREND name

name is the two-character device name.

In bootable handlers, the .DREND macro is invoked twice, once explicitly by the programmer and once implicitly by the .DRBOT macro. When .DRBOT is invoked, it implicitly generates a .DREND macro to close the memory resident part of the handler. You end the boot area with a second .DREND macro.

For examples of the .DREND macro, see the handler listings in Appendix A. The symbols defined by .DREND are shown in Table 1-11.

1.2.7 Pseudodevices

You can write a device handler for a pseudodevice (one that does not interrupt, and is not a mass storage device) to take advantage of the queued I/O system and the fact that handlers can remain memory resident. Examples of handlers for pseudodevices are NL (the null device), MQ (the message queue handler), SL (the single-line command editor), and UB (the UMR handler).

All the executable code of such a handler must appear in the I/O initiation section. The handler should then issue the `.DRFIN` macro call to terminate the operation and return the queue element. Since pseudodevices do not interrupt, the handler needs no interrupt service section and no `.DRAST` macro call.

1.2.8 Handler Data Structures Related to Block 0

The following sections describe data structures that relate to block 0 of the handler file image. The data structure can reside within block 0 or be pointed to by an address contained there.

1.2.8.1 Handler Service Routine Environment

This section describes the handler service routine entry environment and error processing. The routines are defined by the `.DRPTR` macro and located at a file address in the handler (See the `.DRPTR` section in the *RT-11 System Macro Library Manual*).

Handler Service Routine Entry Environment

The following registers and their contents constitute the handler service routine entry environment. These registers (R0 through R5) are set up by RT-11. All registers are available and none needs to be preserved.

Table 1-5: Handler Service Routine Entry Environment

Register	Contents
R0	Contains starting address of the current running handler service routine.
R1	Contains starting address of GETVEC routine if a CTI Bus-based processor. Otherwise, R1 contains the address of a routine that always returns carry set.

Table 1–5 (Cont.): Handler Service Routine Entry Environment

Register	Contents
R2	Contains the value $\$SLOT*2$. That value is the length of the $\$PNAME$ table in bytes. You can use that value to locate information in the handler tables concerning this handler. The following table shows the order in memory and size in bytes, relative to $\$SLOT*2$, of the pertinent handler tables and the contents of those tables:

Table	Size	Contents
$\$OWNER$:	$\langle \$SLOT*2 \rangle *2$	Ownership table; can be removed from (generated out of) monitors
$\$UNAM1$:	$\langle \$SLOT*2 \rangle +4$	Physical name of device table
$\$UNAM2$:	$\langle \$SLOT*2 \rangle +4$	Logical name of device table
$\$PNAME$:	$\$SLOT*2$	Installed handlers table
$\$ENTRY$:	$\langle \$SLOT*2 \rangle +2$	Handler address table. Last word contains value -1 and indicates end of table
$\$STAT$:	$\$SLOT*2$	$DSTATUS$ value table
$\$DVREC$:	$\$SLOT*2$	Handler disk block table
$\$HSIZE$:	$\$SLOT*2$	Handler memory size table
$\$DVSIZ$:	$\$SLOT*2$	Device blocks table
$\$PNAM2$:	$\langle \$SLOT*2 \rangle +2$	Optional physical device name table for extended-unit (single letter) device names. Last word contains default device name, if assigned

You can use that table in the following manner. R5 contains the $\$ENTRY$ table entry address for this device handler. You could find, for example, the name for this handler in the $\$PNAME$ table by subtracting the value for $\$SLOT*2$ from the value contained in R5. Likewise, you could find the $DSTATUS$ value for this handler in the $\$STAT$ table by adding the value of $\langle \$SLOT*2 \rangle +2$ to the value contained in R5.

See the *RT-11 System Internals Manual* for more information about the handler tables.

Table 1–5 (Cont.): Handler Service Routine Entry Environment

Register	Contents																					
R3	Indicates the type of entry. The value in R3 indicates the type of routine that called the handler service routine: <table border="1"><thead><tr><th>Value</th><th>Name</th><th>Meaning</th></tr></thead><tbody><tr><td>0</td><td>HRR.FF</td><td>Entered from .FETCH</td></tr><tr><td>2</td><td>HRR.RE</td><td>Entered from .RELEASE</td></tr><tr><td>4</td><td>HRR.LO</td><td>Entered from the LOAD command</td></tr><tr><td>6</td><td>HRR.UN</td><td>Entered from the UNLOAD command</td></tr><tr><td>10</td><td>HRR.AB</td><td>Entered from a job abort (RELEASE routine)</td></tr><tr><td>12</td><td>HRR.SY</td><td>Entered from a system bootstrap load (LOAD routine)</td></tr></tbody></table>	Value	Name	Meaning	0	HRR.FF	Entered from .FETCH	2	HRR.RE	Entered from .RELEASE	4	HRR.LO	Entered from the LOAD command	6	HRR.UN	Entered from the UNLOAD command	10	HRR.AB	Entered from a job abort (RELEASE routine)	12	HRR.SY	Entered from a system bootstrap load (LOAD routine)
Value	Name	Meaning																				
0	HRR.FF	Entered from .FETCH																				
2	HRR.RE	Entered from .RELEASE																				
4	HRR.LO	Entered from the LOAD command																				
6	HRR.UN	Entered from the UNLOAD command																				
10	HRR.AB	Entered from a job abort (RELEASE routine)																				
12	HRR.SY	Entered from a system bootstrap load (LOAD routine)																				
R4	Contains the address of a read routine you can use to perform I/O to the system device, which has been opened as non-file-structured. You must load the following registers with the following contents to use the read routine: <table border="1"><thead><tr><th>Register</th><th>Contents</th></tr></thead><tbody><tr><td>R0</td><td>Block number to read</td></tr><tr><td>R1</td><td>Number of words to read</td></tr><tr><td>R2</td><td>Buffer address</td></tr></tbody></table> <p>You can read into only the low 28K words of memory. To read into high memory, you must first read into low memory and then move the data. The read routine returns with carry clear if there are no errors; carry bit is set if there are errors.</p>	Register	Contents	R0	Block number to read	R1	Number of words to read	R2	Buffer address													
Register	Contents																					
R0	Block number to read																					
R1	Number of words to read																					
R2	Buffer address																					
R5	Contains a pointer to the \$ENTRY table entry for this handler.																					

Handler Service Routine Error Processing

The following list shows how errors in handler service routines should be processed:

- If no errors occur, exit with carry bit clear.
- If errors occur, exit with carry bit set.

The response from RT–11 to handler service routines that exit with the carry bit set varies according to the following:

- If the handler service routine was called by the .FETCH request, RT–11 refuses to fetch the handler.

You should not depend on this response with handlers that should never be fetched; use the `.DRPTR FETCH=*NO*` parameter instead.

- If the handler service routine was called by the `.RELEASE` request, RT-11 releases the handler.
- If the handler service routine was called by the `LOAD` command, RT-11 refuses to load the handler.

You should not depend on this response with handlers that should never be loaded; use the `.DRPTR LOAD=*NO*` parameter instead.

- If the handler service routine was called by the `UNLOAD` command, RT-11 refuses to unload the handler. Further RT-11 response is determined by the contents of R0:
 - If R0 is returned with value zero, RT-11 displays the error message, `?KMON-F-Unable to unload handler`.
 - If R0 is returned with value other than zero, RT-11 displays the error message located at the address (in low memory) contained in R0.
- If the handler service routine was called by a job abort, RT-11 ignores the carry bit; the job aborts.
- If the handler service routine was called by a system bootstrap load, the handler can do one of the following:
 - Clear the carry bit and continue.
 - Set the carry bit and return.

On UNIBUS and Q-bus processors, RT-11 displays the message, `?BOOT-U-Failure to load system handler`, and the system halts. On CTI Bus-based processors, RT-11 displays code 000013 and the system halts.

- Set the carry bit and send an error message to the console terminal.

The handler sends an error message to the console terminal, using the following code:

```
CODE    =: <200!DEV.xx>
REPORT  =: 672
        JSR      R1, @#REPORT
        .WORD    MSG
        .BYTE    CODE
MSG:    .ASCIZ   "message"
        .EVEN
```

For UNIBUS and Q-bus processors:

- RT-11 ignores the contents of the byte `CODE`.
- RT-11 adds the prefix `"?BOOT-U-"` to `"message"`.
(For the distributed RT-11, `"message"` is `"Failure to load system handler"`.)
- The system halts.

For CTI Bus processors:

- RT-11 ignores the contents of "message".
- RT-11 displays the octal value contained in CODE with no prefix. The value in CODE should be 200!DEV.XX, where DEV.XX is the device id for this handler. You can find DEV.XX for this handler in the \$DVREC: handler table.

(For the distributed RT-11, CODE is 000013.)

- The system halts.

1.2.8.2 Special Function Code Support Table (H.SPFx)

H.SPFx supports both the *list* and *extension table* method for describing those special functions used within the handler. Using .DRSPF to create the table is described in the *RT-11 System Macro Library Manual*.

The .DRSPF macro places the table in octal offsets 22 through 30 in the handler's block 0. Offsets 22 through 26 support the list method and each offset has the same structure and is composed of a low and high byte. Offset 30 is a word pointer to a list of other special functions.

The symbol names for the values in H.SPFx are defined in the .DSPDF macro in the distributed file SYSTEM.MLB.

The following is the structure of offsets 22 through 26.

Bit	Symbol	Meaning
Low Byte	DSP.XN	The low byte, consisting of a bit mask that specifies the supported low-order numbers (xxN):
	Bit	Symbol Meaning
	001	DSP.X0 xx0 bit mask
	002	DSP.X1 xx1 bit mask
	004	DSP.X2 xx2 bit mask
	010	DSP.X3 xx3 bit mask
	020	DSP.X4 xx4 bit mask
	040	DSP.X5 xx5 bit mask
	100	DSP.X6 xx6 bit mask
	200	DSP.X7 xx7 bit mask

Bit	Symbol	Meaning																					
High Byte	DSP.NX	The high byte, made up of a value to specify the type of special function and the high order numbers (NNx). Specifying a type of special function forces the table entry to a single special function.																					
001–004	DSP.TY	The type of special function:																					
		<table border="1"> <thead> <tr> <th>Value</th> <th>Symbol</th> <th>Meaning</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>DSP.UK</td> <td>Unknown type</td> </tr> <tr> <td>1</td> <td>DSP.RD</td> <td>READ type</td> </tr> <tr> <td>2</td> <td>DSP.WR</td> <td>WRITE type</td> </tr> <tr> <td>3</td> <td>DSP.MV</td> <td>MOVEMENT type</td> </tr> <tr> <td>4</td> <td>DSP.RW</td> <td>TRANSFER type</td> </tr> <tr> <td>5–7</td> <td></td> <td>Reserved</td> </tr> </tbody> </table>	Value	Symbol	Meaning	0	DSP.UK	Unknown type	1	DSP.RD	READ type	2	DSP.WR	WRITE type	3	DSP.MV	MOVEMENT type	4	DSP.RW	TRANSFER type	5–7		Reserved
Value	Symbol	Meaning																					
0	DSP.UK	Unknown type																					
1	DSP.RD	READ type																					
2	DSP.WR	WRITE type																					
3	DSP.MV	MOVEMENT type																					
4	DSP.RW	TRANSFER type																					
5–7		Reserved																					
010–200	DSP.NN	Value for the special function's high-order two numbers																					

As an example, assume support for special functions 372, 373, and 377 (no type specified). The contents of the table entry for these would appear in a byte dump as:

```
370 214
```

For an example that includes the TYPE parameter, assume the special function 376 of type WRITE. The contents of the table entry for that would appear in a byte dump as:

```
372 100
```

1.2.8.3 Bad-Block Replacement Geometry Table (H.REPL)

H.REPL stores the geometry of the software (not MSCP) bad-block replacement table. The .DREST macro places a pointer to this table in offset 32₈ in the handler's block 0. The table must be located in block 0.

Of the distributed RT-11 device handlers, H.REPL is found in the RL01/02 and RK06/07 handlers.

The symbol names for the values in H.REPL are defined in the .RGTDF macro in the distributed file SYSTEM.MLB.

The table consists of 1-byte entries and is 6 bytes long.

Offset	Symbol	Contents
0	RGT.FG	A flag in bit 0. If bit 0 is clear, all blocks are replaceable; if set, only some blocks are replaceable. Bits 1–7 are reserved.
1	RGT.PD	A constant for locating the bad sector file. The last addressable block plus this constant is the bad sector file location.
2	RGT.BS	Size in sectors of bad sector file.
3	RGT.TC	Number of tracks per cylinder.
4	RGT.ST	Number of sectors per track.
5	RGT.SB	Half the number of sectors per block, such that two times this number is the sectors per block.
6	RGT.SZ	Size of this table.

1.2.8.4 Bad-Block Replacement Table (HB.BAD)

The bad-block replacement table is stored in the home block of RL01/02 and RK06/07 volumes, beginning at offset 6 (HB.BAD) and ending at offset 200.

The symbol names for the values in HB.BAD are defined in the .BBRDF macro in the distributed file SYSTEM.MLB.

Offset	Name	Meaning
0	BBR.BD	Bad block number.
2	BBR.GD	Replacement block number.
	BBR.SZ	Entry size.

1.2.8.5 Second Handler Status Word (H.STS2)

The following table defines the bits in the second handler status word (H.STS2), which the .DREST macro places in offset 36₈ of block 0.

Bit	Symbol	Meaning
000001	HS2.BI	Handler cannot be installed by the monitor bootstrap.
000002	HS2.KI	Handler cannot be installed by the DCL INSTALL command.
000004	HS2.KL	Handler cannot be loaded by the DCL LOAD command.
000010	HS2.KU	Handler cannot be unloaded by the DCL UNLOAD command.
000020	HS2.MO	Handler supports DCL MOUNT and DISMOUNT commands.

Bit	Symbol	Meaning
000040– 100000		Reserved.

1.2.8.6 Handler SYSGEN Options Byte (H.GEN)

The .DREND macro stores the SYSGEN option bits in H.GEN (byte offset 60₈ of block 0).

The value stored in H.GEN is the values for the SYSGEN options OR'd with the value of the .DREND FORCE= parameter.

The symbol names for the values in H.GEN are defined in the .SGNDF macro in the distributed file SYSTEM.MLB. (Note that only symbols in the range 1–200 can be used.)

Bit	Symbol	Meaning
001	ERLG\$	Handler supports error logging.
002	MMGT\$	Handler supports extended memory.
004	TIMIT\$	Handler supports device timeout.
010	RTEM\$	Handler is running under RTEM–11.
020– 200		Reserved.

1.2.8.7 Handler Internal Data Table and Descriptor Structure (H.TYPE, H.DATA, and H.DLEN)

The structure described in this section is a descriptor table. That is, the structure describes tables located elsewhere in the handler. The contents and location of the structure vary according to the macro that writes it. The structure can be placed in block 0 or an address can be placed in block 0 that points to the structure:

- The .DREST or .DRTAB macro stores the structure in block 0 offsets 70 through 74. The indicated offsets are from location 70.
- The .DRUSE macro stores the structure in the handler file and writes a pointer to the structure in block 0 offset 106. The indicated offsets are from handler file address pointed to by offset 106.

The symbol names for the values in H.TYPE, H.DATA, and H.DLEN are defined in the .DUSDF macro in the distributed file SYSTEM.MLB.

Offset	Symbol	Contents
00	DT.ID	If table generated by .DREST or .DRUSE, contains the RAD50 device type identifier. If table generated by .DRTAB, contains the value -1.
02	DT.PTR	If table generated by .DREST, contains the file address of internal data tables. If table generated by .DRUSE or .DRTAB, contains the file address of the list of data table descriptors.
04	DT.SIZ	Length in bytes of the data table pointed to by this structure.
06	DT.ESZ	When table is generated by .DRUSE only, is the length in bytes of each entry in the table pointed to by this structure.
10	DT.EOL	When table is generated by .DRUSE only, is a null word that signifies the end of the descriptor list.

1.2.8.8 UMR Support and Extended Device-Unit Handlers (H.64UM)

The contents of H.64UM describes the attributes of an extended device-unit handler and the support for UNIBUS Mapping Registers (UMRs).

The .DRDEF macro writes H.64UM in octal location 100 in the handler's block 0.

The symbol names for the values in H.64UM are defined in the .HUMDF macro in the distributed file SYSTEM.MLB.

Bit	Symbol	Meaning
000001– 000004	HUM.PU	Required number of permanent UMRs.
000010	HUM.S6	Handler supports other extended device-unit handlers (used in LD handler).
000020	HUM.DM	Handler uses DMA.
000040	HUM.UB	Handler includes .DRDEF macro DMA=str parameter (argument YES or NO).
000100– 100000	HUM.64	Field containing RAD50 letter for extended device-unit handler.

If HUM.UB bit is clear, bits HUM.UB, HUM.DM, and HUM.PU are reserved.

If HUM.PU bits are nonzero, HUM.DM must be set.

1.2.9 Handler Data Structures Related to Block 1

The following sections describe data structures that relate to block 1 of the handler file image. The data structure can reside within block 1 or be pointed to by an address contained there.

1.2.9.1 Handler Flag Word (H1.FLG)

H1.FLG contains flags that provide information about the handler.

The `.DRBEG` macro writes H1.FLG in octal location 10 of the handler's block 1 (location 1010 of the file image).

The symbol names for the values in H1.FLG are defined in the `.HBFDF` macro in the distributed file `SYSTEM.MLB`.

Bit	Symbol	Meaning
000001– 004000		Reserved.
010000	DV2.DM	Handler supports DMA and is compatible with RT-11 V5.5 (and subsequent) UMR support.
020000	DV2.NL	Handler cannot be loaded by KMON; can only be loaded by BSTRAP (at bootstrap time).
040000	DV2.V2	The first vector table set up by <code>.DRVTB</code> is followed by a second table. The second table is only for display purposes.
100000	DV2.NF	Handler cannot be fetched but instead must be loaded.

1.2.9.2 Handler Service Routine Entry Point Word (H1.NOP)

H1.NOP describes whether entry points to various handler service routines exist. It also defines the existence of a second handler flag word (H1.FG2). The low 5 bits are significant; the other bits are used to construct a NOP instruction and can be disregarded.

The `.DRBEG` macro stores the entry point in H1.NOP (offset 12₈ of block 1).

The symbol names for the values in H1.NOP are defined in the `.HUMDF` macro in the distributed file `SYSTEM.MLB`.

Bit	Symbol	Meaning
000001	HNP.FE	Handler contains entry point to a <code>FETCH</code> service routine.
000002	HNP.RE	Handler contains entry point to a <code>RELEASE</code> service routine.
000004	HNP.LO	Handler contains entry point to a <code>LOAD</code> service routine.
000010	HNP.UN	Handler contains entry point to an <code>UNLOAD</code> service routine.

Bit	Symbol	Meaning
000020	HNP.F2	Handler contains a second flag word (H1.FG2).
000040	HNP.N1	Part of the NOP instruction (disregard).
000100		Reserved.
000200	HNP.N2	Part of the NOP instruction (disregard).
000400– 100000		Reserved.

1.2.9.3 Second Handler Flag Word (H1.FG2)

H1.FG2 contains flags that provide additional information about the handler. If a flag indicates that a location after H1.FG2 is defined, then the preceding locations (to H1.FG2) are also defined.

The `.DRBEG` macro stores the second handler flag word in H1.FG2 (offset 16₈ in the handler's block 1).

The symbol names for the values in H1.FG2 are defined in the `.HF2DF` macro in the distributed file `SYSTEM.MLB`.

Bit	Symbol	Meaning
000001	HF2.SC	Handler code performs special function address checking (therefore H1.SCK exists).
000002	HF2.SD	Handler lists special functions that use DMA (therefore H1.SDF and H1.SCK exist).
000004	HF2.LD	Handler contains pointer to LD translation table (therefore H1.LDT, H1.SDF, and H1.SCK exist).
000010	HF2.ND	Handler contains nonstandard DMA special functions (therefore H1.NDF, H1.LDT, H1.SDF, and H1.SCK exist).
000020– 002000		Restricted.
004000	HF2.SR	Handler requires serial satisfaction of I/O requests.
010000	HF2.DM	Handler performs DMA and is compatible with RT-11 V5.5 UMR support.
020000	HF2.S6	Handler supports other extended device-unit handlers (used in LD handler).
040000	HF2.64	Handler supports extended device-unit requests.
100000	HF2.F3	Handler contains a third flag word.

1.2.10 Skeleton Outline of a Device Handler

The skeleton outline in Figure 1–1 provides the structure for a simple device handler. In the figure, *SK* is the device name.

Figure 1–1: Skeleton Device Handler

```
.Title SK -- Handler Skeleton
; SK DEVICE HANDLER
.IDENT /V05.05/
.SBTTL PREAMBLE SECTION
.MCALL .DRDEF ; Get handler definitions
.MCALL .ASSUME ; Checking macro
.MCALL .EXIT ; To finish run
.MACRO ... ; Define ellipsis (allow
; ellipsis to assemble)
.ENDM
; Generate nonexecutable handler information tables
; containing the following information:
; Handler is SK
; Handler ID is 350 (user-written handler)
; Handler accepts neither .READ nor .WRITE
; Handler accepts .SPFUN requests
; Device is 1 block in size
; Device has a CSR at 176544
; Device has a vector at 20
.DRDEF SK,350,WONLY$,SPFUN$,1,176544,20
; Handler has .Fetch and $LOAD code to be executed:
.DRPTR FETCH=Fetch,LOAD=Load
; Handler is for a "Null" class device
; Handler has a data table called DATABL
; Data table is of the SKL format
.DREST CLASS=DVC.NL,DATA=DATABL,TYPE=SKL
; Handler accepts the following SPFUN codes:
; 372,376,377
.DRSPF <372>,TYPE=T
.DRSPF <376>,TYPE=W
.DRSPF <377>,TYPE=R
; Handler CSR is not to be checked at install,
; but is to be displayed:
.DRINS -SK
; Here is any installation check code
...
RETURN
.ASSUME . LE 400,MESSAGE=<;Installation area overflow>
```

Figure 1–1 (continued on next page)

Figure 1–1 (Cont.): Skeleton Device Handler

```

; Handler accepts SET SK [NO]BONES command:
.DRSET BONES,123456,CORPUS,NO
CORPUS:
COM R3 ; SET SK BONES
NOP ; Pad code
.ASSUME . EQ CORPUS+4,MESSAGE=<;No option code in wrong place>
NOCORP: ; SET SK NOBONES
MOV R3,PICKNT ; Set value in block 1
RETURN
.ASSUME . LE 1000,MESSAGE=<;Set area overflow>

.SBTTL HEADER SECTION
.DRBEG SK ; Handler Queue Manager Entry point
BR START ; Skip data table
DATABL:
.RAD50 "SKL" ; Table ID
WRIST: .BLKW 1 ; Table contents
ANKLE: .BLKW 1 ; ...
;Set up the Vector table:
SK$VTB: .DRVTB SK,SK$VEC,SKINT,0
.DRVTB ,SK$VEC+4,SKINT,1
PICKNT: .BLKW 1 ; Value controlled by Set command
.ASSUME .-2 LE SKSTRT+1000,MESSAGE=<;Set object not in block 1>

.SBTTL I/O INITIATION SECTION
START: ; Executable Queue code
...
RETURN
.SBTTL INTERRUPT SERVICE SECTION
.DRAST SK,4,ABORT ; Interrupt entry point
BCS INT2 ; Interrupt from second vector
...
RETURN
INT2: ; Second interrupt vector code
...
RETURN
.SBTTL I/O COMPLETION SECTION
ABORT: ; Abort entry point
...
.DRFIN SK ; Completion return
; End of memory resident part of handler
.DRBOT SK,ENTRY ; Boot code
ENTRY:
... ; Hard boot code to call read routine
RETURN

```

Figure 1–1 (continued on next page)

Figure 1–1 (Cont.): Skeleton Device Handler

```
READ:
    ...                ; Read routine
    RETURN
.SBTTL HANDLER TERMINATION SECTION
    .DREND SK          ; End of boot code
    .PSECT SETOVR      ; Suggested block aligned Psect
FETCH:
    ...                ; Code executed on FETCH
    RETURN
LOAD:
    ...                ; Code executed on LOAD
    RETURN
RUN:
    ...                ; Code executed on RUN
    .EXIT
    .END RUN
```

1.3 Abort Processing

This section describes the behavior of the resident monitor (RMON) and a device handler when a job abort occurs.

The action taken by RMON in abort processing is determined by three criteria:

- The setting of the ABTIO\$ and HNDLR\$ bits in the device status word (H.STS).
- The action that caused the abort.
- The presence or absence of a current queue element belonging to the aborting job (or job and channel in the case of .ABTIO aborts).

The first two criteria are described in the following sections. Section 1.3.2 contains a table showing the matrix and order of RMON actions based on combinations of all those criteria.

1.3.1 Handler Status Word Bits ABTIO\$ and HNDLR\$

The combination of ABTIO\$ and HNDLR\$, whether set or clear, determines to the following extent how RMON performs abort processing for that handler and other handlers that are loaded in memory:

- If ABTIO\$ is set, the handler is entered by RMON during any type of abort; the status of HNDLR\$ (set or clear) does not matter.
- If ABTIO\$ or HNDLR\$ is set (but not both), the handler is entered by RMON when a .ABTIO request is issued by a program to any handler.

When a program invokes the .ABTIO request for a channel associated with any handler, RMON calls the abort entry point of all in-memory handlers having that bit combination (ABTIO\$ or HNDLR\$ set, but not both). RMON checks each handler for I/O requests that might be internally queued on the channel that is specified in the .ABTIO request. RMON performs abort processing for any outstanding I/O request on the channel being aborted by the .ABTIO request. RMON does not discard the current queue element (*ddCQE*) and whether or not it is satisfied is determined by the handler.

If the hanlder aborts the current queue element, it should clear the queue element's completion routine address (Q.COMP) and issue a .DRFIN to return the queue element to the monitor. All outstanding queue elements that are associated with the aborting job or job and channel are removed from the handler's queue element list.

- If HNDLR\$ is set and ABTIO\$ is clear, RMON does not keep count (in I.IOCT) of the number of outstanding queue elements for that handler.

Some handlers, such as the distributed RT-11 MQ and Ethernet handlers, can post a request without necessarily expecting satisfaction of that request. To allow such handlers to be aborted, RMON is inhibited from keeping a count (in I.IOCT) of all outstanding I/O requests. Such handlers can then be aborted when they still contain outstanding queue elements.

Any user-written internally queued handler that can post an I/O request without requiring satisfaction of that request should be built with HNDLR\$ set and ABTIO\$ clear.

1.3.2 Types of Aborts and Action Taken by RMON

The resident monitor performs abort processing for any of the following actions:

Abort Type	Description
.CHAIN	I/O for the chaining job is allowed to complete.
.EXIT	Job I/O is allowed to complete.
.SRESET	
.HRESET	Hard error condition. Job I/O is stopped. ? <i>MON-F</i> - means an abort caused by a fatal monitor error. < <i>CTRL/C</i> > means a double CTRL/C typed at the keyboard.
? <i>MON-F</i> - < <i>CTRL/C</i> >	
.ABTIO (Handler used by this channel)	A .ABTIO request is issued for a handler that is associated with the aborting job's channel control block.

Abort Type	Description
.ABTIO (All other handlers)	This handler assembled with device status word bit HNDLR\$ set and ABTIO\$ clear, and is entered whenever a .ABTIO request is called for any handler on any channel.

Table 1–6 illustrates RMON abort processing. It not only shows the actions performed by RMON, but also the order in which they are performed. Before the table is a legend that defines and explains the symbols used in the table.

The order of certain symbols in the tables is important. The symbols show the order of abort processing for the type of abort. A note defines the symbols that should be read in order.

Symbol Definitions and Explanations for Table 1–6

Symbol	Definition/Explanation
Abort Type	The action that caused the abort.
A\$=0	The handler is not built with ABTIO\$ (ABTIO\$=0).
A\$=1	The handler is built with ABTIO\$ (ABTIO\$=1).
H\$=0	The handler is not built with HNDLR\$ (HNDLR\$=0).
H\$=1	The handler is built with HNDLR\$ (HNDLR\$=1).
ddCQE	The handler contains a current queue element belonging to the aborting job (or job and channel if .ABTIO). The absence of this symbol in a header indicates the handler has no current queue element associated with the aborting job (or job and channel if .ABTIO).

NOTE

The order of the following symbols in the tables is important. The symbols show the order of abort processing for the type of abort. For example, the symbols EJ show that operation E is performed first and operation J is performed next.

C	RMON removes all queue elements belonging to the job and channel from the queue and decrements I.IOCT one time for each element removed.
C~	RMON removes all queue elements belonging to the job and channel from the queue but does not decrement I.IOCT.
E	RMON calls the handler's abort entry point.

Symbol	Definition/Explanation
J	RMON removes all queue elements belonging to the job from the queue and decrements I.IOCT one time for each element removed.
J~	RMON removes all queue elements belonging to the job from the queue but does not decrement I.IOCT.
Q	RMON waits for all I/O requests for which it expects satisfaction to be satisfied.
S	RMON waits for all I/O requests for which it expects satisfaction to be satisfied and then issues a .ABTIO for every channel associated with the job.
()	RMON performs abort processing only if there is outstanding I/O on the channel.
–	RMON does not perform abort processing on this handler.

Table 1–6: RMON Abort Processing

Abort Type	A\$=0	A\$=0	A\$=0	A\$=0	A\$=1	A\$=1	A\$=1	A\$=1
	H\$=0	H\$=0 ddCQE	H\$=1	H\$=1 ddCQE	H\$=0	H\$=0 ddCQE	H\$=1	H\$=1 ddCQE
.CHAIN	S	S	S	S	S	S	S	S
.EXIT .SRESET	Q	Q	QEJ~	QEJ~	QEJ	QEJ	QEJ	QEJ
.HRESET ?MON-F- <CTRL/C>	J	EJ	EJ~	EJ~	EJ	EJ	EJ	EJ
.ABTIO (Handler used by this channel)	(C)	(EC)	(EC~)	(EC~)	(EC)	(EC)	(EC)	(EC)
.ABTIO (All other handlers)	–	–	(EC~)	(EC~)	(EC)	(EC)	–	–

1.4 Handlers That Queue Internally

A device handler can maintain one or more of its own internal queues of outstanding I/O requests instead of using the usual monitor/handler I/O queue. The purpose of maintaining an internal queue is that it permits several operations to take place on the device simultaneously—that is, the handler can service several requests to access the device at once. Internal queuing might also be useful if a handler needs to perform some type of request ordering based on device-specific criteria.

The distributed RT-11 handlers that control communications, XC, XL, NC, NQ, and NU, use internal queuing to process simultaneous input and output requests. See Figure A-3 for a commented source listing of the XL handler for guidance in implementing internal queuing in your handler.

1.4.1 Implementing Internal Queuing

A handler is entered at its `.DRBEG` code whenever the queue manager places an I/O request queue element on the handler's empty device queue. The handler checks the queue element for validity. An invalid request returns an immediate hard error.

A handler that implements internal queuing decides how to dispose of the current queue element based on whether processing the request requires post-interrupt activity (another interrupt). If the I/O request does not require post-interrupt activity by the handler, the handler processes the queue element immediately and returns, through `.DRFIN`, to the monitor. If processing the request cannot be immediately satisfied, the handler removes the request queue element from the device queue and places it on an internal queue. The device queue is then available for another request.

The internally queued handler has sole responsibility for managing internally queued queue elements; for moving them between the internal queue and the device queue. The handler is also responsible for returning appropriate queue elements to the monitor because of an abort on a channel or job.

1.4.2 Interrupt Service for Handlers That Queue Internally

When an operation completes, the handler is normally entered at its interrupt entry point, `ddINT`. After this, various actions are taken depending on the circumstances. If there is more than one internal queue, the handler determines which request this interrupt involves and, therefore, which internal queue. If the operation is not complete, the handler restarts it or continues it and simply returns to the monitor. If the transfer is complete, the handler returns the request to the monitor by using a fake device queue and modified `.DRFIN` code.

The handler returns the request to the monitor without exiting in order to process any further outstanding requests. The fake device queue is used to avoid any race condition conflict with the monitor over the use of the device queue. The modified form of `.DRFIN` code uses a `CALL` rather than a `JMP` instruction, so that the handler can regain control after the request is returned to the monitor.

The following example illustrates how an internally queued handler returns a queue element to the monitor. In the example, R4 points to the third word of the queue element to be returned.

```

.
.
.
MOV     R4,ddFCQE           ; Make queue element first
MOV     R4,ddFLQE           ; and last on fake device queue
CLR     Q$LINK(R4)         ; Make sure it doesn't link anywhere
MOV     PC,R4               ; R4 -> Fake device queue
ADD     #ddFCQE-.,R4        ; ...
MOV     @#$SYPTR,R5         ; R5 -> $RMON
CALL    @$QCOMP(R5)        ; Return the queue element
.
.
.
; Check the internal queue and start another operation if necessary
.
.
.
RETURN

; Fake device queue
        .WORD    0           ; Required
ddFLQE: .BLKW           ; Fake LQE
ddFCQE: .BLKW           ; Fake CQE

```

1.4.3 Abort Procedures for Handlers That Queue Internally

As explained in Section 1.3, the contents of the handler status word, H.DSTS, determines how a handler and RMON process aborts. In particular, it is the ABTIO\$/HNDLR\$ bit combination in the handler status word. There are some particular considerations with abort processing for a handler that internally queues I/O requests:

- Does the handler expect satisfaction of all outstanding I/O requests?
Setting bit ABTIO\$ and not HNDLR\$ stops RMON from maintaining the count (I.IOCT) of outstanding I/O requests for the handler.
- Do other handlers in the system need to be notified if the handler processes an abort? Conversely, does the handler need to be notified if other handlers on the system process an abort?

All in-memory handlers that are built with either ABTIO\$ or HNDLR\$ set (but not both set) are entered at their abort entry point by RMON whenever a .ABTIO request is issued by a program. Also, RMON checks for internally queued I/O requests on the specified channel. Abort processing is performed on any handler having outstanding I/O requests on the channel being aborted by a .ABTIO request.

Whether or not the current I/O request (*ddCQE*) is satisfied is determined by the handler code. All other queue elements associated with the job or the job and

channel are removed from the handler's queue element list. That is, *ddLQE* and *ddCQE* are set to the same value.

When the handler is entered at the abort entry point, it checks its internal queue for elements belonging to the aborted job. The job number is passed to the handler in R4. Whether the handler aborts all queue elements belonging to that job or only those for a particular channel is determined by the contents of R5. If R5 contains zero, the handler should abort all queue elements assigned to that job. If R5 is nonzero, it points to the first word of a channel control block (the channel status word), and the handler should abort only the queue elements for that channel.

The handler should purge its internal queue of those elements and use the following procedure to reduce the monitor's count of outstanding I/O requests. R0 through R3 must be saved and restored.

1. Remove any internal queue elements that belong to the aborting job or channel. If there are none, simply issue the RETURN instruction.
2. Otherwise, link the removed elements through the element's link word (Q.LINK); the last element's link word must be 0. Set *ddCQE* to point to the last element of this linked list.
3. Clear each aborting queue element's completion routine address (Q.COMP).
4. Issue the .DRFIN macro.

1.5 Set Options

The keyboard monitor SET command permits you to change certain characteristics of a device handler. The handler must exist as a *dd.SYS* file on the system device (*ddX.SYS* for mapped systems), where *dd* is the two-character device name. For example, the following command changes the column width for a printer:

```
SET LP WIDTH=80          (The default is 132 columns)
```

Another type of SET command can enable or disable a function. The following example shows how a SET command can cause the system to send carriage returns to a printer or to refrain from sending them.

```
SET LP CR                (Sends carriage returns; this is the default)
```

```
SET LP NOCR              (Does not send carriage returns)
```

Note that you negate the CR option by adding NO to the start of the option. See the *RT-11 Commands Manual* for more information on the SET options available with existing RT-11 device handlers.

A device handler you write can contain code to implement different options. Follow the format outlined in the following sections to learn how to add SET options to your handler. Adding a SET option affects only the handler file; you need not make any changes to the monitor. Note that SET options are valid for both data and system devices.

1.5.1 How the SET Command Executes

The SET command is driven entirely by a table in block 0 of the handler file and by a set of routines, also in block 0, that modify instructions and data in blocks 0 and 1 of the handler. Remember that block 0 refers to addresses 0 through 776, and that the handler header starts in block 1 at location 1000 in the file.

When you type a SET command at the console terminal, the monitor parses the command line and looks for the handler file on the system device. (The type of handler matches the monitor, such as DU.SYS for unmapped monitors or DUX.SYS for mapped monitors.) The handler need not be installed in the running system. The monitor then reads blocks 0 and 1 of the handler into the USR buffer. It scans the table in block 0 until it finds the table entries for the SET option you specified. From the table entry, it can find the particular routine designed to implement that option and the modifiers permitted by that routine, such as NO or a numeric value. The monitor then executes the routine, which contains instructions that modify code in blocks 0 or 1 of the handler. The code in block 1 is part of the body of the handler and contains the instructions for the default settings of all the SET options. After the code is modified, the monitor writes blocks 0 and 1 back out to the system device. Thus, as a result of the SET command, some instructions or data in the handler file are changed. However, any memory-resident copy of the handler is not affected.

1.5.2 SET Table Format

The table for the SET options consists of a series of four-word entries, with one entry per option. The table begins at location 400 in block 0 of the handler and ends an entry with a word zero. Use the .DRSET macro, described below, to generate the table. Examples of overlaid SET code are located in the example handlers in Appendix A.

The first word of the table is a value to be passed in R3 to the SET routine associated with the option when the monitor processes this option. This word can be a numeric value—such as the default column width for a printer—or it can be an instruction to substitute for another instruction in block 1 of the handler. It must not be 0.

The second and third words of the table are the option name in Radix-50, such as WIDTH or CR. In the table, the characters are left justified and filled with spaces.

The low byte of the fourth word is an offset to the routine that performs the code modification. The high byte indicates the type of SET parameter that is valid. Setting the 100 bit shows that a decimal argument is required. A value of 140 shows that an octal argument is required. Setting the 200 bit means that the NO prefix is valid for this option.

Table 1-7 shows a summary of the SET option table.

Table 1–7: SET Option Table

Offset	Name	Meaning
0	DSE.R3	Value to pass in R3 to the SET routine
2–4	DSE.NA	Radix–50 for option name (two words)
6	DSE.SB	Offset to option routine
7	DSE.PA	Parsing option bits:

Bit	Name	Meaning
0–4		Reserved
5	DSE.8	Set means option has octal value Clear means option has decimal value
6	DSE.NU	Numeric value allowed
7	DSE.NO	NO prefix allowed

DSE.ES	Entry size	
--------	------------	--

1.5.3 .DRSET Macro

Use the .DRSET macro to set up the option table by calling the macro once for each option so that the macro calls appear one after the other. You must invoke the .DRSET macro after .DRDEF and before the .DRBEG macro.

The format for the .DRSET macro is as follows:

.DRSET option, val, rtn[, mode]

option is the name of the SET option, such as WIDTH or CR. The name can be up to six alphanumeric characters long and cannot contain any embedded spaces or tabs.

val is a parameter that will be passed to the routine in R3. It can be a numeric constant, such as the minimum column width, or an entire instruction enclosed in angle brackets to substitute for an existing instruction in block 0 or 1 of the handler. This parameter must not be 0.

rtn is the name of the routine that modifies the code in block 0 or 1 of the handler. The routine must follow the option table in block 0 and not extend above file address 776. If you need more space for SET code, then this lets you overlay the SET code. See the DL example handler in Appendix A.

mode is an optional argument to indicate the type of SET parameter. Enter *NO* to indicate that a NO prefix is valid for the option. Enter *NUM* if a decimal value is required. Enter *OCT* if an octal value is required. Omitting the *mode* argument indicates that the option takes neither a NO prefix nor a numeric argument. You can combine the NO and numeric arguments as follows. The construction <NO,NUM> indicates that both a NO prefix and a decimal value are valid. The construction <NO,OCT> indicates that both a NO prefix and an octal value are valid. Omitting the *mode* argument forces a 0 into the high byte of the last word of the table entry.

See the sections below for examples of the .DRSET macro.

The first .DRSET macro issues an .ASECT directive and sets the location counter to 400 for the start of the table. The macro also generates a zero word for the end of the table. Because the macro leaves the location counter at the end of the table, you should place the routines to modify code immediately after the .DRSET macro calls in your handler. This makes sure that they are located in block 0 of the handler file.

1.5.4 Routines to Modify the Handler

Your handler needs a routine for each SET option. You need only one routine for an option and the NO version of that option. The purpose of the routine is to modify code in the body of the handler based on the SET command typed on the console terminal. One routine can support several SET options. Typically, the value passed in R3 is used to determine which SET option is being performed.

The routines must immediately follow the option table, described above, and they must be located in block 0, after the table and below address 1000. The code in the body of the handler that the routines modify must be in block 1 of the handler, within the first 256₁₀ words.

The name of the routine is its default entry point. This is the entry point for options that take a numeric value, for options that take neither a numeric value nor a NO prefix, and for options that accept a NO prefix but do not currently have it. The entry point for options that allow and have a NO prefix is the default entry point + 4.

On entry to the routine, for all options, the carry bit is clear and registers R0, R1, and R3 contain information for use by the routine and R4 and R5 should be preserved. If numeric values are valid for the option, R0 contains the numeric value from the SET command line. R1 contains the unit number specified as part of the device name; if no unit number was specified, the sign bit is set. R3 contains the *val* word of the SET option table (from .DRSET).

The routine can indicate that a command is illegal by returning with the carry bit set. For example, the printer SET WIDTH option does not allow a width less than 30. If the option routine indicates failure, the monitor prints an error message and does not write out blocks 0 and 1. Thus, the check can be made after the block 1 code is modified.

Once you have added the routines for each option to your handler, you can use the following line of code to make sure you are within the size bounds:

```
.IIF GT,<.-1000>, .ERROR .-1000 ; SET code too big!
```

Then you continue with the rest of the handler code, starting with the `.DRBEG` macro, which implicitly resets the location counter to 1000 and establishes the handler header.

1.5.5 Examples of SET Options

The following examples taken from a printer handler are implementations of SET options.

The examples were chosen to reflect the SET command examples shown at the beginning of this section. The SET commands were as follows:

```
SET LP WIDTH = 80
SET LP CR
SET LP NOCR
```

First, the handler invokes the `.DRSET` macro to set up the option tables for the two options `WIDTH` and `CR`.

The first call indicates that the printer `WIDTH` option is being established, that 30 decimal is a default value of some kind, that `O.WIDTH` is the routine to process the option, and that it takes a numeric argument.

```
.DRSET WIDTH,30.,O.WIDTH,NUM
```

The next call indicates that the printer `CR` option is being established, that `NOP` is to be passed to the routine, that `O.CR` is the name of the routine to process the option, and that the `CR` option can take a `NO` prefix.

```
.DRSET CR,NOP,O.CR,NO
```

The two macro calls generate the following table:

```
.ASECT
. = 400
.WORD 30. ;MINIMUM WIDTH
.RAD50 \WIDTH \ ;OPTION NAME
.BYTE <O.WIDTH-400>/2
.BYTE 100

.WORD NOP ;INSTRUCTION TO PASS
.RAD50 \CR \ ;OPTION NAME
.BYTE <O.CR-400>/2
.BYTE 200
.WORD 0 ;END OF TABLE
```

The routines to process these options immediately follow the end of the table. The following examples show the routines. The body of the code in block 1 of the handler that the routines modify is shown at the end of the section.

```

O.WIDTH:MOV    R0,COLCNT           ;MOVE VALUE FROM USER TO
            MOV    R0,RSTC+2       ;TWO CONSTANTS
            CMP    R0,R3           ;COMPARE NEW VALUE TO
                                     ;MINIMUM WIDTH, 30.
            RTS    PC              ;RETURN; C BIT SET ON ERROR

```

Note in the example above that the instructions in the routine O.WIDTH change data in two locations in block 1 of the handler.

```

O.CR:   MOV    (PC)+,R3           ;ENTRY POINT FOR "CR"; MOVE
                                     ;ADDRESS OF NEXT LINE TO R3
        MOV    R3,CROPT          ;ENTRY POINT FOR
                                     ;"NOCR" (O.CR+4);
                                     ;MOVE EITHER "NOP" OR
                                     ;PREVIOUS LINE TO CROPT
        BEQ    RSTC-CROPT+.       ;A NEW INSTRUCTION
        RTS    PC                ;RETURN

```

NOTE

While executing the routines to process a SET option, R4 and R5 are not available for use.

The routine O.CR has two entry points: for the "CR" option, the routine is entered at O.CR; for the "NOCR" option, the routine is entered at O.CR + 4. Note that (1) the routine substitutes one of two instructions for an instruction located in block 1; (2) a NOP instruction is moved to CROPT if the "NOCR" option is selected; (3) if "CR" is selected, the BEQ RSTC-CROPT+. instruction is moved to CROPT.

The construction of the BEQ instruction is necessary because the branch is being assembled into a location other than the one from which it will be executed. In all the routines, a branch instruction must use the following construction to generate the correct address:

```
BR      A-B+.
```

A is the destination of the branch instruction.

B is the address of the branch instruction.

. is the current location counter.

Generally, only routines for options that accept NO use these branch instructions.

Finally, look at the code in the interrupt service section of the handler that is modified by the routines you have just seen. Remember that the code to be modified must be located in block 1 of the handler, in the first 256₁₀ words.

```

      .
      .
COLCNT: .WORD    COLSIZ                ;# OF PRINTER COLUMNS LEFT
      .
      .
CHRTST: CMPB     R5,#HT                ;IS CHAR TAB?
        BEQ      TABSET                ;YES, RESET TAB
        CMPB     R5,#LF                ;IS IT LINE FEED?
        BEQ      RSTC                  ;YES, RESTORE COLUMN COUNT
        CMPB     R5,#CR                ;IS IT CARRIAGE RETURN?
CROPT:  NOP                               ;"NOP" IF "NOCR" OPTION;
                                           ;ELSE IF "CR" OPTION, USE
                                           ;"BEQ RSTC-CROPT+." FROM
                                           ;SET ROUTINES IN BLOCK 0.
        CMPB     R5,#FF                ;IS IT FORM FEED?
        BNE      IGNORE                ;NO, IT IS NON-PRINTING
RSTC:   MOV      #COLSIZ,COLCNT        ;RE-INIT COLUMN COUNTER

```

From the examples in the first part of this section, you can see how the routines in block 0 can modify data and instructions in block 1 of the handler.

1.6 Device I/O Timeout

The device timeout feature lets a handler assign a completion routine to be executed if an interrupt does not occur within a specified time interval. Thus, the handler can perform the equivalent of a mark time operation without the need for a `.SYNCH` call and its attendant potential delay.

Device timeout is supported by all distributed mapped monitors and is an optional feature on unmapped monitors, available through system generation. (Device timeout support requires monitor timer support, which is included on all distributed monitors except SB.) Device timeout is required by the RT-11 multiterminal monitor and support for it is automatically included when you build that monitor.

Within the handler, you select device timeout by including the system conditional `TIM$IT=1`. RT-11 provides two macros to help you implement device timeout in your handler. The macros, which are described below, are `.TIMIO` and `.CTIMIO`. They are available only to device handlers. If you assemble the handler file with the conditional `TIM$IT` equal to 1, the `.DRDEF` macro issues a `.MCALL` directive for the `.TIMIO` and `.CTIMIO` macros.

All code in your handler that applies strictly to device timeout support should be placed inside conditional assembly directives. These directives should include the device timeout code if the symbol `TIM$IT` is 1, and omit it otherwise. This way, the system parameters select whether or not the device timeout code is included in the handler each time you assemble it.

1.6.1 .TIMIO Macro

Use the `.TIMIO` macro in the handler I/O initiation section to issue the timeout call. You can issue the request anywhere in the handler except at interrupt level. If you need to issue the request at interrupt level, you must issue a `.FORK` macro call first.

The `.TIMIO` request schedules a completion routine to run after the specified time interval has elapsed. The completion routine runs in the context of the job indicated in the timer block. In mapped monitor systems, the completion routine executes with kernel mapping, since it is still a part of the interrupt service routine. (See the *RT-11 System Internals Manual* for more information about interrupt service routines and the mapped monitor environment.) As usual with completion routines, R0 and R1 are available for use. When the completion routine is entered, R0 contains the sequence number of the request that timed out.

Because you must go to fork level (and processor priority 0) to issue a `.TIMIO` or `.CTIMIO` request at interrupt level, your handler must disable device interrupts before issuing the `.FORK`, or must be carefully coded to avoid reentrancy problems. Note that you cannot reuse a timer block until either the timer element expires and the completion routine is entered, or the timer element is canceled successfully.

The format of the macro is as follows:

.TIMIO *tbk,hi,lo*

tbk is the address of the timer block, a seven-word pseudotimer queue element, described below. Note that you must not use a number sign (#) before *tbk*.

hi is a constant specifying the high-order word of a two-word time interval.

lo is a constant specifying the low-order word of a two-word time interval.

The timer block format is shown in Table 1–8.

Table 1–8: Timer Block Format

Offset	Name	Agent	Contents
0	C.HOT	.TIMIO	High-order time word.
2	C.LOT	.TIMIO	Low-order time word.
4	C.LINK	monitor	Link to next queue element; 0 indicates none.
6	C.JNUM	user	Owner's job number; get this from the queue element.
10	C.SEQ	user	Sequence number of timer request. The valid range for sequence numbers is from 177700 through 177377.
12	C.SYS	monitor	–1
14	C.COMP	user	Address of the completion routine to execute if timeout occurs. The monitor zeroes this word when it calls the completion routine, indicating that the timer block is available for reuse.

Although the `.TIMIO` macro moves the high- and low-order time words to the timer block for you, you must take care to specify them properly in the macro call. Express the time interval in ticks. There are 60₁₀ ticks per second if your system is running with 60-cycle power. If your system is running with 50-cycle power, there are 50₁₀

ticks per second. Professional 300 series processors have 60_{10} ticks per second with either line frequency. Time values for 50-cycle power are shown in square brackets ([]) immediately after the 60-cycle figure.

The low-order time word accommodates values of up to 65535_{10} ticks. That is equal to about 1092 [1310] seconds, or about 18.2 [21.8] minutes. If you need to specify a time interval of 18.2 [21.8] minutes or less, place a zero in the *hi* argument, and the number of ticks in the *lo* argument to the .TIMIO macro.

If you need to specify a time interval longer than 18.2 [21.8] minutes, think of the high-order word as a carry word. Each interval of 18.2 [21.8] minutes' duration causes a carry of 1 into the high-order word. So, to specify an interval slightly greater than 18.2 [21.8] minutes, supply a 1 to the *hi* argument, and a 0 to the *lo* argument. To specify 36.4 [43.6] minutes, move 2 to the *hi* argument, 0 to the *lo* argument, and so on. Since the 2-word time permits you to indicate up to 65565 units of 18.2 [21.8] minutes each, the largest time interval you can specify is about 2.3 [2.7] years.

The only words of information you must set up yourself in the timer block are the job number, the sequence number, and the address of the completion routine. You can get the job number from the current queue element, and then move it to the timer block. You assign the sequence number yourself. To ensure a unique number, use a value of $177000+dd\$COD$, where $dd\$COD$ is the device identifier code used in the .DRDEF macro at the beginning of the handler. The job number and sequence number are passed to the completion routine when it is entered. You must move the address of the completion routine to the seventh word of the timer block in a position-independent manner.

The .TIMIO macro expands as follows:

```
.TIMIO   tbk,hi,lo
JSR      R5,@$TIMIT           ;POINTER AT END OF HANDLER
.WORD    tbk - .
.WORD    0                    ;CODE FOR .TIMIO
.WORD    hi                   ;HI ORDER TIME INTERVAL
.WORD    lo                   ;LO ORDER TIME INTERVAL
```

1.6.2 .CTIMIO Macro

When the condition the handler was waiting for occurs, you should issue a cancel timeout call, which disables the completion routine. Use the .CTIMIO macro call in your handler to cancel the timeout request. Execution must be in system state when you issue the call. Be sure to issue a .FORK call first if you use .CTIMIO at interrupt level.

For example, a printer handler could check for an off-line condition. When a program requests an I/O transfer, the handler's I/O initiation section forces an immediate interrupt. The handler's interrupt service section then checks the device error bit. If the bit is set, the printer is not on line and the handler prints a message, sets a 2-minute timer with .TIMIO, and returns to the monitor with a RETURN instruction to wait for another interrupt. The device should not interrupt again until the error condition has been fixed by an operator. If no interrupt occurs within two minutes,

the timer completion routine prints another error message, sets another 2-minute timer, and returns again to the monitor with RETURN to wait for an interrupt. (See Figure 1-2 for a printer handler example.)

In this example, when an interrupt finally occurs and the error bit is clear, the handler issues the .CTIMIO call to cancel the timed wait.

As another example, a disk handler could set a timer before it starts up a seek operation. When the interrupt occurs, the seek is complete, and the handler should then cancel the timer.

If the time interval in any application has already elapsed and the device has, therefore, timed out, the .CTIMIO request fails. Because the completion routine has already been placed in the queue, the .CTIMIO call returns with the carry bit set. You can usually ignore this condition.

The format of the .CTIMIO macro call is as follows:

.CTIMIO *tbk*

tbk is the address of the seven-word timer block described above. Note that this time block you specify in the .CTIMIO call must be the same one already used by the corresponding .TIMIO request.

The .CTIMIO macro expands as follows:

```
.CTIMIO
JSR      R5,@$TIMIT           ;POINTER AT END OF HANDLER
.WORD    tbk - .
.WORD    1                    ;CODE FOR .CTIMIO
```

Note that if a job aborts and your handler is entered at its abort entry point, you must immediately cancel any outstanding timer requests. However, if a timer completion routine has already been entered, you must wait for it to execute.

1.6.3 Device Timeout Applications

Device timeout support is used by RT-11 in only a few instances. However, there are a number of conditions in which timer requests are appropriate. If you are writing a handler for your own device, consider the following sections to determine whether or not timer requests would be useful to you.

1.6.3.1 Multiterminal Service

The resident multiterminal service in RT-11 that supports DZ11 and DZV11 modems uses device timeout to check the status of remote dial-up lines. The bootstrap starts up a polling routine to check each modem for a change in status. If a change occurs, the terminal service takes the appropriate action: it either recognizes a new line or disconnects a line when carrier is lost. Finally, the polling routine issues a .TIMIO call to start a half-second timer. The timer completion routine restarts the polling routine after a half-second elapses.

1.6.3.2 Typical Timer Procedure for a Disk Handler

A disk handler could implement a timer procedure for any disk operation. The purpose of the timer routine is to cancel or restart any operation that takes too long. If an operation does not complete within a reasonable amount of time, chances are good that a disk error of some sort occurred.

The handler's I/O initiation section sets a timer by using the .TIMIO call. Then the handler starts up the operation that a job requested: a read, write, or seek operation. The handler returns to the monitor with a RETURN instruction and waits for a device interrupt.

If an interrupt occurs before the time limit expires, the handler cancels the timer and performs its normal sequence of error checking on the results of the transfer. In general, the handler either drops to fork level to restart an incorrect operation, or exits to the monitor with .DRFIN to remove the current queue element.

If an interrupt does not occur within the time limit, the timer completion routine begins to execute. Its first action should be to simulate an interrupt. This action duplicates the handler environment after a genuine interrupt and makes sure that the stack has the necessary information. Then the timer completion routine acts as though the device interrupted but the transfer was in error. The timer completion routine simply branches to the correct section of code in the interrupt service section of the device handler to finish the processing.

The timer completion routine should use the following instructions to simulate an interrupt and enter system state:

```
MOV      @SP, -(SP)           ;MAKE ROOM ON THE STACK
CLR      2(SP)                ;FAKE INTERRUPT PS = 0
.MTPS    #340                  ;GO TO PRIORITY 7
.INTEN   0,PIC                 ;ENTER SYSTEM STATE
```

After the handler enters system state, it takes the appropriate action as a result of the timeout. The handler can try the operation again. To do this, it decrements the retry count, drops to fork level, and branches to the I/O initiation section. The code in the initiation section sets another timer, restarts the transfer, and returns to the monitor with a RETURN instruction to await another interrupt.

If the handler decides that the timeout indicates a serious error, one that should not be retried, this same procedure can be followed for a transfer whose retry count is used up. In this case, the handler sets the hard error bit in the Channel Status Word and then exits to the monitor with the .DRFIN call to remove the current queue element.

NOTE

Before a handler goes through the .DRFIN routine to remove the current queue element, it must cancel any timer request that has not yet expired.

1.6.3.3 Printer Handler Example

The extended example shown in Figure 1–2 consists of excerpts from a version of the RT–11 parallel interface printer handler modified to use timer support to check for the device off-line condition.

When the handler’s I/O initiation section starts up a transfer, it forces an immediate interrupt, which causes the handler’s interrupt service section to check the error bit in the CSR. If there is an error, control passes to the routine OFFLIN, which issues a .SYNCH call to enter user state, prints an error message on the console terminal, and then sets a 2-minute timer. The handler then returns to the monitor with a RETURN instruction and waits for the device to interrupt.

If the device interrupts, it means that the error condition has been corrected by an operator. The handler cancels the timer and checks the error bit once again to make sure there are no problems. If there is no error, the handler proceeds as usual. If there is an error, the handler loops back to the OFFLIN routine. If an interrupt does not occur within two minutes, the timer completion routine begins to execute. It prints an error message, sets another 2-minute timer, and returns to the monitor with a RETURN instruction to await an interrupt.

Figure 1–2: Printer Handler Example

```
; I/O INITIATION SECTION
      .DRBEG  LP
      MOV    LPCQE,R4          ;R4 POINTS TO CURRENT Q ENTRY
      ASL    6(R4)            ;WORD COUNT TO BYTE COUNT
      BCC    LPERR            ;A READ REQUEST IS ILLEGAL
      BEQ    LPDONE           ;SEEKS COMPLETE IMMEDIATELY
RET:   BIS    #100,@LPS       ;CAUSE AN INTERRUPT, STARTING TRANSFER
      RTS    PC

; INTERRUPT SERVICE SECTION
      .ENABL  LSB

      .DRAST  LP,4,LPDONE
      CLR    @LPS             ;DISABLE INTERRUPTS
      .FORK   FRKBLK
      TST    TICMPL           ;IS A TIMER ELEMENT ACTIVE?
      BEQ    1$              ;NO
      .CTIMIO TIMBLK         ;YES, CANCEL IT
      BCS    1$              ;ERROR
      CLR    TICMPL          ;AND DON'T DO IT AGAIN
1$:   MOV    LPCQE,R4        ;R4 POINTS TO CURRENT QUEUE ELEMENT
      TST    @(PC)+          ;ERROR CONDITION?
LPS:  .WORD  LP$CSR         ;LINE PRINTER STATUS REGISTER
ERROPT: BMI  OFFLIN         ;YES, HANG TILL CORRECTED
      .
      .
      .

; I/O COMPLETION SECTION
LPDONE: CLR    @LPS         ;TURN OFF INTERRUPT
      .DRFIN  LP
      .
      .
      .
```

Figure 1–2 (continued on next page)

Figure 1–2 (Cont.): Printer Handler Example

```

; PRINTER OFF LINE, PRINT WARNING EVERY 2 MINUTES

OFFLIN: MOV     LPCQE,R5             ;POINT TO QUEUE ELEMENT
        MOVB   Q$JNUM(R5),R5       ;GET JOB NUMBER OF CURRENT JOB
        ASR   R5                    ;SHIFT IT
        ASR   R5                    ; RIGHT
        ASR   R5                    ; 3 BITS
        BIC   #^C<16>,R5          ;ISOLATE JOB NUMBER
        MOV   R5,SYJNUM            ;SAVE IT FOR .SYNCH
        MOV   R5,TIJNUM            ;SAVE IT FOR .TIMIO
        .SYNCH SYNBK,PIC           ;GO TO USER STATE
        RTS   PC                   ;SYNCH FAILED, PUNT

1$:     CLR   TICMPL                ;INDICATE THAT WE GOT HERE
        TST  @LPS                  ;IS THERE STILL AN ERROR?
        BPL  2$                    ;NO, QUIT
        MOV  PC,R0                 ;AS COMPLETION ROUTINE, PRINT MESSAGE
        ADD  #MESSAG-. ,R0         ;POINT TO MESSAGE AS PIC
        .PRINT
        MOV  PC,R0                 ;IN A PIC WAY,
        ADD  #1$-. ,R0             ; POINT TO TIMIO COMPLETION ROUTINE
        MOV  R0,TICMPL             ;SAVE IT
        .TIMIO TIMBLK,0,2*60.*60. ;SET A 2-MINUTE TIMER
        RTS   PC

2$:     BIS   #100,@LPS            ;ENABLE INTERRUPTS
        RTS   PC                   ;RETURN LATER

TIMBLK: .WORD  0                   ;TIMER BLOCK: HI ORDER TIME
        .WORD  0                   ;LO ORDER TIME
        .WORD  0                   ;LINK
TIJNUM: .WORD  0                   ;JOB NUMBER
        .WORD  177700+LP$COD       ;SEQUENCE NUMBER
        .WORD  0                   ;MONITOR PUTS -1 HERE
TICMPL: .WORD  0                   ;ADDRESS OF COMPLETION ROUTINE
SYNBK:  .WORD  0                   ;SYNCH BLOCK
SYJNUM: .WORD  0                   ;JOB NUMBER
        .WORD  0,0,0,-1,0         ;OTHER
.FRKBLK: .BLKW 4                   ;FORK BLOCK
MESSAG: .ASCIZ  /?LP-W-LP off line - please correct/
        .EVEN
        .DREND LP

```

1.7 Error Logging

Error logging is an optional feature of RT–11 designed to help you monitor the reliability of your system. Device handlers that include support for error logging call the error logger after each I/O transfer. The error logger creates a historical record of the device's I/O activity that you can use to check its reliability.

You must perform a system generation to select error logging. Error logging is supported in all environments. If your system has the capability to run system jobs, the error logger runs as a system job; on FB systems, as an ordinary foreground job; on single-job systems, as a handler.

The system generation conditionals for error logging are as follows:

ERL\$G If this value = 1, it indicates that error logging is enabled for this system.

ERL\$S This condition defines the number of 256-word blocks to use for the internal logging buffer with single job monitors.

ERL\$U This represents the maximum number of individual device units for which the error logger collects statistics. The default value is 10, and the absolute maximum number is 30. Each unit adds seven words to the error logger. One slot is required for each unit. (For example, two slots are required for a system with an RK05 with two units.) Your response to a system generation dialogue question establishes the value of this variable.

You should consider your time and memory requirements before deciding to use error logging because error logging creates a certain minimal amount of overhead for each I/O transfer, and the error logger itself uses almost 2K words of memory. However, the error logger does not have to run constantly, so that the memory it requires can be made available to your programs when necessary, and calls that your handler makes to the error logger return immediately. The most efficient way to use the error logging system is as a check when you suspect device reliability problems, which means using it only when necessary.

The following sections describe how to implement error logging in your device handler and what information you should log. They also show you how to add headings for your device to the error reporting program. See the *RT-11 System Utilities Manual* for more information on the entire error logging system and how to use it.

All code in your handler that applies strictly to error logging should be placed inside conditional assembly directives. These directives should include the error logging code if the symbol **ERL\$G** is 1, and omit it otherwise. This way, the system parameters select whether or not the error logging code is included in the handler each time you assemble it.

1.7.1 When and How to Call the Error Logger

A handler calls the error logger after each I/O transfer, whether the transfer was successful or not. If the transfer was in error, the handler calls the error logger once for each retry of the transfer.

Since calls to the error logger must be serialized, the handler can issue them only during I/O initiation or following a **.FORK** call.

The handler must set up registers before it issues the call to the error logger. The register assignments for the three kinds of calls are described in the following sections.

1.7.1.1 To Log a Successful Transfer

Set up **R4** and **R5** as described below before calling the error logger after each successful transfer.

R5 must point to the third word (**BLKN**) of the current queue element.

R4 contains two bytes of information: the high byte is the device-identifier byte, *dd\$COD*; the low byte is -1 .

1.7.1.2 To Log a Hard Error

Set up R2 through R5 as described below before calling the error logger after a hard error has occurred. Generally, hard errors are those that are not recoverable. Examples of hard errors are device off line or not powered up, device write-locked, and so forth. Further, a soft error that has exhausted its allotted number of retries is considered a hard error.

R5 must point to the third word (BLKN) of the current queue element.

R4 contains two bytes of information: the high byte is the device identifier byte, *dd\$COD*; the low byte is 0.

R3 contains two bytes of information: the high byte contains the total number of retries allotted for this transfer; the low byte contains the number of device registers whose contents should appear in the error report.

R2 is a pointer to a buffer in the handler that contains the device registers to be logged.

1.7.1.3 To Log a Soft Error

Set up R2 through R5 as described below before calling the error logger after a soft error has occurred. Generally, soft errors are those that are recoverable and can possibly be corrected by retrying the transfer. Examples of soft errors include timing errors and hardware read or write errors.

Initialize a counter in your handler with the total number of retries allotted for each transfer. Decrement the count as each retry for a soft error is performed. When the count reaches zero, the error logger considers the error to be a hard error. On soft error, the error report prints a separate entry for each retry of a given transfer.

All retries are printed in the report even if the registers are identical. The report does not distinguish between hard or soft immediate errors. It prints only the contents of the registers at the time of the error and the value of the retry count. An immediate hard error can be recognized in the output since it will appear with a retry count of 0 with no immediately previous errors on that device and unit (with a retry count greater than 0).

R5 must point to the third word (BLKN) of the current queue element.

R4 contains two bytes of information: the high byte is the device identifier byte, *dd\$COD*; the low byte is the current value of the retry counter. (This value should decrease with each retry until it reaches 0, at which point the error is considered a hard error.)

R3 contains two bytes of information: the high byte contains the total number of retries allotted for this transfer; the low byte contains the number of device registers whose contents should appear in the error report.

R2 is a pointer to a buffer in the handler that contains the device registers to be logged.

1.7.1.4 Differences Between Hard and Soft Errors

The error logger itself does not differentiate between hard and soft errors and records the same information in both cases. However, by examining the report, you can determine if a hard error occurred, because a transfer that has exhausted all its retries will have records in the report for each of these retries, including one with a retry count of 0. It is therefore up to you to interpret the error.

In some circumstances, user-correctable errors, such as device off line or write-locked, should not call the error logger. Usually disk and tape hardware errors are the only ones reported, since these are the errors that reflect device reliability.

1.7.1.5 To Call the Error Logger

Once the required registers are set up, call the error logger as follows:

```
CALL @$ELPTR
```

\$ELPTR contains a pointer into the Resident Monitor. The .DREND macro allocates space in the handler for this pointer. The pointer is filled in at bootstrap time (for the system device) or at .FETCH or LOAD time (for a data device). If the error logger is not running, the monitor returns immediately to the handler. If the error logger is running, a link word in RMON contains its entry point. The following lines of code from RMON show how the call to the error logger is accomplished.

```
$ERLOG: MOV      (PC)+, -(SP)          ;ENTER HERE FROM HANDLER
;PUSH NEXT WORD ON STACK
$ELHND: .WORD    0                    ;0 IF ERROR LOGGER NOT RUNNING;
;ELSE CONTAINS ERROR
;LOGGER ENTRY POINT
        BNE     1$                    ;BRANCH IF LOADED
        TST    (SP)+                  ;PURGE STACK
1$:     RTS     PC                     ;INVOKES ERROR LOGGER OR
;RETURNS TO HANDLER
```

On return from the error logger call, R0 through R3 are preserved and R4 and R5 are indeterminate.

1.7.2 Error Logging Examples

See the handler listings in Appendix A for examples of error logging.

1.7.3 How to Add a Device to the Reporting Program

After you implement error logging in your device handler, the next step is to modify the reporting system so that the name of your device will appear in the report headings and the registers will be printed properly. The file ERRTXT.MAC contains the information for report headings for the devices supported by the RT-11 error logging reporting utility ERROUT. To include your device, edit this file, reassemble it, and relink it.

Use the following commands to reassemble and relink ERRTXT.MAC:

```
.MACRO/LIST ERRTXT
.LINK ERROUT, ERRTXT
```

ELBLDR Macro

Use the ELBLDR macro to add a new device to the error log reporting system. Edit the file ERRTXT.MAC to add the ELBLDR macro call for your device. The format of the call is as follows:

ELBLDR *xx*,<type>,C1,C2,<C3>

- xx* is the device-identifier byte, *dd*\$COD, that you specified in the .DRDEF macro. It must be a value between 0 and 377 octal.
- type* is any ASCII string you want to print on the report as the device type. It can be up to 59 characters long. Remember to enclose it in angle brackets.
- C1* is one of the two strings *DISK* or *TAPE*. It identifies the device general classification.
- C2* is the 2-character device name. You must specify exactly two characters.
- <*C3*> is a list of device register mnemonics (minus the first two characters) representing the registers that the handler logs. Separate the mnemonics with commas. Remember to use the angle brackets (<>).

Assembly errors result if you do not specify the parameters to ELBLDR correctly.

None of the parameters for the ELBLDR call is optional.

For example, the ELBLDR call for the RK handler is as follows:

```
ELBLDR 0 ,<RK11/RK05> ,DISK ,RK ,<DS ,ER ,CS ,WC ,BA ,DA ,DB>
```

This example shows that the device is the RK11/RK05 disk, its 2-character name is RK, its device-identifier byte is 0, and the registers its handler logs are RKDS, RKER, RKCS, RKWC, RKBA, RKDA, and RKDB.

The default input file name for ERRROUT is ERRLOG.DAT. However, you can save previous ERRLOG.DAT files by renaming or copying them. Thus, ERRROUT can operate on any file with the same format as ERRLOG.DAT. The name is not important; the format is. The internal format of the data in this file is documented in the *RT-11 Volume and File Formats Manual*.

1.8 Special Functions

Handlers use special functions to perform device-specific actions for which there are no corresponding RT-11 programmed requests. Chapter 2 describes those special functions supported by the distributed RT-11 device handlers.

The .SPFUN programmed request initiates special functions. When a program issues a .SPFUN request, it supplies a special function code as one of the arguments. It is the handler's responsibility to process the special function.

1.8.1 .SPFUN Programmed Request

The format of the .SPFUN programmed request is as follows:

```
.SPFUN area,chan,func,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]
```

See the *RT-11 System Macro Library Manual* for a description of the .SPFUN programmed request. See Chapter 2 for many special function examples within distributed handlers.

To use special function calls in your handler, you define the interface between the programmed request and the device handler. Thus, the meanings of the *buf*, *wcnt*, and *blk* parameters depend on the particular special function the request invokes; their meaning is dependent on the handler.

Note, however, the following:

- Although the monitor checks to make sure that *buf* is a valid address within the job area, it does not make sure that *buf* plus *wcnt* is still within the job area. It is therefore your responsibility to specify valid values if you use the .SPFUN request to transfer data.
- When using a mapped monitor and therefore a virtual address for *buf*, the buffer address must be mapped before the request is issued. Once the request is issued and the EMT returns, address translation has been performed and the buffer address can be unmapped. In the case of a read (input) operation, if the buffer address is subsequently unmapped, the address must be remapped before data can be accessed from the buffer.
- As previously mentioned, the *buf*, *blk*, and *wcnt* parameters can have any meaning that is supported by the particular handler. You could, therefore, pass an address as an argument.

Of those parameters, the RT-11 monitor performs address translation for only *buf*. Therefore, if you pass a mapped address in *blk* or *wcnt*, you must not unmap that address while the request is outstanding or active; that is:

- For nonwait, noncompletion I/O, until a .WAIT request succeeds on the channel.
- For wait mode I/O, until the request returns.
- For completion mode I/O, until the completion routine is entered.

If the special function call is to return a single value, *buf* should be a one-word buffer area. You are free to interpret *wcnt* and *blk* as anything you choose. They can be specification words of some sort, pointers to more buffers, and so on, as long as the handler interprets them according to the special function code. Note that the monitor does not alter these values in any way when it passes them to the handler. For example, it does not change the word count from positive to negative.

1.8.2 How to Support Special Functions in a Device Handler

Do the following to implement support for special function calls in your handler:

- Specify `SPFUN$` as one of the bits in the `.DRDEF stat` parameter argument. This indicates that the handler can accept special functions.
- Use the `.DRSPF` macro to list the supported special functions.
- Define symbols in the handler to represent the types of special functions the handler can perform. For example, the `DY` diskette handler defines the following special function codes:

```
SIZ$FN = 373           ;GET DEVICE SIZE
WDD$FN = 375           ;WRITE WITH DELETED DATA MARK
WRT$FN = 376           ;WRITE ABSOLUTE SECTOR
RED$FN = 377           ;READ ABSOLUTE SECTOR
```

Note that all special function codes must be negative byte values (that is, they must be in the range 200 through 377₈). Consult Chapter 2 for those symbols and codes already defined by RT-11. For the sake of consistency across devices, it is advisable to have each special function code represent the same operation on all devices. So, check first to see if a code for your function already exists and use it if it does. If there is no existing code for your particular function, assign codes starting with 200 and work toward 377 from there. (For extended device-unit handlers, the range is 360–377.) This policy should avoid conflicts with future RT-11 codes.

When the handler is entered for an I/O transfer, it should check the fourth word of the queue element to see if this is a request for a special function. `Q.FUNC`, which is the low byte of the fourth word of the I/O queue element, contains the special function code. On standard I/O requests for read, write, and seek operations, this byte is 0. For special function calls, this value is the negative special function code. Ignore any special function code that is not valid for your device.

If this is a request for a special function, the handler should initiate that function and return with a `RETURN` instruction. In the interrupt service section the handler should, as usual, check for errors and determine whether the operation is complete. The handler returns either data or words of status information to the calling program in the user buffer.

Since you are implementing the special functions for a particular device, you can establish the calling convention for that function in the `.SPFUN` programmed request as well as the return convention from the handler. Be sure the handler treats the arguments appropriately for each different special function call.

For a good example of a handler that implements special functions, see the `DX` handler in Appendix A.

1.8.3 Variable Size Volumes

A handler can control a device that permits volumes with two or more different sizes to be used. Examples of such handlers are the `DM` handler—which can service both `RK06` and `RK07` disks through a single controller—and the `DY` handler—which can service either a single-density or a double-density diskette in a single device unit. A handler for a device that supports volumes of different sizes should pass the size, in

blocks, of the smallest volume in the *size* parameter of the `.DRDEF` macro. This is the value that is returned to a running program when it issues the `.DSTAT` programmed request.

If it is important that a running program know the size of the volume that is currently mounted, the program can issue a special function to return the volume size. The handler must be able to respond to the request by returning the actual volume size in a one-word buffer area. The handler must have implemented support for special functions, as described above. The standard special function code for returning the actual volume size is 373.

1.8.4 Bad-Block Replacement

If your handler is to support bad-block replacement (BBR) by using a replacement table in the home block, you must implement the BBR special function codes as they are implemented for the DL and DM handlers. See Chapter 2 for more information.

1.9 Devices with Special Directories

The RT-11 monitor can interface to file-structured devices having nonstandard (that is, non-RT-11) directories. Magtapes are an example of special devices. Their handlers set bit 12 (SPECL\$) of the device status word. The USR processes directory operations for RT-11 directory-structured devices; for special devices, the handler must process directory operations such as `.CLOSE`, `.DELETE`, `.LOOKUP`, `.ENTER`, `.RENAME`, `.PURGE`, informational (`.GFxxx`, `.SFxxx`, and `.FPROT`), and `.CLOSZ`, as well as data transfers. See the *RT-11 System Macro Library Manual* for information on those requests.

The monitor requests a special directory operation by placing a positive, nonzero value in the function code byte (Q.FUNC) of the queue element. The positive function codes are standard for all devices. The symbol names are defined in the distributed file, `SYSTEM.MLB`, and are as follows:

Code	Name	Function
1	CLOS	Close
2	DELE	Delete
3	LOOK	Lookup
4	ENTR	Enter
5	RENM	Rename
7	INFO	<code>.GFxxx</code> , <code>.SFxxx</code> , and <code>.FPROT</code> operations
10	CLOZ	Close with size operation

In a queue element for a special directory operation, word 5 (Q.BUFF) of the queue element contains a pointer to the file descriptor block containing the device name, file name, and file type in Radix-50.

Software errors (such as file not found, or directory full) occurring in special directory device handlers during directory operations are returned to the monitor, processed, and appear in byte 52 as the standard, documented error codes. Hardware errors are returned in the usual manner by setting bit 0 in the Channel Status Word pointed to by the second word of the queue element.

Programmed requests for directory operations to special directory devices are handled by the standard programmed requests. When a .LOOKUP is issued, for example, the monitor checks the device status word for the special device bit. If the device has a special directory structure, the proper function code is inserted into the queue element and the element is directly queued to the handler, bypassing any processing by the USR. Device independence is maintained, since .LOOKUP, .ENTER, .CLOSE, and .DELETE operations are transparent to the user.

For a special device .LOOKUP, the file length is returned in word 6 of the queue element (Q.WCNT). For a .ENTER, word 6 returns the length of the new file.

1.10 Device Handlers in Mapped Systems

Device handlers for unmapped system environments require a few changes to work properly in mapped systems. Before describing the environment for a handler in a mapped system, the following sections outline the nomenclature conventions. The final sections explain how a handler communicates with a user buffer in extended memory.

1.10.1 Naming Conventions and the System Conditional

When you write a device handler, write a common source file called *dd*.MAC, where *dd* is the 2-character device. That source file is then assembled with the correct monitor conditional file such as XM.MAC and the system generation conditional file, such as SYSGEN.CND. This procedure ensures that the system generation features that the handler supports match those of the monitor.

The system generation conditional that represents extended memory support is MMG\$T, which has a value of 0 if extended memory support is not selected and a value of 1 if extended memory support is selected. The system conditional MMG\$T is correctly set in the distributed monitor conditional files. This means that the extended memory code is only assembled when the value of the conditional MMG\$T is 1. The assembly produces *ddX*.OBJ for mapped systems, or *dd*.OBJ for unmapped systems.

All code in your handler that applies strictly to memory management support should be placed inside conditional assembly directives. These directives should include the memory management code if the symbol MMG\$T is 1, and omit it otherwise. This way, the system parameters select whether or not the memory management code is included in the handler each time you assemble it.

1.10.2 Mapped Monitor Environment

In a mapped monitor system, at least the handler's root must reside within the low 28K words of physical memory. Typically the entire handler is written to reside in low memory.

The distributed mapped monitors support the .FETCH request, so usually your handler need not be continually loaded in memory. All Digital-supplied handlers for mapped monitors are fetchable with the exception of those few listed in the handler restrictions section of the *RT-11 System Release Notes*.

When handlers are entered, they run with kernel mapping, which permits access to the lower 28K words of memory plus the device I/O page (see Chapter 3 in the *RT-11 System Internals Manual*). The program that requests the I/O transfer, however, need not have the same mapping as kernel mapping. In fact, the program can fall into one of three valid categories:

- A privileged job whose mapping is identical to kernel mapping.
- A privileged job that maps to physical memory addresses above 28K words.
- A virtual job or completely virtual job with any kind of mapping.

Just as RT-11 supplies macros to ease the writing of parts of a device handler, so too does it provide monitor routines that simplify managing mapped systems. RT-11 distributes subroutines that perform the address conversion for you.

The program requesting an I/O transfer supplies a 16-bit virtual buffer address in the programmed request, although that portion of the user's virtual addressing space may be mapped somewhere else in physical memory. The handler must therefore find the actual 18- or 22-bit physical address of the user data buffer before moving information to it or from it. The monitor verifies that the user buffer area occupies contiguous locations in physical memory.

The fact that in a mapped system, locations in physical memory are expressed as 18- or 22-bit addresses, is important when you need to specify an address within the handler itself as a buffer address. If, for example, the handler contains a string of zeroes that it writes to a device as part of initialization, the handler sets up the device write operation, specifying the address of the string in the handler as the buffer address. Since the handler is located within the lower 28K words of physical memory, its physical address can be expressed as its virtual 16-bit address plus extra mapping bits (bits 16 and 17 of an 18-bit address, or bits 16-21 of a 22-bit address), which must be 0.

The *RT-11 System Internals Manual* describes memory mapping in detail.

The RT-11 monitor provides routines for handlers to use to access the real user data buffer in physical memory. The following sections describe these routines and the situations in which they are useful.

1.10.3 Address Translation

RT-11 provides the following two routines for performing address translation for the address passed in Q.BUFF.

- **\$MPMEM**

Call \$MPMEM to return the physical address to be used for MOV operations.

- **\$MPPHY**

Call \$MPPTR (which in turn points to \$MPPHY in RMON) to perform address translation for I/O DMA operations.

1.10.3.1 \$MPMEM Routine

The \$MPMEM subroutine uses queue element offsets Q.MEM (and Q.BUFF) to perform the PAR1 offset mapping.

