

XLTTM86

USER'S GUIDE

DIGITAL RESEARCH™

XLT86^{T.M.}

8080 to 8086 Assembly Language Translator
USER'S GUIDE

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FOREWORD

XLT86™ is a Digital Research software product that aids in the translation of 8080 assembly language programs to equivalent 8086 programs. XLT86 takes the CP/M® and MP/M™ environment into account, so that translated programs operate properly under both CP/M-86™ and MP/M-86™. XLT86 can also be used as a teaching tool by examining the output when XLT86 is applied to existing 8080 programs. Unlike other 8086 translators, XLT86 uses global data flow analysis techniques to determine 8080 register usage and reduce the number of generated 8086 instructions.

The XLT86 translator is available for operation under CP/M and MP/M for the 8080, 8085, and Z80® microprocessors with a minimum 40K Transient Program Area (TPA). XLT86 requires a 64K CP/M system to effectively translate any significant 8080 programs. Using a 4Mhz Z80 microprocessor, XLT86 translates programs at approximately 120 to 150 lines per minute, depending upon backup storage access speed. XLT86 is written in PL/I-80™ and thus can be adapted for use on computer systems that support Subset G. Specifically, XLT86 is available for cross-development on the Digital Equipment Corporation VAX 11/750 or 11/780 minicomputer, operating with the standard DEC VMS software. However, programs are supplied in machine code form, so it is not necessary to own PL/I-80 or any of its subsystems to operate XLT86.

The XLT86 system components, including the files XLT86.COM, XLT00.OVL, and XLT01.OVL, are distributed in IBM-compatible single density disk form. Before operating XLT86, copy these system components to a working disk and save the distribution disk for archive purposes. If the working disk medium can be dismounted, it must be marked with the notice shown below to properly comply with the Software License Agreement:

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This User's Guide presents the overall translation process, along with operator interface and command syntax. This manual also describes the format of the translated program, including the details of the 8080 to 8086 operation code translation.

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SECTION 1

THE TRANSLATION PROCESS

1.1 Input and Output Files

XLT86 reads an 8080 program from a file with type ASM and produces a file of type A86 containing the equivalent translated 8086 assembly language program. The filename for the 8080 source program, as well as filenames for all output files from XLT86, is taken from the command line typed by the operator. For example, the console command:

XLT86 DUMP

executes the XLT86 program using the file "DUMP.ASM" as input. The translation produces the output file "DUMP.A86".

The 8080 source program must be in a form acceptable to the standard Digital Research assembly language translators ASM, MAC, or RMAC. XLT86 processes conditional assembly statements, and produces an output program that results from evaluation of the particular conditions included in the 8080 program. However, macro definitions, macro invocations, and repeat loops are not altered in the translation. To properly translate programs that include macros or repeat loops, first assemble the programs under MAC or RMAC to produce a printer listing file of type PRN. Rename this PRN file to type ASM and edit the file to remove the beginning column positions, resulting in a file acceptable as input to XLT86. The A86 output file is now in a form acceptable to the Digital Research ASM86 assembler, requiring little or no modification for execution under CP/M-86 and MP/M-86.

XLT86 produces two additional files: a PRN file and a \$\$\$ file. A file of type PRN contains error lines and messages along with optional listing and trace information. The PRN file is in a form suitable for listing on the system printer and contains embedded form-feed and tab characters. A temporary file of type \$\$\$ is also created during translation. This temporary file is automatically deleted upon normal completion of XLT86.

The XLT86 program consists of a "root module" called XLT86.COM, which is loaded and executed when you enter the XLT86 command line shown above. There are two additional "overlays" called XLT00.OVL and XLT01.OVL that must be present on your default disk drive. These two overlays are automatically loaded and executed at the appropriate time during the translation.

1.2 Translation Phases

The translation itself takes place in five phases. Each phase has a specific name that appears at the console during translation so that the operator can monitor the progress of XLT86. Table 1-1 lists the phase names.

Table 1-1. XLT86 Translation Phases

Phase	Meaning
Symbol Setup	determines the location of each symbol in the 8080 source program.
Setup Blocks	determine the "Basic Blocks" necessary for the data flow analysis.
Join Blocks	construct a "Directed Graph" connecting each basic block, corresponding to program flow of control.
List Blocks	produce an optional list of Basic Blocks following flow analysis showing register and flag usage for each 8080 instruction.
Translate-86	translates the 8080 instructions to 8086 form, using the information gathered by the flow analysis.