\$MPMEM is located at an address 22(decimal) bytes below the entry address of monitor routine \$P1EXT. \$P1EXT is pointed to by RMON fixed offset P1\$EXT (432).

Before the call, R5 must point to Q.BUFF, the fifth word of the queue element.

On return from the call:

- The first word of the stack, (SP), contains the low-order 16 bits of the physical buffer address.
- The second word of the stack, 2(SP), contains the high-order bits of the physical buffer address. The bit positions for an 18-bit address are 4 and 5; those for a 22-bit address are 4 through 9.

The following code fragment illustrates using \$MPMEM. (In code preceding the fragment, R4 was pointed to Q.BLKN, the third word in the queue element.)

```
MOV    @#$SYPTR,R3      ; Get start of RMON
MOV    P1$EXT(R3),R3    ; R3 --> $P1EXT
MOV    R4,R5           ; Make R5 --> 5th word (Q.BUFF) of
CMP    (R5)+,(R5)+     ; queue element
CALL   $MPMEM(R3)      ; Map KT-11 virtual to physical
MOV    (SP)+,R2        ; R2 = low 16 bits physical address
MOV    (SP)+,R3        ; R3 = high 2 (or 6) bits physical
                          ; address
      .
      .
      .
```

See also Sections 1.10.6 and 1.10.7.2.

\$MPMEM uses Q.MEM rather than Q.PAR because in the case of UMR on UNIBUS processors, the value stored in Q.PAR can diverge from the value stored in Q.MEM.

1.10.3.2 \$MPPHY Routine

Call the \$MPPHY routine to find the user buffer in physical memory to perform DMA I/O operations. \$MPPHY uses the Q.PAR and Q.BUFF queue element offsets to create the correct 18- or 22-bit address for the user buffer.

The format of the call for the \$MPPHY routine is as follows:

CALL @\$MPPTR

\$MPPTR contains a pointer to the \$MPPHY routine in the Resident Monitor. The .DREND macro allocates space for this pointer at the end of the handler. The pointer is filled in at bootstrap time (for the system device) or at LOAD time (for a data device).

Before the call:

R5 must point to Q.BUFF, the fifth word in the queue element.

After the call:

(SP), the first word on the stack, contains the low-order 16 bits of the physical buffer address.

2(SP), the second word on the stack, contains the high-order bits of the physical buffer address in bit positions 4 and 5, if it is an 18-bit address, or in bit positions 4 through 9, if it is a 22-bit address.

R5 points to Q.WCNT, the sixth word in the queue element. The value is not changed.

The following example is from the RK handler.

```

      CMP      (R5)+,(R5)+          ;Advance to bufr addr in queue elt
      CALL    @$MPPTR              ;Convert user virtual addr to physical
      MOV     (SP)+,-(R4)          ;Put low 16 bits in RKBA,
                                   ; High bits on stack
      MOV     (R5)+,-(R4)          ;Put word count into RKWC
      BEQ     7$                   ;0 Count = SEEK
      BMI     5$                   ;Negative = WRITE, So
                                   ; all set up
      NEG     @R4                  ;Positive = READ,
                                   ;Fix count for controller
      MOV     #CSIE!FNREAD!CSGO,R3 ;Function is READ
5$:    BIS     (SP)+,R3            ;Merge high order address
                                   ; bits into function
      MOV     R3,-(R4)            ;Start the operation
6$:    RTS     PC                  ;Await interrupt
```

1.10.4 Character Devices: \$GETBYT and \$PUTBYT Routines

The handlers for character-oriented devices, such as printers, must transfer the data from the device to the user buffer area themselves. The transfer is usually one byte at a time. The device itself uses registers in the I/O page to store one character at a time. The handler can use two monitor routines—\$GETBYT and \$PUTBYT—to move data between the I/O page and the user buffer area.

1.10.4.1 \$GETBYT Routine

A handler can use the \$GETBYT monitor routine to move a byte from the user buffer in physical memory to the stack. The handler can then move the character into the device data buffer register in the I/O page and initiate an I/O transfer.

The format of the call for the \$GETBYT routine is as follows:

CALL @\$GTBYT

\$GTBYT contains a pointer to the \$GETBYT routine in the Resident Monitor. The .DREND macro allocates space for this pointer at the end of the handler. The pointer

is filled in at bootstrap time (for the system device) or at LOAD time (for a data device).

Before the call:

R4 must point to *Q.BLKN*, the third word in the queue element.

After the call:

(*SP*), the first word on the stack, contains the next byte from the user buffer in the low byte. The contents of the high byte are not defined.

R4 is unchanged.

The following example from the XL handler shows how the handler gets a byte from the user buffer and outputs it.

```
GNXTCH: MOV      XOCQE,R4          ; R4->current output queue element
         BEQ      10$              ; None available...
         .
         .
         .
         TST      Q$WCNT(R4)       ; Any characters left to output?
         BEQ      20$              ; Nope, this request is complete
         INC      Q$WCNT(R4)       ; Yes, now there is one less to do
         CALL     @$GTBYT          ; Get the byte to output
         MOV      (SP)+,R5
```

The buffer address (*Q.BUFF*) in the queue element is updated by 1. If a mapping overflow occurs, the monitor routine subtracts 100 from the value in *Q.BUFF* and adds 1 to the value in *Q.PAR* and *Q.MEM*. Mapping overflow occurs if *Q.BUFF* is 20100 or more.

1.10.4.2 \$PUTBYT Routine

After a successful data transfer, a handler can get a character from the device data buffer register in the I/O page and push it onto the stack. It can then use the \$PUTBYT monitor routine to move a byte from the stack to the user buffer in physical memory.

The format of the call for the \$PUTBYT routine is as follows:

CALL @\$PTBYT

\$PTBYT contains a pointer to the \$PUTBYT routine in the Resident Monitor. The .DREND macro allocates space for this pointer at the end of the handler. The pointer is filled in at bootstrap time (for the system device) or at LOAD time (for a data device).

Before the call:

R4 must point to *Q.BLKN*, the third word in the queue element.

The byte to transfer to the user buffer must be on the top of the stack. The character must be in the low byte of the stack's first word. The high byte is unpredictable.

After the call:

The word containing the character to transfer is removed from the stack.

R4 is unchanged.

The buffer address (Q.BUFF) in the queue element is updated by 1. If a mapping overflow occurs, the monitor routine subtracts 100 from the value in Q.BUFF and adds 1 to the value in Q.PAR and Q.MEM. Mapping overflow occurs if Q.BUFF is 20100 or more.

The following example from the XL handler shows how the handler gets a character and moves it to the user buffer.

```
30$:  
.  
.  
.  
    MOVB    R5,-(SP)           ; Put character here for PUTBYT  
    CALL    @$PTBYT           ; Call the routine  
    DEC     Q$WCNT(R4)        ; Is transfer complete? (z-bit=1 if so)
```

1.10.5 Any Device: \$PUTWRD Routine

The monitor routine, \$PUTWRD, is similar to \$PUTBYT, except that \$PUTWRD moves a word to the user buffer in physical memory instead of a byte. This routine is useful when the handler needs to transfer a word of status information to the user buffer, rather than a data character from a device. Handlers for any kind of device can use \$PUTWRD.

The format of the call for the \$PUTWRD routine is as follows:

CALL @\$PTWRD

\$PTWRD contains a pointer to the \$PUTWRD routine in the Resident Monitor. The .DREND macro allocates space for this pointer at the end of the handler. The pointer is filled in at bootstrap time (for the system device) or at LOAD time (for a data device).

Before the call:

R4 must point to Q.BLKN in the queue element.

The word to transfer to the user buffer must be on the top of the stack.

After the call:

The word to transfer is removed from the stack.

R4 is unchanged.

The buffer address (Q.BUFF) in the queue element is updated by 1. If a mapping overflow occurs, the monitor routine subtracts 100 from the value in Q.BUFF and adds 1 to the value in Q.PAR and Q.MEM. Mapping overflow occurs if Q.BUFF is 20100 or more.

The following example from the DY handler shows the handler responding to a special function call that requests the size of the currently mounted volume. In this

case, the larger of two possible diskettes is mounted. The handler uses \$PUTWRD to move the size of the volume to the user buffer area.

```
MOV      #DDNBLK, -(SP)          ;Push size in blocks onto stack
MOV      DYCQE, R4              ;Point R4 to Q.BLKN
CALL     @$PTWRD                ;Call the routine
```

1.10.6 Mapping Directly to the User Buffer

Some situations call for combinations of the procedures described in the previous sections. Others require more effort on the handler's part to accomplish a transfer. Some handlers cannot make good use of monitor routines and must access the user buffer directly.

The DM handler for the RK06 disk, for example, normally uses the \$MPPHY monitor routine to convert mapped addresses to physical addresses. However, when a Cyclic Redundancy Check (CRC) error occurs, the handler performs its own mapping to the user buffer and then applies the correction for the error before continuing the transfer. The procedure for a handler to map to the user buffer is as follows.

Devices such as the RX01 diskette transfer data one sector at a time between the disk itself and an internal disk data buffer called a silo. Monitor routines for character-oriented devices available to a silo device are too slow for the RX01. So, the handler for the RX01 diskette maps to the user buffer in physical memory and then performs the I/O operation as though it were a simple transfer between memory and the device. The handler implements this mapping by borrowing kernel PAR1.

The handler does this mapping through kernel PAR1. Handlers map to the user buffer through the monitor routine \$P1EXT¹.

\$P1EXT copies from the handler to the monitor stack the instructions necessary to transfer the data, thereby removing the instructions from possible PAR1 space. \$P1EXT next sets the proper PAR1 value and then executes the instructions copied to the stack. When finished, \$P1EXT restores PAR1, clears the monitor stack, and returns to the handler at the word following the instruction list. Upon return, all registers are unchanged except as modified by the instruction list.

Call the routine \$P1EXT with a JSR R0 followed by a word containing the number of bytes+2 to copy to the monitor stack, a series of instructions to perform the data transfer, and the PAR1 value (Q.PAR) from the queue element. The following instructions from the DX handler illustrate this technique. R1 is the byte count to transfer, R2 points to the user buffer, R4 points to the RX01 CSR, and R5 points to the RX01 data register. P1\$EXT is a monitor fixed offset containing a pointer to the routine \$P1EXT.

¹ Because all relevant code is executed outside the PAR1 area, the interrupt service in the PAR1 area is handled in mapped monitors by a vector forwarding technique that is transparent to the handler.

```

        MOV     @#SYSPTR,R4           ;R4 -> monitor base
        MOV     P1$EXT(R4),(PC)+     ;Get addr of externalization routine
$P1EXT: .WORD   P1$EXT               ;Pointer to externalization routine
        .
        .
;--- Remove two lines below if not memory management
        JSR     R0,$P1EXT           ;Let monitor execute the following code
        .WORD   PARVAL-           ;Number of bytes + 2 to copy
;---
2$:     TSTB    @R4                 ;Test transfer ready flag
        BPL     2$                 ;Wait till ready
3$:     MOVB   (R2)+,@R5           ;Move a char from user bufr to RX01
        TSTB    @R4                 ;Set CSR for next time
        DECB   R1                 ;Check transfer count
        BNE     2$                 ;If not 0, more to transfer

;--- If memory management, terminate list with PAR1 value
PARVAL: .WORD   0                 ;Remove if not memory management
;---

;Continue with normal processing from here on.

```

The following restrictions apply to the instruction list passed to \$P1EXT:

- No instruction in the list can reference any location in the handler, except for relative-address references within the list itself.
- The instruction list can use the stack for temporary storage, but it cannot remove any previous values from the stack or leave any values on the stack after it is done.
- If used in the instruction list, R0 must be saved and restored.
- Instruction lists of more than 32 words are not recommended because of stack space limitations.

If your handler must access the user buffer directly, it is important that you understand how PAR1 maps to the user area. Figure 1–3 shows a virtual job in a typical mapped system with the user buffer located in physical memory above the 28K-word boundary. The user program is mapped to the buffer through PAR6. The handler calls \$P1EXT, which borrows kernel PAR1, puts the Q.PAR value from the queue element there, and then uses the Q.BUFF value from the queue element to access the user buffer.

PAR1 maps to physical memory in units of 32-word decimal blocks and at most can map an area 4K words long. (Note that the page length of PDR1 is always set to map the entire page.) If the user buffer starts at a location in physical memory that is not an even multiple of 32 words, PAR1 maps to the first 32-word boundary below the start of the buffer. The PAR1 mapping area can start at any address in physical memory whose low-order two octal digits are 0. Thus, with a particular PAR1 mapping, as much as 4K words or 4K minus 31 decimal words of the user buffer will be mapped. Figure 1–4 shows how this mapping works.

Figure 1–4 shows a buffer area located at 331724 in physical memory with the application program mapped to the buffer through PAR6. The buffer is 24 octal

Figure 1–3: Device Handler Mapping to User Buffer Area

bytes above 331700, which is a 32-word boundary. \$P1EXT puts the Q.PAR value, 3317, into PAR1, replacing the default PAR1 value of 0200. This causes PAR1 to map to a 4K-word area in physical memory starting at address 331700. As a result, when the handler refers to kernel virtual addresses in the range 20000 through 37776, it accesses physical memory locations 331700 through 351676. Since the value in Q.BUFF is 20024, by using that value, the handler can access the start of the user buffer area at location 331724.

If the amount of data to be transferred is large, you may need to advance the buffer pointer and adjust the mapping to account for it. There are two ways to advance the buffer pointer. The easier way is to modify PAR1 as you go. For example, for every 32 words you advance through the buffer, add 1 to the PAR1 value and subtract 64 from the offset. The DX handler example just described transfers 64 words at a time, adding 2 to PAR1 (and subtracting 128 from the offset) after each transfer to avoid mapping overflow.

Another way to advance the buffer pointer is to modify the value of Q.BUFF by modifying the value in the queue element itself. To adjust the mapping, step through the following procedures, thinking in terms of 4K-word units. First, after you modify the value of Q.BUFF, compare the new value to 40000. If the value is greater than or equal to 40000, subtract 20000 from it, and add 200 to Q.PAR. These procedures take care of not only adjusting the mapping, but also avoid mapping overflow.

Figure 1–4: PAR1 Mapping

Finally, here are steps to follow to access any location in the user buffer area, if you are given a byte offset from the beginning of the buffer. Essentially, you must determine the number of 32-word units in the offset by dividing the 16-bit byte offset by 100 octal and adding the quotient to PAR1 and the remainder to Q.BUFF. Then you will be able to access the correct location in the buffer.

For example, suppose you needed to access the byte at offset 12345 from the start of the buffer shown in Figure 1–4. Dividing 12345 by 100 yields a quotient of 123 and a remainder of 45. Adding 123 to the current value of Q.PAR, which is 3317, yields 3442 for the new PAR1 value. Adding 45 to the value of Q.BUFF, which is 020024, gives 020071 as the new buffer address. (Note that this is a byte address.)

1.10.7 Extended Memory Subroutines

This section describes a set of subroutines that allow you to perform the following extended memory operations:

- Move data from one place to another in extended memory.
- Obtain a specified amount of memory from the free memory list maintained by RT–11.

- Find a specified global region.
- Convert a user virtual address into a 22-bit physical address.

The entry points for the subroutines that perform these operations are located directly below P1EXT.

1.10.7.1 Converting a Virtual Address into a Physical Address (\$JBREL)

The \$JBREL subroutine returns the physical address that corresponds to a virtual address for the job number you supply. Your program must be in Kernel mode when it calls \$JBREL. If your program is in User mode, use the .CALLK request to transfer control to Kernel mode.

\$JBREL is located at an address 26_{10} bytes below the entry address of monitor routine \$P1EXT. \$P1EXT is pointed to by RMON fixed offset P1\$EXT (432).

You supply a job number and virtual address to \$JBREL in the following registers:

Register	Contents		
R0	The virtual address to be translated into a physical address.		
R1	The job number, addressing mode, and space-type (instruction or data) for which the virtual address applies. You can determine the job's number from the .GTJB request. R1 contains the following information, none of which is validated for accuracy:		
	Bits	Value	Meaning
	0	0	Reserved
	1–3	0–7	Job Number
	4–7	0	Reserved
	8–9		Addressing mode:
		00	User
		01	Supervisor
		10	Reserved
		11	Reserved
	10		Address space:
		0	Data space (if enabled)
		1	Instruction space
	11–15	0	Reserved
R3	The size in 32-word chunks.		

\$JBREL passes the job number and virtual address to the monitor. The monitor performs the address translation and returns to \$JBREL. If the specified virtual address is not mapped to a virtual job, the equivalent kernel-mapped address is returned.

On return, if the carry bit is clear, \$JBREL provides the following information:

Register	Contents
R1	The PAR1 relocation bias.
R2	The PAR1 displacement.
R3	The amount of contiguously mapped memory that begins at the returned PAR1 bias and displacement, in 32-word chunks. If the value returned is less than that specified in R3 as input to \$JBREL. The V-bit (overflow) is set; otherwise it is cleared.

If carry is set on return, R1 and R2 contain random data.

The following example code assumes you are running in User mode and, therefore, require the .CALLK request to transfer control to virtual mapping in Kernel mode:

```
.MCALL .CALLK, .PRINT, .EXIT
.LIBRARY "SYSTEM.MLB"
.MCALL .SYCDF, .FIXDF, .PIXDF
.NLIST BEX

MMG$T   =: 1

.SYCDF           ; Define system logicals:
.FIXDF           ; $SYPTR - base of fixed area
.PIXDF           ; P1$EXT - offset of $P1EXT
               ; $CJVPT - routine offset from $P1EXT

VIRTAD   =: 0      ; Virtual address to be translated
JOBNUM   =: 16    ; Job number of virtual address

START:   MOV      #VIRTAD,R0      ; Virtual address to translate
         MOV      #JOBNUM,R1     ; Job number for translation
               ; virtual address is user mode
               ; and data space (if enabled)
         MOV      #5,R3          ; Check that 5 64-byte chunks
               ; are contiguously mapped
         MOV      @#$SYPTR,R2    ; R2 = RMON Base
         MOV      P1$EXT(R2),-(SP) ; Stack pointer to $P1EXT routine
         ADD      #$CJVPT,@SP    ; Make it point to $JBREL for .CALLK
         .CALLK                    ; Enter KERNEL mode
               ; Execute $JBREL
               ; Return to USER mode
         BCS     10$             ; Branch if error occurred
         BVS     20$             ; Branch if less than 5 64-byte
               ; chunks are contiguously mapped
         MOV     R1,PAR1BS      ; Store returned PAR1 value
         MOV     R2,PAR1OF     ; Store returned PAR1 offset
         BR      DONE          ; Branch to program exit

10$:    .PRINT   #ERROR1      ; Report the error
         BR      DONE          ; Branch to program exit

20$:    .PRINT   #ERROR2      ; Report the error
DONE:   .EXIT                    ; Done with example

PAR1BS: .WORD   0              ; Physical address's PAR1 value
PAR1OF: .WORD   0              ; Physical address's PAR1 offset
ERROR1: .ASCIIZ /Error: Check for invalid job number./
ERROR2: .ASCIIZ /Error: Not all of requested address block is /
         .ASCIIZ /contiguously mapped./
```

1.10.7.2 Moving Data Within Extended Memory (\$BLKMOV)

The \$BLKMOV subroutine moves the contents of memory from one place in 22-bit physical memory to another. The entry point is \$P1EXT-2.

In the following example, R0 contains the address of \$P1EXT, and BLKMOV equals -2. \$BLKMOV moves the data from the specified input buffer to the specified output buffer.

```
MOV      #input_buffer_par1,R1
MOV      #input_buffer_par1offset,R2
MOV      #output_buffer_par1,R3
MOV      #output_buffer_par1offset,R4
MOV      #word_count,R5
CALL     BLKMOV(R0)
```

1.10.7.3 Obtaining Free Memory (XALLOC)

The XALLOC subroutine obtains a specified amount of memory from the free memory list maintained by RT-11. The size argument passed in R2 is in units of 32_{10} words. To allocate 32000_{10} words, specify 1000. as the size passed to R2. The entry point for the subroutine is \$P1EXT-6.

In the following example, R0 contains the address of \$P1EXT, and XALLOC equals -6.

```
MOV      #required_size,R2
CALL     XALLOC(R0)
```

If the required amount of memory is not available, the carry bit will be set on return. In this event, R2 contains the size of the largest amount available.

If the required amount of memory is available, the carry bit will be reset on return. In this event, the memory has been removed from the free list, and R1 contains the region address divided by 32_{10} .

XALLOC uses R3 and destroys the contents of this register.

1.10.7.4 Returning Memory to the Free List (XDEALC)

The XDEALC subroutine returns a specified section of extended memory to the free memory list maintained by RT-11. The entry point for XDEALC is \$P1EXT-18₁₀. \$P1EXT is pointed to by RMON fixed offset P1\$EXT (432).

The address and size of the section of extended memory to be returned are specified in units of 32_{10} words. Load R1 with the starting address divided by 32_{10} and R2 with the size of the region in units of 32_{10} words.

In the following example, R0 contains the address of \$P1EXT, and \$XDEPT is -18₁₀.

```
MOV      #region_address,R1            ; Address in units of 32. words
MOV      #region_size,R2              ; Size in units of 32. words
CALL     $XDEPT(R0)
```

On return from XDEALC, the carry bit is clear if the memory was returned. If the carry bit is set, the memory was not returned because the free memory has become too fragmented.

XDEALC destroys the contents of R1 and R2. If you want to preserve the contents of those registers across the call, you must save them.

1.10.7.5 Finding a Global Region (FINDGR)

The FINDGR subroutine finds a global region that has a specific name. The entry point for this subroutine is \$P1EXT-10.

In the following example, R4 contains the address of \$P1EXT, and FINDGR equals -10_{10} .

```
MOV      #rad50_name_area,R5
CALL    FINDGR(R4)
```

where *rad50_name_area* is the address of a 2-word area containing the RAD50 name of the region to search for.

If the specified region is found, the carry bit is clear on return. In this event, R1 points to the size word of the associated global region control block.

If no region by the specified name is found, the carry bit is set on return. In this event, R1 points to the size word of the next available global region control block.

If no more global region control blocks are available, R1 is returned with a zero value.

1.10.7.6 Converting a Virtual Address into a Physical Address (\$USRPH)

The \$USRPH subroutine converts a user virtual address in the current running job into a 22-bit physical address.

NOTE

No job number is specified. Ensuring that the current running job is also the job for which the address translation is intended is quite difficult. Therefore, unless you have a very good reason for using this routine, Digital recommends you instead use the \$JBREL routine, for which you can specify the job number.

The entry point for this subroutine is \$P1EXT-14₁₀.

In the following example, R5 contains the address of \$P1EXT, and CVAPHY equals -14.

```
MOV      #virtual_address,R0
CALL    CVAPHY(R5)
```

On return, R1 will contain the high-order address bits, and R2 will contain the low-order address bits.

1.11 System Device Handlers and Bootstraps

In these sections, a description of monitor files precedes an explanation of how to create a system device handler or modify an existing handler to use as a system device. Within the main body of this explanation, details are given on the primary driver and on various bootstrap routines. The final sections provide background information on the DUP procedures for bootstrapping a new system device.

1.11.1 Monitor Files

A monitor file must reside on your system device and can have any name you choose, but its required file type is .SYS. If you create a monitor through the system generation process, its name is RT11xx.SYG. You must rename the monitor to .SYS before you use it.

Blocks 1 through 4 of each monitor file contain the secondary bootstrap. The secondary bootstrap loads the system device handler and the monitor into memory. It also modifies the monitor tables to connect the monitor with the device handler and assigns the default DK and SY names.

Each device handler that can be used as a system device handler has a special block of device-specific code in it called the *primary driver* that is used by the secondary bootstrap to read the system device handler file and the monitor file from the system device. The secondary bootstrap has room in its own block 0 to store the primary driver.

1.11.2 Creating a System Device Handler

To create a system device handler, you must add the primary driver to a standard handler for a data device. As described in the following sections, the .DRBOT macro does much of that work for you.

1.11.2.1 .DRBOT Macro

Use the .DRBOT macro to help you set up the primary driver. It also invokes the .DREND macro to mark the end of the handler so that the primary driver will not be loaded into memory during normal operations. In general, the code in the primary driver does not have to be Position-Independent. However, any non-PIC reference must be expressed relative to *dd*BOOT::. Note also that locations 60₈ through 116₈ are not available for your use.

The format for the .DRBOT macro is as follows:

.DRBOT name,entry,read

name is the 2-character device name.

entry is the entry point of the software bootstrap routine.

read is the entry point of the bootstrap read routine.

The .DRBOT macro puts a pointer to the start of the primary driver into location 62 of the handler file. It puts the length, in bytes, of the primary driver into location 64. The primary driver, including the error routine supplied by .DREND, must not

exceed 1000₈ bytes. Location 66 contains the offset from the start of the primary driver to the start of the bootstrap read routine.

Issue the .DRBOT macro call before the .DREND macro call. Then put the primary driver code between .DRBOT and .DREND, remembering that the primary driver must be one block or less in size—that is, it must be 1000₈ bytes long or less, including the error routine and the locations from 60₈ through 116₈. The .DREND macro is called twice in a system device handler: once by .DRBOT, and once when you use it at the very end of the primary driver. The first occurrence of .DREND closes out the nonsystem section of the device handler and sets up a table of pointers into the monitor, among other things. The second .DREND call, the one you issue yourself, creates the BIOERR bootstrap error routine, instead of repeating the pointer table.

If you use the BOOT command to bootstrap the new device, DUP passes the system unit number to the primary driver in location 4722 and in R0. If you bootstrap the device with a hardware bootstrap or some non-RT-11 utility program, the primary driver must determine the device unit number that was booted and save it in location 4722 and in R0.

1.11.2.2 Primary Driver

The primary driver you add to a standard handler for a data device consists of four parts:

- Entry routine
- Software bootstrap
- Bootstrap read routine
- Bootstrap error routine

The primary driver works together with the RT-11 bootstrap, BSTRAP, to boot the new system device. The primary driver is contained entirely within the p-sect *dd*BOOT, where *dd* is the 2-character device name. The code is loaded and executes, beginning at location 0 in physical memory.

For examples of the primary driver, see the handler listings in Appendix A.

1.11.2.3 Entry Routine

The entry point for the primary driver is *dd*BOOT::. This location must contain only two instructions, and these must follow the Digital standard bootstrap sequence. These instructions are a NOP and a branch to the start of the software bootstrap. If the start of the software bootstrap is too far away for a branch, you can branch to a JMP instruction that starts the software bootstrap. The entry routine for the RK handler is as follows (BOOT1 is defined in the primary driver):

```
RKBOOT:: NOP
          BR      BOOT1
```

Any hardware bootstrap causes the code in p-sect *dd*BOOT to load into memory at location 0. It also starts execution at *dd*BOOT::.

1.11.2.4 Software Bootstrap

The DUP utility executes the software bootstrap as the result of a jump or branch from the entry routine. Upon entry, all registers are available for use in the software bootstrap. The software bootstrap performs the following functions in the order shown:

1. Sets up the stack at location 10000.
2. Saves the number of the device unit from which the system was just bootstrapped. The method you use to find the unit number varies depending on the device; some unit numbers are passed in R0, and others must be extracted from the CSR. Save the unit number on the stack, and elsewhere in memory, if necessary.
3. Calls the bootstrap read routine to read in the rest of the bootstrap.
4. Puts a pointer in B\$READ to the bootstrap read routine.
5. Puts the Radix-50 value for "B\$DNAM" in B\$DEVN.
6. Stores the device unit number in B\$DEVU.
7. Jumps to B\$BOOT in RT-11's bootstrap to continue.

The software bootstrap should be located in the primary driver immediately below location *ddBOOT* + 664. (Locations 664 through 776 contain the error routine created by .DREND.)

1.11.2.5 Bootstrap Read Routine

The purpose of the bootstrap read routine (the primary bootstrap) is to read the volume in the device unit from which the system was just bootstrapped. It is called by both the RT-11 bootstrap (BSTRAP, the secondary bootstrap) and by DUP (the software bootstrap), as described in the previous section.

The interface through which the other routines pass information to the bootstrap read routine is as follows:

R0 contains the block number to read.

R1 contains the word count to read.

R2 contains the memory buffer address into which to store the data.

All registers are available for use in the bootstrap read routine, as is the stack.

The bootstrap read routine normally is a noninterrupt routine, used to read the volume according to the parameters passed in R0 through R2. On error, the routine should jump to BIOERR. If there are no errors, it should return with a RETURN instruction, with the carry bit clear.

The bootstrap read routine should be located in your primary driver at location *ddBOOT* + 120. (Location 120 is the lowest address at which the read routine can be located.)

1.11.2.6 Bootstrap Error Routine

The bootstrap error routine starts at location BIOERR::. The code in this routine is supplied completely by the .DREND macro, which you place at the end of the primary driver.

1.11.3 DUP and the Bootstrap Process

This section shows how DUP carries out three commands related to bootstrapping. The commands are as follows:

```
BOOT ddn:filnam
COPY/BOOT xxn:filnam ddm:
BOOT ddn:
```

1.11.3.1 BOOT *ddn:filnam*

Use the *BOOT ddn:filnam* command to perform a software bootstrap of a specific monitor file on a specific device. In the command line, *dd* represents the 2-character device name; *n* is its unit number. Both the new monitor file and the new device handler must be present on device *dd*.

As soon as this command is issued, DUP first checks that device *dd* is a random-access device. Next, it locates the monitor file *filnam.SYS* on the device. (The .SYS file type is both the default and the required file type.) Then DUP reads blocks 1 through 4 into a memory buffer. These blocks contain the secondary bootstrap for the monitor.

The next-to-last word in block 4 contains the suffix for the handlers associated with this monitor. DUP uses this to build the file name of the device handler, usually *dd.SYS* or *ddX.SYS*. DUP reads block 0 of the device handler file into a memory buffer, using the contents of locations 62 and 64 to locate the primary driver, and reads it into a memory buffer.

Next, DUP copies the primary driver into a buffer at the beginning of the secondary bootstrap, which is also in a memory buffer. It loads the information shown in Table 1–9 for the primary driver and the secondary bootstrap.

Table 1–9: DUP Information

Offset from Start of Memory Buffer	Contents
4722	Booted unit number (B\$DEVU)
4724–4726	Booted file name in Radix–50 (B\$DNAM)
5000	Date at which booted
5002–5004	Time at which booted

DUP then copies the primary driver and secondary bootstrap from the memory buffer into memory locations 0 through 5004. Then it jumps to location 1000 to start the secondary bootstrap at its DUP entry point so that the secondary bootstrap can load the monitor and the system device handler into memory.

Figure 1–5 illustrates the procedure.

Figure 1–5: BOOT *ddn:filnam* Procedure

1.11.3.2 COPY/BOOT *xxn:filnam ddm:*

Use the *COPY/BOOT xxn:filnam ddm:* to copy the secondary bootstrap from the monitor file on device *xx* to blocks 2, 3, 4, and 5 of device *dd*. In the command line, *xx* represents the device on which the monitor file is stored; *n* is its unit number; *dd* represents the 2-character name of the device that is to receive the bootstrap; *m* is its unit number.

As soon as this command is issued, DUP checks that devices *xx* and *dd* are random-access devices. Next, it locates the monitor file *filnam.SYS* on the *xxn:* device. It reads blocks 1 through 4 into a memory buffer. These blocks contain the secondary bootstrap for the monitor.

DUP locates the appropriate handler file on device *dd*. DUP then reads block 0 of the device handler file into a memory buffer, using the contents of locations 62 and 64 to locate the primary driver, and reads it into a memory buffer.

The handler for the system device *dd* must already be located on *dd* before you can copy the bootstrap to the device. DUP loads two words of Radix-50 for *filnam* into locations 4724 and 4726 of the memory buffer. Next, DUP copies the primary driver into block 0 of device *dd*. Finally, DUP writes the secondary bootstrap to blocks 2 through 5 of device *dd*.

Figure 1-6 illustrates the procedure.

Figure 1-6: COPY/BOOT *xxn:filnam ddm*: Procedure

1.11.3.3 BOOT *ddn*:

Use the *BOOT ddn:* command to perform a software bootstrap of a specific device that already has a specific monitor secondary bootstrap in blocks 2, 3, 4, and 5 (placed there by the COPY/BOOT command). In the command line, *dd* represents the 2-character name of the device to be booted; *n* is its unit number. Both the new monitor file and the new device handler must be present on device *dd*.

As soon as this command is issued, DUP first checks that device *dd* is a random-access device. Then it reads blocks 2, 3, 4, and 5 into a memory buffer. These blocks contain the secondary bootstrap for the monitor. The primary driver is already in locations 0 through 776.

DUP locates the appropriate handler file on device *dd*. This procedure is a check that the volume has a system device handler stored on it so that it can be validly bootstrapped.

DUP then extracts the file name of the monitor file from locations 724 and 726 of block 4 and locates the monitor file on the device to make sure that it really exists.

Next, DUP loads the information shown in Table 1–10 for the primary driver and the secondary bootstrap.

Table 1–10: DUP Information

Offset from Start of Memory Buffer	Contents
4722	Booted unit number
5000	Date booted
5002–5004	Time booted

DUP then copies the primary driver and secondary bootstrap from the device into memory locations 0 through 4777. Then it jumps to location 1000 to start the secondary bootstrap at its DUP entry point so that the secondary bootstrap can load the monitor and the system device handler into memory.

If the `/FOREIGN` option is used, DUP reads in block 0 and jumps to location 0.

Figure 1–7 illustrates the procedure.

1.12 Including Support for Multiterminal Handler Hooks

Including handler hooks support in a multiterminal monitor and in your handler lets the handler use any serial line on the system. The distributed LS and XL handler source files contain conditionalized support for multiterminal handler hooks. In this section, the XL handler is used to provide example code. A copy of the XL handler with extended comments is located in Appendix A.

This section provides information on including support for multiterminal handler hooks in your handler. Chapter four in the *RT-11 System Internals Manual* contains a section that describes how the monitor supports such handlers. You should read that section before you read this one, as that section describes the basic monitor/handler protocol. It also describes the monitor data structures that your handler writes and accesses and the interrupt service routines your handler uses to read and write data.

Figure 1–7: BOOT ddn: Procedure

Support for multiterminal handler hooks should be included in at least the following places. Each item is described in detail with example code.

- Installation code following `.DRINS`, Section 1.12.1
- Set code for the supported `SET` command conditions at `.DRSET`, Section 1.12.2
- Establish the monitor hooks at installation or `LOAD/FETCH` code, Section 1.12.3
- Handler hook interrupt processing during execution of interrupt service code, Section 1.12.4
- Remove handler hooks connection with the monitor at `UNLOAD/RELEASE` code Section 1.12.5

1.12.1 Installation Support

The handler does the following at installation:

- Determines if the handler should use the handler hooks monitor support.
If not required, the handler can install for nonmultiterminal support.
If required but not available, the handler refuses to install.
- Assuming the proper conditions are met, the handler accepts the installation.

The following code is from the installation section of the XL handler source. R0 contains the contents of RMON fixed offset 54, \$SYPTR, and \$THKPT has been defined as 472:

```

        TSTB    I$MTTY                ;Are handler hooks needed?
        BEQ     20$                    ;Nope...
        TST     $THKPT(R0)            ;Yes, is the support available?
        BEQ     40$                    ;Nope, reject the installation
        BR      30$                    ;Yes, nothing to do until fetch/load

20$:    .
        .
        .
        .
30$:    TST     (PC)+                  ;Accept the installation (carry=0)
40$:    SEC                           ;Reject the installation (carry=1)
        RETURN
        .
        .
        .
I$MTTY: .BYTE   -1                    ; : Install-time 'hooks required' flag
        .BYTE                       ;reserved
```

1.12.2 SET Command Support

Two SET command conditions should be supported by a handler that has been built with hooks support:

- SET dd LINE=n

Support for this condition is included so that the handler can change the serial line to which it will attach. The default line number can be established during system generation.

- SET dd [NO]MTTY

Support for this condition is included when a handler is built to support both a standard interface and the multiterminal monitor hooks. In such a case, by default the handler assumes connection to the standard interface until the command is issued.

When the MTTY condition is specified, the handler should clear the installation CSR (found in handler file image 176). The handler should also clear the vector information in the handler header (handler file image offset 1002). The original contents of these words can be built into words elsewhere in the handler from which they can be restored when the NOMTTY condition is specified.

The code to support those SET command conditions for the XL handler follows:

The following is in block 0, following the installation code:

```

I$MTTY: .BYTE    -1                ; : Install-time 'hooks required' flag
        .BYTE                ;reserved
VECSAV: .WORD    100000+<<XL$VTB-H1.VEC>/2-1> ; : Vector info for SET NOMTTY
CSRSV:  .WORD    XL$CSR            ; : CSR info for SET NOMTTY
        .
        .
.DRSET  LINE     16.                O.LINE  NUM      ;LINE=n
.DRSET  MTTY     -1                O.MTTY  NO       ;[NO]MTTY
        .
        .
; SET XL LINE=line_number

O.LINE: CMPB     R0,R3              ;Is line number valid?
        BHI     O.ERR              ;Nope...
        MOVB    R0,O$LINE          ;Yes, set line number to use
        BR      O.NOR

; SET XL [NO]MTTY

O.MTTY: BR       10$               ;Entry point for MTTY
        NOP                ;placekeeper
        CLR     R0              ;Entry point for NOMTTY
        MOV     CSRSV,INSCSR     ;Nope, restore install-time CSR
        MOV     VECSAV,H1.VEC    ; and vector information
        BR      20$

10$:   CLR     INSCSR              ;Reset install-time CSR and
        CLR     H1.VEC            ; vector so handler installs
20$:   MOVB    R0,O$MTTY          ;Set/Reset MTTY hooks use flag
        MOVB    R0,I$MTTY        ; and inform install code of setting
        BR      O.NOR
        .
        .
O.NOR: TST     (PC)+              ;Success (carry=0)
O.ERR: SEC                ;Failure (carry=1)
        RETURN
        .
        .

```

The following is in the executable portion of the handler (block 1 and beyond):

```

; *** SET ***
O$MTTY: .BYTE    -1                ;Default to hooks used
        .Assume <O$MTTY-XLSTRT> LE 1000 MESSAGE=<Code to set not in block 1>

; *** SET ***
O$LINE: .BYTE    XL$LUN            ;Default line to use
        .Assume <O$LINE-XLSTRT> LE 1000 MESSAGE=<Code to set not in block 1>

```

1.12.3 Establish Hooks Connection with Monitor

The handler establishes hooks connection with the monitor at the LOAD/FETCH code. The code should do the following:

- Determine if handler hooks are required and if not, proceed with nonmultiterminal hooks code (connect to standard interface) so long as the CSR and vector do not conflict with any TCB in the multiterminal configuration.

The handler installation code should determine if support exists in the monitor for handler hooks. Therefore, because the handler is installed, support for handler hooks can be assumed.

- From RMON fixed offset \$THKPT, the handler should access the monitor data structure THOOKS and store the addresses of the hooks support routines in the in-memory image of the handler.
- Conduct the following tests:
 1. Determine which serial line is to be used and verify its validity.
Compare the requested line number with the maximum supported in THOOKS.
 2. Determine if the line is available.
From THOOKS data, access the TCB for the line and determine if the T.CSR word exists (showing the interface is present on the system) and if the value is correct.
 3. Verify that the line is not the console line.
Check the CONSL\$ bit in the T.STAT word of the TCB.
 4. Verify that the line is not already owned.
Check that the T.OWNER word of the TCB is zero.
- If the tests above are passed, the handler should determine its physical name and place it in the handler at the word just before the handler hooks routine.
- If the tests above are not passed, the handler should report a LOAD/FETCH error by setting the PSW carry bit and return.
- The handler then performs the following operations in the indicated order (to avoid any race condition):
 1. Store the address of the TCB to which it is attached in memory.
 2. Place the address of its handler hooks entry point in the T.OWNER word of the TCB.
That address must reside in the low 28K-words of memory in Kernel mode and Instruction address space.
 3. Set the HANMT\$ bit in the T.STAT word of the TCB.
 4. If you handler needs to monitor modem control signals, set the HANMC\$ bit in the T.STAT word of the TCB. Otherwise, modem control is handled by the multiterminal monitor as described in the remote terminal section of the *RT-11 System Internals Manual*

The following code from the XL handler source illustrates connection between the handler and the monitor.

```

;+
;
; LOAD
;
; This routine is entered on FETCH or LOAD of the XL handler
; and is used 1) to verify use of the handler in the specific
; configuration and, if needed, 2) to establish the required
; connections between the handler and the interrupt service of
; a monitor with support for multiterminal handler hooks.
;-

        .ENABL  LSB

FETCH::
LOAD::
        MOV     R5,ENTRY$           ; Save entry point
        MOV     R2,SLOT$           ; and table size
        MOV     @R5,R5             ; R5 -> Base of handler (in memory)
        MOV     @##$SYPTR,R0       ; R0 -> Base of RMON
        TSTB   <O$MTTY-XLLQE>(R5) ; Terminal hooks to be used?
        BEQ    20$                 ; Then use normal DL
        MOV     $THKPT(R0),R1      ; R1 -> Multiterminal handler hooks
        ; data structure in RMON
        BEQ    60$                 ; Monitor doesn't have the support...
        TSTB   (R1)+               ; Bypass structure size byte
        MOVVB  (R1)+,R2            ; R2 = Number of LUNs on system
        MOV    (R1)+,R3            ; R3 -> TCB list
        MOV    (R1)+,<MTOENX-XLLQE>(R5) ; Set pointer to output enable routine
        MOV    (R1)+,<MTYBRX-XLLQE>(R5) ; Set pointer to Break control routine
        MOV    (R1)+,<MTYCTX-XLLQE>(R5) ; Set pointer to Control routine
        MOV    (R1)+,<MTYSTX-XLLQE>(R5) ; Set pointer to Status routine
        MOVVB  <O$LINE-XLLQE>(R5),R0 ; R0 = Line to attach to
        BMI    60$                 ; Must be a positive number
        CMPB   R0,R2               ; Is line in this configuration?
        BGE    60$                 ; Nope, invalid line number
        ASL    R0                  ; Shift for word offset into TCB list
        ADD    R0,R3               ; R3 -> TCB list entry
        MOV    @R3,R3              ; R3 -> TCB for LUN
        TST    T.CSR(R3)           ; Is the line present in hardware?
        BEQ    60$                 ; Nope...
        TST    T.STAT(R3)         ; Is the line a console?

        .Assume CONSL$ EQ 100000
        BMI    60$                 ; Yes...
        MOV    R5,R0               ; R0 -> Handler hook routine
        ADD    #<XLHOOK-XLLQE>,R0  ; ...
        TST    T.OWNR(R3)          ; Is the line already attached?
        BEQ    10$                 ; Nope...
        CMP    R0,T.OWNR(R3)       ; Yes, to this handler?
        BNE    60$                 ; Nope...
10$:    MOV    ENTRY$,R1            ; R1 -> $ENTRY entry
        SUB    SLOT$,R1            ; R1 -> $PNAME ENTRY
        MOV    @R1,-2(R0)          ; Inform handler of its physical name,
        MOV    R3,<TCBADX-XLLQE>(R5) ; link the handler to the TCB
        BIS    #<HANMT$!HANMC$>,T.STAT(R3) ; declare line owned by handler
        ; and that handler will process modem,
        MOV    R0,T.OWNR(R3)       ; finally link the TCB to the handler
        BR     50$
        .ENDC ;NE XL$MTY

20$:    BIT    #MTTY$,$SYSGE(R0)   ; Is this a multiterminal monitor?
        BEQ    50$                 ; Nope, then there can't be a conflict
        .ADDR  #MTAREA,R0          ; R0 -> .MTSTAT EMT area
        .ADDR  #MTSTAT,R1         ; R1 -> Status block
        .MTSTA R0,R1               ; Get info about multiterminal system
        BCS    60$                 ; Errors?
        MOV    @##$SYPTR,R0       ; R0 -> $RMON
        MOV    MTSTAT,R1           ; R1 -> First TCB in system
        ADD    R0,R1               ; ...
        MOV    MTSTAT+MST.LU,R2    ; R2 = Highest LUN on the system
        ; (Number_of_LUNs - 1)

30$:    TST    T.CSR(R1)           ; Is this a configured line?
        BEQ    40$                 ; Nope...
        CMP    <XIS-XLLQE>(R5),T.CSR(R1) ; Will use of the CSR conflict?
        BEQ    60$                 ; Yes, reject the load
        CMP    <XL$VTB-XLLQE>(R5),T.VEC(R1) ; Will use of the VECTOR conflict?
        BEQ    60$                 ; Yes, reject the load

40$:    ADD    MTSTAT+MST.ST,R1    ; On to next TCB
        DEC    R2                  ; More TCB's to check?

```

```

        BGE      30$                ; Yep...
        .BR      50$                ; Nope, use of interface won't conflict

50$:    TST      (PC)+              ; Success return
60$:    SEC                      ; Error return
        RETURN

ENTRY$: .BLKW                      ; : -> $ENTRY table entry
SLOT$:  .BLKW                      ; : Size of a monitor handler table

MTAREA: .BLKW   3                   ; : EMT area for .MTSTAT
MTSTAT: .BLKW   8                   ; : Status block from .MTSTAT

```

1.12.4 Handler Hook Interrupt Processing

The handler hook interrupt entry point is called by the monitor whenever an interrupt occurs on the line to which the handler is attached.

When an input interrupt occurs, the monitor calls the handler hook entry point with the character in R5 and the TH.PIC function code in R0. The handler processes the character, preserving the registers, and returns.

When an output interrupt occurs, the monitor calls the handler hook entry point with the TH.GOC function code in R0. The handler returns the next output character in R5 and the PS carry bit is clear. If the handler has no character for output, it returns the PS carry bit set. All registers are preserved.

The multiterminal interrupt service controls character output. A handler cannot send output directly to an interface, but must instead indicate it has output by calling the MTOENB routine.

The following code from the XL handler source file illustrates the process.

```

; The following byte indicates whether the handler should make use
; of the multiterminal hooks during FETCH/LOAD and during operation.

; *** SET ***
O$MTTY: .BYTE   -1                  ;Default to hooks used
.Assume <O$MTTY-XLSTRT> LE 1000 MESSAGE=<Code to set not in block 1>

; *** SET ***
O$LINE: .BYTE   XL$LUN              ;Default line to use
.Assume <O$LINE-XLSTRT> LE 1000 MESSAGE=<Code to set not in block 1>

ISPND:  .BYTE   -1                  ; : Input suspend flag
OSPND:  .BYTE   -1                  ; : Output suspend flag
        .EVEN
        .
        .
;+
;
; XLHOOK
;   Entered from multiterminal input or output interrupt service.
;
; Call (TH.GOC):
;   R0 = Function code
;
; Return (TH.GOC):
;   PSW<C> = 0, R5 = character
;   PSW<C> = 1, no character available
;
; Call (TH.PIC):
;   R0 = Function code
;   R5 = character
;-

        .ENABL  LSB

```

```

XLPNAM: .Rad50 /XL/ ; : Rad50 physical name
; loaded by FETCH/LOAD code

XLHOOK:
.Assume <XLHOOK-XLPNAM> EQ 2 MESSAGE=<XLPNAM must precede XLHOOK>

MOV R4,-(SP) ;Save register for awhile

; Function code = 1 = TH.GOC
; (Get Output Character)

CMP R0,#TH.GOC ;'Get Output Character' request?
BNE 10$ ;Nope...
TSTB OSPND ;Is output suspended?
BNE 20$ ;Yep...
CALL HOINT ;Yes, hook handler output service
BR 30$

; Function code = 2 = TH.PIC
; (Put Input Character)

10$: CMP R0,#TH.PIC ;'Put Input Character' request?
BNE 20$ ;Nope...
TSTB ISPND ;Is input suspended?
BNE 20$ ;Yep...
CALL HIINT ;Yes, hook handler input service
BR 30$

20$: SEC ;Return failure
30$: MOV (SP)+,R4 ;*C* Restore previously saved register
RETURN

```

1.12.5 Remove Handler Hooks Connection to Monitor at UNLOAD/RELEASE

Upon UNLOAD or RELEASE, the handler should perform the following operations in the indicated order (to avoid any race conditions):

1. Clear the HANMT\$ and HANMC\$ bits in T.STAT in the TCB.
2. Clear the T.OWNER word of the TCB.

The following code from the XL source file illustrates the procedure:

```

UNLOAD::
MOV @R5,R5 ; R5 -> Handler entry point (XLLQE)
TST <STATFG-XLLQE>(R5) ; Is handler in use?
BNE 10$ ; Nope, it can be unloaded...
MOV <XICQE-XLLQE>(R5),-(SP) ; Check internal queues
BIS <XOCQE-XLLQE>(R5),(SP)+ ; ...
BEQ RELEAS ; They're empty...
.ADDR #NOUNLO,R0 ; R0 -> Error message string
; (KMON reports error)
SEC ; Indicate error
RETURN ; and return to KMON

RELEAS::
MOV @R5,R5 ; R5 -> Handler entry point (XLLQE)
10$: TSTB <O$MTTY-XLLQE>(R5) ; Terminal hooks in use?
BEQ 20$ ; Nope...
MOV <TCBADX-XLLQE>(R5),R1 ; R1 -> TCB we're hooked to
BEQ 30$ ; We're not...
CALL <DISINI-XLLQE>(R5) ; Disable input
CALL <DISOUI-XLLQE>(R5) ; and output interrupts
CLR R0 ; Deassert all modem control bits
CALL <SETSTT-XLLQE>(R5) ; ...
CLR T.OWNER(R1) ; Disconnect TCB from handler
BIC #<HANMT$!HANMC$>,T.STAT(R1) ; ...
BR 30$

20$: ; Perform UNLOAD/RELEASE operations
; for a nonhooked handler
.
.
.
30$: CLC
RETURN

```

```
NOUNLO: .NLCSI TYPE=I,PART=PREFIX
        .ASCIZ "F-Handler may not be unloaded while in use"
```

1.13 Including Extended Device-Unit Support

When modifying a user-written handler to enable extended device-unit support, you should be aware of how an extended-unit handler interacts with two RMON tables (\$OWNER and \$PNAM2) and the functions specific to an extended-unit handler in the following:

- .DRDEF and .DRBEG macros
- LOAD/FETCH routine
- UNLOAD/RELEASE routine
- Q.FUNC byte of an I/O queue element

See the *RT-11 System Internals Manual* for a description of the \$PNAM2 table. See the following sections for a description of the changes to the \$OWNER table, the macros, routines, and byte.

1.13.1 .DRDEF and .DRBEG Macros

The .DRDEF and .DRBEG macros generate required code for the preamble of a device handler. Specify the UNIT64=YES parameter to the .DRDEF macro to place the 1-letter extended-unit handler name and ownership table in blocks 0 and 1 of the handler.

The format for calling .DRDEF is:

```
.DRDEF name,code,stat,size,csr,vec,[UNIT64=YES]
```

Note the *name* parameter to the .DRDEF call. The macro .DRDEF uses the first letter of the name parameter as the 1-letter physical device name of an extended-unit device.

The name parameter defines both the dd\$NAM constant (the traditional 2-letter physical device name) and the dd\$PN2 constant (the 1-letter device name).

The macro .DRDEF places the RAD50 representation of the 1-letter device name followed by two blanks in offset H.64UM (100₈) of block 0 of an extended-unit handler. It also places the location of the extended-ownership table in offset H.UNIT (76₈) of block 0 of the handler and indicates generation of the table in the last 32-bytes of memory-resident code. (If the monitor under which the handler is running does not support extended device units, those last 32 bytes are not loaded into memory.)

.DREND creates the extended-ownership table in the memory resident portion of the handler because UNIT64=YES was specified in the previous call to the .DRDEF macro. The extended-ownership table (dd\$U64) is always 16₁₀ words (64₁₀ nibbles) long.

1.13.2 LOAD/FETCH and UNLOAD/RELEASE Routines

You place the new LOAD/FETCH and UNLOAD/RELEASE routines in the extended-unit handler. You place those routines in the handler SETOVR PSECT and order the PSECTs in the handler source code such that SETOVR is the last. You then link the handler as illustrated in the following example command:

```
.LINK/NOBITMAP/EXE:BIN:SPX.SYG/BOUNDARY:512. OBJ:SPX  
Boundary? SETOVR  
.
```

1.13.2.1 LOAD/FETCH Routine

If the running RT-11 monitor has extended device-unit and device ownership support, then the LOAD/FETCH routine:

1. Places a pointer to the handler's extended-ownership table in the second word of the handler's entry in the monitor's \$OWNER table.
2. Sets the first word of the handler's entry in the \$OWNER table to a value of 2. This value, and any nonzero even value in the \$OWNER unit 0 nibble, is a flag (the \$XUNIT flag) indicating both that the handler supports up to 64 units and that the second word of the handler's entry in the \$OWNER table points to a separate list holding the \$OWNER nibbles. The \$XUNIT flag is filled in at bootstrap/install time.

The definition of a nibble (4 contiguous bits) in the \$OWNER table for a nonextended-unit device handler is that its value is either 0 or the job number + 1. Therefore, an \$OWNER nibble of a nonextended-unit device handler is always 0 or odd, since job numbers are always even.

Figure 1-8 shows the handler entry in the \$OWNER table pointing to the extended-ownership table in the handler.

1.13.2.2 UNLOAD/RELEASE Routine

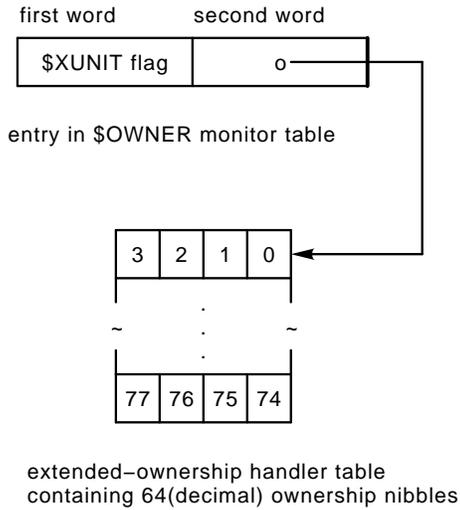
If extended-unit support is enabled in the running RT-11 monitor, the UNLOAD/RELEASE code of an extended-unit handler clears the second word of the handler's entry in the monitor's \$OWNER table, since the extended-ownership table (along with the handler itself) is being removed from memory.

1.13.2.3 Example LOAD/FETCH and UNLOAD/RELEASE Routines

The following example LOAD/FETCH and UNLOAD/RELEASE routines would be appropriate for extended device-unit handlers:

```
.IF NE,dd$N64  
.PSECT SETOVR  
.SBTTL LOAD - Load/Fetch code for extended device-unit handler
```

Figure 1-8: Relationship of \$OWNER Table to Extended-Ownership Table



```

;+
; Example LOAD/FETCH routine for a extended device-unit Handler.
;
; INPUT
;
;      R0          -> handler routine being called
;      R1          -> GETVEC routine
;      R2          $SLOT*2
;      R3          type code
;      0 -- .FETCH
;      2 -- .RELEASE
;      4 -- $LOAD
;      6 -- $UNLOAD
;      10-- Job Abort
;      12-- BSTRAP
;      R4          -> read routine
;      R5          -> $ENTRY word as above
;
;
; BSTRAP or KMON INSTALL modifies $PNAME, $PNAM2, and $OWNER+0
; for an extended device-unit handler. You need to insert only the
; address of the extended-ownership table into $OWNER+2 here.
;
;
; OUTPUT
;      Registers need not be saved by the handler code
;      Carry clear, unless an error was detected by
;      $SYS or the handler code.
;      If an I/O error occurred, R0 is cleared and Carry set.
;      If the handler returns with Carry set, R0 is passed,
;      as it was returned by the handler.
;-

.LIBRARY "SRC:SYSTEM" ;Indicates SYSTEM.MLB

.MCALL .CF3DF        ; CF3.64 definition
.MCALL .FIXDF        ; $CNFG3 and $PNPTR definitions
.MCALL .SYCDF        ; $SYPTR definition

.CF3DF
.FIXDF
.SYCDF

```

```

LOAD:
FETCH:
      MOV     @#$SYPTR,R0           ; R0 -> Base of RMON
      BIT     #CF3.64,$CNFG3(R0)   ; Extended unit support in monitor?
      BEQ     10$                   ; Branch if not, done.
      BIT     #CF3.OW,$CNFG3(R0)   ; Owner table support in monitor?
      BEQ     10$                   ; Branch if not, done.
20$:   MOV     R2,R3                 ; R3 = $SLOT*2
      ASL     R3                     ; *4
      ASL     R3                     ; *8
      ADD     R2,R3                 ; R3 = $SLOT*10
      CALL    FIXOWN                ; Insert extended ownership table addr
                                      ; into $OWNER word #2.
                                      ; R1 -> $OWNER+2
10$:   CLC
      RETURN                        ; Done.

.SbTtl UNLOAD - Unload/release code for a extended device-unit handler

;+
; Example UNLOAD/RELEASE routine for a extended device-unit Handler.
;
; INPUT
;
; R0          -> handler routine being called
; R1          -> GETVEC routine
; R2          $SLOT*2
; R3          type code
;             0 -- .FETCH
;             2 -- .RELEASE
;             4 -- $LOAD
;             6 -- $UNLOAD
;             10-- Job Abort
;             12-- BSTRAP
; R4          -> read routine
; R5          -> $ENTRY word as above
;
;
; This routine should zero the $OWNER+2 pointer to the extended ownership
; table of an extended device-unit handler.
;
; OUTPUT
;
; Registers need not be saved by the handler code
; Carry clear, unless an error was detected by $SYS or the handler code
; If an I/O error occurred, R0 will be cleared and Carry set.
; If the handler returns with carry set, R0 will be passed
; as it was returned by the handler.
;-

RELEASE:
UNLOAD:
      MOV     @#$SYPTR,R0           ; R0 -> base of RMON
      BIT     #CF3.64,$CNFG3(R0)   ; extended device-unit support
                                      ; in monitor?
      BEQ     10$                   ; Branch if not
      BIT     #CF3.OW,$CNFG3(R0)   ; Owner table support in monitor?
      BEQ     10$
      MOV     R2,R3                 ; R3 = $SLOT*2
      ASL     R3                     ; *4
      ASL     R3                     ; *8
      ADD     R2,R3                 ; R3 = $SLOT*10.
      CALL    FIXOWN                ; R1 -> $OWNER+2
      CLR     @R1
10$:   CLC
      RETURN

.SBTTL FIXOWN - insert pointer to extended ownership table into $OWNER

```

```

;+
; FIXOWN - insert pointer to extended ownership table into second word
; of $OWNER table (64 UNITS ONLY!!!)
;
; INPUT
;   R2 = $SLOT*2
;   R3 = $SLOT*10.
;   R5 -> $ENTRY entry for this handler
;   dd$X64: extended ownership table
;
; OUTPUT
;   $OWNER+2 points to extended ownership table
;   R1 points to $OWNER+2
;-

FIXOWN: MOV    @#SYPTR,R1          ; R1 -> $RMON
        MOV    $PNPTR(R1),-(SP)
        ADD    R1,@SP            ; @SP -> beginning of $PNAME
        ADD    R2,@SP            ; @SP -> beginning of $ENTRY
        MOV    R5,R1             ; R1 -> $ENTRY entry for this handler
        SUB    (SP)+,R1          ; R1 = byte offset into $ENTRY
        ADD    R5,R1             ; R1 = $ENTRY + double-word index
        SUB    R3,R1             ; R1 -> $OWNER of this handler + 8.
        CMP    -(R1),-(R1)       ; R1 -> $OWNER of this handler + 4
        MOV    @R5,-(R1)         ; move addr of ddLQE into $OWNER+2
        ADD    #dd$X64-ddLQE,@R1 ; make $OWNER (pic) to point to
                                   ; extended ownership table

        RETURN

.ENDC ;NE dd$N64

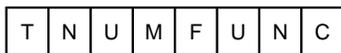
```

1.13.3 Q.FUNC Definition

The Q.FUNC byte of an I/O queue element passed to an extended-unit handler is different from the Q.FUNC byte of an 8-unit handler. However, the Q.FUNC byte passed to an 8-unit handler is unchanged to allow upward compatibility and to allow extended-unit handlers to function properly for units 0-7, when extended-unit support is not included in the running RT-11 monitor.

Q.FUNC is the low byte of the fourth word of the I/O queue element passed to a handler in an I/O request. Q.FUNC contains the special function code and the high 3 bits of the handler unit number.

The following diagram shows the bit layout of the Q.FUNC byte for an extended-unit handler:



T means the TYPE of I/O request.

NUM means the UNIT NUMBER.

FUNC means the FUNCTION.

The I/O request can be one of two types:

- On standard I/O requests or requests for special directory operations, the T bit is 0. In this case:
 - NUM is the high 3 bits of the handler unit number.
 - FUNC is a value 0000 through 1111. (The value 0000 specifies a read, write, or seek operation; 0001 through 1111 specifies a special directory operation.)

- On special function (SPFUN) I/O requests, the T bit is 1. In this case:
 - NUM is one's complement of the high 3 bits of the handler unit number.
 - FUNC is a value 0000 through 1111 that specifies an SPFUN operation from SPFUN 360 (0000) to SPFUN 377 (1111).

1.13.4 Programmed Requests of Extended-Unit Handlers

You must modify programs that assemble device specifications from physical device names and unit numbers for those programs to support extended-unit handlers. You can do this in conjunction with use of the .CSTAT programmed request, which reports the device on which a file is located.

For an extended-unit handler, .CSTAT returns the 2-letter device name from the \$PNAME table if the device unit specified falls in the 0-7 range. If the device unit specified is greater than 7, .CSTAT returns the 1-letter device name found in the new \$PNAM2 table.

1.14 How to Assemble, Link, and Install a Device Handler

Assembling, linking, and installing a new device handler are simple procedures described in detail in the following sections.

1.14.1 Assembling a Device Handler

The command you use to assemble your handler can include the following elements:

- Your MACRO-11 source file should be named *dd*.MAC, where *dd* is the 2-character device name.
- You can use the /SHOW:MEB assembler option to print the expansions of macros such as .DRBEG and .DRAST in the assembly listing.
- Each monitor has a corresponding conditional source file, such as XM.MAC, which defines the basic features of that monitor.
- SYSGEN.CND is the default name of the SYSGEN conditional file and is a product of the system generation process. Omit this file if you are assembling a device handler that will run with a distributed RT-11 monitor.

If your handler is to be used with a monitor that was produced through the system generation process, you must use the SYSGEN conditional file with which you assembled that monitor so that the handler conditionals will match the monitor conditionals and the handler will operate in the correct environment. You can specify the name of the SYSGEN conditional file by requesting an answer file during the SYSGEN process as the .CND file takes the file name of the answer file.

- If you have used symbol names from the distributed system library SYSTEM.MLB, you should assemble your handler with that library. The default device for SYSTEM.MLB is SRC, so you should assign SRC to that device on which SYSTEM.MLB resides and include SRC in the full file specification (SRC:SYSTEM.MLB).

Include the line `.LIBRARY "SRC:SYSTEM.MLB"` early in your program to call that library.

To assemble a handler for an unmapped system, use the following command where *mon* is the distributed monitor conditional source file:

```
.MACRO/CROSSREFERENCE/SHOW:MEB/LIST mon+SYSGEN.CND+SYSTEM.MLB/LIBRARY+dd/OBJECT
```

To assemble a handler for a mapped system, use the following command, where *mon* is the distributed monitor conditional source file:

```
.MACRO/CROSSREFERENCE/SHOW:MEB/LIST mon+SYSGEN.CND+SYSTEM.MLB/LIBRARY+dd/OBJECT:ddX
```

1.14.2 Linking a Device Handler

Once your source file assembles without errors, you are ready to link it. To link a handler for an unmapped system, use the following command:

```
.LINK/NOBITMAP/EXECUTE:dd.SYS dd
```

To link a handler for a mapped system, use the following command:

```
.LINK/NOBITMAP/EXECUTE:ddX.SYS ddX
```

If the handler requires block alignment of some code, use the following command where *nnn* is the block alignment boundary for PSECT *psect*:

```
.LINK/NOBITMAP/EXECUTE:ddX.SYS/BOUNDARY:nnn ddX  
psect
```

1.14.3 Installing a Device Handler

Before you can use your new handler, you must inform the monitor that the handler is present and you want it installed. Add the monitor information about it to the monitor device tables described in Chapter 2 of the *RT-11 System Internals Manual*. The process of adding a new device is called installation. There are two separate routines in the RT-11 system that can install a device handler: the bootstrap and the monitor `INSTALL` command. Both routines require a device's hardware to be present on the system before they install the device handler. (Section 1.14.3.6 describes a way to circumvent this restriction if you need to install a handler for a nonexistent device.)

The following sections describe the various ways to install device handlers in an RT-11 system.

1.14.3.1 Using the Bootstrap to Install Handlers Automatically

The bootstrap routine first locates the system device handler on the device from which you booted the system and installs it. Then it scans the rest of the handler files on the system device and tries to install the corresponding handler for each hardware device it finds on the system. If the hardware is not present, the bootstrap does not install the device.

The only difficulty with this procedure occurs when there are more handler files than device slots. A distributed monitor reserves one device slot for each device RT-11 supports. A monitor you create through system generation reserves one slot for each device you request. In addition, it provides the number of empty slots you specify.

A slot is considered to be reserved for a particular device if the \$PNAME monitor table has an entry for that device. A slot is empty if \$PNAME has a zero word.

The automatic device installation routine in the bootstrap has a set of priorities to determine which handlers to install when there are more handlers than slots. If all slots are empty, the bootstrap installs the system device handler plus the first handlers it encounters on the system device whose device hardware is present. For example, if a system has eight slots, all empty, the bootstrap installs the system device handler and the first seven legitimate handlers it finds on the system device.

If one or more slots are reserved for specific devices (that is, the devices have entries in the \$PNAME table), the bootstrap reserves those slots for the corresponding handlers until it can verify the presence of the appropriate hardware. If the hardware exists, the bootstrap installs its device handler. If the hardware is not present, the bootstrap clears its \$PNAME entry, thus creating an empty slot.

Figure 1–9 summarizes the algorithm the bootstrap uses to install device handlers.

As you can see, handlers with entries in the \$PNAME table have higher priority at boot time. If the handler file is on the system device and the device hardware exists, the bootstrap always installs the handler.

When you write a device handler yourself, you should have no problem installing it in your RT–11 system because you can rely on the bootstrap to install the handler for you if the handler resides on the system device, if its hardware is present, and if there is an empty slot in the monitor tables. If your system has no free slot, you can create one or more by simply storing fewer device handler files on your system device and rebooting the system. You can also use the monitor INSTALL command (described in Section 1.14.3.2) to install a new handler without rebooting the system. (This new handler may be one that the bootstrap could not install due to lack of free slots, or it may be a new handler that you just created or just copied to the system device.) Or, if you created your system through system generation, you can use the DEV macro (described in Section 1.14.3.3) to reserve a slot for a new device handler and give it priority for installation at bootstrap time. Figure 1–10 summarizes the ways you can install a new device handler.

1.14.3.2 Using the INSTALL Command to Install Handlers Manually

Before using the INSTALL command to install a handler manually, use the SHOW command to see if there are any empty device slots on your system. If there are none, use the REMOVE command to remove a device you do not need and make room for your new device, which you then add by using the INSTALL command. The formats of these commands are documented in the *RT–11 Commands Manual*.

If a device slot was already available, your device will install automatically the next time you bootstrap the system. If you used REMOVE and INSTALL to add your new device to the system, you must reissue the commands after each bootstrap. To

Figure 1–9: Bootstrap Algorithm for Installing Device Handlers

Figure 1–10: Installing a New Device Handler

install the new device automatically at each bootstrap, put **REMOVE** and **INSTALL** commands in your system's startup indirect file. This saves you the trouble of typing the commands yourself. In addition, it gives the device the appearance of being permanently installed.

1.14.3.3 Using the DEV Macro to Aid Automatic Installation

If you created your system through a system generation, you can edit a system **MACRO-11** source file to add a new device to the **\$PNAME** table, thus giving it preference in the automatic handler installation procedure. The file you edit is **SYSGEN.TBL**, one of the files you assemble to create a monitor file.

Use the **DEV** macro in the file **SYSGEN.TBL** to add a new device to the system permanently. The format of the **DEV** macro is as follows:

DEV name,s

name is the 2-character device name.

s represents the device status word (leave this argument blank).

The following examples are taken from the SYSGEN.TBL file:

```
DEV      RK                ;INSTALLS THE RK DISK
DEV      LP                ;INSTALLS THE LINE PRINTER
DEV      MT                ;INSTALLS MAGTAPE
```

After you edit SYSGEN.TBL to add the DEV macro call for your device, you must reassemble it. Use the following command:

```
.MACRO/OBJECT:TBxx mon+SYSGEN.CND+SYSGEN.TBL
```

xx represents the monitor type, such as SB, FB, XM, or another of the mapped monitors.

mon represents the monitor conditional source file, such as XM.MAC, which defines the basic features of that monitor. Once the assembly is complete, relink the object files to create your new monitor. Follow the commands in the command file that resulted from your system generation procedure to build the modified system.

1.14.3.4 Installing Devices Whose Hardware Is Present

Both routines in RT-11 that can install a device handler—the bootstrap code and the monitor INSTALL command code—install handlers only for those devices whose hardware is present on the current system configuration. The routines look at location 176 in block 0 of the handler and test the address that 176 contains, which is normally the base CSR for the device. If the hardware for the device is not present on the system, a bus timeout occurs, causing a trap to 4, which the installation routines field. As a result, neither the bootstrap routine nor the INSTALL command will install the device handler. In addition, the INSTALL command prints the *?KMON-F-Invalid device installation* message.

The installation routines think the device's hardware is present if its CSR responds on the bus. However, this simple test is not sufficient to determine, in some cases, which hardware device is present. For example, some devices are assigned the same addresses in the I/O page for one or more of their status registers. If RT-11 just tested a "shared" I/O page address, it still does not know which of two devices is really present and therefore which handler to install. The RX01 and RX02 diskette devices, for example, have the same bus address and the same number of status registers in the I/O page. When RT-11 attempts to install the DX handler, it must be able to determine whether or not hardware is present, and whether or not it is the RX01 device. Clearly, it should not install the DX handler when the hardware is really the RX02 device.

There is almost always some difference between two or more devices that is discernible from their registers in the I/O page. Each handler for one of the hard-to-identify devices can test for this difference and inform the RT-11 installation routine whether or not it should install the device handler it is currently considering.

1.14.3.5 Writing an Installation Verification Routine

RT-11 handlers for devices with shared I/O page addresses all contain an installation verification routine to distinguish which hardware device is actually present and to permit or inhibit installation of the current handler. If you write a device handler yourself, you can include your own installation verification routine.

In general, the installation verification routines distinguish which hardware is present based on one of the three following conditions:

- Of the two devices that share some registers, one device has more registers than the other.
- If two devices share addresses for all their registers, and if they have the same number of registers, sometimes one device has a read/write bit where the other device has a read-only bit.
- Sometimes a device has a unique identification bit or byte.

The installation verification routines, then, determine which device is present based on the results of testing one of the distinguishing conditions. Once this determination has been made, the routine signals to the RT-11 installation routine whether or not to install the current handler and then returns to the monitor with the carry bit set to prevent installation and with the carry bit clear to permit installation.

Note that your installation verification routine can use all registers.

Entry Points of the Installation Verification Routine

An installation verification routine that you write in your own handler starts at location 200 in block 0 of the handler. It must not extend beyond location 360, unless you link your handler with the /NOBITMAP option – in which case location 376 is the limit. Location 200 is the entry point that the bootstrap code uses to install a data device. The INSTALL monitor code always enters here, as well.

Location 202 is the entry point that the bootstrap code uses to install the system device. The INSTALL monitor code never enters here.

If you do not care whether your handler is installed as the system device or as a data device, put a NOP instruction at location 200. If your handler must be installed as the system device handler, use the following instructions to prevent its installation under any other circumstances:

```
      . = 200                ;NON-SYSTEM ENTRY POINT
      BR      ERROR          ;BRANCH TO ERROR ROUTINE
      .
      .
      .
      ; Code to execute when installed as system device
      .
      .
      .
ERROR: SEC                ;SET CARRY TO PREVENT INSTALLATION
      RTS      PC          ;AND RETURN
```

The `.DRINS` macro sets up the installation code area in block 0 of a device handler. `.DRINS` defines symbols for the installation verification code entry points and for the installation CSR. After `.DRINS` is called, the location counter is set to 200, the address of the data device installation entry point.

.DRINS Macro—Use the `.DRINS` macro near the beginning of your device handler, before the header section. The `.DRINS` macro is described in Section 1.2.1.3.

If the Hardware for This Handler Has an Extra Register

If this handler is for a device that shares an I/O page address with another device, you can identify which device is present if the two devices have a different number of registers. When the device for the current handler has one more register than the other device, use the following instructions to test for the extra register:

```

MOV      176,R0          ;GET THE SHARED CSR
TST      n(R0)          ;TEST THE EXTRA REGISTER AT OFFSET n
                        ;THE SHARED CSR
RTS      PC             ;RETURN (WITH CARRY SET
                        ;IF WRONG DEVICE)

```

This routine tests the extra register. If there is no device configured there, the bus times out, causes a trap to 4, and sets the carry bit. The installation verification routine returns to the monitor with the carry bit set, indicating that the correct hardware for the current handler is not present, and that this handler should not be installed.

On the other hand, if the extra register responds to the test, the `TST` instruction returns with the carry bit clear, which means that the correct hardware for this device handler is present, and that RT-11 should install the handler.

If the Hardware for This Handler Has Fewer Registers

If the hardware for the other device that shares an I/O address with the device for this handler has more registers, this handler can test for the absence of the extra register. If the extra register is not found, RT-11 should install the current handler.

The following instructions take care of this situation:

```

MOV      176,R0          ;GET THE SHARED CSR
TST      n(R0)          ;TEST THE EXTRA REGISTER AT OFFSET n
                        ;FROM 176. IS A DEVICE HERE?
BCC      1$             ;YES, OTHER DEVICE IS HERE.
CLC                          ;NO, CLEAR CARRY
RTS      PC             ;INSTALL CURRENT HANDLER

1$:      SEC             ;SET CARRY
RTS      PC             ;DO NOT INSTALL CURRENT HANDLER

```

Essentially, this routine checks for the presence of the other device's extra register. If it is not present, the routine instructs RT-11 to install the current handler.

If an Identification Bit or Byte Exists

If the devices that share an I/O page address also share an identification bit or byte, an installation verification routine can check the bit or byte and determine which hardware is present. It can then permit or inhibit the installation of the current handler based on that information.

In RT-11, for example, the RX01 and RX02 devices share the CSR. Bit 11, called CSRX02, is clear if the device is an RX01, and set if the device is an RX02. The following example is from the DY device handler, which should only be installed if RX02 hardware is present.

```

.ASECT
. = 200
NOP
BIT #CSRX02,@176
BEQ 1$

TST (PC)+
1$: SEC
RTS PC
;VERIFICATION ROUTINE GOES HERE
;SAME CHECK FOR SYSTEM AND NON-SYSTEM
;IS RX02 BIT ON?
;NO, THIS IS AN RX01.
;DON'T INSTALL THIS
;DY HANDLER.
;CLEAR CARRY, SKIP SEC INSTRUCTION.
;WE HAVE AN RX02, INSTALL DY HANDLER
;SET CARRY, DON'T INSTALL DY HANDLER
;RETURN TO MONITOR

```

If One Device Has a Read/Write Bit

If one of the devices that share an I/O page address has a read/write bit in the CSR where the other device has a read-only bit, the verification routine can determine which hardware is present by following a general procedure to check the bit and permit or inhibit the installation of the current handler based on the results. The routine should read the bit, toggle it, and write it back to the CSR. Then the routine should read the bit again. If the value of the bit changed, the device with the read/write register is present. If the value remained constant, the device with the read-only register is present. The routine can set the carry bit appropriately and return to the monitor. If carry is set, RT-11 does not install this handler. If carry is clear, RT-11 does install this handler.

1.14.3.6 Overriding the Hardware Restriction

If for any reason you need to install a device handler whose hardware is not present in your current system configuration, you can circumvent the checks in the bootstrap and INSTALL routines by running SIPP and patching the handler. You clear location 176 in the handler file's block 0, then use the INSTALL command or reboot the system to install the device handler.

1.15 How to Test and Debug a Device Handler

Once your new handler is assembled, linked, and installed, you are ready to begin testing it. Remember during debugging that you must remove the old handler and install the new one each time you create a new version of *dd(X).SYS*.

Test the handler in three stages, according to these guidelines:

1. Use the hardware version (SDH.SYS or SDHX.SYS) of DBG-11 to observe the handler as it processes a data transfer.

If for some reason you would rather use ODT or VDT to observe the handler as it processes a data transfer, see Sections 1.15.2 and 1.15.3. However, debugging is significantly easier when using a symbolic debugger, so look closely at using DBG-11 before choosing ODT or VDT.

2. Test the handler with keyboard monitor commands, with system utility programs, and with FORTRAN, C, or another programming language. Try the COPY command, for example, to copy data to and from the device, or run PIP to do the same thing. Try using the handler with FORTRAN READ or WRITE statements, or with BASIC-PLUS INPUT or PRINT statements. If your handler sets the bit in the device status word that indicates that the handler is for an RT-11 directory-structured device, DUP will operate correctly on the device with no further modifications. That is, you should be able to use DUP to initialize the device (through the INITIALIZE command) and to consolidate free space (through the SQUEEZE command). The RESORC program needs no modification to recognize the new device and will include it in its SHOW DEVICES report.
3. Give the handler an extended workout with an application program that uses wait-mode I/O, asynchronous I/O, and completion routines.

When the handler passes all the tests successfully, you can begin using it as part of your regular RT-11 system.

1.15.1 Using DBG-11 to Test a Handler

Chapter 5 of the *DBG-11 Symbolic Debugger User's Guide* describes using DBG-11 to debug a device handler. If you have not used DBG-11 previously, you should work through the examples in the manual before testing and debugging your handler.

1.15.2 Using ODT to Test a Handler

The easiest way to use ODT to test a handler is to run ODT as the foreground job. If you normally use only a single-job monitor, it is worthwhile to switch to a multi-job monitor just for debugging.

Since you will be doing some careful debugging work, Digital also recommends that you be the sole user during this time. Bring up your system from a hardware bootstrap. Do not start any system jobs or load any handlers.

Link ODT for the foreground with the following command:

```
.LINK/MAP/BACKGROUND ODT
```

Next, load the device handler you need to debug:

```
.LOAD dd[X]
```

Now, issue a SHOW D command. Note the address given for the device handler that you are debugging. For this example, assume the value is 131634. Subtract 6 (in octal) from this address to get the base address of the handler. In this case,

```
131634
-     6
-----
131626
```

Start ODT as the foreground job:

```
.FRUN ODT
```

ODT V01.04

*

Set relocation register 0 to the value computed from the address given by the SHOW D command:

```
131626;0R
```

You can step through the handler in memory as you follow the instructions in your assembly listing. The first five words are the header; the first executable instruction is the sixth word. Set your first breakpoint at the sixth word:

```
0,12;0B
```

Set other breakpoints at various points in the handler that you want to examine during debugging. Another critical place is the interrupt entry point. You can find its location by checking the handler's MACRO-11 listing. Remember, the interrupt entry point is called *ddINT*; you should be able to find it easily and set a breakpoint there.

When you have finished setting breakpoints in the handler, exit from ODT:

```
0;G
```

Now try using the handler. You could try using DUP to initialize the device, or PIP to copy data to the device. Or, run a test program that you have designed especially for this purpose. When execution reaches the first breakpoint in the handler, ODT takes control. Use ODT as usual to examine locations and check their values, or to modify instructions. Note that the default priority of ODT is 7; this prevents other interrupts from disturbing your debugging session. Since you are the only user on the system, ODT's high priority should cause no problem. (Note, however, that the system clock will lose time, and that ODT usually cannot debug race conditions.)

When you are satisfied with the handler's performance, remove the breakpoints from it and proceed with the remainder of execution through the handler:

```
;B  
;P
```

Be careful not to unload the foreground job (ODT) while there are still breakpoints set in the handler.

1.15.3 Using ODT in a Mapped Environment

By following a few special guidelines, you can use ODT to debug a device handler in the mapped environment.

Carefully select a place for ODT in memory. You can link it with an application program, or link it so it resides somewhere in memory where it will not be destroyed. If a breakpoint is to be taken in kernel mode, ODT must not reside in the PAR1 area (locations 20000 through 37776). The safest place to put ODT is in the foreground partition, as described in Section 1.15.2.

When you are debugging with ODT, the I/O page must always be mapped.

Setting breakpoints also requires care. As soon as you enter ODT, look at the breakpoint trap vector (BPT) at locations 14 and 16 in low memory. When you set a breakpoint, you must manually set the current mode bits, bits 14 and 15, of the PS at location 16. Set them to the mode you expect at the time the breakpoint occurs. The values are 11 for User Mode, 01 for Supervisor, and 00 for Kernel. (RT-11 utility programs such as PIP and DUP run in User Mode and expect the mode bits to be set to 11.)

After setting breakpoints, type 0;G to exit from ODT. This causes an .EXIT request to be performed, which destroys the BPT vector. So, after you exit from ODT, you must manually reconstruct the contents of the vector by using the Deposit command, as follows:

```
D 14=(correct contents of 14),(correct contents of 16)
```

Make sure no other jobs are running when you do this, since context switching causes this technique to fail.

1.16 Contents of .SYS Image of a Device Handler

Table 1-11 shows the layout of the .SYS image of a handler after assembly and linking. Tables 1-3 and 1-4 contain more information about blocks 0 and 1 of the device handler file image. Locations not otherwise identified are reserved for future use by Digital.

Table 1-11: Device Handler .SYS Image

Location	Contents
000000	Handler identifier in RAD50
000002	Pointer to a FETCH service routine (file address)
000004	Pointer to a RELEASE service routine (file address)
000006	Pointer to a LOAD service routine (file address)
000010	Pointer to an UNLOAD service routine (file address)
000012-000016	Reserved
000020	Device classification
000021	Device classification modifier
000022	First special function (index method) list
000024	Second special function (index method) list
000026	Third special function (index method) list
000030	Pointer to further special functions (extension table method)
000032	Pointer to bad-block replacement table
000034	Reserved

Table 1–11 (Cont.): Device Handler .SYS Image

Location	Contents
000036	Second status word
000040-000050	Reserved
000052	Handler size (<i>ddEND–ddSTRT</i>)
000054	Device size (<i>ddDSIZ</i>)
000056	Device status word (<i>ddSTS</i>)
000060	Result of FORCE= parameter; byte contains device SYSGEN options for SET <i>dd SYSGEN</i>
000061	Reserved
000062	Pointer to the primary bootstrap (file address)
000064	Bootstrap size in bytes
000066	Pointer to the bootstrap read routine (file address)
000070	Varies with the handler; see Table 1–3
000072	Varies with the handler; see Table 1–3
000074	Size in bytes of total list of handler data table descriptors
000076	Pointer to extended device-unit ownership table (file address)
000100	Letter name of extended device-unit handler and device characteristics for UMR support
000102	Reserved
000104	Pointer to further block 0 type information not included in block 0 (file address)
000106	Pointer to the handler data descriptor table (file address)
000110-000173	Information written by .MODULE and .AUDIT. (The CSR table begins at 176 and expands downward to a zero word.)
000174-000xxx	'Display' CSRs (DISCSR) read by RESORC. If more than one, each written into previous location
000176	'Installation' CSR (INSCSR); beginning of CSR table
000200	Data device installation entry point (INSDAT)
000202	System device installation entry point (INSSYS)
000204-000377	Installation code; must link with /NOBITMAP option or else range is 204-357
000400-000777	SET code/tables
001000	Either the device vector or an offset to the table of vectors (<i>ddSTRT</i>)

Table 1–11 (Cont.): Device Handler .SYS Image

Location	Contents
001002	Offset to the interrupt service entry point
001004	Priority (340)
001006	Pointer to the last queue element (<i>ddLQE</i>)
001010	Pointer to the current queue element (<i>ddCQE</i>) in the handler memory image
001010	Flag word (in handler file image)
001012	NOP instruction OR'd with flags
001014	Branch instruction (optional)
001016	Second flag word (optional)
001020	Pointer to SPFUN address check routine (optional)
001022	pointer to DMA SPFUN table (optional)
001024	Pointer to LD translation table (optional)
1012	Handler entry point
n	Abort entry point (from .DRAST; may be above 1777)
n+2	Interrupt entry point (from .DRAST; may be above 1777)
1776	High limit of area modifiable by SET code
<i>dd\$END</i>	\$RLPTR: (from .DREND) \$MPPTR: (from .DREND) \$GTBYT: (from .DREND) \$PTBYT: (from .DREND) \$PTWRD: (from .DREND) \$ELPTR: (from .DREND) \$TIMIT: (from .DREND) \$INPTR: (from .DREND) \$FKPTR: (from .DREND) <i>ddEND=.</i> (from .DREND)
<i>ddEND:</i>	
<i>ddBOOT:</i>	NOP; Start of primary bootstrap (from .DRBOT)

Table 1-11 (Cont.): Device Handler .SYS Image

Location	Contents
	BR <i>entry</i> ; Label <i>entry</i> from .DRBOT
entry-14	020; (from .DRBOT) This byte identifies the type of CPU. A value of 20 indicates a PDP-11.
entry-12	Controller types; (from .DRBOT) This byte indicates the type of controllers that the operating system supports for this device. Its value in RT-11 V5 can be the OR'd result of the following codes: 101 Non-MSCP UNIBUS controller 102 Non-MSCP LSI-11 buscontroller 110 MSCP UNIBUS controller 120 MSCP LSI-11 bus controller
entry-10	020; (from .DRBOT) This byte identifies the type of file structure on the disk. A value of 20 indicates RT-11 file structure.
entry-6	checksum; (from .DRBOT) The checksum byte is a checksum of the previous three bytes. It is computed as the complement of the sum of the bytes.
entry-4	0; (from .DRBOT)
entry-2	diskette type; (from .DRBOT) This byte contains a bootstrap identification number in bits 0-6 and a flag to indicate single- or double-sided diskettes in bit 7. The values can be: Bit 7 = 0, Single-sided diskette Bit 7 = 1, Double-sided diskette
entry:	BR .+2 or BMI .+2 (from .DRBOT) Digital suggests that <i>entry</i> be located above location 120 in the bootstrap block. This will avoid conflict with vectors and the monitor SYSCOM area as the monitor is bootstrapped.
	Start of primary bootstrap read routine
662	High limit of primary bootstrap
664	Start of bootstrap error code
776	End of bootstrap error code

Programming for Specific Devices

This chapter provides information on device handlers that have special device-dependent characteristics. Read this chapter if you need to program specifically for one of the following devices:

- DL (RL01/RL02 disk handler)
- DM (RK06/RK07 disk handler)
- DU (MSCP disk handler)
- DW (CTI Bus-based disk handler)
- DX and DY (RX01/RX02 diskette handlers)
- DZ (Diskette handler)
- LD (Logical disk handler)
- MM, MS, and MT (Magtape handlers)
- MU (TMSCP magtape handler)
- NL (Null handler)
- NC, NQ, and NU (Ethernet handlers)
- UB (UNIBUS mapping register handler)
- VM (Virtual Memory handler)
- XC and XL (Communications Port handlers)

Much of the information in this chapter is based on other information in the RT-11 documentation set. You should be familiar with pertinent information found elsewhere rather than relying only on the information in this chapter. For example, much of the description of special functions as they apply to particular device handlers in this chapter assumes you know and understand the description of special functions (the .SPFUN request) in the *RT-11 System Macro Library Manual*.

You should look at the on-line index utility, INDEX, or in the printed *RT-11 Master Index* for other information in the RT-11 documentation set that pertains to a particular device handler or to various RT-11 features as they apply to that device handler.

Device handler operations are often controlled by various special functions. In this manual, you will be presented with both a code number and name for a special function. You can use the code in the particular special function call (.SPFUN) as

documented. You can use the name (rather than the code) if you include in your program a macro call for the appropriate macro in the file, SYSTEM.MLB.

The following macros in SYSTEM.MLB define the names for the indicated type of special functions:

Name	Device type
.SFDDF	Disk device handlers
.SFMDF	Magtape device handlers
.SFNDF	Ethernet device handlers
.SFXDF	VTCOM device handlers
.SFODF	'Other' device handlers, such as PI

You could, for example, include the following code in your program or handler to define the names of all the disk type special functions in this manual. Then you could use the special function name, rather than the more cryptic function code. (Be sure that the volume that contains SYSTEM.MLB is also assigned the logical name SRC.)

```
.LIBRARY "SRC:SYSTEM.MLB"  
.MCALL .SFDDF  
.SFDDF
```

2.1 DL (RL01/RL02 Disk Handler)

This section provides specific programming information for RL01 and RL02 disks.

2.1.1 Support for Special Functions

The RL01/RL02 disk handler supports the following special functions. The device-specific parameter arguments are the same as for DM; see Section 2.2 for information.

Code	Name	Action
377	SF.ARD	Read operation without doing bad-block replacement; returns definitive error data.
376	SF.AWR	Write operation without doing bad-block replacement; returns definitive error data.
374	SF.BBR	Re-read the bad-block replacement table in the handler (the program changed it).
373	SF.SIZ	Determine the size, in 256_{10} -word blocks, of a particular volume.

2.1.2 Support for Bad-Block Replacement

Bad-block replacement for the RL01 and RL02 is similar to the bad-block support for the RK06/RK07 (DM). However, the RL01 and RL02 generate neither the bad sector error (BSE) nor the header validity error (HVRC). Therefore, the handler must check the bad-block replacement table for each I/O transfer. Since the table is always in memory as part of the DL handler, the I/O delay is not significant.

The last track of the RL01 and RL02 disks contains a table of the bad sectors that were discovered during manufacture of the disk. The 10_{10} blocks preceding this table (the last 10_{10} blocks in the second-to-last track) are set aside for bad-block replacements. The maximum number of bad blocks (10_{10}) is defined in the handler.

As with the RK06 and RK07, you determine at initialization time whether to cover bad blocks with .BAD files or create a replacement table for them and substitute good blocks during I/O transfers. The advantage of using bad-block replacement is that it makes a disk with some bad blocks appear to have none. On the other hand, covering bad blocks with .BAD files fragments the disk. Because RT-11 files must be stored in contiguous blocks, this fragmentation limits the size of the largest file that can be stored.

The monitor file cannot reside on a block that contains a replaced block if you are using bad-block replacement. If this condition occurs, a boot error results when you bootstrap the system. In this case, move the monitor so that it does not reside on a block with an error.

If you specify the /REPLACE option during initialization of an RL01 or RL02 disk, DUP scans the disk for bad blocks. It merges the scan information with the

manufacturing bad sector table, allocates a replacement for each bad block, and writes a table of the bad blocks and their replacements in the first 20 words of block 1 of the disk. Block 1 is a table of two-word entries. The first word is the block number of a bad block; the second word is its allocated replacement. The last entry in the table is 0. The entries in the table are in order by ascending bad block number. A sample table is shown in Figure 2–1.

Figure 2–1: Bad-Block Replacement Table

Bad block	12	Entry 1
Its replacement	10210	
Bad block	37	Entry 2
Its replacement	10211	
Bad block	553	Entry 3
Its replacement	10212	
End of list	0	Entry 4

The handler contains space to hold a resident copy of the bad block table for each unit. The amount of space allocated is defined by the SYSGEN conditional DL\$UN, which represents the number of RL01/RL02 units to be supported. The value defaults to 2 if it is not defined. The handler reads the disk copy of the table into its resident area under the following three conditions:

- If a request is passed to the handler and the table for that unit has not been read since the handler was loaded into memory.
- If a request is passed to the handler and the handler detects Volume Check drive status. This status indicates that the drive spun down and spun up again, which means that the disk was probably changed.
- If an SF.BBR request is passed to the handler. This special function is used by DUP when it initializes the disk table to ensure that the handler has a valid resident copy.

2.2 DM (RK06/RK07 Disk Handler)

This section provides specific programming information for RK06 and RK07 disks.

2.2.1 Support for Special Functions

The RK06/RK07 disk handler supports the following special functions:

Code	Name	Action
377	SF.ARD	Read operation without doing bad-block replacement; returns definitive error data.
376	SF.AWR	Write operation without doing bad-block replacement; returns definitive error data.
374	SF.BBR	Reread the bad-block replacement table in the handler (the program changed it). SF.BBR uses no parameters.
373	SF.SIZ	Determine the size, in 256_{10} -word blocks, of a particular volume.

The special function (.SPFUN) request has the following general form, with the *area* and *chan* parameters and the optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*, and the other parameter arguments as described below:

Macro Call: .SPFUN *area,chan,func,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]*

func is the code for the function to be performed or the name of the function if the program has been assembled with the distributed module SYSTEM.MLB.

buf For SF.ARD and SF.AWR, the buffer size must be one word larger than required for the data. The first word of the buffer contains any returned error information. The remaining words in the buffer contain the data transferred. The error codes and information are as follows:

Code	Name	Meaning
100000	ES.SUC	The I/O operation is successful.
100001	ES.ECC	An ECC error is corrected.
100002	ES.RTY	An error was recovered on a retry.
100004	ES.UFF	An error was recovered through an offset retry.
100010	ES.RCL	An error was recovered after recalibration.
100200	ES.BBR	A bad block is detected (BSE error).
1774xx	ES.ERR	An error was not recovered.

For SF.BBR, *buf* should be 0.

For SF.SIZ, *buf* is a 1-word buffer where the .SPFUN request returns the size of the volume in 256_{10} -word blocks.

wcnt For SF.BBR, *wcnt* should be 0.

For SF.SIZ, *wcnt* should be 1.

blk For SF.BBR, *blk* should be 0.

For SF.SIZ, *blk* should be 0.

2.2.2 Support for Bad-Block Replacement

The last cylinder of the RK06 and RK07 disks is used for bad-block replacement and error information. RT-11 supports a maximum of 32_{10} replaceable bad blocks on these disks. The bad-block information is stored in block 1 on track 0, cylinder 0, of the disk. The replacement blocks are stored on tracks 0 and 1 of the last cylinder. A bad-block replacement table is created in block 1 of the disk by the DUP utility program when the disk is initialized. When a bad block is encountered and the table is not present in the handler from the same volume, the DM handler reads a replacement table from block 1 of the disk and stores it in the handler.

When a bad sector error (BSE) or header validity error (HVRC) is detected during a read or write, the DM handler replaces the bad block with a corresponding good block from the replacement tracks. The bad-block replacement feature of RT-11 requires blocks 0 through 5 and tracks 0 and 1 of the last cylinder to be good. This procedure causes an I/O delay since the read/write heads must move from their present position on the disk to the replacement area, and back again.

If this I/O delay cannot be tolerated, the disk can be initialized without bad-block replacement. In this case, bad blocks are covered by .BAD files. Neither the bad blocks nor the replacement tracks will be accessed.

You determine at volume initialization time whether to cover bad blocks with .BAD files or to create a replacement table for them and substitute good blocks during I/O transfers. The advantage of using bad-block replacement is that it makes a disk with some bad blocks appear to have none. On the other hand, covering bad blocks with .BAD files fragments the disk. Because RT-11 files must be stored in contiguous blocks, this fragmentation limits the size of the largest file that can be stored.

Only BSE and HVRC errors trigger the DM handler's bad block replacement mechanism. If a bad block develops that is not a BSE or HVRC error, the disk must be reformatted to have this new block included in the replacement mechanism. Reformatting should detect the new bad block. Mark it so that it generates a BSE or HVRC error and add the block number to the bad-block information on the disk. The disk should then be initialized to add the bad block to the replacement table.

The monitor file cannot reside on a block that contains a BSE error if you are using bad-block replacement. If this condition occurs, a boot error results when you bootstrap the system. In this case, move the monitor so that it does not reside on a block with a BSE error. Further, the monitor file (and any handler files) must reside in physically contiguous blocks—none of the blocks can be in the replacement table.

2.3 DU (MSCP Disk Handler)

This section provides specific programming information for MSCP disk devices.

The DU handler for RT-11 supports any disk system using the Mass Storage Communications Protocol (MSCP) interface. All disks using MSCP appear the same to the host computer. Thus, a single RT-11 DU handler can access any kind of MSCP disk.

2.3.1 Support for Special Functions

The DU handler supports the following special functions:

Code	Name	Section	Action
377	SF.ARD	2.3.5.2	Read operation without doing bad-block replacement; returns definitive error data.
376	SF.AWR	2.3.5.2	Write operation without doing bad-block replacement; returns definitive error data.
373	SF.SIZ	2.3.2	Determine the size, in 256_{10} -word blocks, of a particular volume.
	SF.S16		<i>blk</i> argument for SF.SIZ to indicate 16-bit starting block.
	SF.S32		<i>blk</i> argument for SF.SIZ to indicate 32-bit starting block.
372	SF.TAB	2.3.6	Returns the MSCP translation table.
371	SF.OBY		Obsolete; replaced by SF.BYP (360).
367	SF.R32	2.3.5.3	Read with 32-bit block number.
366	SF.W32	2.3.5.3	Write with 32-bit block number.
360	SF.BYP	2.3.7	Provides direct MSCP access.

2.3.2 Determining Volume Size (SF.SIZ), Code 373

Special function SF.SIZ returns the volume size in the word pointed to by the *buf* parameter argument. For DU, this special function is enhanced over that provided in the DL, DM, DY, and LD handlers. SF.SIZ for DU can return a 32-bit value for the device volume size and is, therefore, appropriate for use with device volumes that contain more than 65K blocks.

The volume size returned by the enhanced SF.SIZ is determined by any partition mapping. If a partition is mapped to the unit to which the channel is opened, the returned volume size is calculated from the base of the mapped partition to the usable end of the volume. If, for example, you have mapped unit DU1 to partition 1, an SF.SIZ for DU1 returns a volume size from the base of partition 1 to the usable end of the volume. If you reference the first partition on the volume, SF.SIZ returns the usable size of the entire volume.

The following description of parameters lists any differences between those for returning a 16-bit volume size and those for the 32-bit volume size.

Macro Call: `.SPFUN area,chan,#SF.SIZ,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

- area* is the address of a 6-word EMT argument block
- chan* is the channel opened on the unit for which you want the volume size
- SF.SIZ* is code 373 or the name SF.SIZ if the program has been assembled with the distributed module SYSTEM.MLB
- buf* For 16-bit value, is the address of a 1-word buffer in which volume size is returned
For 32-bit value, is the address of a 4-word buffer that on return contains a 32-bit value for the volume size followed by a 32-bit value for the MSCP logical block number from which the volume size was calculated.
The low-order base bits contain the value 0 and the high-order base bits contain a value indicating the partition to which this unit is currently mapped. If the unit does not exist, SF.SIZ returns a hard error and the contents of *buf* are undefined
- wcnt* For 16-bit volume size, is 1. For 32-bit volume size, is 4
- blk* For 16-bit volume size, is 0, indicating subcode SF.S16. For 32-bit volume size, is 1, indicating subcode SF.S32

2.3.3 Obtaining the DU Device Status (STATU\$)

DU has a status word containing information about the last operation performed by the handler. The status word is called STATU\$ and is located at an offset from the base of DU. See Table 2–1. The offset is stored in the handler as an entry in the table set up by the .DRTAB macro. The first word of the 2-word table entry is the RAD50 characters *UMS*, followed by the value of STATU\$. Using .DRTAB is described in the *RT-11 System Macro Library Manual*. The low 5 bits of STATU\$ contain the status information. All other bits are reserved.

Table 2–1: STATU\$ Status Information

Octal Value	Meaning
00	Success
01	Invalid command
02	Command aborted
03	Unit off line
04	Unit available
05	Medium format error
06	Write-protected medium

Table 2–1 (Cont.): STATU\$ Status Information

Octal Value	Meaning
07	Compare error
10	Data error
11	Host buffer access error
12	Controller error
13	Drive error

Use DBG–11, ODT/VDT, Console ODT, or the E keyboard command to examine the contents of STATU\$. You will need to perform customization patch 2.7.32 located in the *RT–11 Installation Guide* to use the E command. Use the SHOW MEMORY command display to find the base of the DU handler and add the offset to that base.

You can obtain the information returned in STATU\$ from within a program by calling the sytem subroutine, IGTDUS, as described in the *RT–11 System Subroutine Library Manual*.

2.3.4 Support for Bad-Block Replacement

All MSCP (DU) hard-disk systems support bad-block replacement (BBR), performed either by the disk controller or as a feature of the DU handler. For those MSCP hard disks for which BBR is provided by the controller, no support is required by the DU handler; bad-block replacement is transparent to RT–11.

In MSCP systems that use an RQDX1, RQDX2, or RQDX3 controller, BBR is performed by the controller. In those systems, BBR is done automatically by the hardware and does not require bad-block support in the DU handler.

In MSCP systems that use a KDA50, UDA50, KLESI–QA, or KLESI–UA controller, BBR can be performed by the DU handler.

Table 2–2 lists the MSCP controllers and drives supported by RT–11 and indicates whether bad-block replacement (BBR) is performed by the controller or the DU handler. (There is no BBR support for RX50 devices or write-only media.)

Table 2–2: MSCP Bad-Block Replacement (BBR)

MSCP Controller	Bad Block Replaced by:	MSCP Drive
RQDX1	controller	Supported RD-type drives
RQDX2	controller	Supported RD-type drives
RQDX3	controller	Supported RD-type drives
KLESI–QA	handler	Supported RC-type drives

Table 2–2 (Cont.): MSCP Bad-Block Replacement (BBR)

MSCP Controller	Bad Block Replaced by:	MSCP Drive
KLESI-UA	handler	Supported RC-type drives
UDA50	handler	Supported RA-type drives
KDA50	handler	Supported RA-type drives

The distributed DU for mapped monitors (DUX.SYS) supports handler BBR. If you are going to use an unmapped monitor with MSCP disks that require handler BBR, you should perform a system generation for that monitor and request support for DU handler bad-block replacement. Once you have generated such support, you can change monitors and continue DU handler bad-block replacement.

The following is general information on BBR as performed by DU:

- Bad-block replacement is a technique in which substitute blocks are provided for blocks that have caused a read or write error. The replacement blocks appear to occupy the disk positions of the original blocks, and the disk appears to contain only good blocks. You can force bad-block replacement on a device by performing a read and verify operation on all blocks. You perform such a read/verify operation by issuing a `FORMAT/VERIFY:ONLY` command for the device.
- Whether bad-block replacement is performed by the controller or the handler, it has the effect of making a disk appear to be error free. In certain cases, however, an I/O operation, a verification procedure, or a bad-block search may report the presence of bad blocks on a disk with replaced blocks. In such cases, any block identified as a bad block should be considered to be a good block with bad data. This means that the controller or handler provided a replacement block for a defective block but was unable to recover the data it contained.
- You can force MSCP class devices to clear bad blocks that contain soft errors by coupling the `DUP /H` option with the `/B` or `/K` option. The `/H` option is not available as a `KMON` command. You should use only the `DUP /H/B` or `/H/K` command options with blank media or a volume you have just backed up.
- If the DU handler is unable to replace a block on a device, DU displays the following error message:

```
?DU-E-Replace command failure or inconsistent RCT.  
?DU-E-Software write protecting volume.
```

If you receive that message, you should immediately back up that volume. Then check any file you had open for lost data. You cannot write to that volume again without first taking it off line and then placing it on line.

2.3.5 Non-File-Structured Read and Write Operations

DU supports three methods for performing non-file-structured read and write operations.

2.3.5.1 JREAD and JWRITE

You can perform absolute (non-file-structured access) reads and writes to any MSCP device, using the JREAD and JWRITE system subroutines. JREAD and JWRITE use a 32-bit starting block number, which lets you read and write to any block on any DU device. See the *RT-11 System Subroutine Library Manual* for details on JREAD and JWRITE.

2.3.5.2 Special Functions SF.AWR and SF.ARD

DU supports special functions SF.AWR (code 376) and SF.ARD (code 377). SF.AWR and SF.ARD are appropriate for devices that contain no more than 65K blocks. If the DU device contains more than 65K blocks, see Section 2.3.5.3. For DU, SF.AWR performs a write to the specified sector, and SF.ARD performs a read from the specified sector. Those writes and reads are not absolute; bad-block replacement and block vectoring remain in force.

Special functions SF.AWR and SF.ARD are especially useful because they return status information in the first word of the return buffer. Status information includes any occurrence of a bad-block error, forced error, or drive error. No discrimination for such errors is returned by a .WRITE or .READ request.

DU support for SF.AWR and SF.ARD is the same as DM with the following exceptions:

- DU supports an additional error code:

Code	Name	Meaning
140000	ES.FRC	A forced error occurred. If the device is a disk drive that supports BBR, the device controller or DU handler discovered bad data on a good (replaced) block. (Bad-block replacement was performed but no data was recovered.) If the device does not support BBR, this is an unexpected condition.

- For DU, bad-block replacement and block vectoring remain in force.

2.3.5.3 Special Functions SF.R32 and SF.W32

DU supports two special functions that perform non-file-structured block reads (SF.R32, code 367) and writes (SF.W32, code 366) on devices that contain more than 65K blocks. Because these special functions perform non-file-structured operations, they should generally not be used to perform operations on any device partition that contains a file structure.

Special functions SF.W32 and SF.R32 perform the same operations as the JWRITE and JREAD functions; JWRITE and JREAD use special functions SF.W32 and SF.R32. JWRITE and JREAD are described in the *RT-11 System Subroutine Library Manual*.

CAUTION

SF.W32 can write data to the reserved blocks on your DU device, which can render your DU device useless, because those blocks contain the replacement control table (RCT). You should, therefore, always issue a special function SF.SIZ (373) to a DU device to determine the volume size, because SF.SIZ returns the size at the boundary between the usable logical blocks and the RCT. Writing data only up to the volume size returned by SF.SIZ ensures you will not write data into the RCT.

The format for these special functions is:

Macro Call: .SPFUN *area,chan,func,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]*

- area* is the address of a 6-word EMT argument block.
- chan* is a channel number for I/O in the range 0 to 376₈.
- func* is the symbol or numeric code value for the function to be performed:

Code	Name	Meaning
366	SF.W32	32-bit non-file-structured block write
367	SF.R32	32-bit non-file-structured block read

- buf* is the buffer address.
- wcnt* is the number of words to transfer. Valid values are 0 through 077777₈.

blk is the address of a 4-word argument block:

blk+0 is a 2-word (32-bit) starting block number for this request. The first word contains the low-order bits. The second word contains the high-order bits.

The correspondence between the starting block number and a particular block on a device is determined by any partitioning and unit mapping of the device:

If the device has not been partitioned, starting block 0 specifies physical (and logical) block 0 — the start of the device. Any starting block number is offset from physical block 0.

If the device has been partitioned, logical block 0 of partition 0 continues to contain physical block 0. However, the starting block 0 of this request, because of device partitioning, corresponds to logical block 0 of the unit opened on this channel. Any starting block number is offset from logical block 0 of the partition mapped to the unit. For example, if the channel is opened for a non-file-structured operation to unit DU1 and DU1 is mapped to partition 1 (block 200000₈, starting block 0 corresponds to physical block 200000₈ of this device).

If, for example, your device contains an RT–11 file structure in partition 0, which is mapped to DU0, you could ensure the integrity of that file structure by always performing non-file-structured operations above partition 0 on the device.

blk+4 on return, contains the number of words actually transferred

blk+6 is reserved

2.3.6 DU Translation Table (SF.TAB), Code 372

The DU translation table defines the correspondence between RT–11 unit numbers and MSCP unit numbers, ports, and partitions. The format of the table is given in Figure 2–2.

Special function SF.TAB (code 372) interacts with the translation table from an address contained in the *buf* argument of the SF.TAB call. You can read the contents of the translation table to the buffer or write the contents of the buffer to the table. Whether the SF.TAB request is a read or write operation is determined by the *wcnt* parameter argument. This procedure is explained in this section.

For RT–11 V5.4, changes were made in the structure of the DU handler translation table. The names of the offsets in the table and the size of the table was changed. All programs you write to access the information contained in the table should use the following offsets. All programs you have written should be changed to use the following offset names.

Beginning with RT–11 V5.5, you can build a DU handler that supports more than eight units. That affects the size of the translation table.

Figure 2–2: DU Handler Translation Table

RT–11 Unit 0	MSCP Unit Number	
	Port	Partition
RT–11 Unit 1	MSCP Unit Number	
	Port	Partition
⋮	⋮	⋮

Whenever an I/O request is passed to the DU handler, DU uses the RT–11 unit number as an index into this table, extracts the MSCP unit number, port, and partition that have been assigned to that RT–11 unit, and uses the information to access the proper disk.

Size of the Translation Table

The size of the DU translation table in the DU handler is related to the number of device units supported by DU. The DU handler can support up to 64₁₀ units. Therefore, the translation table can contain up to 64 table entries.

Structure of the Translation Table

The DU unit translation table consists of a table header followed by table entries. Previously, the DU unit translation table had no header. Now, the DU unit translation table has a header starting at offset DU.ID, which is a word containing the Radix–50 value for the characters DU.

DU.ID is followed by DU.NUM. The low byte of DU.NUM contains the number of entries in the table. The high byte of DU.NUM is reserved.

The structure of the rest of the table remains as before. However, the offset names you should use to specify elements of the table have changed. The following is the structure of the table with the changed offset names:

Table 2–3: MSCP (DU) Translation Table Header

Offset	Name	Meaning
0	DU.ID	Radix–50 value for characters DU
2	DU.NUM	Byte containing number of entries in table
3		Reserved
4	DU.ENT	The offset of the first table entry

Each table entry consists of 4 bytes. Digital recommends you use the symbol DU.ESZ to represent the 4-byte size of each entry.

Table 2–4: MSCP (DU) Translation Table Entry

Offset	Name	Meaning
0	DU.UNI	Physical MSCP unit number. The symbol DU\$Uxx=nnnnnn is the initial value for the translation table when the handler is assembled. In the symbol, <i>xx</i> is the octal RT–11 DU unit number (0-7 or 0-77) and <i>nnnnnn</i> is the MSCP unit number. The SET Dxx UNIT=nnnnnn command can subsequently change the value.
2	DU.PAR	Byte containing partition number. The symbol DU\$Axx=nnn is the initial value for the translation table when the handler is assembled. In the symbol, <i>xx</i> is the octal RT–11 DU unit number (0-7 or 0-77) and <i>nnn</i> is the partition number. The SET DU PART=nnn command can subsequently change the value.
3	DU.POR	Byte containing MSCP port (controller) number. The symbol DU\$Oxx=nnn is the initial value for the translation table when the handler is assembled. In the symbol, <i>xx</i> is the octal RT–11 DU unit number (0-7 or 0-77) and <i>nnn</i> is the MSCP port number. The SET DU PORT=nnn command can subsequently change the value.

Accessing the Translation Table

Before Version 5.5, the translation table access special function code SF.TAB (372) supported only eight units. The *wcnt* parameter for SF.TAB accepted two arguments, SF.TRD (1) to indicate a read of the table and SF.TWR (–1) to indicate a write to the table. The size of the table was fixed at eight entries. If the DU handler on your system continues to support only eight DU devices, you continue to read and write to the translation table as before.

However, if the DU handler on your system supports more than eight units, the SF.TAB special function accepts other values for the *wcnt* parameter to support the extended device units. For DU handlers that implement the extended device-unit feature, you indicate both a read or write operation and the size of the table you are reading and writing by specifying a positive or negative numeric argument for the *wcnt* parameter. A positive numeric argument indicates a read operation of the specified number of words from the DU translation table to the buffer. A negative number indicates a write operation of the specified number of words from the buffer to the DU translation table.

You can use the following procedure to read the translation table from a DU handler that supports extended device units into a buffer and write the translation table from a buffer to DU. The procedure assumes you want to verify or do not currently know the number of entries in the table.

1. A translation table entry is created for each supported unit. You can determine the number of entries by doing a read SF.TAB to return the table entry DU.NUM. DU.NUM is the low byte of the second word in the table and contains the octal number of table entries. Therefore, for the *wcnt* parameter, supply the argument +2, and for the *buf* parameter, point to a 2-word buffer.

2. The translation table header and each entry continue to contain two words. Therefore, you can then read the entire DU handler extended device-unit translation table by supplying the value $\text{HEADER}+(2*\text{DU.NUM})$ for the *wcnt* parameter. For example, if *DU.NUM* indicated 16 entries, the value to specify for *wcnt* to read the entire table would be $+(2+(2*16))$. The *buf* parameter would point to a buffer of the same size.
3. You could write the contents of the buffer to the DU handler by specifying the value $-(2+(2*16))$ for the *wcnt* parameter.

You can avoid the calculation process by specifying a buffer of 130_{10} words, which can hold the largest translation table.

2.3.7 Special Function Bypass (SF.BYP), Code 360

Special function SF.BYP bypasses all unit number translations and allows direct access to the MSCP port. For DU, SF.BYP (direct MSCP access) serves the same purpose as the MU handler's SF.BYP (direct TMSCP access).

The request syntax and parameter argument definitions for SF.BYP are as follows:

Macro Call: `.SPFUN area,chan,#SF.BYP,buf,wcnt,blk`

area is the address of a 6-word EMT argument block.

chan is a channel number in the range 0 to 376_8 .

SF.BYP is code 360 or the name SF.BYP if the program has been assembled with the distributed module SYSTEM.MLB.

buf is the address of the 52_{10} -word TMSCP area.

wcnt when nonzero, is the virtual address of a data buffer to send to the handler. That virtual address is translated to a physical address and placed in the buffer of the TMSCP area.

when zero, the buffer address in the TMSCP area is not altered

blk indicates whether the handler should perform retries:

1 = specifies retries

0 = specifies no retries

The buffer address in special function SF.BYP must point to a 52-word area in the user's job. The first 26 words are used to hold:

- A response packet length in bytes
- A virtual circuit identifier
- An end packet when the command is complete

The second 26 words are set up by the caller and contain:

- A length word (length of command)
- A virtual circuit identifier (must have octal 1 (001) in high byte)

- A valid MSCP command (48-byte command buffer)

Except for port initialization, the user program must do all command packet sequencing, error handling, and reinitialization when the bypass operations are complete. The format of the control block is shown below:

Word	Contents
0	Response Packet Length
1	Virtual Circuit ID (from UDA or QDA controller)
2	MSCP Response Buffer (24 words)
26	Command Packet Length (48 bytes)
27	Virtual Circuit ID (from host)
28	MSCP Command (24 words)
51	Last Word of MSCP Command Packet

2.3.8 Addressing an MSCP Disk

You identify an MSCP disk to the DU handler by specifying:

- The MSCP unit number, in the range 0 through 253
- The controller port number, in the range 0 through 3
- The disk partition number, in the range 0 through 255

As DU is distributed, you address a disk—DU0 through DU7, as desired—and the DU handler references the disks that have been assigned to those RT-11 unit numbers. You can perform a system generation and request extended device-unit support for DU, which lets you address up to 64₁₀ disks. See the *RT-11 System Generation Guide* for information.

The default port number is 0, the default partition number is 0, and the default unit numbers correspond to the RT-11 unit numbers. Thus, if no modifications or SET commands are made to the DU handler, an MSCP disk will be referenced exactly like any other RT-11 disk; DU0 will refer to disk unit 0, DU1 will refer to disk unit 1, and so on. However, the names DU0 through DU7 can be reassigned to the MSCP disks of your choice by specifying MSCP unit, port, and partition numbers. Each of these parameters is described below.

2.3.8.1 MSCP Unit Numbers

Traditionally, there has always been a one-to-one correspondence between a physical disk drive unit number and an RT-11 disk unit number. This one-to-one correspondence does not necessarily apply to disks using the MSCP interface. Neither is an MSCP disk controller limited to eight units, nor are the unit identifying numbers limited to the range 0 through 7. The MSCP unit number of a disk is defined by the unit number plug of the disk drive. Although MSCP disks on most RT-11 systems may never have a unit number plug greater than 7, MSCP unit

numbers can be in the range 0 through 253. The DU handler supports a 16-bit MSCP unit number, if required by the system configuration.

The relationship between an RT-11 unit number and an MSCP disk unit number is defined within the DU handler. Typically, any necessary assignments are made at system installation time by using a SET command in the following form:

```
SET DUn UNIT=x
```

For example, you might issue the SET command

```
SET DU7 UNIT=21
```

Any references to DU7 would then go to MSCP unit number 21.

2.3.8.2 Controller Port Numbers

The controller port number provides a way of logically identifying the vector/CSR pair of a particular MSCP controller when your system has more than one.

You can access a second MSCP controller through the DU handler in one of two ways. One way is to create a second copy of the handler, as described in Section 2.3.9. You can then use the original DU handler to access disks connected through the first controller port, and the new copy of the handler to access disks connected through the second controller port. Although this procedure requires two copies of the handler, it allows totally independent operation of the two ports, giving maximum I/O throughput.

The second way is to configure the DU handler for multiple ports by defining the conditional assembly parameter DU\$PORTS=*n*. If memory space is at a premium, this may be your best choice. However, the ports will not operate independently and I/O throughput may be slower. If a request is pending for a disk interfaced through port 0, any requests for a disk interfaced through port 1 must wait for the port 0 I/O to complete. The DU handler supports up to four ports, numbered 0 through 3. CSR and vector values for each port can be assigned with SET commands in the following form:

```
SET DU VECTOR=nnnnnn  
SET DU VECx=nnnnnn
```

```
SET DU CSR=mmmmmm  
SET DU CSRx=mmmmmm
```

The value for *x* can be 2, 3, or 4.

If you configure the DU handler for multiple ports, you must specify the port number when you assign an RT-11 unit number to a disk interfaced through a port other than 0. You can do this with a SET command in the following form:

```
SET DUn PORT=x
```

For example, you might issue the SET command:

```
SET DU7 PORT=1
```

This command might be combined with an MSCP unit number assignment:

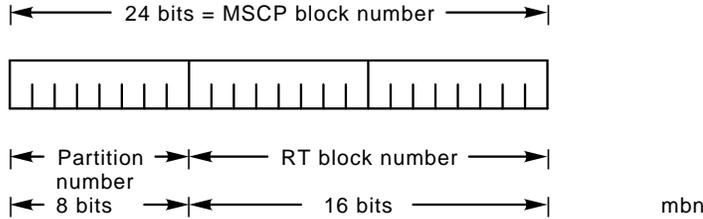
```
SET DU7 UNIT=21,PORT=1
```

You can perform a system generation and request support for multiport booting, as described in Section 2.3.10.

2.3.8.3 Disk Partition Numbers

Disk partition numbers allow RT-11 to use disks having more than 65,535 blocks. The disk partition number can be thought of as a high-order block number, as shown in Figure 2-3.

Figure 2-3: MSCP Disk Block Number



If a disk has more than 65,535 blocks, the DU handler divides the disk into logical partitions of 65,535¹ blocks each. The DU handler supports up to 256₁₀ disk partitions. Therefore, the largest disk DU can access has 256*65,535 blocks. To an RT-11 user, such a disk would appear to be 256 separate 65,535-block disks, each disk having its own directory.

Because the DU handler stores the partition numbers as bytes, DU supports an MSCP block number of no more than 24 bits, even though full MSCP supports block numbers of up to 32 bits. However, the partition number entries in the DU handler's translation table could be expanded to word entries if desired and 32-bit block numbers supported with no particular difficulty. Refer to Section 2.3.1 for details of the format of the DU handler's translation table.

Partition numbers are assigned with a SET command in the following form:

```
SET DUn PART=x
```

For example, you might issue the SET command

```
SET DU3 PART=1
```

This command could be combined with unit and port assignments as well:

```
SET DU3 UNIT=2, PORT=0, PART=1
```

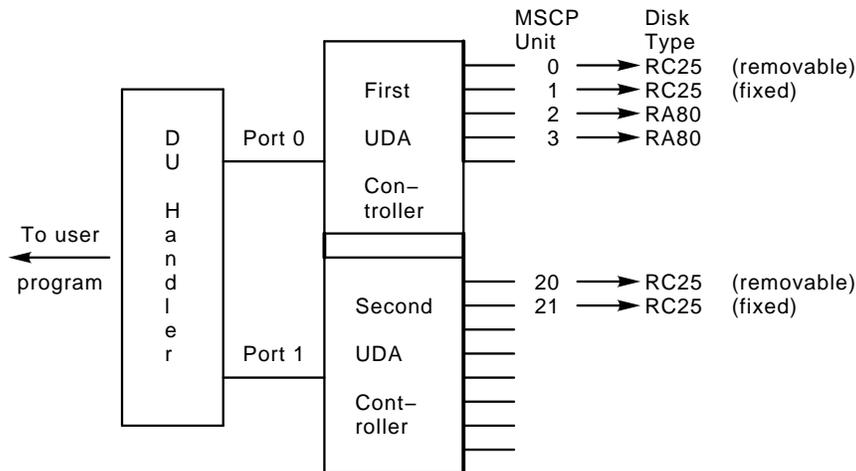
¹ Although RT-11 block numbers can be 0 through 177777₈, or a total of 65,536₁₀ blocks (200000₈, or 000000 in 16 bits since the 17th bit is lost), the size of a partition is defined as 65,535₁₀ blocks (177777₈), with RT-11 block numbers 0 through 177776. This avoids the problem of 16-bit overflow when dealing with the partition size. Because the partition number is added onto the left of the RT-11 block number to give the MSCP block number, one block between each partition is unused. Refer to the list below for the block numbers of the first three partitions:

Partition	Block Numbers
0	000000-177776, block 177777 unused
1	200000-377776, block 377777 unused
2	400000-577776, block 577777 unused

The mnemonic DU3 will then refer to the MSCP disk with unit plug 2 interfaced through port 0, beginning at block 65,536 of the disk (partition 1).

An example using several disks may help to clarify these concepts. Consider the example of a system with two UNIBUS Disk Adaptor (UDA) controllers interfaced to six disks, shown in Figure 2-4.

Figure 2-4: Two-Port DU Handler



The user of the system illustrated issues the following SET commands:

```
SET DU0 UNIT=0,PORT=0,PART=0
SET DU1 UNIT=1,PORT=0,PART=0
SET DU2 UNIT=2,PORT=0,PART=0
SET DU3 UNIT=2,PORT=0,PART=1
SET DU4 UNIT=3,PORT=0,PART=0
SET DU5 UNIT=3,PORT=0,PART=1
SET DU6 UNIT=20,PORT=1,PART=0
SET DU7 UNIT=21,PORT=1,PART=0
```

These commands assign DU0 to the first (removable) disk of the RC25 with MSCP unit number 0, and DU1 to the fixed disk of the RC25, identified as MSCP unit number 1. The disk unit with MSCP unit number 2 is an RA80, which has more than 65,535 blocks. Therefore, the next commands assign DU2 and DU3 to partition 0 and partition 1 of this disk, respectively. DU4 and DU5 are assigned in similar fashion to partitions 0 and 1 of the RA80 with MSCP unit number 3. Another RC25, interfaced to the second port of the UDA controller, is identified by MSCP units 20 and 21. The last two SET commands assign DU6 and DU7 to the two disks of this RC25 disk system. See Table 2-3 for information on setting up the default settings.

2.3.9 Creating a Second DU Handler

You can create a second DU handler under all monitors. The procedure is different for unmapped or mapped monitors.

2.3.9.1 Under Unmapped Monitors

You cannot run multiple DU handlers through the same MSCP controller; each handler must have a separate controller. Copy the handler to another file name and then modify the new file. Use the handler SET commands to change the vector and CSR of the copy to the values for the second port. For example, you could copy DU.SYS to DA.SYS and use the following SET commands to change the CSR and vector of the DA file:

```
SET DA VEC=nnnnnn  
SET DA CSR=mmmmmm
```

The variables *nnnnnn* and *mmmmmm* are the vector and CSR addresses of the second port.

2.3.9.2 Under Mapped Monitors

You cannot run multiple DU handlers through the same MSCP controller; each handler must have a separate controller. You can use the following procedure to create a second DU handler that can be used together with the distributed DU handler under all mapped monitors:

1. If you intend to perform a system generation to build a DU handler with support for extended device units or for any other reason, you must do that before creating a second DU handler. You must also preserve the system generation work files.
2. The second DU handler must be assigned a name that does not conflict with any distributed handler. If the second DU handler will be assembled for extended device-unit support, the first letter of the second DU handler cannot be D or L. For the purpose of this procedure, the second DU handler is named BU. Therefore, copy the DU source file to BU:

```
.COPY DU.MAC BU.MAC RET
```

3. Unprotect BU.MAC and open it with the editor.
4. Perform a search operation for the symbol DU\$NAM and on the other side of the equal sign, change the string <^RDU > to <^RBU >, so that the entire line of code resembles the following:

```
.IIF NDF DU$NAM,          DU$NAM = <^RBU >
```
5. Exit from the editor.
6. If you used the system generation procedure to build the DU handler, use the following procedure to assemble and link BU.MAC. If you did not build the DU handler and are using the distributed DU, proceed to step 7.
 - a. Copy the device-build (.DEV) command file that was created during the system generation to a file named BU.DEV.

- b. Open the file BU.DEV on the editor.
- c. Perform a search operation for the string +DU. The search places the cursor near the end of the first of three command lines that pertain to DU. The three command lines begin with MACRO, LINK, and SETOVR.

By placing an exclamation mark (!) character at the beginning of each line, comment out all command lines except the initial commands that assign device logical names and the three command lines that apply to DU.

- d. On the command lines that assemble and link DU, change all references from DU to BU, by replacing the *D* with *B*.
- e. Exit from the editor.
- f. Issue the following command to run BU.DEV as a command file:

```
.$@BU.DEV 
```

BU.DEV builds the file BUX.SYG.

- g. When BU.DEV has completed, copy the file BUX.SYG to BUX.SYS.
- h. Determine the current CSR and vector addresses for DU, using the following command:

```
.SHOW DEV:DU 
```

The MSCP port characteristics, such as CSR and vector addresses, for DU and BU cannot overlap. Specify addresses for BU that do not conflict with DU by using appropriate SET commands.

- 7. If you did not build DU by using the system generation process, issue the following commands to assemble and link BU. In the commands, *ddn* represents that device on which the distributed system conditional file (such as XM.MAC), the created file, BU.MAC, and the system library SYSTEM.MLB reside:

```
.ASSIGN ddn: SRC   
.MACRO/OBJ:BUX ddn:(XM+BU)   
.LINK/NOBITMAP/EXE:BUX.SYS/BOUNDARY:512. DK:BUX   
Boundary? SETOVR 
```

- 8. Determine the current CSR and vector addresses for DU, using the following command:

```
.SHOW DEV:DU 
```

Specify addresses for BU that do not conflict with DU by using appropriate SET commands.

2.3.10 Multiport Booting

During system generation, you can select an option for the DU handler that will let you boot RT-11 from any DU port. If you do not specify DU multiport booting during SYSGEN, you can boot RT-11 from DU port 0 only. Use the following procedure to enable multiport DU booting:

1. Use the SET DUn commands to map the particular DU device to the MSCP unit, port, and partition numbers. For example:

```
.SET DU3 UNIT=0, PORT=1   
.SET DU4 UNIT=1, PORT=1   
.SET DU5 UNIT=2, PORT=1 
```

For the SET commands to take effect, you must UNLOAD and then LOAD the handler if it is a data device or reboot it if it is a system device.

2. Copy the resulting DU handler to the port on the DU devices you want to be able to boot. For example:

```
.COPY DUX.SYS DU3: 
```

3. To hard-boot the DU unit on a new port, use the COPY/BOOT command to copy the bootstrap to the volume on the desired port. The DU unit on that port will also support the soft-boot BOOT DUn: command.

2.4 DW (CTI Bus-based Disk Handler)

This section provides specific programming information for the hard disks on CTI Bus-based computers.

2.4.1 Support for Special Functions

The DW handler supports the following special functions:

Code	Name	Action
377	SF.ARD	Read
376	SF.AWR	Write
373	SF.SIZ	Return device size

The special function (.SPFUN) request has the following general form, with the *area* and *chan* parameters and the optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual* and the other parameter arguments as described below:

Macro Call: `.SPFUN area,chan,func,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

<i>func</i>	is the special function code or the name if the program is assembled with the distributed file SYSTEM.MLB.
<i>buf</i>	For SF.ARD and SF.AWR, is the address of a 256 ₁₀ -word buffer. For SF.SIZ, is the address of a one-word buffer in which the size of the volume is returned.
<i>wcnt</i>	For SF.ARD and SF.AWR, is the track to read or write.
<i>blk</i>	For SF.ARD and SF.AWR, is the logical block (rather than physical block) to be read or written. Because the physical block number for DW is one less than the logical block number, address physical block 0 as logical block -1. For SF.SIZ, should be set to 0.

2.5 DX and DY (Diskette Handlers)

This section provides specific programming information for RX01 and RX02 diskettes.

As distributed, DX and DY support one controller that supports two drives. Each DX and DY handler can support two controllers (and therefore four drives). For example, if the RX01 handler is created through system generation to support two controllers, it will support four devices: DX0, DX1, DX2, and DX3. DX0 and DX1 are drives 0 and 1 of the standard diskette at CSR 177170 and vector 264. DX2 and DX3 are drives 0 and 1 of the other controller (standard alternate address CSR 177150 and vector 270). Note that only one I/O process can be active at one time, even though there are two controllers. Overlapped I/O to the handler is not permitted.

Data is stored on DX and DY diskettes in sectors. Double-density diskette sectors are 128 words long. RT-11 normally reads and writes them in groups of two sectors. Single-density diskette sectors are 64 words long. RT-11 reads and writes them in groups of four sectors. However, special function requests for absolute reads and writes can access sectors individually.

2.5.1 Support for Special Functions

The DX and DY handlers support the following special functions:

Code	Name	Action
377	SF.ARD	Read absolute sector
376	SF.AWR	Write absolute sector
375	SF.WDD	Write absolute sector with deleted data mark
373	SF.SIZ	Return device size, in 256 ₁₀ -word blocks (DY only)

A request to write absolute blocks should not write anything in track 0 if you want to use DUP or the COPY/DEVICE command to back up the volume. DUP does not copy data in track 0. Also, be sure you specify a valid buffer address and word count. The monitor checks that the *buf* parameter argument is in the job area, but it does not check the validity of *buf+(2*wcnt)-1*.

The special function (.SPFUN) request has the following general form, with the *area* and *chan* parameters and the optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual* and the other parameter arguments as described below:

Macro Call: `.SPFUN area,chan,func,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

func is the code for the function to be performed, or the name of the function if the program has been assembled with the distributed module SYSTEM.MLB.

- buf* For SF.ARD, SF.AWR, and SF.WDD, is the location of a 129-word buffer (for double-density diskettes) or a 65-word buffer (for single-density diskettes). The first word of the buffer, the flag word, is normally set to 0.
- The flag word set to 1 indicates a read on a physical sector containing a deleted data mark. The data area of the buffer extends from the second word to the end of the buffer.
- buf* for SF.SIZ is the location of a one-word buffer in which 494 is returned by single-density diskettes and 988 is returned by double-density diskettes.
- wcnt* For SF.ARD, SF.AWR, and SF.WDD, is the absolute track number, 0 through 76, to be read or written.
- wcnt* for SF.SIZ is reserved and should be set to 1
- blk* For SF.ARD, SF.AWR, and SF.WDD, is the absolute sector number, 1 through 26, to be read or written.
- blk* for SF.SIZ is reserved and should be set to 0.

The diskette should be opened with a non-file-structured .LOOKUP. The following example performs a synchronous sector read from track 0, sector 7, into a 65-word area called BUFF.

```
.SPFUN #RDLIST,#SF.ARD,#BUFF,#0,#7,#0
```

2.6 DZ (Diskette Handler)

This section provides specific programming information for diskettes on CTI Bus-based computers.

2.6.1 Support for Special Functions

The DZ handler supports the following special functions:

Code	Name	Action
377	SF.ARD	Read absolute sector
376	SF.AWR	Write absolute sector

The special function (.SPFUN) request has the following general form, with the *area* and *chan* parameters and the optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual* and the other parameter arguments as described below:

Macro Call: `.SPFUN area,chan,func,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

func is the code for the function or the name of the function if the program is assembled with the distributed file SYSTEM.MLB.

wcnt is the track to be written.

blk is the sector.

buf is the address of a 256₁₀-word buffer.

The .SPFUN requests do not interleave sectors. RX50 diskettes have 80 tracks, and the .SPFUN requests wrap to track 0 after track 79.

2.7 LD (Logical Disk Handler)

This section provides specific programming information for logical disks.

The Logical Disk handler implements logical disk support. The LD handler accepts I/O requests just like any other disk handler. By means of embedded translation tables, the LD handler determines which physical disk and which starting block offset should be used for each LD I/O request. When the proper physical disk and block number are determined, the LD handler updates the block number and unit number in the I/O queue element so that they correspond to the values for the assigned physical disk. The LD handler then places the queue element on the I/O queue for the physical disk so that the actual I/O can take place.

In addition to operating as outlined above, the LD handler can also be run as a program. When run, the LD handler accepts CSI command lines and switches to initialize, assign, verify, write-enable, or write-lock logical disk units.

2.7.1 Support for Special Functions

The logical disk handler supports the following special functions:

Code	Name	Action
372	SF.TAB	Access the translation tables
373	SF.SIZ	Return unit size. The parameter arguments for SF.SIZ for LD are the same as for DM. See Section 2.2 for information

2.7.2 LD Translation Tables (SF.TAB), Code 372

Special function SF.TAB (code 372) interacts with the translation tables from an address contained in the *buf* parameter argument of the SF.TAB call. You can read the contents of the translation tables to the buffer or write the contents of the buffer to the tables. Whether the SF.TAB request is a read or write operation is determined by the *wcnt* parameter argument. This procedure is explained in this section.

For RT-11 V5.4, changes were made to the structure of the LD translation tables. All programs you write to access the information contained in those tables should reflect the changes. All programs you have written to access LD translation tables should be changed to reflect the changes.

The tables start at a header; previously they started at a label. Following the 2-word header are four LD translation tables. That is unchanged. However, the names of offsets you use to reference the tables have changed. Some table contents have also changed.

Further, you can now build support for up to 64 logical disk units, which affects how you use the tables.

Size of the Translation Tables

The size of the LD translation tables in the LD handler is related to the number of logical disk units supported by LD. Beginning with Version 5.5, you can use

the system generation procedure (SYSGEN) to build an LD handler that supports extended device units. By default, SYSGEN builds support for 16₁₀ logical disk units when you request extended device-unit support. You can request up to 64₁₀ units. Of those 64 units, 32 can be mounted and 32 are reserved to Digital.

Structure of the Translation Tables

The LD translation tables consist of a 2-word header followed by four LD translation tables. The LD translation tables start at header LD.ID. Header LD.ID is a 1-word table identifier and contains the Radix-50 value for the characters LD. Header LD.ID is followed by LD.NUM, a 1-byte count of the number of entries in the table. As LD is distributed, the value in LD.NUM is 10₈, indicating eight table entries. If LD is built for extended device-unit support, the value in LD.NUM can contain a value up to 100₈, indicating support for up to 64 logical disk units. LD.NUM is the low-order byte of the word LD.ID+2. The high-order byte of LD.ID+2 is reserved.

The four LD translation table offset names, location, and contents are:

LD.FLG (LD.ID+4) The table beginning at offset LD.FLG is the table previously at the label HANDLR. LD.FLG contains one word for each LD unit number. The count of LD unit numbers is stored in LD.NUM. The bits in each word of LD.FLG have the following meaning:

Bits	Name	Meaning
0-5	LD.NDX	An index to the handler tables in RMON for the physical device corresponding to the LD unit number.
6	LD.UNX	A flag that signals the index entry (bits 0-5) may be inaccurate and should be updated. LD sets LD.UNX for all units if, upon entry, the LDREL\$ bit in RMON fixed offset CONFG2 is set.
7	LD.UOF	A flag that signals the entry in the LD.OFS table for that LD unit may be inaccurate. LD.UOF is set whenever a volume is squeezed. LD checks LD.UOF each time it uses an LD unit; if set, LD verifies that unit's LD.OFS table entry before proceeding.
8-13	LD.UNT	Contain the unit number of the physical disk assigned to the logical disk unit.
14	LD.RDO	Is the write-lock bit. If LD.RDO set, the LD unit is read only.
15	LD.ACT	Is the allocation bit. If LD.ACT set, the LD unit is assigned. If LD.ACT clear, the LD unit is not assigned.

LD.OFS (LD.FLG+<2*Contents of LD.NUM>) The second translation table starts at the offset LD.OFS and contains one word for each LD unit number. The count of LD unit numbers is stored in LD.NUM. Each word in LD.OFS contains the offset in blocks from the beginning of the assigned physical disk to the start of the area on that physical disk assigned to that LD unit number.

LD.SIZ (LD.FLG+<4*Contents of LD.NUM>) The third translation table starts at offset LD.SIZ and contains one word for each LD unit number. The count of LD unit numbers is stored in LD.NUM. Each word in LD.SIZ contains the size in blocks of the area on the physical disk assigned to that logical disk unit.

LD.NAM (LD.FLG+<6*Contents of LD.NUM>) The fourth translation table starts at the label LD.NAM and contains four words for each LD unit number. The count of LD unit numbers is stored in LD.NUM.

The first word of each 4-word entry contains the Radix-50 2-character name of the physical disk that is assigned to that logical disk unit. That Radix-50 word must be the physical (not logical) device name without any unit number. DL is a valid physical device name; DK and DL1 are not valid.

The second, third, and fourth words of each entry contain the Radix-50 file name and file type assigned as the logical disk.

Accessing the Translation Tables

Before Version 5.5, the translation table access special function code SF.TAB (372) supported only eight units. The *wcnt* parameter for SF.TAB accepted two arguments, +1 to indicate a read of the table and -1 to indicate a write to the table. The size of each LD translation table was fixed at eight entries. Beginning with Version 5.5, if the LD handler on your system continues to support only eight logical disk units, you continue to read and write to the translation tables as before.

However, if the LD handler on your system supports more than eight units, the SF.TAB special function provides additional values for the *wcnt* parameter to support the extended device units. For LD handlers that implement the extended device-unit feature, you indicate both a read or write operation and the size of the table you are reading and writing by specifying a positive or negative numeric argument for the *wcnt* parameter. A positive numeric argument indicates a read operation of the specified number of words from the LD translation tables to the buffer. A negative number indicates a write operation of the specified number of words from the buffer to the translation tables. For example, a *wcnt* parameter argument of +16 reads 16 words, and an argument of -16 writes 16 words.

You can use the following procedure to read the translation tables from an LD handler that supports extended device units into a buffer and write the translation table from a buffer to LD. The procedure assumes you do not currently know (or want to verify) the number of entries in the table.

1. Entries are reserved in each translation table for the total number of logical disk units supported by the handler. The offset at which each table starts is determined by the number of supported units. Therefore, to determine the starting offset for each table within the four translation tables, you first determine how many logical disk units are supported by the handler.
2. You can determine the number of entries by doing a read SF.TAB to return the table entry LD.NUM. LD.NUM is the low byte of the second word in the table and contains the number of table entries. Therefore, for the *wcnt* parameter, supply the argument +2, and for the *buf* parameter, point to a 2-word buffer.

3. Once you have determined the number of supported logical disk units, you can use that value to perform read/write operations for the tables.
4. You can read the LD translation tables into memory by performing a single SF.TAB read operation. The number of words in the LD translation tables is two for the header (LD.ID plus LD.NUM), the value in LD.NUM for each of the first three tables and four times the value in LD.NUM for the fourth table:

$$2+7*(LD.NUM)$$

For example, if LD.NUM indicated 100_8 entries, the value to specify for *wcnt* to read the entire table would be $+450_{10}$. The *buf* parameter would point to a buffer of the same size.

You could write the contents of the buffer to the LD handler by specifying the value -450_{10} for the *wcnt* parameter.

2.7.3 Other Bits Used by the LD Handler

The LD handler uses bit 4 (LDREL\$) in CONFIG2, monitor fixed offset 370. This bit is set whenever a handler is unloaded or released. The LD handler checks this bit to see if a handler assigned to an LD unit has been removed from memory since it was last used. If the bit is set, the LD handler sets bit 7 in all the entries in the HANDLR table, then clears the LDREL\$ bit. When the LD handler begins to process an I/O request, the LD handler checks bit 7 for the requested LD unit. If bit 7 is set, the LD handler verifies that the handler for the disk assigned to that LD unit number is in memory, then clears the bit. The LD handler checks and clears bit 7 for a unit only when an I/O request is sent to that unit. Checking only when absolutely necessary ensures that the LD handler will not waste time verifying units that may never be used by a particular user program.

2.8 MM, MS, and MT (Magtape Handlers)

This section provides specific programming information for reel-type magnetic tape devices.

Magnetic tape (magtape) has a sequential (not random-access) file structure. There is no directory at the beginning of each tape. RT-11 magtape handlers support a file structure that is compatible with ANSI tape labels and format, giving you full access to the tape controller without concern for the specifics of the device. See *RT-11 Volume and File Formats Manual* for more information on the format of magtapes and tape labels.

NOTE

Support for RT-11 magtape file structure is compatible only among systems that support DEC and ANSI standards for tape labels and file formats. DOS-formatted tapes cannot be read or written.

See the *RT-11 Commands Manual* for SET command conditions for each of the magtape handlers. Those conditions can set the number of tracks, the density, the parity of the tape drive, and the CSR and vector addresses.

See also the *RT-11 Master Index* under *Magtape* and the individual magtape handlers for more information.

2.8.1 File Structure Module (FSM)

The File Structure Module (FSM) creates the file structure on magtapes written by the distributed magtape handlers. The FSM is a discrete module (FSM.MAC) that is assembled with the magtape hardware handlers when handlers are built; it is included in the distributed magtape handlers. The FSM uses a protocol that is understood by RT-11 utilities and described in the *RT-11 Volume and File Formats Manual*.

When you issue a call for a file-oriented operation, the monitor (and perhaps the USB) builds a queue element and passes it to the FSM. The FSM processes the operation by manipulating the magtape drive.

Through the system generation procedure, you can build each of the magtape handlers without the FSM; a hardware-only version of each handler. A hardware magtape handler is smaller and requires less memory, but does not contain any routines that define a file structure. It does contain routines that manipulate the magtape drive. See Section 2.8.7.

Further, unless you write your own file structure module that duplicates the functionality of the FSM, RT-11 utilities do not understand whatever protocol you use to manage the magtape.

Therefore, Digital recommends that you use the distributed magtape handlers (unless special circumstances indicate that a handler without the FSM is appropriate), since only the handlers that contain the FSM can communicate with the RT-11 system utility programs.

This section uses some magtape-specific abbreviations:

BOT beginning-of-tape

EOF end-of-file

EOT physical end-of-tape

LEOT logical end-of-tape

LEOT consists of an EOF1 label (which includes one tape mark) followed by two tape marks.

2.8.2 Compatibility of Magtape Operations with the FSM

As briefly explained above, the distributed magtape handlers contain the basic magtape hardware handler, which is assembled with a file structure module (FSM). As shown in the following tables, some magtape operations are intercepted by the FSM and some operations bypass the FSM and are processed directly by the basic magtape hardware handler.

Although the distributed magtape handlers can process all the magtape operations described in this section, performing hardware-oriented operations that are incompatible with the FSM disrupts the magtape's file structure and can make the magtape unsuitable for further file-oriented operations. In other words, to preserve the file-oriented nature of a magtape volume, perform only file-oriented operations on that volume or other operations that are compatible with the FSM.

The operations you can perform on a magtape can be divided into three classes:

- Operations that use the FSM. These are file-structured operations that require the distributed handlers.
- Operations that bypass the FSM but are compatible with the FSM. These are non-file-structured operations that the FSM understands.
- Operations that bypass the FSM and produce a magtape that is incompatible with the FSM. You can perform these operations with the distributed handlers but the resulting magtape is not compatible with the FSM or any RT-11 utilities.

The following tables list magtape operations and their compatibility with the FSM. The tables list where more information can be found for each operation.

Table 2-5: Magtape Operations That Use the FSM

Operation	Section	Description
FSM Search by Sequence Number	2.8.4.1	Search for a file on a magtape based on file's sequence number.
FSM Search by File Name	2.8.4.2	Search for a file on a magtape based on the file name.
.ENTER	2.8.4.3	Open a file.

Table 2–5 (Cont.): Magtape Operations That Use the FSM

Operation	Section	Description
.LOOKUP	2.8.4.4	Find a file.
.READx	2.8.4.5	Read from a file.
.WRITx	2.8.4.6	Write to a file.
.CLOSE	2.8.4.8	Close a file.
.PURGE	2.8.4.9	Delete entry and close channel.

Table 2–6: Magtape Operations That Are Compatible with the FSM

Operation	Code	Section	Description
NFS .LOOKUP	N/A	2.8.5.1	Open a channel to a device (non-file-structured .LOOKUP operation). Required before any special function.
SF.USR	354	2.8.5.2	After NFS .LOOKUP, can be used in the following ways: Perform asynchronous directory operations that do not require the USR. Emulate a file-structured .LOOKUP or .ENTER to gain access to a file for further special function operations.
SF.MRD	370	2.8.5.3	After initial NFS .LOOKUP and SF.USR, perform read operations of variable length blocks.
SF.MWR	371	2.8.5.4	After initial NFS .LOOKUP and SF.USR, perform write operations of variable length blocks.
SF.MST	367	2.8.5.7	After initial NFS .LOOKUP, stream TS05 (MS only).
.CLOSE	N/A	2.8.5.6	Close channel and make device available.

Table 2–7: Magtape Operations That Are Not Compatible with the FSM

Operation	Code	Section	Description
SF.MOR	372	2.8.6.1	Rewind and place drive off line.
SF.MRE	373	2.8.6.2	Rewind.
SF.MWE	374	2.8.6.3	Write with extended gap.
SF.MBS	375	2.8.6.4	Backspace.
SF.MFS	376	2.8.6.5	Forward space.
SF.MTM	377	2.8.6.6	Write tapemark.

Table 2–7 (Cont.): Magtape Operations That Are Not Compatible with the FSM

Operation	Code	Section	Description
NFS .READx	N/A		Obsolete. Non-file-structured read operation (use SF.MRD).
NFS .WRITx	N/A		Obsolete. Non-file-structured write operation (use SF.MWR).

2.8.3 Spacing Error Recovery

Any errors detected during spacing operations abort the recovery attempt, and generate a hard (position) error.

Magtape handlers both with or without the FSM perform the following operations if a read parity error is detected.

1. Backspaces over the block and rereads. When unsuccessful, the procedure is repeated until five read commands have failed.
2. Backspaces five blocks, spaces forward four blocks, then reads the record.
3. Repeats steps 1 and 2 eight times or until the block is read successfully.

The handler performs the following operations upon detection of a read after write (RAW) parity error.

1. Backspaces over one block.
2. Erases 3 inches of tape and rewrites the block. In no case is an attempt made to rewrite the block over the bad spot, since, even if the attempt succeeds, the block could be unreliable and cause problems later.
3. Repeats steps 1 and 2 if the read after write still fails. When 25 feet of erased tape have been written, a hard error is given.

2.8.4 Magtape Operations That Use the FSM

The following magtape operations, listed in Table 2–5, use the FSM. The distributed magtape handlers support these operations.

2.8.4.1 FSM Searching by Sequence Number

The FSM can search for files on tape based on their sequence number. It uses the relationship between the current tape position and the desired new position to find the desired file according to the following algorithm:

1. When the file sequence number for the desired file is greater than the number of the current position, the handler moves the tape forward.

For example, if the tape is currently positioned at file sequence number 1, and the desired file is number 2, the tape moves forward from its position at the tape mark after file number 1 to the tape mark at the start of file number 2.

2. When the file sequence number for the desired file is less than the number of the current position, the handler optimizes its seek time by moving the tape backward or forward, depending on the location of the file. In practice, the handler almost always rewinds the tape and then searches forward.

For example, assume the number of the current position is 2 and the desired file has sequence number 1. The tape leaves its position at the tape mark for file 2 and rewinds to the beginning of the volume. It then moves forward to the tape mark at the start of file 1. As another example, assume the current position is 9 and the desired file has sequence number 6. The tape rewinds to the beginning of the volume and the search proceeds in the forward direction.

If you release the handler through the UNLOAD command or the .RELEASE programmed request, the file position is lost. In this situation the tape moves backward until the handler locates BOT or a label from which it can determine the tape's position.

2.8.4.2 FSM Searching by File Name

The FSM can search for files on tape based on their file names. The routine to match file names uses an algorithm that enables the handler to recognize file names and file types used by other Digital operating systems. The FSM uses the file identifier field, translating the contents to a recognizable file name. This file name is matched to a file name stored in Radix-50 format. The format is as follows:

filnam.typ

filnam is a valid RT-11 file name left-justified in a six-character field. Unused character positions are not padded.

typ is a file type left-justified in a 3-character field.

The algorithm the handler uses is backward compatible across all versions of the operating system. RT-11 tapes can be detected by the presence of *RT11* in character positions 64 through 67 of the HDR1 label. The algorithm is as follows:

1. Clear the character count (CC).
2. Check the next character in the file name. If it is a dot, do the following:
 - a. Mark a dot found.
 - b. When $CC < 6$, insert spaces and increment the CC until it equals 6.
 - c. When $CC > 6$, delete characters and decrement the CC until it equals 6.
3. If $CC = 6$ and if *RT11* is found in character positions 64 through 67 of the system code field, insert a dot in the translated name, mark the dot found, and increment CC.
4. Move the character into the translated file name and point to the next character.
5. Increment the CC.
6. When $CC < 10_{10}$ go back to step 2.

7. Check the dot-found indicator. If no dot was found, back up four characters and insert .DAT for the file type.
8. Perform a character-by-character comparison between the desired file name and the file name that was just translated from the file identifier field in the HDR1 label. When they match exactly, consider the file found.

2.8.4.3 .ENTER Programmed Request

The .ENTER programmed request opens a file on a magtape by writing a HDR1 label and tape mark on the tape and leaving the tape positioned after the tape mark. The request initializes some internal tables and makes entries for the last block written and current block number. (The last block or file on tape is always the most recent one written.) Table 2–8 shows the sequence number values for .ENTER requests.

The .ENTER programmed request has the following format, with the *area*, *chan*, and *dblk* parameters as described in the *RT-11 System Macro Library Manual*. The *seqnum* parameter is described below.

Macro Call: .ENTER area,chan,dblk,,seqnum

Table 2–8: Sequence Number Values for .ENTER Requests

Seqnum Argument	File Name	Action Taken	Tape Position
>0	not null	Position at file sequence number and perform a .ENTER.	Found: ready to write. Not found: at LEOT; LEOT is an EOF1 label followed by two tape marks. LEOT is different from the physical end-of-tape.
0	not null	Rewind tape and search tape for file name. If found then give error. If not found then enter the file.	Found: before file. Not found: ready to write.
-1	not null	Position tape at LEOT and enter file.	Ready to write.
-2	not null	Rewind tape and search tape for file name. Enter file at found file or LEOT, whichever comes first.	Ready to write.
0	null	Perform a non-file-structured .LOOKUP.	Tape is rewound.

The .ENTER request returns the errors shown in Table 2–9.

Table 2–9: .ENTER Errors

Byte 52 Code	Meaning
0	Channel in use.

Table 2–9 (Cont.): .ENTER Errors

Byte 52	
Code	Meaning
1	Device full. EOT was detected while writing HDR1. Tape is positioned after the first tape mark following the last EOF1 label on the tape. No such job exists (system job support only).
2	Device already in use. Magtape already has a file open on that unit.
3	File exists, cannot be deleted.
4	File sequence number not found. Tape is positioned the same as for device full.
5	Invalid argument error. A <i>seqnum</i> argument in the range –3 through –32767 was detected. A null file name was passed to .ENTER.

The .ENTER request issues a directory hard error if errors occur while entering the file.

2.8.4.4 File-Structured .LOOKUP Programmed Request

A file-structured .LOOKUP request finds a file by searching for a specific HDR1 label. Upon finding and reading the HDR1 label, the tape is positioned before the first data block of the file.

The .LOOKUP request has the following format, with the *area*, *chan*, and *dblk* parameters as described in the *RT-11 System Macro Library Manual*. The *seqnum* parameter argument values are shown in Table 2–10:

Macro Call: .LOOKUP *area,chan,dblk,seqnum*

Table 2–10: Sequence Number Values for File-Structured .LOOKUP Requests

Seqnum Argument	File Name	Action Taken	Tape Position
>0	null	Perform a file-structured .LOOKUP on the file sequence number.	Found: ready to read first data block. Not found: at LEOT.
0	not null	Rewind to the beginning of tape, then use file name to perform a file-structured .LOOKUP.	Found: ready to read first data block. Not found: at LEOT.
–1	not null	Do not rewind; perform a file-structured .LOOKUP for a file name.	Found: ready to read first data block from the current position. Not found: at LEOT.
>0	not null	Position at file sequence number and perform a file-structured .LOOKUP. If file name does not match file name given, return error.	Found: ready to read first data block. Not found: at the beginning of the file specified by the sequence number.

The file-structured `.LOOKUP` returns the errors shown in Table 2–11.

Table 2–11: `.LOOKUP` Errors

Byte 52 Code	Meaning
0	Channel in use.
1	File not found. Tape is positioned after the first tape mark following the last EOF1 on the tape.
2	Device in use. Magtape already has a file open.
5	Invalid argument error. A <i>seqnum</i> argument in the range <code>-2</code> through <code>-32767</code> was detected. A <code>.LOOKUP</code> request must have a positive sequence number.
6	Invalid unit number.

The `.LOOKUP` request issues a directory hard error if errors occur while entering the file.

2.8.4.5 `.READx` Programmed Requests

In this section, the term `.READx` refers to the `.READ`, `.READC`, and `.READW` group of programmed requests. Further, `.READx` requests are described for files that have been opened with the `.ENTER` and file-structured `.LOOKUP` requests.

The `.READx` requests read data from magtape in blocks of 512 bytes each. If a request is issued for fewer than 512 bytes, the handler reads the correct number of bytes. If the request is for more than 512 bytes, the handler performs the request with multiple 512-byte transfers (the last request may be for fewer than 512 bytes).

The `.READx` requests are valid in a file opened with a `.LOOKUP` request. They are also valid in a file opened with an `.ENTER` request, provided the block number requested does not exceed the last block written. (Exceeding the last block written returns code 0.)

If a tape mark is read, the routine repositions the tape so that another request causes the tape mark to be read again. When a `.CLOSE` is issued to a file opened by an `.ENTER` request, the tape position is left unchanged. Because magtape is sequentially accessed, a reposition in a file (a backup) without subsequently positioning to the end of the file (before a `.CLOSE`) causes data loss.

The guidelines for block numbers are as follows:

1. When a `.LOOKUP` is used (to search the file) with this request, the handler tries to position the tape at the indicated block number. When it cannot, a 0 (EOF code) is issued, and the tape is positioned after the last block on the file.
2. On an entered file, `.READx` checks to determine if the block requested is past the last block in the file. If it is, the tape is not moved and the 0 error code is issued.

The `.READx` request has the following format, with the *area*, *chan*, *buf*, *wcnt*, *blk* and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.READx area,chan,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

Table 2–12 shows the errors the `.READx` requests return.

Table 2–12: .READx Errors

Byte 52 Code	Meaning
0	Attempt to read past a tape mark; also generated by block that is too large.
1	Hard error occurred on channel.
2	Channel not open.

2.8.4.6 .WRITx Programmed Requests

In this section, the term `.WRITx` refers to the `.WRITE`, `.WRITC`, and `.WRITW` group of programmed requests. Further, `.WRITx` requests are described for files that have been opened with the `.ENTER` and file-structured `.LOOKUP` requests.

The `.WRITx` requests write data to magtape in blocks of 512 bytes. If a request is issued for fewer than 512 bytes, the handler forces the writing of 512 bytes from the buffer address. If a request is issued for more than 512 bytes, the handler performs multiple 512-byte transfers.

The `.WRITx` requests are valid in a file opened with an `.ENTER`. Once a file is opened, `.WRITx` determines if the requested block is past the last block in the file. If it is, the tape is not moved and the 0 error code is issued.

The `.WRITx` request has the following format, with the *area*, *chan*, *buf*, *wcnt*, *blk* and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.WRITx area,chan,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

Table 2–13 shows the errors the `.WRITx` requests return.

Table 2–13: .WRITx Errors

Byte 52 Code	Meaning
0	End-of-tape. The data for the last write was not written, but the previous block is valid. Also issued if the block number is too large.
1	Hard error occurred on channel.
2	Channel not open.

After a write operation, the rest of the tape is undefined (see Figure 2–5).

Figure 2–5: Operations Performed After the Last Block Written on Magtape

In example 1 in Figure 2–5, blocks A, B, and C are written on the tape with the head positioned in the gap immediately following block C. Any forward operation of the tape drive except by write commands (that is, write, erase gap and write, or write tape mark) yields undefined results due to hardware restrictions.

In example 2 in Figure 2–5, the head is shown positioned at BOT after a rewind operation so that successive read operations can read blocks A, B, and C. The head is left positioned as shown in example 3. Note that this is the same condition as shown in example 1, and all restrictions indicated in example 1 are applicable.

2.8.4.7 .CLOSZ, .DELETE, .GFxxx, .RENAME, and .SFxxx Programmed Requests

These requests are invalid operations on magtape, and any attempt to execute them returns an invalid operation code (code 2) in byte 52.

2.8.4.8 .CLOSE Programmed Request

The action of the .CLOSE request depends on how the file was opened.

- When a file is opened with an .ENTER request, the file is closed by writing a tape mark, an EOF1 label, and three more tape marks. In this operation, the tape is left positioned just before the second tape mark at LEOT. Note that the rest of the tape is no longer readable.

- When a file is opened with a file-structured `.LOOKUP`, the tape is positioned after the tape mark following the EOF1 label for that file.

The `.CLOSE` request has the following format, with the *chan* parameter as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.CLOSE chan`

This request issues a directory hard error if a malfunction is detected. The error can be recovered with the `.SERR` request.

2.8.4.9 `.PURGE` Programmed Request

The action performed by a `.PURGE` request is determined by the following:

- If the magtape channel has been opened by a `.ENTER` request, a `.PURGE` request deletes the current entry by a series of `BACKUP` and `WRITE-TAPE-MARK` operations, leaving the magtape positioned just before the second tape mark at `LEOT`.
- If the magtape channel has been opened with a file-structured or non-file-structured `.LOOKUP`, the `.PURGE` request frees the unit table entry for the handler, closes the channel, and makes the handler available for other operations.

The `.PURGE` request has the following format, with the *chan* parameter as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.PURGE chan`

2.8.5 Magtape Operations That Are Compatible with the FSM

The following magtape operations (as listed in Table 2-6), bypass the FSM but are compatible with the FSM. The distributed magtape handlers support these operations and a magtape that is manipulated by these functions is supported by RT-11 utilities.

2.8.5.1 Non-File-Structured `.LOOKUP` Programmed Request

You must issue a non-file-structured `.LOOKUP` request to open a channel to the device before starting any I/O operations. The non-file-structured `.LOOKUP` request causes the handler's hardware level to mark the drive busy so that no other channel can be opened to that drive until a `.CLOSE` is issued.

The `.LOOKUP` request has the following format, with the *area*, *chan*, and *dblk* parameters as described in the *RT-11 System Macro Library Manual*. The values for the *seqnum* parameter argument are described in Table 2-14:

Macro Call: `.LOOKUP area,chan,dbl,seqnum`

Table 2–14: Sequence Number Values for Non-File-Structured .LOOKUP Requests

Seqnum Argument	File Name	Action Taken	Tape Position
0	null	Perform a non-file-structured .LOOKUP.	Rewound.
-1	null	Perform a non-file-structured .LOOKUP.	Not moved.

Table 2–15 shows the errors that can be returned by the non-file-structured .LOOKUP request.

Table 2–15: Non-File-Structured .LOOKUP Errors

Byte 52 Code	Meaning
0	Channel in use; channel already open.
1	File not found; no such job.
2	Device in use. The drive being accessed is already attached to another channel.
5	Argument is invalid; for example, magtape file sequence number.
6	Invalid unit number.

2.8.5.2 Asynchronous Directory Operations (SF.USR), Code 354

SF.USR must be preceded by a non-file-structured .LOOKUP and can be used to perform two operations:

- SF.USR can perform asynchronous directory operations without the USR, which makes it useful for long tape searches. It is particularly useful in multi-job environments, because the search operation locks the USR during directly issued .ENTER and .LOOKUP requests.
- SF.USR allows an emulation of the .ENTER and file-structured .LOOKUP requests to be issued after a non-file-structured .LOOKUP assigns a channel to the magtape handler.

The special function SF.USR has the following format, with the *area* and *chan* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.SPFUN area,chan,#SF.USR,buf,blk`

SF.USR is the code 354 or the name SF.USR if the program has been assembled with the distributed file SYSTEM.MLB.

buf

is the address of a 7-word block with the following format:

Word	Meaning
0–2	Radix–50 representation of the file name.
3	One of the following codes: 3 for .LOOKUP 4 for .ENTER
4	Sequence number value. See the corresponding sections for .LOOKUP or .ENTER for complete information on the interpretation of this value.
5,6	Reserved.

blk

is the address of a 4-word error and status block used for returning .LOOKUP and .ENTER errors that are normally reported in byte 52. See Section 2.8.5.5. Only the first word of *blk* is used by this request. The other three words are reserved for future use and must be zero. If the value of *blk* is 0, no error information is returned. Figure 2–6 shows a programming example.