The command line:

XLT86 DUMP

activates the XLT86 translator using the DUMP.ASM program as input. The default action of XLT86 prints the name of each phase at the console as the translation proceeds, as shown below.

```
Symbol Setup
Setup Blocks
Join Blocks
List Blocks
Translate-86
```

The files processed by the "XLT86 DUMP" command are shown in Figure 1-1, below.

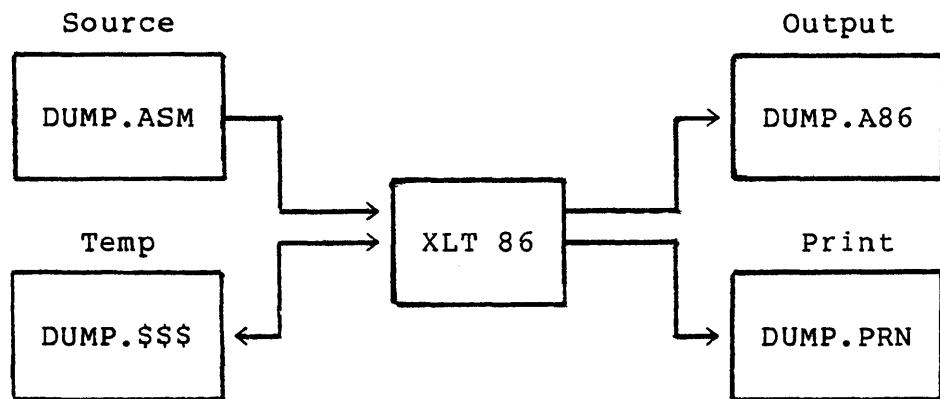


Figure 1-1. Processed Files

All files are placed on the drive specified by the operator as the prefix on the source filename. In the above example, all files are placed on the current default drive, which must also contain the XLT86 program along with its overlays. An alternative form:

XLT86 B:DUMP

overrides the default drive and obtains the source file from drive B. XLT86 creates the output, temporary, and print files on drive B as well. When several drives are available, it may be advantageous to place the various files on separate disks. In this case, you must use XLT86 "parameters," described in the following section, to override the default values.

SECTION 2

TRANSLATION PARAMETERS

2.1 Parameter Syntax

Several XLT86 parameters can be included in the command line by the operator or embedded within the 8080 source program to control the translation process. Parameters are grouped together into a parameter list enclosed within square brackets:

[p₁ p₂ ... p_n]

where p₁ through p_n denote one or more parameters optionally separated by blanks. When included on the command line, the XLT86 invocation appears as follows:

XLT86 filename [p₁ p₂ ... p_n]

When included within the source program, the opening bracket of the parameter list must begin in the first column position. The parameters denoted by p₁ through p_n are one or two character sequences in upper- or lower-case, with optional intervening blanks, as listed in Table 2-1, below.

Table 2-1. Translation Parameters

Parameter	Meaning
Ax	Place the A86 file on drive x where x = A, B, ..., P.
B	Produce a list of Basic Blocks in the PRN file.
C	Assume the 8086 "compact model" for execution.
J	Translate conditional jumps to short conditionals.
L	Send the PRN file directly to the system printer.
N	Show the line and statement number being processed.
P	Place the PRN file on drive x where x = A, B, ..., P.
R	Assume all flags active at subroutine returns.

Table 2-1. (continued)

Parameter	Meaning
S	Assume non-overlapping 8086 code and data segments.
T	Place the \$\$\$ file on drive x where x = A, B, ..., P.
80	Create an 8080 assembly listing in the PRN file.
86	Create an 8086 line and statement listing.

2.2 The B (Block Trace) Parameter

The A (A86), P (PRN), and T (TMP) parameters allow you to select alternate disk drives for use during the translation process when only limited disk space is available on each drive. Otherwise, disk drives are selected as described above.