Figure 2–6: Asynchronous Directory Operation Example

```
.TITLE Asynchronous Directory Operation Example
.ENABLE LC           ; Print lower case
.NLIST BEX          ; Don't list text storage
.MCALL .LOOKUP, .SPFUN, .CLOSE, .PRINT, .EXIT

; Definitions
SF.USR = -20.       ; Asynchronous request
LOOKUP = 3          ; Lookup code for async request
ENTER = 4           ; Enter code for async request
CHAN = 0            ; Use channel 0
FNF = 1             ; 1 = File not found error
FSN = 0             ; Use 0 as file sequence number

;Example assumes that magtape handler is loaded.
```

Figure 2–6 (continued on next page)

Figure 2–6 (Cont.): Asynchronous Directory Operation Example

```

START:  .LOOKUP #AREA,#CHAN,#NFSBLK,#0      ; Open a channel
        ; for the next request
        BCS     LOOKER      ; Branch if error occurred
        .SPFUN #AREA,#CHAN,#SF.USR,#COMBLK,#ERRBLK
        ; Do a lookup
        BCC     FILFND      ; Branch if file found
        CMP     #FNF,ERRBLK ; File not found error?
        BEQ     NOTFND      ; Branch if yes
        MOV     #ASYERR,R0   ; No, some other error
        BR      CLOSE
LOOKER: MOV     #LOOERR,R0   ; NFS Lookup error
        BR      CLOSE
FILFND: MOV     #OK,R0      ; Report success
        BR      CLOSE
NOTFND: MOV     #FNFERR,R0   ; Report file not found
CLOSE:  .PRINT                ; Print error pointed to
        ; by R0
        .CLOSE #CHAN        ; Clean up...
        .EXIT                ; and return to monitor

;Data area
AREA:   .BLKW    5          ; EMT argument block
NFSBLK: .RAD50   /MT /      ; Use this to open
        .WORD    0          ; magtape in non-file-
        .WORD    0          ; structured mode
        .WORD    0
COMBLK: .RAD50   /FILNAMTYP/ ; This is the file name
        ; we're looking for
        .WORD    LOOKUP     ; This is the asynch op
        ; code for lookup
        .WORD    FSN        ; This is file sequence
        ; number for the lookup
        .WORD    0,0        ; Reserved (must be 0)
ERRBLK: .WORD    1          ; Set first word non-0
        .WORD    0,0,0      ; so errors return here

;Messages
LOOERR: .ASCIZ   /Non-file-structured lookup failed/
OK:     .ASCIZ   /File found, lookup successful/
FNFERR: .ASCIZ   /File not found/
ASYERR: .ASCIZ   /Error in asynchronous request/
        .EVEN
        .END     START

```

2.8.5.3 Read Physical Blocks (SF.MRD), Code 370

After an NFS .LOOKUP request (and optionally after an SF.USR), the SF.MRD request reads blocks of any size.

The special function SF.MRD has the following format, with the *area*, *chan*, *buf*, *wcnt*, and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: .SPFUN *area,chan,#SF.MRD,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]*

SF.MRD is the code 370 or the name SF.MRD if the program is assembled with the distributed file SYSTEM.MLB.

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the errors shown in Table 2–16. Additional qualifying information for these errors is returned in the first two words of the *blk* parameter argument status block. See Section 2.8.5.5.

Table 2–16: SF.MRD (Code 370) Errors

Byte 52 Code	First Word Code	Qualifying Information
EOF (Value = 0)	1	Tape before EOF only (tape mark detected).
	2	Tape before EOT only (no tape mark detected).
	3	Tape before EOF and EOT (tape mark detected).
Hard error (Value = 1)	0	No additional information (consult documentation for your particular tape drive for all possible error conditions).
	1	Tape drive not available.
	2	The controller lost the tape position.
	3	Nonexistent memory accessed.
	4	Tape is write locked.
	5	The last block read had more information. The MM handler returns (in the second status word) the number of words not read.
	6	A short block was read. The second status word contains the difference between the number of words requested and the number read.

2.8.5.4 Write Physical Blocks (SF.MWR), Code 371

After an NFS .LOOKUP request and optionally after an SF.USR, the SF.MWR request writes blocks of any size.

The special function SF.MWR has the following format, with the *area*, *chan*, *buf*, *wcnt* and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: .SPFUN *area,chan,#SF.MWR,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]*

SF.MWR is the code 371 or the name SF.MWR if the program is assembled with the distributed file SYSTEM.MLB.

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the errors shown in Table 2-17.

Table 2-17: SF.MWR (Code 371) Errors

Byte 52 Code	First Word Code	Qualifying Information
EOF (Value = 0)	1	Tape before EOF only (tape mark detected).
	2	Tape before EOT only (no tape mark detected).
	3	Tape before EOF and EOT (tape mark detected).
Hard error (Value = 1)	0	No additional information (consult documentation for your particular tape drive for all possible error conditions.)
	1	Tape drive not available.
	2	The controller lost the tape position.
	3	Nonexistent memory accessed.
	4	Tape is write locked.

NOTE

The TJU16 tape drive can return a hard error if a write request with a word count less than 7 is attempted.

2.8.5.5 Exception (Error and Status) Reporting

Special function requests report end-of-file and hard error conditions through byte 52 in the system communication area. You can also receive additional information about those two error conditions. You can specify an address in the special function's *blk* parameter that points to a 4-word error and status block which returns that information.

Specify #0 for *blk* if you do not want exception reporting.

Although all four words in the error and status block must be initialized to 0 before the first special function is called, only words 1 and 2 of the status block return information. Words 3 and 4 are reserved and not written and therefore need only be initialized once (remain as set to 0).

The meaning of the error and status block contents is tied to the contents of byte 52 in the system communications area. The program should therefore check the state of the carry bit and byte 52 before attaching importance to the contents of the error and status block.

End-of-File Condition Exception Reporting

Besides an actual EOF, the magtape handler's hardware level returns an end-of-file condition when the handler encounters an EOT, tape mark, or BOT. An end-of-file condition produces the following:

- Sets the carry bit and byte 52 is zero.
- The first word of the error and status block is shown in Table 2–18.
- The second word contains the number of blocks not spaced when a tape mark is detected during a spacing operation.

Table 2–18: End-of-File Qualifying Information

First Word	Meaning
1	Tape before EOF only (tape mark detected).
2	Tape before EOT only (no tape mark detected).
3	Tape before EOT and EOF (tape mark detected).
4	Tape before BOT (no tape mark detected).

Hard Error Condition Exception Reporting

A hard error condition:

- Sets the carry bit and byte 52 is 1.
- Returns in the first word the qualifying information shown in Table 2–19.

Table 2–19: Hard Error Qualifying Information

First Word	Meaning
0	No additional information (includes parity error and all others not listed below. Consult documentation for your particular tape drive for all possible error conditions.)
1	Tape drive not available.

Table 2–19 (Cont.): Hard Error Qualifying Information

First Word	Meaning
2	The controller lost the tape position. When this error occurs, rewind or backspace the tape to a known position.
3	Nonexistent memory was accessed.
4	Tape is write locked.
5	The last block read had more information. The MM handler returns (in the second status word) the number of words not read.
6	A short block was read. The second status word contains the difference between the number of words requested and the number of words read.

2.8.5.6 .CLOSE Programmed Request

The magtape handler at the hardware level accepts the .CLOSE request and causes the handler to mark the drive as available; the channel becomes free.

The .CLOSE request has the following format, with the *chan* parameter as described in the *RT–11 System Macro Library Manual*:

Macro Call: .CLOSE *chan*

2.8.5.7 Enabling 100ips Streaming on a TS05/TSU05/TSV05 (SF.MST), Code 367

The SF.MST special function places the TS05 drive in 100ips streaming mode.

The special function SF.MST has the following format, with the *area* and *chan* parameters as described in the *RT–11 System Macro Library Manual*:

Macro Call: .SPFUN *area,chan,#SF.MST,buf,,blk*

SF.MST is the code 367 or the name SF.MST if the program is assembled with the distributed file SYSTEM.MLB.

buf is a word which enables or disables streaming.

If *buf* contains a 1, streaming is enabled.

If *buf* contains a 0, streaming is disabled.

blk is a pointer to a 4-word error block. (See Section 2.8.5.5.)

Streaming is automatically turned off when a .CLOSE is issued on a channel open on magtape, when an abort occurs, or if there is a magtape I/O error.

This special function is valid only for a TS05 using the MS handler. An SF.MST call is ignored if it is used with any other magtape handler or if it is used with the MS handler running a TS11 magtape.

If you want to run a TS05 in streaming mode, you must also use double-buffered I/O so that there is always a request pending in the magtape I/O queue. If there is not, there will be too much delay between I/O requests and the streaming will not work properly.

2.8.6 Magtape Operations That Are Not Compatible with the FSM

The magtape operations listed in Table 2-7 and described below bypass the FSM and are incompatible with the file structure produced by the FSM. The operations are direct hardware calls to the magtape handler. The distributed magtape handlers accept these operations, but a magtape that is manipulated by these functions is no longer ANSI-compatible or supported by RT-11 utilities.

When any of the following operations is called, the stored file sequence number and block number information are erased and are not reinitialized until a .CLOSE and another file-opening command have been performed. Note that the .CLOSE moves and, in the case of the file opened with .ENTER, writes the tape regardless of any commands that have been issued since the file was opened. When the file is closed, the magtape handler cannot write the size of the file because the file size is lost to the handler. It writes a zero in its place. The file sequence number field will be correct.

You initiate operations and use these special functions in the same manner as those that are compatible with the FSM:

1. Open a channel to the device by issuing a non-file-structured .LOOKUP.
2. You can optionally open a file on the magtape volume by issuing an SF.USR.
3. Issue the special functions to read, write, or position the magtape.
4. Close the channel.

If you are going to be using the operations in this section consistently, you should investigate performing a system generation and building a magtape handler that does not contain the FSM; a hardware-level-only handler. Such a handler is appropriate for the operations in this section and has a much smaller memory image. See Section 2.8.7 and the *RT-11 System Generation Guide* for information.

2.8.6.1 Rewinding and Going Off Line (SF.MOR), Code 372

This request is the same as rewind, except that it takes the tape drive off line and then rewinds to BOT. The handler is free to accept commands after the rewind is initiated.

The special function SF.MOR has the following format, with the *area*, *chan*, and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: .SPFUN *area,chan,#SF.MOR,,,blk[,crtn][,BMODE=str][,CMODE=str]*

SF.MOR is the code 372 or the name SF.MOR if the program is assembled with the distributed file SYSTEM.MLB.

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the same error code and qualifying information as the rewind request.

2.8.6.2 Rewinding (SF.MRE), Code 373

The SF.MRE request rewinds the tape to BOT. The MT and MM handlers cannot accept other requests until the rewind operation is complete; the MS handler can.

The special function SF.MRE has the following format, with the *area*, *chan*, and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.SPFUN area,chan,#SF.MRE,,,blk[,crtn][,BMODE=str][,CMODE=str]`

SF.MRE is the code 373 or the name SF.MRE if the program is assembled with the distributed file SYSTEM.MLB.

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the error shown in Table 2–20.

Table 2–20: SF.MRE (Code 373) Errors

Byte 52 Code	First Word Code	Qualifying Information
Hard error (Value = 1)	0	No additional information (consult documentation for your particular tape drive for all possible error conditions).
	1	Tape drive not available.

2.8.6.3 Writing with Extended Gap (SF.MWE), Code 374

This request permits you to write on tapes that have bad spots. The call syntax is identical to the SF.MWR request except for its function code, which is 374. The errors are explained in Table 2–21.

Table 2–21: SF.MWE (code 374) Errors

Byte 52 Code	Meaning
0	The EOT marker has been detected.
1	Hard error occurred on channel.
2	Channel not open.

Additional qualifying information for these errors is returned in the first two words of the status block. See Section 2.8.5.5.

2.8.6.4 Spacing Backward (SF.MBS), Code 375

The SF.MBS request spaces the magtape backward block-by-block or until a tape mark is detected.

You should note that because magtape is sequentially accessed, an SF.MBS operation in a file without a subsequent positioning to the end of the file (before a .CLOSE) causes data loss.

The special function SF.MBS has the following format, with the *area*, *chan*, *wcnt* and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: .SPFUN *area,chan,#SF.MBS,,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]*

SF.MBS is the code 375 or the name SF.MBS if the program is assembled with the distributed file SYSTEM.MLB.

wcnt is the number of blocks to space past (must not exceed 65534₁₀).

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the errors shown in Table 2–22.

Table 2–22: SF.MBS (Code 375) Errors

Byte 52 Code	First Word Code	Qualifying Information
EOF (Value = 0)	1	Tape before EOF only (tape mark detected).
	2	Tape before EOT only (no tape mark detected).
	3	Tape before EOF and EOT (tape mark detected).
	4	Tape before BOT (no tape mark detected). The second word in the status block contains the number of blocks requested to be spaced <i>wcnt</i> , minus the number of blocks spaced if a tape mark or BOT is detected. Otherwise, its value is not defined.
Hard error (Value = 1)	0	No additional information (consult documentation for your particular tape drive for all possible error conditions).
	1	Tape drive not available.
	2	The controller lost the tape position.

2.8.6.5 Spacing Forward (SF.MFS), Code 376

The SF.MFS request spaces the magtape forward block-by-block or until a tape mark is detected. When a tape mark is detected, the handler reports it along with the number of blocks not skipped. These commands can be used to issue a space-to-tape-mark command by passing a number greater than the maximum number of blocks on a tape. The tape is left positioned after the tape mark or the last block passed. The two spacing requests have the following forms.

The special function SF.MFS has the following format, with the *area*, *chan* and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.SPFUN area,chan,#SF.MFS,,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

SF.MFS is the code 376 or the name SF.MFS if the program is assembled with the distributed file SYSTEM.MLB.

wcnt is the number of blocks to space past (must not exceed 65534₁₀).

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the errors shown in Table 2–23.

Table 2–23: SF.MFS (Code 376) Errors

Byte 52 Code	First Word Code	Qualifying Information
EOF (Value = 0)	1	Tape at EOF only (tape mark detected).
	2	Tape at EOT only (no tape mark detected).
	3	Tape at EOF and EOT (tape mark detected). The second word in the status block contains the number of blocks requested to be spaced (<i>wcnt</i>), minus the number of blocks spaced if a tape mark or BOT is detected. (A tape mark is counted as a block.) Otherwise, its value is not defined. The tape will be positioned after the tape mark on forward spacing and before the tape mark on backward spacing.
Hard error (Value = 1)	0	No additional information (consult documentation for your particular tape drive for all possible error conditions).
	1	Tape drive not available.
	2	The controller lost the tape position.

NOTE

Due to hardware restrictions, Digital recommends that no forward space commands be issued if the reel is positioned past the EOT marker.

2.8.6.6 Writing a Tape Mark (SF.MTM), Code 377

The SF.MTM request writes a tape mark.

The special function SF.MTM has the following format, with the *area*, *chan* and optional *crtn*, *BMODE=str*, and *CMODE=str* parameters as described in the *RT-11 System Macro Library Manual*:

Macro Call: `.SPFUN area,chan,#SF.MTM,,,blk[,crtn][,BMODE=str][,CMODE=str]`

SF.MTM is the code 377 or the name SF.MTM if the program is assembled with the distributed file SYSTEM.MLB.

blk is the address of a 4-word error and status block used for returning the exception conditions. See Section 2.8.5.5.

This request returns the errors shown in Table 2–24. Additional qualifying information for these errors is returned in the first two words of the *blk* argument status block. See Section 2.8.5.5.

Table 2–24: SF.MTM (Code 377) Errors

Byte 52 Code	First Word Code	Qualifying Information
EOF (Value = 0)	1	Tape before EOF only (tape mark detected).
Hard error (Value = 1)	0	No additional information (consult documentation for your particular tape drive for all possible error conditions).
	1	Tape drive not available.
	2	The controller lost the tape position.
	4	Tape is write locked.

2.8.7 Hardware Magtape Handler

The hardware magtape handlers are identical to the distributed handlers except they are not built with the FSM. Therefore, the hardware magtape handlers accept only hardware requests. These are applicable in I/O operations where no file structure exists. Any file structure request you make to the hardware handler results in a monitor directory I/O error. The hardware handler is a subset of the file structure magtape handler. It can perform I/O operations on physical blocks, position the tape, and recover from errors.

Any file-structured request causes the hardware handler to issue a hard error. The hardware handler accepts only the non-file-structured .LOOKUP, .CLOSE, or special function requests.

If you do not need the file structure support, use the hardware handlers. You must perform a SYSGEN (see the *RT-11 System Generation Guide*) to get the hardware magtape handlers, then you must rename them in order to use them. Use a series of monitor commands similar to the following, which replace the file structure MS handler with the hardware MS handler.

1. Remove the distributed handler:

```
.REMOVE MS [RET]
```

2. Save the distributed handler:

```
.RENAME/SYS MS[X].SYS MS[X]FS.SYS [RET]
```

3. Replace the distributed handler with the hardware handler you built during SYSGEN:

```
.RENAME/SYS MS[X]HD.SYG MS[X].SYS [RET]
```

4. Install the hardware handler:

```
.INSTALL MS RET
```

2.8.8 Transporting Tapes to RT-11

RT-11 can read files written on other computer systems that support the ANSI standard labels. The following sections give a few examples of how to write ANSI tapes on some common Digital PDP-11 operating systems. Keep in mind that there are other factors involved in addition to the label and format compatibility, including density, parity, and number of tracks. Consult the appropriate system documentation for complete information on using magtapes under the different operating systems. (See the *RT-11 Volume and File Formats Manual* and the *RT-11 System Utilities Manual* for information on transporting tapes from RT-11 to other systems.)

2.8.8.1 From RSTS/E

RSTS/E supports two types of magtape format, DOS-11 and ANSI. In the following examples, *dd* represents the magtape handler name. To ensure that an ANSI file structure is written, issue the following commands:

Examples

1. ASSIGN ddn: .ANSI

Allocates the device to the job and ensures that an ANSI file structure is used.

2. RUN \$PIP
ddn:xxxxxxx/ZE

PIP initializes the tape; *xxxxxx* is the volume ID.

3. Really zero ddn:? YES

PIP prompts before initializing the tape.

4. PIP ddn:=TEST1.MAC,TEST2.MAC

PIP copies files to the tape.

5. DEASSIGN ddn:

Deallocates the device.

2.8.8.2 From RSX-11M

RSX-11M needs the following commands to access a magtape:

Examples

1. ALL ddn:

Allocates a drive.

2. INI ddn:RT11

Initializes the tape and gives the name *RT11* as the volume identification.

3. MOU ddn:RT11

Mounts the tape volume.

4. PIP ddn:=[13,14]TEST1.MAC,TEST2.MAC

Copies files to the tape.

5. DMO ddn:RT11

Dismounts the tape volume.

6. DEA ddn:

Deassigns the drive.

2.8.8.3 From RSX-11D and IAS

Use the following commands to write an ANSI tape on RSX-11D or IAS:

Examples

1. INI ddn:RT11

Initializes the tape and gives the name *RT11* as the volume identification.

2. MOU ddn:RT11

Mounts the tape volume.

For RSX-11D, use PIP to write files to the tape; for IAS, use the COPY command.

Examples

1. DMO ddn:RT11

Dismounts the tape volume.

The contents of files written under the RSX-11D, RSX-11M, and IAS systems do not necessarily correspond to those types of data files under RT-11. For example, under RT-11, text files consist of stream ASCII data (carriage return and line feed characters are embedded in the text); the other operating systems use a different type of character storage. Be sure to pay attention to the contents of the files you need to transfer.

When you write files to be read under RT-11, the only valid block size the utility programs use is 512 characters per block. However, the DIR program will list the directory of any ANSI compatible tape.

2.8.8.4 From VMS

Creating a magtape on a VAX processor running the VMS operating system for subsequent transfer to a PDP-11 running RT-11 is described in the *RT-11 Volume and File Formats Manual*. Look there for the procedure.

2.8.9 Seven-Track Magnetic Tape

Seven-track tapes contain six data tracks and one parity track, so a maximum of six data bits can be contained in one tape character. With seven-track tapes, the MT handler operates in either six-bit mode or core dump mode.

Six-bit mode is not compatible with the data normally created by PDP-11 systems; it is provided for transferring data to or from other systems. In addition, file structure operations cannot be performed in this mode. With the density set at 200 or 556 bpi, the magtape always operates in six-bit mode. When reading in six-bit mode, the handler places each six-bit tape character right-justified in a PDP-11 byte; the high-order two bits of the byte are set to 0. When writing in six-bit mode, the handler writes the low-order six bits of a PDP-11 byte as the six data bits of a tape character; the high-order two bits of the PDP-11 byte are not transferred or affected.

Core dump mode is compatible with PDP-11 systems. At 800 bpi, seven-track tape transfers can occur in either six-bit mode (SET MT: DENSE=807) or core dump mode (the default). Figure 2-7 illustrates the differences between six-bit mode and core dump mode.

In core dump mode, each PDP-11 byte is split into two tape characters. In writing to the tape, the handler writes the low-order four bits of a PDP-11 byte as the low-order four bits of the first tape character and the high-order four bits of the PDP-11 byte as the low-order four bits of the next tape character. The high-order two bits of each tape character are set to 0.

In reading from the tape, the reverse process occurs. The low-order four bits of the first tape character become the low-order four bits of the PDP-11 byte; the low-order four bits of the next tape character become the high-order four bits of the PDP-11 byte.

The high-order two bits of each tape character are not involved in the transfer, although they are included in the parity calculation. Thus, in core dump mode, the actual number of tape characters read or written is twice the number of PDP-11 bytes requested to be transferred; this conversion is performed by the magtape controller.

Figure 2–7: Seven-Track Tape

2.9 MU (TMSCP Magtape Handler)

This section provides specific programming information for TMSCP magtapes.

The MU handler supports magtape systems that use the tape mass storage communication protocol (TMSCP).

NOTE

The MU handler contains the same basic structure and provides the same support for programmed requests and special functions as described in Section 2.8 except as explicitly stated in this section. Therefore, this section describes only how the MU handler is different from the MM, MS, and MT handlers.

2.9.1 Support for Special Functions

The following special functions are either not supported by the reel-type magtape handlers or are supported in a different manner.

The SF.MTB and SF.BYP special functions are not affected by the presence (or absence) of the File Structure Module (FSM), as they are not concerned with operations on magtape volumes. Rather, they are concerned with data structures within the handler itself or the handler's controller.

Code	Name	Function
352	SF.MTB	Magtape data table access
	SF.TRD	<i>wcnt</i> argument for a read from the table; specified with a +1
	SF.TWR	<i>wcnt</i> argument for a write to the table; specified with a -1
360	SF.BYP	Direct TMSCP access; special function bypass
374	SF.MWE	Not Supported; writes with extended file gap executes as a write (SF.MWR) operation

2.9.1.1 TMSCP Translation Tables (SF.MTB), Code 352

Whenever an I/O request is passed to the MU handler, MU uses the RT-11 unit number as an index into the translation tables. MU then extracts the TMSCP unit number and port that have been assigned to that RT-11 unit, and uses the information to access the proper magtape drive.

You can read or write (modify) the memory-resident contents of the translation tables by using SF.MTB.

Size of the Translation Tables

The size of the translation tables is determined by the number of device units supported by DU. The distributed MU supports one unit; you can build an MU

that supports up to four units. You can determine the number of supported units for a particular handler by reading the MU.NUM field, as explained further.

Structure of the Translation Tables

As shown in Tables 2–25 and 2–26, the MU unit translation tables consist of a table header followed by table entries. The header starts at offset MU.ID, which is a word containing the Radix–50 value for the characters MU.

The MU.ID offset is followed by MU.NUM. The low byte of MU.NUM contains the number of entries in the table (and therefore the number of supported units). The high byte of MU.NUM is reserved.

The next offset is MU.ENT, which contains a pointer to the first table entry.

Table 2–25: TMSCP (MU) Translation Table Header

Offset	Name	Meaning
0	MU.ID	Radix–50 value for characters MU
2	MU.NUM	Byte containing number of entries in table
3		Reserved
4	MU.ENT	The offset of the first table entry

Each table entry is 4 bytes, and Digital recommends you use the symbol MU.ESZ to represent the 4-byte size of each entry.

Table 2–26: TMSCP (MU) Translation Table Entry

Offset	Name	Meaning
0	MU.UNI	Physical TMSCP unit number. The symbol MU\$Ux=nnnnnn is the initial value for the translation table when the handler is assembled. In the symbol, <i>x</i> is the octal RT–11 MU unit number (0–3) and <i>nnnnnn</i> is the TMSCP unit number. The SET MUx UNIT=nnnnnn command can subsequently change the value.
2	MU.JOB	Byte containing the number of the job connected to this TMSCP unit.
3	MU.POR	Byte containing the TMSCP port (controller) number. The symbol MU\$Ox=nnn is the initial value for the translation table when the handler is assembled. In the symbol, <i>x</i> is the octal RT–11 MU unit number (0–3) and <i>nnn</i> is the TMSCP port number. The SET MU PORT=nnn command can subsequently change the value.
4	MU.ESZ	Size of an entry (4 bytes)

Accessing the Translation Tables

Special function SF.MTB can read or write the TMSCP translation tables. Whether a read or write operation is performed is determined by the *wcnt* argument. Specify +1 (SF.TRD) for *wcnt* to read the tables; -1 (SF.TWR) to write the tables.

The translation tables are read from or written to a buffer, which is pointed to by the *buf* parameter.

2.9.1.2 Special Function Bypass (SF.BYP), Code 360

Special function SF.BYP bypasses all unit number translation and allows direct access to the TMSCP port. For MU, SF.BYP (direct TMSCP access) serves the same purpose as the DU handler's SF.BYP (direct MSCP access).

The request syntax and parameter argument definitions for SF.BYP are as follows:

Macro Call: `.SPFUN area,chan,#SF.BYP,buf,wcnt,blk`

<i>area</i>	is the address of a 6-word EMT argument block.
<i>chan</i>	is a channel number in the range 0 to 376 ₈ .
<i>SF.BYP</i>	is code 360 or the name SF.BYP if the program has been assembled with the distributed module SYSTEM.MLB.
<i>buf</i>	is the address of the 52 ₁₀ -word TMSCP area.
<i>wcnt</i>	when nonzero, is the virtual address of a data buffer to send to the handler. That virtual address is translated to a physical address and placed in the buffer of the TMSCP area. when zero, the buffer address in the TMSCP area is not altered.
<i>blk</i>	indicates whether the handler should perform retries: 1 = specifies retries 0 = specifies no retries

The buffer address in special function SF.BYP must point to a 52-word area in the user's job. The first 26 words are used to hold:

- A response packet length in bytes
- A virtual circuit identifier
- An end packet when the command is complete

The second 26 words are set up by the caller and contain:

- A length word (length of command)
- A virtual circuit identifier (must have octal 1 (001) in high byte)
- A valid TMSCP command (48₁₀-byte command buffer)

Except for port initialization, the user program must do all command packet sequencing, error handling, and reinitialization when the bypass operations are complete.

2.9.2 Unit Support, CSR and Vectors

The distributed MU handler supports one unit. Using the system generation procedure, you can build an MU handler that supports up to four units. Each unit requires a separate controller and you can only boot RT-11 from unit MU0, which must be installed at CSR address 774500 and vector address 260. The addresses for MU1 through MU3 float; they depend on what other devices are on the bus. The default CSR and vector addresses are as follows:

CSR	Vector
774500	260
774504	340
774510	344
774514	350

2.10 NL (Null Handler)

The null handler accepts all read and write requests. On output operations, this handler acts as a data sink. When a program calls NL, the handler returns immediately to the monitor indicating that the output is complete. The handler returns no errors and causes no interrupts. On input operations, NL returns an immediate EOF indication for all requests; no data is transferred. Hence, the contents of the input buffer are unchanged.

2.11 NC, NQ, NU (Ethernet Handlers)

RT-11 includes three Ethernet handlers that provide support for Ethernet class controllers. The NC Ethernet handler supports the DECNA controller for CTI Bus-based processors. The NQ Ethernet handler supports the DELQA and DEQNA Ethernet controllers for Q-bus processors. The NU Ethernet handler supports the DELUA and DEUNA controllers for UNIBUS processors.

Each handler supports only one controller and a maximum of eight units. These unit numbers are used as a logical connection between a user program and an address /protocol pair to be recognized by the Ethernet hardware.

2.11.1 Restrictions

Observe the following Ethernet handler restrictions:

- The handlers run only under mapped monitors.
- The handlers cannot be fetched and must be loaded.
- Programs that call the Ethernet handlers must be written to perform with the following elements in the order indicated:
 1. Use the .LOOKUP programmed request to open a channel to the device unit.
 2. Allocate the unit using .SPFUN 200.
 3. Perform the Ethernet operation or operations.
 4. Deallocate the unit using .SPFUN 200.
 5. Use the .CLOSE programmed request to close the channel to the specified device unit.

2.11.2 Support for Special Functions

The Ethernet handlers support the following special functions. The special function names are from the .NALDF macro in the distributed file SYSTEM.MLB.

Code	Name	Section	Function
200	SF.NAL	2.11.2.1	Allocate/Deallocate unit
201	SF.PRO		Reserved
202	SF.NPR	2.11.2.2	Enable/Disable protocol type
203	SF.NMU	2.11.2.3	Enable/Disable multicast address
204	SF.NWR	2.11.2.4	Transmit Ethernet frame
205	SF.NRD	2.11.2.5	Receive Ethernet frame

Successful completion of a .SPFUN request clears the carry bit. Completion with error sets the carry bit, and the status word in the buffer contains an error code.

2.11.2.1 Allocate/Deallocate Unit (SF.NAL), Code 200

The allocate unit special function allocates a unit of the Ethernet handler for a job's exclusive use.

The deallocate unit special function deallocates the unit so it can be used by another job.

2.11.2.1.1 Allocate Unit

The following is the form of the special function allocate unit:

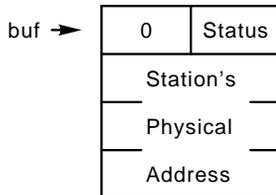
Macro Call: `.SPFUN area,chan,#SF.NAL,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.

chan is a channel number in the range 0 to 376₈.

SF.NAL is code 200 or the name SF.NAL if the program is assembled with the distributed file SYSTEM.MLB.

buf is the address of a 4-word buffer containing the status word and space for the station's physical address. The buffer contents are returned by the allocate unit special function.



The high byte of the status word contains a 0. Allocate unit returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success
2	Controller error while attempting to initialize the network interface (controller).
3	No resources (unit in use).
11	Reserved.

wcnt is #0.

blk is #1.

2.11.2.1.2 Deallocate Unit

The following is the form of the special function deallocate unit:

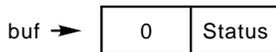
Macro Call: `.SPFUN area,chan,#SF.NAL,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.

chan is a channel number in the range 0 to 376₈.

SF.NAL is code 200 or the name SF.NAL if the program is assembled with the distributed file SYSTEM.MLB.

buf is the address of a 1-word buffer containing the status word.



The high byte of the status word contains a 0. Deallocate unit returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).
11	Unit still active.

wcnt is #0.

blk is #0.

2.11.2.2 Enable/Disable Protocol Type (SF.NPR), Code 202

The enable protocol type special function adds a protocol type to the list of those to be recognized by the unit. Only one protocol type can be specified for each unit. At least one protocol type must be enabled to receive Ethernet frames.

The disable protocol type special function removes the protocol type from the list of those recognized by the unit.

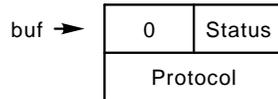
2.11.2.2.1 Enable Protocol Type

The following is the form of the special function enable protocol type:

Macro Call: `.SPFUN area,chan,#SF.NPR,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.

chan is a channel number in the range 0 to 376₈.
func is code 202 or the name SF.NPR if the program is assembled with the distributed file SYSTEM.MLB.
buf is the address of a 2-word buffer that contains the status word followed by the protocol type word.



The high byte of the status word contains a 0. Enable protocol type returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).
3	No resources (unit's protocol table is full).
6	Reserved.
10	Protocol type in use.

The protocol type is specified by the user.

wcnt is #0.
blk is #1.

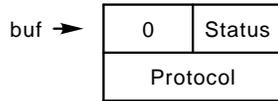
2.11.2.2.2 Disable Protocol Type

The following is the form of the special function disable protocol type:

Macro Call: `.SPFUN area,chan,#SF.NPR,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.
chan is a channel number in the range 0 to 376₈.
SF.NPR is code 202 or the name SF.NPR if the program is assembled with the distributed file SYSTEM.MLB.

buf is the address of a 2-word buffer that contains the status word, followed by the protocol type word.



The high byte of the status word contains a 0. Disable protocol returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).

wcnt is #0.

blk is #0.

2.11.2.3 Enable/Disable Multicast Address (SF.NMU), Code 203

The enable multicast address special function adds a multicast address to the list of those to be recognized by that unit. You need not specify the unit's physical or broadcast address. RT-11 supports only one multicast address per handler unit.

The disable multicast address special function removes a multicast address from the list of those to be recognized by the unit.

2.11.2.3.1 Enable Multicast Address

The following is the form of the special function enable multicast address:

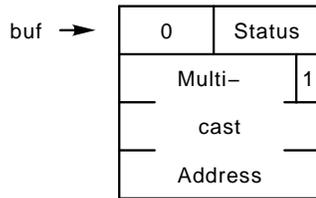
Macro Call: `.SPFUN area,chan,#SF.NMU,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.

chan is a channel number in the range 0 to 376₈.

func is code 203 or the name SF.NMU if the program is assembled with the distributed file SYSTEM.MLB.

buf is the address of a 4-word buffer that contains the status word, followed by the 3-word multicast address. The low-order bit of the first address word should be a 1.



The high byte of the status word contains a 0. Enable multicast address returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).
3	No resources (unit's address table is full, or hardware address table is full).

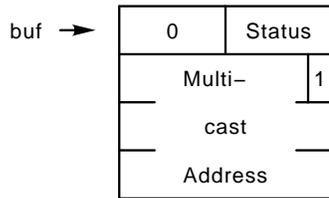
wcnt is #0.
blk is #1.

2.11.2.3.2 Disable Multicast Address

The following is the form of the special function disable multicast address:

Macro Call: `.SPFUN area,chan,#SF.NMU,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.
chan is a channel number in the range 0 to 376₈.
func is code 203 or the name SF.NMU if the program is assembled with the distributed file SYSTEM.MLB.
buf is the address of a 4-word buffer that contains the status word, followed by the 3-word multicast address. The low-order bit at the first address word should be a 1.



The high byte of the status word contains a 0. Disable multicast address returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).

wcnt is #0.
blk is #0.

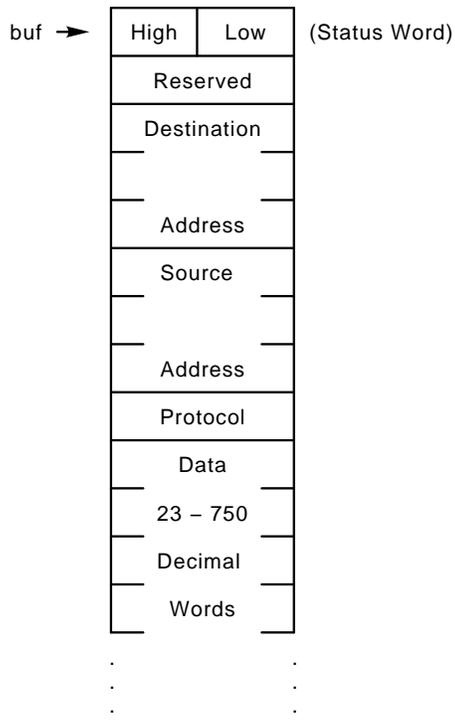
2.11.2.4 Transmit Ethernet Frame (SF.NWR), Code 204

The special function transmit Ethernet frame transmits the Ethernet frame pointed to in the *buf* parameter argument. If the source address field of the frame is nonzero, it is kept and used. If the source field of the frame is zero, the unit's physical address is inserted in the source field before transmission.

The following is the form of the special function transmit Ethernet frame:

Macro Call: `.SPFUN area,chan,#SF.NWR,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.
chan is a channel number in the range 0 to 376₈.
func is code 204 or the name SF.NWR if the program is assembled with the distributed file SYSTEM.MLB.
buf is the address of a variable-size buffer containing a word for returning status, a reserved word, and up to 757₁₀ words comprising the Ethernet frame to be transmitted.



Transmit Ethernet frame returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).
13	Transmit failed. When status code 13 is returned in the low byte of the status word, transmit Ethernet frame returns one of the following octal status subcodes in the high byte of the status word: 1 = Invalid frame length. 2 = Excessive collisions. 3 = Carrier check failed.

wcnt is determined by the variable size of the user buffer (including the status and reserved words). The packet size (including the status and reserved words) can vary between 32_{10} and 759_{10} words.

blk is #0.

2.11.2.5 Receive Ethernet Frame (SF.NRD), Code 205

The receive Ethernet frame special function returns the next Ethernet packet with the desired unit address and protocol type to the buffer. The function does not return Ethernet frames that are received with errors.

The following is the form of the special function receive Ethernet frame:

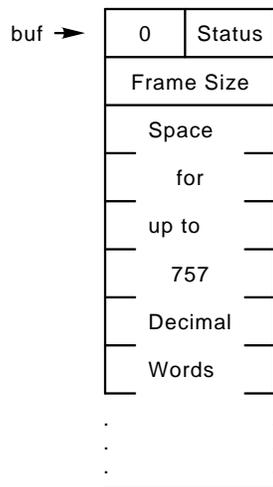
Macro Call: `.SPFUN area,chan,#SF.NRD,buf,wcnt,blk[,crtn][,BMODE=str][,CMODE=str]`

area is the address of a 6-word EMT argument block.

chan is a channel number in the range 0 to 376_8 .

func is code 205 or the name SF.NRD if the program is assembled with the distributed file SYSTEM.MLB.

buf is the address of a variable-size buffer containing a word for returned status, a word for returned frame size, and up to 757_{10} words to receive the Ethernet frame. The buffer contents are returned by the receive Ethernet frame special function.



The high byte of the status word contains a 0. The receive Ethernet frame special function returns one of the following octal status codes in the low byte of the status word:

Code	Meaning
0	Success.
1	Unknown unit. The specified unit was not opened by the job issuing the request.
2	Controller error while attempting to initialize the network interface (controller).

wcnt is the size of the user buffer including the status and frame size words. The maximum value allowed for the argument is 759₁₀; the minimum is 32₁₀.

blk is #0.

2.11.3 Example of Allocating an Ethernet Unit

The following example allocates a unit of the Ethernet handlers.

```

CONF2 = 370 ;Config word 2
          ; (RMON fixed offset)
PROS$ = 020000 ;RT is running on a PRO-3xx
BUS$ = 000100 ;Q-bus/UNIBUS processor
.
.
.
.GVAL #AREA,#CONF2 ;Get contents of Config word 2
MOV #<^RNC >,DBLK ;Assume PRO
BIT #PROS$,R0 ;Correct assumption?
BNE 10$ ;yes...
MOV #<^RNQ >,DBLK ;No, so assume Q-bus
BIT #BUS$,R0 ;Correct assumption?
BNE 10$ ;yes...
MOV #<^RNU >,R0 ;Nope, must be
          ; UNIBUS after all
10$: .GTJB #AREA,#JOB DAT ;Get info on this job
MOV JOB DAT,R0 ;R0 = job number (*2)
ASR R0 ;Convert to job number 0-7
ADD #<^R 0>,R0 ;Make it final RAD50 digit
ADD R0,DBLK ; and add it to
          ; the device name
.LOOKUP #AREA,#0,#DBLK ;Open a channel to Ethernet
.
. ;.LOOKUP error processing
.
.SPFUN #AREA,#0,#200,#BUFFER,#0,#1
          ;Allocate the unit to this job
.

```

```

.                                     ;.SPFUN error processing
.
AREA:  .BLKW  3
JOBDAT: .BLKW 12.
DBLK:  .WORD  0,0,0,0
BUFFER: .BLKW  4
.
.
.                                     ;END OF EXAMPLE
.
```

2.12 PI (CTI Bus-Based Processor Interface System Support Handler)

This section contains specific information about the PI system support handler and using RT-11 with CTI Bus-based processors. PI is called a *system support handler* because RT-11 requires PI to provide certain necessary connections with the computer hardware. At bootstrap time, the monitor loads PI before binding with the system device handler file on the system volume.

2.12.1 Support for Special Functions

The PI handler supports the following special functions which are used only with the GIDIS graphics package, as described in the *RT-11 System Subroutine Library Manual*:

Code	Name	Action
371	SF.PWR	Send command packet to GIDIS.
370	SF.PRD	Get status from GIDIS.

2.12.2 PI Keyboard Support

PI supports the keyboard in normal mode or function key mode.

2.12.2.1 Normal Mode

PI supports the following keys in normal mode:

- All keys on the main keypad.
- All keys on the numeric keypad.
- Cursor control (arrow) keys on the editing keypad.
- The following special function keypad keys: HOLD SCREEN (F1), PRINT SCREEN (F2), SETUP (F3), ESCAPE (F11), BACK SPACE (F12), and LINE FEED (F13).

PRINT SCREEN (F2) prints a copy of the text from your terminal screen directly on your printer. PRINT SCREEN cannot be used to print graphics. You must be running the transparent spooling package (SPOOL) under a mapped monitor to use PRINT SCREEN.

SETUP (F3) clears a locked keyboard and turns off the WAIT light when pressed. Note that the SETUP key has nothing to do with the setup utility.

The following keys do not function in normal mode:

- Special function keys F4 through F10, F14, HELP (F15), DO (F16), and F17 through F20.
- Editing keypad keys FIND, INSERT HERE, REMOVE, SELECT, PREV SCREEN, and NEXT SCREEN. Editing functions under RT-11 use the numeric keypad (see the *PDP-11 Keypad Editor User's Guide*.)

2.12.2.2 Function Key Mode (DECFKM)

Programs written for the PI handler can place the terminal in function key mode. In function key mode, each special function key sends an assigned control sequence to the processor. The control sequence is not assigned a specific function, but software can be programmed to recognize the control sequence.

A program places the terminal in function key mode by sending the 7-bit escape sequence:

```
ESC[?39h (transmitted as octal 033 133 077 063 071 150)
```

A program returns the terminal to normal key mode by sending the 7-bit escape sequence (note the lower-case l (?39l)):

```
ESC[?39l (transmitted as octal 033 133 077 063 071 154)
```

The following table lists control sequences for the special function keys:

Key	Control Sequence	Key	Control Sequence
F1	ESC[11~	DO (F16)	ESC[29~
F2	ESC[12~	F17	ESC[31~
F3	ESC[13~	F18	ESC[32~
F4	ESC[14~	F19	ESC[33~
F5	ESC[15~	F20	ESC[34~
F6	ESC[17~	COMPOSE CHARACTER	ESC[10~
F7	ESC[18~	FIND	ESC[1~
F8	ESC[19~	INSERT HERE	ESC[2~
F9	ESC[20~	REMOVE	ESC[3~
F10	ESC[21~	SELECT	ESC[4~
F11	ESC[23~	PREV SCREEN	ESC[5~
F12	ESC[24~	NEXT SCREEN	ESC[6~
F13	ESC[25~		
F14	ESC[26~		
HELP (F15)	ESC[28~		

2.12.3 Video Terminal Support

PI supports the CTI Bus-based processor's video terminal in the following manner:

2.12.3.1 Advanced Video Option Emulation

The PI handler supports a limited emulation of the VT100 implementation of the advanced video option, and uses the same escape sequences as the VT100 terminal. The limited emulation supports all VT100 character renditions (attributes) except BLINK; BLINK displays as BOLD. BOLD is not supported in 132-column mode, and 132-column mode is supported only by the mapped monitors.

2.12.3.2 Text Cursor Mode (DECTCEM)

Text cursor mode lets a program control whether the cursor is displayed on the video screen. Enabling text cursor mode displays the cursor and is the default. Text cursor mode is necessary when working with text because the cursor shows where the next character will be displayed.

A program places the terminal in text cursor mode by sending the 7-bit escape sequence:

```
[ESC][?25h (transmitted as octal 033 133 077 062 065 150)
```

A program takes the terminal out of text cursor mode by sending the 7-bit escape sequence (note the lower-case l (?25l)):

```
[ESC][?25l (transmitted as octal 033 133 077 062 065 154)
```

The cursor display can also be controlled using the SETUP CURSOR and SETUP NOCURSOR commands described in the *RT-11 Commands Manual*.

2.12.3.3 Device Attributes (DA)

A program uses the device attributes request/reply exchange to ask the terminal, "what are you?". The response sent by the terminal to the program can identify the terminal as a specific VT100 terminal (the default) or as a nonspecific member of the VT100 series of terminals. The SETUP modes VT100 and GENERIC100 (see the *RT-11 Commands Manual*) determine which of the two responses the terminal sends the program. Digital recommends that all programs recognize both the VT100 and the GENERIC100 device attributes reply.

A program can request information on two levels. The primary level DA requests basic compatibility information. The secondary level DA requests the specific version and edit level of the PI handler.

The terminal reply to primary and secondary DA requests gives this information, and also tells the program which monitor the system is running. The following is a complete DA interchange:

A program requests primary DA by sending the 7-bit escape sequence:

```
[ESC][c (transmitted as octal 033 133 143)
```

- If the terminal is SETUP VT100, it responds by sending the 7-bit escape sequence:
 - When running under an unmapped monitor:


```
ESC[?1;1c (transmitted as octal 033 133 077 061 073 061 143)
```
 - When running under a mapped monitor:


```
ESC[?1;3c (transmitted as octal 033 133 077 061 073 063 143)
```
- If the terminal is SETUP GENERIC100 without 132-column capability (running under an unmapped monitor), it responds by sending the 7-bit escape sequence:


```
ESC[?61c (transmitted as octal 033 133 077 066 061 143)
```
- If the terminal is SETUP GENERIC100 with 132-column capability (running under a mapped monitor), it responds by sending the 7-bit escape sequence:


```
ESC[?61;1c (transmitted as octal 033 133 077 066 061 073 061 143)
```

A program requests the secondary DA by sending the 7-bit escape sequence:

```
ESC[>c (transmitted as octal 033 133 076 143)
```

- If the terminal is operating under an unmapped monitor, it responds by sending the 7-bit escape sequence:


```
ESC[>7;VVnnc (transmitted as octal 033 133 076 067 073 V V n n 143)
```

where *VV* is the version number, and *nn* is the edit level of the PI handler.
- If the terminal is operating under a mapped monitor, it responds by sending the 7-bit escape sequence:


```
ESC[>8;Vnnc (transmitted as octal 033 133 076 070 073 V V n n 143)
```

where *VV* is the version number, and *nn* is the edit level of the PI handler.

2.13 UB (UNIBUS Mapping Register (UMR) System Support Handler)

This section describes the UB handler that provides support for the UNIBUS mapping registers on UNIBUS processors. The UB handler provides DMA (direct memory access) support for 22-bit memory addressing during I/O operations.

UB is called a *system support handler* because RT-11 requires UB to provide certain necessary connections with the computer hardware. At bootstrap time, the monitor loads UB before binding with the system device handler file on the system volume. Therefore, UB cannot be installed with the INSTALL command. Instead, UB is automatically installed and loaded in memory on UNIBUS processors with the following configuration:

- The processor is running a mapped monitor.
- The processor contains more than 256K-bytes of memory.
- The processor contains UNIBUS Mapping Registers at addresses 170200 through 170400 to support 40₈ 2-word UMRs.
- At least one device handler on the system uses DMA in performing I/O operations. All distributed RT-11 handlers that can perform DMA are so marked.

Section 2.13.3 describes how to provide UMR support in a user-written DMA handler.

- All installed user-written (not distributed) device handlers are compatible with RT-11 support for UB. All installed device handlers must be marked as compatible with UB, whether or not they perform DMA operations.

Section 2.13.2 describes how to make a non-DMA user-written device handler compatible with RT-11 UB support.

UNIBUS Mapping Registers function in a manner that is similar to the Memory Management Unit (MMU) registers that provide 22-bit address translation for the CPU. The UMRs provide address translation (mapping) from the 18-bit UNIBUS to the 22-bit memory bus.

2.13.1 UMR Support with Distributed Handlers

On supported UNIBUS system configurations, UB is automatically installed and loaded when the processor is booted. At that point, DMA I/O operations are handled transparently by the processor UMR hardware and the RT-11 operating system. Programs that use distributed RT-11 device handlers require no modification to support DMA access to a peripheral device.

The aspects of UMR support that apply to distributed handlers are:

- Permanent UMR allocation.

Because of internal buffers, some RT-11 device handlers, such as DL, DM, DU, NU, and the various magtape handlers, require a preallocation of one or more permanent UMRs. RT-11 preallocates those permanent UMRs when the device handlers are installed at system boot. RT-11 reserves those permanent UMRs

for those handlers when they are loaded. See Table 2–27. You can regain any preallocated permanent UMRs for handlers that install but you are not using, by renaming the device handler. Such a renamed handler does not install at the next system boot.

Contiguous permanent UMRs are allocated from the list of reserved permanent UMRs when handlers are loaded and returned to reserved status when handlers are unloaded. After numerous load/unload operations, the list of reserved permanent UMRs can become fragmented. A symptom of this condition is the inability to load a device handler that requires multiple permanent UMRs even when sufficient reserved permanent UMRs exist. Two courses of action are available if that condition occurs. You can reboot your system, or you can issue the SHOW UMR command and unload the device handlers that are displayed as occupying slots between the available reserved permanent UMRs. The system device handler resides at the top of the list. You should consolidate the list from the base upward.

- Temporary UMR allocation.

Many distributed device handlers require one or more UMRs on a temporary basis to process I/O requests. RT–11 allocates temporary UMRs as the need occurs. Each processor contains 31₁₀ accessible UMRs, and the allocation of UMRs can be displayed by the command SHOW UMR.

- Serialization of I/O request satisfaction.

When UB is loaded in memory, RT–11 no longer always satisfies I/O requests in serial order.

Of the distributed RT–11 device handlers, only DU and the magtape handlers (MM, MS, MT, and MU) require that I/O requests are satisfied in serial order. The guarantee of I/O request serialization is internal to those handlers and requires no user intervention.

However, RT–11 does not guarantee that I/O requests for other device handlers are satisfied in serial order. Rather, I/O requests are satisfied in the quickest manner possible, which might or might not be serial. For example, an I/O request that requires four UMRs might be queued for a time waiting for UMR allocation, while a subsequent I/O request requiring fewer UMRs is satisfied. However, if required, you can force serialized I/O request processing, using the SET UB SERIAL=n command, described in the *RT–11 Commands Manual*.

You can control other aspects of UMR support by specifying conditions for the SET UB command. Other than those conditions, UMR support is totally transparent when using the distributed RT–11 device handlers.

Table 2–27: Distributed Handler Support for UMRs

Device Handler	DMA=	PERMUMR=
DL	YES	1
DM	YES	1
DU	YES	2
DW	NO	
DY	YES	1
MM	YES	If support for FSM included, requires 1 if no support for FSM, requires 0
MS	YES	If support for FSM included, requires 1 if no support for FSM, still requires 1
MT	YES	If support for FSM included, requires 1 if no support for FSM, requires 0
MU	YES	If support for FSM included, requires 3 if no support for FSM, requires 2
NU	YES	3
RK	YES	0
VM	NO	

2.13.2 Including Required UB Support in User-Written Non-DMA Handlers

All installed device handlers, including those that perform no DMA operations, must be modified for compatibility with UB. Otherwise, the RT–11 monitor bootstrap does not load UB and the system then operates with only the low 256K words of memory accessible to DMA operations.

You must explicitly specify whether each user-written device handler supports DMA, using the `.DRDEF` macro's `DMA=str` parameter. If a device handler does not perform DMA operations and, therefore, does not require UMR allocation, specify `DMA=NO`.

2.13.3 Including UMR Support in User-Written DMA Handlers

UMR support is appropriate for a device handler that performs I/O operations and is capable of DMA. Including UMR support in such a device handler lets the handler access computer memory beyond the 18-bit 256K-byte boundary during I/O operations.

The following paragraphs describe elements of the new UMR support that must be considered before you include UMR support in a device handler. Each element is either described when listed or you are pointed to the appropriate section of this manual where you will find the element description.

Including UMR support in any device handler requires that you understand the following items:

- The handler should not perform DMA operations from within its own install code. If a handler must be written to perform DMA from within its install code, you must turn off UB (SET UB NOINSTALL), reboot the system, and then install the handler.
- The handler must use the .DRDEF macro and include one or more of the parameters, *DMA=str*, *PERMUMR=n*, and *SERIAL=str*, as described in the *RT-11 System Macro Library Manual*.
- If the handler uses the .QELDF macro to define queue elements, you should read about the offset, Q.MEM, as described in *RT-11 System Macro Library Manual*.
- RMON automatically allocates temporary UMRs for all .READx and .WRITx requests to handlers that are marked as DMA=YES. RMON also automatically releases all such temporary UMRs. Both operations are completely transparent to the handler.
- If the handler previously used queue element offsets Q.PAR and Q.BUFF to calculate non-DMA I/O virtual addresses, it must now use the new offset Q.MEM in conjunction with Q.BUFF. Q.MEM is described in Section 1.2.1.1.2. The handler now uses Q.PAR to calculate only DMA I/O virtual addresses.
- If the handler previously used extended memory subroutines \$GETBYT, \$PUTBYT, \$PUTWRD, or \$MPPHY, read the paragraphs *Changes to extended memory subroutines for UMR support*, in *RT-11 System Release Notes*.
- You should examine the new RMON fixed offsets, \$QHOOK, \$H2UB, and the bits defined for UB in \$CNFG3. They are described in *RT-11 System Internals Manual*.
- You should decide if I/O requests for the handler or the job must be satisfied in serial order. Once UB is loaded in memory, I/O requests are not guaranteed to be satisfied in serial order by default.

If the handler requires serialized I/O request satisfaction, you must specify the .DRBEG macro *SERIAL=YES* parameter argument when you build the handler. See the .DRDEF macro information in the *RT-11 System Macro Library Manual*.

If the job requires serialized I/O request satisfaction, see the SET UB SERIAL=n command described in the *RT-11 Commands Manual*.

- The device handler must use permanent or temporary UMRs for each special function that performs a DMA I/O operation.

The handler uses permanent UMRs for processing special functions that result in a DMA I/O operation to the handler internal buffer.

If the handler contains internal buffers that store command packets and responses, the handler has to use the ALLUMR routine to explicitly obtain at least one permanent UMR. The handler must explicitly release all permanent UMRs when it unloads, using the RLSUMR routine. Obtaining and releasing permanent UMRs is described in Sections 2.13.3.3 and 2.13.3.4.

The handler allocates at least one temporary UMR for each special function that performs a DMA I/O operation to the user buffer. The temporary UMRs are allocated either implicitly or explicitly.

Special functions (.SPFUNs) used by the handler are categorized as standard or nonstandard. A standard special function uses the .SPFUN *buf* parameter as the read/write buffer address and the *wcnt* parameter as the operation word count. Temporary UMRs for standard special functions are allocated implicitly. Defining standard special functions is described in Section 2.13.3.1.

A nonstandard special function does not use *buf* as the read/write buffer address or *wcnt* as the operation word count. The handler must explicitly obtain temporary UMRs for nonstandard special functions, requiring additional processing by UB. Processing nonstandard special functions is described in Section 2.13.3.3.

2.13.3.1 Defining Special Functions for Implicit UMR Allocation

The device handler should implicitly allocate UMRs for special functions that do the following:

- Perform DMA operations.
- Use the *buf* and *wcnt* parameters in the documented manner; are standard special functions.

The handler supports implicit UMR allocation for standard special functions by using the .DRBEG *SPFUN=spsym* parameter and a list of those functions. The *spsym* argument is the label of the list of those functions. The list is structured in the same manner as that used for the .DRSPF *extension table* method. However, unlike the .DRSPF macro, no pointer to the list resides in block 0 of the handler and the concept of special function *type* has no meaning and is not included.

The list of standard special functions must continuously reside in the low-memory portion of the handler whenever the handler is loaded. For all special functions in the list, RMON performs the UMR allocation and the address translation.

Defining special functions for implicit UMR allocation is illustrated in the example program in this section.

2.13.3.2 Explicitly Allocating Permanent UMRs (ALLUMR)

If the device handler contains internal buffers that store command packets and responses, you must allocate at least one permanent UMR to the device handler.

RT-11 allows up to 22₁₀ UMRs to be permanently allocated to handlers and one UMR is permanently allocated to the I/O page. When the system is booted, RT-11 allocates the one UMR to the system's I/O page and then reserves permanent UMRs for requesting device handlers as each handler is installed. Therefore, unless the 23₁₀ limit is reached, RT-11 reserves sufficient permanent UMRs to support all installed device handlers that request permanent UMR allocation. However, reserved permanent UMRs are not allocated to a device handler until it is loaded. Unallocated reserved permanent UMRs are available for explicit allocation, using

the ALLUMR routine. You can determine the current UMR allocation on your system by issuing the SHOW UMR command.

The ALLUMR routine, which resides in UB, is called to permanently allocate UMRs. If the handler requires UMRs for a single, contiguous chunk of memory, you need call ALLUMR only once. If the handler requires UMRs for noncontiguous chunks of memory, repeatedly call ALLUMR to allocate UMRs for each chunk.

You reference the UB entry vector through the \$H2UB fixed offset (460) in RMON. The ALLUMR routine is offset 1 word (\$H2UB+2) from the address pointed to by \$H2UB.

Use the following procedure to allocate permanent UMRs:

1. Calculate the number of permanent UMRs you need for each contiguous chunk of memory. One permanent UMR is required for each 4096 words of contiguous internal buffer space.
2. Specify the total number of permanent UMRs the handler requires in the *PERUMR=n* parameter of the .DRDEF macro in your handler source code. The RT-11 monitor bootstrap (BSTRAP) uses that information to reserve the number of UMRs you permanently allocate to the handler.
3. Before calling ALLUMR to allocate permanent UMRs for an internal buffer space, set up the following registers:

Register	Contents
R0	Number of permanent UMRs to be allocated for this contiguous chunk of internal buffer space. If you request more than one permanent UMR, the address of the first is defined by R1 and R2, and each subsequent UMR is offset by a value of 20000 ₈ .
R1	Bits 0–15 of the 22-bit physical memory base address (word aligned) of the internal buffer.
R2	Bits 16–21 of the 22-bit physical memory base address of the internal buffer.
R4	The address of a 1-word location in low memory that contains two RAD50 identifying characters. The SHOW UMR command displays these characters to identify this permanent UMR allocation. (In distributed handlers, is the device handler name.) The monitor must have continuous access to the specified memory location. If ALLUMR is called more than once for this handler, R4 in subsequent calls must contain a different address in low memory for each call. The 1-word location contents can be, but do not need to be, the same two RAD50 characters.

The contents of R3 and R5 are not defined or preserved across the call.

4. Within the device handler FETCH/LOAD code, call the ALLUMR routine. On return from ALLUMR:

If the carry bit is clear:

- R1 contains bits 0–15 of the 18-bit UNIBUS virtual address of the internal buffer.
- R2 contains bits 16 and 17 of the 18-bit UNIBUS virtual address of the internal buffer.
- The handler uses the address returned by ALLUMR (or some offset from that address) to program the device for DMA I/O to/from the handler internal buffer.

If the carry bit is set, insufficient UMRs are available for allocation and the handler must fail its load code.

Once you have successfully called and returned from ALLUMR, your handler code should confirm that the FETCH/LOAD succeeded. If the fetch/load operation fails after successfully returning from ALLUMR, you must call RLSUMR to free the allocated UMRs.

2.13.3.3 Explicitly Obtaining Temporary UMRs (GETUMR)

Device handlers that support nonstandard .SPFUN I/O DMA operations to or from a user buffer must call GETUMR to explicitly obtain temporary UMRs to service those requests. The temporary UMRs are automatically released after the request is serviced. The handler uses the GETUMR routine, described in this section, to obtain the UMRs. Be sure to call GETUMR before removing the queue element from the handler's current queue element (xxCQE) list.

The handler supports explicit UMR allocation for nonstandard special functions by using the `.DRBEG NSPFUN=nspsym` parameter and a list of those functions. The *nspsym* argument is a unique symbol name that is the same as the label at the list of those functions. The list is structured in the same manner as that used for the `.DRSPF extension table` method. However, unlike the `.DRSPF` macro, no pointer to the list resides in block 0 of the handler and the concept of special function *type* has no meaning and is not included.

The list of nonstandard special functions must continuously reside in the low-memory portion of the handler whenever the handler is loaded. Also, the handler must call GETUMR (with a word count of zero) even when a listed nonstandard special function performs no I/O and no UMRs are needed.

Defining special functions for explicit UMR allocation is illustrated in the example program in this section.

The handler calls the GETUMR routine, which resides in UB, to obtain temporary UMRs. You reference the UB entry vector through the \$H2UB fixed offset (460) in RMON. The GETUMR routine is located at the address pointed to by \$H2UB (offset 0).

Use the following procedure to explicitly obtain temporary UMRs:

1. Before calling GETUMR, set up the following registers:

Register	Contents
R0	Number of words to be transferred; the word count. If no DMA I/O is to be performed by this request, R0=0.
R1	Contents determined by R3: R3 = 0 R1 contains the Q.PAR value that is calculated by the handler. RMON cannot calculate the Q.PAR value because the special function's <i>buf</i> parameter contains a nonstandard argument. R3 = 1 R1 contains bits 0-15 of the 22-bit physical memory base address (word aligned).
R2	Contents determined by R3: R3 = 0 R2 is unused. R3 = 1 R2 contains bits 16-21 of the 22-bit physical memory base address.
R3	Contents indicate the type of address being specified: R3 = 0 Address is PAR value, specified in R1. R2 is not used. R3 = 1 Address is 22-bit physical address, specified in R1 and R2.
R4	Queue element offset Q.BLKN.

The contents of all unused registers are not defined or preserved across the call.

2. Within the device handler code that processes nonstandard special functions, call the GETUMR routine. On return from GETUMR:
 - If the carry bit is clear, the contents on return for R1 and R2 are defined by the contents of R3 when GETUMR was called. If GETUMR is called with R3 = 0, on return, R1 contains the new Q.PAR equivalent value and R2 is not defined. If GETUMR is called with R3 = 1, on return, R1 contains bits 0–15 and R2 contains bits 16 and 17 of the 18-bit UNIBUS virtual address.
 - If the carry bit is set, UB is unable to immediately allocate the requested UMRs for the queue element and the handler should simply return to the monitor.

2.13.3.4 Explicitly Releasing Permanent UMRs (RLSUMR)

All permanent UMRs that are allocated by a handler must be explicitly released by the handler when the handler is unloaded. A corresponding RLSUMR routine must be called for each ALLUMR routine that was called.

The RLSUMR routine, which resides in UB, releases permanent UMRs. You reference the UB entry vector through the \$H2UB fixed offset (460) in RMON. The RLSUMR routine is offset 2 words (\$H2UB+4) from the address pointed to by \$H2UB.

Use the following procedure to explicitly release permanent UMRs:

1. Before calling RLSUMR, set up the following register:

Register	Contents
R1	The address of the 2-character RAD50 device handler name specified in R4 of the corresponding ALLUMR routine. (The contents of RLSUMR R1 match the contents of corresponding ALLUMR R4.)

The contents of R0 and R2–R5 are not defined or preserved across the call.

2. Within the device handler RELEASE/UNLOAD code, call the RLSUMR routine.

On return from RLSUMR, all UMRs that were permanently allocated to the handler by the corresponding ALLUMR routine are released.

2.13.4 Example (Skeletal) Handler

The following example skeletal handler illustrates the macros and routines required to support UMRs.

```
.SBTTL  CONDITIONAL ASSEMBLY SUMMARY
;+
;COND
;
;      MMG$T = 1           Std conditional (XM only)
;      TIM$T           Std conditional (no code effects)
;      ERL$G           Std conditional (no code effects)
;-

.MACRO  ...
.ENDM

.MCALL .DRDEF .ASSUME .ADDR .DRSPF
.LIBRARY "SRC:SYSTEM"
.MCALL .SYCDF .FIXDF .HANDF .UBVDF .PIXDF

.SYCDF
.FIXDF
.HANDF
.UBVDF
.PIXDF

; UB Definitions
;  XB internal DMA buffer equates

      BUFSIZ  =: 20000           ; Size of XB internal DMA buffer
      NOUMRS  =: <BUFSIZ+7777/10000> ; Number of permanent UMRs required

; Special function definitions
; All special functions are DMA except for FN$SIZ and FN$MPM.
; FN$WRT AND FN$RED go in UBTAB. FN$REP uses a permanent UMR.
; FM$NSP is nonstandard so it goes in UBNTAB.
```

```

FN$MPM =: 370 ; Illustrate use of $MPMEM (not DMA)
FN$NSP =: 371 ; Nonstandard SPFUN (DMA to
; user buffer)
FN$SIZ =: 373 ; Get device size (not DMA)
FN$REP =: 374 ; Force reread of replacement table
FN$WRT =: 376 ; Absolute write (no bad block)
FN$RED =: 377 ; Absolute read (replacement)

.DRSPF <FN$MPM> ; Illustrate use of $MPMEM
.DRSPF <FN$NSP> ; Nonstandard SPFUN (DMA to
; user buffer)
.DRSPF <FN$SIZ> ; Get device size
.DRSPF <FN$REP> ; Force reread of replacement table
.DRSPF <FN$WRT> ; Absolute write (no bad block)
.DRSPF <FN$RED> ; Absolute read (replacement)

; DRDEF'S serial argument must be set equal to yes since XB calls
; GETUMR and depends on receiving queue elements from RMON in serial order.
; Calls to GETUMR can interfere with the serial ordering of queue elements
; unless "SERIAL = YES" is specified here.

.DRDEF XB,0,SPFUN$,0,0,0,DMA=YES,PERMUMR=NOMURS,SERIAL=YES
.DRPTR FETCH=FETCH,LOAD=FETCH,RELEASE=RELEAS,UNLOAD=RELEAS
.DREST CLASS=DVC.NL

; Start of handler

.DRBEG XB,SPFUN=UBTAB,NSPFUN=UBNTAB
XBBASE=XBSTRT+6
BR BEGIN ; Branch around data area

; Data area
$ENTPT: .WORD 0 ; Pointer to $ENTRY table
$PNMPT: .WORD 0 ; Pointer to $PNAME table
H2UB: .WORD 0 ; Pointer to UBVECT
XBSLOT: .WORD 0 ; XB'S offset in device tables
XBENT: .WORD 0 ; XB'S $ENTRY table entry pointer
XBPNA: .WORD 0 ; XB'S $PNAME table entry pointer

;+
; Definition of the handler internal buffer and the words that are
; used to program DMA devices that transfer data to and from it.
;-

XBDBUF: .WORD BUFSIZE ; XB DMA buffer - it is
; mapped by permanent UMRS
BUFADH: .WORD 0 ; Bits 0-15 of UNIBUS virtual
; Pointer to XBDBUF
BUFADL: .WORD 0 ; Bits 16-18 of UNIBUS virtual
; Pointer to XBDBUF

; Table of standard DMA SPFUNs that do DMA transfers to areas of
; memory not mapped by XB's permanent UMRS. UB will intercept these requests
; and assign temporary UMRS to them in the same manner as for .READx and
; .WRITx requests.
UBTAB: .DRSPF -,<FN$WRT> ; Absolute write, no bad block
.DRSPF -,<FN$RED> ; Absolute read (replacement)
.WORD 0 ; Table terminator

; Table of nonstandard DMA SPFUNs that do DMA transfers to areas of
; memory not mapped by XB's permanent UMRS. XB MUST explicitly allocate
; UMRS for the nonstandard SPFUNs listed here by calling UB's GETUMR
; routine. If no DMA transfer will take place (because of error, for
; example) XB should call GETUMR with a word count of 0. IF XB processes
; a nonstandard DMA SPFUN listed in UBNTAB without calling GETUMR,
; the job's I/O stream will hang.
UBNTAB: .DRSPF -,<FN$NSP> ; DMA to user buffer
.WORD 0 ; Table terminator

```

```

BEGIN:  MOV    XBCQE,R4                ; Point to current queue element
        MOVB   Q$FUNC(R4),R2          ; Get function code / unit number
        CMPB   R2,#FN$MPPM           ; Dispatch to function routine
        BEQ    FNMPM
        CMPB   R2,#FN$NSP
        BEQ    FNNSP
        CMPB   R2,#FN$SIZ
        BEQ    FNSIZ
        CMPB   R2,#FN$REP
        BEQ    FNREP
        CMPB   R2,#FN$WRT
        BEQ    FNWRT
        CMPB   R2,#FN$RED
        BEQ    FNRED
        TST    R2                    ; Normal request?
        BNE    XBEXIT                ; No, unknown SPFUN
        BR     XBRDWR                ; Yes, process read,write
        ...

;      Routines to perform SPFUN operations
;      at entry, R4 -> queue element

FNNSP:  MOV    #4000,R0                ; R0 = word count
        MOV    Q$PAR(R4),R1          ; Get address from QEL
        MOV    @$$SYPTR,R3          ; Get start of RMON
        MOV    $H2UB(R3),R5         ; R5 = UB entry vector
        CLR    R3                    ; Address type is PAR value
        CALL   UB.GET(R5)           ; Try to get UMRS
        ; (Note that at time of call, the
        ; Queue element must be on xxCQE)
        BCS    RETURN              ; Unable to get UMRs-do simple RETURN
        ...                        ; Got UMRs, initiate transfer
        BR     XBEXIT              ; DRFIN because this is an example
        ; Handler and there are really no
        ; Interrupts associated with it.
        ; If there were, the DRFIN would be
        ; Issued at interrupt time when
        ; The DMA transfer is finished.
        ; This is true for the other SPFUN
        ; Routines below, as well.

FNMPM:  ;
;      This routine illustrates how to call $MPMEM. $MPMEM is used
;      to map KT-11 virtual addresses (as described by Q.MEM and Q.BUFF
;      offsets in the queue element) to 18 or 22-bit physical addresses.
;      $MPMEM must be used for this purpose instead of $MPPHY when the
;      handler has DMA = YES. (When DMA = NO, the handler may use
;      either $MPMEM or $MPPHY.)
;
;      At entry:      R4 -> Q.BLKN offset in queue element
        MOV    @$$SYPTR,R3          ; Get start of RMON
        MOV    P1$EXT(R3),R3        ; R3 -> $P1EXT
        MOV    R4,R5                ; Make R5 -> 5TH word (Q.BUFF) of
        CMP    (R5)+,(R5)+         ; Queue element
        CALL   $MPMEM(R3)          ; Map KT-11 virtual to physical
        MOV    (SP)+,R2             ; R2 = low 16 bits physical address
        MOV    (SP)+,R3            ; R3 = HIGH 2 (OR 6) bits physical
        ; address
        ...                        ; Fall through to DRFIN

FNSIZ:
FNREP:
FNWRT:
FNRED:
XBRDWR:

XBEXIT: .DRFIN  XB                ; Return to monitor, done with
; queue element

RETURN: RETURN                    ; Return to monitor, not done with
; queue element

XBINT:
...

.DREND  XB

```

```

.SBTTL  FETCH/LOAD CODE
;+
;      FETCH
;
;      ENTRY:  R0 = Starting address of this handler service routine.
;              R1 = Address of GETVEC routine.
;              R2 = Value $SLOT*2. (length of the $PNAME table in bytes.)
;              R3 = Type of entry.
;              R4 = Address of SY read routine.
;              R5 -> $ENTRY slot for this handler.
;
;
;-

FETCH:  MOV     R5,R1                ; Save PTR to XB'S $ENTRY slot
        MOV     @R1,R0              ; Get address of XBLQE
        MOV     @$$SYPTR,R4         ; Get start of RMON
        MOV     $H2UB(R4),R3        ; R3 = UBVECT pointer
        MOV     R3,<H2UB-XBBASE>(R0) ; H2UB = address of UBVECT
        MOV     $PNPTR(R4),R3       ; R3 = RMON offset to PNAME table
        ADD     R4,R3               ; R3 -> PNAME table address
        MOV     R3,<$PNMPT-XBBASE>(R0) ; $PNMPT -> PNAME table address
        ADD     R2,R3               ; R3 -> $ENTRY table
        MOV     R3,<$ENTPT-XBBASE>(R0) ; $ENTPT -> $ENTRY table
        MOV     R5,<XBENT-XBBASE>(R0) ; XBENT -> XB'S $ENTRY table entry
        SUB     R2,R5               ; R5 -> XB'S $PNAME table entry
        MOV     R5,<XBPNA-XBBASE>(R0) ; XBPNA -> XB'S $PNAME table entry

;+
;      Allocate permanent UMRs to point into XB's internal DMA buffers,
;      XBDBUF and XBFILL, and get the UNIBUS virtual address.
;-

        MOV     #<XBDBUF-XBBASE>,R1 ; R1 = LOW 16 bits of DMABUF address
        ADD     R0,R1                ;
        CLR     R2                    ; R2 = HIGH 6 Bits of DMABUF address
        MOV     <XBPNA-XBBASE>(R0),R4 ; R4 -> PNAME entry for XB
        MOV     <H2UB-XBBASE>(R0),R5 ; Get UB entry address
        MOV     R0,-(SP)              ; Save XB starting address
        MOV     #NOUMRS,R0           ; R0 = number of UMRs required
        CALL    UB.ALL(R5)           ; Call ALLUMR
        MOV     (SP)+,R0             ; Restore XB starting address
        BCS     30$                  ; Couldn't get UMR, fail the load
        MOV     R1,<BUFADL-XBBASE>(R0) ; Store UNIBUS virtual address low
        MOV     R2,<BUFADH-XBBASE>(R0) ; Store UNIBUS virtual address high
        CLC                               ; Load succeeded
30$:    RETURN

;+
;      RELEAS
;
;      Routine to unload XB
;
;      Entry:  same as for load.
;
;-

.ENABL  LSB
RELEAS:
        MOV     R5,R1                ; R1 = $ENTRY slot for DM
        SUB     R2,R1                ; R2 -> $PNAME SLOT for DM
        MOV     @$$SYPTR,R4         ; Get start of RMON
        MOV     $H2UB(R4),R5        ; R5 = UB entry vector
        CALL    UB.RLS(R5)          ; Release UMRs
        RETURN                       ; And exit

.END

```

2.14 VM (Virtual Memory Handler)

This section contains specific programming information for the VM device. The *Introduction to RT-11* contains complete information on using the VM device. You should read the VM chapter in the *Introduction to RT-11* first.

The VM handler installation code determines the size of memory when the handler is installed. After determining the size of memory, the handler installation code reserves all extended memory above the handler's base address. The handler does not need to perform this operation each time it is loaded, thereby speeding the handler load process.

If you do not want to use VM and do not want VM to reserve memory for its own use, you have several options. You can remove the VM handler from your system disk so that it will not be installed when you bootstrap your system. You can set the base address above the high limit of available memory, which will prevent handler installation. Or, you can put a command in your startup command file to remove the VM handler from your system after the bootstrap has installed it. Otherwise, the VM handler installation code will always reserve extended memory for its own use, thereby making it unavailable to your program.

The base address (n) used in the SET VM BASE= n command is the desired base address in octal, divided by 100_8 . For example, the value 1600 sets the base address at the 28K-word address boundary, or 10000 sets the base address at the 128K-word address boundary; any other value between 1600 and the physical memory high limit is also acceptable. Lowering the value at which you set the VM base increases the region size. The table below gives a list of some K-word memory sizes and corresponding values for n .

K-words	N
28	1600
32	2000
64	4000
96	6000
128	10000
256	20000
512	40000
1024	100000

Figure 2-8 shows a 22-bit system with a VM base address of 10000 (128K words).

If you are using a mapped monitor and your hardware does not have 22-bit addressing, the default VM handler will not install; you will have to change the base address to a lower value before using VM with your mapped system. You can

Figure 2–8: VM Handler in a 22-Bit System

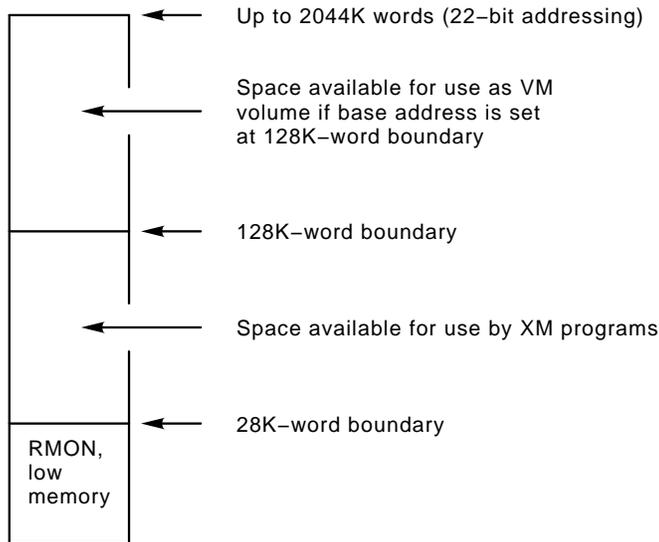
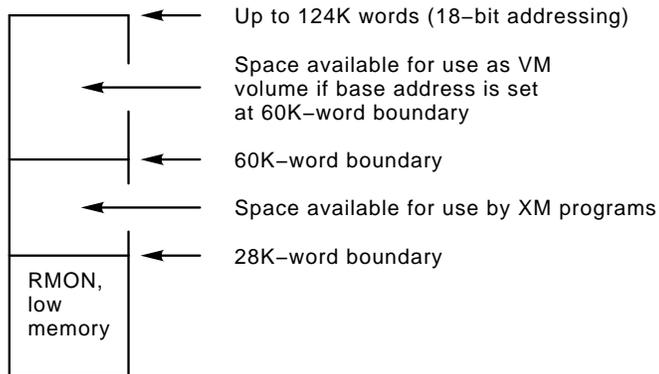


Figure 2–9: VM Handler in an 18-Bit System



still use extended memory for both an extended memory program and a VM volume, but the space available for one will be reduced by the space occupied by the other. Refer to Figure 2–9, showing an 18-bit system with the VM base address set to 3600 (60K words).

2.15 XC and XL (Communication Port (VTCOM) Handlers)

XC and XL are non-file-structured communications handlers. They support the virtual terminal communication package, VTCOM. However, their design does not preclude their use in other communication programs. The XC handler supports the CTI Bus-based computer communication port. The XL handler supports a variety of ports. See the RT-11 Software Product Description (SPD), included with your documentation set, for a list of supported ports.

XC or XL (depending on your system) is required when you use VTCOM.

XC and XL support the VTCOM utility, using `.READx`, `.WRITx`, and `.SPFUN` programmed requests.

2.15.1 `.READx` and `.WRITx` Support

The XC and XL handlers support the `.READ`, `.READC`, `.READW`, `.WRITE`, `.WRITC`, and `.WRITW` requests. You use the `.READx` and `.WRITx` requests with XC and XL handlers as described in the *RT-11 System Macro Library Manual*. Note, however, the following additional information:

- You should specify the value 0 in the *blk* argument for the first request to XC or XL. All subsequent calls should specify a nonzero value for the *blk* argument.
- NULL characters are ignored by XC and XL during both `.READs` from and `.WRITEs` to the handlers.
- XC and XL pass only 7-bit data. The eighth (high-order) bit is stripped from each byte.

2.15.2 Special Functions (`.SPFUN`) Support

In general, the XC and XL handlers support the `.SPFUN` request as described in the *RT-11 System Macro Library Manual*. Note, however, the following general information:

- You should specify the value 0 in the *blk* argument for the first request to XC or XL. All subsequent calls should specify a nonzero value for the *blk* argument.
- NULL characters are ignored by the XC and XL handlers; NULL characters are not stored or sent. However, `SF.SRD` (code 203) uses a NULL character to signal the end of available data (see `SF.SRD` in Table 2-28).
- XC and XL pass only 7-bit data. The eighth (high-order) bit is stripped from each byte.

The XC and XL handlers support the following special function codes. Specific information about using each special function is included in the description for that request.

Table 2–28: XC/XL Special Function Codes

Code	Name	Description
201	SF.CLR	<p>Resets the internal flag, indicating a received XOFF. Then sends an XON to the host.</p> <p>Example:</p> <pre>.SPFUN #area,#chan,#SF.CLR,#buf,#wcnt,#blk[,#crtn][,BMODE=str][,CMODE=str]</pre>
202	SF.BRK	<p>Sets or resets the state of the BREAK bit in the serial interface. Transition of the BREAK bit from 0 to 1 to 0 can get the attention of certain communications devices, such as terminal concentrators.</p> <p>The <i>wcnt</i> argument is a flag that indicates whether the BREAK bit should be set or reset. Specify a value of 1 for the <i>wcnt</i> argument to set the BREAK bit; specify 0 to reset it. Digital recommends you use some time delay between turning the bit on and turning it off; do that by sending one or two characters.</p> <p>Examples:</p> <p>To turn on (set) the BREAK bit:</p> <pre>.SPFUN #area,#chan,#SF.BRK,#buf,#1,#blk[,#crtn][,BMODE=str][,CMODE=str]</pre> <p>To turn off (reset) the BREAK bit:</p> <pre>.SPFUN #area,#chan,#SF.BRK,#buf,#0,#blk[,#crtn][BMODE=str][,CMODE=str]</pre>
203	SF.SRD	<p>Performs a special read from the handler. The <i>wcnt</i> argument specifies the number of bytes to be read. The read is completed when one of the following conditions is met:</p> <ul style="list-style-type: none"> • The number of bytes specified in the <i>wcnt</i> argument have been transferred. • The available characters have been transferred, when the number of available characters was less than the value specified in the <i>wcnt</i> argument. • One character has been transferred, when no characters were available when the request was issued. <p>The byte following the last transferred character contains a NULL. You must allow for that NULL byte in your buffer.</p> <p>Example:</p> <p>The following example reads no more than six (but at least one) characters from XC or XL and places them in the buffer RCVBUF. RCVBUF must be at least seven bytes in length to receive the six characters and the NULL byte.</p> <pre>.SPFUN #area,#chan,#SF.SRD,RCVBUF,#6,#blk[,#crtn][,BMODE=str][,CMODE=str]</pre>

Table 2–28 (Cont.): XC/XL Special Function Codes

Code	Name	Description
204	SF.STS	<p>Returns the driver status in the first word of the specified buffer. SF.STS always returns one word.</p> <p>The high byte of the returned word contains the driver support level. The driver support level number will be updated as support is changed in the XC and XL handlers. Programs should verify operation with an established driver support level. The current (V5.6) driver support level is 18₁₀.</p> <p>The low byte contains the status of two internal flags and a modem control signal. The significant bits of the low byte are:</p>

Bit	Meaning
-----	---------

0	Set if an XOFF has been sent to the host.
1	Set if an XOFF has been received from the host.
2	Set if the CLEAR TO SEND line is set.
3	Set if Carrier Detect is high (on); clear if Carrier Detect is low (off).
4	Set if Ring Indicator is high (on); clear if Ring Indicator is low (off).
5-7	Reserved.

Example:

The following example returns the driver support level in the high byte and the status of internal flags in the low byte of the 1-word buffer STATUS.

```
.SPFUN #area,#chan,#SF.STS,#STATUS,#1,#blk[ ,#crtn][ ,BMODE=str][ ,CMODE=str]
```

205	SF.OFF	Sets a flag that disables interrupts when the program exits. Digital recommends you issue .SPFUN SF.OFF before your program exits.
-----	--------	--

Example:

```
.SPFUN #area,#chan,#SF.OFF,#buf,#wcnt,#blk[ ,#crtn][ ,BMODE=str][ ,CMODE=str]
```

Table 2–28 (Cont.): XC/XL Special Function Codes

Code	Name	Description
206	SF.DTR	<p>Sets or resets the state of the DTR modem control signal. Setting (asserting) DTR can cause modems to answer an incoming call. Resetting (deasserting) DTR can cause modems to terminate a current call. DTR can also get the attention of certain communications devices, such as the Mini-Exchange. Specify a value of 1 for the <i>wcnt</i> argument to set the DTR control signal; specify 0 to reset the DTR control signal.</p> <p>Not all interfaces support the DTR control signal. On interfaces that do not support DTR, the setting or resetting of DTR has no effect.</p> <p>Example:</p> <p>The following example sets the DTR control signal:</p> <pre>.SPFUN #area,#chan,#SF.DTR,#buf,#1,#blk[,#crtn][,BMODE=str][,CMODE=str]</pre> <p>The following example resets the DTR control signal:</p> <pre>.SPFUN #area,#chan,#SF.DTR,#buf,#0,#blk[,#crtn][,BMODE=str][,CMODE=str]</pre>

2.15.3 EOF (End-of-File) Detection

A CTRL/Z within data being read is treated as end-of-file (EOF) by the .READ request. At least two .READ requests are necessary to return the EOF error (carry bit set and byte 52 containing error code 0). The first .READ request transfers into your buffer all data up to (but not including) the CTRL/Z. The rest of the buffer is padded with nulls. A second .READ request is required to get the EOF error. Subsequent .READ requests can return additional characters.

Appendix A

DX, DL, and XL Device Handlers

This appendix contains annotated assembly listings of the commented DX, DL, and XL device handler source files. Besides showing good handler writing practice and demonstrating the various device handler macros, each listing illustrates certain specific device handler features:

- DX illustrates a fairly simple serial device handler.
- DL illustrates software bad block replacement.
- XL illustrates internal queuing and multiterminal handler hooks.

Each device handler was assembled with both SYSMAC.SML and SYSTEM.MLB.

Figure A-1: DX Diskette Handler

DX - RX01 Floppy Disk Handler MACRO V05.05 Tuesday 26-Feb-91 14:15

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7-	1	DRIVER REQUEST ENTRY POINT
8-	1	START TRANSFER OR RETRY
9-	1	SILOFE - FILL OR EMPTY THE SILO
10-	1	TABLES, FORK BLOCK, END OF DRIVER
11-	1	BOOTSTRAP DRIVER

1 000001 mng\$t= 1

```
.MCALL .MODULE
2 000000 .MODULE DX,VERSION=17,COMMENT=<RX01 Floppy Disk Handler>,AUDIT=YES
3
4 ;
5 ;          COPYRIGHT (c) 1989 BY
6 ;          DIGITAL EQUIPMENT CORPORATION, MAYNARD, MASS.
7 ;          ALL RIGHTS RESERVED
8 ;
9 ;THIS SOFTWARE IS FURNISHED UNDER A LICENSE AND MAY BE USED AND COPIED
10 ;ONLY IN ACCORDANCE WITH THE TERMS OF SUCH LICENSE AND WITH THE
11 ;INCLUSION OF THE ABOVE COPYRIGHT NOTICE. THIS SOFTWARE OR ANY OTHER
12 ;COPIES THEREOF MAY NOT BE PROVIDED OR OTHERWISE MADE AVAILABLE TO ANY
13 ;OTHER PERSON. NO TITLE TO AND OWNERSHIP OF THE SOFTWARE IS HEREBY
14 ;TRANSFERRED.
15 ;
16 ;THE INFORMATION IN THIS SOFTWARE IS SUBJECT TO CHANGE WITHOUT NOTICE
17 ;AND SHOULD NOT BE CONSTRUED AS A COMMITMENT BY DIGITAL EQUIPMENT
CORPORATION.
```

```

18          ;
19          ;DIGITAL ASSUMES NO RESPONSIBILITY FOR THE USE OR RELIABILITY OF ITS
20          ;SOFTWARE ON EQUIPMENT THAT IS NOT SUPPLIED BY DIGITAL.

```

CONDITIONAL ASSEMBLY SUMMARY

```

1          .SBTTL  CONDITIONAL ASSEMBLY SUMMARY
2          ;+
3          ;COND
4          ;      DXT$O  (0)          Two controller support
5          ;              0          support 1 controller
6          ;              1          support 2 controllers
7          ;
8          ;      DX$CSR (177170)     primary CSR
9          ;      DX$CS2 (177174)     second CSR
10         ;
11         ;      DX$VEC (264)         primary Vector
12         ;      DX$VC2 (270)         second Vector
13         ;
14         ;      MMG$T                 std conditional
15         ;      TIM$IT                 std conditional (no code effects)
16         ;      ERL$G                 std conditional
17         ;-

```

Preamble Section

```

1          .SBTTL  DEFINITIONS
2
3          .ENABL  LC
4

```

Monitor offsets and SYSCOM locations are defined with mnemonics so that references to them can be found easily:

```

5          ; SOME RT-11 MACROS WE WILL USE
6
7          .MCALL .DRDEF .ASSUME .BR .ADDR
8
9          000342 .DSTATUS=:342          ;EMT code for .DSTATUS
10         000375 .READ  =:375          ;EMT code for .READ
11         000010 ..READ  =:010         ; subcode for .READ
12         000375 .WRITE =:375          ;EMT code for .WRITE
13         000011 ..WRIT =:011         ; subcode for .WRITE
14
15         000017 SYSCHN =:17           ; system channel
16
17         ; RT-11 SYSCOM LOCATIONS
18
19         000044 JSW    =:44           ;JOB STATUS WORD
20         000054 SYSPTR =:54          ;POINTER TO BASE OF RMON
21         000432 PLEXT =: 432         ;OFFSET FROM $RMON TO EXTERNAL ROUTINE
22

```

If DXT\$O=1, there are two controllers:

```

23         ; RX01 CONTROLLER DEFAULTS
24
25         .IIF NDF DXT$O, DXT$O=0     ;DEFAULT TO ONLY ONE CONTROLLER
26
27         .IIF NDF DX$CS2, DX$CS2 == 177174 ;2ND CONTROLLER CSR
28         .IIF NDF DX$VC2, DX$VC2 == 270   ;2ND CONTROLLER VECTOR
29

```

The .DRDEF macro (with macro expansion):

```

30 000000          .DRDEF  DX,22,FILST$!SPFUN$!DX$COD,494.,177170,264,DMA=NO
                  .MCALL  .DRAST, .DRBEG, .DRBOT, .DREND, .DREST, .DRFIN, .DRFMS, .DRFMT
                  .MCALL  .DRINS, .DRPTR, .DRSET, .DRSPF, .DRTAB, .DRUSE, .DRVTE
                  .MCALL  .FORK, .QELDF
                  .IIF NDF RTE$M RTE$M=0
                  .IIF NE RTE$M RTE$M=1
                  .IIF NDF TIM$IT TIM$IT=0
                  .IIF NE TIM$IT TIM$IT=1
                  .IIF NDF MMG$T MMG$T=0
000001          .IIF NE MMG$T MMG$T=1
                  .IIF NDF ERL$G ERL$G=0
                  .IIF NE ERL$G ERL$G=1
                  .IIF NE TIM$IT, .MCALL  .TIMIO, .CTIMI
000000          .QELDF
                  .IIF NDF MMG$T,MMG$T=1
000001          .IIF NE MMG$T,MMG$T=1
000000          Q.LINK=:0
000002          Q.CSW=:2.
000004          Q.BLKN=:4.
000006          Q.FUNC=:6.
000007          Q.JNUM=:7.
000007          Q.UNIT=:7.
000010          Q.BUFF=:^o10
000012          Q.WCNT=:^o12
000014          Q.COMP=:^o14
                  .IRP   X,<LINK,CSW,BLKN,FUNC,JNUM,UNIT,BUFF,WCNT,COMP>
                  Q$'X=:Q.'X-^o4
                  .ENDR
DX - RX01 Floppy Disk Handler  MACRO V05.05  Tuesday 26-Feb-91 19:46  Page 4-1
DEFINITIONS

```

```

177774  Q$LINK=:Q.LINK-^o4
177776  Q$CSW=:Q.CSW-^o4
000000  Q$BLKN=:Q.BLKN-^o4
000002  Q$FUNC=:Q.FUNC-^o4
000003  Q$JNUM=:Q.JNUM-^o4
000003  Q$UNIT=:Q.UNIT-^o4
000004  Q$BUFF=:Q.BUFF-^o4
000006  Q$WCNT=:Q.WCNT-^o4
000010  Q$COMP=:Q.COMP-^o4
                  .IF  EQ MMG$T
                  Q.ELGH=:^o16
                  .IFF
000016  Q.PAR=:^o16
000020  Q.MEM=:^o20
                  .IRP   X,<PAR,MEM>
                  Q$'X=:Q.'X-^o4
                  .ENDR
000012  Q$PAR=:Q.PAR-^o4
000014  Q$MEM=:Q.MEM-^o4
000024  Q.ELGH=:^o24
                  .ENDC
000001  HDERR$=:1
020000  EOF$=:^o20000
000400  VARSZ$=:^o400
001000  ABTIO$=:^o1000
002000  SPFUN$=:^o2000
004000  HNDLR$=:^o4000
010000  SPECL$=:^o10000
020000  WONLY$=:^o20000
040000  RONLY$=:^o40000
100000  FILST$=:^o100000
000756  DXDSIZ=:494.
000022  DX$COD=:22
102022  DXSTS=:<22>!<FILST$!SPFUN$!DX$COD>
                  .IIF NDF DX$VEC,DX$VEC=264
                  .GLOBL DX$VEC

```

The .DRPTR macro with no parameters:

```

31 000200          .DRPTR

```

The .DREST macro to define handler class and class modifier:

The .DRSPF macro to define supported special functions:

```

33 000076          .DRSPF <377>          ;Read Absolute
34 000032          .DRSPF <376>          ;Write Absolute
35 000032          .DRSPF <375>          ;Write Deleted
36
37                ; CONTROL AND STATUS REGISTER BIT DEFINITIONS
38
39 000001          CSGO    ==:    1          ;INITIATE FUNCTION
40 000020          CSUNIT ==:    20         ;UNIT BIT
41 000040          CSDONE ==:    40         ;DONE BIT
42 000100          CSINT  ==:    100        ;INTERUPT ENABLE
43 000200          CSTR   ==:    200        ;TRANSFER REQUEST
44 004000          CSRX02 ==:    4000       ;CONTROLLER IS RX02 (ALWAYS 0)
45 040000          CSINIT ==:    40000     ;RX11 INITIALIZE
46 100000          CSERR  ==:100000       ;ERROR
47
48                ; CSR FUNCTION CODES IN BITS 1-3
49
50 000000          CSFBUF ==:0*2          ;0 - FILL SILO (PRE-WRITE)
51 000002          CSEBUF ==:1*2          ;1 - EMPTY SILO (POST-READ)
52 000004          CSWRT  ==:2*2          ;2 - WRITE SECTOR
53 000006          CSRD   ==:3*2          ;3 - READ SECTOR
54                ;4 - UNUSED
55 000012          CSRDST ==:5*2          ;5 - READ STATUS
56 000014          CSWRTD ==:6*2          ;6 - WRITE SECTOR WITH DELETED DATA
57 000016          CSMAIN ==:7*2          ;7 - MAINTENANCE
58
59 000002          CSREAD ==:CSEBUF&CSRD&CSRDST&CSMAIN
60
61 000032          .ASSUME CSRD&2          NE 0    ;2 BIT MUST BE ON IN READ
62 000032          .ASSUME CSWRT&2        EQ 0    ;2 BIT MUST BE OFF IN WRITE
63 000032          .ASSUME CSWRTD&2       EQ 0    ;2 BIT MUST BE OFF IN WRITE
64
65                ; ERROR AND STATUS REGISTER BIT DEFINITIONS
66
67 000001          ESCRC  ==:    1          ;CRC ERROR
68 000002          ESPAR  ==:    2          ;PARITY ERROR
69 000004          ESID   ==:    4          ;INITIALIZE DONE
70 000100          ESDD   ==:    100        ;DELETED DATA MARK
71 000200          ESDRY  ==:    200        ;DRIVE READY
72
73                ; ERROR LOG VALUES
74
75 000003          DXNREG ==:3            ;# OF REGISTERS TO READ FOR ERROR LOG.
76 000010          RETRY  ==:8.          ;RETRY COUNT
77
78 100000          SPFUNC ==:100000       ;SPECIAL FUNCTIONS FLAG
79                ; ( IN COMMAND WORD)
80
81                ; GENERAL COMMENTS:
82                ;
83                ; THIS HANDLER SERVES AS THE STANDARD RT-11 RX01 DEVICE HANDLER AS
84                ; BOTH THE SYSTEM DEVICE HANDLER AND NON-SYSTEM HANDLER. IT ALSO PRO-
85                ; VIDES THREE SPECIAL FUNCTION CAPABILITIES TO SUPPORT PHYSICAL I/O
86                ; ON THE FLOPPY AS A FOREIGN VOLUME. THE SPECIAL FUNCTIONS ARE:
87                ;   CODE   ACTION
88                ;   377   ABSOLUTE SECTOR READ. WCNT=TRACK, BLK=SECTOR, BUFFER=65
89                ;         WORD BUFFER OF WHICH WORD 1 IS DELETED DATA FLAG.
90                ;   376   ABSOLUTE SECTOR WRITE. ARGUMENTS SAME AS READ.
91                ;   375   ABSOLUTE SECTOR WRITE WITH DELETED DATA. 1ST WORD
92                ;         OF 65 WORD BUFFER ALWAYS SET TO 0.
93                ;
94                ; IN STANDARD RT-11 MODE A 2:1 INTERLEAVE IS USED ON A SINGLE TRACK AND
95                ; A 6 SECTOR SKEW IS USED ACROSS TRACKS. TRACK 0 IS LEFT ALONE FOR
96                ; PROPOSED ANSI COMPATABILITY.

```

Installation checks:

```

1          .SBTTL  INSTALLATION CHECKS
2
3          .IF EQ DXT$O
4 000032      .DRINS  DX
5          .IFF
6          .DRINS  DX,<DX$CS2>
7          .ENDC ;EQ DXT$O
8
9 000200 000240      NOP          ;SAME CHECK FOR SYSTEM AND NON-SYSTEM HANDLER
10 000202 032777     BIT          #CSRX02,@INSCSR ;IS THE RX02 BIT ON?
11          004000
12          177766
11 000210 001561     BEQ          O.GOOD        ;NOPE, IS AN RX01, INSTALL IT
12 000212 000561     BR           O.BAD         ;YES, AN RX02, DON'T INSTALL IT
13
14          ; Routine to find the entry for DX in the monitor device tables
15
16 000214      FINDRV:
17 000214          .ADDR  #DEVNAM,R0          ;R0->DEVICE NAME
18 000222          .ADDR  #DAREA+1,-(SP)      ;(SP)->.DSTATUS INFO AREA(+physical)
19 000230 104342     EMT          .DSTATUS     ;*** (.DSTAT #DAREA+1,#DEVNAM) ***
20 000232 103551     BCS          O.BAD        ;IN CASE IT'S NOT KNOWN
21 000234 016701     MOV          DAREA+4,R1   ;RETURN THE ENTRY POINT
22          000010
22 000240 001145     BNE          O.GOOD
23 000242 000545     BR           O.BAD        ;UNLESS HANDLER'S NOT LOADED
24
25 000244      DAREA: .BLKW  4                ;.DSTAT INFORMATION BLOCK
26 000254 016300     DEVNAM: .RAD50 /DX /     ;DEVICE NAME
27
28          ; The emt area for reads/writes of the handler is placed here
29          ; to leave room for code for the set options
30
31 000256      017 BAREA: .BYTE  SYSCHN,..READ ;CHANNEL 17, READ
32          000257  010
33 000260          .BLKW                    ;BLOCK NUMBER
34 000262          .BLKW                    ;BUFFER
35 000264 000400     .WORD  256.            ;WORD COUNT
36 000266 000000     .WORD  0              ;COMPLETION (WAIT)
37
38          ; NOW ALTER THE CODE WHICH WILL BE WRITTEN BACK TO DISK
39 000270      X.WP:
40 000270          .ADDR  #DXWPRO,R0          ;R0-> THE WRITE PROTECT TABLE
41 000276 060300     ADD          R3,R0        ; POINT TO ENTRY
42 000300 112710     MOV          (PC)+,(R0)    ; AND SET IT THE WAY THE USER WANTS IT
43 000302      O.WPF: .BLKW  1
44
45          ; NOW TO ALTER THE IN-CORE COPY OF THE PROTECTION TABLE
46
47 000304 004767     CALL          FINDRV        ;IS THE HANDLER LOADED?
48          177704
48 000310 103521     BCS          O.GOOD        ;NOPE...
49 000312 023701     CMP          @#SYSPTR,R1    ;is this the system handler?
50          000054
50 000316 101003     BHI          10$                ; no, then leave 1-shot as is
51 000320 012761     MOV          #100000,DXW1-DXLQE(R1) ; yes, set it
52          100000
53          000076
52 000326      10$:
53 000326 060301     ADD          R3,R1                ;ADD IN UNIT OFFSET
54 000330 116761     MOV          O.WPF,DXWPRO-DXLQE(R1) ;SET THE WRITE-PROTECT STATUS
55          177746
56          000010
55 000336 000506     BR           O.GOOD
56
57          .IIF GT,<.-376> .ERROR ;INSTALLATION CODE IS TOO LARGE;

```

The DX handler supports several SET options. Immediately following the installation code, the .DRSET macro is used to define the parameter table for each SET option:

```

1          .SBTTL SET OPTIONS
2
3          ; The write-protect/enable SET option makes use of the new
4          ; calling convention, i.e. the unit number (DXn, n=0 if a space)
5          ; passed in R1.
6
7 000340          .DRSET CSR,      160000, O.CSR,  OCT
8 000412          .DRSET VECTOR,  500,      O.VEC,  OCT
9
10         .IF NE DXT$O
11         .DRSET CSR2,      160000, O.CSR2, OCT
12         .DRSET VEC2,      500,      O.VEC2, OCT
13         .ENDC;NE DXT$O
14
15 000422          .DRSET RETRY,  127.,      O.RTRY,  NUM
16
17         .IF NE ERL$G
18         .DRSET SUCCES,  -1,      O.SUCC,  NO
19         .ENDC;NE ERL$G
20
21 000432          .DRSET WRITE,  1,      O.WP,  NO
22
23         002256          BTCSR = <DXEND-DXSTRT>+<BOTCSR-DXBOOT>+1000
24

```

The code to process each SET options follows the .DRSET macro calls. Normally, SET options change only the disk-resident copy of a handler, not the memory-resident copy. The DX handler SET options include special code to modify both the memory-resident and the disk-resident copy of the handler.

```

25 000442 020003 O.CSR: CMP      R0,R3          ;IS CSR IN RANGE? (>160000)
26 000444 103444 BLO      O.BAD          ;NOPE...
27 000446 010067 MOV      R0,INSCSR      ;YES, INSTALLATION CODE NEEDS IT
28 000452 010067 MOV      R0,DISCSR     ;FILL IN DISPLAY CSR
29
30         ; When the csr for units 0 and 1 is changed, the bootstrap must
31         ; be altered such that it will use the correct controller.
32
33         ;R1->READ/WRITE EMT AREA
34 000456          .ADDR  #BAREA+4,R1          ; (BUFFER ADDRESS WORD)
35         ;BUILD ADDRESS OF BUFFER
36 000464          .ADDR  #1000,R2           ; (WHICH WILL OVERWRITE CORE
37         ; COPY OF BLOCK 1)
38 000472 010211 MOV      R2,(R1)          ;SET THE BUFFER ADDRESS
39 000474 012741 MOV      #BTCSR/1000,-(R1)       ;SET TO BLOCK NUMBER TO READ/WRITE
40         ; (BOOT BLOCK THAT NEEDS MODIFICATION)
41 000500 005741 TST      -(R1)          ;R1->EMT AREA
42 000502 010003 MOV      R0,R3          ;SAVE CSR ELSEWHERE, EMT NEEDS R0
43 000504 010100 MOV      R1,R0          ;R0->EMT AREA FOR READ
44 000506 104375 EMT      .READ          ; *** (.READW) ***
45 000510 103422 BCS      O.BAD          ;
46 000512 010362 MOV      R3,<BTCSR&777>(R2)       ;SET THE NEW CSR
47 000516 010100 MOV      R1,R0          ;R0->EMT AREA FOR WRITE
48 000520          .ASSUME ..READ+1 EQ ..WRIT
49 000520 105260 INCB     1(R0)          ;BUMP FROM 'READ' TO 'WRITE'
50 000524 104375 EMT      .WRITE          ; *** (.WRITW) ***
51 000526 103415 BCS      O.SYWL         ; SY: write-locked
52 000530 010100 MOV      R1,R0          ;R0->EMT AREA
53 000532          .ASSUME ..WRIT-1 EQ ..READ
54 000532 105360 DECB     1(R0)          ;CHANGE FROM 'WRITE' TO 'READ'
55 000536 012760 MOV      #1,2(R0)          ; OF BLOCK 1 OF HANDLER
56 000544 104375 EMT      .READ          ; *** (.READW) ***
57 000546 103403 BCS      O.BAD          ;
58
59         .IF EQ DXT$O
60 000550 010367 MOV      R3,RXCSA
61 000504'

```

```

61          .IFF
62          MOV     R3,DXCSR
63          .ENDC ;EQ DXT$O
64
65 000554 005727 O.GOOD: TST     (PC)+          ;GOOD RETURN (CARRY CLEAR)
66 000556 000261 O.BAD:  SEC          ;ERROR RETURN (CARRY SET)
67 000560 000207          RETURN
68
69 000562          O.SYWL:
70 000562 011600          MOV     @SP,R0          ; copy return address
71 000564 005200          INC     R0              ; point to opcode at return
72 000566 122720          CMPB   #BR/400,(R0)+   ; is it a BR xxx?
73          000001
73 000572 001371          BNE    O.BAD           ; NO, old style SET
74 000574 010016          MOV    R0,@SP         ; use alternate return (RET+2)
75 000576 000767          BR     O.BAD           ; with carry set
76
77 000600 020003 O.VEC:  CMP     R0,R3          ;VECTOR IN RANGE?
78 000602 103365          BHIS   O.BAD           ;NOPE...
79 000604 032700          BIT    #3,R0          ;YES, BUT ON A VECTOR BOUNDRY?
80          000003
80 000610 001362          BNE    O.BAD           ;NOPE...
81
82          .IF EQ DXT$O
83 000612 010067          MOV    R0,DXSTRT      ;YES, SET IT IN ENTRY AREA
84          000000'
85          .IFF
86          MOV     R0,DX$VTB          ;PLACE IT IN MULTI-VECTOR TABLE
87          .ENDC ;NE DXT$O
88 000616 000756          BR     O.GOOD
89
90          .IF NE DXT$O
91          O.CSR2:  CMP     R0,R3          ;CSR IN RANGE?
92          BLO    O.BAD           ;NOPE...
93          MOV    R0,DXCSR2        ;YES, PLACE IT IN CODE
94          MOV    R0,DISCS2        ;SET DISPLAY CSR
95          BR     O.GOOD
96
97          O.VEC2:  CMP     R0,R3          ;VECTOR IN RANGE?
98          BHIS   O.BAD           ;NOPE...
99          BIT    #3,R0          ;YES, BUT IS IT ON A VECTOR BOUNDARY?
100         BNE    O.BAD           ;NOPE...
101         MOV    R0,DX$VTB+6       ;YES, PLACE IN MULTI-VECTOR TABLE
102         BR     O.GOOD
103         .ENDC ;NE DXT$O
104
105 000620 020003 O.RTRY:  CMP     R0,R3          ;ASKING FOR TOO MANY?
106 000622 101355          BHI    O.BAD           ;YES, USER IS BEING UNREASONABLE
107 000624 010067          MOV    R0,DRETRY      ;NOPE, SO TELL THE HANDLER
108          000034'
108 000630 001351          BNE    O.GOOD         ;OKAY IF NON-ZERO
109 000632 000751          BR     O.BAD           ;CAN'T ASK FOR NO RETRIES
110
111         .IF NE ERL$G
112         O.SUCC:  MOV     #0,R3          ;'SUCCESS' ENTRY POINT
113         ; (MUST BE TWO WORDS)
114         N.SUCC:  MOV     R3,SCSFLG      ;'NOSUCCES' ENTRY POINT
115         ;.ASSUME O.SUCC+4 EQ N.SUCC
116         BR     O.GOOD
117         .ENDC ;NE ERL$G
118
119 000634 000240 O.WP:  NOP          ;'WRITE' ENTRY POINT
120 000636 005003          CLR    R3              ;CLEAR FLAG
121 000640          N.WP:  ;'NOWRITE' ENTRY POINT
122 000640          .ASSUME O.WP+4 EQ N.WP
123 000640 010367          MOV    R3,O.WPF       ;SAVE THE USER'S SELECTION
124          177436
124 000644 010103          MOV    R1,R3          ; save unit number
125 000646 020327          CMP    R3,#DXT$O*2+1 ;IS IT A VALID UNIT
126          000001
126 000652 101341          BHI    O.BAD           ;NOPE...
127 000654 000167          JMP    X.WP           ; go to rest of the code
127          177410
128

```

All of the code to process SET options must fit within the first block of the handler. The following line tests to make sure that this condition is satisfied:

```
129                .IIF GT,<.-1000> .ERROR ;SET CODE IS TOO LARGE;
```

Header Section

```
1                .SBTTL DRIVER REQUEST ENTRY POINT
2
3                .ENABL LSB
4
```

The .DRBEG macro:

```
5 000660        .DRBEG DX
```

I/O Initiation Section

```
6 000014 000401    BR        DXENT        ;BRANCH AROUND PROTECTION TABLE
7
8 000016                DXWPRO:
9                000001    .REPT    DXT$O+1
10               .BYTE    0,0
11               .ENDR
12 000020                .ASSUME . LE DXSTRT+1000
13
14               .IF NE ERL$G
15               SCSFLG: .WORD    0        ; :SUCCESSFUL LOGGING FLAG (DEFAULT=YES)
16               ; =0 - LOG SUCCESSES,
17               ; <>0 - DON'T LOG SUCCESSES
18               .ASSUME . LE DXSTRT+1000
19               .ENDC ;NE ERL$G
20
21               .IF NE DXT$O
22               .DRVTB DX,DX$VEC,DXINT
23               .DRVTB ,DX$VC2,DXINT
24               .ENDC ;NE DXT$O
25
26 000020                DXENT:
27               .IF NE MMG$T
28 000020 013704        MOV        @#SYSPTR,R4        ; R4 -> MONITOR BASE
29               000054
30 000024 016427        MOV        PLEXT(R4),(PC)+ ; GET ADDRESS OF EXTERNALIZATION ROUTINE
31               000432
32 000030 000432        $PLEXT: .WORD    PLEXT        ; POINTER TO EXTERNALIZATION ROUTINE
33               .ENDC ;NE MMG$T
34
35 000032 012727        MOV        (PC)+,(PC)+ ;INITIALIZE RETRY COUNT
36 000034 000010        DRETRY: .WORD    RETRY        ; :RETRY MAXIMU
37               000036 .ASSUME . LE DXSTRT+1000
38 000036 000000        RXTRY: .WORD    0        ; :CURRENT RETRY COUNT
```

The following instructions assemble the controller function to start up an operation and sort out special functions.

```
37
38 000040 016703        MOV        DXCQE,R3        ;GET POINTER TO QUEUE ELEMENT
39               177744
40 000044 012305        MOV        (R3)+,R5        ;GET BLOCK NUMBER
41 000046 012704        MOV        #CSRD!CSGO,R4   ;GUESS THAT CONTROLLER FUNCTION IS READ
42               000007
43 000052                .ASSUME Q$BLKN+2 EQ Q$FUNC
44 000052 112301        MOVB       (R3)+,R1        ;PICK UP SPECIAL FUNCTION CODE (SIGN EXTENDED)
45 000054                .ASSUME Q$FUNC+1 EQ Q$UNIT
46 000054 112300        MOVB       (R3)+,R0        ;PICK UP THE UNIT NUMBER
47 000056 106200        ASRB       R0            ;SHIFT IT TO CHECK FOR ODD UNIT
48 000060 103002        BCC        1$            ;BRANCH IF EVEN UNIT
49 000062 052704        BIS        #CSUNIT,R4    ;SELECT ODD UNIT FOR TRANSFER
50               000020
51 000066                1$:
52 000066                .IF EQ DXT$O ;ONE CONTROLLER
```

```

50 000066 132700 BITB #6/2,R0 ;ANY UNITS BUT 0 OR 1?
    000003
51 000072 001163 BNE RXERR ;BRANCH IF YES, ERROR
52 .IFF
53 MOV (PC)+,(SP) ;ASSUME FIRST DX CONTROLLER
54 DXCSR = .
55 .WORD DX$CSR
56 .ASSUME . LE DXSTRT+1000
57 ASRB R0 ;SHIFT UNIT TO CHECK FOR SECOND CONTROLLER
58 BCC 2$ ;NOPE, FIRST CONTROLLER
59 MOV (PC)+,(SP) ;CHANGE CSR TO USE SECOND CONTROLLER
60 DXCSR2 = .
61 .WORD DX$CSR2
62 .ASSUME . LE DXSTRT+1000
63 2$: MOV (SP)+,RXCSA
64 ASRB R0 ;BUT WAS IT UNIT 4 TO 7?
65 BCS RXERR ;ERROR IF SO
66 .ENDC ;EQ DXT$O
'' 67 000074 .ASSUME Q$UNIT+1 EQ Q$BUFF
68 000074 012300 MOV (R3)+,R0 ;GET THE USER'S BUFFER ADDRESS
69 000076 .ASSUME Q$BUFF+2 EQ Q$WCNT
70 000076 012302 MOV (R3)+,R2 ;GET WORD COUNT
71 000100 100017 BPL 3$ ;POSITIVE MEANS READ, SO ALL SET UP
72
73 ; HERE TO CHECK IF UNIT IS WRITE-PROTECTED
74
75 000102 006327 ASL (PC)+ ; CHECK WRITE ANYWAY ONE-SHOT
76 000104 000000 DXW1: .WORD -. ; 100000 MEANS WRITE ANYWAY
77 000106 .ASSUME . LE DXSTRT+1000
78 000106 103412 BCS 33$ ; SKIP TEST IF WRITE ANYWAY
79 000110 005046 CLR -(SP) ;SET TO GET UNIT
80 000112 .ASSUME Q$WCNT+2 EQ Q$COMP
81 000112 116316 MOVB Q$UNIT-Q$COMP(R3),(SP) ;GET IT (PLUS OTHER CRUFT
    177773
82 000116 042716 BIC #<^C3>,(SP) ; WHICH WE DISCARD NOW
    177774
83 ;ADD ADDRESS OF WRITE-PROTECT TABLE
84 000122 .ADDR #DXWPRO,(SP),ADD; TO UNIT OFFSET
85 000130 105736 TSTB @(SP)+ ;CHECK UNIT WRITE STATUS
86 000132 001143 BNE RXERR ;IT'S WRITE-PROTECTED, USER CAN'T DO THIS
87 000134 .ASSUME CSRD-2 EQ CSWRT
88 000134 124444 33$: CMPB -(R4),-(R4) ;CHANGE CSRD (3*2) TO CSWRT (2*2) FOR WRITE

```

Ensure that a write equals a read code minus 2:

```

89 000136 .ASSUME CSWRT EQ CSRD-2
90 000136 005402 NEG R2 ; AND MAKE WORD COUNT POSITIVE
91 000140 006301 3$: ASL R1 ;DOUBLE THE SPECIAL FUNCTION CODE
92 000142 060701 ADD PC,R1 ;FORM PIC REFERENCE TO CHGTBL

```

The codes for read and write operations stay the same. If the operation is for a special function, this routine sets the sign bit of the function code word, and modifies the function:

```

93 000144 066104 ADD CHGTBL-(R1),R4 ;MODIFY THE CODE, SET SIGN BIT IF SPFUN
    000740
94 000150 010467 MOV R4,RXFUN2 ;SAVE THE FUNCTION CODE AND SPFUN FLAG
    000320
95 000154 100435 BMI 7$ ;IF SPFUN, GO DO SPECIAL SETUP
96
97 ; NORMAL I/O, CONVERT TO TRACK AND SECTOR NUMBER AND INTERLEAVE
98

```

FILLCT indicates whether a multiple of four sectors has been written. If not, the handler will later zero-fill to reach a multiple of four.

```

99 000156 110267      MOVB   R2,FILLCT      ;SAVE WORD COUNT IN CASE WE HAVE TO FILL
000537
100 000162 105367      DECB   FILLCT        ; EXTRA SECTORS ON WRITE
000533
101 000166 006302      ASL    R2            ;MAKE WORD COUNT UNSIGNED BYTE COUNT
102 000170 006305      ASL    R5            ;NORMAL READ/WRITE. COMPUTE REAL SECTOR NUMBER
103 000172 006305      ASL    R5            ; AS BLOCK*4
104 000174 012704      MOV    (PC)+,R4      ;LOOP COUNT FOR 8 BIT DIVISION
105 000176      371      .BYTE  -7,-26.      ;COUNT BECOMES 1, -26 IN HIGH BYTE FOR LATER
000177      346
106 000200 022705 4$:      CMP    #26.*200,R5  ;DOES 26 GO INTO DIVIDEND?
006400
107 000204 101002      BHI    5$           ;BRANCH IF NOT, C CLEAR
108 000206 062705      ADD    #-26.*200,R5 ;SUBTRACT 26 FROM DIVIDEND, SET C
171400
109 000212 006105 5$:      ROL    R5           ;SHIFT DIVIDEND AND QUOTIENT
110 000214 105204      INCB   R4           ;DECREMENT LOOP COUNT
111 000216 003770      BLE    4$           ;BRANCH UNTIL DIVIDE DONE
112 000220 110501      MOVB   R5,R1        ;COPY TRACK NUMBER 0:75, ZERO EXTEND
113 000222 060405      ADD    R4,R5        ;BUMP TRACK TO 1-76, MAKE SECTOR<0
114 000224 010104      MOV    R1,R4        ;COPY TRACK NUMBER
115 000226 006301      ASL    R1           ;MULTIPLY
116 000230 060401      ADD    R4,R1        ; BY
117 000232 006301      ASL    R1           ; 6
118 000234 162701 6$:      SUB    #26.,R1      ;REDUCE TRACK NUMBER * 6 MOD 26
000032
119 000240 003375      BGT    6$           ; TO FIND OFFSET FOR THIS TRACK, -26:0
120 000242 010167      MOV    R1,TRKOFF   ;SAVE IT
000132
121 000246 000412      BR     8$           ;GO SAVE PARAMETERS AND START
122
123          ; SPECIAL FUNCTION REQUEST, SET TRACK AND SECTOR AND BYTE COUNT
124

```

The routine passes a 65-word buffer. The first word is 0 if there is no deleted data mark.

```

125 000250 000305 7$:      SWAB   R5           ;PUT PHYSICAL SECTOR IN HIGH BYTE
126 000252 150205      BISB   R2,R5        ; AND PHYSICAL TRACK IN LOW BYTE
127 000254 012702      MOV    #128.,R2     ;SET THE BYTE COUNT TO 128
000200
128
129          .IF EQ MMG$T
130          CLR    (R0)+ ;CLEAR DELETED DATA FLAG WORD, BUMP USER ADDR
131          .IFF
132 000260 016704      MOV    DXCQE,R4    ;POINT TO QUEUE ELEMENT AT Q.BLKN
177524
133 000264 005046      CLR    -(SP)        ;STACK A ZERO AND STORE IT IN FIRST WORD OF
134 000266 004777      CALL  @$PTWRD      ; BUFFER. NOTE THAT Q.BUFF GETS BUMPED BY 2
000634
135 000272 005720      TST    (R0)+        ;ADD 2 TO OUR COPY OF USER BUFFER ADDRESS
136          .ENDC ;EQ MMG$T
137
138          ; MERGE HERE TO START OPERATION

```

Save the user virtual buffer address, the track, the byte count, and the PAR1 value for mapped systems:

```

139
140 000274 010027 8$:      MOV    R0,(PC)+    ;SAVE BUFFER ADDRESS
141 000276 000000  BUFRAD: .WORD  0    ; : USER VIRTUAL BUFFER ADDRESS
142 000300 010567      MOV    R5,TRACK    ;SAVE IT FOR STARTING I/O
000126
143 000304 010227      MOV    R2,(PC)+    ; AND BYTE COUNT.
144 000306 000000  BYTCNT: .WORD  0    ; : BYTE COUNT FOR TRANSFER
145
146          .IF NE MMG$T
147 000310 005723      TST    (R3)+        ;SKIP THE COMPLETION ROUTINE ADDRESS
148 000312 011367      MOV    @R3,PARVAL  ;SAVE THE PAR1 VALUE FOR MAPPING USER BUFFER
000542
149          .ENDC ;NE MMG$T
150
151 000316          .BR     RXINIT    ;GO TO FORK LEVEL AND START IT UP
152

```

The calculations are done; the routine can now start an operation or a retry. Before it starts, however, it arranges transfer routines for interrupt entry. To get to the ready state, force one interrupt, then return to 1\$:

```

1          .SBTTL  START TRANSFER OR RETRY
2
3          .ENABL  LSB
4
5 000316  012767  RXINIT: MOV    #100000,RXIRTN  ;SET RETURN AFTER INITIAL INTERRUPT
                100000
                000172
6 000324  016704          MOV    RXCSA,R4          ;ENSURE THAT WE POINT TO THE CSR
                000154
7 000330  000441          BR     RXIENB          ;GO INTERRUPT, RETURN TO 1$ LATER
8
9 000332  032700  1$:     BIT    #CSREAD,R0      ;READ OR WRITE FUNCTION?
                000002
10 000336  001005          BNE   3$              ;IF READ, GO FILL THE SILO FROM DISK
11 000340  004067  2$:     JSR    R0,SILOFE      ;WRITE, LOAD THE SILO FROM THE USER BUFFER
                000440

```

Parameters for SIOFE routine:

```

12 000344  000001          .WORD  CSFBUF!CSGO      ; FILL BUFFER COMMAND
13 000346  112215          MOVB  (R2)+,@R5        ; MOVB TO BE PLACED IN-LINE IN SILOFE
14 000350  010115          MOV   R1,@R5         ; ZERO-FILL INSTRUCTION FOR SHORT WRITES

```

The following routine changes a sector number to an interleaved sector number:

```

15 000352  116702  3$:     MOVB  SECTOR,R2      ;GET THE SECTOR NUMBER
                000055
16 000356  003014          BGT   5$              ;POSITIVE MEANS SPFUN, DON'T INTERLEAVE
17 000360  162702          SUB   #-14.,R2       ;ADD 14 TO DO INTERLEAVING
                177762
18 000364  003003          BGT   4$              ;IF > 0, MAP -13:-1 TO 2:26, NOTE C=0
19 000366  062702          ADD   #12.,R2       ; ELSE MAP -26:-14 TO 1:25
                000014
20 000372  000261          SEC                      ;ADD 1 WHEN DOUBLING
21 000374  006102  4$:     ROL   R2              ;DOUBLE AND INTERLEAVE, SECTOR 1:26
22 000376  062702          ADD   (PC)+,R2      ;ADD IN THE TRACK OFFSET, SECTOR -25:26
23 000400  000000  TRKOFF: .WORD  0          ; : TRACK OFFSET = TRACK*6 MOD 26, RANGE -26:0
24 000402  003002          BGT   5$              ;NO MODULUS PROBLEMS
25 000404  062702          ADD   #26.,R2       ;FIX TO PUT SECTOR IN 1:26 RANGE
                000032
26 000410  010014  5$:     MOV   R0,@R4          ;SET THE FUNCTION IN THE FLOPPY CONTROLLER
27 000412  105714  6$:     TSTB  @R4              ;WAIT FOR
28 000414  001776          BEQ   6$              ; TRANSFER READY
29 000416  100161          BPL  RXRTRY         ;TRANSFER DONE WITHOUT TRANSFER READY, ERROR
30 000420  110215          MOVB  R2,@R5        ;SET SECTOR NUMBER
31 000422  105714  7$:     TSTB  @R4              ;WAIT AGAIN FOR
32 000424  001776          BEQ   7$              ; TRANSFER READY
33 000426  100155          BPL  RXRTRY         ;TRANSFER DONE WITHOUT TRANSFER READY, ERROR
34 000430  112715          MOVB  (PC)+,@R5    ;SET THE TRACK NUMBER
35 000432          000  TRACK:  .BYTE  0          ;TRACK NUMBER
36 000433          000  SECTOR: .BYTE  0          ;SECTOR NUMBER, KEPT < 0 UNLESS SPFUN

```

Start the operation and return to the monitor:

```

37 000434  052714  RXIENB: BIS    #CSINT,@R4      ;SET IE TO CAUSE AN INTERRUPT WHEN DONE IS UP
                000100
38 000440  000207          RETURN          ;RETURN, WE'LL BE BACK WITH AN INTERRUPT
39
40 000442  016704  RXERR:  MOV    DXCQE,R4        ;R4 -> CURRENT QUEUE ELEMENT
                177342
41 000446  052754          BIS    #HDERR$,@-(R4) ;SET HARD ERROR IN CSW
                000001
42 000452  000524          BR     13$       ;EXIT ON HARD ERROR
43

```

Interrupt Service Section

The .DRAST macro:

```
44 000454          .DRAST  DX,5,RXABRT      ;AST ENTRY POINT TABLE
```

Drop to fork level rather than device priority because the routine is lengthy and it needs all the registers.

```
45 000464          .FORK   DXFBLK        ;REQUEST FORK LEVEL IMMEDIATELY
```

Load registers; if the transfer is successful, this routine dispatches to the appropriate section for this interrupt. The three possibilities are: the first interrupt occurred; a read operation completed; a write operation completed. (A seek operation is treated as a zero-length read.)

```
46 000472 012700      MOV     (PC)+,R0      ;GET A VERY USEFUL FLAG WORD
47 000474 000000  RXFUN2: .WORD  0      ; : READ OR WRITE COMMAND ON CORRECT UNIT
48 000476 012703      MOV     #128.,R3      ;LOAD A HANDY CONSTANT
                                000200
49 000502 012704      MOV     (PC)+,R4      ;GET ADDRESS OF RX CONTROLLER
50 000504 177170  RXCSA: .WORD  DX$CSR      ; : ADDRESS OF CONTROLLER
51 000506                                .ASSUME . LE DXSTRT+1000
52 000506 010405      MOV     R4,R5      ;POINT R5 TO RX DATA BUFFER
53 000510 005725      TST     (R5)+      ;CHECK FOR ERROR, R5 -> DX REGISTER WITH ERROR
54 000512 100523      BMI     RXRTRY      ;ERROR, PROCESS IT
55 000514 006327      ASL     (PC)+      ;NO ERROR, DISPATCH AFTER INTERRUPT
56 000516 000000  RXIRTN: .WORD  0      ;OFFSET TO INTERRUPT CONTINUATION
57 000520 103704      BCS     1$      ;FIRST INTERRUPT, START I/O
58 000522 032700      BIT     #CSREAD,R0      ;READ OR WRITE?
                                000002
59 000526 001442      BEQ     10$      ;WRITE, DON'T EMPTY SILO
60 000530 005700      TST     R0      ;READ, IS THIS A SPECIAL FUNCTION?
```

The silo is a 128-byte (decimal) storage area in the diskette logic.

```
61 000532 100033      BPL     9$      ;NO, SIMPLY EMPTY THE SILO THAT WAS JUST READ
62 000534 032715      BIT     #ESDD,@R5      ;IF SPFUN READ, IS DELETED DATA FLAG PRESENT?
                                000100
63 000540 001430      BEQ     9$      ;NOPE, JUST EMPTY THE SILO
64
```

This routine puts a 1 in the first word of the user buffer if a deleted data mark was present on a special function read operation.

```
65          .IF EQ  MMG$T
66          MOV     BUFRAD,R2      ;GET ADDRESS OF USER BUFFER AREA
67          INC     -(R2)      ;SET FLAG WORD TO 1 TO INDICATE DELETED DATA
68          .IFF
69 000542 010401      MOV     R4,R1      ;SAVE R4
70 000544 016704      MOV     DXCQE,R4      ;POINT TO QUEUE ELEMENT
                                177240
71 000550 012746      MOV     #1,-(SP)      ;STACK A 1 TO PUT INTO FLAG WORD
                                000001
72 000554 162764      SUB     #2,Q$BUFF(R4)      ;MOVE BUFFER POINTER BACK TO FIRST WORD.
                                000002
                                000004
73 000562 026427      CMP     Q$BUFF(R4),#20000 ;POINTER OUT OF THIS PAR'S RANGE?
                                000004
                                020000
74 000570 103011      BHS     85$      ;NOPE...
75 000572 062764      ADD     #20000,Q$BUFF(R4) ;YES, GET IT BACK IN RANGE
                                020000
                                000004
76 000600 162764      SUB     #200,Q$PAR(R4) ; IN THE PREVIOUS PAR
                                000200
                                000012
77 000606 162764      SUB     #200,Q$MEM(R4) ; IN THE PREVIOUS PAR
```

```

000200
000014
78 000614 004777 85$: CALL @$PTWRD ;STORE IN 1ST WORD. Q.BUFF IS AGAIN ORIGINAL+2
000306
79 000620 010104 MOV R1,R4 ;RESTORE R4.
80 .ENDC ;EQ MMG$T
81
82 000622 004067 9$: JSR R0,SILOFE ;FOR READ, MOVE THE DATA FROM SILO TO BUFFER
000156
83 000626 000003 .WORD CSEBUF!CSGO ; EMPTY BUFFER COMMAND
84 000630 111522 MOV @R5,(R2)+ ; MOVB TO BE PLACED IN LINE IN SILOFE
85 000632 011502 MOV @R5,R2 ; DATA SLUFFER TO BE USED FOR SHORT READ

```

This point marks the successful completion of one sector for a read or write operation. The next routine increments the pointers for the next interleaved sector.

```

86 000634 105267 10$: INCB SECTOR ;RETURN HERE AFTER WRITES. BUMP SECTOR NUMBER
177573
87 000640 001012 BNE 11$ ;NOT OFF END OF TRACK YET
88 000642 062767 ADD #-26.*400+1,TRACK ;RESET SECTOR, BUMP TO NEXT TRACK
163001
177562
89 000650 062767 ADD #6,TRKOFF ;BUMP TRACK OFFSET VALUE
000006
177522
90 000656 003403 BLE 11$ ;OK IF STILL IN RANGE -25:0
91 000660 162767 SUB #26.,TRKOFF ;RESET TO PROPER RANGE MOD 26
000032
177512

```

The following routine increments the buffer address by 128 bytes, and reduces the byte count by 128. If the operation is not complete, it transfers another sector.

```

92 000666 11$:
93 .IF EQ MMG$T
94 ADD R3,BUFRAD ;UPDATE BUFFER ADDRESS
95 .IFF
96 000666 062767 ADD #2,PARVAL ;CHANGE MAP TO BUMP ADDRESS FOR NEXT TIME
000002
000164
97 .ENDC ;EQ MMG$T
98
99 000674 160367 SUB R3,BYTCNT ;REDUCE THE AMOUNT LEFT TO TRANSFER
177406
100 000700 101214 BHI 1$ ;LOOP IF WE ARE NOT DONE

```

The transfer is done. The routine sets the byte count to 0, and goes to 12\$ if this was a read or a special function operation.

```

101 000702 005067 CLR BYTCNT ;FIX BYTE COUNT SO THAT WRITES ARE ALL 0-FILLS
177400
102 000706 032700 BIT #CSREAD!SPFUNC,R0 ;READ OR SPECIAL FUNCTION OPERATION?
100002
103 000712 001004 BNE 12$ ;IF SO, NO ZERO-FILLING, SO WE'RE DONE

```

The operation was a write. The routine may need to be zero-filled up to three sectors (see FILLCT above).

```

104 000714 062727 ADD #040000,(PC)+ ;CHECK ORIGINAL WORD COUNT FOR # OF SECTORS
040000
105 000720 000 .BYTE 0 ; FILLER
106 000721 000 FILLCT: .BYTE 0 ; : ORIGINAL WORD COUNT LOW BYTE IN HIGH BYTE
107 000722 103206 BCC 2$ ;YES, LOOP FOR ZERO-FILLING ON WRITE
108 000724 12$: ;AHH, A SUCCESSFUL TRANSFER IS DONE
109 .IF NE ERL$G

```

Log a successful transfer:

```

110                                TST     SCSFLG           ;LOGGING SUCCESSFUL TRANSFERS?
111                                BNE     13$              ;NOPE...
112                                MOV     #DX$COD*400+377,R4 ;SET UP R4 = ID/-1
113                                MOV     DXCQE,R5        ; AND R5 -> CURRENT QUEUE ELEMENT
114                                CALL    @$ELPTR         ;CALL ERROR LOGGER TO REPORT SUCCESS
115                                .ENDC ;EQ ERL$G
116
117 000724 005077 13$: CLR     @RXCSA           ;DISABLE FLOPPY INTERRUPTS
177554

```

I/O Completion Section

```

118 000730          14$: .DRFIN DX           ;GO TO I/O COMPLETION
119

```

The abort routine:

```

120                                ; ABORT TRANSFER
121
122 000746 012777 RXABRT: MOV     #CSINIT,@RXCSA ;PERFORM AN RX11 INITIALIZE
040000
177530
123 000754 005067          CLR     DXFBLK+2       ;CLEAR FORK BLOCK TO AVOID A DISPATCH
000130

```

Go to .DRFIN if no error:

```

124 000760 000763          BR     14$           ; AND FINISH UP THIS I/O
125

```

If error logging was built:

```

126                                .DSABL LSB
127
128                                ; TRANSFER ERROR HANDLING
129
130 000762          RXRTRY:
131                                .IF NE ERL$G
132                                .ADDR  #DXRBUF,R3        ;R3 -> LOCATION TO STORE REGISTER INFO.
133                                MOV     R3,R2           ;SAVE IN R2 FOR LATER
134                                MOV     @R4,(R3)+       ;STORE RXCS
135                                MOV     @R5,(R3)+       ;STORE STATUS RXES
136                                MOV     #CSMAIN!CSGO,@R4 ;READ ERROR REGISTER (NO INTERRUPTS)
137                                1$: BIT     #CSDONE,@R4 ;WAIT FOR READ COMPLETION
138                                BEQ     1$
139                                MOV     @R5,@R3         ;STORE IN BUFFER
140                                MOV     DRETRY,R3
141                                SWAB   R3
142                                ADD     #DXNREG,R3      ;R3 = MAX RETRIES/# OF REGS
143                                MOV     #DX$COD*400,R4   ;R4 = DEVICE ID IN HIGH BYTE
144                                BISB   RXTRY,R4        ; AND CURRENT RETRY COUNT IN LOW BYTE
145                                DECB   R4             ; -1 FOR THIS ERROR
146                                MOV     DXCQE,R5        ;R5 -> QUEUE ELEMENT
147                                CALL    @$ELPTR         ;CALL ERROR LOGGER
148                                MOV     RXCSA,R4        ;RESTORE R4 = RXCS ADDRESS
149                                .ENDC ;NE ERL$G
150

```

See if a retry is allowed:

```

151 000762 005367          DEC     RXTRY           ;SHOULD WE TRY AGAIN?
177050
152 000766 003002          BGT     2$              ;YES
153 000770 000167          JMP     RXERR         ;NOPE, REPORT AN ERROR
177446
154
155 000774 012714 2$:     MOV     #CSINIT,@R4        ;START A RECALIBRATE
040000

```

Retry the operation:

```

156 001000 000167      JMP      RXINIT          ;EXIT THROUGH START OPERATION CODE
      177312
1      .SBTTL  SILOFE - FILL OR EMPTY THE SILO
2      ;+
3      ; SILOFE - FILL OR EMPTY THE SILO, DUMPING OR ZERO-FILLING IF NEEDED
4      ;
5      ;      R3 = 128.
6      ;      R4 -> FLOPPY CSR
7      ;      JSR      R0,SILOFE
8      ;      COMMAND: CSFBUF!CSGO FOR FILL (WRITE)
9      ;      CSEBUF!CSGO FOR EMPTY (READ)
10     ;      FILL/EMPTY INSTRUCTION: (R2 -> USER BUFFER, R5 -> RXDB)
11     ;      MOV B (R2)+,@R5 FOR FILL (WRITE)
12     ;      MOV B @R5,(R2)+ FOR EMPTY (READ)
13     ;      SLUFF INSTRUCTION: (R1 = 0, R5 -> RXDB)
14     ;      CLRB @R5      FOR FILL (WRITE)
15     ;      MOV B @R5,R2 FOR EMPTY (READ)
16     ;      R1 = RANDOM
17     ;      R2 = RANDOM
18     ;
19     ; NOTE: 1. THIS ROUTINE ASSUMES ERROR CAN NOT COME UP DURING A FILL OR EMPTY!!
20     ;      2. SEEK DOES A SILO EMPTY, A TIME WASTER
21     ;-
22     .ENABL  LSB

```

The diskette deals only in units of 128 decimal bytes. If a request to read is for fewer than 128 bytes, the handler reads 128 bytes and sloughs the extra bytes. If a request to write is for fewer than 128 bytes, the handler zero-fills to reach 128 bytes.

```

23 001004 012014 SILOFE: MOV      (R0)+,@R4      ;INITIATE FILL OR EMPTY BUFFER COMMAND
24 001006 012067      MOV      (R0)+,3$      ;PUT CORRECT MOV INSTRUCTION IN FOR FILL/EMPTY
      000036
25 001012 012067      MOV      (R0)+,5$      ;PUT IN INSTRUCTION TO SLUFF DATA
      000052
26 001016 016701      MOV      BYTCNT,R1      ;GET BYTE COUNT
      177264
27 001022 001417      BEQ      4$      ;IF ZERO, WE ARE SEEKING OR ZERO FILLING
28 001024 020103      CMP      R1,R3      ;IS THE BYTE COUNT <= 128?
29 001026 101401      BLOS    1$      ;OK IF SO
30 001030 010301      MOV      R3,R1      ;DO ONLY 128 BYTES AT A TIME
31 001032 016702 1$:  MOV      BUFPRD,R2      ;GET USER VIRTUAL BUFFER ADDRESS IN R2
      177240

```

The following section of code can be executed in two different ways. If the handler is assembled for an unmapped monitor, the code between the symbols 2\$ and PARVAL is simply executed in-line. If the handler is assembled for a mapped monitor, the JSR to PIEXT and the word PARVAL are included. In this situation, the routine P1EXT copies the code between 2\$ and PARVAL to the monitor stack, uses the value passed in PARVAL to map to the user buffer, and executes the code from the monitor stack. This is done to ensure that the code is not in the PAR1 area when it is executed, since PAR1 is used to map to the user buffer.

```

32      .IF NE  MMG$T
33 001036 004077      JSR      R0,@$PIEXT      ;Let the monitor execute the following code.
      176766
34 001042 000016      .WORD    PARVAL-.      ;Number of instructions in bytes plus 2.
35      .ENDC ;NE MMG$T
36 001044 105714 2$:  TSTB    @R4      ;**EXT** TRY FOR THE TRDY
37 001046 100376      BPL     2$      ;**EXT** TRANSFER READY
38 001050 000000 3$:  HALT     ;**EXT** INSTRUCTION TO MOV OR SLUFF DATA FROM
39 001052 105714      TSTB    @R4      ;**EXT** TOUCH THE CSR TO GET IT READY
40 001054 105301      DECB    R1      ;**EXT** CHECK FOR COUNT DONE
41 001056 001372      BNE     2$      ;**EXT** STILL MORE TO TRANSFER
42      .IF NE  MMG$T
43 001060 000000      PARVAL: .WORD    0      ;using this value for the PAR 1 bias.
44      .ENDC ;NE MMG$T

```

The slough routine:

```
45 001062 105714 4$:      TSTB   @R4           ;WAIT FOR TRANSFER READY OR TRANSFER DONE
46 001064 003003          BGT    6$           ;TDNE UP WITH NO TRDY, SO ALL DONE
47 001066 001775          BEQ    4$           ;LOOP
48 001070 000000 5$:      HALT                    ;TRANSFER READY, SO SLUFF DATA
49 001072 000773          BR     4$           ;LOOP TO SLUFF MORE
50 001074 000200 6$:      RTS     R0           ;RETURN
51                          .DSABL  LSB

1                          .SBTTL  TABLES, FORK BLOCK, END OF DRIVER
2
3                          ; CHANGES TO CSR CODE FOR SPECIAL FUNCTIONS
4
5 001076 100006          .WORD  CSWRD-CSRD+SPFUNC      ;375: READ+GO -> WRITE DELETED+GO
6 001100 077776          .WORD  CSWRD-CSRD+SPFUNC      ;376: READ+GO -> WRITE+GO
7 001102 100000          .WORD  CSRD-CSRD+SPFUNC      ;377: READ+GO -> READ+GO
8 001104 000000  CHGTBL: .WORD  0           ; READ/WRITE STAY THE SAME
9
10 001106 000000  DXFBLK: .WORD  0,0,0,0      ;DX FORK QUEUE ELEMENT
    001110 000000
    001112 000000
    001114 000000
11
12                          .IF NE  ERL$G
13  DXRBUF: .BLKW  DXNREG          ;ERROR LOG STORAGE
14                          .ENDC ;NE ERL$G
```

Bootstrap driver

```
1                          .SBTTL  BOOTSTRAP DRIVER
2
```

The .DRBOT macro:

```
3 001116                  .DRBOT  DX,BOOT1,READ
```

Termination Section

The .DREND macro generated by .DRBOT (the macro expansion):

```
001116                  .DREND  DX,0,
    .IF B <>
001116                  .PSECT  DXDVR
    .IFF
    .PSECT
    .ENDC
    .IF NDF DX$END,DX$END::
    .IF EQ  .-DX$END
    .IF NE MMG$T!<0&2.>
001116 000000 $RLPTR: .WORD  0
001120 000000 $MPPTR: .WORD  0
001122 000000 $GTBYT: .WORD  0
001124 000000 $PTBYT: .WORD  0
001126 000000 $PTWRD: .WORD  0
    .ENDC
    .IF NE ERL$G!<0&1>
    $ELPTR: .WORD  0
    .ENDC
    .IF NE TIM$IT!<0&4.>
    $TIMIT: .WORD  0
    .ENDC
001130 000000 $INPTR: .WORD  0
001132 000000 $FKPTR: .WORD  0
    .IF NDF ...V22 ...V22=0
    .IF NE ...V22&^o40000
    DX$X64 =:
    .REPT  16.
    .WORD  0
    .ENDR
    .ENDC
    .GLOBL  DXSTRT
```

The following line marks the end of the loadable portion of the handler. It is used to determine the handler's length.

```

001134' DXEND==.
.IFF
.PSECT DXBOOT
.IIF LT <DXBOOT-.^o664>,.ERROR;?SYSMAC-E-Primary boot too large;
.=DXBOOT+^o664
BIOERR: JSR R1,REPORT
.WORD IOERR-DXBOOT
REPORT: MOV #BOOTF-DXBOOT,R0
MOV #30002$-DXBOOT,R2
CALL @R2
MOV @R1,R0
CALL @R2
MOV #CRLFLF-DXBOOT,R0
CALL @R2
30001$: HALT
BR 30001$
30002$: TSTB @#TPS
BPL 30002$
MOVB (R0)+,@#TPB
BNE 30002$
RETURN
BOOTF: .ASCIZ <CR><LF>"?BOOT-U-"
IOERR: .ASCII "I/O error"
CRLFLF: .ASCIZ <CR><LF><LF>
.EVEN
.IIF NDF ...V7,...V7=-1
.REPT 4.
.WORD ...V7
.ENDR
DXBEND::
.ENDC
.IIF NDF TPS,TPS:^o177564
.IIF NDF TPB,TPB:^o177566
000012 LF:^o12
000015 CR:^o15
001000 B$BOOT:^o1000
004716 B$DEVN:^o4716
004722 B$DEVU:^o4722
004730 B$READ:^o4730
.IF NDF B$DNAM
.IF EQ MMG$T
B$DNAM:^RDY
.IFF
B$DNAM:^RDXX
.ENDC ; EQ MMG$T
.ENDC ; NDF B$DNAM
000062 .ASECT
.=^o62
000062 000000' .WORD DXBOOT,DXBEND-DXBOOT,READ-DXBOOT
000064 001000
000066 000224
000000 .PSECT DXBOOT
000000 000240 DXBOOT: NOP
000002 000413 BR BOOT1-2.
000100 ...V2:^o100
.IRP X <UBUS,QBUS>
...V3=0
.IIF IDN <X> <UBUS> ...V3=1.
.IIF IDN <X> <QBUS> ...V3=2.
.IIF IDN <X> <CBUS> ...V3=4.
.IIF IDN <X> <UMSCP> ...V3:^o10
.IIF IDN <X> <QMSCP> ...V3:^o20
.IIF IDN <X> <CMSCP> ...V3:^o40
.IIF EQ ...V3 .ERROR;?SYSMAC-E-Invalid C O N T R O L, found - UBUS,QBUS;
...V2=...V2!...V3
.ENDR
000000 ...V3=0
000001 .IIF IDN <UBUS> <UBUS> ...V3=1.
.IIF IDN <UBUS> <QBUS> ...V3=2.
.IIF IDN <UBUS> <CBUS> ...V3=4.
.IIF IDN <UBUS> <UMSCP> ...V3:^o10
.IIF IDN <UBUS> <QMSCP> ...V3:^o20
.IIF IDN <UBUS> <CMSCP> ...V3:^o40

```

```

.IIF EQ ...V3 .ERROR;?SYSMAC-E-Invalid C O N T R O L, found - UBUS,QBUS;
000101 ...V2=...V2!...V3
000000 ...V3=0
.IIF IDN <QBUS> <UBUS> ...V3=1.
000002 .IIF IDN <QBUS> <QBUS> ...V3=2.
.IIF IDN <QBUS> <CBUS> ...V3=4.
.IIF IDN <QBUS> <UMSCP> ...V3=^o10
.IIF IDN <QBUS> <QMSCP> ...V3=^o20
.IIF IDN <QBUS> <CMSCP> ...V3=^o40
.IIF EQ ...V3 .ERROR;?SYSMAC-E-Invalid C O N T R O L, found - UBUS,QBUS;
000103 ...V2=...V2!...V3
000026' .=BOOT1-6.
000026 020 .BYTE ^o20,...V2,^o20,^o^C<20+...V2+20>
000027 103
000030 020
000031 234
.IF EQ <1-1>
000032 000400 BR BOOT1
.IFF
.IF EQ <1-2.>
BMI BOOT1
.IFF
.ERROR;?SYSMAC-E-Invalid S I D E S, expecting 1/2, found - 1;
.ENDC
.ENDC
4
5 000014' . = DXBOOT+14
6 000014 000120 .WORD READS-DXBOOT
7 000016 000340 .WORD 340
8 000020 000070 .WORD WAIT-DXBOOT
9 000022 000340 .WORD 340
10

```

Locations 34 through 52 are reserved for Digital.

```

11 000034' . = DXBOOT+34 ;34-52 USEABLE
12 000034 116067 BOOT1: MOVB UNITRD-DXBOOT(R0),RDCMD ;SET READ FUNCTION FOR CORRECT UNIT
000056
000066
13 000042 011706 REETRY: MOV @PC,SP ;INIT SP WITH NEXT INSTRUCTION
14 000044 012702 MOV #200,R2 ;AREA TO READ IN NEXT PART OF BOOT
000200
15 000050 005000 CLR R0 ;SET TRACK NUMBER
16 000052 000446 BR B2$ ;OUT OF ROOM HERE, GO TO CONTINUATION
17
18 000056' . = DXBOOT+56
19 000056 007 UNITRD: .BYTE CSGO+CSRD ;READ FROM UNIT 0, SETS WEIRD BUT OK PS
20 000057 027 .BYTE CSGO+CSRD+CSUNIT;READ FROM UNIT 1
21
22 000070' . = DXBOOT+70 ;PAPER TAPE VECTORS
23 000070 005714 WAIT: TST @R4 ;IS TR, ERR, DONE UP? INT ENB CAN'T BE
24 000072 001776 BEQ WAIT ;LOOP TILL SOMETHING
25 000074 100762 BMI REETRY ;START AGAIN IF ERROR
26 000076 000002 RTIRET: RTI ;RETURN
27
28 000120' . = DXBOOT+120
29 000120 012704 READS: MOV (PC)+,R4 ;R4 -> RX STATUS REGISTER
30 000122 177170 BOTCSR: .WORD DX$CSR
31 000124 010405 MOV R4,R5 ;R5 WILL POINT TO RX DATA BUFFER
32 000126 012725 MOV (PC)+,(R5)+ ;INITIATE READ FUNCTION
33 000130 000000 RDCMD: .WORD 0 ;GETS FILLED WITH READ COMMAND
34 000132 000004 IOT ;CALL WAIT SUBROUTINE
35 000134 010315 MOV R3,@R5 ;LOAD SECTOR NUMBER INTO RXDB
36 000136 000004 IOT ;CALL WAIT SUBROUTINE
37 000140 010015 MOV R0,@R5 ;LOAD TRACK NUMBER INTO RXDB
38 000142 000004 IOT ;CALL WAIT SUBROUTINE
39 000144 012714 MOV #CSGO+CSEBUF,@R4;LOAD EMPTY BUFFER FUNCTION INTO RXCS
000003
40 000220 BROFFS = READF-. ;USE FOR COMPUTING BR OFFSET
41 000150 000004 RDX: IOT ;CALL WAIT SUBROUTINE
42 000152 105714 TSTB @R4 ;IS TRANSFER READY UP?
43 000154 100350 BPL RTIRET ;BRANCH IF NOT, SECTOR MUST BE LOADED
44 000156 111522 MOV @R5,(R2)+ ;MOVE DATA BYTE TO MEMORY
45 000160 005301 DEC R1 ;CHECK BYTE COUNT
46 000162 003372 BGT RDX ;LOOP AS LONG AS WORD COUNT NOT UP
47 000164 005002 CLR R2 ;KLUDGE TO SLUFF BUFFER IF SHORT WD CNT

```

```

48 000166 000770      BR      RDX          ;LOOP
49
50 000170 010601 B2$:  MOV      SP,R1      ;SET TO BIG WORD COUNT
51 000172 005200      INC      R0          ;SET TO ABSOLUTE TRACK 1
52 000174 011703      MOV      @PC,R3     ;ABSOLUTE SECTOR 3 FOR NEXT PART
53 000176      .ASSUME BPT EQ 3

      .IF EQ <<BPT>>-<<3>>
      .IFF
      .IF B <>
      .ERROR;?SYSMAC-W-"BPT EQ 3" is not true;
      .IFF
      .ERROR ;?SYSMAC-;
      .ENDC
      .ENDC

54 000176 000003      BPT          ;CALL READS SUBROUTINE
55      ;SECTOR 2 OF RX BOOT
56 000200 122323 BOOT2:  CMPB      (R3)+,(R3)+ ;BUMP TO SECTOR 5
57 000202 000003      BPT          ;CALL READS SUBROUTINE
58 000204 122323      CMPB      (R3)+,(R3)+ ;BUMP TO SECTOR 7
59 000206 000003      BPT          ;CALL READS SUBROUTINE
60 000210 032767      BIT      #CSUNIT,RDCMD ;CHECK UNIT ID
      000020
      177712
61 000216 001173      BNE      BOOT      ;BRANCH IF BOOTING UNIT 1, R0=1
62 000220 005000      CLR      R0          ;SET TO UNIT 0
63 000222 000571      BR      BOOT      ;NOW WE ARE READY TO DO THE REAL BOOT
64
65 000224 012737 READ:  MOV      (PC)+,@(PC)+ ;MODIFY READ ROUTINE
66 000226 000167      .WORD    167
67 000230 000150      .WORD    RDX-DXBOOT
68 000232 012737      MOV      (PC)+,@(PC)+
69 000234 000214      .WORD    READF-RDX-4
70 000236 000152      .WORD    RDX-DXBOOT+2
71 000240 012737      MOV      #READ1-DXBOOT,@#B$READ ;CALLS TO B$READ WILL GO TO READ1
      000300
      004730
72 000246 012737      MOV      #TRWAIT-DXBOOT,@#20 ;LETS HANDLE ERRORS DIFFERENTLY
      000416
      000020
73 000254 005037      CLR      @#JSW      ;CLEAR JSW SINCE THE DX BOOT IN SYSCOM AREA
      000044
74 000260 005767      TST      HRDBOT     ;DID WE REACH HERE VIA A HARDWARE BOOT?
      000346
75 000264 001405      BEQ      READ1     ;YES, DON'T SET UP UNIT NUMBER
76 000266 013703      MOV      @#B$DEVU,R3 ;NO, SET UP UNIT NUMBER
      004722
77 000272 116367      MOVB     UNITRD-DXBOOT(R3),RDCMD ;STORE UNIT NUMBER
      000056
      177630
78 000300 006300 READ1:  ASL      R0          ;CONVERT BLOCK TO LOGICAL SECTOR
79 000302 006300      ASL      R0          ;LSN=BLOCK*4
80 000304 006301      ASL      R1          ;MAKE WORD COUNT BYTE COUNT
81 000306 010046      1$:  MOV      R0,-(SP)    ;SAVE LSN FOR LATER
82 000310 010003      MOV      R0,R3      ;WE NEED 2 COPIES OF LSN FOR MAPPER
83 000312 010004      MOV      R0,R4
84 000314 005000      CLR      R0          ;INIT FOR TRACK QUOTIENT
85 000316 000402      BR      3$         ;JUMP INTO DIVIDE LOOP
86
87 000320 162703      2$:  SUB      #23.,R3     ;PERFORM MAGIC TRACK DISPLACEMENT
      000027
88 000324 005200      3$:  INC      R0          ;BUMP QUOTIENT, STARTS AT TRACK 1
89 000326 162704      SUB      #26.,R4     ;TRACK=INTEGER(LSN/26)
      000032
90 000332 100372      BPL      2$         ;LOOP - R4=REM(LSN/26)-26
91 000334 022704      CMP      #-14.,R4    ;SET C IF SECTOR MAPS TO 1-13
      177762
92 000340 006103      ROL      R3          ;PERFORM 2:1 INTERLEAVE
93 000342 162703      4$:  SUB      #26.,R3     ;ADJUST SECTOR INTO RANGE -1,-26
      000032
94 000346 100375      BPL      4$         ;(DIVIDE FOR REMAINDER ONLY)
95 000350 062703      ADD      #27.,R3     ;NOW PUT SECTOR INTO RANGE 1-26
      000033
96 000354 000003      BPT          ;CALL READS SUBROUTINE
97 000356 012600      MOV      (SP)+,R0    ;GET THE LSN AGAIN
98 000360 005200      INC      R0          ;SET UP FOR NEXT LSN
99 000362 005701      TST      R1          ;WHAT'S LEFT IN THE WORD COUNT
100 000364 003350      BGT      1$         ;BRANCH TO TRANSFER ANOTHER SECTOR
101 000366 000207      RETURN

```

```

102
103 000370 005714 READF: TST @R4 ;ERROR, DONE, OR TR UP?
104 000372 001776 BEQ READF ;BR IF NOT
105 000374 100533 BMI BIOERR ;BR IF ERROR
106 000376 105714 TSTB @R4 ;TR OR DONE?
107 000400 100011 BPL READFX ;BR IF DONE
108 000402 111522 MOVB @R5,(R2)+ ;MOVE DATA BYTE TO MEMORY
109 000404 005301 DEC R1 ;CHECK BYTE COUNT
110 000406 003370 BGT READF ;LOOP IF MORE
111 000410 012702 MOV #1,R2 ;SLUFF BUFFER IF SHORT WD CNT
    000001
112 ;DON'T DESTROY LOC 0
113 000414 000765 BR READF ;LOOP
114
115 000416 005714 TRWAIT: TST @R4 ;ERROR, DONE, OR TR UP?
116 000420 100521 BMI BIOERR ;HARD HALT ON ERROR
117 000422 001775 BEQ TRWAIT ;BR IF NOT
118 000424 000002 READFX: RTI
119
120 000606' . = DXBOOT+606
121 000606 012706 BOOT: MOV #10000,SP ;SET STACK POINTER
    010000
122 000612 010046 MOV R0,-(SP) ;SAVE THE UNIT NUMBER
123 000614 012700 MOV #2,R0 ;READ IN SECOND PART OF BOOT
    000002
124 000620 012701 MOV #<4*400>,R1 ;EVERY BLOCK BUT THE ONE WE ARE IN
    002000
125 000624 012702 MOV #1000,R2 ;INTO LOCATION 1000
    001000
126 000630 005027 CLR (PC)+ ;CLEAR TO SHOW HARDWARE BOOT
127 000632 000001 HRDBOT: .WORD 1 ;INITIALLY SET TO 1
128 000634 004767 CALL READ ;GO READ IT IN
    177364
129 000640 012737 MOV #READ1-DXBOOT,@#B$READ ;STORE START LOCATION FOR READ ROUTINE
    000300
    004730
130 000646 012737 MOV #B$DNAM,@#B$DEVN ;STORE RAD50 DEVICE NAME
    016330
    004716
131 000654 012637 MOV (SP)+,@#B$DEVU ;STORE THE UNIT NUMBER
    004722
132 000660 000137 JMP @#B$BOOT ;START SECONDARY BOOT
    001000
133
134 000664 .DREND DX
    001134 .IF B <>
    .PSECT DXDVR
    .IFF
    .PSECT
    .ENDC
    .IF NDF DX$END,DX$END::
    .IF EQ .-DX$END
    .IF NE MMG$T!<0&2.>
    $RLPTR: .WORD 0
    $MPPTR: .WORD 0
    $GTBYT: .WORD 0
    $PTBYT: .WORD 0
    $PTWRD: .WORD 0
    .ENDC
    .IF NE ERL$G!<0&1>
    $ELPTR: .WORD 0
    .ENDC
    .IF NE TIM$IT!<0&4.>
    $TIMIT: .WORD 0
    .ENDC
    $INPTR: .WORD 0
    $FKPTR: .WORD 0
    .IF NDF ...V22 ...V22=0
    .IF NE ...V22&^o40000
    DX$X64 =: .
    .REPT 16.
    .WORD 0
    .ENDR
    .ENDC
    .GLOBL DXSTRT
    DXEND==.
    .IFF
    000664 .PSECT DXBOOT

```

```

.IIF LT <DXBOOT-.+^o664>,.ERROR: ?SYSMAC-E-Primary boot too large;
.=DXBOOT+^o664
000664 004167 BIOERR: JSR    R1,REPORT
000002
000670 000753          .WORD  IOERR-DXBOOT
000672 012700 REPORT: MOV    #BOOTF-DXBOOT,R0
000740
000676 012702          MOV    #30004$-DXBOOT,R2
000722
000702 004712          CALL   @R2
000704 011100          MOV    @R1,R0
000706 004712          CALL   @R2
000710 012700          MOV    #CRLFLF-DXBOOT,R0
000764
000714 004712          CALL   @R2
000716 000000 30003$: HALT
000720 000776          BR     30003$
000722 105737 30004$: TSTB   @#TPS
177564
000726 100375          BPL    30004$
000730 112037          MOVB  (R0)+,@#TPB
177566
000734 001372          BNE    30004$
000736 000207          RETURN
000740 015          BOOTF:  .ASCIZ  <CR><LF>" ?BOOT-U-"
000741 012
000742 077
000743 102
000744 117
000745 117
000746 124
000747 055
000750 125
000751 055
000752 000
000753 111 IOERR:  .ASCII  "I/O error"
000754 057
000755 117
000756 040
000757 145
000760 162
000761 162
000762 157
000763 162
000764 015 CRLFLF: .ASCIZ  <CR><LF><LF>
000765 012
000766 012
000767 000

          .EVEN
.IIF NDF ...V7,...V7=-1
000004 .REPT  4.
          .WORD  ...V7
          .ENDR
000770 177777          .WORD  ...V7
000772 177777          .WORD  ...V7
000774 177777          .WORD  ...V7
000776 177777          .WORD  ...V7
001000 DXBEND: :
          .ENDC

135
136 000001          .END

```

Symbol table

ABTIO\$	001000		DVC.VT	000015		O.GOOD	000554	
BAREA	000256		DVM.DM	000002		O.RTRY	000620	
BIOERR	000664R	003	DVM.DX	000001		O.SYWL	000562	
BOOT	000606R	003	DVM.NF	000200		O.VEC	000600	
BOOTF	000740R	003	DVM.NS	000001		O.WP	000634	
BOOT1	000034R	003	DV2.V2	040000		O.WPF	000302	
BOOT2	000200R	003	DXBEND	001000RG	003	PARVAL	001060R	002
BOTCSR	000122R	003	DXBOOT	000000RG	003	P1EXT	000432	
BROFFS=	000220		DXCQE	000010RG	002	Q\$BLKN	000000	
BTCR\$ =	002256		DXDSIZ	000756		Q\$BUFF	000004	
BUFRAD	000276R	002	DXEND =	001134RG	002	Q\$COMP	000010	
BYTCNT	000306R	002	DXENT	000020R	002	Q\$CSW	177776	
B\$BOOT	001000		DXFLBK	001106R	002	Q\$FUNC	000002	
B\$DEVN	004716		DXINT	000456RG	002	Q\$JNUM	000003	
B\$DEVU	004722		DXLQE	000006RG	002	Q\$LINK	177774	
B\$DNAM	016330		DXNREG	000003		Q\$MEM	000014	
B\$READ	004730		DXSTRT	000000RG	002	Q\$PAR	000012	
B2\$	000170R	003	DXSTS	102022		Q\$UNIT	000003	
CHGTBL	001104R	002	DXSYS	000006RG	002	Q\$WCNT	000006	
CR	000015		DXT\$O =	000000		Q.BLKN	000004	
CRLFLF	000764R	003	DXWPRO	000016R	002	Q.BUFF	000010	
CSDONE	000040		DXW1	000104R	002	Q.COMP	000014	
CSEBUF	000002		DX\$COD	000022		Q.CSW	000002	
CSERR	100000		DX\$CSR=	177170 G		Q.ELGH	000024	
CSFBUF	000000		DX\$CS2=	177174 G		Q.FUNC	000006	
CSGO	000001		DX\$END	001116RG	002	Q.JNUM	000007	
CSINIT	040000		DX\$NAM=	016300		Q.LINK	000000	
CSINT	000100		DX\$VC2=	000270 G		Q.MEM	000020	
CSMAIN	000016		DX\$VEC=	000264 G		Q.PAR	000016	
CSRD	000006		EOF\$	020000		Q.UNIT	000007	
CSRST	000012		ERL\$G =	000000		Q.WCNT	000012	
CSREAD	000002		ESCRC	000001		RDCMD	000130R	003
CSRX02	004000		ESDD	000100		RDX	000150R	003
CSTR	000200		ESDRY =	000200 G		READ	000224R	003
CSUNIT	000020		ESID	000004		READF	000370R	003
CSWRT	000004		ESPAR	000002		READFX	000424R	003
CSWRD	000014		FILLCT	000721R	002	READS	000120R	003
DAREA	000244		FILST\$	100000		READ1	000300R	003
DEVNAM	000254		FINDRV	000214		REENTRY	000042R	003
DISCSR	000174		HDERR\$	000001		REPORT	000672R	003
DRETRY	000034R	002	HNDLR\$	004000		RETRY	000010	
DVC.CT	000006		HRDBOT	000632R	003	RONLY\$	040000	
DVC.DE	000010		HS2.BI	000001		RTE\$M =	000000	
DVC.DK	000004		HS2.KI	000002		RTIRET	000076R	003
DVC.DL	000012		HS2.KL	000004		RXABRT	000746R	002
DVC.DP	000011		HS2.KU	000010		RXCSA	000504R	002
DVC.LP	000007		HS2.MO	000020		RXERR	000442R	002
DVC.MT	000005		INSCSR	000176		RXFUN2	000474R	002
DVC.NI	000013		INSDAT	000200		RXIENB	000434R	002
DVC.NL	000001		INSSYS	000202		RXINIT	000316R	002
DVC.PS	000014		IOERR	000753R	003	RXIRTN	000516R	002
DVC.SB	000020		JSW	000044		RXRTRY	000762R	002
DVC.SI	000016		LF	000012		RXTRY	000036R	002
DVC.SO	000017		MMG\$T =	000001		SECTOR	000433R	002
DVC.TP	000003		N.WP	000640		SILOFE	001004R	002
DVC.TT	000002		O.BAD	000556		SPECL\$	010000	
DVC.UK	000000		O.CSR	000442		SPFUNC	100000	
SPFUN\$	002000		\$MPPTR	001120RG	002	...V15=	000340	
SYSCHN	000017		\$PTBYT	001124RG	002	...V16=	000000	
SYSPTR	000054		\$PTWRD	001126RG	002	...V17=	000000	
TIM\$IT=	000000		\$P1EXT	000030R	002	...V18=	000001	
TPB	177566		\$RLPTR	001116RG	002	...V19=	000000	
TPS	177564		.AUDIT	107123 G		...V2 =	000103	
TRACK	000432R	002	.DSTAT	000342		...V20=	000000	
TRKOFF	000400R	002	.DX	000021 G		...V21=	000000	
TRWAIT	000416R	003	.READ	000375		...V22=	000000	
UNITRD	000056R	003	.WRITE	000375		...V27=	000000	
VARSZ\$	000400		..READ	000010		...V28=	000270	
WAIT	000070R	003	..WRIT	000011		...V3 =	000002	
WONLY\$	020000		...V10=	000040		...V4 =	000000	
X.WP	000270		...V11=	000370		...V5 =	000114	
\$FKPTR	001132RG	002	...V12=	000370		...V6 =	000270	
\$GTBYT	001122RG	002	...V13=	000000		...V7 =	177777	
\$INPTR	001130RG	002	...V14=	000000		...V9 =	000000	

```

. ABS. 000660 000 (RW,I,GBL,ABS,OVR)
      000000 001 (RW,I,LCL,REL,CON)
DXDVR 001134 002 (RW,I,LCL,REL,CON)
DXBOOT 001000 003 (RW,I,LCL,REL,CON)

```

Figure A-2: DL Disk Handler

In the interests of clarity, code from the DL handler that does not apply to PDP-11 processors has been removed. Further, the contents of some of the macro expansions has been removed when those contents served no instructive purpose. In both cases, the removed lines are indicated by ellipses.

```
DL - RL01/RL02 Disk Handler      MACRO V05.05  Thursday 28-Feb-91 15:01
```

Table of contents

```

CONDITIONAL ASSEMBLY SUMMARY
MACROS AND DEFINITIONS
*** THIS HANDLER SUPPORTS 2 UNITS ***
HANDLER MACROS
HARDWARE DEFINITIONS
INSTALLATION CODE
SET OPTIONS
REQUEST ENTRY POINT
INITIALIZE FOR TRANSFER, SET FUNCTION CODE, FIX WORD COUNT
COMPUTE DISK ADDRESS AND START TRANSFER
ENSURE THAT DISK IS ON TRACK BEFORE TRANSFER
DLXFER - START AN I/O TRANSFER
DLINT - INTERRUPT ENTRY POINT
HANDLE THE ERRORS
FINISH SUCCESSFUL OPERATION
GET DEVICE SIZE
DLXCT - FUNCTION EXECUTION ROUTINES
DLSQUE - SETUP PSEUDO QUEUE ELEMENT
DATA AREAS
BOOTSTRAP DRIVER
BOOTSTRAP READ ROUTINE
BOOTSTRAP CONTINUED
FETCH/LOAD CODE

```

Mapped monitor conditional:

```

1          000001  MMG$T = 1
           .MCALL .MODULE
2 000000    .MODULE DL,VERSION=42,COMMENT=<RL01/RL02 Disk Handler>,AUDIT=YES
3
4          ;
5          ;           COPYRIGHT 1989, 1990 BY
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7          ;           ALL RIGHTS RESERVED
8          ;
9          ;THIS SOFTWARE IS FURNISHED UNDER A LICENSE AND MAY BE USED AND COPIED
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18         ;CORPORATION.
19         ;
           ;DIGITAL ASSUMES NO RESPONSIBILITY FOR THE USE OR RELIABILITY OF ITS

```

Conditional Assembly Summary

```

1          .SBTTL  CONDITIONAL ASSEMBLY SUMMARY
2          ;+
3          ;COND
4          ;
5          ;
6          ;
7          ;
8          ;
9          ;      DL$UN  (2)          unit to support (additive only)
10         ;          1-4          valid range
11         ;
12         ;      EIS$I  (MMG$T)     use SOB instruction (no code effects!)
13         ;          0            simulate SOB
14         ;          2            use SOB
15         ;
16         ;      DL$CSR  (174400)    CSR
17         ;      DL$VEC  (160)      Vector

19         ;
20         ;      MMG$T          std conditional
21         ;      TIM$IT        std conditional (no code effects)
22         ;      ERL$G          std conditional
23         ;-

```

Preamble Section

Each macro you use in the handler requires the `.MCALL` statement, as line 6 shows. Note that `.DRDEF` issues many of the `.MCALL` statements for you so you need not explicitly call them.