The B (Block Trace) parameter provides a trace in the PRN file showing register usage information collected by the data flow analyzer. This parameter is not normally selected since the trace information is of no particular value unless you are interested in detailed register usage. The B parameter trace consists of a sequence of register usage tables for each Basic Block in the form shown below.

```
Block At 011E (subr), A86 = 083F
Entry Active: B-D-HL-AOZSPI Exit Active: BCDEHL-----
-----|stmt#| opcode uses | op | v1 | v2 | opcode kills | live regs |
-----| 23 | -----AOZSPI | PUSH | PSW| |----- |B-D-HL-A----|
-----| 24 | -----A----- | MOV | E | A |---E----- |B-DEHL-----|
-----| 25 | -----MVI | C | 05 |---C----- |BCDEHL-----|
-----| 26 | -----CALL|0005| |----- |BCDEHL-----|
-----
```

The Basic Block address in the original 8080 program is listed and the type of block is identified. The block type is "subr" for subroutines, "code" for main-line code, and "data" for data blocks. The A86 address is an approximation of the corresponding 8086 address used to determine short and long branch jump ranges. The remaining information shows register and flag use at block entry and at each instruction within the block. The registers and flags are displayed as a vector of letters and hyphens, where each letter represents the presence of a register or flag in the display, and each hyphen signifies that the corresponding register or flag is

absent in the vector. Given that all registers and flags are present, the display appears as follows:

BCDEHLMАЗOSPI

Table 2-2 lists the letter denotations of the above display.

Table 2-2. Letter Denotations for Registers and Flags

Letter	Meaning
B	Register B, or high(BC)
C	Register C, or low (BC)
D	Register D, or high(DE)
E	Register E, or low (DE)
H	Register H, or high(HL)
L	Register L, or low (HL)
M	Register M, memory operand
A	Register A, 8-bit Accumulator
O	Overflow Flag, carry or borrow
Z	Zero Flag
S	Sign Flag
P	Even Parity Flag
I	Interdigit Carry Flag

The registers active upon entry are listed first. In the example shown above, the data flow analysis has determined that the B, D, and HL registers, along with all flag registers, are in use upon entry to the block. The active registers following this block are then listed, consisting of the BC, DE, and HL register pairs. Then each instruction in the Basic Block is given, with a preceding statement number that can be cross-referenced with the 8080 source program. The instruction itself is listed with the hexadecimal values of its two optional parameters.

The "opcode uses" field shows the register set used by the operation code, while the "opcode kills" field lists the registers destroyed by the operation. The "live registers" field provides the information used by the Translate-86 phase to minimize the generated code. This field lists the registers and flags that are referenced following the instruction and is derived by examining the Directed

Graph corresponding to the 8080 source program. Again, the information collected by the flow analyzer is optionally displayed using the B parameter. This display is not required for normal operation of the translator.

2.3 The C (Compact) Parameter

The C (Compact) parameter causes XLT86 to generate 8086 machine code using the "Compact Memory Model" described in the CP/M-86 System Guide. Under normal circumstances, XLT86 assumes the "8080 Memory Model" where code and data segments overlap. To accomplish this overlap of segments, the program is analyzed to determine Basic Blocks that contain code and data. The program is assumed to begin with a code segment and, if a data segment is encountered as defined by a sequence of DS, DB, or DW statements, XLT86 produces the following statements that provide the proper transition:

```
L@n      EQU      $  
DSEG  
ORG      Offset L@n
```

Similarly, the transition from a data segment back to a code sequence is marked by the generated statements:

```
L@n      EQU      $  
CSEG  
ORG      Offset L@n
```

where L@n is a sequentially generated label. The labels are generated as required by XLT86, taking the form:

```
L@1      L@2      L@3      L@4      ....      L@32767
```

Enabling the "C" parameter prevents the code and data segments from overlapping. In this case, the transition from code to data and data to code is marked by either

DSEG

or

CSEG

respectively. See also the description of the S (Segments) parameter. When enabled, the S parameter completely overrides the C parameter.

2.4 The J (Jump) Parameter

The J (Jump) parameter enables the short jump analysis option of XLT86. When enabled, XLT86 translates 8080 conditional jumps to either short conditional jumps or negated short conditional jumps followed by short unconditional jumps, depending upon the byte count to the target of the jump. That is, a "JZ x" instruction becomes either

```
JZ x
...
x:
```

or

```
JNZ L@n
JMPS x
L@n:
```

The first case results if the label "x" is within the range of a short jump, while the second form results from a target label beyond the range of a short jump. The J parameter is enabled by default, and should be disabled using the NOJ form described below only if you want to manually edit your conditional jumps following program translation.