Macros and Definitions

```

1          .SBTTL  MACROS AND DEFINITIONS
2
3
4          .ENABL  LC
5
6          .MCALL  .DRDEF, .MFPS, .MTPS, .ASSUME, .ADDR, .BR

```

A call is made to a macro (`.UBVDF`) in the system definition library `SYSTEM.MLB`. `SYSTEM.MLB` is always found on logical device `SRC`:

```

7
8          .LIBRARY "SRC:SYSTEM.MLB"
9          .MCALL  .UBVDF
10 000000  .UBVDF
11

```

Various monitor offsets and locations are defined with mnemonics so that references to them can be found easily:

```

12         ; VECTOR DEFINITIONS
13
14         000004      NXM.V  == 4          ;NON-EXISTENT MEMORY TRAP VECTOR
15         000020      IOT.V  == 20       ;IOT TRAP VECTOR
16
17         ;
18         ;
19         ;
20         ;
21         ;
22         ;
23         ;
24         ;
25         ;
26         ;
27         ;
28         ;
29
30         ; SYSTEM GENERATION OPTION
31
32         .IIF NDF DL$UN, DL$UN == 2      ;NUMBER OF UNITS SUPPORTED
33         .IIF GT DL$UN-4, DL$UN == 4    ;CAN'T HAVE MORE THAN 4 UNITS
34         .IIF LE DL$UN, DL$UN == 1     ;CAN'T HAVE NO UNITS
35
36         .IRP      X,<\DL$UN>

```

Handler Unit Support

```

37          .SBTTL  *** THIS HANDLER SUPPORTS X UNITS ***
38          .ENDR
39
40          ; SPECIAL FUNCTION DEFINITIONS
41          ; ALL SPECIAL FUNCTIONS ARE DMA EXCEPT FOR FN$SIZ AND FN$GET
42          ; FN$WRT AND FN$RED GO IN UBTAB. FN$REP USES A PERMANENT UMR
43
44          FN$GET =: 370          ;GET DEVICE STATUS
45          000373 FN$SIZ =: 373          ;GET DEVICE SIZE
46          000374 FN$REP =: 374          ;FORCE RE-READ OF REPLACEMENT TABLE

```

Use the replacement table with:

```

47          000376 FN$WRT =: 376          ;ABSOLUTE WRITE (NO BAD BLOCK)
48          000377 FN$RED =: 377          ;ABSOLUTE READ (REPLACEMENT)
49          ;NOTE: if you add a SPFUN code also add it to .DRSPF
50
51          ; ERROR LOGGING DEFINITIONS
52
53          000010 DLRCNT =: 8.          ;ERROR RETRY COUNT
54          000006 DLREG =: 6          ;REGISTERS TO LOG ON ERROR
55
56          ; RL11/RL01 PARAMETERS
57          ; GEOMETRY: 256 CYLINDERS (512 ON RL02)
58          ; 2 TRACKS PER CYLINDER
59          ; 20 BLOCKS PER TRACK
60          ; 2 128-WORD SECTORS PER BLOCK
61
62          000024 DLBPT =: 20.          ;NUMBER OF BLOCKS PER TRACK
63          012000 DLWPT =: 256.*DLBPT    ;WORDS PER TRACK
64          000012 DLNBAD =: 10.         ;NUMBER ALLOWABLE BAD BLOCKS PER DISK
65          023742 DLSIZE =: <256.*2-1>*DLBPT-DLNBAD ;BLOCKS PER RL01 (LESS BSF)
66          047742 DLSIZ2 =: <512.*2-1>*DLBPT-DLNBAD ;BLOCKS PER RL02 (LESS BSF)
67          000052 DLTSIZ =: DLNBAD*4.+2 ;SIZE OF BAD BLOCK TABLE
68          ; (PLUS END OF TABLE FENCE)
69
70          ; UB DEFINITIONS
71          ; FIXED OFFSETS EQUATES (.FIXDF)
72
73          000404 $PNPTR =: 000404 ;RMON OFFSET OF PNAME TABLE
74          000432 P1$EXT =: 000432 ;RMON OFFSET OF $P1EXT ADDRESS
75          000460 $H2UB =: 000460 ;RMON OFSET OF UB ENTRY VECTOR PTR
76
77          ; EXTENDED MEMORY SUBROUTINE OFFSETS FROM $P1EXT (.PIXDF)
78
79          177752 $MPMEM =: -22. ;OFFSET TO MAP KT-11 VIRTUAL TO PHYSICAL
80
81          ; UB ENTRY VECTOR EQUATES (.UBVDF)
82
83          ; UB.IDV =: 0 ; IDENTIFICATION WORD
84          ; UB.VDV =: <^rUBV> ; IDENTIFICATION WORD VALUE
85          ; UB.GET =: 2 ; JUMP TO GETUMR
86          ; UB.ALL =: 6 ; JUMP TO ALLUMR
87          ; UB.RLS =: 12 ; JUMP TO RLSUMR
88
89          ; DL INTERNAL DMA BUFFER EQUATES
90
91          000054 BUFSIZ =: 54*2*DL$UN ; SIZE OF DL INTERNAL DMA BUFFER
92          ; WORD SIZE OF DLBBUF*DL$UN
93          000001 NOUMRS =: <BUFSIZ+7777/10000> ; NUMBER OF PERMANENT UMRS REQUIRED
94
95
96

```

The .DRDEF performs much of the work of the preamble section. It is called with different parameters depending on whether or not the handler supports memory mapping (MMG\$T=1). The following includes much of the macro expansion:

Handler Macros

```

1          .SBTTL  HANDLER MACROS
2
4          .IF    EQ    MMG$T

```

The .DRDEF macro (with macro expansion) for unmapped monitors:

```

5          .DRDEF  DL,5,FILST$!SPFUN$!VARSZ$,DLSIZE,174400,160,DMA=NO
          .MCALL  .DRAST, .DRBEG, .DRBOT, .DREND, .DREST, .DRFIN, .DRFMS, .DRFMT
          .MCALL  .DRINS, .DRPTR, .DRSET, .DRSPF, .DRTAB, .DRUSE, .DRVTB
          .MCALL  .FORK, .QELDF
          .IIF  NDF RTE$M RTE$M=0
          .IIF  NE RTE$M RTE$M=1
          .IIF  NDF TIM$IT TIM$IT=0
          .IIF  NE TIM$IT TIM$IT=1
          .IIF  NDF MMG$T MMG$T=0
          .IIF  NE MMG$T MMG$T=1
          .IIF  NDF ERL$G ERL$G=0
          .IIF  NE ERL$G ERL$G=1
          .IIF  NE TIM$IT, .MCALL  .TIMIO, .CTIMI
000000    .QELDF
          .IIF  NDF MMG$T,MMG$T=1
          .IIF  NE MMG$T,MMG$T=1
000000    Q.LINK=:0
000002    Q.CSW=:2.
000004    Q.BLKN=:4.
000006    Q.FUNC=:6.
000007    Q.JNUM=:7.
000007    Q.UNIT=:7.
000010    Q.BUFF=:^o10
000012    Q.WCNT=:^o12
000014    Q.COMP=:^o14
          .IRP   X,<LINK,CSW,BLKN,FUNC,JNUM,UNIT,BUFF,WCNT,COMP>
          Q$'X=:Q.'X-^o4
          .ENDR
177774    Q$LINK=:Q.LINK-^o4
177776    Q$CSW=:Q.CSW-^o4
000000    Q$BLKN=:Q.BLKN-^o4
000002    Q$FUNC=:Q.FUNC-^o4
000003    Q$JNUM=:Q.JNUM-^o4
000003    Q$UNIT=:Q.UNIT-^o4
000004    Q$BUFF=:Q.BUFF-^o4
000006    Q$WCNT=:Q.WCNT-^o4
000010    Q$COMP=:Q.COMP-^o4
          .IF  EQ MMG$T
000016    Q.ELGH=:^o16
          .IFF
          Q.PAR=:^o16
          Q.MEM=:^o20
          .IRP   X,<PAR,MEM>
          Q$'X=:Q.'X-^o4
          .ENDR
          Q.ELGH=:^o24
          .ENDC
000001    HDERR$=:1
020000    EOF$=:^o20000
000400    VARSZ$=:^o400
001000    ABTIO$=:^o1000
002000    SPFUN$=:^o2000
004000    HNDLR$=:^o4000
010000    SPECL$=:^o10000
020000    WONLY$=:^o20000
040000    RONLY$=:^o40000
100000    FILST$=:^o100000
023742    DLDSIZ=:DLSIZE
000005    DL$COD=:5
102405    DLSTS=:<5>!<FILST$!SPFUN$!VARSZ$>
          .IIF  NDF DL$VEC,DL$VEC=160
          .GLOBL DL$VEC
          .
          .
          .

```

The .DRPTR macro with no parameters:

```

6          .DRPTR
          .
          .
7          .IFF      ;EQ MMG$T

```

The `.DRDEF` macro (with macro expansion) for mapped monitors. The handler is defined for the RL01; if it is for an RL02, the size is changed later. Note that handler supports UMRs.

```

8 000000          .DRDEF  DL,5,FILST$!SPFUN$!VAR SZ$,DLSIZE,174400,160,DMA=YES,PERMUMR=NOMURS

```

The `.DRPTR` macro with parameters:

```

9 000200          .DRPTR  FETCH=FETCH,LOAD=FETCH,RELEASE=RELEAS,UNLOAD=RELEAS
10          .ENDC    ;EQ      MMG$T

```

The `.DREST` macro (with macro expansion). Argument `REPLACE=RTABLE` shows DL does a software bad-block replacement—see installation code:

```

11 000022          .DREST  CLASS=DVC.DK,REPLACE=RTABLE
          000000 DVC.UK  =:0
          000001 DVC.NL  =:1
          000002 DVC.TT  =:^o2
          000003 DVC.TP  =:^o3
          000004 DVC.DK  =:^o4
          000005 DVC.MT  =:^o5
          000006 DVC.CT  =:^o6
          000007 DVC.LP  =:^o7
          000010 DVC.DE  =:^o10
          000011 DVC.DP  =:^o11
          000012 DVC.DL  =:^o12
          000013 DVC.NI  =:^o13
          000014 DVC.PS  =:^o14
          000015 DVC.VT  =:^o15
          000016 DVC.SI  =:^o16
          000017 DVC.SO  =:^o17
          000020 DVC.SB  =:^o20

          000001 DVM.NS  =:1
          000001 DVM.DX  =:1
          000002 DVM.DM  =:^o2
          000200 DVM.NF  =:^o200

          040000 DV2.V2  =:^o40000

          000001 HS2.BI  =:1
          000002 HS2.KI  =:^o2
          000004 HS2.KL  =:^o4
          000010 HS2.KU  =:^o10
          000020 HS2.MO  =:^o20
          .
          .
          .

```

Point to special functions for UNIBUS mapping register support:

```

18          .IF NE MMG$T
19 000076          .DRSPF  +UBTAB          ;SPFUN FOR UB GOES IN TABLE UBTAB
20          .ENDC    ;NE MMG$T
21

```

Define special functions:

```

22 000032          .DRSPF
23          .DRSPF  <FN$GET>          ;GET DEVICE STATUS
24 000032          .DRSPF  <FN$SIZ>    ;GET DEVICE SIZE
25 000032          .DRSPF  <FN$REP>    ;FORCE RE-READ OF REPLACEMENT TABLE
26 000032          .DRSPF  <FN$WRT>    ;ABSOLUTE WRITE (NO BAD BLOCK)
27 000032          .DRSPF  <FN$RED>    ;ABSOLUTE READ (REPLACEMENT)
28
29          .IIF    NDF    EIS$I    EIS$I = MMG$T
30          .IIF    EQ     EIS$I    .MCALL SOB

```

Define hardware offsets:

```

1          .SBTTL  HARDWARE DEFINITIONS
2
3          ; RL11 DEVICE REGISTER OFFSETS
4
5          ;DEFINE THE OFFSETS
6          000000      RLCS    ==: 0          ;CONTROL STATUS REGISTER
7          000002      RLBA    ==: 2          ;BUS ADDRESS REGISTER
8          000004      RLDA    ==: 4          ;DISK ADDRESS REGISTER
9          000006      RLMP    ==: 6          ;MULTI-PURPOSE REGISTER
10         000010      RLBAE   ==: 10         ;BUS ADDRESS REGISTER (EXTENDED)
11         .
12         .
13         .
14         .
15         .
16         .
17         .
18
19         ; RLCS BIT ASSIGNMENTS
20
21         100000      CSERR    ==: 100000     ;ERROR SUMMARY
22         040000      CSDE     ==: 040000     ;DRIVE ERROR
23         036000      CSERRC   ==: 036000     ;ERROR CODE MASK
24         020000      CSNXM    ==: 020000     ;NON-EXISTENT MEMORY
25         010000      CSDLT    ==: 010000     ;DATA LATE
26         010000      CSHNF    ==: 010000     ;HEADER NOT FOUND
27         004000      CSDCRC   ==: 004000     ;DATA CRC ERROR
28         004000      CSHCRC   ==: 004000     ;HEADER CRC ERROR
29         002000      CSOPI    ==: 002000     ;OPERATION INCOMPLETE
30         001400      CSDS01   ==: 001400     ;DRIVE SELECT BITS 0 AND 1
31         000400      CSDS0    ==: 000400     ;DRIVE SELECT BIT 0
32         000200      CSCRDY   ==: 000200     ;CONTROLLER READY
33         000100      CSIE     ==: 000100     ;INTERRUPT ENABLE
34         000040      CSBA17   ==: 000040     ;BUS ADDRESS BIT 17
35         000020      CSBA16   ==: 000020     ;BUS ADDRESS BIT 16
36         000016      CSFUN    ==: 000016     ;FUNCTION CODE
37         000001      CSDRDY   ==: 000001     ;DRIVE READY
38
39         ; RLCS FUNCTION CODE VALUES
40
41         000000      FNNOP    ==: 0*2        ;NO OPERATION
42         000002      FNWCHK   ==: 1*2        ;WRITE CHECK
43         000004      FNGSTS   ==: 2*2        ;GET DRIVE STATUS
44         000006      FNSEEK   ==: 3*2        ;SEEK
45         000010      FNRDH    ==: 4*2        ;READ HEADERS
46         000012      FNWRITE  ==: 5*2        ;WRITE DATA
47         000014      FNREAD   ==: 6*2        ;READ DATA
48         000016      FNRDNH   ==: 7*2        ;READ DATA WITH NO HEADER CHECK
49
50         ; RLMP GET STATUS RETURNED BIT ASSIGNMENTS
51
52         100000      STWDE    ==: 100000     ;WRITE DATA ERROR
53         040000      STCHE    ==: 040000     ;CURRENT HEAD ERROR
54         020000      STWL     ==: 020000     ;WRITE LOCK STATUS
55         010000      STSKTO   ==: 010000     ;SEEK TIMEOUT ERROR
56         004000      STSP     ==: 004000     ;SPEED ERROR
57         002000      STWGE    ==: 002000     ;WRITE GATE ERROR
58
59         001000      STVC     ==: 001000     ;VOLUME CHECK
60         000400      STDSE    ==: 000400     ;DRIVE SELECT ERROR
61         000200      STDT     ==: 000200     ;DRIVE TYPE
62         000100      STHS     ==: 000100     ;HEAD SELECT STATUS
63         000040      STCO     ==: 000040     ;COVER OPEN
64         000020      STHO     ==: 000020     ;HEADS HOME
65         000010      STBH     ==: 000010     ;BRUSHES HOME
66         000007      STST     ==: 000007     ;STATE BIT MASK
67         000005      STSLM    ==: 000005     ;DRIVE IN SEEK-LINEAR MODE STATE
68
69         ; RLDA BIT VALUES FOR SEEK COMMANDS
70
71         077600      SKCADF   ==: 077600     ;CYLINDER ADDRESS DIFFERENCE
72         000200      SKCA0    ==: 000200     ;CYLINDER ADDRESS DIFFERENCE BIT 0
73         000020      SKHS     ==: 000020     ;HEAD SELECT (SURFACE 0 OR 1)
74         000004      SKDIR    ==: 000004     ;DIRECTION (0 => OUTWARD, 1 => INWARD)
75         000001      SKMARK   ==: 000001     ;MARK BIT MUST BE 1 TO INDICATE A SEEK
76
77         ; RLDA BIT VALUES FOR I/O COMMANDS
78         077600      IOCA     ==: 077600     ;CYLINDER ADDRESS

```

```

79      000200      IOCA0  =: 000200      ;CYLINDER ADDRESS BIT 0
80      000100      IOHS   =: 000100      ;HEAD SELECT
81      000077      IOSA   =: 000077      ;SECTOR ADDRESS MASK
82
83      ; RLDA BIT VALUES FOR GET STATUS COMMAND
84
85      000010      GSRST  =: 000010      ;RESET DRIVE
86      000002      GSGS   =: 000002      ;GET STATUS INDICATOR MUST BE 1
87      000001      GSMARK =: 000001      ;THIS MUST BE 1 TO INDICATE GET STATUS
88

```

More RMON references:

```

89      ; RMON REFERENCES
90
91      000054      SYSPTR =: 54          ; SYSCOM pointer to RMON
92      000370      CONFIG2 =: 370       ; second configuration word
93      000100      BUS$   =: 000100     ;
94      020000      PROS$  =: 020000     ;
95      020100      BUS$M  =: BUS$!PROS$ ;Mask for type bits
96      020100      BUS$X  =: BUS$!PROS$ ;Strange (busless) KXJ
97      020000      BUS$C  =: PROS$     ;CTI bus
98      000100      BUS$Q  =: BUS$      ;QBUS
99      000000      BUS$U  =: 0         ;UNIBUS
100
101
102      000375      .READ  =: 375       ; EMT code for .READ
103      000010      ..READ =: 010      ; subcode for .READ
104      000375      .WRITE =: 375       ; EMT code for .WRITE
105      000011      ..WRIT =: 011      ; subcode for .WRITE
106
107      000017      SYSCHN =: 17        ; system channel

```

Installation checks (RL01/02 run on UNIBUS or Q-bus only):

```

3      .SBTTL  INSTALLATION CODE
4
5      000032      .DRINS  DL
6
7      000200      000401      BR      10$      ;Data device installation check
8      000202      .ASSUME . EQ INSSYS
9      000202      000414      BR      20$      ;System device installation check (none)
10
11      000204      013700      10$:  MOV     @#SYSPTR,R0      ; get address of RMON
12      000210      016000      MOV     CONFIG2(R0),R0      ;Get configuration word for BUS check
13      000214      042700      BIC     #^C<BUS$M>,R0      ;Isolate bus bits
14      000220      022700      CMP     #<BUS$X>,R0      ;Running on KXJ?
15      000224      001404      BEQ     30$      ;Yes, don't install
16      000226      022700      CMP     #<BUS$C>,R0      ;CTI?
17      000232      001401      BEQ     30$      ;Yes, don't install
18      000234      005727      20$:  TST     (PC)+      ; clear carry, skip setting carry
19      000236      000261      30$:  SEC      ; set carry
20      000240      000207      RETURN

```

The following is SET code. If there is insufficient room in the SET code area, some code can be moved up into the installation code area.

```

21
22      000242      O.SYWL:
23      000242      011600      MOV     @SP,R0          ; copy return address
24      000244      005200      INC     R0            ; point to opcode at return
25      000246      122720      CMPB   #BR/400,(R0)+   ; is it a BR xxx?
26      000252      001135      BNE   O.BAD         ; NO, old style SET
27      000254      010016      MOV   R0,@SP        ; use alternate return (RET+2)
28      000256      000533      BR    O.BAD         ; with carry set
29

```

The following sets up the table for software bad-block replacement:

```

30 000260      002 RTABLE: .BYTE  2,10.,5.,2.,40.,1.      ; Replacement factors table
    000261      012
    000262      005
    000263      002
    000264      050
    000265      001

```

All blocks can be replaced. This defines the geometry of the disk:

```

31                                     ; all replacable
32                                     ; 10. blocks to skip
33                                     ; 5. sectors of bad sector file
34                                     ; 2. tracks per cylinder
35                                     ; 40. sectors per track
36                                     ; 2**1 sectors per block
37

```

Installation code area size check:

```

38 000266      .Assume . LE 400,MESSAGE=<;Install code overflow>

```

The DL handler supports several SET command conditions:

Set Options

```

2          .SBTTL  SET OPTIONS
3
4 000266      .DRSET  CSR,      160000, O.CSR,  OCT
5 000412      .DRSET  VECTOR,  500,    O.VEC,  OCT
6
7 000422      .DRSET  RETRY,   127.,    O.RTRY, NUM
8
9          .IF NE ERL$G
10         .DRSET  SUCCES,  -1,    O.SUCC, NO
11         .ENDC ;NE ERL$G
12
13         004124      BTCSR   = <DLEND-DLSTRT>+<BOTCSR-DLBOOT>+1000
14
15         ; SET DL CSR=address
16
17 000432  020003  O.CSR:  CMP    R0,R3          ;CSR IN RANGE?
18 000434  103444      BLO    O.BAD            ;NOPE...
19 000436  010067      MOV    R0,INSCSR        ;YES, INSTALLATION CODE NEEDS IT
    177534
20 000442  010067      MOV    R0,DISCSR        ;AND RESORC DOES TOO
    177526
21
22         ; When the CSR is changed, we must also alter the bootstrap so
23         ; that it will use the correct CSR.
24
25
26 000446      .ADDR   #BAREA+4,R1            ;R1->READ/WRITE EMT AREA
    ; (BUFFER ADDRESS WORD)
27         ;R2->BUFFER
28 000454      .ADDR   #1000,R2              ; (OVERWRITES CORE COPY OF BLOCK 1)
29 000462  010211      MOV    R2,(R1)         ;SET THE BUFFER ADDRESS
30 000464  012741      MOV    #BTCSR/1000,-(R1) ; THE BLOCK TO READ/WRITE
    000004
31         ; (BOOT BLOCK THAT NEEDS ALTERING)
32 000470  005741      TST    -(R1)           ;R1->EMT AREA
33 000472  010003      MOV    R0,R3          ;SAVE CSR ELSEWHERE, EMT NEEDS R0
34 000474  010100      MOV    R1,R0         ;R0->EMT AREA FOR READ
35 000476  104375      EMT    .READ         ; *** (.READW) ***
36 000500  103422      BCS    O.BAD
37 000502  010362      MOV    R3,<BTCSR&777>(R2) ;SET THE NEW CSR
    000124
38 000506  010100      MOV    R1,R0         ;R0->EMT AREA FOR WRITE
39 000510      .ASSUME ..READ+1 EQ ..WRIT
40 000510  105260      INCB   1(R0)         ;CHANGE FROM 'READ' TO 'WRITE'
    000001
41 000514  104375      EMT    .WRITE         ; *** (.WRITW) ***
42 000516  103651      BCS    O.SYWL
43 000520  010100      MOV    R1,R0         ;R0->EMT AREA (LAST TIME, HONEST)

```

```

44 000522                .ASSUME ..WRIT-1 EQ ..READ
45 000522 105360          DECB  1(R0)                ;CHANGE FROM 'WRITE' TO 'READ'
                                000001
46 000526 012760          MOV   #1,2(R0)            ; OF HANDLER BLOCK 1
                                000001
                                000002
47 000534 104375          EMT   .READ                ; *** (.READW) ***
48 000536 103403          BCS   O.BAD
49 000540 010367          MOV   R3,DLCSR            ;TELL HANDLER ABOUT NEW CSR
                                000032'
50 000544 005727 O.GOOD: TST   (PC)+                ;GOOD RETURN (CARRY CLEAR)
51 000546 000261 O.BAD:  SEC                      ;ERROR RETURN (CARRY SET)
52 000550 000207          RETURN
53
54                        ; SET DL VECTOR=address
55
56 000552 020003 O.VEC:  CMP   R0,R3                ;VECTOR IN RANGE? (<500)
57 000554 103374          BHIS  O.BAD                ;NOPE...
58 000556 032700          BIT   #3,R0                ;YES, BUT ON A VECTOR BOUNDRY?
                                000003
59 000562 001371          BNE   O.BAD                ;NOPE...
60 000564 010067          MOV   R0,DLSTRT            ;TELL HANDLER ABOUT NEW VECTOR
                                000000'
61 000570 000765          BR    O.GOOD
62
63                        ; SET DL RETRY=count
64
65 000572 020003 O.RTRY:  CMP   R0,R3                ;Test retry limits
66 000574 101364          BHI   O.BAD                ;Branch if out of bounds
67 000576 010067          MOV   R0,DRETRY            ;Store the user selected retry count
                                000742'
68 000602 001761          BEQ   O.BAD                ;Zero retries not allowed
69 000604 000757          BR    O.GOOD                ;Otherwise, good
70
71                        .IF NE ERL$G
72
73                        ; SET DL [NO]SUCCES
74
75 O.SUCC: MOV   #0,R3                ;'SUCCESS' ENTRY POINT
76                                ; (MUST BE TWO WORDS)
77                                MOV   R3,SCSFLG            ;'NOSUCCESS' ENTRY POINT
78                                BR    O.GOOD
79                                .ENDC ;NE ERL$G
80
81 000606 017 BAREA:  .BYTE  SYSCHN,..READ            ;CHANNEL 17, READ
                                000607 010
82 000610                .BLKW                                ;BLOCK NUMBER
83 000612                .BLKW                                ;BUFFER ADDRESS
84 000614 000400                .WORD  256.                ;WORD COUNT
85 000616 000000                .WORD  0                ;COMPLETION (WAIT)
86

```

SET code overflow check:

```

87 000620                .Assume . LE 1000,MESSAGE=<;Set area overflow>
88                                .ENDC

```

Header Section

Request Entry Point

```

1                                .SBTTL  REQUEST ENTRY POINT
2
3                                .ENABL  LSB
4
5                                .IF EQ MMG$T

```

The .DRBEG macro for unmapped monitors:

```

6                                .DRBEG  DL
7                                .IFF ;EQ MMG$T

```

The .DRBEG macro for mapped monitors:

```

8 000620          .DRBEG DL,SPFUN=UBTAB
9                .ENDC ;EQ MMG$T
10

```

I/O Initiation Section

```

11          000006' DLBASE=DLSTRT+6
12
13 000024 016705          MOV     DLCQE,R5          ;POINT TO CURRENT QUEUE ELEMENT
          177760
14 000030 012704          MOV     (PC)+,R4          ;POINT TO CONTROLLER CSR
15 000032          .ASSUME .-DLSTRT LT 1000
16 000032 174400 DLCSR:  .WORD  DL$CSR          ;ADDRESS OF CONTROLLER
17 000034 016500          MOV     Q$FUNC(R5),R0          ;GET FUNCTION CODE / UNIT NUMBER
          000002
18 000040 110002          MOVB   R0,R2          ;GET SPECIAL FUNCTION CODE
          .
          .
          .
24 000042 120227          CMPB   R2,#FN$SIZ          ;.SPFUN LESS THAN 373 (SIGNED BYTE)
          000373
25 000046 002403          BLT    5$          ;YES, .SPFUN 200 THRU 372 INVALID
26 000050 120227          CMPB   R2,#FN$REP+1          ;IS THIS .SPFUN 375
          000375
27 000054 001002          BNE    10$          ;NO, HAVE VALID SPFUN REQUEST
28 000056 000167          5$:   JMP    DLQCOM          ;DISMISS QUEUE REQUEST
          001572
29
30 000062          10$:
32 000062 042700          BIC    #^C<7*400>,R0          ;ISOLATE UNIT NUMBER BITS
          174377
33 000066 020027          CMP    R0,#DL$UN*400          ;DO WE SUPPORT THIS UNIT?
          001000
34 000072 103136          BHIS   DLELNK          ;NO, ERROR NOW
35
36 000074 010067          MOV    R0,DLUNIT          ;SAVE UNIT NUMBER
          001062
37 000100          .ASSUME CSDS01 EQ      3*400
38 000100 012767          MOV    #FNREAD!CSIE,DLCODE          ;ASSUME READ (FOR TABLE)
          000114
          001050
39
40          .IF NE MMG$T
41 000106 120227          CMPB   R2,#FN$SIZ          ;SEE IF .SPFUN GET SIZE
          000373
42 000112 001407          BEQ    15$          ;YES -- DON'T CHANGE Q.BUFF AND Q.PAR
43 000114          .ASSUME Q$BLKN+4 EQ Q$BUFF
44 000114 022525          CMP    (R5)+,(R5)+          ;POINT TO Q.BUFF IN QUEUE ELEMENT
45 000116          .ASSUME Q$BUFF+2 EQ Q$WCNT ; done by MPPTR
46 000116 004777          CALL   @$MPPTR          ;CONVERT ADDRESS TO 18 BIT PHYSICAL
          002160
47 000122          .ASSUME Q$WCNT-2 EQ Q$BUFF
48 000122 012645          MOV    (SP)+,-(R5)          ;REPLACE Q.BUFF WITH BITS <15:00>
49 000124          .ASSUME Q$BUFF-4 EQ Q$BLKN
50 000124 024545          CMP    -(R5),-(R5)          ;FIX QUEUE ELEMENT POINTER
51 000126 012665          MOV    (SP)+,Q$PAR(R5)          ;SAVE BITS <21:16> IN Q.PAR WORD
          000012
52          .ENDC ;NE MMG$T
53

```

The software bad-block replacement table is named *DLBBUF*:

```

54 000132          15$:
55 000132          .ADDR  #DLBBUF-<DLTSIZ+2>,R3          ; GET BIASED ADDRESS OF TABLE BUFFER
56 000140 000300          SWAB   R0          ;GET UNIT NUMBER
57 000142 062703          20$:  ADD    #DLTSIZ+2,R3          ;POINT TO NEXT UNIT'S TABLE
          000054
58 000146 005300          DEC    R0          ; REDUCE UNIT NUMBER
59 000150 100374          BPL    20$          ; ALL GONE?
60 000152 010327          MOV    R3,(PC)+          ;SAVE POINTER TO UNIT'S
61 000154 000000          DLCC:  .WORD  0          ; CURRENT CYLINDER TABLE (LOW ADDR)
62 000156 005723          TST    (R3)+          ;POINT TO REPLACEMENT TABLE
63 000160          .ASSUME .+4 EQ DLUSIZ

```

```

65 000160 012727      MOV      #DLSIZE,(PC)+      ;ASSUME RL01
      023742
      .
      .
      .

```

Test for RL01 or RL02; select correct size:

```

70 000164 000000 DLUSIZ: .WORD 0
72 000166 004767      CALL      DLGST      ;GET DISK STATUS
      001546
73 000172 105701      TSTB     R1      ;SINGLE DENSITY?
74 000174 100003      BPL      25$      ;IF ZERO, RL01 SINGLE DENSITY
75 000176 012767      MOV      #DLSIZ2,DLUSIZ ;IF SET, RL02 DOUBLE DENSITY
      047742
      177760
76 000204 005700 25$:    TST      R0      ;Now, error in get status?
77 000206 100403      BMI      30$      ;Yes, invalidate everything
78 000210 032701      BIT      #STVC,R1 ;IS THERE A NEW DISK IN THIS DRIVE?
      001000
79 000214 001403      BEQ      35$      ;NO, SAME AS LAST TIME
80 000216 012743 30$:    MOV      #-1,-(R3) ;INVALIDATE CURRENT CYLINDER
      177777
81 000222 012313      MOV      (R3)+,@R3 ; AND INVALIDATE REPLACEMENT TABLE
82
      .IFF
83
      CMPB     R2,#FN$GET ;SEE IF .SPFUN GET SPECIAL STATUS
84
      BEQ      DLGSTA ;YES, GO DO IT!
85
      CALL      DLGST ;GET DISK STATUS (NORMAL)
86
      TST      R0      ;Now, error in get status?
87
      BMI      30$      ;Yes, invalidate everything
88
      CALL      INVVC ;INVALIDATE IF VOLUME CHECK ON
89
      BR       35$      ;SKIP NEXT
90
91
      30$:    CALL      INVAL ;UNCONDITIONAL INVALIDATION
92
      .ENDC
93

```

Following code decides if we use bad-block replacement table (only for special functions). DLSQUE, DLADDR, and DLEXFR are used for replacement table read.

```

94 000224 120227 35$:    CMPB     R2,#FN$REP ;CHECK OUT THE SPECIAL FUNCTION
      000374
95 000230 002002      BGE      40$      ;BRANCH IF NOT 'GET SIZE'
96
      ; (NOTE SIGNED COMPARE)
97 000232 000167      JMP      DLGSIZ ;GO DO 'GET SIZE'
      001434
98
99 000236 001410 40$:    BEQ      50$      ;GO READ BAD-BLOCK REPLACEMENT TABLE
100 000240 101045      BHI      55$      ;GO DO ABSOLUTE BLOCK READ/WRITE
101 000242 005765      TST      Q$WCNT(R5) ;NORMAL REQUEST, SEEK?
      000006
102 000246 001002      BNE      45$      ;BRANCH IF NOT
103 000250 000167 DLFLNK: JMP      DLQCOM ;.DRFIN TIME
      001400
104
105 000254 005713 45$:    TST      @R3      ;IS TABLE IN MEMORY YET?
106 000256 100046      BPL      DLTRAN ;YES, WE CAN GO DO THE TRANSFER

```

Reread the replacement table.

1. Read replacement table into memory if it's not there.
 - a. Save current queue element.
 - b. Build pseudoqueue element to read the replacement table (DLSQUE).
 - c. Allow transfer to start (DLADDR).
 - d. Eventually, the request gets to the end of the I/O initiation section and returns to monitor.

- e. Request is completed and returns to interrupt entry (.DRAST).
- f. Continues down to DLEXFR to determine if we were rereading the table and dismiss the the pseudoqueue element if we were. The queue element that prompted the reading of the replacement table still exists. It can now be processed.

2. Replacement table already in memory—use it. Go to DLTRAN to use it.

```

107          ;
108          ; WE ALWAYS COME HERE TO REREAD THE REPLACEMENT TABLE
109          ;
110 000260    50$:
111          .IF NE MMG$T
112 000260    010346      MOV     R3,-(SP)          ;SAVE R3
113 000262          .ADDR  #DLBBUF,R3          ;R3=PIC ADDRESS OF START OF DLBBUF
114 000270    016701      MOV     DLCC,R1          ;R1=START ADDRESS FOR THIS UNIT
115          177660
116 000274    160301      SUB     R3,R1          ;R1=OFFSET INTO DLBBUF FOR THIS UNIT
117 000276    062701      ADD     #2,R1          ;POINT TO REPLACEMENT TABLE
118          000002
119 000302    016702      MOV     BUFADH,R2        ;GET HI ORDER DLBBUF ADDRESS
120          001760
121 000306    016703      MOV     BUFADL,R3        ;GET LOW ORDER DLBBUF ADDRESS
122          001756
123 000312    060103      ADD     R1,R3          ;R3=THIS UNIT'S START ADDR IN UMR
124 000314    103002      BCC     52$          ;BRANCH IF NO CARRY
125 000316    062702      ADD     #CSBA16,R2       ;ADD CARRY TO HI ORDER ADDR
126          000020
127 000322    010267    52$:  MOV     R2,DLBPAR        ;PUT HI ORDER ADDR INTO PSEUDO QEL
128          001550
129          ;
130          ; MOV     Q$MEM(R5),DLBMEM        ;PUT Q$MEM INTO PSEUDO QEL (NOT NEEDED)
131          .ENDC ;NE MMG$T
132 000326    012701      MOV     #1,R1          ;TABLE IS IN BLOCK 1
133          000001
134 000332    012702      MOV     #DLTSIZ/2,R2       ;WORDS TO READ (TABLE SIZE)
135          000025

```

Build queue element to read table.

```

127 000336    004767      CALL    DLSQUE          ;SET UP REST OF PSEUDO QUEUE ELEMENT
128          001470
129
130          .IF NE MMG$T
131 000342    012603      MOV     (SP)+,R3        ;RESTORE R3 (ADDR FOR MOV'S)
132          .ENDC ;NE MMG$T
133 000344    012713      MOV     #-1,@R3        ;FLAG THAT THERE IS NO TABLE IN MEMORY
134          177777
135 000350    011343      MOV     @R3,-(R3)      ;VOID CURRENT CYLINDER, TOO

```

Read in the table.

At DLADDR, pseudoqueue element is processed to read in replacement table. I/O initiation will start transfer and return to the monitor. When transfer is complete, the .DRAST section is entered to dismiss the pseudoqueue element.

```

135 000352    000512      BR     DLADDR          ;COMPUTE DISK ADDRESS AND START THE
136          ; TABLE READ
137

```

```

138 000354 105202 55$: INCB R2 ;ABSOLUTE BLOCK READ?
139 000356 ;.ASSUME FN$RED EQ 377
140 000356 001510 BEQ DLADDR ;YES, WE ARE ALL SET UP
141 000360 012767 MOV #FNWRITE!CSIE,DLCODE ;SET WRITE FUNCTION CODE
      000112
      000570
142 000366 000504 BR DLADDR ;GO DO IT
143
144 ;.IF NE ERL$G
145 .ASSUME .-DLSTRT LT 1000
146 SCSFLG: .WORD 0 ; :SUCCESS LOGGING FLAG (DEFAULT=YES)
147 ; =0 - LOG SUCCESSES
148 ; <0 - DON'T LOG SUCCESSES
149 .ENDC ;NE ERL$G
150
151 .DSABL LSB
      .
      .
      .
74
75 000370 000167 DLELNK: JMP DLEROR ;LINK TO FATAL ERROR
      001212

```

Set up and perform I/O:

```

1 ;SBTTL INITIALIZE FOR TRANSFER, SET FUNCTION CODE, FIX WORD COUNT
2
3 ;+
4 ; SET READ OR WRITE FUNCTION CODE
5 ; IF TRANSFER HAS REPLACED BLOCKS IN IT, BREAK IT INTO PIECES AND
6 ; SEND EACH PIECE TO DLADDR SEPARATELY FOR I/O
7 ; NOTE: ALL PIECES EXCEPT THE FIRST ARE BLOCK MULTIPLES
8 ;
9 ; R4 -> CSR
10 ; R5 -> USER QUEUE ELEMENT
11 ;-
12
13 .ENABL LSB
14
15 000374 005765 DLTRAN: TST Q$WCNT(R5) ;READ OR WRITE OPERATION?
      000006
16 000400 100005 BPL 1$ ;READ...
17 ; (NOTE: THIS FAILS 2ND TIME THROUGH)
18 000402 005465 NEG Q$WCNT(R5) ;WRITE, MAKE WORD COUNT POSITIVE
      000006
19 000406 012767 MOV #FNWRITE!CSIE,DLCODE ;SET WRITE FUNCTION CODE
      000112
      000542
20 000414 016502 1$: MOV Q$WCNT(R5),R2 ;MAYBE, DETERMINE LENGTH OF
      000006
21 000420 010203 MOV R2,R3 ;TRANSFER IN BLOCKS
22 000422 062703 ADD #255.,R3
      000377
23 000426 105003 CLRB R3
24 000430 000303 SWAB R3
25 000432 ;.ASSUME Q$BLKN EQ 0
26 000432 061503 ADD @R5,R3 ;COMPUTE FIRST BLOCK AFTER TRANSFER
27 000434 026703 CMP DLUSIZ,R3 ;DOES OPEATION EXTEND INTO REPLACEMENT
      177524

```

Checking if bad-block replacement is needed:

```

28
29 000440 103753      BLO  DLLELNK      ;BLOCKS ?
30 000442 016700      MOV  DLCC,R0      ;YES, NOT ALLOWED W READ/WRITE
                          177506
                          ;POINT TO REPLACEMENT TABLE - 2
31 000446 005760      TST  4(R0)        ;IS THE FIRST REPLACEMENT BLOCK = 0?
                          000004
32 000452 001452      BEQ  DLADDR      ;YES, THEN INVALID TABLE (FILES-11)
33 000454 005720      2$: TST  (R0)+      ;SKIP OVER REPLACEMENT BLOCK NUMBER
34 000456 012001      MOV  (R0)+,R1    ;GET NEXT BLOCK NUMBER TO REPLACE
35 000460 001447      BEQ  DLADDR      ;END OF TABLE, NO REPLACEMENT, DO IO
36 000462                .ASSUME Q$BLKN EQ 0
37 000462 020115      CMP  R1,@R5      ;THIS BAD BLOCK PART OF TRANSFER?
38 000464 103773      BLO  2$          ;NOPE, BELOW, IGNORE IT
39 000466 020103      CMP  R1,R3       ;BAD BLOCK WITHIN TRANSFER?
40 000470 103043      BHIS DLADDR      ;NOPE, BEYOND, WHOLE TRANSFER GOOD
41 000472 011001      MOV  @R0,R1      ;YES, PICK UP REPLACEMENT BLOCK NUMBER
42 000474 014000      MOV  -(R0),R0    ;GET BAD BLOCK NUMBER
43 000476                .ASSUME Q$BLKN EQ 0
44 000476 161500      SUB  @R5,R0      ;COMPUTE DISTANCE OF BAD BLOCK
45
46 000500 001004      BNE  3$          ; INTO TRANSFER
47
48
                          ; NOT THE FIRST BLOCK,
                          ; GO DO GOOD FIRST PART

```

The replacement table is being used. Pseudoqueue elements are built to break-up the transfer.

```

49
50
51
52 000502 005200      INC  R0           ;SET BLOCK COUNT TO BE 1 BLOCK
53 000504 000302      SWAB R2          ;IS THE REAL COUNT > 1 BLOCK
54
55 000506 001403      BEQ  5$          ; HI BYTE>0?
56 000510 000401      BR   4$          ;COUNT < 256. WORDS, FIX AND USE IT
57
58
59
60
61
62 000512                .ASSUME Q$BLKN EQ 0
63 000512 011501      3$: MOV  @R5,R1    ;START BLOCK OF PARTIAL=ORIGINAL BLOCK
64 000514 010002      4$: MOV  R0,R2    ;COPY BLOCK COUNT OF TRANSFER
65 000516 000302      5$: SWAB R2       ; MULTIPLY BY 256. TO GET WORD COUNT
66 000520 016503      MOV  Q$BUFF(R5),R3 ;GET ORIGINAL BUFFER ADDRESS
                          000004
67
68 000524                .ASSUME Q$BLKN EQ 0
69 000524 060015      ADD  R0,@R5      ;UPDATE BLOCK NUMBER BY PARTIAL
70
71 000526 160265      SUB  R2,Q$WCNT(R5) ;FIX WORD COUNT IN USER QUEUE ELEMENT
                          000006
72 000532 010200      MOV  R2,R0       ;COPY THE WORD COUNT
73 000534 006300      ASL  R0           ;CHANGE WORD COUNT TO BYTE COUNT
74 000536 060065      ADD  R0,Q$BUFF(R5) ;UPDATE USER BUFFER ADDRESS
                          000004
75
76
77 000542 016567      .IF NE MMG$T
                          MOV  Q$PAR(R5),DLBPAR ;*C*SET HI ADDR BITS IN PSEUDO QUEUE
                          000012
                          001326
78 000550 016567      MOV  Q$MEM(R5),DLBMEM ;*C*SET HI ADDR BITS IN PSEUDO QUEUE
                          000014
                          001322
79 000556 103006      BCC  6$          ;NO OVERFLOW
80 000560 062765      ADD  #CSBA16,Q$PAR(R5) ;OVERFLOW ORIGINAL ADDRESS INTO
                          000020
                          000012
81
82 000566 062765      ADD  #CSBA16,Q$MEM(R5) ; HIGH BITS
                          000020
                          000014
                          ; OVERFLOW ORIGINAL ADDRESS INTO
83
84 000574                6$:
85
86
87 000574 004767      CALL DLSQUE      ; FILL IN REST OF PSEUDO QUEUE

```

```

001232
88 000600      .BR      DLADDR      ;COMPUTE ADDRESS AND DO I/O
89
90      .DSABL  LSB

1
2
3      ;+
4      ;      R4 -> CSR
5      ;      R5 -> QUEUE ELEMENT (USER OR PSEUDO)
6      ;-
7
8      .ENABL  LSB
9
10 000600 010527 DLADDR: MOV      R5,(PC)+      ;SAVE POINTER TO QUEUE ELEMENT
11                                     ; WE ARE USING
12 000602 000000 DLQPTR:  .WORD  0
13 000604                                     .ASSUME Q$BLKN EQ 0
14 000604 011502      MOV      @R5,R2      ;GET BLOCK NUMBER
15 000606 100670      BMI      DLELNK      ;NO NEGATIVE BLOCK NUMBERS!
16 000610 012701      MOV      #DLBPT,R1     ;GET NUMBER OF BLOCKS ON ONE TRACK
    000024
17 000614                                     .ASSUME DLBPT  EQ 20.
18 000614 005000      CLR      R0      ;INITIALIZE I/O DISK ADDRESS TO 0
19 000616 000410      BR       2$      ;ENTER DIVIDE LOOP
20
21 000620 010203 1$:  MOV      R2,R3      ;COPY DIVIDEND
22 000622 042702      BIC      #^C<17>,R2    ;COMPUTE DIV = 16Q + R
    177760
23 000626 040203      BIC      R2,R3      ; AND GET 16Q TO WORK WITH
24 000630 060300      ADD      R3,R0      ;RESULT <- RESULT + IOHS/4
25 000632                                     .ASSUME IOHS/2/2 EQ 16.
26 000632 006203      ASR      R3      ;COMPUTE 8Q
27 000634 006203      ASR      R3      ; THEN 4Q
28 000636 160302      SUB      R3,R2      ;NEW DIVIDEND = R - 4Q
29 000640 020201 2$:  CMP      R2,R1      ;DONE? (NUMBER NOW < DLBPT)
30 000642 103366      BHIS    1$      ;NOPE...
31 000644 006300      ASL      R0      ;YES, QUOTIENT*IOHS/4 => QUO*IOHS/2
32 000646 050200      BIS      R2,R0      ;MERGE BLOCK NUMBER WITH TRACK
33 000650 006300      ASL      R0      ;*2 FOR TWO 128. WORD SECTORS/BLOCK
34 000652 103646      BCS      DLELNK      ;OVERFLOW MEANS BEYOND END OF DEVICE
35 000654 100004      BPL      3$      ;POSITIVE IS OK FOR EITHER RL01/02
36 000656 026727      CMP      DLUSIZ,#DLSIZ2 ;NEGATIVE IS OK FOR RL02 ONLY
    177302
    047742
37 000664 001241      BNE      DLELNK      ; BUT NOT OK FOR RL01
38 000666 010067 3$:  MOV      R0,DLDA     ;SAVE STARTING DISK ADDRESS
    000256
39 000672 160201      SUB      R2,R1      ;CALCULATE BLOCKS LEFT ON TRACK
40 000674 000301      SWAB    R1      ;CONVERT TO WORDS LEFT ON TRACK
41 000676 010167      MOV      R1,DLWTRK    ;SAVE THAT NUMBER
    000224
42
43      .IF NE  ERL$G
44      MOV      Q$WCNT(R5),DLWC      ;SET WORD COUNT FOR EL
45      .ENDC ;NE ERL$G
46
47 000702 012727      MOV      #1,(PC)+      ;CLEAR RETRY COUNT
    000001
48                                     ; (THESE ARE FATAL ERRORS)
49 000706 000000 DLRTY:  .WORD  0
50 000710 004767      CALL    DLRST      ;RESET DRIVE

```

```

001062
51 000714 004767 CALL DLGST ;AND GET STATUS
001020
52 000720 100434 BMI DLERJM ;ERROR HERE IS FATAL
53 000722 006200 ASR R0 ;IS THE DRIVE READY?
54 000724 .ASSUME CSDRDY EQ 1
55 000724 103032 BCC DLERJM ;NO, FATAL UNRETRYABLE ERROR
56 000726 042701 BIC #STWL!STHS!STDT,R1 ;IGNORE WRITE LOCK, HEAD SELECT,
020300
57 ; DRIVE TYPE
58 000732 022701 CMP #STHO!STBH!STSLM,R1 ;HEADS, BRUSHES AND STATE OK?
000035
59 000736 001025 BNE DLERJM ;NO, FATAL ERROR
60
61 000742' DRETRY = .+2
62 000740 .ASSUME DRETRY-DLSTRT LT 1000
63 000740 012767 MOV #DLRCNT,DLRTY ;SET REAL RETRY COUNT
000010
177740
64 000746 .BR DLTRAK ;GET ON TRACK
65
66 .DSABL LSB

1 .SBTTL ENSURE THAT DISK IS ON TRACK BEFORE TRANSFER
2
3 ;+
4 ; CALCULATE THE DIFFERENCE WORD FOR THE SEEK.
5 ; TRY 16 TIMES TO READ A HEADER.
6 ; IF ALL FAIL, LOG AN ERROR AND ISSUE A REVERSE SEEK (SEEK -1 TRACK)
7 ; AND A READ HEADER TO CAUSE AN INTERRUPT.
8 ;
9 ; R4 -> CSR
10 ; R5 -> QUEUE ELEMENT
11 ;-
12
13 .ENABL LSB
14
15 000746 005027 DLTRAK: CLR (PC)+ ;RESET REVERSE SEEK FLAG
16 000750 000000 DLREV: .WORD 0
17 000752 017701 MOV @DLCC,R1 ;GET CURRENT CYLINDER
177176
18 000756 022701 CMP #-1,R1 ;IS IT VALID?
177777
19 000762 001015 BNE 2$ ;YES, USE IT TO START WITH
20 ;***ACTION*** OLD CODE HAS ANOTHER RETRY VALUE
21 000764 016702 MOV DRETRY,R2 ;SET READ HEADER RETRY COUNT
177752
22 000770 006302 ASL R2 ; (DLRCNT*2)
23 000772 012701 1$: MOV #FNRDH,R1 ;SET CODE FOR READ HEADERS FUNCTION
000010
24 000776 004767 CALL DLXCT ;EXECUTE THE FUNCTION
001006
25 001002 100005 BPL 2$ ;FUNCTION EXECUTED OK
26 001004 077206 SOB R2,1$ ; any retries left?
27 001006 105267 INCB DLREV ;SET REVERSE SEEK FLAG
177736
28 001012 000167 DLERJM: JMP DLERRH ;RETRY OPERATION
000452
29
30 001016 016700 2$: MOV DLDA,R0 ;RETRIEVE STARTING DISK ADDRESS
000126
31 001022 012702 MOV #IOSA,R2 ;MASK OUT
000077
32 001026 040200 BIC R2,R0 ;SECTOR BITS FROM DESIRED ADDRESS
33 001030 040201 BIC R2,R1 ; AND FROM CURRENT ADDRESS
34 001032 020001 CMP R0,R1 ;DO WE NEED TO DO A SEEK?
35 001034 001427 BEQ DLXFER ;NOPE, ALREADY ON CYLINDER AND HEAD
36 001036 010003 MOV R0,R3 ;YES, SAVE DESIRED CYLINDER AND HEAD
37 001040 005202 INC R2 ;GET MASK FOR HEAD SELECT
001042 .ASSUME IOHS EQ IOSA+1
39 001042 040200 BIC R2,R0 ;STRIP HEAD SELECT BIT FROM
40 ; DESIRED ADDRESS
41 001044 040201 BIC R2,R1 ; AND FROM CURRENT ADDRESS
42 001046 160001 SUB R0,R1 ;COMPUTE DISTANCE FROM DESIRED
43 ; TO ACTUAL CYLINDER
44 001050 .ASSUME SKCADF EQ IOCA
45 001050 103003 BHIS 3$ ;DESIRED <= ACTUAL, MOVE TOWARD EDGE

```

```

46 001052 005401      NEG    R1                ;DESIRED > ACTUAL, MOVE TOWARD SPINDLE
47 001054 052701      BIS    #SKDIR,R1        ; (SET DIRECTION BIT)
    000004

48 001060 005201 3$:  INC    R1                ;SET MARKER BIT
49 001062                .ASSUME  SKMARK  EQ    1
50 001062 030203      BIT    R2,R3            ;DO WE WANT TO USE SURFACE 1?
51 001064 001402      BEQ    4$                ;NO
52 001066 052701      BIS    #SKHS,R1        ;YES, SET SURFACE 1 BIT
    000020
53 001072 012777 4$:  MOV    #-1,@DLCC        ;VOID KNOWLEDGE OF CURRENT CYLINDER
    177777
    177054
54 001100 004767      CALL  DLSEEK            ;EXECUTE THE SEEK
    000622
55 001104 100571      BMI    DLERRH          ;OOPS, ERROR EXECUTING SEEK
56 001106 016777      MOV    DLDA,@DLCC      ;SET CURRENT CYLINDER
    000036
    177040

57 001114                .BR    DLXFER            ;NOW DO THE TRANSFER
58
59                .DSABL  LSB

1                .SBTTL  DLXFER - START AN I/O TRANSFER
2
3                ;+
4                ;      R4 -> CSR
5                ;      R5 -> QUEUE ELEMENT
6                ;-
7
8                .ENABL  LSB
9 001114      DLXFER:

11 001114 062704      ADD    #RLMP,R4        ;POINT TO RLMP IN CONTROLLER
    000006

13 001120 062705      ADD    #Q$WCNT,R5     ;POINT TO WORD COUNT IN QUEUE ELEMENT
    000006

14 001124 012703      MOV    (PC)+,R3        ;GET NUMBER OF WORDS LEFT ON TRACK
15 001126 000000      DLWTRK: .WORD 0
16 001130 020315      CMP    R3,@R5          ;COMPARE AGAINST TOTAL TRANSFER
17 001132 101401      BLOS  1$                ;<=, USE REMAINDER OF TRACK
18 001134 011503      MOV    @R5,R3          ;>, USE TOTAL TRANSFER COUNT
19 001136 010327 1$:  MOV    R3,(PC)+        ;SAVE TRANSFER COUNT FOR LATER
20 001140 000000      DLWC:  .WORD 0        ; : TRANSFER COUNT
21 001142 005403      NEG    R3                ;MUST BE 2'S COMPLEMENT

23 001144 010314      MOV    R3,@R4          ;LOAD WORD COUNT INTO CONTROLLER
24 001146 012744      MOV    (PC)+,-(R4)     ;LOAD STARTING DISK ADDRESS
25 001150 000000      DLDA:  .WORD 0
26 001152 014544      MOV    -(R5),-(R4)     ;SET BUS ADDRESS
    .
    .
35 001154 012700      MOV    (PC)+,R0        ;GET FUNCTION CODE
36 001156 000000      DLCODE: .WORD 0        ;READ OR WRITE CODE
37 001160 052700      BIS    (PC)+,R0        ;ADD IN UNIT SELECT BITS
38 001162 000000      DLUNIT: .WORD 0        ;UNIT NUMBER IN BITS 8-9
39
40                .IF NE  MMG$T
41 001164 000416 $RLV1A: BR    10$                ;IF NO RLV12...
42                ; (CHANGED TO 'NOP' IF USING RLV12)
43 001166 016546      MOV    Q$PAR-Q$BUFF(R5),-(SP) ;SAVE Q22 HIGH-ORDER BITS
    000006
44 001172 006216      ASR    (SP)            ;SHIFT THEM TO THEIR CORRECT POSITIONS
45 001174 006216      ASR    (SP)
46 001176 006216      ASR    (SP)
47 001200 006216      ASR    (SP)
48 001202 012664      MOV    (SP)+,RLBAE-RLBA(R4) ;SET THE HIGH-ORDER BITS
    000006
49 001206 016546      MOV    Q$PAR-Q$BUFF(R5),-(SP) ;SAVE HIGH-ORDER BUS ADDRESS
    000006
50 001212 042716      BIC    #<^C60>,(SP)    ;STRIP TO HIGH-ORDER BITS<17:16>
    177717
51 001216 052600      BIS    (SP)+,R0        ; AND MERGE WITH COMMAND WORD

```

```

52 001220 000410      BR      30$
53
54 001222 032765 10$:  BIT      #1700,Q$PAR-Q$BUFF(R5) ;22-BIT ADDRESS SPECIFIED?
    001700
    000006
55 001230 001402      BEQ      20$ ;NOPE, THEN ADDRESS IS OKAY TO USE
56 001232 000167      JMP      DLEROR ;YES, CAN'T BE USED ON NON RLV12
    000350
57
58 001236 056500 20$:  BIS      Q$PAR-Q$BUFF(R5),R0 ;MERGE EXTENDED ADDRESS BITS INTO
    000006 ; COMMAND WORD
59
60 001242          30$:
61      .ENDC ;NE MMG$T
62
64 001242 010044      MOV      R0,-(R4) ;LOAD FUNCTION AND GO
68 001244 000207      RETURN ;WAIT FOR AN INTERRUPT
69
70      .DSABL LSB

```

Interrupt Service Section

```

1      .SBTTL DLINT - INTERRUPT ENTRY POINT
2
3          ; INTERRUPTS ENTER THE HANDLER HERE
4
5      .ENABL LSB

```

The `.DRAST` macro:

When a function is completed, the device interrupts, and the handler is entered here to dismiss the interrupt and the queue element.

```

6
7 001246      .DRAST DL,5
    .
    .
    .

```

Drop to fork level rather than device priority because the routine is lengthy and it needs all the registers.

```

14 001256      .FORK  DLFBLK ;GO TO FORK LEVEL

```

Load the registers.

```

15 001264 016704      MOV      DLCSR,R4 ;POINT TO CSR ADDRESS
    176542
16 001270 016705      MOV      DLQPTR,R5 ;POINT TO QUEUE ELEMENT
    177306
17 001274 105767      TSTB    DLREV ;REVERSE SEEK IN PROGRESS?
    177450
18 001300 001222      BNE      DLTRAK ;YES, GO RETRY THE REAL TRANSFER
20 001302 005714      TST      @R4 ;CHECK RLCS
25 001304 100471      BMI      DLERRH ;IF ERROR, GO DIAGNOSE IT
26 001306          .ASSUME CSERR EQ 100000
27 001306 016703      MOV      DLWC,R3 ;GET WORD COUNT OF THIS TRANSFER
    177626
28 001312 160365      SUB      R3,Q$WCNT(R5) ;CALCULATE WORDS REMAINING TO TRANSFER
    000006
29 001316 001036      BNE      2$ ;MORE TO DO, USE NEXT TRACK
30 001320 026727      CMP      DLCODE,#FNWRITE!CSIE ;WAS THE LAST FUNCTION A WRITE?
    177632
    000112
31 001326 001030      BNE      11$ ;NO, DONE WITH THIS (PARTIAL) ELEMENT
33 001330 032764      BIT      #1,RLDA(R4) ;GOT A SECTOR TO WRITE YET?
    000001
    000004

```

```

38 001336 001424      BEQ    11$                ;NO, DON'T ZERO FILL
39 001340 005265      INC    Q$WCNT(R5)          ;SET WORD COUNT TO 1
    000006
40
    ; (CONTROLLER FILLS 127.)
41
    .IF EQ MMG$T
42      .ADDR    #DLFILL,-(SP)      ;GET THE BUFFER ADDRESS
43      .IFF ;EQ MMG$T
44 001344 016765      MOV    BUFADH,Q$PAR(R5)      ;GET HI ADDR OF UMR
    000716
    000012
45 001352 016746      MOV    BUFADL,-(SP)          ;GET LO ADDR OF UMR
    000712
46 001356 062716      ADD    #<BUFEND-DLBBUF>,@SP    ;POINT TO DLFILL
    000130
47 001362 103003      BCC    100$                ;IF NO OVERFLOW, BRANCH
48 001364 062765      ADD    #CSBA16,Q$PAR(R5)      ;UPDATE HI ORDER ADDRESS BITS
    000020
    000012
49 001372 100$:
50      .ENDC ;EQ MMG$T
51 001372 012665      MOV    (SP)+,Q$BUFF(R5)      ;SET THE BUFFER ADDRESS
    000004
53 001376 016467      MOV    RLDA(R4),DLDA         ; AND THE DISK ADDRESS
    000004
    177544
57
58 001404 000167 1$:   JMP    DLTRAK                ;GO DO IT (DLWTRK > 1 = Q$WCNT)
    177336
59
60 001410 000167 11$:  JMP    DLEXFR                ;GO FINISH TRANSFER
    000204
61
62 001414 006303 2$:   ASL    R3                    ;CHANGE WORD COUNT TO BYTE COUNT
63 001416 060365      ADD    R3,Q$BUFF(R5)        ;UPDATE USER BUFFER ADDRESS
    000004
64
65      .IF NE MMG$T
66 001422 103003      BCC    3$                    ;NO OVERFLOW
67 001424 062765      ADD    #CSBA16,Q$PAR(R5)    ;UPDATE HIGH ORDER ADDRESS BITS
    000020
    000012
68 001432 3$:
69      .ENDC ;EQ MMG$T
70
71 001432 052767      BIS    #77,DLDA            ;UPDATE SURFACE/CYLINDER ADDRESS
    000077
    177510
72 001440 005267      INC    DLDA                  ; TO FIRST SECTOR, NEXT HEAD/CYLINDER
    177504
73 001444 001460      BEQ    DLEROR                ;OVERFLOWED DEVICE !!!
74 001446 100004      BPL    301$                 ;OK FOR EITHER RL01/02
75 001450 026727      CMP    DLUSIZ,#DLSIZ2      ;MINUS OK ONLY FOR RL02
    176510
    047742
76 001456 001053      BNE    DLEROR                ; VERY BAD IF RL01 !!!
77 001460 012767 301$:  MOV    #DLWPT,DLWTRK        ;SAVE NUMBER OF WORDS ON A WHOLE TRACK
    012000
    177440
78 001466 000746 4$:   BR     1$                    ;GO CONTINUE TRANSFER ON NEXT TRACK

1
2      .SBTTL  HANDLE THE ERRORS
3 001470      DLERRH:
4      .IF EQ  ERL$G
6 001470 011403      MOV    @R4,R3                ;GET RLCS CONTENTS WITH ERROR BITS

```

```

10      .IFF
11      MOV      R4,R1          ;GET CSR ADDRESS
12      .ADDR   #DLRBLK,R2    ; CALCULATE ADDRESS OF REGISTER BUFFER
13      MOV      R2,R3          ;SAVE BUFFER ADDRESS
14      MOV      (R1)+,(R3)+   ;TRANSFER RLCS
15      MOV      (R1)+,(R3)+   ;TRANSFER RLBA
16      MOV      (R1)+,(R3)+   ;TRANSFER RLDA
17      MOV      (R1)+,(R3)+   ;TRANSFER RIMP
18      CALL    DLGST          ;GET THE DRIVE STATUS INFO
19      MOV      R1,(R3)+      ; AND SAVE IT FOR ERROR LOGGER
20      COM      R1            ;COMPLEMENT
21      BIT      #STWL,R1      ;Write lock error?
22      BEQ      5$           ;Yes, don't log it
23                                     ; (reversed logic due to COM above)
24      MOV      DLDA,(R3)+     ;SAVE THE DISK ADDRESS THAT WE USED
25      $RLV1B: BR      10$     ;IF NO RLV12...
26                                     ; (CHANGED TO 'NOP' IF USING RLV12)
27      MOV      RLBAE(R4),(R3)+ ;TRANSFER RLBAE
28
29      10$:    MOV      DRETRY,R3
30      SWAB    R3
31      ADD     #DLREG,R3      ;R3= MAX RETRIES/ NUMBER OF REGISTERS
32
33      $RLV1C: BR      20$     ;IF NO RLV12...
34                                     ; (CHANGED TO 'NOP' IF USING RLV12)
35      INC     R3            ;BUMP FOR EXTRA REGISTER ON RLV12
36
37      20$:    JSR      R4,FIXWC          ;GET Q$WCNT SET RIGHT, PUSH OLD VALUE
38      MOV      DLRTY,R4          ;GET NUMBER OF RETRIES LEFT
39      ADD     #DL$COD*400-1,R4    ;SET DEVICE ID FLAG, COUNT=COUNT-1
40                                     ; (report retries remaining, not
41                                     ; current retry number)
42      CALL    @$ELPTR          ;LOG THE ERROR
43      MOV      (SP)+,Q$WCNT(R5) ;RESET WORD COUNT
44      5$:    MOV      DLCSR,R4      ;POINT TO CSR AGAIN
45      MOV      DLRBLK,R3        ;GET RLCS AT TIME OF FAILURE
46      .ENDC ;EQ ERL$G
47
48      001472 012777          MOV      #-1,@DLCC          ;INVALIDATE CURRENT CYLINDER
49                                     ; (FORCE READ HEADER)
50      001500 004767          CALL    DLRST          ;RESET DRIVE
51      001504 105767          TSTB   DLREV          ;REVERSE SEEK REQUIRED?
52      001510 001415          BEQ     6$           ;NO, GO SEE IF WE CAN RETRY
53      001512 105267          51$:  INCB   DLREV          ;SET REVERSE SEEK FLAG IF RETRY
54                                     ; FROM DRIVE N:002
55      001516 012701          MOV     #177600!SKMARK,R1 ;Reverse seek to cylinder zero
56      001522 004767          CALL    DLSEEK        ;EXECUTE THE SEEK
57      001526 100427          BMI    DLEROR        ;SEEK FAILED, CALL IT FATAL
58      001530 016700          MOV     DLUNIT,R0     ;GET UNIT NUMBER TO USE
59      001534 052700          BIS    #CSIE!FNRDH,R0 ;ADD CODE FOR READ HEADER
60      000110
61      001540 010014          MOV     R0,@R4        ;LOAD FUNCTION AND GO
62
63      65      001542 000207          RETURN          ;WAIT FOR THE INTERRUPT
64
65      67      001544 106203          6$:  ASRB   R3          ;AT TIME OF FAILURE, WAS DRIVE READY?
66      001546 103361          BCC    51$          ;NO, REVERSE SEEK UNTIL IT IS
67      001550          .ASSUME CSDRDY EQ 1
68      001550 006303          ASL    R3          ;SHIFT TO GET DRIVE ERROR BIT IN CARRY
69      001552 006303          ASL    R3          ; AND NXM BIT IN SIGN
70      001554 100414          BMI    DLEROR        ;FATAL IF NON-EXISTENT MEMORY
71      001556          .ASSUME CSNXM EQ 020000
72      001556 103010          BCC    7$           ;GO RETRY IF NOT DRIVE ERROR
73      001560          .ASSUME CSDE EQ 040000
74      001560 004767          CALL    DLGST        ;DRIVE ERROR, GO GET DRIVE STATUS
75      001564 032701          BIT    #STWGE,R1    ;WRITE GATE ERROR?
76      002000
77      001570 001406          BEQ    DLEROR        ;FATAL IF NOT

```

```

79 001572 032701      BIT      #STWL,R1              ;YES, WRITE GATE WITH WRITE LOCK?
      020000
80 001576 001003      BNE      DLEROR              ;YES, FATAL
81 001600 005367      7$:    DEC      DLRTY              ;ANY RETRIES LEFT?
      177102
82 001604 003330      BGT      4$                  ;YES, GO DO ONE
83 001606 016705      DLEROR: MOV     DLCQE,R5        ;GET QUEUE ELEMENT POINTER
      176176
84 001612
85 001612 052755      BIS      #HDERR$,@-(R5)      .ASSUME Q$BLKN-2 EQ Q$CSW
      000001              ;FLAG CHANNEL ERROR
86 001616 000416      BR       DLQCOM              ;FINISH-UP
87
88              .DSABL  LSB

1              .SBTTL  FINISH SUCCESSFUL OPERATION
2
3              .ENABL  LSB
4
5 001620 016705      DLEXFR: MOV     DLCQE,R5        ;GET ORIGINAL QUEUE ELEMENT POINTER
      176164
6 001624 020567      CMP      R5,DLQPTR          ;PSEUDO QUEUE IN USE?
      176752

```

Test if we're dismissing a queue element for a replacement table reread or if we're doing a partial transfer using replacement. If a partial transfer, go back and get the rest before we dismiss the original queue element.

```

7 001630 001411      BEQ      1$                  ;NO, THIS IS THE END OF THE REQUEST
8 001632 126527      CMPB     Q$FUNC(R5),#FN$REP  ;WAS FUNCTION A FORCE TABLE RE-READ?
      000002
      000374
9 001640 001405      BEQ      1$                  ;YES, WE ARE NOW DONE
10 001642 005765      TST      Q$WCNT(R5)         ;IS THERE ANYTHING LEFT TO TRANSFER?
      000006
11 001646 001402      BEQ      1$                  ;NOPE, ALL DONE
12 001650 000167      JMP      DLTRAN              ;GO DO NEXT PART OF BROKEN TRANSFER
      176520
13
14 001654          1$:
15              .IF NE  ERL$G
16              JSR      R4,FIXWC          ;FIX WORD COUNT FOR READ/WRITE
17              TST      (SP)+           ;DUMP STACKED OLD VALUE
18              TST      SCSFLG          ;LOGGING SUCCESSES?
19              BNE      DLQCOM          ;NOPE...
20              MOV      #DL$COD*400+377,R4 ;FLAG SUCCESS FOR EL
21              CALL     @$ELPTR          ;CALL THE ERROR LOG HANDLER
22              .ENDC ;NE ERL$G
23

```

I/O Completion Section

Dismiss the queue element.

```

24 001654          DLQCOM: .DRFIN  DL              ;COMPLETE I/O OPERATION
25
26              .DSABL  LSB

1              .SBTTL  GET DEVICE SIZE
2
3              ; SPECIAL FUNCTION TO GET VOLUME SIZE:
4              ; READ THE DRIVE TYPE BIT FOR THE SELECTED DRIVE. THEN RETURN THE
5              ; DRIVE'S SIZE, IN BLOCKS, IN THE FIRST WORD OF THE USER'S BUFFER.
6
7 001672          DLGSIZ:
8              .IF EQ  MMG$T
9              MOV      DLUSIZ,@Q$BUFF(R5) ;PUT SIZE IN BUFFER
10             .IFF
11 001672 016746      MOV      DLUSIZ,-(SP)        ;SET SIZE ON STACK
      176266
12 001676 010504      MOV      R5,R4              ;COPY QUEUE POINTER FOR PUTWORD
13 001700 004777      CALL     @$PTWRD            ;PUT SIZE IN BUFFER
      000404

```

```

14          .ENDC ;EQ MMGST
15 001704 005700      TST      R0          ;Was there an error (no drive?)
16                                     ;R0 should be CSR from DLGST
17 001706                                     .Assume CSEERR EQ 100000
18 001706 100362      BPL      DLQCOM      ;Branch if not
19 001710 032700      BIT      #CSERRC,R0  ;Is there an error code?
20                                     036000
21 001714 001334      BNE      DLEROR      ;Branch if yes
22 001716 032701      BIT      #STVC,R1    ;Is it a volume check error?
23 001722 001731      BEQ      DLEROR      ;If not, report hard error
24 001724 000753      BR       DLQCOM

1          .SBTTL DLXCT - FUNCTION EXECUTION ROUTINES
2
3          ;+
4          ; EXECUTE A GET DRIVE STATUS OR ANY NON-INTERRUPT FUNCTION
5          ; AND WAIT FOR COMPLETION
6          ;
7          ; INPUTS:
8          ;     R1 = FUNCTION CODE          IF DLXCT
9          ;     R4 = SEEK DIFFERENCE WORD  IF DLSEEK
10         ;     R4 -> CSR
11         ;
12         ; OUTPUTS:
13         ;     FUNCTION EXECUTED
14         ;
15         ;     R0 = CSR CONTENTS
16         ;     R1 = MP CONTENTS
17         ;     N = 1 IF ERROR
18         ;-
19
20         .ENABL  LSB
21
22 001726          DLSEEK:
24 001726 010164      MOV      R1,RLDA(R4)  ;LOAD DIFFERENCE WORD IN CONTROLLER
25         000004
28 001732 012701      MOV      #FNSEEK,R1  ;ISSUE SEEK COMMAND
29         000006
30 001736 000424      BR       DLXCT
31 001740          DLGST:
33 001740 012764      MOV      #GSGS!GSMARK,RLDA(R4) ;TELL DRIVE TO GET STATUS
34         000003
35         000004
37 001746 004767      CALL     1$          ;EXECUTE THE GET STATUS
38         000032
39 001752 100026      BPL      4$          ;NO ERROR SO EXIT
40 001754 005764      TST      RLBA(R4)    ;ERROR -- IS IT AFTER BUS INIT?
41         000002
45 001760 001023      BNE      4$          ;NO -- LOG THE ERROR
46 001762 004767      CALL     DLRST      ;YES -- DO A RESET
47         000010
48 001766 012764      MOV      #GSGS!GSMARK,RLDA(R4) ;AND TRY THE GET STATUS AGAIN
49         000003
50         000004
52 001774 000403      BR       1$          ;BUT ONLY TRY IT ONCE!
53
54 001776          DLRST:
56 001776 012764      MOV      #GSRST!GSGS!GSMARK,RLDA(R4) ;GET DRIVE RESET COMMAND
57         000013
58         000004
60 002004 012701 1$:  MOV      #FNGSTS,R1    ;GET 'GET STATUS' FUNCTION CODE
61         000004
62 002010 056701 DLXCT: BIS      DLUNIT,R1  ;ADD IN UNIT SELECT BITS
63         177146
63 002014 010114      MOV      R1,@R4      ;GIVE IT TO DRIVER
64 002016 105714 2$:  TSTB     @R4        ;WAIT FOR FUNCTION TO BE ACCEPTED

```

```

70 002020 100376      BPL      2$
71 002022              3$:
73 002022 016401      MOV      RLMP(R4),R1      ;GET RETURNED STATUS WORD
      000006
74 002026 011400      MOV      @R4,R0      ; AND CSR VALUE (SET N-BIT IF ERROR)
      .
      .
      .
83 002030 000207 4$:  RETURN
84
85                      .DSABL  LSB
      .
      .
      .
44                      .DSABL  LSB
45

```

DLSQUE is used to read the bad-block replacement table into memory and to break up a transfer that uses the table.

```

1                      .SBTTL  DLSQUE - SETUP PSEUDO QUEUE ELEMENT
2
3                      ;+
4                      ; SET UP THE PSEUDO QUEUE FOR BAD BLOCK TABLE READS OR PARTIAL TRANSFERS
5                      ;
6                      ; INPUTS:
7                      ;     R1 = STARTING BLOCK NUMBER OF PARTIAL TRANSFER
8                      ;     R2 = WORD COUNT
9                      ;     R3 -> BUFFER
10                     ;     R5 -> USER QUEUE ELEMENT
11                     ;
12                     ; OUTPUTS:
13                     ;     R0 = RANDOM
14                     ;     R5 -> PSEUDO QUEUE ELEMENT
15                     ;-
16
17 002032      DLSQUE:
18 002032      .ADDR  #DLBWCT,R0      ; POINT TO PSEUDO QUEUE ELEMENT
19 002040 010210      MOV      R2,@R0      ;STORE WORD COUNT
20 002042 010340      MOV      R3,-(R0)      ;STORE BUFFER ADDRESS
21 002044 016540      MOV      Q$FUNC(R5),-(R0)      ;COPY UNIT NUMBER AND
      000002
22
23 002050 010140      MOV      R1,-(R0)      ; SPECIAL FUNCTION BYTE
24 002052      ;STORE BLOCK NUMBER
25 002052 014560      MOV      -(R5),-2(R0)      .ASSUME Q$BLKN-2 EQ Q$CSW
      177776      ;STORE POINTER TO CSW
26 002056 010005      MOV      R0,R5      ;POINT R5 AT PSEUDO QUEUE
27 002060 000207      RETURN
28
29                      .IF NE  ERL$G

1                      .SBTTL  FIXWC - FIX WORD COUNT FOR LOGGER
2
3                      ;+
4                      ; FIX WORD COUNT IN QUEUE ELEMENT FOR ERROR LOGGER
5                      ;
6                      ; INPUTS:
7                      ;     R5 -> QUEUE ELEMENT
8                      ;     DLWC = WORD COUNT USED FOR I/O
9                      ;
10                     ; OUTPUTS:
11                     ;     R4 = RANDOM
12                     ;     @SP = OLD VALUE OF Q$WCNT TO RESTORE
13                     ;     Q$WCNT(R5) = DLWC (NEGATED IF WRITE)
14                     ;-
15
16 FIXWC:  MOV      Q$WCNT(R5),@SP      ;SAVE OLD COUNT ON STACK
17         MOV      DLWC,Q$WCNT(R5)      ;SET THE CORRECT VALUE
18         CMP      DLCODE,#FNWRITE!CSIE      ;WAS IT A WRITE?
19         BNE     1$      ;NO
20         NEG     Q$WCNT(R5)      ;YES, FIX ELEMENT VALUE
21 1$:     JMP      @R4      ;RETURN
22
23                      .ENDC ;NE ERL$G

```

```

1          .SBTTL DATA AREAS
2
3          ; PSEUDO QUEUE ELEMENT
4
5 002062 177777 .WORD -1 ;ADDRESS OF CSW
6 002064 177777 .WORD -1 ;BLOCK NUMBER
7 002066 000 .BYTE 0 ;SPECIAL FUNCTION BYTE
8 002067 377 .BYTE -1 ;UNIT NUMBER
9 002070 177777 DLBADD: .WORD -1 ;BUFFER ADDRESS
10 002072 177777 DLBWCT: .WORD -1 ;WORD COUNT
11
12          .IF NE MMG$T
13 002074 000000 .WORD 0 ;COMPLETION ADDRESS
14 002076 177777 DLBPAR: .WORD -1 ;PAR VALUE
15 002100 177777 DLBMEM: .WORD -1 ;MEM VALUE
16 002102 000000 .WORD 0 ; (RESERVED)
17          .ENDC ;NE MMG$T
18
19
20          ; BAD BLOCK REPLACEMENT TABLE BUFFER AND CURRENT CYLINDER WORD
21          ;
22          ; CONSISTS OF ONE WORD AND ONE TABLE FOR EACH UNIT.
23          ; EACH TABLE CONSISTS OF TWO WORD ENTRIES. WORD 1
24          ; IS BAD BLOCK AND WORD 2 IS IT'S REPLACEMENT. A
25          ; TABLE IS ENDED BY A ZERO ENTRY.
26          ;
27          ; THIS TABLE WILL BE MAPPED INTO HIGH MEMORY WITH UB SUPPORT
28          ;
29

```

This is the bad-block replacement table:

```

30 002104          DLBBUF:
31          000002 .REPT DL$UN ;ONE TABLE PER UNIT
32          .WORD -1 ;CURRENT CYLINDER NUMBER (-1=UNKNOWN)
33          .WORD -1 ;INDICATES TABLE NOT READ YET
34          .BLKB DLTSIZ-2 ;THE TABLE
35          .ENDR
36 002234          BUFEND:
37
38          ; DLFILL ALSO USES THE PERMANENT UMR
39
40 002234 000000 DLFILL: .WORD 0 ;MUST BE 0 TO ZERO-FILL BUFFER
41
42 002236 000000 DLFBLK: .WORD 0,0,0,0 ;FORK QUEUE BLOCK
43 002240 000000
44 002242 000000
45 002244 000000
46
47          .IF NE ERL$G
48          DLRBLK: .BLKW DLREG+1 ;DL STATUS REGISTERS FOR CALL
49          ; TO ERROR LOGGERS (+1 FOR RLBAE)
50          .ENDC ;NE ERL$G
51          .IF NE MMG$T
52          ;+
53          ; DL INTERNAL VARIABLE DEFINITIONS.
54          ;-
55 002246 000000 $ENTPT: .WORD 0 ; POINTER TO $ENTRY TABLE
56 002250 000000 $PNMPT: .WORD 0 ; POINTER TO $PNAME TABLE
57 002252 000000 H2UB: .WORD 0 ; POINTER TO UBVECT
58 002254 000000 DLSLOT: .WORD 0 ; DL'S OFFSET IN DEVICE TABLES
59 002256 000000 DLENT: .WORD 0 ; DL'S $ENTRY TABLE ENTRY POINTER
60 002260 000000 DLPNA: .WORD 0 ; DL'S $PNAME TABLE ENTRY POINTER
61 002262 000000 DLILQE: .WORD 0 ; DL INTERNAL QUEUE LAST QEL POINTER
62 002264 000000 DLICQE: .WORD 0 ; DL INTERNAL QUEUE FIRST QEL POINTER
63
64          ;+
65          ; DEFINITION OF THE HANDLER INTERNAL BUFFER AND THE WORDS THAT ARE
66          ; USED TO PROGRAM DMA DEVICES THAT TRANSFER DATA TO AND FROM IT.
67          ;-
68
69 002266 000000 BUFADH: .WORD 0 ; BITS 0-15 OF UNIBUS VIRTUAL POINTER TO DLBBUF
70 002270 000000 BUFADL: .WORD 0 ; BITS 16-21 OF UNIBUS VIRTUAL POINTER TO DLBBUF

```

```

71
72          ; TABLE OF STANDARD DMA SPFUNS THAT DO NOT HAVE A PERMANENT UMR
73          ; ALLOCATED TO THEM
74
75 002272      UBTAB:  .DRSPF  -, <FN$WRT>          ;ABSOLUTE WRITE, NO BAD BLOCK
76 002274      .DRSPF  -, <FN$RED>          ;ABSOLUTE READ (REPLACEMENT)
77 002276 000000      .WORD    0              ;TABLE TERMINATOR
78
79          .ENDC ;NE MMG$T

```

Bootstrap Driver

```

1          .SBTTL  BOOTSTRAP DRIVER
2

```

The .DRBOT macro:

```

3 002300      .DRBOT  DL,BOOT1,B.READ
177777      ...V7=-1
            .IIF IDN NO,YES,...V7=0

```

Termination Section

The .DREND macro generated by .DRBOT (the macro expansion):

```

001770      .DREND  DL,0,
            .IF B <>
001770      .PSECT  DLDVR
            .IFF
            .PSECT
            .ENDC
            .IIF NDF DL$END,DL$END::
            .IF EQ  .-DL$END
            .IF NE MMG$T!<0&2.>
            $RLPTR:: .WORD  0
            $MPPTR:: .WORD  0
            $GTBYT:: .WORD  0
            $PTBYT:: .WORD  0
            $PTWRD:: .WORD  0
            .ENDC
            .IF NE ERL$G!<0&1>
            $ELPTR:: .WORD  0
            .ENDC
            .IF NE TIM$IT!<0&4.>
            $TIMIT:: .WORD  0
            .ENDC
001770 000000  $INPTR:: .WORD  0
001772 000000  $FKPTR:: .WORD  0
            .IIF NDF ...V22 ...V22=0
            .IF NE ...V22&^o40000
            DL$X64 =:
            .REPT  16.
                .WORD  0
            .ENDR
            .ENDC
            .GLOBL  DLSTRT

```

The following line marks the end of the loadable portion of the handler. It is used to determine the handler's length in memory.

```

001774' DLEND==.
            .IFF
            .PSECT  DLBOOT
            .IIF LT <DLBOOT-.^o664>,.ERROR;?SYSMAC-E-Primary boot too large;
            .=DLBOOT+^o664
            BIOERR: JSR    R1,REPORT
                .WORD  IOERR-DLBOOT
            REPORT: MOV    #BOOTF-DLBOOT,R0
                MOV    #30002$-DLBOOT,R2
                CALL   @R2
                MOV    @R1,R0
                CALL   @R2
                MOV    #CRLF-LF-DLBOOT,R0
                CALL   @R2
            30001$: HALT

```

```

BR      30001$
30002$: TSTB  @#TPS
        BPL   30002$
        MOVB  (R0)+,@#TPB
        BNE   30002$
        RETURN
BOOTF:  .ASCIZ <CR><LF>"?BOOT-U-"
IOERR:  .ASCII  "I/O error"
CRLFLF: .ASCIZ <CR><LF><LF>
        .EVEN
        .IIF NDF ...V7,...V7=-1
        .REPT 4.
        .WORD ...V7
        .ENDR
DLBEND::
        .ENDC
        .IIF NDF TPS,TPS=:^o177564
        .IIF NDF TPB,TPB=:^o177566
000012  LF=:^o12
000015  CR=:^o15
001000  B$BOOT=:^o1000
004716  B$DEVN=:^o4716
004722  B$DEVU=:^o4722
004730  B$READ=:^o4730
        .IF NDF B$DNAM
        .IF EQ MMG$T
        B$DNAM=^RDL
        .IFF
        B$DNAM=^RDLX
        .ENDC ; EQ MMG$T
        .ENDC ; NDF B$DNAM
000062  .ASECT
        .=^o62
000062  000000' .WORD DLBOOT,DLBEND-DLBOOT,B.READ-DLBOOT
000064  001000
000066  000210
000000  .PSECT DLBOOT
000000  000240 DLBOOT::NOP
000002  000415 BR BOOT1-2.
000100  ...V2=^o100
        .IRP X <UBUS,QBUS>
        ...V3=0
        .IIF IDN <X> <UBUS> ...V3=1.
        .IIF IDN <X> <QBUS> ...V3=2.
        .IIF IDN <X> <CBUS> ...V3=4.
        .IIF IDN <X> <UMSCP> ...V3=^o10
        .IIF IDN <X> <QMSCP> ...V3=^o20
        .IIF IDN <X> <CMSCP> ...V3=^o40
        .IIF EQ ...V3 .ERROR;?SYSMAC-E-Invalid C O N T R O L, found - UBUS,QBUS;
        ...V2=...V2!...V3
        .ENDR
        ...V3=0
000000  ...V3=0
000001  .IIF IDN <UBUS> <UBUS> ...V3=1.
        .IIF IDN <UBUS> <QBUS> ...V3=2.
        .IIF IDN <UBUS> <CBUS> ...V3=4.
        .IIF IDN <UBUS> <UMSCP> ...V3=^o10
        .IIF IDN <UBUS> <QMSCP> ...V3=^o20
        .IIF IDN <UBUS> <CMSCP> ...V3=^o40
        .IIF EQ ...V3 .ERROR;?SYSMAC-E-Invalid C O N T R O L, found - UBUS,QBUS;
000101  ...V2=...V2!...V3
000000  ...V3=0
        .IIF IDN <QBUS> <UBUS> ...V3=1.
000002  .IIF IDN <QBUS> <QBUS> ...V3=2.
        .IIF IDN <QBUS> <CBUS> ...V3=4.
        .IIF IDN <QBUS> <UMSCP> ...V3=^o10
        .IIF IDN <QBUS> <QMSCP> ...V3=^o20
        .IIF IDN <QBUS> <CMSCP> ...V3=^o40
        .IIF EQ ...V3 .ERROR;?SYSMAC-E-Invalid C O N T R O L, found - UBUS,QBUS;
000103  ...V2=...V2!...V3
000032' .=BOOT1-6.
000032  020 .BYTE ^o20,...V2,^o20,^o^C<20+...V2+20>
000033  103
000034  020
000035  234
        .IF EQ <1-1>
000036  000400 BR BOOT1
        .IFF
        .IF EQ <1-2.>

```

```

                                BMI      BOOT1
                                .IFB
                                .ERROR;?SYSMAC-E-Invalid S I D E S, expecting 1/2, found - 1;
                                .ENDC
                                .ENDC
4
5                                000040' . = DLBOOT+40                                ;PUT THE JUMP BOOT INTO SYSCOM AREA
6 000040 000137 BOOT1: JMP      @#BOOT-DLBOOT                                ;START THE BOOTSTRAP
                                000600

1                                .SBTTL  BOOTSTRAP READ ROUTINE
2
3                                .ENABL  LSB
4
5                                000210' . = DLBOOT+210
6 000210 005004 B.READ: CLR      R4                                ;CLEAR TRACK COUNTER
7 000212 162700 1$:  SUB      #DLBPT,R0                                ;COUNT DOWN ANOTHER WHOLE TRACK
                                000024
8 000216 103403      BLO      2$                                ;IF OVERFLOW, DONE
9 000220 062704      ADD      #IOHS,R4                                ;ADD IN ANOTHER TRACK
                                000100
10 000224 000772      BR       1$                                ;LOOP FOR MORE
11
12 000226 062700 2$:  ADD      #DLBPT,R0                                ;CORRECT TRACK COUNTER
                                000024
13 000232 006300      ASL      R0                                ;CONVERT REMAINDER TO SECTOR IN TRACK
14 000234 050400      BIS      R4,R0                                ;MERGE SECTOR WITH TRACK/CYL
16 000236 016705      MOV      BOTCSR,R5                                ;GET ADDRESS OF CONTROLLER
                                000344
17 000242 062705      ADD      #RLDA,R5                                ;POINT TO DISK ADDRESS REGISTER
                                000004
18 000246 016567      MOV      RLCS-RLDA(R5),B.DLCS                    ;GET CURRENT CSR VALUE
                                177774
                                000174
19 000254 042767      BIC      #^C<CSDS01>,B.DLCS                    ;ISOLATE CURRENT UNIT NUMBER
                                176377
                                000166
21 000262 004767      CALL     B.SEEK                                ;SEEK TO PROPER TRACK
                                000066
22 000266 005401      NEG      R1                                ;NEGATE WORD COUNT
24 000270 010265      MOV      R2,RLBA-RLDA(R5)                            ;SET BUS ADDRESS
                                177776
.
.
.
28 000274      DLREAD:
30 000274 010165      MOV      R1,RLMP-RLDA(R5)                            ;SET WORD COUNT
                                000002
31 000300 010015      MOV      R0,@R5                                ;SET DISK ADDRESS
.
.
.
36 000302 004067      JSR      R0,B.XCT                                ;EXECUTE THE READ
                                000136
37 000306 000014      .WORD   FNREAD                                ;READ FUNCTION CODE
38 000310 000241      CLC                                ;ENSURE CARRY=0 BEFORE RETURN
39 000312 100053      BPL      5$                                ;SUCCESS, EXIT
41 000314 011503      MOV      @R5,R3                                ;GET LAST DISK ADDRESS
.
.
.
46 000316 042703      BIC      #^C<IOSA>,R3                                ;CLEAR ALL BUT SECTOR ADDRESS
                                177700
47 000322 022703      CMP      #DLBPT*2,R3                                ;TRACK OVERRUN?
                                000050
48 000326 001156      BNE      BIOERR                                ;IF NOT, REAL ERROR, EXIT

```

```

50 000330 011503      MOV      @R5,R3          ;GET DISK ADDRESS
      .
      .
54 000332 160003      SUB      R0,R3          ;COMPUTE SECTORS TRANSFERRED
55 000334 000303      SWAB    R3              ;CONVERT SECTORS TO WORD COUNT
56 000336 006203      ASR     R3              ;
57 000340 060301      ADD     R3,R1          ;REMOVE WORDS TRANSFERRED
59 000342 011500      MOV     @R5,R0        ;GET DISK ADDRESS
      .
      .
63 000344 062700      ADD     #IOHS-<DLBPT*2>,R0 ;INCREMENT SURFACE/TRACK
      000030
64 000350 012746      MOV     #DLREAD-DLBOOT,-(SP) ;CALL TO SEEK NEXT TRACK, THEN READ IT
      000274
65 000354              .BR     B.SEEK        ;SEEK NOW
66
67 000354 004067 B.SEEK: JSR     R0,B.XCT ;EXECUTE READ HEADERS
      000064
68 000360 000010      .WORD  FNRDH          ;READ HEADER FUNCTION CODE
70 000362 016503      MOV     RLMP-RLDA(R5),R3 ;GET CURRENT DISK TRACK AND SURFACE
      000002
      .
      .
74 000366 042703      BIC     #IOHS!IOSA,R3   ;CLEAR SURFACE/SECTOR TO GET
      000177
75
76 000372 010004      MOV     R0,R4          ;CURRENT TRACK
77 000374 042704      BIC     #IOHS!IOSA,R4   ;COPY DESIRED DISK ADDRESS
      000177
78
79 000400 160403      SUB     R4,R3          ;CLEAR SURFACE/SECTOR TO GET
80 000402 103003      BCC    3$             ; DESIRED TRACK
81
82 000404 005403      NEG     R3             ;SUBTRACT DESIRED FROM CURRENT TRACK
83
84 000406 052703      BIS     #SKDIR,R3      ;IF CURRENT >= DESIRED,
      000004 ; SEEK OUTWARD BY DIFF
85 000412 032700 3$: BIT     #IOHS,R0      ;MAKE POSITIVE DIFFERENCE OF
      000100 ; DELTA POSITION
86 000416 001402      BEQ    4$             ;INDICATE MOVE TOWARD SPINDLE
87 000420 052703      BIS     #SKHS,R3      ;DO WE DESIRE SURFACE 1?
      000020 ; NO, LEAVE SURFACE SELECT 0
88 000424 005203 4$: INC     R3             ;SET BIT TO SELECT SURFACE 1
89
90 000426 010315      MOV     R3,@R5        ;SET MARKER BIT
      .
      .
94 000430 004067      JSR     R0,B.XCT      ;LOAD DIFFERENCE WORD
      000010 ; EXECUTE A SEEK
95 000434 000006      .WORD  FNSEEK        ;EXECUTE THE FUNCTION IN R3 AND RETURN ERROR STATUS
96 000436 100512      BMI    BIOERR        ;SEEK FUNCTION CODE
      .
98 000440 010015      MOV     R0,@R5        ;IF PL, OK
      .
102 000442 000207 5$: RETURN ;SET ACTUAL DISK ADDRESS
103
104
105
106
107 000444 012003 B.XCT: MOV     (R0)+,R3 ;RETURN
      .
      .

```

```

109 000446 052703      BIS      (PC)+,R3          ;ADD UNIT BITS TO FUNCTION CODE
110 000450 000000 B.DLCS: .WORD 0          ;BOOTED UNIT NUMBER
111 000452 010365      MOV      R3,RLCS-RLDA(R5) ;EXECUTE FUNCTION
      177774
112 000456 032765 6$:      BIT      #CSERR!CSCRDY,RLCS-RLDA(R5) ;WAIT FOR COMPLETION OR ERROR
      100200
      177774
      .
      .
118 000464 001774      BEQ      6$              ;NEITHER, LOOP
119 000466 000200      RTS      R0              ;RETURN WITH N=1 IF ERROR
      .
      .
      . = DLBOOT+600
6      000600'
8      000600 012706 BOOT:    MOV      #10000,SP          ;SET STACK POINT
      010000
10     000604 013746      MOV      @(PC)+,-(SP)
11     000606 174400 BOTCSR: .WORD DL$CSR
12     000610 042716      BIC      #^C1400,@SP      ;STRIP TO UNIT NUMBER
      176377
13     000614 000316      SWAB     @SP              ;MOVE TO BITS 0-1
15     000616 012700      MOV      #2,R0            ;READ IN SECOND PART OF BOOT
      000002
16     000622 012701      MOV      #4*256.,R1       ;FOUR BLOCKS TO READ
      002000
17     000626 012702      MOV      #1000,R2         ;INTO LOCATION 1000
      001000
18     000632 004767      CALL     B.READ           ;READ THE REST OF THE BOOT
      177352
19     000636 012737      MOV      #B.READ-DLBOOT,@#B$READ ;STORE START LOCATION OF READ ROUTINE
      000210
      004730
20     000644 012737      MOV      #B$DNAM,@#B$DEVN ;STORE RAD50 DEVICE NAME
      015370
      004716
22     000652 012637      MOV      (SP)+,@#B$DEVU   ;SET THE UNIT NUMBER IN THE BOOT
      004722
      .
      .
      .
26     000656 000137      JMP      @#B$BOOT         ;GO DO THE BOOT WORK
      001000
27
28     000662      .DREND DL
29
      .IF B <>
001774 .PSECT DLDVR
      .IFF
      .PSECT
      .ENDC
      .IF NDF DL$END,DL$END::
      .IF EQ -.DL$END
      .IF NE MMG$T!<0&2.>
$RLPTR: .WORD 0
$MPPTR: .WORD 0
$GTBYT: .WORD 0
$PTBYT: .WORD 0
$PTWRD: .WORD 0
      .ENDC
      .IF NE ERL$G!<0&1.>
$ELPTR: .WORD 0
      .ENDC
      .IF NE TIM$IT!<0&4.>
$TIMIT: .WORD 0
      .ENDC
$INPTR: .WORD 0
$FKPTR: .WORD 0
      .IF NDF ...V22 ...V22=0
      .IF NE ...V22&^o40000
DL$X64 =: .
      .REPT 16.
      .WORD 0
      .ENDR

```

```

        .ENDC
        .GLOBL DLSTRT
        DLEND==.
        .IFF
000662      .PSECT DLBOOT
        .IIF LT <DLBOOT-.^o664>,.ERROR;?SYSMAC-E-Primary boot too large;
        .=DLBOOT+^o664
000664      000664' 004167      BIOERR: JSR      R1,REPORT
        000002
000670      000753      .WORD      IOERR-DLBOOT
000672      012700      REPORT: MOV      #BOOTF-DLBOOT,R0
        000740
000676      012702      MOV      #30004$-DLBOOT,R2
        000722
000702      004712      CALL     @R2
000704      011100      MOV     @R1,R0
000706      004712      CALL     @R2
000710      012700      MOV     #CRLFLF-DLBOOT,R0
        000764
000714      004712      CALL     @R2
000716      000000      30003$: HALT
000720      000776      BR      30003$
000722      105737      30004$: TSTB  @#TPS
        177564
000726      100375      BPL     30004$
000730      112037      MOVB   (R0)+,@#TPB
        177566
000734      001372      BNE     30004$
000736      000207      RETURN
000740      015      BOOTF:  .ASCIZ  <CR><LF>"?BOOT-U-"
000741      012
000742      077
000743      102
000744      117
000745      117
000746      124
000747      055
000750      125
000751      055
000752      000
000753      111      IOERR:  .ASCII  "I/O error"
000754      057
000755      117
000756      040
000757      145
000760      162
000761      162
000762      157
000763      162
000764      015      CRLFLF: .ASCIZ  <CR><LF><LF>
000765      012
000766      012
000767      000

        .EVEN
        .IIF NDF ...V7,...V7=-1
000004      .REPT  4.
        .WORD  ...V7
        .ENDR
000770      177777      .WORD  ...V7
000772      177777      .WORD  ..V7
000774      177777      .WORD  ...V7
000776      177777      .WORD  ...V7
001000      DLBEND::
        .ENDC
31          .IF NE MMG$T

```

```

1          .SBTTL  FETCH/LOAD CODE
2          ;+
3          ;      FETCH
4          ;
5          ;      ENTRY:  R0 = STARTING ADDRESS OF THIS HANDLER SERVICE ROUTINE.
6          ;                  R1 = ADDRESS OF GETVEC ROUTINE.
7          ;                  R2 = VALUE $SLOT*2. (LENGTH OF THE $PNAME TABLE IN BYTES.)
8          ;                  R3 = TYPE OF ENTRY.
9          ;                  R4 = ADDRESS OF SY READ ROUTINE.
10         ;                  R5 -> $ENTRY SLOT FOR THIS HANDLER.
11         ;
12         ; -
13
14
15 001000 010501  FETCH:  MOV     R5,R1          ; SAVE PTR TO DL'S $ENTRY SLOT
16 001002 011505          MOV     @R5,R5          ; GET ADDRESS OF DLLQE
17 001004 016504          MOV     DLCSR-DLLQE(R5),R4      ; GET CSR FOR DL
18         000024
19 001010 005046          CLR     -(SP)          ; SPACE FOR RETURN VALUE
20 001012          .MFPS          ; GET PROCESSOR STATUS
21 001024          .MTPS    #340          ; RAISE PROCESSOR PRIORITY LEVEL TO 7
22 001044 013746          MOV     @#NXM.V+2,-(SP)      ; ;SAVE CURRENT NXM TRAP PSW
23         000006
24 001050 013746          MOV     @#NXM.V,-(SP)      ; ;SAVE CURRENT NXM TRAP VECTOR
25         000004
26 001054          .ADDR   #DLNXM,-(SP)      ; ;BUILD ADDRESS TO OUR TRAP ROUTINE
27 001062 012637          MOV     (SP)+,@#NXM.V      ; ;SET UP THE NXM VECTOR
28         000004
29 001066 012737          MOV     #340,@#NXM.V+2     ; ;SET UP THE NXM PSW
30         000340
31         000006
32 001074 005764          TST     RLBAE(R4)         ; ;BAE REGISTER EXIST?
33         000010
34 001100 012637          MOV     (SP)+,@#NXM.V      ; ;MAYBE, FIRST RESTORE NXM VECTOR
35         000004
36 001104 012637          MOV     (SP)+,@#NXM.V+2     ; ;      AND NXM PSW
37         000006
38 001110 006166          ROL     2(SP)             ; ;SAVE THE CARRY BIT
39         000002
40 001114          .MTPS          ; RESTORE PREVIOUS PRIORITY LEVEL
41 001126 006026          ROR     (SP)+             ; RESTORE CARRY
42 001130 103403          BCS     20$              ; NOT AN RLV12... BR IS NOT AN ERROR
43
44
45         .IF EQ ERL$G
46 001132 012765          MOV     #NOP,$RLV1A-DLLQE(R5) ; Change BR to NOP for RLV12
47         000240
48         001156
49
50         .IFF
51 001140          MOV     #NOP,R0           ; R0="NOP"
52 001142          MOV     R0,$RLV1A-DLLQE(R5) ; PATCH 'BR' TO 'NOP' FOR RLV12
53 001144          MOV     R0,$RLV1B-DLLQE(R5) ; PATCH ERROR LOGGING CODE SO IT KNOWS
54 001146          MOV     R0,$RLV1C-DLLQE(R5) ; ABOUT EXTRA REGISTER (RLBAE)
55         .ENDC ;EQ ERL$G
56
57 001140 000241 20$:  CLC
58
59         ;+
60
61         ;      LOAD UP LOCAL VARIABLES WITHIN DL
62         ; -
63
64 001142 010105          MOV     R1,R5           ; RESTORE PTR TO DLLQE TO R5
65 001144 011100          MOV     @R1,R0          ; GET ADDRESS OF DLLQE
66 001146 013704          MOV     @#SYSPTR,R4     ; GET START OF RMON
67         000054
68 001152 016403          MOV     $H2UB(R4),R3    ; R3 = UBVECT POINTER
69         000460
70 001156 010360          MOV     R3,<H2UB-DLBASE>(R0) ; H2UB = ADDRESS OF UBVECT
71         002244
72 001162 016403          MOV     $PNPTR(R4),R3   ; R3 = RMON OFFSET TO PNAME TABLE
73         000404
74 001166 060403          ADD     R4,R3           ; R3 -> PNAME TABLE ADDRESS
75 001170 010360          MOV     R3,<$PNMPT-DLBASE>(R0) ; $PNMPT -> PNAME TABLE ADDRESS
76         002242
77 001174 060203          ADD     R2,R3           ; R3 -> $ENTRY TABLE
78 001176 010360          MOV     R3,<$ENTPT-DLBASE>(R0) ; $ENTPT -> $ENTRY TABLE
79         002240
80 001202 010560          MOV     R5,<DLENT-DLBASE>(R0) ; DLENT -> DL'S $ENTRY TABLE ENTRY

```

```

002250
60 001206 160205      SUB    R2,R5          ; R5 -> DL'S $PNAME TABLE ENTRY
61 001210 010560      MOV    R5,<DLPNA-DLBASE>(R0) ; DLPNA -> DL'S $PNAME TABLE ENTRY
002252
62

```

Allocate permanent UMRs if using UNIBUS mapping registers:

```

63          ;+
64          ;          ALLOCATE PERMANENT UMRs TO POINT INTO DL'S INTERNAL DMA BUFFERS,
65          ;          DLBBUF AND DLFILL, AND GET THE UNIBUS VIRTUAL ADDRESS.
66          ;-
67
68 001214 012701      MOV    #<DLBBUF-DLBASE>,R1    ; R1 = LOW 16 BITS OF DMABUF ADDRESS
002076
69 001220 060001      ADD    R0,R1              ;
70 001222 005002      CLR    R2                ; R2 = HIGH 6 BITS OF DMABUF ADDRESS
71 001224 016004      MOV    <DLPNA-DLBASE>(R0),R4 ; R4 -> PNAME ENTRY FOR DL
002252
72 001230 016005      MOV    <H2UB-DLBASE>(R0),R5  ; GET UB ENTRY ADDRESS
002244
73 001234 010046      MOV    R0,-(SP)          ; SAVE DL STARTING ADDRESS
74 001236 012700      MOV    #NOUMRS,R0        ; R0 = NUMBER OF UMRS REQUIRED
000001
75 001242 004765      CALL   UB.ALL(R5)        ; CALL ALLUMR
000002
76 001246 012600      MOV    (SP)+,R0          ; RESTORE DL STARTING ADDRESS
77 001250 103411      BCS   30$              ; COULDN'T GET UMR, FAIL THE LOAD
78 001252 010160      MOV    R1,<BUFADL-DLBASE>(R0) ; STORE UNIBUS VIRTUAL ADDRESS LOW
002262
79 001256 006302      ASL   R2                ; SHIFT
80 001260 006302      ASL   R2                ; HI BITS LEFT 4
81 001262 006302      ASL   R2                ; TO GET THEM INTO THE
82 001264 006302      ASL   R2                ; CORRECT PLACE
83 001266 010260      MOV    R2,<BUFADH-DLBASE>(R0) ; STORE UNIBUS VIRTUAL ADDRESS HIGH
002260
84 001272 000241      CLC                    ; LOAD SUCCEEDED
85 001274 000207      30$: RETURN
86
87
88          ;
89
90 001276 052766      $DLNXM: BIS    #1,2(SP)          ;SET THE CARRY BIT
000001
000002
91 001304 000002      RTI
92

```

Routine to unload DL and release any UMRs:

```

1          ;+
2          ;          RELEAS
3          ;
4          ;          ROUTINE TO UNLOAD DL
5          ;
6          ;          ENTRY:  SAME AS FOR LOAD.
7          ;
8          ;-
9
10         .ENABL  LSB
11 001306      RELEAS::
12 001306 010501      MOV    R5,R1          ; R1 = $ENTRY SLOT FOR DM
13 001310 160201      SUB    R2,R1          ; R2 -> $PNAME SLOT FOR DM
14 001312 013704      MOV    @#SYSPTTR,R4    ; GET START OF RMON
000054
15 001316 016405      MOV    $H2UB(R4),R5    ; R5 = UB ENTRY VECTOR
000460
16 001322 004765      CALL   UB.RLS(R5)      ; RELEASE UMRS
000004
17 001326 000207      RETURN          ; AND EXIT
18
19         .ENDC ;NE MMG$T
21
22         000001      .END

```

Symbol Table From Assembly

ABTIO\$	001000		DLBWCT	002072R	002	DOC\$UN=	000000	
BAREA	000606		DLCC	000154R		DRETRY=	000742R	002
BIOERR	000664R	003	DLCODE	001156R	002	DVC.CT	000006	
BOOT	000600R	003	DLCQE	000010RG	002	DVC.DE	000010	
BOOTF	000740R	003	DLCSR	000032R	002	DVC.DK	000004	
BOOT1	000040R	003	DLDA	001150R	002	DVC.DL	000012	
BOTCSR	000606R	003	DLDSIZ	023742		DVC.DP	000011	
BTCSR =	004124		DLELNK	000370R	002	DVC.LP	000007	
BUFADH	002266R	002	DLEND =	002316RG	002	DVC.MT	000005	
BUFADL	002270R	002	DLENT	002256R	002	DVC.NI	000013	
BUFEND	002234R	002	DLERJM	001012R	002	DVC.NL	000001	
BUFSIZ	000054		DLEROR	001606R	002	DVC.PS	000014	
BUS\$	000100		DLERRH	001470R	002	DVC.SB	000020	
BUS\$C	020000		DLEXFR	001620R	002	DVC.SI	000016	
BUS\$M	020100		DLFBLK	002236R	002	DVC.SO	000017	
BUS\$Q	000100		DLFILL	002234R	002	DVC.TP	000003	
BUS\$U	000000		DLFLNK	000250R	002	DVC.TT	000002	
BUS\$X	020100		DLGSIZ	001672R	002	DVC.UK	000000	
B\$BOOT	001000		DLGST	001740R	002	DVC.VT	000015	
B\$DEVN	004716		DLICQE	002264R	002	DVM.DM	000002	
B\$DEVU	004722		DLILQE	002262R	002	DVM.DX	000001	
B\$DNAM	015370		DLINT	001250RG	002	DVM.NF	000200	
B\$READ	004730		DLLQE	000006RG	002	DVM.NS	000001	
B.DLCS	000450R	003	DLNBAD	000012		DV2.V2	040000	
B.READ	000210R	003	DLPNA	002260R	002	EIS\$I =	000001	
B.SEEK	000354R	003	DLQCOM	001654R	002	EOF\$	020000	
B.XCT	000444R	003	DLQPTR	000602R	002	ERL\$G =	000000	
CONFG2	000370		DLRCNT	000010		FETCH	001000R	003
CR	000015		DLREAD	000274R	003	FILST\$	100000	
CRLFLF	000764R	003	DLREG	000006		FIX\$ED=	000001	
CSBA16	000020		DLREV	000750R	002	FNGSTS	000004	
CSBA17	000040		DLRST	001776R	002	FNNOP	000000	
CSCRDY	000200		DLRTY	000706R	002	FNRDH	000010	
CSDCRC	004000		DLSEEK	001726R	002	FNRDNH	000016	
CSDE	040000		DLSIZE	023742		FNREAD	000014	
CSDLT	010000		DLsiz2	047742		FNSEEK	000006	
CSDRDY	000001		DLslot	002254R	002	FNWCHK	000002	
CSDS0	000400		DLsqUE	002032R	002	FNWRIT	000012	
CSDS01	001400		DLSTRT	000000RG	002	FN\$RED	000377	
CSERR	100000		DLSTS	102405		FN\$REP	000374	
CSERRC	036000		DLsys	000006RG	002	FN\$SIZ	000373	
CSFUN	000016		DLTRAK	000746R	002	FN\$WRT	000376	
CSHCRC	004000		DLTRAN	000374R	002	GSgs	000002	
CSHNF	010000		DLTsiZ	000052		GSMARK	000001	
CSIE	000100		DLUNIT	001162R	002	GSRST	000010	
CSNXM	020000		DLUSIZ	000164R	002	HDERR\$	000001	
CSOPI	002000		DLWC	001140R	002	HNDLR\$	004000	
DISCSR	000174		DLWPT	012000		HS2.BI	000001	
DLADDR	000600R	002	DLWTRK	001126R	002	HS2.KI	000002	
DLBADD	002070R	002	DLXCT	002010R	002	HS2.KL	000004	
DLBASE=	000006R	002	DLXFER	001114R	002	HS2.KU	000010	
DLBBUF	002104R	002	DL\$COD	000005		HS2.MO	000020	
DLBEND	001000RG	003	DL\$CSR=	174400 G		H2UB	002252R	002
DLBEMEM	002100R	002	DL\$END	002300RG	002	INSCSR	000176	
DLBOOT	000000RG	003	DL\$NAM=	015340		INSDAT	000200	
DLBPAR	002076R	002	DL\$UN =	000002 G		INSSYS	000202	
DLBPT	000024		DL\$VEC=	000160 G		IOCA	077600	
IOCA0	000200		RELEAS	001306RG	003	\$ENTPT	002246R	002
IOERR	000753R	003	REPORT	000672R	003	\$FKPTR	002314RG	

Figure A-3: XL Communications Handler

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RESSTT - Reset line state bits
SETSTT - Set line state bits
GETC - Input a character
PUTC - Output a character
INPUT BUFFER AREA
LOAD - Handler FETCH/LOAD code
UNLOAD - UNLOAD/.RELEASE CODE

```
1                   .MCALL .MODULE
2 000000            .MODULE XL,VERSION=36,COMMENT=<Communications Driver>
3
4                   ;
5                   ;                    COPYRIGHT 1989, 1990, 1991 BY
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22
23                    .ENABL LC
24
25                   ;+
26                   ;
27                   ; FACILITY:        RT-11 Device driver
28                   ;
```

```

9          ; FUNCTIONAL DESCRIPTION:
10         ;
11         ; This driver aids in the writing of virtual terminal software. It
12         ; supports the XON/XOFF protocol in that if receives too many chars
13         ; it will transmit a CTRL/S and send a CTRL/Q when it again has room.
14         ; It will also stop transmitting if it receives a CTRL/S and resume
15         ; on a CTRL/Q. Normal RT-11 READ/WRITE commands can be done to the
16         ; plus various special functions. On any data transfer, chars are
17         ; striped to seven bits and chars of value zero are ignored. On output
18         ; the character following a carriage return is not output.
19         ;

```

CONDITIONAL ASSEMBLY SUMMARY

```

1          .SBTTL Conditional assembly summary
2
3
4          000001 XL$LUN = 1
5          000001 XL$MTY = 1
6          000001 XL$PDP = 1
7          ;
8          ;
9          000001 MMG$T = 1
10         000001 TIM$IT = 1
11         ;+
12         ;COND
13         ;
14         ; XL$DVE (0) support for DLV11E
15         ; 0 no support
16         ; 1 support
17         ;
18         ; XL$PC (0) support for PRO300 series
19         ; 0 no support
20         ; 1 support
21         ;
22         ; XL$SBC (0) support for SBC-11/21[+] and MXV SLUs
23         ; 0 no support
24         ; 1 support
25         ;
26         ; Exactly one of XL$PC, XL$DVE and XL$SBC
27         ; may be specified.
28         ;
29         ; XL$PDT (0) support PDT lights
30         ; 0 no support
31         ; 1 support
32         ;
33         ; XL$PDT is ignored if XL$PC is 1
34         ;
35         ; XL$PRI (4) interrupt priority
36         ; (5) if XL$SBC is 1
37         ; 4-7 range
38         ;
39         ; XL$CSR (176500) CSR address
40         ; (173300) if XL$PC is 1
41         ;
42         ; XL$VEC (300) Vector address
43         ; (210) if XL$PC is 1
44         ;
45         ; XL$MTY (0) No support multiterminal handler hooks
46         ; 1 Support for multiterminal handler hooks
47         ;
48         ; XL$MTY may be 1 only when XL$PC is 0.
49         ;
50         ; XL$LUN (1) Line number to use in multiterminal
51         ;
52         ; MMG$T std conditional
53         ; TIM$IT std conditional (no code effect)
54         ; ERL$G std conditional (no code effect)
55         ;
56         ;-

```

MACROS AND DEFINITIONS

```

1          .SBTTL  MACROS AND DEFINITIONS
2
3          .LIBRARY      "SRC:SYSTEM.MLB"
4

```

Prepare for using standard definitions:

```

5          ; Declare the RT system macros we'll be using
6
7          .MCALL .DRDEF .MTPS .INTEN
8          .MCALL .ASSUM .ADDR .BR
9          .MCALL .MTSTA
10
11         ; Define and verify some conditionals
12
13         .IIF NDF XL$DVE XL$DVE = 0      ;Default to non DLV11-E interface
14
15         .IIF NDF XL$PC XL$PC = 0      ;Default to non PRO-3xx support
16
17         .IIF NDF XL$PDT XL$PDT = 0    ;Default to no PDT lights display
18
19         .IIF NDF XL$SBC XL$SBC = 0    ;Default to non SBC-11 interface
20
21         .IIF NDF XL$MTY XL$MTY = 0    ;Default to no support for MTY hooks
22         .IIF NDF XL$LUN XL$LUN = 1    ;Default to LUN 1
23
24 000000      .Assume <XL$PC & <XL$DVE ! XL$SBC>> EQ 0 MESSAGE=<Conflicting options>
25 000000      .Assume <XL$DVE + XL$SBC> LE 1 MESSAGE=<Conflicting options>
26 000000      .Assume <XL$PC & XL$MTY> EQ 0 MESSAGE=<Conflicting options>
27
28         ; Set the audit trail
29
30         000000      .XLGEN = XL$PC ! <XL$DVE * 2> !
<XL$SBC * 4> ! <XL$PDT * 10>
31         000020      .XLGEN = .XLGEN ! <XL$MTY * 20>
32
33 000000      .AUDIT .XL                      ;The handler
34 000110 107123      .WORD .AUDIT
35 000112 000044      .WORD .XL
36 000114 177777      .WORD -1
37 000000      .AUDIT .XLGEN                  ; and the conditionals
38 000114 000020      .WORD .XLGEN
39 000116 177777      .WORD -1
40
41         ; Define the device
42         ; o Entered on all aborts
43         ; o handles .SPFUN system call
44
45         .IF NE XL$PC
46         XL$CSR = 173300                    ;Force these for a PRO
47         XL$VEC = 210
48         .ENDC ;NE XL$PC
49
50         .IIF NDF XL$PRI XL$PRI = 4      ;Interrupt processing level
51
52         .IF NE XL$SBC
53         XL$PRI = 5                        ;Force this for SBC-11/12[+] and MXV
54         .ENDC ;NE XL$SBC
55

```

The .DRDEF macro with expansion:

```

51 000000      .DRDEF XL,57,<ABTIO$!HNDLR$!SPFUN$>,0,176500,300,DMA=NO
52 000100 000040      .WORD 40
53 000176 176500      .WORD XL$CSR
54
55         .IF EQ XL$PC
56

```

The .DRPTR macro with expansion:

```

54 000200          .DRPTR  FETCH=FETCH,LOAD=LOAD,UNLOAD=UNLOAD,RELEAS=RELEAS
    000000 031066  .RAD50   "HAN"
    000002 002704'  .WORD   FETCH
    000004 003312'  .WORD   RELEAS
    000006 002704'  .WORD   LOAD
    000010 003256'  .WORD   UNLOAD
    000021    000   .BYTE   0
55                    .IFF ;EQ XL$PC
56                    .DRPTR  UNLOAD=UNLOAD,RELEAS=RELEAS
57                    .ENDC ;EQ XL$PC
58

```

The .DREST macro with expansion:

```

59 000022          .DREST  CLASS=DVC.VT
    000000 031066  .RAD50   "HAN"
    000020    015  .BYTE   DVC.VT
    000021    000  .BYTE   0
    000032 000000  .WORD   0
    000036 000000  .WORD   0
    000072 000000  .WORD   0
    000074 000000  .WORD   0
60

```

Support the following special functions (.DRSPF):

```

61 000076          .DRSPF  <201>          ;Reset 'received XOFF from host' flag
    000022    176  .BYTE   176
    000023    200  .BYTE   200
    000024    000  .BYTE   0
    000025    000  .BYTE   0
    000026    000  .BYTE   0
    000027    000  .BYTE   0
    000030 000000  .WORD   000000
62                    ; and send XON to host
63 000032          .DRSPF  <202>          ;Set/clear BREAK
    000022    176  .BYTE   176
    000023    200  .BYTE   200
    000024    000  .BYTE   0
    000025    000  .BYTE   0
    000026    000  .BYTE   0
    000027    000  .BYTE   0
    000030 000000  .WORD   000000
64                    ; word count <> 0, BREAK
65                    ; word count = 0, end BREAK
66 000032          .DRSPF  <203>          ;Special read. Word count is maximum
    000022    176  .BYTE   176
    000023    200  .BYTE   200
    000024    000  .BYTE   0
    000025    000  .BYTE   0
    000026    000  .BYTE   0
    000027    000  .BYTE   0
    000030 000000  .WORD   000000
67                    ; number of bytes to read. Terminates
68                    ; when number of bytes specified have
69                    ; been read or when the input buffer
70                    ; is empty. Always reads at least one
71                    ; byte even if buffer is empty when
72                    ; the read is issued.
73 000032          .DRSPF  <204>          ;Returns driver status in first word
    000022    176  .BYTE   176
    000023    200  .BYTE   200
    000024    000  .BYTE   0
    000025    000  .BYTE   0
    000026    000  .BYTE   0
    000027    000  .BYTE   0
    000030 000000  .WORD   000000
74                    ; of buffer. High byte = driver edit
75                    ; level. Low byte = XOFF status and
76                    ; some modem signals.
77 000032          .DRSPF  <205>          ;Sets a flag which will cause
    000022    176  .BYTE   176
    000023    200  .BYTE   200
    000024    000  .BYTE   0
    000025    000  .BYTE   0
    000026    000  .BYTE   0

```

```

000027 000 .BYTE 0
000030 000000 .WORD 000000
78 ; interrupts to be turned off on
79 ; program exit
80 000032 .DRSPF <206> ;Sets/Resets DTR
000022 176 .BYTE 176
000023 200 .BYTE 200
000024 000 .BYTE 0
000025 000 .BYTE 0
000026 000 .BYTE 0
000027 000 .BYTE 0
000030 000000 .WORD 000000
81 ; word count
<> 0, set DTR
82 ; word count = 0, reset DTR
83
84 ; Handler version number given to VTCOM in INIT message
85
86 000022 $$$VER == 18. ;VTCOM and XL must be a matched set
87
88 ; RT-11 System communications area
89

```

The following macros (through .TSTDF) use the standard definitions from SYSTEM.MLB:

```

90 .MCALL .SYCDF
91 000032 .SYCDF ;Define system communications area
92
93 ; RMON Fixed offset area
94
95 .MCALL .FIXDF .CF1DF .CF2DF
96 .MCALL .SGNDF
97
98 000032 .FIXDF ;Define RMON fixed offsets
1 000032 .CF1DF ;Define config word 1 bits
2 000032 .CF2DF ;Define config word 2 bits
3 000032 .SGNDF ;Define SYSGEN features word bits
4
5 ; Multiterminal status block
6
7 .MCALL .MSTDF
8
9 000032 .MSTDF ;Define .MTSTA status block
10
11 ; Handler header definitions
12
13 .MCALL .HBGDF
14 000032 .HBGDF ;Define handler header
15
16 ; Handler hooks related definitions
17
18 .MCALL .THKDF .TCBDF .TSTDF
19 000032 .TCBDF ;Define TCB offsets
20 000032 .THKDF ;Define handler hooks data structure
21 000032 .TSTDF ;Define T.STAT word bits
22
23 ; Input buffer definitions
24
25 000100 BUFSIZ = 64. ;Size of input buffer (in bytes)
26 000020 STPSIZ = BUFSIZ/4 ;Low-water mark (when to send XOFF)
27 000060 RSTSIZ = BUFSIZ*3/4 ;High-water mark (when to send XON)
28
29 ; Control Characters
30
31 000012 C.LF = 12 ;Line feed
32 000015 C.CR = 15 ;Carriage return
33 000021 C.CTLQ = 21 ;XON (^Q)
34 000023 C.CTSL = 23 ;XOFF (^S)
35 000032 C.CTLZ = 32 ;End-of-file (^Z)
36
37 ; .SPFUN codes supported by driver
38
39 000201 CLRDRV = 201 ;Reset 'received XOFF from host' flag
40 ; and send XON to host
41 000202 BRKDRV = 202 ;Set/clear BREAK

```

```

42                                     ; word count <> 0, BREAK
43                                     ; word count = 0, end BREAK
44      000203      SRDDRV = 203      ;Special read. Word count is maximum
45                                     ; number of bytes to read. Terminates
46                                     ; when number of bytes specified have
47                                     ; been read or when the input buffer
48                                     ; is empty. Always reads at least one
49                                     ; byte even if buffer is empty when
50                                     ; the read is issued.
51      000204      STSDRV = 204      ;Returns driver status in first word
52                                     ; of buffer. High byte = driver edit
53                                     ; level. Low byte =
54      000001      ST.XFH = 000001    ;XOFF sent to host
55      000002      ST.XOF = 000002    ;XOFF received from host
56      000004      ST.CTS = 000004    ;Dataset: Clear To Send asserted
57      000010      ST.CD  = 000010    ;Dataset: Carrier Detect asserted
58      000020      ST.RI  = 000020    ;Dataset: Ring Indicate asserted
59
60      000205      OFFDRV = 205      ;Sets a flag which will cause
61                                     ; interrupts to be turned off on
62                                     ; program exit
63      000206      DTRDRV = 206      ;Sets/Resets DTR
64                                     ; word count <> 0, set DTR
65                                     ; word count = 0, reset DTR
66      ;NOTE: if you add special function code, add them to .DRSPF too!
67
68      ; Interface bit definitions
69
70      040000      RC.RI  = 040000    ;Ring indicator
71      020000      RC.CTS = 020000    ;Clear to send
72      010000      RC.CD  = 010000    ;Carrier detect
73      000100      RC.IE  = 000100    ;Interrupt enable
74      000004      RC.RTS = 000004    ;Request to send
75      000002      RC.DTR = 000002    ;Data terminal ready
76
77      000100      XC.IE  = 000100    ;Transmitter: interrupt enable
78
79      .IF NE XL$DVE
80      XC.SMK = 170000    ;Speed mask
81      XC.SCE = 004000    ;Speed change enable
82      .ENDC ;NE XL$DVE
83
84      .IF NE XL$SBC
85      XC.SMK = 000070    ;Speed mask
86      XC.SCE = 000002    ;Speed change enable
87      .ENDC ;NE XL$SBC
88
89      000001      XC.BRK = 000001    ; BREAK
90
91      .IF NE XL$PC
92
93      ; PRO-3xx Interrupt controller registers
94
95      ICODR = 173200    ;Interrupt controller 0 data register
96      ICOCR = ICODR+2  ;Interrupt controller 0 csr register
97
98      ; PRO-3xx Communications port registers
99
100     XL$BUF = XL$CSR    ;Recv/Xmit buffer register
101     XL$CSA = XL$CSR+2  ;CSR register A
102     XL$CSB = XL$CSR+6  ;CSR register B
103     XL$MC0 = XL$CSR+10 ;Modem control register 0
104     XL$MC1 = XL$CSR+12 ;Modem control register 1
105     XL$BAU = XL$CSR+14 ;Baud rate control register
106
107     ; CSRA Write/Read register bit definitions
108
109     RPT.R0 = 000      ;Write/Read register 0
110     CRC.TR = 300     ; Reset transmit underrun/end of message latch
111     CMD.RE = 020     ; Reset external/status interrupts
112     CMD.CR = 030     ; Channel reset
113     CMD.RT = 050     ; Reset transmitter interrupt pending
114     CMD.ER = 060     ; Reset error latches
115     CMD.EI = 070     ; End of interrupt
116     RPT.R1 = 001     ;Write/Read register 1
117     W1.RIE = 030     ; Receiver interrupt enable
118     ; (Int. on rec. char or special (no parity))
119     W1.TIE = 002     ; Transmitter interrupt enable

```

```

120          RPT.R2 = 002          ;Write/Read register 2
121          RPT.R3 = 003          ;Write register 3
122          RCL.8 = 300          ; Receiver character length (8 bits)
123          W3.RXE = 001          ; Receiver enable
124          RPT.R4 = 004          ;Write register 4
125          CLK.16 = 100          ; 16x rate multiplier
126          STP.1 = 004          ; 1 stop bit
127          W4.EVN = 002          ; Even parity
128          W4.PEN = 001          ; Parity enable
129          RPT.R5 = 005          ;Write register 5
130          TCL.8 = 140          ; Transmit character length (8 bits)
131          W5.SB = 020          ; Send break
132          W5.TXE = 010          ; Transmitter enable
133
134          ; CSRB Write/Read register bit definitions
135
136          RPT.R1 = 001          ;Write/Read register 1
137          W1.REQ = 004          ; MUST be loaded with 004
138          RPT.R2 = 002          ;Write/Read register 2
139          W2.REQ = 000          ; MUST be loaded with 000
140          R2.IMK = 034          ; Interrupt vector mask
141          IMK.BE = 020          ; Transmit buffer empty
142          IMK.ES = 024          ; External/Status change
143          IMK.CA = 030          ; Received character available
144          IMK.SR = 034          ; Special receiver condition
145
146          ; Modem control Register bit definitions
147
148          CLK.BG = 000          ; Rx = RBRG, Tx = TBRG      ->MD = none
149          M0.DTR = 020          ; Data terminal ready
150          M0.RTS = 010          ; Request to send
151          M1.RI = 100          ; Ring indicator
152          M1.CTS = 040          ; Clear to send
153          M1.CD = 020          ; Carrier detect
154          .ENDC ;NE XL$PC
155
156          ; Baud rate mask definitions (PRO-3xx, DLV11-E,F and MXV11-B)
157
158          .IF NE <XL$PC ! XL$DVE>
159          B.50 = 000          ; 50 baud
160          B.75 = 001          ; 75 baud
161          B.110 = 002          ; 110 baud
162          B.134 = 003          ; 134.5 baud
163          B.150 = 004          ; 150 baud
164          B.300 = 005          ; 300 baud
165          B.600 = 006          ; 600 baud
166          B.1200 = 007          ; 1200 baud
167          B.1800 = 010          ; 1800 baud
168          B.2000 = 011          ; 2000 baud
169          B.2400 = 012          ; 2400 baud
170          B.3600 = 013          ; 3600 baud
171          B.4800 = 014          ; 4800 baud
172          B.7200 = 015          ; 7200 baud
173          B.9600 = 016          ; 9600 baud
174          B.192K = 017          ; 19.2k baud
175          .ENDC ;NE <XL$PC ! XL$DVE>
176
177          ; Baud rate mask definitions [SBC-11 only]
178
179          .IF NE XL$SBC
180          B.300 = 000          ; 300 baud
181          B.600 = 001          ; 600 baud
182          B.1200 = 002          ; 1200 baud
183          B.2400 = 003          ; 2400 baud
184          B.4800 = 004          ; 4800 baud
185          B.9600 = 005          ; 9600 baud
186          B.192K = 006          ; 19.2K baud
187          B.384K = 007          ; 38.4k baud
188          .ENDC ;NE XL$SBC
189
190          ; Miscellaneous definitions
191
192          177776          PS          =: 177776          ; Processor status word
193          000007          UNITMK      =: 007          ; Q$UNIT unit number mask
194          000370          JOBMK       =: 370          ; Q$JNUM job number mask
195
196          ; Macro to define LSB of bit field
197

```

```

198 .MACRO LSBDF SYMBOL,VALUE
199     SYMBOL = VALUE & <-VALUE>
200 .ENDM ;LSBDF

```

Block 0 of handler file

```

1 .SBTTL Block 0 of handler file
2

```

The SPEED table is placed low in block 0 without conflicting with audit trail:

```

3 000032 .ASECT
4 000120 . = 120
5
6 .IF NE <XL$PC ! XL$DVE ! XL$SBC>
7 ; SPEED table. Mask for given speed is same as word offset into table.
8 ; To select 134.5 bps, specify 134 in the SET command.
9
10 SPEEDT:
11 .IF NE <XL$DVE ! XL$PC>
12 .WORD 50., 75., 110., 134., 150., 300.
13 .WORD 600., 1200., 1800., 2000., 2400., 3600.
14 .WORD 4800., 7200., 9600., 19200.
15 .ENDC ;NE <XL$DVE ! XL$PC>
16
17 .IF NE XL$SBC
18 .WORD 300., 600., 1200., 2400., 4800., 9600.
19 .WORD 19200., 38400.
20 .ENDC ;NE XL$SBC
21
22 .WORD 0 ;Table fence
23
24 .ENDC ;NE <XL$PC ! XL$DVE ! XL$SBC>
25

```

We must ensure that 0 fence for display CSRs is not overwritten:

```

26 000120 .Assume . LE DISCSR-2 MESSAGE=<Code before installation code too large>

```

INSTALLATION CODE

```

1 .SBTTL INSTALLATION CODE
2
3 .ENABL LSB
4
5 .IF EQ XL$MTY
6 .DRINS XL
7 .IFF ;EQ XL$MTY

```

Ensure that install-time CSR is zero when defaulting to MTTY, so the handler always installs:

```

8 000120 .DRINS -XL
9 000172 000000 .WORD 0
10 000174 176500 DISCSR: .WORD <-XL$CSR>
11 000176 000000 INSCSR: .WORD 0
12 .ENDC ;EQ XL$MTY
13
14 000200 000401 BR 10$ ;Install as a data device
15 000202 000416 BR 40$ ; never as a system device
16
17 000204 013700 10$: MOV @#$SYPTR,R0 ;R0->$RMON
18 000054
19 000210 032760 BIT #PRO$$,$CNFG2(R0) ;Installing on a PRO-3xx?
20 020000
21 000370
22
23 .IF EQ XL$PC
24 000216 001010 BNE 40$ ;Yes, then reject the installation
25 .IF NE XL$MTY
26 000220 105767 TSTB I$MTTY ;Are handler hooks needed?
27 000020

```

```

21 000224 001404      BEQ      20$           ;Nope...
22 000226 005760      TST      $THKPT(R0)       ;Yes, is the support available?
                                000000G
23 000232 001402      BEQ      40$           ;Nope, reject the installation

```

Hooks cannot be established until handler is in memory, which doesn't happen until Fetch/Load:

```

24 000234 000400      BR       30$           ;Yes, nothing to do until fetch/load
25
26 000236           20$:
27                                .ENDC ;NE XL$MTY
28                                .IFB ;EQ XL$PC
29                                BEQ      40$           ;Nope, then reject the installation
30                                .ENDC ;EQ XL$PC
31
32                                .IF EQ XL$PC
33                                .IF NE <XL$DVE ! XL$SBC>
34                                MOV      INSCSR,R0       ;R0->Receiver CSR

```

Speed set at install-time:

```

35                                MOV      ISPEED,4(R0)       ;Set the speed (in transmitter CSR)
36                                .ENDC ;NE <XL$DVE ! XL$SBC>
37                                .IFB ;EQ XL$PC
38                                MOV      ISPEED,@XL$BAU   ;Set the XMIT/RECV baud rate
39
40                                ; Things to do through csr A
41
42                                MOV      #XL$CSA,R0       ;R0->csr A
43                                MOV      #CMD.CR,@R0      ;Reset channel A
44                                MOV      #CRC.TR,@R0      ;Reset transmitter underrun latch
45
46                                MOV      #RPT.R4,@R0      ;Select csr A, write register 4
47                                MOV      #<CLK.16!STP.1>,@R0 ; set clock rate xl6, 1 stop bit
48
49                                MOV      #RPT.R3,@R0      ;Select csr A, write register 3
50                                MOV      #<W3.RXE!RCL.8>,@R0 ; set receiver enable, 8-bit chars
51
52                                MOV      #RPT.R5,@R0      ;Select csr A, write register 5
53                                MOV      #<W5.TXE!TCL.8>,@R0 ; set transmitter enable, 8-bit chars
54
55                                MOV      #RPT.R2,@R0      ;Select csr A, write register 2
56                                MOV      #0,@R0          ; *** must be loaded with 0 ***
57                                MOV      #CMD.RE,@R0      ;Reset external/status interrupts
58
59                                ; Things to do through csr B
60
61                                MOV      #XL$CSB,R0       ;R0->csr B
62                                MOV      #CMD.CR,@R0      ;Reset channel B
63
64                                MOV      #RPT.R2,@R0      ;Select csr B, write register 2
65                                MOV      #W2.REQ,@R0      ; *** ensure base vector of 0 ***
66
67                                MOV      #RPT.R1,@R0      ;Select csr B, write register 1
68                                MOV      #W1.REQ,@R0      ; *** ensure correct vector info ***
69
70                                ; Now we play with the interrupt controller
71
72                                MOV      #<30!3>,@#ICOCR   ;Enable comm port interrupts
73
74                                ; And finally, the modem
75
76                                MOV      #CLK.BG,@#XL$MCO   ;Set modem clock
77                                .ENDC ;EQ XL$PC
78
79 000236 005727 30$:  TST      (PC)+       ;Accept the installation (carry=0)
80 000240 000261 40$:  SEC              ;Reject the installation (carry=1)
81 000242 000207      RETURN
82
83                                .DSABL  LSB
84
85                                .IF NE <XL$PC ! XL$DVE ! XL$SBC>
86                                .IF NE <XL$DVE ! XL$SBC>
87                                LSBDF  ...,XC.SMK       ;Determine lowest bit of speed mask
88                                .ENDC ;NE <XL$DVE ! XL$SBC>

```

```

89             ISPEED:
90             .IF NE XL$PC
91             .WORD  <B.1200 * 20> + B.1200 ;Default to 1200 baud RECV and XMIT
92             .ENDC ;NE XL$PC
93
94             .IF NE <XL$DVE ! XL$SBC>
95             .WORD  <B.1200 * ...> ! XC.SCE ;Default to 1200 baud RECV and XMIT
96             .ENDC ;NE <XL$DVE ! XL$SBC>
97             .ENDC ;NE <XL$PC ! XL$DVE ! XL$SBC>
98
99             .IF NE XL$MTY

```

Default flag to MTTY if built for hooks support:

```

100 000244    377 ISMTTY: .BYTE  -1                ; : Install-time 'hooks required' flag
101 000245    000         .BYTE                   ;reserved

```

Duplicate code from .DRBEG to restore pointer to vector table when SET XL NOMTTY is issued:

```

102 000246    000000C VECSAV: .WORD  100000+<<XL$VTB-H1.VEC>/2-1> ; : Vector info for SET NOMTTY
103 000250    176500 CSRSVAV: .WORD  XL$CSR                ; : CSR info for SET NOMTTY
104                                     .ENDC ;NE XL$MTY
105
106 000252                                     .Assume . LE 400 MESSAGE=<Installation code too large>

```

SET OPTION PARAMETER TABLE

```

1             .SBTTL  SET OPTION PARAMETER TABLE
2
3             ;      Option  Data                Routine Syntax
4             ;      -----  ----                -
5
6             .IF EQ 1
7             .DRSET BIT8    <^c177>            O.BIT8  NO      ;[NO]BIT8
8             .ENDC ;EQ 1
9
10            .IF EQ XL$PC
11 000252     .DRSET  CSR      160012            O.CSR   OCT     ;CSR=n
12 000400     160012     160012
13 000402     012712     .RAD50  \CSR\
14 000406     021        .BYTE   <O.CSR-^o400>/2.
15 000407     140        .BYTE   ...V2
16 000410     000000     .WORD    0
17 000412     .DRSET  VECTOR  477              O.VEC   OCT     ;VECTOR=n
18 000410     000477     477
19 000412     105113     .RAD50  \VECTOR\
20 000414     077552     .WORD    0
21 000416     046        .BYTE   <O.VEC-^o400>/2.
22 000417     140        .BYTE   ...V2
23 000420     000000     .WORD    0
24             .ENDC ;EQ XL$PC
25
26            .IF EQ XL$PC
27            .IF NE XL$PDT
28            .DRSET  LIGHTS  -1                O.LGHT  NO      ;[NO]LIGHTS
29            .ENDC ;NE XL$PDT
30            .ENDC ;EQ XL$PC
31
32            .IF NE XL$MTY
33 000422     .DRSET  LINE     16.              O.LINE  NUM     ;LINE=n
34 000420     000020     16.
35 000422     046166     .RAD50  \LINE\
36 000424     017500
37 000426     056        .BYTE   <O.LINE-^o400>/2.
38 000427     100        .BYTE   ...V2
39 000430     000000     .WORD    0
40 000432     .DRSET  MTTY    -1                O.MTTY  NO      ;[NO]MTTY
41 000430     177777     -1
42 000432     052164     .RAD50  \MTTY\
43 000434     116100
44 000436     063        .BYTE   <O.MTTY-^o400>/2.
45 000437     200        .BYTE   ...V2
46 000440     000000     .WORD    0

```

```

24             .ENDC ;NE XL$MTY
25
26             .IF NE <XL$PC ! XL$DVE ! XL$SBC>
27 .DRSET  SPEED  NOP                O.SPEE  NUM      ;SPEED=n
28             .ENDC ;NE <XL$PC ! XL$DVE ! XL$SBC>

```

SET OPTION PROCESSING ROUTINES

```

1             .SBTTL  SET OPTION PROCESSING ROUTINES
2
3             .IF EQ 1
4             ; SET XL [NO]BIT8
5
6 O.BIT8:  CLRB   R3                    ;Ensure high bit is left alone
7             NOP                        ;placekeeper
8             MOV   R3,CHMASK            ;Save character alteration mask
9             RETURN
10            .ENDC ;EQ 1
11
12            .IF EQ XL$PC
13
14            ; SET XL CSR=octal_address
15

```

When SET XL MTTY in effect, cannot alter install-time CSR (176); must save it for restore when SET XL NOMTTY issued:

```

16 000442      O.CSR:
17             .IF NE XL$MTY
18 000442 010067  MOV   R0,CSRSVAV        ;Yes, update saved CSR for SET NOMTTY
19             177602
20 000446 105767  TSTB   I$MTTY            ;Are we set MTTY?
21             177572
22 000452 001002  BNE    20$                          ;Yep, don't set install-time word
23             .ENDC ;NE XL$MTY
24
25 000454 010067 10$:  MOV   R0,INSCSR        ;Let installation code know
26             177516
27 000460 010067 20$:  MOV   R0,DISCSR        ;Fill in display CSR
28             177510
29 000464      .ADDR  #XIS,R1                ;R1 -> Where to put CSR info
30 000464 010701  MOV   PC,R1
31 000466 062701  ADD   #XIS-. ,R1
32             177444'
33 000472 012702  MOV   #4,R2                    ;R2 = Count of words to set
34             000004
35 000476 010021 30$:  MOV   R0,(R1)+          ;Set a table entry
36 000500 062700  ADD   #2,R0                    ;Prepare for next entry
37             000002
38 000504 005302  DEC   R2                            ;More to do?
39 000506 003373  BGT   30$                          ;Yep...
40 000510 020003  CMP   R0,R3                          ;Was address specified in range?
41 000512 000207  RETURN                                ; c-bit=0 if so, =1 if not
42
43             ; SET XL VECTOR=octal_address
44
45 000514 010067 O.VEC: MOV   R0,XL$VTB        ;Save the new input interrupt vector
46             000142'
47 000520      ADD   #4,R0
48             000004
49 000524 010067  MOV   R0,XL$VTB+6            ; and output interrupt vector
50             000150'
51 000530 020300  CMP   R3,R0                          ;Was address specified in range?
52 000532 000207  RETURN                                ; c-bit=0 if so, =1 if not
53
54             .IF NE XL$PDT
55
56             ; SET XL [NO]LIGHTS
57
58 O.LGHT:  CLR   R3                    ;LIGHTS entry point
59             NOP                        ; (padding)
60             COM   R3                    ;NOLIGHTS entry point
61             MOV   R3,LitFlg            ;Set/Reset lights flag
62             BR    O.NOR
63             .ENDC ;NE XL$PDT

```

```

52
53             .IF NE XL$MTY
54
55             ; SET XL LINE=line_number
56
57 000534 120003 O.LINE: CMPB   R0,R3             ;Is line number valid?
58 000536 101027             BHI    O.ERR             ;Nope...
59 000540 110067             MOVB   R0,O$LINE         ;Yes, set line number to use
60 000544 000423             BR     O.NOR
61
62             ; SET XL [NO]MTTY
63
64 000546 000411 O.MTTY: BR     10$                 ;Entry point for MTTY
65 000550 000240             NOP                    ;placekeeper
66 000552 005000             CLR    R0                 ;Entry point for NOMTTY
67 000554 016767             MOV    CSRSV,INSCSR        ;Nope, restore install-time CSR
68 000562 016767             MOV    VECSV,H1.VEC         ; and vector information
69 000570 000404             BR     20$
70
71 000572 005067 10$: CLR    INSCSR                 ;Reset install-time CSR and
72 000576 005067             CLR    H1.VEC             ; vector so handler installs
73 000602 110067 20$: MOVB   R0,O$MTTY             ;Set/Reset MTTY hooks use flag
74 000606 110067             MOVB   R0,I$MTTY         ; and inform install code of setting
75 000612 000400             BR     O.NOR
76             .ENDC ;NE XL$MTY
77             .ENDC ;EQ XL$PC
78
79             .IF NE <XL$PC ! XL$DVE ! XL$SBC>
80
81             ; SET XL SPEED=decimal_speed

```

Setting speed alters the on-disk image, but also takes immediate effect:

```

82
83             O.SPEE:

```

Can't use when MTTY is in effect because not all lines have programmable baud rate:

```

84             .IF NE XL$MTY
85             TSTB   I$MTTY             ;Handler hooks in use?
86             BNE   O.ERR             ;Yes, can't touch the CSR
87             .ENDC ;NE XL$MTY
88
89             .ADDR  #SPEEDT,R1         ;R1 -> Baud rate table
90             10$: TST   @R1             ;End of table?
91             BEQ   O.ERR             ;Yes, speed requested is invalid
92             CMP   R0,(R1)+           ;Nope, request match this entry?
93             BNE   10$               ;Nope, try another speed entry
94             SUB   PC,R1             ;Yes, determine speed mask
95             SUB   #<SPEEDT+2-.>,R1   ; ...
96
97             .IF NE XL$PC
98             ASR   R1                 ;Convert from byte to word offset
99             MOVB  R1,-(SP)           ;Save the receive speed mask
100            ASL   R1                 ;And make transmit speed match
101            ASL   R1                 ; by shifting
102            ASL   R1                 ; it to the
103            ASL   R1                 ; high nibble
104            BISB  (SP)+,R1           ;OR in the receive speed mask
105            MOVB  R1,@#XL$BAU        ; and change the speed now
106            .ENDC ;NE XL$PC
107
108            .IF NE XL$DVE
109            SWAB  R1                 ;Move to high byte
110            ASL   R1                 ; then shift mask to where
111            ASL   R1                 ; it should be for

```

```

112 ASL R1 ; a DLV11-E
113 .ENDC ;NE XL$PC
114
115 .IF NE XL$SBC
116 ASL R1 ;Shift mask to where it
117 ASL R1 ; should be for SBC or MXV SLU
118 .ENDC ;NE XL$SBC
119
120 .IF NE <XL$DVE ! XL$SBC>
121 BIS #XC.SCE,R1 ;Set the 'speed change enable' bit
122 MOV INSCSR,R0 ;R0->Receiver CSR
123 .ENDC ;NE <XL$DVE ! XL$SBC>
124
125 MOV R1,ISPEED ;Save new speed for installation
126
127 .IF NE <XL$DVE ! XL$SBC>
128 MOV ISPEED,4(R0) ;Set the speed (in transmitter CSR)
129 .ENDC ;NE <XL$DVE ! XL$SBC>
130 .BR O.NOR
131 .ENDC ;NE <XL$PC ! XL$DVE ! XL$SBC>
132
133 000614 005727 O.NOR: TST (PC)+ ;Success (carry=0)
134 000616 000261 O.ERR: SEC ;Failure (carry=1)
135 000620 000207 RETURN
136
137 000622 .Assume . LE 1000 MESSAGE=<Set code too large>

```

DRIVER ENTRY

```

1 .SBTTL DRIVER ENTRY
2
3 ; The handler gets entered here each time the monitor places a new
4 ; request on the device queue. The handler either processes the
5 ; request immediately and returns it to the monitor or the request
6 ; is removed from the device queue and placed on one of the internal
7 ; queues. There is one internal queue for input and one for output.
8 ;
9 ; Because of the separate queues, simultaneous input and output may
10 ; be performed.
11
12 .ENABL LSB
13
14 .IF EQ XL$MTY
15 .DRBEG XL

```

Following code is for hooks support. Ensures vector word is zero so handler loads without affecting any vectors when XL is SET MTTY. Restored with SET XL NOMTTY.

```

16 .IFF ;EQ XL$MTY
17 000622 .DRBEG XL,0 ;Default to use handler hooks
000052 002704 .WORD <XLEND-XLSTRT>
000054 000000 .WORD XLDSIZE
000056 007057 .WORD XLSTS
000060 000006 .WORD ^o<ERL$G+<MMG$T*2>+<TIM$IT*4>+<RTE$M*10>>
000000 000000 .WORD 0&^C3.
000002 001120 .WORD XLINT-.,^o340
000004 000340
000006 000000 XLLQE:: .WORD 0
000010 000000 XLCQE:: .WORD 0
000012 000257 .WORD 257
18 .ENDC ;EQ XL$MTY
19
20 000014 016704 MOV XLCQE,R4 ;R4->Current queue element
177770
21 000022' STATFG = <. + 2>
22 000020 006227 ASR #1 ;First call since .FETCH/LOAD or
000001
23 ; last shutdown?
24 000024 103013 BCC 40$ ;Nope...
25
26 .IF EQ XL$PC
27 000026 004767 CALL ENAINI ;Turn on receiver interrupts
002166

```

```

28 000032 012700      MOV     #<RC.RTS!RC.DTR>,R0      ;Assert DTR
      000006
29 000036 004767      CALL    SETSTT                    ; ...
      002412
30 000042 012767      MOV     #-2,SNDS                  ;Indicate we must send an XON
      177776
      001070
31 000050 004767      CALL    ENAOUI                    ;Enable output interrupts
      002216
32                                     .IF NE XL$PDT
33                                     CALL    SETLIT                      ;Set the lights to indicate state
34                                     .ENDC ;NE XL$PDT
35                                     .IFF ;EQ XL$PC
36                                     MOV     #RPT.R1,@CSRA                ;Select csr A, write register 1
37                                     BIS     #<W1.RIE!W1.TIE>,SSRAW1      ;Turn on RECV and XMIT interrupts
38                                     MOV     SSRAW1,@CSRA                ; (update from software register)
39                                     BIS     #<M0.DTR!M0.RTS>,@MCR0      ;Force DTR and RTS
40                                     MOVB   #C.CTLQ,@DBUF              ;First thing we send is an XON
41                                     .ENDC ;EQ XL$PC
42
43 000054 116405 40$:      MOVB   Q$FUNC(R4),R5              ;Get the function code
      000002
44 000060 001040      BNE     SPFUN                      ;If non-zero, we have a .SPFUN
45 000062 006364      ASL    Q$WCNT(R4)                ;Convert word count to byte count
      000006
46 000066 103406      BCS    WRITE                       ;If negative, write request
47                                     ; otherwise, read
48 000070 004567 READ:   JSR     R5,XLENQ                ;Queue the read request
      002002

```

Internal input queue:

```

49 000074 000000 XOCQE:   .WORD 0                          ; : address of first element on queue
50 000076 000000 XOLQE:   .WORD 0                          ; : address of last element on queue
51 000100 000167      CALLR  XIIN                        ;Process any input already received,
      001436
52                                     ; read will be completed via
53                                     ; interrupts
54
55 000104 005267 WRITE:  INC     QCHG                       ;Set 'queue being modified' flag
      001066
56 000110 004567      JSR     R5,XLENQ                ;Queue the write request
      001762

```

Internal output queue:

```

57 000114 000000 XOCQE:   .WORD 0                          ; : address of first element on queue
58 000116 000000 XOLQE:   .WORD 0                          ; : address of last element on queue
59 000120 005067      CLR    QCHG                       ;Reset 'queue being modified' flag
      001052
60
61                                     .IF EQ XL$PC
62 000124 004767      CALL    ENAOUI                    ;Enable output interrupts
      002142
63                                     .IFF ;EQ XL$PC
64                                     CALL    GNXTCH                      ;Get a character for output
65                                     BEQ    50$                          ;None available...
66                                     MOVB   R5,@DBUF                    ;Now prime the interrupt pump
67                                     .ENDC ;EQ XL$PC
68
69 000130 000207 50$:      RETURN
70
71                                     .DSABL  LSB

```

REGISTERS AND VECTOR TABLES

```

1          .SBTTL  REGISTERS AND VECTOR TABLES
2
3          .IF EQ XL$PC
4          ; *** Begin Critical Ordering ***
5 000132 176500  XIS:  .WORD  XL$CSR           ; : Receiver status register
6 000134 176502  XIB:  .WORD  XL$CSR+2       ; : Receiver buffer register
7 000136 176504  XOS:  .WORD  XL$CSR+4       ; : Transmitter status register
8 000140 176506  XOB:  .WORD  XL$CSR+6       ; : Transmitter buffer register
9          ; *** End Critical Ordering ***
10         .IFF ;EQ XL$PC
11         DBUF:  .WORD  XL$BUF           ; : Input/Output buffer register
12         CSRA:  .WORD  XL$CSA          ; : Control/Status register A
13         CSRB:  .WORD  XL$CSB          ; : Control/Status register B
14         MCR0:  .WORD  XL$MCO          ; : Modem control/status register 0
15         MCR1:  .WORD  XL$MC1          ; : Modem control/status register 1
16         BAUD:  .WORD  XL$BAU          ; : Baud rate control register
17         .ENDC ;EQ XL$PC
18
19         ; Now for some software registers
20
21         .IF NE XL$PC
22         SSRAW1: .WORD  0                ;Software status A, write register 1
23         SSRAW5: .WORD  <W5.TXE!TCL.8>  ;Software status A, write register 5
24         .ENDC ;NE XL$PC
25
26         ; Define the interrupt vectors
27
28         .IF EQ XL$PC
29 000142         .DRVTB  XL,XL$VEC,XIINT      ;Input interrupt servicer
30 000142 000300  .WORD  XL$VEC&^C3.,XIINT-. ,^o340!0,^o100000
31 000144 001214
32 000146 000340
33 000150 100000
34
35 000152         .DRVTB  ,XL$VEC+4,XLINT      ;Output interrupt servicer
36 000150 000304  .WORD  XL$VEC+4&^C3.,XLINT-. ,^o340!0,^o100000
37 000152 000750
38 000154 000340
39 000156 100000
40
41         .IFF ;EQ XL$PC
42         .DRVTB  XL,XL$VEC,XLINT      ;Input/Output interrupt servicer
43         .DRVTB  ,XL$VEC+4,XLINT
44         .ENDC ;EQ XL$PC
45
46 000160 177600  CHMASK: .WORD  ^C177      ;Character mask
47
48         .IF EQ XL$PC
49         .IF NE XL$PDT

```

LIGHTS ROUTINE FOR PDT-11'S

```

1          .SBTTL  LIGHTS ROUTINE FOR PDT-11'S
2
3          ;+
4          ;
5          ; Sets PDT lights to indicate XON/XOFF state.
6          ;
7          ; LED 1 on if PDT has sent XOFF
8          ; LED 2 on if PDT has received XOFF
9          ;
10         ;-
11
12         SETLIT: TST  (PC)+                ;SET XL LIGHTS in effect?
13         LITFLG: .WORD  0                  ; : lights flag (0 = no, <0 = yes)
14         BEQ  30$                          ;Nope...
15         MOV  #040000,R5                   ;Default to lights off
16         TST  SNDS                          ;XOFF sent to host?
17         BLE  10$                          ;Nope...
18         BIS  #000100,R5                   ;Yes, turn on LED 1
19         10$: TST  RECS                     ;XOFF received from host?
20         BEQ  20$                          ;Nope...
21         BIS  #000200,R5                   ;Yes, turn on LED 2
22         20$: MOV  R5,#177420              ;Force the new lights setting
23         30$: RETURN
24
25         .ENDC ;NE XL$PDT
26         .ENDC ;EQ XL$PC