2.5 The L (List) Parameter

The L (List) parameter sends the listing file directly to the system printer, thus avoiding the intermediate PRN file. The system printer, or printer driver, must handle form-feeds (ctl-L) and tabs (ctl-I) to every eighth column position. If your printer does not properly support these characters, you can leave the L parameter disabled and use the CP/M PIP utility command form:

```
PIP LST:=filename.PRN[T8F]
```

where the PIP parameter "T8" expands tabs to blanks at every eighth column, and the "F" parameter deletes the form-feed character on transmission.

2.6 The N (Number) Parameter

The N (Number) parameter displays the current line and statement number on CRT-type console devices as the translation proceeds. Each line and statement number is displayed with an intervening carriage return, without a line feed, so that each successive display overwrites the previous value. In this way, you can easily monitor the progress of XLT86 as it proceeds through the source program during the translation.

2.7 The R (Return) Parameter

The R (Return) parameter overrides the default assumptions about register usage at the end of a subroutine. XLT86, by default, assumes that all registers are in use at the end of a subroutine in the absence of additional information. This is a safe, but possibly restrictive, assumption that might cause more 8086 code to be generated near the return statements of each subroutine. If you know that the entire 8080 program being translated contains subroutines that do not return flag registers, then you should include the R parameter in the command line to reduce the amount of generated code.

Alternatively, you can precede the return statements of various subroutines with "[R]" parameters when they do not return flag registers, as long as balancing "[NOR]" parameters, described below, are included to return to the default assumptions, where necessary.

2.8 The S (Segment) Parameter

The S (Segment) parameter informs XLT86 that the original source program contains embedded CSEG and DSEG directives that delimit the code and data segments. In this case, XLT86 makes no attempt to derive the code and data segment information and, instead, assumes that the CSEG and DSEG directives passed through to the 8086 program correctly define the appropriate segments. The S parameter is automatically set when the source program contains ASEG, CSEG, or DSEG directives, and completely overrides the effect of the C (Compact) parameter.

2.9 The 80 Parameter

The 80 parameter causes XLT86 to produce a pseudo-assembly listing of the original 8080 source program, giving the source line and statement number along with the assembled machine code location. If the B parameter is simultaneously enabled, additional Basic Block information precedes each straight-line code segment. When both 80 and B are enabled, the trace appears as shown below:

```
----- Basic Block (2) 011E
| Predecessors: 0119 0111 0105 0100
| Successors : 0125
| Reg's Killed: -C-E-----
| Reg's Used : -----AOZSPI
22    22 011E pr:
23    23 011E      push   psw
24    24 011F      mov    e,a
25    25 0120      mvi    c,1st
26    26 0122      call   bdos
```

Each Basic Block of the listing is preceded by the Basic Block Header consisting of the location (011E in the example above), a set of predecessor blocks where the program flow of control comes from (0119, 0111, 0105, and 0100), and a set of successor blocks where program flow could continue (0125, above). The set of registers killed are listed, along with the set of registers used by the operation codes within the block. No global data flow information is displayed in this trace (see the B parameter described earlier).

2.10 The NO Parameter

The two character sequence "NO" preceding the B, C, J, L, N, R, S, 80, and 86 parameters negates the effect of the parameter once it has been set. Further, the A, P, and T parameters are ignored when they occur within the source file and are effective only on the command line. The parameters B, C, J, L, N, R, S, 80, and 86 parameters, along with their negated forms, can occur in the command line or within the source program. When they occur within the program, they apply to the segment of code following their occurrence. Assuming that the default drive is d, where d is a valid drive code A, B, ..., P, the default values assumed for each parameter are identical to the complete, but redundant, command line shown below:

```
XLT86 d:filename [Ad NOB NOC J NOL NON Pd NOS NOR Td NO80 NO86]
```


SECTION 3

TRANSLATED PROGRAM FORMAT AND CONTENT

3.1 Translated Program Format

XLT86 constructs the 8086 program from the original 8080 program by first analyzing the program register usage. Then, using the collected information, XLT86 translates each label, operation code, and operand expression into an equivalent 8086 program segment. In performing the translation, XLT86 uses as many program fragments from the original 8080 source program as possible. These program fragments include labels, expressions, and comment fields. Due to differences in assembly language formats, however, labels and expressions might be altered somewhat to maintain their original meaning.

The translation occurs line-by-line, where each 8080 source line may contain several statements delimited by exclamation symbols. XLT86, however, always generates a single statement per output line. The output line includes an optional label in column one, followed by a single tab character. The translated operation code field is placed immediately following the tab character. If the operation code has one or two operand fields, another tab character is included and the operand fields are inserted. The operand fields themselves are constructed by either translating 8080 registers to their 8086 equivalents, or through the construction of an expression that is the translation of the original form. If a comment field is present in the source program, it is copied to the 8086 program intact with sufficient leading tabs to position the comment to column forty, if that position has not already been reached. Comments beginning in column one are reproduced without leading tab characters. Further, comments that begin in column one with the character "*" are started, instead, by the two character sequence ";"* to maintain compatibility with ASM-86.

For pseudo-assembly purposes, the assumed origin of the 8080 program is 0100H, corresponding to the base of the TPA under CP/M. This assumed origin resolves label addresses during pseudo-assembly and does not normally affect the translation process. However, if an ORG statement is encountered at the beginning of the program before any code or data is encountered, the program origin is set to the value given in the operand field of the ORG statement.

Program-relative operand references, along with absolute addresses, are allowed in the source program. In this case, XLT86 generates a label of the form "L@n" at the target location. For example below, the 8080 instruction sequence shown to the left results in the 8086 program shown to the right:

NOP	L@1: NOP
NOP	NOP
JMP \$-2	JMPS L@1

Similarly, the absolute 8080 assembly language shown to the left below results in the program shown to the right:

ORG 300H	ORG 300H
NOP	L@1: NOP
NOP	NOP
JMP 300H	JMP L@1

In this case, the ORG statement is necessary to override the default assumption.

From this last example, it appears that XLT86 is capable of translating 8080 programs produced through disassembly. Unfortunately, disassemblers cannot generally distinguish between code and data areas. If the code and data sections can be separated into distinct areas, where the code is disassembled with absolute address operands and the data areas consist of DS, DB, and DW operations, then XLT86 performs the translation.

Operand fields are translated according to their context and, for notational purposes, we make the following definitions.

Table 3-1. Operand Field Abbreviations

Abbreviation	Definition
ib	immediate byte operand (MVI A,ib)
iw	immediate word operand (LXI H,iw)
mb	byte in memory (STA x)
mw	word in memory (LHLD x)
mn	near memory (CALL x)
rb	byte in register (ADD B)
rw	word in register (DAD B)

The translation of an expression is denoted by a prime ('') following the expression type. Thus, ib is translated to ib', iw to iw', and so forth. Register translation takes place according to the following table.

Table 3-2. Register Translation

8080 Register (rb)	8086 Register (rb')
A	AL
B	CH
C	CL
D	DH
E	DL
H	BH
L	BL

The M (Memory) register has no direct equivalent in the 8086 environment, so XLT86 produces an "equate" statement in the following form at the beginning of each program.

```
M EQU Byte Ptr 0[BX]
```

Thus, the M register remains unchanged in the translation with the assumption that the BX register contains the offset to the proper memory location.

The 16-bit register pair translation occurs as shown in Table 3-3, below.

Table 3-3. 16-Bit Register Translation

8080 Register (rw)	8086 Register (rw')
PSW	AX
B	CX
D	DX
H	BX
SP	SP

The 8080 PSW and 8086 AX register have a loose correspondence depending upon register usage at the time of translation. The exact correspondence is defined below under the PUSH and POP operators.

3.2 Translated Program Content

Expressions are normally composed of literal constants, data variable references, program label references, and register references. XLT86 computes the type of each expression as the translation proceeds, resulting in one of the following expressions.

Table 3-4. Expressions

Expression	Meaning
constant	consists only of literal constants
variable	consists of zero or more constants and one or more variable references
label	consists of zero or more constants or variables, and one or more labels
register	consists of zero or more constants, variables, or labels and one or more register references

The translation of ib, iw, mb, mw, and mn is described in Table 3-5, below. This translation takes place after XLT86 scans the expression to determine its type, as described above.