```

SPFUN PROCESSING

```

1          .SBTTL  SPFUN PROCESSING
2
3          ; This section of code gets jumped to. It expects that the address of the
4          ; queue element is in R4 and the address of the special function code to
5          ; be executed is in R5.
6

```

Special read may require post-interrupt processing, so it must be internally queued:

```

7 000162 120527 SPFUN: CMPB   R5,#SRDDRV      ;Special read request?
                        000203
8 000166 001740      BEQ    READ          ; Yes, go queue it
9 000170 120527      CMPB   R5,#BRKDRV      ;[end]BREAK request?
                        000202
10 000174 001423     BEQ    20$           ; Yes...
11 000176 120527     CMPB   R5,#CLRDRV      ;Clear driver flags request?
                        000201
12 000202 001440     BEQ    40$           ; Yes...
13 000204 120527     CMPB   R5,#STSDRV      ;Status request?
                        000204
14 000210 001445     BEQ    50$           ; Yes...
15 000212 120527     CMPB   R5,#OFFDRV      ;Shutting us down?
                        000205
16 000216 001502     BEQ    100$          ;Yes...
17 000220 120527     CMPB   R5,#DTRDRV      ;DTR set/reset?
                        000206
18 000224 001514     BEQ    110$          ;Yes...
19
20 000226          10$:  .DRFIN  XL          ;Unknown .SPFUN, ignore
                        ;Inform monitor of completion

```

SPFUN routines can be processed without post-interrupt processing, so they are handled without being moved to internal queue and returned to RT-11:

```

000226 010704      MOV    PC,R4
000230 062704      ADD    #XLCQE-.,R4
                        177560
000234 013705      MOV    @#^o54,R5
                        000054
000240 000175      JMP    @^o270(R5)
                        000270
21
22          ; [end]BREAK processing
23          ;   Word count indicates operation
24          ;   (0 = end break, non-zero = break)
25
26 000244 005764 20$:  TST    Q$WCNT(R4)      ;Break or end-break?
                        000006
27 000250 001406     BEQ    30$           ;If zero, end-break...
28 000252 012767     MOV    #1,BRKFLG      ;Break, set 'break in progress' flag
                        000001
                        000652
29
30          .IF EQ XL$PC
31 000260 004767     CALL   SETBRK        ;Turn on break
                        002064
32          .IFF ;EQ XL$PC
33          MOV    #RPT.R5,@CSRA      ;Select csr A, write register 5
34          BIS    #W5.SB,SSRAW5      ;Turn on break
35          MOV    SSRAW5,@CSRA      ; (update from software register)
36          .ENDC ;EQ XL$PC
37
38 000264 000760     BR    10$
39
40 000266          30$:
41          .IF EQ XL$PC
42 000266 004767     CALL   RESBRK        ;Turn off break
                        002030
43          .IFF ;EQ XL$PC
44          MOV    #RPT.R5,@CSRA      ;Select csr A, write register 5
45          BIC    #W5.SB,SSRAW5      ;Turn off break
46          MOV    SSRAW5,@CSRA      ; (update from software register)
47          .ENDC ;EQ XL$PC

```

```

48
49 000272 005067 CLR BRKFLG ;Reset the 'break in progress' flag
000634
50
51 .IF EQ XL$PC
52 000276 004767 CALL ENAOUI ;Make sure output is running
001770
53 .ENDC ;EQ XL$PC
54
55 000302 000751 BR 10$
56
57 ; Clear driver flags request
58 ; resets received XOFF flag
59 ; sends XON to host
60
61 000304 005067 40$: CLR RECS ;Reset the 'received XOFF' flag
000660
62
63 .IF EQ XL$PC
64 000310 012767 MOV #-2,SNDS ;Indicate we want an XON sent
017776
000622
65 000316 004767 CALL ENAOUI ;Make sure output is running
001750
66 .IF NE XL$PDT
67 CALL SETLIT ;Update lights display
68 .ENDC ;NE XL$PDT
69 .IFF ;EQ XL$PC
70 CLR SNDS ;Indicate that an XON has been
71 MOVB #C.CTLQ,@DBUF ; sent
72 .ENDC ;EQ XL$PC
73
74 000322 000741 BR 10$
75
76 ; Get Status request
77 ; returns handler version in high byte
78 ; returns XON/XOFF state in low byte
79 ; bit 0 on if host has been XOFF'd
80 ; bit 1 on if host has XOFF'd us
81 ; bit 2 on if CTS is asserted
82 ; bit 3 on if CD is asserted
83 ; bit 4 on if RI is asserted
84
85 000324 012705 50$: MOV $$$VER*400,R5 ;High byte = handler version
011000
86 000330 005767 TST SNDS ;Have we XOFF'd host?
000604
87 000334 003401 BLE 60$ ;Nope...
88
89 000336 .ASSUME ST.XFH EQ 1
90 000336 005205 INC R5 ;Yes, set the indicator
91 000340 005767 60$: TST RECS ;Have we been XOFF'd?
000624
92 000344 001402 BEQ 70$ ;Nope...
93 000346 052705 BIS #ST.XOF,R5 ;Yes, set the indicator
000002
94 000352 70$:
95 .IF EQ XL$PC
96 000352 004767 CALL GETSTT ;Get current status
002022
97 000356 032700 BIT #RC.CTS,R0 ;Is 'Clear To Send' asserted?
020000
98 .IFF ;EQ XL$PC
99 BIT #M1.CTS,@MCR1 ;Is 'Clear To Send' asserted?
100 .ENDC ;EQ XL$PC
101
102 000362 001402 BEQ 80$ ;Nope...
103 000364 052705 BIS #ST.CTS,R5 ;Yes, set an indicator
000004
104
105 000370 80$:
106 .IF EQ XL$PC
107 000370 032700 BIT #RC.CD,R0 ;Is 'Carrier Detect' asserted?
010000
108 .IFF ;EQ XL$PC
109 BIT #M1.CD,@MCR1 ;Is 'Carrier Detect' asserted?
110 .ENDC ;EQ XL$PC
111

```

```

112 000374 001402      BEQ      82$                ;Nope...
113 000376 052705      BIS      #ST.CD,R5          ;Yes, set an indicator
      000010

114
115 000402              82$:
116                      .IF EQ XL$PC
117 000402 032700      BIT      #RC.RI,R0          ;Is 'Ring Indicator' asserted?
      040000

118                      .IFF ;EQ XL$PC
119                      BIT      #M1.RI,@MCR1          ;Is 'Ring Indicator' asserted?
120                      .ENDC ;EQ XL$PC
121

122 000406 001402      BEQ      84$                ;Nope...
123 000410 052705      BIS      #ST.RI,R5          ;Yes, set an indicator
      000020

124
125 000414              84$:
126                      .IF EQ MMG$T
127                      MOV      R5,@Q$BUFF(R4)          ;Return the status word
128                      .IFF ;EQ MMG$T
129 000414 010546      MOV      R5,-(SP)          ;Return the status word
130 000416 004777      CALL    @$PTWRD           ; ...
      002252

131                      .ENDC ;EQ MMG$T
132

133 000422 000701      BR       10$
134
135                      ; Shut down driver request (OFFDRV)
136                      ; Sets a flag such that when VTCOM exits, interrupts will
137                      ; not be re-enabled. STATFG is used as the once-only,
138                      ; interrupt startup flag.
139

140 000424 116446 100$:  MOVB     Q$JNUM(R4),-(SP)    ;Save Q$JNUM
      000003
141 000430 042716      BIC     #^C<JOBMK>,@SP      ;Isolate job number issuing request
      177407
142 000434 006216      ASR     @SP                 ;Shift for abort code check
143 000436 006216      ASR     @SP
144 000440 006216      ASR     @SP
145 000442 112667      MOVB     (SP)+,JNUM         ;Save it for later check
      000040
146 000446 012767      MOV     #1,STATFG          ;Reset us to pre-start state
      000001
      177346

147 000454 000664      BR       10$
148
149                      ; Set/Reset DTR (DTRDRV)
150                      ; Sets or resets DTR based on word count
151                      ; (0 = DTR off, <>0 = DTR on)
152
153 000456              110$:
154                      .IF EQ XL$PC
155 000456 004767      CALL    GETSTT             ;Get current state
      001716
156 000462 042700      BIC     #<RC.RTS!RC.DTR>,R0 ;Assume DTR is desired off
      000006

157                      .IFF ;EQ XL$PC
158                      MOVB     @MCR0,R0                ;Get current state
159                      BIC     #<M0.DTR!M0.RTS>,R0       ;Assume DTR is desired off
160                      .ENDC ;EQ XL$PC
161

162 000466 005764      TST     Q$WCNT(R4)         ;Correct assumption?
      000006

163 000472 001402      BEQ     115$                ;Yep...
164
165                      .IF EQ XL$PC
166 000474 052700      BIS     #<RC.RTS!RC.DTR>,R0 ;Nope, turn it on
      000006

167                      .IFF ;EQ XL$PC
168                      BIS     #<M0.DTR!M0.RTS>,R0       ;Nope, turn it on
169                      .ENDC ;EQ XL$PC
170

171 000500              115$:
172                      .IF EQ XL$PC
173 000500 004767      CALL    SETSTT             ;Assert desired bits
      001750

174                      .IFF ;EQ XL$PC
175                      MOVB     R0,@MCR0                ;Set desired state

```

```

176                                     .ENDC ;EQ XL$PC
177
178 000504 000650          BR          10$
179
180 000506                JNUM:      .BLKW                ; :Job number which issued OFFDRV
181
182                                     .IF NE XL$PC

```

INTERRUPT SERVICE/DISPATCHER

```

1                                     .SBTTL  INTERRUPT SERVICE/DISPATCHER
2
3                                     ;+
4                                     ;
5                                     ; Interrupt entry point for input and output interrupts. The interrupt
6                                     ; type is determined by bits <04:02> in RR2 of CSR B. The four defined
7                                     ; types of interrupts are:
8                                     ;
9                                     ;      1) Transmitter buffer empty      (^B100xx)
10                                    ;      2) External/status change      (^B101xx)
11                                    ;      3) Received character available (^B110xx)
12                                    ;      4) Special receiver condition  (^B111xx)
13                                    ;
14                                    ;-
15
16                                    .DRAST  XL,4,XLDONE
17
18                                    MOV      #RPT.R2,@CSRB          ;Select csr B, read register 2
19                                    MOV      @CSRB,-(SP)            ;Get the interrupt type
20                                    BIC      #^C<R2.IMK>,@SP        ;Strip the uninteresting stuff
21                                    ASR      @SP                    ;Shift for word table offset
22                                    .ADDR   #INTTAB,@SP,ADD         ;Add address of start of table
23                                    MOV      @(SP),@SP              ;Get the table entry
24                                    ADD     PC,@SP                  ;Convert to address
25 INTDSP: JMP      @(SP)+          ;Dispatch the interrupt
26
27 ESINT: MOV      #CMD.RE,@CSRA          ;Reset external/status interrupts
28 IECOM: MOV     #CMD.EI,@CSRA          ;Declare end of interrupt
29 RETURN
30
31 SRINT: MOV     #CMD.ER,@CSRA          ;Reset error latches
32 JMP      XIINT                        ; then handle as received character
33
34 INTTAB: .WORD  IECOM-INTDSP          ;unknown interrupt
35         .WORD  IECOM-INTDSP          ;unknown interrupt
36         .WORD  IECOM-INTDSP          ;unknown interrupt
37         .WORD  IECOM-INTDSP          ;unknown interrupt
38         .WORD  XOINT-INTDSP          ;Transmitter buffer empty
39         .WORD  ESINT-INTDSP          ;External/Status change
40         .WORD  XIINT-INTDSP          ;Received character available
41         .WORD  SRINT-INTDSP          ;Special receiver interrupt
42         .ENDC ;NE XL$PC
43
44                                     .IF NE XL$MTY

```

MULTITERMINAL HANDLER HOOKS SUPPORT DATA

```

1                                     .SBTTL  Multiterminal Handler Hooks Support Data
2
3                                     ; The following byte indicates whether the handler should make use
4                                     ; of the multiterminal hooks during FETCH/LOAD and during operation.
5
6                                     ; *** SET ***

```

Set/reset by SET XL [NO]MTTY:

```

7 000510 377 O$MTTY: .BYTE -1          ;Default to hooks used
8 000511          .Assume <O$MTTY-XLSTRT> LE 1000 MESSAGE=<Code to set not in block 1>
9
10                                     ; *** SET ***

```

Set/reset by SET XL LINE=n:

```

11 000511      001 O$LINE: .BYTE  XL$LUN           ;Default line to use
12 000512      .Assume <O$LINE-XLSTRT> LE 1000 MESSAGE=<Code to set not in block 1>
13
14 000512      377 ISPND:  .BYTE  -1             ; : Input suspend flag
15 000513      377 OSPND:  .BYTE  -1             ; : Output suspend flag
16
                .EVEN

```

XLHOOK - Multiterminal Handler Hooks Hook Routine

```

1
2                .SBTTL  XLHOOK  - Multiterminal Handler Hooks Hook Routine
3
4                ;+
5                ;
6                ; XLHOOK
7                ;       Entered from multiterminal input or output interrupt service.
8                ;
9                ; Call (TH.GOC):
10               ;       R0 = Function code
11               ;
12               ; Return (TH.GOC):
13               ;       PSW<c> = 0, R5 = character
14               ;       PSW<c> = 1, no character available
15               ;
16               ; Call (TH.PIC):
17               ;       R0 = Function code
18               ;       R5 = character
19               ;-
20
21               .ENABL  LSB
22

```

The following line must reside before hook entry point:

```

23 000514  113740  XLPNAM: .Rad50  /XL/           ; : Rad50 physical name
24                                     ; loaded by FETCH/LOAD code
25
26 000516      XLHOOK:
27 000516      .Assume <XLHOOK-XLPNAM> EQ 2 MESSAGE=<XLPNAM must precede XLHOOK>
28
29 000516  010446      MOV      R4,-(SP)           ;Save register for awhile
30
31               ; Function code = 1 = TH.GOC
32               ;       (Get Output Character)
33
34 000520  020027      CMP      R0,#TH.GOC       ;'Get Output Character' request?
35               000001
36 000524  001006      BNE      10$              ;Nope...
37 000526  105767      TSTB    OSPND            ;Is output suspended?
38               177761
39 000532  001014      BNE      20$              ;Yep...
40 000534  004767      CALL    HOINT            ;Yes, hook handler output service
41               000370
42 000540  000412      BR      30$
43
44               ; Function code = 2 = TH.PIC
45               ;       (Put Input Character)
46
47 000542  020027  10$:  CMP      R0,#TH.PIC     ;'Put Input Character' request?
48               000002
49 000546  001006      BNE      20$              ;Nope...
50 000550  105767      TSTB    ISPND            ;Is input suspended?
51               177736
52 000554  001003      BNE      20$              ;Yep...
53 000556  004767      CALL    HIINT            ;Yes, hook handler input service
54               000604
55 000562  000401      BR      30$
56
57 000564  000261  20$:  SEC                      ;Return failure
58 000566  012604  30$:  MOV      (SP)+,R4       ;*C* Restore previously saved register
59 000570  000207      RETURN
60
61               .DSABL  LSB

```

PREMTY - Prepare for multiterminal hook

```

1          .SBTTL  PREMTY  - Prepare for multiterminal hook
2
3          ;+
4          ;
5          ; PREMTY
6          ; Prepares for use of a multiterminal hook.
7          ;
8          ; Return:
9          ; R3 -> TCB
10         ;
11         ; Note:
12         ; *** Co-routine ***
13         ; Saves R3
14         ;
15         ;-
16
17 000572 105767 PREMTY: TSTB   O$MTTY           ;Terminal hooks in use?
18         177712
19 000576 001410      BEQ    10$                ;Nope...
20 000600 010346      MOV    R3,-(SP)          ;Save some registers for awhile
21 000602 016703      MOV    TCBADX,R3        ;R3 -> TCB hooked to us
22         002046
23 000606 016646      MOV    2(SP),-(SP)       ;Restack the return address
24         000002
25 000612 004736      CALL   @(SP)+            ;Co-routine back to caller
26 000614 012603      MOV    (SP)+,R3         ;Restore previously saved register
27 000616 005726      TST    (SP)+           ;Discard old return address
28         ; to return to callers caller
29 000620 000207 10$: RETURN
30
31         .ENDC ;NE XL$MTY

```

DRIVER RESET ENTRY

```

1          .SBTTL  DRIVER RESET ENTRY
2
3          ;+
4          ;
5          ; This routine is entered on the abort of a job or an HRESET. It
6          ; dequeues and tells RT that all I/O requests by a job are done. It
7          ; expects to be entered with the number of the aborting job in R4.
8          ;
9          ; Entered with:
10         ; R4 = Job number {aborting | issuing .ABTIO}
11         ; R5 = 0 if abort by job
12         ; -> Channel Control Block (CCB) if abort by channel (.ABTIO)
13         ;
14         ;-
15
16         .ENABL  LSB
17
18 000622 010046 XLDONE: MOV    R0,-(SP)           ;Save R0 for awhile
19
20         .IF EQ XL$PC
21 000624 004767      CALL   DISINI           ;Turn off input interrupts
22         001342
23
24         .IFF ;EQ XL$PC
25 000624 004767      MOV    #RPT.R1,@CSRA        ;Select csr A, write register 1
26 000624 004767      BIC    #W1.RIE,SSRAW1      ;Turn off input interrupts
27 000624 004767      MOV    SSRAW1,@CSRA        ; (update from software register)
28         .ENDC ;EQ XL$PC
29
30 000630 004467      JSR    R4,50$              ; while we remove entries from the
31         000110
32         .WORD  XICQE-60$-Q$LINK             ; input queue
33
34 000636 120467      CMPB   R4,JNUM              ;Is aborting job same as one which
35         177644
36
37         ; issued OFFDRV call?
38 000642 001003      BNE    5$                ;No, so interrupts should still be on
39 000644 005767      TST    STATFG             ;Should we turn interrupts back on?
40         177152
41 000650 001002      BNE    10$                ;Nope...

```

```

36 000652          5$:
37                .IF EQ XL$PC
38 000652 004767    CALL     ENAINI                ;Turn input interrupts back on
                001342
39                .IFF ;EQ XL$PC
40                MOV      #RPT.R1,@CSRA        ;Select csr A, write register 1
41                BIS      #W1.RIE,SSRAW1       ;Turn input interrupts back on
42                MOV      SSRAW1,@CSRA        ; (update from software register)
43                .ENDC ;EQ XL$PC
44
45 000656 005267 10$: INC      QCHG                ;Set the 'queue being modified' flag
                000314
46 000662 004467    JSR      R4,50$                ; while we remove entries from the
                000056
47 000666 177144    .WORD   XOCQE-60$-Q$LINK       ; output queue
48 000670 005067    CLR      QCHG                ;Reset the 'queue being modified' flag
                000302
49
50 000674 120467    CMPB    R4,JNUM                ;Is aborting job same as one which
                177606
51                ;issued OFFDRV call?
52 000700 001003    BNE     15$                ;No, so interrupts should still be on
53 000702 005767    TST    STATFG                ;Again, interrupts back on?
                177114
54 000706 001002    BNE     30$                ;Nope...
55 000710          15$:
56                .IF EQ XL$PC
57 000710 004767    CALL     ENAOUI                ;Turn output interrupts back on
                001356
58                .IFF ;EQ XL$PC
59                MOV      R5,-(SP)            ;Save R5 for awhile
60                CALL    GNXTCH                ;Get a character for output
61                BEQ     20$                ;None available...
62                MOVB   R5,@DBUF                ;Now prime the interrupt pump
63                20$: MOV      (SP)+,R5        ;Restore R5
64                .ENDC ;EQ XL$PC
65
66 000714 012600 30$: MOV      (SP)+,R0        ;Restore R0
67 000716 005767    TST    XLCQE                ;Anything to return to RT?
                177066
68 000722 001001    BNE     40$                ;Yes...

```

Use RETURN if no internally-queued elements are being aborted:

```

69 000724 000207    RETURN                ;Nope, just return
70

```

Use .DRFIN if any abortable queue elements have been placed on the device queue.

NOTE

Only abortable queue elements should be placed on the device queue.

```

71 000726          40$: .DRFIN  XL
72 000726 010704    MOV      PC,R4
73 000730 062704    ADD     #XLCQE-. ,R4
74                177060
75 000734 013705    MOV     @#^o54,R5
76                000054
77 000740 000175    JMP     @^o270(R5)
                000270
72
73                ; The following code scans the internal queue for queue elements which
74                ; match the abort criteria (job number for job abort, channel if abort
75                ; by channel). It then dequeues them from the internal queue, returning
76                ; them to the device queue.
77

```

Internal queuing code. Used to remove abortable queue elements from internal queues:

```

78          000004          SP.CCB = 4          ;Stacked CCB pointer
79          000006          SP.JOB = 6          ;Stacked job number
80
81 000744 010546 50$: MOV R5,-(SP)          ;Save CCB pointer
82 000746 012405          MOV (R4)+,R5          ;Pick up the displacement and
83 000750 010446          MOV R4,-(SP)          ; store the return address
84 000752 060705          ADD PC,R5          ;Calculate actual address
85                                     ; (60$ must follow this)
86 000754 010546 60$: MOV R5,-(SP)          ;Save the Q header address
87 000756 016504 70$: MOV Q$LINK(R5),R4      ;Link to the next entry
      177774
88 000762 001450          BEQ 120$          ;If zero, no more
89 000764 005766          TST SP.CCB(SP)      ;Abort by channel (.ABTIO) ?
      000004
90 000770 001405          BEQ 80$          ;Nope, aborting job...
91 000772 026466          CMP Q$CSW(R4),SP.CCB(SP) ;Yes, this qelement for that channel?
      177776
      000004
92 001000 001037          BNE 110$          ;Nope...
93 001002 000412          BR 90$          ;Yes, go remove it
94
95 001004 116400 80$: MOV Q$JNUM(R4),R0      ;Get number of job being aborted
      000003
96 001010 006200          ASR R0          ; and
97 001012 006200          ASR R0          ; shift
98 001014 006200          ASR R0          ; to
99 001016 042700          BIC #^C<37>,R0      ; isolate job bits
      177740
100 001022 020066          CMP R0,SP.JOB(SP)   ;Job own this queue element?
      000006
101 001026 001024          BNE 110$          ;Nope...
102 001030 016465 90$: MOV Q$LINK(R4),Q$LINK(R5) ;Yes, unlink it from the list
      177774
      177774
103 001036 005064          CLR Q$LINK(R4)      ;Make sure it doesn't link anywhere
      177774
104 001042 005767          TST XLCQE          ;Anything on the queue?
      176742
105 001046 001005          BNE 100$          ;Yes, then link it in at the end
106 001050 010467          MOV R4,XLCQE        ;Otherwise, make it the first
      176734
107 001054 010467          MOV R4,XLLQE        ; and only
      176726
108 001060 000736          BR 70$          ;Check for more elements to abort
109
110 001062 016700 100$: MOV XLLQE,R0          ;R0->element at end of queue
      176720
111 001066 010460          MOV R4,Q$LINK(R0)   ;Link it to this new one
      177774
112 001072 010467          MOV R4,XLLQE        ; and make the new one last
      176710
113 001076 000727          BR 70$          ;Check for more elements to abort
114
115          ; Here if element is not part of the aborting job
116
117 001100 010405 110$: MOV R4,R5          ;Skip this element
118 001102 000725          BR 70$          ;Check for more elements to abort
119
120          ; DeQueue is done, record the new end of the queue
121
122 001104 012604 120$: MOV (SP)+,R4          ;R4->Queue header
123 001106 010564          MOV R5,Q$LINK+2(R4) ;Set the new end of queue
      177776
124 001112 012604          MOV (SP)+,R4          ;Recover the return address
125 001114 012605          MOV (SP)+,R5          ;Restore CCB pointer
126 001116 000204          RTS R4
127
128          .DSABL LSB

```

OUTPUT INTERRUPT SERVICER

```

1          .SBTTL  OUTPUT INTERRUPT SERVICER
2
3          .IF EQ XL$PC
4 001120    .DRAST  XL,XL$PRI,XL$DONE
001120    000640    BR      XLDONE
001122    004577    XLINT::JSR    R5,@$INPTR
001552
001126    000140    .WORD   ^C<XL$PRI*^o40>&^o340
5          .IFF ;EQ XL$PC
6          XOINT:
7          .ENDC ;EQ XL$PC
8
9          .ENABL  LSB
10

```

Hook output interrupt entry point:

```

11 001130    HOINT:          ;Output interrupt hook point
12
13 001130    005727    TST      (PC)+          ;Is break in progress?
14 001132    000000    BRKFLG: .WORD 0          ; : 'break in progress' flag (0=no)
15 001134    001030    BNE     30$          ;Yes, then don't do any output
16 001136    005727    TST      (PC)+          ;Need to send an XON or XOFF?
17 001140    000000    SNDS:   .WORD 0          ; : send XON/XOFF flag
18                                     ; -2 = XON should be sent
19                                     ; -1 = XOFF should be sent
20                                     ; 0 = XON has been sent
21                                     ; 1 = XOFF has been sent
22 001142    100011    BPL     10$          ;Neither...
23 001144    112705    MOVB    #C.CTLQ,R5    ;Assume we are to send an XON
000021
24 001150    062767    ADD     #2,SNDS       ;Are we correct? (SNDS = 0 if yes)
000002
0017762
25 001156    001414    BEQ     20$          ;Yes, go send it
26 001160    112705    MOVB    #C.CTLS,R5    ;No, we must send an XOFF
000023
27 001164    000411    BR      20$          ;Now go send it
28
29 001166    005727    10$:   TST      (PC)+          ;Have we been XOFF'd?
30 001170    000000    RECS:  .WORD 0          ; : received XOFF flag
31 001172    001011    BNE     30$          ;Yes, then don't do any output
32 001174    005727    TST      (PC)+          ;No, are output queues being modified?
33 001176    000000    QCHG:  .WORD 0          ; : 'queues being modified' flag
34 001200    001006    BNE     30$          ;Yes, then don't do any output
35 001202    004767    CALL    GNXTCH       ;Go get a character to output
000024
36 001206    001403    BEQ     30$          ;None available...
37
38 001210    20$:
39 001210    004767    CALL    PUTC          ;Output the character
001276
40
41          .IF EQ XL$PC
42          .IF NE XL$PDT
43          CALL    SETLIT          ;Update the PDT lights display
44          .ENDC ;NE XL$PDT
45 001214    000403    BR      40$
46          .ENDC ;EQ XL$PC
47
48 001216    30$:
49          .IF EQ XL$PC
50 001216    004767    CALL    DISOUI       ;Turn off output interrupts
001022
51          .IFF ;EQ XL$PC
52          MOV     #CMD.RT,@CSRA    ;Reset transmitter interrupt pending
53          MOV     #CMD.EI,@CSRA    ;Declare end of interrupt
54          .ENDC ;EQ XL$PC
55
56 001222    000401    BR      50$
57
58 001224    005727    40$:   TST      (PC)+
59 001226    000261    50$:   SEC
60 001230    000207    60$:   RETURN
61
62          .DSABL  LSB

```

GNXTCH - Get next output character

```

1          .SBTTL  GNXTCH  - Get next output character
2
3          ;+
4          ;
5          ; GNXTCH
6          ;   Obtains the next character from the output queue and returns
7          ;   it in R5.
8          ;
9          ; CALL:
10         ;   CALL   GNXTCH
11         ;
12         ; RETURNS:
13         ;   z-bit = 0, R5 contains character to be output
14         ;   z-bit = 1, no characters available to output
15         ;
16         ; NOTES:
17         ;   As requests are completed, the associated queue elements are
18         ;   returned to RT-11.
19         ;
20         ;-
21
22         .ENABL  LSB
23
24 001232 016704 GNXTCH: MOV   XOCQE,R4          ;R4->current output queue element
25         176656
26 001236 001426      BEQ   10$                ;None available...
27
28         .IF EQ MMG$T
29         ADD   #Q$WCNT,R4          ;R4->word count
30         TST   @R4                 ;Any characters left to output?
31         BEQ   20$                ;Nope, this request is complete
32         INC   @R4                 ;Yes, now there is one less to do
33         MOVB  @-(R4),R5           ;Get the byte to output
34         INC   @R4                 ;bump pointer to next byte
35 001240 005764      TST   Q$WCNT(R4)        ;Any characters left to output?
36         000006
37 001244 001424      BEQ   20$                ;Nope, this request is complete
38 001246 005264      INC   Q$WCNT(R4)        ;Yes, now there is one less to do
39         000006
40 001252 004777      CALL  @$GTBYT          ;Get the byte to output
41         001412
42 001256 112605      MOVB  (SP)+,R5
43         .ENDC ;EQ MMG$T
44
45 001260 046705      BIC   CHMASK,R5        ;Strip the undesired bits
46         176674
47 001264 001762      BEQ   GNXTCH          ; and nulls are not to be suffered
48 001266 006227      ASR   (PC)+          ;Was last character a <CR>?
49 001270 000000      CRFLG: .WORD 0        ; : <CR> flag
50 001272 103003      BCC   5$            ;Nope...
51 001274 120527      CMPB  R5,#C.LF        ;Yes, is this character a <LF>?
52         000012
53 001300 001754      BEQ   GNXTCH          ;Yes, then suppress it...
54 001302 120527      5$:  CMPB  R5,#C.CR    ;Is this character a <CR>?
55         000015
56 001306 001002      BNE   10$            ;Nope...
57 001310 005267      INC   CRFLG          ;Yes, set the flag
58         177754
59 001314 000207      10$:  RETURN
60
61 001316 005267      20$:  INC   QCHG        ;Set the 'queue being modified' flag
62         177654
63
64         .IF EQ XL$PC
65 001322 004767      CALL  DISOUI          ;Shut off the output
66         000716
67
68         .ENDC ;EQ XL$PC
69
70 001326 016704      MOV   XOCQE,R4        ;R4->Current output queue element
71         176562
72 001332 016467      MOV   Q$LINK(R4),XOCQE ;Replace top of output queue with
73         177774
74         176554

```

```

62                                     ; next element
63 001340 004767          CALL    XLFIN          ;Return the element to RT
        000572
64 001344 005067          CLR      QCHG          ;Reset the 'queue being modified' flag
        177626
65
66                                     .IF EQ XL$PC
67 001350 004767          CALL    ENAOUI         ;Restart the output
        000716
68                                     .ENDC ;EQ XL$PC
69
70 001354 000726          BR      GNXTCH
71
72                                     .DSABL LSB

```

INPUT INTERRUPT SERVICER

```

1                                     .SBTTL INPUT INTERRUPT SERVICER
2
3                                     ; This is the input interrupt servicer. Input interrupts are always enabled
4                                     ; once this driver is called for the first time. Only a "RstDrv" SPFUN
5                                     ; request will shut off its interrupt enable.
6
7                                     .IF EQ XL$PC
8 001356          .DRAST  XI,XL$PRI
001356 000207          RETURN
001360 004577 XIINT::JSR    R5,@$INPTR
        001314
001364 000140          .WORD  ^C<XL$PRI*^o40>&^o340
9                                     .IFF ;EQ XL$PC
10 XIINT:
11                                     .ENDC ;EQ XL$PC
12
13                                     .ENABL LSB
14

```

Hook input interrupt entry point:

```

15 001366          HIINT:          ;Input interrupt hook point
16
17 001366 004767          CALL    GETC          ;Get an input character
        001104
18 001372 046705          BIC      CHMASK,R5          ;Strip the undesired bits
        176562
19 001376 001406          BEQ     10$          ; and nulls are not to be suffered
20 001400 120527          CMPB   R5,#C.CTSL          ;Are we being XOFF'd?
        000023
21 001404 001004          BNE     20$          ;Nope...
22 001406 012767          MOV     #1,RECS          ;Yes, set the 'received XOFF' flag
        000001
        177554
23 001414          10$:
24                                     .IF EQ XL$PC
25                                     .IF NE XL$PDT
26 CALL    SETLIT          ;Update the PDT lights display
27                                     .ENDC ;NE XL$PDT
28                                     .IFF ;EQ XL$PC
29 MOV     #CMD.EI,@CSRA          ;Declare end of interrupt
30                                     .ENDC ;EQ XL$PC
31
32 001414 000207          RETURN
33
34 001416 120527          20$:          CMPB   R5,#C.CTLQ          ;Are we being XON'd?
        000021
35 001422 001005          BNE     30$          ;Nope...
36 001424 005067          CLR     RECS          ;Yes, reset the 'received XOFF' flag
        177540
37
38                                     .IF EQ XL$PC
39 001430 004767          CALL    ENAOUI         ;Get the output going again
        000636
40                                     .IFF ;EQ XL$PC
41 CLR     SNDS          ;Indicate that an XON has been
42 MOVB   #C.CTLQ,@DBUF          ; sent
43                                     .ENDC ;EQ XL$PC
44

```

```

45 001434 000767      BR      10$
46
47                ; Here for characters other than XON (^Q) and XOFF (^S)
48
49 001436 005767 30$:  TST      XIBFRE                ;Any room in the input buffer?
                    001170
50 001442 001427      BEQ      50$                ;Nope, go force an XOFF to the host
51
52                ; We have room, so store the character in the ring buffer. It will
53                ; be processed at FORK level.
54
55 001444 016704      MOV      XIBIN,R4                ;Yes, R4=offset into buffer
                    001156
56 001450                .ADDR   #XIBUF,R4,ADD            ;Add address of start of buffer
                    001450 060704      ADD      PC,R4
                    001452 062704      ADD      #XIBUF-.,R4
57 001456 110514      MOVB    R5,@R4                ;Store the character
58 001460 005367      DEC      XIBFRE                ;Buffer has one less free byte now
                    001146
59 001464 005267      INC      XIBIN                ;Bump the offset for next time
                    001136
60 001470                CMP      XIBIN,#BUFSIZ            ;Time to wrap?
                    026727
                    001132
                    000100
61 001476 103402      BLO      40$                ;Nope...
62 001500 005067      CLR      XIBIN                ;Reset the buffer offset
                    001122
63
64                ; Here to check for 'low-water' mark (running out of buffer space)
65
66 001504 026727 40$:  CMP      XIBFRE,#STPSIZ            ;Crossed the 'low-water' mark yet?
                    001122
                    000020
67 001512 101010      BHI      60$                ;Nope, then go process some input
68 001514 005767      TST      SNDS                ;Yes, have we already sent an XOFF?
                    177420
69 001520 003005      BGT      60$                ;Yes, so go process some input
70
71                ; Here to send an XOFF to the host
72
73 001522                50$:
74                .IF EQ XL$PC
75 001522 012767      MOV      #-1,SNDS            ;Request an XOFF to be sent
                    177777
                    177410
76 001530                CALL     ENAOUI                ;Turn on output to make sure
                    000536
77                .IFF ;EQ XL$PC
78                MOV      #1,SNDS            ;Indicate that an XOFF has been
79                MOVB    #C.CTLS,@DBUF      ; sent
80                .ENDC ;EQ XL$PC
81
82                ; Here to process some input
83
84 001534 005767 60$:  TST      XICQE                ;Any requests to satisfy?
                    176334
85 001540 001725      BEQ      10$                ;No, so just return
86
87                .IF NE XL$PC
88                MOV      #CMD.EI,@CSRA      ;Declare end of interrupt
89                .ENDC ;NE XL$PC
90
91 001542                .BR      XIIN
92
93                .DSABL  LSB

```

PROCESS INPUT RECEIVED FROM INTERRUPT SERVICER

```

1          .SBTTL  PROCESS INPUT RECEIVED FROM INTERRUPT SERVICER
2
3          .ENABL  LSB
4
5          ; This routine runs at fork level. It's purpose is to remove characters
6          ; from the ring buffer and use them to satisfy input requests.
7
8 001542 005267 XIIN:  INC    INPRC          ;Did someone beat us to this routine?
          000254
9 001546 001124      BNE    110$          ;Yes..
10
11 001550 004767      CALL   SAV30
          000250
12
13          ; We have the routine. Now we loop to process as much of the input as we can.
14          ; Clear flag to say we own routine and no others can come in. This can be
15          ; done because we are going to check to see if anything is in the input buffer
16          ; after clearing the flag.
17
18 001554      5$:::  CLR    INPRC          ;We're now the owner of this routine
19
20
21
22
23 001554 026727      CMP    XIBFRE,#RSTSIZ          ;Crossed the 'high-water' mark yet?
          001052
          000060
24 001562 103410      BLO    10$          ;Nope...
25 001564 005767      TST    SNDS          ;Yes, have we already sent an XON?
          177350
26 001570 001405      BEQ    10$          ;Yes...
27
28          .IF EQ XL$PC
29 001572 012767      MOV    #-2,SNDS          ;No, then request an XON to be sent
          177776
          177340
30 001600 004767      CALL   ENAOUI          ;Turn on output to make sure
          000466
31          .IFF ;EQ XL$PC
32          CLR    SNDS          ;Now indicate that an XON has been
33          MOVB   #C.CTLQ,@DBUF    ; sent
34          .ENDC ;EQ XL$PC
35
36 001604 016704 10$:  MOV    XICQE,R4          ;Any input requests to satisfy?
          176264
37 001610 001500      BEQ    100$          ;Nope...
38 001612 006227      ASR    (PC)+          ;Time to return an EOF?
39 001614 000000 CTZFPLG: .WORD 0          ; : EOF flag (^Z)
40 001616 103472      BCS    90$          ;Yes...
41 001620 026727      CMP    XIBFRE,#BUFSIZ          ;Anything in the buffer?
          001006
          000100
42 001626 001471      BEQ    100$          ;Nope...
43
44          ; Here to remove a character from the input ring buffer
45
46 001630 016705      MOV    XIBOUT,R5          ;R5=Offset into buffer for next char.
          000774
47 001634      .ADDR   #XIBUF,R5,ADD          ;Add address of start of buffer
          001634 060705      ADD    PC,R5
          001636 062705      ADD    #XIBUF-.,R5
          000670
48 001642 111505      MOVB   @R5,R5          ;Get a character from the ring buffer
49 001644 005267      INC    XIBFRE          ;Buffer has one more free byte
          000762
50 001650 005267      INC    XIBOUT          ;Bump offset for next time
          000754
51 001654 026727      CMP    XIBOUT,#BUFSIZ          ;Time to wrap?
          000750
          000100
52 001662 103402      BLO    20$          ;Nope...
53 001664 005067      CLR    XIBOUT          ;Yes, reset the buffer offset
          000740
54 001670 105764 20$:  TSTB   Q$FUNC(R4)          ;Special function read?
          000002
55 001674 001003      BNE    30$          ;Yes..

```

```

56 001676 120527      CMPB   R5,#C.CTLZ      ;No, is character a ^Z?
    000032
57 001702 001420      BEQ    40$             ;Yes, handle it specially
58
59 001704          30$:
60      .IF EQ MMG$T
61      ADD    #Q$WCNT,R4      ;R4->Word count
62      MOVB   R5,@-(R4)      ;Return the character
63      INC    (R4)+          ;Bump the buffer pointer
64      DEC    @R4            ;Is transfer complete? (z-bit=1 if so)
65      .IFF ;EQ MMG$T
66 001704 110546      MOVB   R5,-(SP)        ;Return the character
67 001706 004777      CALL  @$PTBYT        ; ...
    000760
68 001712 005364      DEC    Q$WCNT(R4)     ;Is transfer complete? (z-bit=1 if so)
    000006
69      .ENDC ;EQ MMG$T
70
71 001716 001422      BEQ    70$             ;Yes...
72 001720 026727      CMP    XIBFRE,#BUFSIZ ;Anything left in buffer?
    000706
    000100
73 001726 001312      BNE    5$             ;Yes, go process it
74 001730 016704      MOV    XICQE,R4      ;R4->Input request queue element
    176140
75 001734 105764      TSTB   Q$FUNC(R4)    ;Special request?
    000002
76 001740 001705      BEQ    5$             ;Nope, process some more input
77 001742 000402      BR     50$           ;Yes, then request is done
78
79 001744 005267 40$:  INC    CTZFLG        ;Set the EOF flag
    177644
80
81 001750          50$:
82      .IF EQ MMG$T
83      ADD    #Q$WCNT,R4      ;R4->word count
84      CLRB   @-(R4)        ;Return a zero byte
85      INC    (R4)+          ;Bump the buffer pointer
86      DEC    @R4            ;Is the transfer complete?
87      BNE    60$           ;Nope...
88      .IFF ;EQ MMG$T
89 001750 105046      CLRB   -(SP)         ;Return a zero byte
90 001752 004777      CALL  @$PTBYT        ; ...
    000714
91 001756 005364      DEC    Q$WCNT(R4)     ;Is the transfer complete?
    000006
92 001762 001372      BNE    50$           ;Nope...
93      .ENDC ;EQ MMG$T
94
95 001764 016704 70$:  MOV    XICQE,R4      ;R4->Current input queue element
    176104
96 001770 016467      MOV    Q$LINK(R4),XICQE ;Replace top of input queue with
    177774
    176076
97      ; next queue element
98 001776 004767 80$:  CALL  XLFIN          ;Return the element to RT
    000134
99 002002 000664      BR     5$             ;And check for more input
100
101 002004 052754 90$:  BIS    #EOF$,@-(R4)   ;Indicate EOF
    020000
102 002010 000765      BR     70$           ;And declare queue element done
103
104 002012          100$:
105      ;;      DEC    INPRC        ;Did anything else come in while
106      ;;      ; we were otherwise occupied?
107      ;;      BPL    5$             ;Yes, then go process it
108 002012 012767      MOV    #-1,INPRC     ;Release the input processing routine
    177777
    000002
109 002020 000207 110$:  RETURN              ;Nope, then we'll retire for awhile
110
111      ; This flag is -1 when no one is executing the XIIN routine.
112      ; It is zero when someone is executing in the XIIN routine.
113      ; It becomes greater than zero to indicate that more input has come
114      ; in while someone was executing the XIIN routine.
115
116 002022 177777  INPRC: .WORD  -1

```

```

117
118             .DSABL LSB
119
120             ; The following routine is used by XIIN to simulate the effects of a
121             ; FORK (saving of registers 0-3 and lowering of priority)
122
123 002024 010046 SAV30: MOV     R0,-(SP)           ;Save some registers
124 002026 010146      MOV     R1,-(SP)           ; ...
125 002030 010246      MOV     R2,-(SP)           ; ...
126 002032 010346      MOV     R3,-(SP)           ; ...
127 002034 016646      MOV     10(SP),-(SP)       ;Restack the return address
           000010
128 002040             .MTPS  #0                 ;Lower our priority
           002040 005046      CLR     -(SP)
           002042 112716      MOV     #0,(SP)
           000000
           002046 013746      MOV     @#^o54,-(SP)
           000054
           002052 062716      ADD     #^o360,(SP)
           000360
           002056 004736      CALL   @(SP)+
129 002060 004736      CALL   @(SP)+           ;Co-routine back to caller
130 002062 012603      MOV     (SP)+,R3           ;Restore the registers
131 002064 012602      MOV     (SP)+,R2           ; ...
132 002066 012601      MOV     (SP)+,R1           ; ...
133 002070 012600      MOV     (SP)+,R0           ; ...
134 002072 005726      TST     (SP)+           ;Discard old return address
135 002074 000207      RETURN          ; and return to caller's caller

```

XLENQ - Place Qelement on internal queue

```

1             .SBTTL XLENQ - Place Qelement on internal queue
2
3             ;+
4             ;
5             ; XLENQ
6             ;
7             ; Removes the current Qelement from the device queue and places
8             ; it on an internal queue. It is presumed (by virtue of the way
9             ; RT works) that there will be only one Qelement in the device
10            ; queue.
11            ; Call:
12            ;
13            ; R4 -> Qelement to be queued
14            ;
15            ; JSR     R5,Q
16            ; .BLKW           ;CQE pointer of internal queue
17            ; .BLKW           ;LQE pointer of internal queue
18            ; Return:
19            ;
20            ; Qelement has been removed from the device queue and placed on
21            ; the specified internal queue.
22            ;-
23

```

Internal queuing code; moves queue element to an internal queue:

```

24 002076 005067 XLENQ: CLR     XLCQE           ;Ensure there are no Qelements
           175706
25 002102 005067      CLR     XLLQE           ;
           175700
26 002106 005715      TST     @R5                 ;Is our internal queue empty?
27 002110 001003      BNE     10$                 ;Nope...
28 002112 010425      MOV     R4,(R5)+           ;Yes, so make it the first
29 002114 010425      MOV     R4,(R5)+           ; and last element
30 002116 000205      RTS     R5
31
32 002120 005725 10$: TST     (R5)+           ;Bump to last element pointer
33 002122 010446      MOV     R4,-(SP)           ;Save address of new element
34 002124 011504      MOV     @R5,R4             ;R4->Last queue element
35 002126 011664      MOV     @SP,Q$LINK(R4)    ;Link it to the new element
           177774
36 002132 012625      MOV     (SP)+,(R5)+       ; and make the new element the last
37 002134 000205      RTS     R5

```

XLFIN - Internal Queue Element Completion

```

1          .SBTTL XLFIN - Internal Queue Element Completion
2
3          ;+
4          ;
5          ; XLFIN
6          ; Used to inform RT-11 of a Qelement which has completed.
7          ;
8          ; Call:
9          ; R4 -> Completed Qelement
10         ;
11         ; Return:
12         ; Qelement has been returned to RT-11
13         ;
14         ; Note:
15         ; o All registers except R4 are preserved
16         ; o Fake device queue is used to return the Qelement to
17         ; RT-11 to avoid race conditions with the real device
18         ; queue.
19         ; o A CALL to monitor completion is used because there may
20         ; be more to do at this time, we don't want to lose control
21         ; to the monitor yet.
22         ;
23         ;-
24

```

Internal queuing code; returns queue element, using fake device queue:

```

25 002136 010467 XLFIN: MOV R4,XLFCQE ;Queue element we are returning will
000520
26 002142 010467 MOV R4,XLFLQE ; become first and last element
000512
27 002146 005064 CLR Q$LINK(R4) ;Unlink it from everything else
177774
28 002152 .ADDR #XLFCQE,R4 ;R4 -> Fake device queue for passing
002152 010704 MOV PC,R4
002154 062704 ADD #XLFCQE--,R4
000506
29 ; to DRFIN
30 002160 013705 MOV @#$SYPTR,R5 ;R5->$RMON
000054

```

Modified form of .DRFIN, used to return queue element to monitor and gain control when monitor is done. Required for hooks support if queue element completes as a result of call from multiterminal service:

```

31 002164 004775 CALL @$QCOMP(R5) ;Inform monitor of I/O completion
000270
32 002170 000207 RETURN
33
34 .IF EQ XL$PC

```

DISINI - Disable input interrupts ENAINI - Enable input interrupts

```

1          .SBTTL DISINI - Disable input interrupts
2          .SBTTL ENAINI - Enable input interrupts
3
4 002172          DISINI:
5          .IF NE XL$MTY
6 002172 105767 TSTB O$MTTY ;Terminal hooks in use?
176312
7 002176 001404 BEQ 10$ ;Nope...
8 002200 112767 MOVB #-1,ISPND ;Disable input interrupt processing
177777
176304
9 002206 000403 BR 20$
10         .ENDC ;NE XL$MTY
11
12 002210 042777 10$: BIC #RC.IE,@XIS ;Turn off input interrupts
000100

```

```

13 002216 175714 000207 20$: RETURN
14
15 002220 ENAINI:
16 .IF NE XL$MTY
17 002220 105767 TSTB O$MTTY ;Terminal hooks in use?
176264
18 002224 001403 BEQ 10$ ;Nope...
19 002226 105067 CLRB ISPND ;Enable input interrupt processing
176260
20 002232 000403 BR 20$
21 .ENDC ;NE XL$MTY
22
23 002234 052777 10$: BIS #RC.IE,@XIS ;Turn input interrupts back on
000100
175670
24 002242 000207 20$: RETURN

```

DISOUI - Disable output interrupts

ENAOUI - Enable output interrupts

```

1 .SBTTL DISOUI - Disable output interrupts
2 .SBTTL ENAOUI - Enable output interrupts
3
4 002244 DISOUI:
5 .IF NE XL$MTY
6 002244 105767 TSTB O$MTTY ;Terminal hooks in use?
176240
7 002250 001404 BEQ 10$ ;Nope...
8 002252 112767 MOVB #-1,OSPND ;Disable output interrupt processing
177777
176233
9 002260 000403 BR 20$
10 .ENDC ;NE XL$MTY
11
12 002262 042777 10$: BIC #XC.IE,@XOS ;Disable output interrupts
000100
175646
13 002270 000207 20$: RETURN
14
15 002272 ENAOUI:
16 .IF NE XL$MTY
17 002272 004767 CALL PREMTY ;Prepare for hook
176274
18 002276 001405 BEQ 10$ ;Terminal hooks not active...
19 002300 105067 CLRB OSPND ;Enable output interrupt processing
176207
20 002304 004777 CALL @MTOENX ; and then enable output interrupts
000334
21 002310 000403 BR 20$
22 .ENDC ;NE XL$MTY
23
24 002312 052777 10$: BIS #XC.IE,@XOS ;Enable output interrupts
000100
175616
25 002320 000207 20$: RETURN

```

RESBRK - Turn off BREAK
SETBRK - Turn on BREAK

```

1          .SBTTL RESBRK - Turn off BREAK
2          .SBTTL SETBRK - Turn on BREAK
3
4 002322      RESBRK:
5              .IF NE XL$MTY
6 002322 004767      CALL    PREMTY          ;Prepare for hook
              176244
7 002326 001404      BEQ     10$           ;Terminal hooks not active...
8 002330 005000      CLR     R0            ;Deassert BREAK
9 002332 004777      CALL    @MTYBRX       ; ...
              000310
10 002336 000403     BR      20$
11              .ENDC ;NE XL$MTY
12
13 002340 042777 10$: BIC     #XC.BRK,@XOS   ;Deassert BREAK
              000001
              175570
14 002346 000207 20$: RETURN
15
16 002350      SETBRK:
17              .IF NE XL$MTY
18 002350 004767      CALL    PREMTY          ;Prepare for hook
              176216
19 002354 001405      BEQ     10$           ;Terminal hooks not active...
20 002356 012700      MOV     #XC.BRK,R0     ;Assert BREAK
              000001
21 002362 004777      CALL    @MTYBRX       ; ...
              000260
22 002366 000403     BR      20$
23              .ENDC ;NE XL$MTY
24
25 002370 052777 10$: BIS     #XC.BRK,@XOS   ;Assert BREAK
              000001
              175540
26 002376 000207 20$: RETURN

```

GETSTT - Get line status
RESSTT - Reset line state bits
SETSTT - Set line state bits

```

1          .SBTTL GETSTT - Get line status
2          .SBTTL RESSTT - Reset line state bits
3          .SBTTL SETSTT - Set line state bits
4
5          ;+
6          ;
7          ; GETSTT
8          ; Returns the current line status
9          ;
10         ; Call:
11         ; none
12         ;
13         ; Return:
14         ; R0 = Line status
15         ;
16         ; Note:
17         ; R3 is altered
18         ;
19         ;-
20
21 002400      GETSTT:
22              .IF NE XL$MTY
23 002400 004767      CALL    PREMTY          ;Prepare for hook
              176166
24 002404 001403      BEQ     10$           ;Terminal hooks not active...
25 002406 004777      CALL    @MTYSTX       ;Get current line status
              000240
26 002412 000402     BR      20$
27              .ENDC ;NE XL$MTY
28
29 002414 017700 10$: MOV     @XIS,R0        ;R0 = Current line status
              175512
30 002420 000207 20$: RETURN
31
32         ;+

```

```

33          ;
34          ; RESSTT
35          ;      Deasserts line state bits
36          ;
37          ; Call:
38          ;      R0 = Bits to deassert
39          ;
40          ; Return:
41          ;      R0 = Updated line status
42          ;
43          ; Note:
44          ;      o R3 is altered
45          ;
46          ;      o Unlike SETSTT, which sets the bits as specified,
47          ;      this routine first reads the status and then
48          ;      deasserts the undesired bits.
49          ;
50          ;-
51
52 002422 010046 RESSTT: MOV    R0,-(SP)          ;Save bits to deassert
53 002424 004767 CALL    GETSTT          ;Get current status
                    177750
54 002430 042600 BIC    (SP)+,R0        ;deassert the desired bits
55
56          .IF NE XL$MTY
57 002432 004767 CALL    PREMTY          ;Prepare for hook
                    176134
58 002436 001403 BEQ    10$              ;Terminal hooks not active...
59 002440 004777 CALL    @MTYCTX          ;Yes, set new line state
                    000204
60 002444 000402 BR     20$              ;
61          .ENDC ;NE XL$MTY
62
63 002446 010077 10$:  MOV    R0,@XIS          ;Set new line status
                    175460
64 002452 000207 20$:  RETURN
65
66          ;+
67          ;
68          ; SETSTT
69          ;      Asserts line state bits
70          ;
71          ; Call:
72          ;      R0 = Bits to assert
73          ;
74          ; Return:
75          ;      R0 = Updated line status
76          ;
77          ; Note:
78          ;      o R3 is altered
79          ;
80          ;      o Unlike RESSTT, which first reads the status and
81          ;      deasserts the undesired bits, this routine simply
82          ;      asserts the desired bits.
83          ;
84          ;-
85
86 002454          SETSTT:
87          .IF NE XL$MTY
88 002454 004767 CALL    PREMTY          ;Prepare for hook
                    176112
89 002460 001403 BEQ    10$              ;Terminal hooks not active...
90 002462 004777 CALL    @MTYCTX          ;Yes, set desired bits
                    000162
91 002466 000402 BR     20$              ;
92          .ENDC ;NE XL$MTY
93
94 002470 050077 10$:  BIS    R0,@XIS          ;Set new line status
                    175436
95 002474 000207 20$:  RETURN
96
97          .ENDC ;EQ XL$PC

```

GETC - Input a character
PUTC - Output a character

```

1          .SBTTL  GETC    - Input a character
2          .SBTTL  PUTC    - Output a character
3
4          ;+
5          ;
6          ; GETC
7          ;      Gets a character from the interface.
8          ;
9          ; Return:
10         ;      R5 = Character
11         ;
12         ; Note:
13         ;      In the case of call during multiterminal hook operation,
14         ;      the character is already in R5 due to the multiterminal
15         ;      input interrupt service code.
16         ;
17         ;-
18
19 002476      GETC:
20         .IF NE XL$MTY
21 002476 105767  TSTB    O$MTTY          ;Terminal hooks in use?
22         176006
23 002502 001002  BNE     10$            ;Yep, bypass normal DL input
24         .ENDC ;NE XL$MTY
25
26 002504 117705  MOVB    @XIB,R5          ;R5 = Character
27         175424
28         .IFF ;EQ XL$PC
29         MOVB    @DBUF,R5          ;Get a character from input
30         .ENDC ;EQ XL$PC
31 002510 000207 10$:  RETURN
32
33         ;+
34         ;
35         ; PUTC
36         ;      Puts a character to the interface.
37         ;
38         ; Call:
39         ;      R5 = Character
40         ;
41         ; Note:
42         ;      In the case of call during multiterminal hook operation,
43         ;      the character is already in R5 due to the multiterminal
44         ;      input interrupt service code.
45         ;
46         ;-
47
48 002512      PUTC:
49         .IF NE XL$MTY
50 002512 105767  TSTB    O$MTTY          ;Terminal hooks in use?
51         175772
52 002516 001002  BNE     10$            ;Yep, bypass normal DL output
53         .ENDC ;NE XL$MTY
54
55 002520 110577  MOVB    R5,@XOB          ;Output the character
56         175414
57         .IFF ;EQ XL$PC
58         MOVB    R5,@DBUF          ;Output the character
59         .ENDC ;EQ XL$PC
60 002524 000207 10$:  RETURN

```

INPUT BUFFER AREA

```

1          .SBTTL  INPUT BUFFER AREA
2

```

Internal receive buffer:

```

3           ; Reserve space for the input buffer and data to manage the input buffer
4
5 002526      XIBUF:  .BLKB  BUFSIZ           ;Input buffer
6 002626 000000 XIBIN:  .WORD  0           ;'Next Character In' offset
7 002630 000000 XIBOUT: .WORD  0           ;'Next Character Out' offset
8 002632 000100 XIBFRE: .WORD  BUFSIZ       ;Number of free bytes in buffer
9
10          ; Define areas for fork blocks used by the interrupt servicers
11
12 002634 000000 DQFBLK: .WORD  0,0,0,0
13          002636 000000
14          002640 000000
15          002642 000000
16
17          .IF NE XL$MTY
18
19
20

```

Handler hooks code; pointers loaded by LOAD code, used to reach hooks routines in multiterminal monitor:

```

16          ; Multiterminal handler hooks pointers
17
18 002644      MTOENX: .BLKW           ; : -> Output enable routine
19 002646      MTYBRX: .BLKW           ; : -> Break control routine
20 002650      MTYCTX: .BLKW           ; : -> Line control routine
21 002652      MTYSTX: .BLKW          ; : -> Line status routine
22 002654      TCBADX: .BLKW          ; : -> TCB we're attached to
23          .ENDC ;NE XL$MTY
24
25          ; Fake queue header for returning completed Qelements
26

```

Internal queuing—fake device queue. Zero word required to simulate non-held handler:

```

27 002656 000000          .WORD  0
28 002660          XLFLQE: .BLKW
29 002662          XLFCQE: .BLKW
30
31 002664          .DREND  XL
32          002664 000000 $RLPTR: .WORD  0
33          002666 000000 $MPPTR: .WORD  0
34          002670 000000 $GTBYT: .WORD  0
35          002672 000000 $PTBYT: .WORD  0
36          002674 000000 $PTWRD: .WORD  0
37          002676 000000 $TIMIT: .WORD  0
38          002700 000000 $INPTR: .WORD  0
39          002702 000000 $FKPTR: .WORD  0
40
41          .IF EQ XL$PC
42

```

LOAD - Handler FETCH/LOAD code

```

1           .SBTTL  LOAD    - Handler FETCH/LOAD code
2
3           ;+
4           ;
5           ; LOAD
6           ; This routine is entered on FETCH or LOAD of the XL handler
7           ; and is used 1) to verify use of the handler in the specific
8           ; configuration and, if needed, 2) to establish the required
9           ; connections between the handler and the interrupt service of
10          ; a monitor with support for multiterminal handler hooks.
11          ;
12          ;-
13
14          .ENABL  LSB
15
16 002704      FETCH::
17 002704      LOAD::
18 002704 010567      MOV     R5,ENTRY$           ;Save entry point
19          000314
19 002710 010267      MOV     R2,SLOT$           ; and table size
20          000312

```

```

20 002714 011505      MOV    @R5,R5          ;R5 -> Base of handler (in memory)
21 002716 013700      MOV    @#$SYPTR,R0    ;R0 -> Base of RMON
      000054
22

```

Hooks code. Establishes linkages between handler and TCB:

```

23      .IF NE XL$MTY
24 002722 105765      TSTB   <O$MTTY-XLLQE>(R5) ;Terminal hooks to be used?
      000502
25 002726 001463      BEQ    20$            ;Then use normal DL
26 002730 016001      MOV    $THKPT(R0),R1   ;R1 -> Multiterminal handler hooks
      000000G
      ; data structure in RMON
28 002734 001531      BEQ    60$            ;Monitor doesn't have the support...
29 002736 105721      TSTB   (R1)+          ;Bypass structure size byte
30 002740 112102      MOV    (R1)+,R2       ;R2 = Number of LUNs on system
31 002742 012103      MOV    (R1)+,R3       ;R3 -> TCB list
32 002744 012165      MOV    (R1)+,<MTOENX-XLLQE>(R5) ;Set pointer to output enable routine
      002636
33 002750 012165      MOV    (R1)+,<MTYBRX-XLLQE>(R5) ;Set pointer to Break control routine
      002640
34 002754 012165      MOV    (R1)+,<MTYCTX-XLLQE>(R5) ;Set pointer to Control routine
      002642
35 002760 012165      MOV    (R1)+,<MTYSTX-XLLQE>(R5) ;Set pointer to Status routine
      002644
36 002764 116500      MOV    <O$LINE-XLLQE>(R5),R0 ;R0 = Line to attach to
      000503
37 002770 100513      BMI    60$            ;Must be a positive number
38 002772 120002      CMP    R0,R2          ;Is line in this configuration?
39 002774 002111      BGE    60$            ;Nope, invalid line number
40 002776 006300      ASL    R0              ;Shift for word offset into TCB list
41 003000 060003      ADD    R0,R3          ;R3 -> TCB list entry
42 003002 011303      MOV    @R3,R3         ;R3 -> TCB for LUN
43 003004 005763      TST    T.CSR(R3)     ;Is the line present in hardware?
      000016
44 003010 001503      BEQ    60$            ;Nope...
45 003012 005763      TST    T.STAT(R3)    ;Is the line a console?
      000014
46
47 003016      .Assume CONSL$ EQ 100000
48 003016 100500      BMI    60$            ;Yes...
49 003020 010500      MOV    R5,R0          ;R0 -> Handler hook routine
50 003022 062700      ADD    #<XLHOOK-XLLQE>,R0 ; ...
      000510
51 003026 005763      TST    T.OWNR(R3)    ;Is the line already attached?
      000012
52 003032 001403      BEQ    10$            ;Nope...
53 003034 020063      CMP    R0,T.OWNR(R3) ;Yes, to this handler?
      000012
54 003040 001067      BNE    60$            ;Nope...
55 003042 016701 10$: MOV    ENTRY$,R1       ;R1 -> $ENTRY entry
      000156
56 003046 166701      SUB    SLOT$,R1       ;R1 -> $PNAME ENTRY
      000154
57 003052 011160      MOV    @R1,-2(R0)     ;Inform handler of its physical name,
      177776
58 003056 010365      MOV    R3,<TCBADX-XLLQE>(R5) ; link the handler to the TCB
      002646

```

HANMC\$ disables RT-11 processing of modem control; handler will process modem:

```

59 003062 052763      BIS    #<HANMT$!HANMC$>,T.STAT(R3) ; declare line owned by handler
      000000C
      000014
60      ; and that handler will process modem,
61 003070 010063      MOV    R0,T.OWNR(R3) ; finally link the TCB to the handler
      000012
62 003074 000450      BR     50$
63      .ENDC ;NE XL$MTY
64

```

The following code protects against vector corruption. Won't allow use of handler in NOMTTY mode if CSR or vector conflicts with a line in multiterminal configuration:

```

65 003076 032760 20$: BIT #MTTY$, $SYSGE(R0) ;Is this a multiterminal monitor?
    020000
    000372
66 003104 001444 BEQ 50$ ;Nope, then there can't be a conflict
67 003106 .ADDR #MTAREA,R0 ;R0 -> .MTSTAT EMT area
    003106 010700 MOV PC,R0
    003110 062700 ADD #MTAREA-. ,R0
    000120
68 003114 .ADDR #MTSTAT,R1 ;R1 -> Status block
    003114 010701 MOV PC,R1
    003116 062701 ADD #MTSTAT-. ,R1
    000120
69 003122 .MTSTA R0,R1 ;Get info about multiterminal system
    003122 012710 MOV #31.*^o400+8.,@R0
    017410
    003126 010160 MOV R1,2.(R0)
    000002
    003132 005060 CLR 4.(R0)
    000004
    003136 104375 EMT ^o375
70 003140 103427 BCS 60$ ;Errors?
71 003142 013700 MOV @#$SYPTR,R0 ;R0 -> $RMON
    000054
72 003146 016701 MOV MTSTAT,R1 ;R1 -> First TCB in system
    000064
73 003152 060001 ADD R0,R1 ; ...
74 003154 016702 MOV MTSTAT+MST.LU,R2 ;R2 = Highest LUN on the system
    000062
75 ; (Number_of_LUNs - 1)
76 003160 005761 30$: TST T.CSR(R1) ;Is this a configured line?
    000016
77 003164 001410 BEQ 40$ ;Nope...
78 003166 026561 CMP <XIS-LLQE>(R5),T.CSR(R1) ;Will use of the CSR conflict?
    000124
    000016
79 003174 001411 BEQ 60$ ;Yes, reject the load
80 003176 026561 CMP <XL$VTB-LLQE>(R5),T.VEC(R1) ;Will use of the VECTOR conflict?
    000134
    000020
81 003204 001405 BEQ 60$ ;Yes, reject the load
82 003206 066701 40$: ADD MTSTAT+MST.ST,R1 ;On to next TCB
    000032
83 003212 005302 DEC R2 ;More TCB's to check?
84 003214 002361 BGE 30$ ;Yep...
85 003216 .BR 50$ ;Nope, use of interface won't conflict
86
87 003216 005727 50$: TST (PC)+ ;Success return
88 003220 000261 60$: SEC ;Error return
89 003222 000207 RETURN
90
91 003224 ENTRY$: .BLKW ; : -> $ENTRY table entry
92 003226 SLOTS$: .BLKW ; : Size of a monitor handler table
93
94 003230 MTAREA: .BLKW 3 ; : EMT area for .MTSTAT
95 003236 MTSTAT: .BLKW 8. ; : Status block from .MTSTAT
96
97 .DSABL LSB
98
99 .ENDC ;EQ XL$PC

```

UNLOAD - UNLOAD/.RELEASE CODE

```

1 .SBTTL UNLOAD - UNLOAD/.RELEASE CODE
2
3 ;+
4 ; UNLOAD
5 ; On entry due to unload command, verifies interrupts have been
6 ; disabled unless the handler is still in use, indicated by
7 ; non-empty internal queues.
8 ;
9 ; On entry due to .RELEASE directive,disable interrupts
10 ;
11 ;-
12
13 .ENABL LSB
14

```

Prevents unload if internal queues are not empty:

```

15 003256          UNLOAD::
16 003256 011505    MOV     @R5,R5           ;R5 -> Handler entry point (XLLQE)
17 003260 005765    TST    <STATFG-XLLQE>(R5) ;Is handler in use?
                                000014
18 003264 001013    BNE    10$             ;Nope, it can be unloaded...
19 003266 016546    MOV    <XICQE-XLLQE>(R5),-(SP) ;Check internal queues
                                000066
20 003272 056526    BIS    <XOCQE-XLLQE>(R5),(SP)+ ; ...
                                000106
21 003276 001405    BEQ    RELEAS          ;They're empty...
22 003300          .ADDR  #NOUNLO,R0      ;R0 -> Error message string
003300 010700    MOV    PC,R0
003302 062700    ADD    #NOUNLO-.,R0
                                000106
23                                ; (KMON reports error)
24 003306 000261    SEC
25 003310 000207    RETURN                ;Indicate error
                                ; and return to KMON
26
27 003312          RELEAS::
28 003312 011505    MOV     @R5,R5           ;R5 -> Handler entry point (XLLQE)
29 003314          10$:
30                                .IF EQ XL$PC
31                                .IF NE XL$MTY

```

Handler hooks code; disconnects TCB and handler:

```

32 003314 105765    TSTB   <O$MTTY-XLLQE>(R5) ;Terminal hooks in use?
                                000502
33 003320 001420    BEQ    20$             ;Nope...
34 003322 016501    MOV    <TCBADX-XLLQE>(R5),R1 ;R1 -> TCB we're hooked to
                                002646
35 003326 001426    BEQ    30$             ;We're not...
36 003330 004765    CALL   <DISINI-XLLQE>(R5) ;Disable input
                                002164
37 003334 004765    CALL   <DISOUI-XLLQE>(R5) ; and output interrupts
                                002236
38 003340 005000    CLR    R0              ;Deassert all modem control bits
39 003342 004765    CALL   <SETSTT-XLLQE>(R5) ; ...
                                002446
40 003346 005061    CLR    T.OWNR(R1)      ;Disconnect TCB from handler
                                000012
41 003352 042761    BIC    #<HANMT$!HANMC$>,T.STAT(R1) ; ...
                                000000C
                                000014
42 003360 000411    BR     30$
43 003362          20$:
44                                .ENDC ;NE XL$MTY
45 003362 016501    MOV    <XIS-XLLQE>(R5),R1 ;R1->Device register base
                                000124
46 003366 042711    BIC    #RC.IE,@R1      ;Turn off input and
                                000100
47 003372 042761    BIC    #XC.IE,4(R1)    ;Output interrupts
                                000100
                                000004
48 003400 042711    BIC    #RC.DTR,@R1     ;Now turn off DTR
                                000002
49                                .IFF ;EQ XL$PC
50                                MOV    #RPT.R1,@#XL$CSA ;Select csr A,write register 1
51                                CLR    @#XL$CSA ;Turn off input and output interrupts
52                                BIC    #<M0.DTR>,@#XL$MCO ;Now turn off DTR
53                                .ENDC ;EQ XL$PC
54
55 003404 000241 30$: CLC
56 003406 000207    RETURN
57
58 003410          NOUNLO: .NLCSI  TYPE=I, PART=PREFIX
003410 077          .ASCII  "?XL-"
.
.
.
59 003414 106          .ASCIZ  "F-Handler may not be unloaded while in use"
.
.
.

```

```

60
61             .DSABL  LSB
62
63             000001 .END

```

Symbol table

```

ABTIO$ 001000          DVM.NS 000001          JNUM   000506R   002
BATCH$ 000010          DV2.V2 040000          JOBMK   000370
BRKDRV= 000202          DZ11$  010000          KT11$  010000
BRKFLG 001132R        002 EIS$  000400          KW11P$ 040000
BUFSIZ= 000100          ENAINI 002220R        002 KXCPU$ 004000
BUS$    000100          ENAOUI 002272R        002 LDREL$ 000020
BUS$C   020000          ENTRY$ 003224R        002 LIGHT$ 000010
BUS$M   020100          EOF$   020000          LKCS$  020000
BUS$Q   000100          ERLG$  000001          LOAD   002704RG   002
BUS$U   000000          ERL$G = 000000          LS11L$ 004000
BUS$X   020100          FBMON$ 000001          MMTG$  000002
CACHE$  000001          FETCH  002704RG        002 MMGT$ = 000001
CHMASK 000160R        002 FILL$ 000001          MPTY$S$ 001000
CIS$    000200          FILST$ 100000          MPTY$  000002
CLK50$  000040          FIX$ED= 000001          MST.CT 000002
CLOCK$  100000          FJOB$  000200          MST.LU 000004
CLRDRV= 000201          FPUL1$ 000400          MST.ST 000006
CONSL$  100000          GETC   002476R        002 MST.SZ 000020
CRFLG   001270R        002 GETSTT 002400R        002 MST.1T 000000
CSRSV   000250          GNXTCH 001232R        002 MTAREA 003230R   002
CTRLC$  040000          GSCCA$ 010000          MTOENX 002644R   002
CTRLU$  000002          GTLNK$ 000400          MTSTAT 003236R   002
CTZFLG 001614R        002 HANMC$= ***** GX  MTTY$  020000
C.CR    = 000015          HANMT$= ***** GX  MTYBRX 002646R   002
C.CTLQ= 000021          HDERR$ 000001          MTYCTX 002650R   002
C.CTLS= 000023          HIINT  001366R        002 MTYSTX 002652R   002
C.CTLZ= 000032          HNDLR$ 004000          NOUNLO 003410R   002
C.LF   = 000012          HNGUP$ 004000          OFFDRV= 000205
DBGSY$  002000          HOINT  001130R        002 OSPND  000513R   002
DH11$  020000          HS2.BI 000001          O$LINE 000511R   002
DISCSR 000174          HS2.KI 000002          O$MTTY 000510R   002
DISINI 002172R        002 HS2.KL 000004          O.CSR  000442
DISOUI 002244R        002 HS2.KU 000010          O.ERR  000616
DOC$UN= 000000          HS2.MO 000020          O.LINE 000534
DQFBLK 002634R        002 HWDSP$ 000004          O.MTTY 000546
DTACH$  000020          HWFPU$ 000100          O.NOR  000614
DTRDRV= 000206          H1.ABT 001002          O.VEC  000514
DVC.CT  000006          H1.BR  001014          PAGE$  000200
DVC.DE  000010          H1.CQE 001010          PDP60$ 100000
DVC.DK  000004          H1.FG2 001016          PDP70$ 040000
DVC.DL  000012          H1.FLG 001010          PREMTY 000572R   002
DVC.DP  000011          H1.HLD 001004          PROS$  020000
DVC.LP  000007          H1.LDT 001024          PS     177776
DVC.MT  000005          H1.LQE 001006          PUTC   002512R   002
DVC.NI  000013          H1.NDF 001026          P1$EXT 000432
DVC.NL  000001          H1.NOP 001012          QCHG   001176R   002
DVC.PS  000014          H1.SCK 001020          QUEUE$ 002000
DVC.SB  000020          H1.SDF 001022          Q$BLKN 000000
DVC.SI  000016          H1.VEC 001000          Q$BUFF 000004
DVC.SO  000017          INCV$  000400          Q$COMP 000010
DVC.TP  000003          INEXP$ 000100          Q$CSW  177776
DVC.TT  000002          INPRC  002022R        002 Q$FUNC 000002
DVC.UK  000000          INSCSR 000176          Q$JNUM 000003
DVC.VT  000015          INSDAT 000200          Q$LINK 177774
DVM.DM  000002          INSSYS 000202          Q$MEM  000014
DVM.DX  000001          ISPND  000512R        002 Q$PAR  000012
DVM.NF  000200          I$MTTY 000244          Q$UNIT 000003
Q$WCNT 000006          THK.OE 000004          XIBUF  002526R   002
Q.BLKN  000004          THK.ST 000012          XICQE  000074R   002
Q.BUFF  000010          THK.SZ 000014          XIIN   001542R   002
Q.COMP  000014          THK.TC 000002          XIINT  001360RG   002
Q.CSW   000002          TH.GOC 000001 G     XILQE  000076R   002
Q.ELGH 000024          TH.PIC 000002 G     XIS    000132R   002
Q.FUNC  000006          TIMER$ 002000          XITSW$ 000040
Q.JNUM  000007          TIMIT$ 000004          XLCQE  000010RG   002
Q.LINK  000000          TIM$IT= 000001          XLDONE 000622R   002
Q.MEM   000020          TSXP$  100000          XLDSIZ 000000
Q.PAR   000016          TTBFI  000206          XLEND = 002704RG   002
Q.UNIT  000007          TTBFI$O 000050          XLENQ  002076R   002

```

Q.WCNT	000012		T.CNFG	000000		XLFCQE	002662R	002
RC.CD =	010000		T.CNF2	000002		XLFIN	002136R	002
RC.CTS=	020000		T.CSR	000016		XLFLQE	002660R	002
RC.DTR=	000002		T.FCNT	000005		XLHOOK	000516R	002
RC.IE =	000100		T.ICTR	000044		XLINT	001122RG	002
RC.RI =	040000		T.IGET	000046		XLLQE	000006RG	002
RC.RTS=	000004		T.IPUT	000042		XLPNAM	000514R	002
READ	000070R	002	T.IRNG	000040		XLSTRT	000000RG	002
RECS	001170R	002	T.ITOP	000050		XLSTS	007057	
RELEAS	003312RG	002	T.JOB	000024		XLSYS	000006RG	002
RESBRK	002322R	002	T.LPOS	000011		XL\$COD	000057	
RESSTT	002422R	002	T.NFIL	000026		XL\$CSR=	176500 G	
RONLY\$	040000		T.OCHR	000010		XL\$DVE=	000000	
RSTSIZ=	000060		T.OCTR	000262		XL\$END	002664RG	002
RTEM\$	000010		T.OGET	000264		XL\$LUN=	000001	
RTE\$M =	000000		T.OPUT	000260		XL\$MTY=	000001	
SAV30	002024R	002	T.OTOP	000266		XL\$NAM=	113740	
SETBRK	002350R	002	T.OWNR	000012		XL\$PC =	000000	
SETSTT	002454R	002	T.PRI	000022		XL\$PDP=	000001	
SHARE\$	002000		T.PTTI	000027		XL\$PDT=	000000	
SLEDI\$	000020		T.PUN	000025		XL\$PRI=	000004	
SLKMO\$	000002		T.STAT	000014		XL\$SBC=	000000	
SLOT\$	003226R	002	T.TCTF	000030		XL\$SPC=	000001	
SNDS	001140R	002	T.TFIL	000004		XL\$VEC=	000300 G	
SPECL\$	010000		T.TID	000032		XL\$VTB	000142RG	002
SPFUN	000162R	002	T.TNFL	000031		XOB	000140R	002
SPFUN\$	002000		T.TTLC	000036		XOCQE	000114R	002
SP.CCB=	000004		T.VEC	000020		XOLQE	000116R	002
SP.JOB=	000006		T.WID	000006		XOS	000136R	002
SRDDRV=	000203		UNITMK	000007		\$BLKEY	000256	
STASK\$	040000		UNLOAD	003256RG	002	\$CHKEY	000260	
STATFG=	000022R	002	USR\$	001000		\$CNFG1	000300	
STPSIZ=	000020		VARSZ\$	000400		\$CNFG2	000370	
STSDRV=	000204		VECSAV	000246		\$CNFG3	000466	
ST.CD =	000010		VIRTV\$	105372		\$CNTXT	000320	
ST.CTS=	000004		VS6\$0	001000		\$CSW	000004	
ST.RI =	000020		WONLY\$	020000		\$DATE	000262	
ST.XFH=	000001		WRITE	000104R	002	\$DECNT	000474	
ST.XOF=	000002		WRWT\$	000040		\$DFLG	000264	
SWREG\$	000004		XC.BRK=	000001		\$DWTYP	000440	
TCBADX	002654R	002	XC.IE =	000100		\$ELTIM	000422	
THK.BK	000006		XIB	000134R	002	\$EMTRT	000400	
THK.CT	000010		XIBFRE	002632R	002	\$ERRBY	000052	
THK.LE	000000		XIBIN	002626R	002	\$ERRCN	000356	
THK.NU	000001		XIBOUT	002630R	002	\$ERRLE	000376	
\$EXTIN	000416		\$QHOOK	000456		\$USRRB	000053	
\$E16LS	000316		\$RLPTR	002664RG	002	\$USRRP	000042	
\$FKPTR	002702RG	002	\$RMON	000000		\$USRTO	000050	
\$FORK	000402		\$RM2CO	000472		\$VIRT	000000	
\$GETVE	000436		\$RTSPC	000464		\$VIRTO	000002	
\$GTBYT	002670RG	002	\$SCROL	000302		\$WILDD	000454	
\$GTVEC	000354		\$SLOT2	000502		\$XTTPB	000500	
\$HSUFF	000412		\$SPSIZ	000504		\$XTTPS	000476	
\$H2CA	000462		\$SPSTA	000414		\$\$\$VER=	000022 G	
\$H2UB	000460		\$SPUSR	000272		.AUDIT	107123 G	
\$IFMXN	000377		\$STATW	000366		.XL	000044 G	
\$IMPLO	000446		\$SYCOM	000040		.XLGEN=	000020 G	
\$INCH	000007		\$SYIND	000364		...V1 =	000003	
\$INCL	000006		\$SYNCH	000324		...V10=	000100	
\$INDDV	000426		\$SYPTR	000054		...V11=	000200	
\$INDST	000417		\$SYSCH	000244		...V12=	000200	
\$INPTR	002700RG	002	\$SYSGE	000372		...V13=	000000	
\$JOBNU	000322		\$SYSUP	000277		...V14=	000000	
\$JOBS	000455		\$SYSVE	000276		...V15=	000176	
\$JSW	000044		\$SYUNI	000274		...V16=	000000	
\$JSX	000004		\$TCFIG	000424		...V17=	000000	
\$KMONI	000450		\$THKPT=	***** GX		...V18=	000000	
\$LOWMA	000326		\$TIMIT	002676RG	002	...V19=	000000	
\$MAXBL	000314		\$TRPLS	000434		...V2 =	000000	
\$MEMPT	000430		\$TRPSE	000442		...V20=	000000	
\$MEMSZ	000420		\$TTFIL	000056		...V21=	000000	
\$MFPS	000362		\$TTKB	000306		...V22=	000000	
\$MONAM	000406		\$TTKS	000304		...V27=	000000	
\$MPPTR	002666RG	002	\$TTNFI	000057		...V28=	001714	
\$MTPS	000360		\$TTPB	000312		...V3 =	000170	
\$NULJB	000444		\$TTPS	000310		...V4 =	000000	
\$PNPTR	000404		\$TT2RM	000470		...V5 =	000116	
\$PROGD	000452		\$UFLOA	000046		...V6 =	001714	

```
$PROGF 000453          $USRAR 000374          ...V9 = 000017
$PTBYT 002672RG      002 $USRLC 000266          ...V97= 000014
$PTWRD 002674RG      002 $USRLO 000352          ...V98= 000000
$QCOMP 000270          $USRPC 000040          ...V99= 177777

. ABS. 000622 000 (RW,I,GBL,ABS,OVR)
        000000 001 (RW,I,LCL,REL,CON)
XLDVR 003467 002 (RW,I,LCL,REL,CON)
Errors detected: 0
```