Table 3-5. Operand Field Translation

Operand Field	Translation
ib and iw	<p>ib' and iw' are constructed from the original ib and iw by first determining the expression type. If the type is "constant," the expression ib or iw remains unchanged in the translation.</p> <p>Otherwise, for each variable, label, dollar sign (\$), or register reference in the expression, XLT86 changes the reference, denoted by x, to "(Offset x)" so that the resulting expression ib' or iw' represents a CS or DS relative offset computation.</p>
mb	<p>The resulting expression mb' is constructed from the original expression mb according to the type of mb. If mb is "constant" then mb' becomes "Byte Ptr .mb" denoting a single byte operand located at a literal constant address relative to DS or CS. Otherwise, the expression mb' becomes "Byte Ptr mb" denoting a byte variable or label address.</p>
mw	<p>Similar to mb, mw' becomes "Byte Ptr .mw" if mw is "constant" and "Byte Ptr mw" otherwise.</p>
mn	<p>The expression mn' is the same as the original mn unless there is no literal label at the target address. In this latter case, a label of the form "L&n" is created at the target address, which becomes the value of mn'.</p>

Due to differences in 8080 and 8086 program formulation requirements, not all valid 8080 expressions can be successfully converted to valid 8086 expressions. Thus, you must be aware that additional editing is required if your translated program produces errors during assembly with ASM-86. In particular, expressions that use arbitrary operations upon constants, variables, labels, and

registers are unlikely to assemble correctly under ASM-86, or any other assembler that uses the Intel conventions.

In the translation table given below, the 8080 operation code is shown to the left, with the translated 8086 code sequence shown to the right. In many cases, the registers that are live at the point where the 8080 operation code occurs determine the exact sequence of code that is generated. In these cases, the alternative forms are given separately. Conditional assembly notation specifies the alternative forms, with the introduction of the following two pseudo-functions:

`live(r1,r2, ..., rn)`

and

`short(mn')`

The "live" function takes a variable number of register arguments and results in a TRUE value if one or more of these registers is live at the point of translation. Otherwise, the "live" function results in a FALSE value. In the Section 2 example for the B parameter, statement 24 (MOV E,A) has the live register set given by the vector:

B-DEHL-----

so that

`live(B,C,D) = TRUE and live(A,O) = FALSE`

The "short" function is used in the translation of conditional jump instructions where the value of `short(mn')` is TRUE if the target of the translated jump address `mn'` is within the range of a conditional jump, or if the "J" parameter is enabled. Otherwise, `short(mn')` results in a FALSE value. XLT86 also uses the notation in Section 2 for label generation. The form "`L@n`" represents labels produced sequentially, starting at `n = 1`, used in the translation of conditional calls, returns, and conditional jumps outside the range of an 8086 conditional transfer. The CC (Call if Carry) operator, for example, translates to a jump conditional to a generated label followed by a direct call. The generated label is then inserted, as shown in the expansion of the 8080 instruction CC SUBR:

```
JNB      L@1
CALL    SUBR
L@1:
```

Table 3-6 gives the translation of each operation code. Note in particular that the following BDOS entry operations:

`CALL 0 CALL 5 JMP 0 JMP 5`

are treated as special cases that are translated to Interrupt 224, reserved by Intel Corporation for entry to CP/M-86 and MP/M-86.

Table 3-6. Translation Table

Operation Code	Translation
ACI ib	ADC AL,ib'
ADC rb	ADC AL,rb'
ADD rb	ADD AL,rb'
ADI ib	ADD AL,ib'
ANA rb	AND AL,rb'
ANI rb	AND AL,rb'
CALL 0	MOV CL,0 MOV DL,0 INT 224
CALL 5	INT 224
CALL mn	CALL mn'
CC mn	JNB L@n CALL mn' L@n:
CM mn	JNS L@n CALL mn' L@n:
CMA	NOT AL
CMC	CMC
CMP rb	CMP AL,rb'
CNC mn	JNAE L@n CALL mn' L@n:
CNZ mn	JZ L@n CALL mn' L@n:

Table 3-6. (continued)

Operation Code	Translation
CP mn	JS L@n CALL mn' L@n:
CPE mn	JNP L@n CALL mn' L@n:
CPI ib	CMP AL,ib'
CPO mn	JP L@n CALL mn' L@n:
CZ mn	JNZ L@n CALL mn' L@n:
DAA	DAA
DAD rw	IF rw = H SHL BX,1 ELSE IF live(O) AND NOT live(Z,S,P,I) ADD BX,rw' ELSE IF NOT live(O) AND live(Z,S,P,I) LAHF ADD BX,rw' SAHF ELSE LAHF ADD BX,rw' RCR SI,1 SAHF RCL SI,1 ENDIF ENDIF ENDIF
DEC rb	DEC rb'
DCX rw	DEC rw'
DI	CLI
EI	STI

Table 3-6. (continued)

Operation Code	Translation
HLT	HLT
IN ib	IN AL,ib'
INR rb	INC rb'
INX rw	IF NOT live(Z,S,P,I) INC rw' ELSE LAHF INC rw' SAHF ENDIF
JC mn	IF short(mn') JB mn' ELSE JNB L@n JMPS mn' L@n: ENDIF
JM mn	IF short(mn') JS mn' ELSE JNS L@n JMPS mn' L@n: ENDIF
JMP 0	MOV CL,0 MOV DL,0 INT 224 RET
JMP 5	INT 224 RET
JMP mn	JMPS mn'
JNC mn	IF short(mn') JNB mn' ELSE JNAE L@n JMPS mn' L@n: ENDIF

Table 3-6. (continued)

Operation Code	Translation
JNZ mn	IF short(mn') JNZ mn' ELSE JZ L@n JMPS mn' L@n: ENDIF
JP mn	IF short(mn') JNS mn' ELSE JS L@n JMPS mn' L@n: ENDIF
JPE mn	IF short(mn') JPE mn' ELSE JNP L@n JMPS mn' L@n: ENDIF
JPO mn	IF short(mn') JPO mn' ELSE JP L@n JMPS mn' L@n: ENDIF
JZ mn	IF short(mn') JZ mn' ELSE JNZ L@n JMPS mn' L@n: ENDIF
LDA mb	MOV AL,mb'
LDAX rw	MOV SI,rw' MOV AL,[SI]
LHLD mw	MOV BX,mw'
LXI rw,iw	MOV rw',iw'
MOV rbl,rb2	MOV rbl',rb2'

Table 3-6. (continued)

Operation Code	Translation
MVI rb,ib	MOV rb',ib'
NOP	NOP
ORA rb	OR AL,rb'
ORI ib	OR AL,rb'
OUT ib	OUT ib',AL
PCHL	JMP BX
POP rw	POP rw' IF rw = PSW AND live (O,Z,S,P,I) XCHG AL,AH SAHF ELSE IF rw = PSW AND live(A) XCHG AL,AH ENDIF ENDIF
PUSH rw	IF rw = PSW AND live(A) LAHF XCHG AL,AH PUSH AX XCHG AL,AH ELSE if rw = PSW LAHF XCHG AL,AH PUSH AX ELSE PUSH rw' ENDIF ENDIF
RAL	RCL AL,1
RAR	RCR AL,1
RC	JNB L@n RET L@n:
RET	RET
RLC	ROL AL,1

Table 3-6. (continued)

Operation Code	Translation
RM	JNS L@n RET L@n:
RNC	JNAE L@n RET L@n:
RNZ	JZ L@n RET L@n:
RP	JS L@n RET L@n:
RPE	JNP L@n RET L@n:
RPO	JP L@n RET L@n:
RRC	ROR AL,1
RST ib	INT ib'
RZ	JNZ L@n RET L@n:
SBB rb	SBB AL,rb'
SBI ib	SBB AL,ib'
SHLD mw	MOV mw',BX
SPHL	MOV SP,BX
STA mb	MOV mb',AL
STAX rw	MOV DI,rw' MOV [DI],AL
STC	STC
SUB rb	SUB AL,rb'
SUI ib	SUB AL,ib'

Table 3-6. (continued)

Operation Code	Translation
XCHG	XCHG BX,DX
XRA rb	XOR AL,rb'
XRI ib	XOR AL,ib'
XTHL	MOV BP,SP XCHG BX,[BP]

SECTION 4
XLT86 ERROR MESSAGES

4.1 Psuedo-assembly Process Error Messages

XLT86 issues error messages that fall into two categories: those produced by the pseudo-assembly process, and those produced during translation. Errors in the first category are not considered fatal, but are simply annotated in the source listing file following the line in which the error occurs. If errors are present, the message:

Number of Errors: n

is displayed at the console following the pseudo-assembly. Examine the PRN file to determine if the errors are significant. Error messages take the form:

** Error: e **, Near t

where e is one of the error codes, and t is a program element near the position where the error occurred. Table 4-1 lists the error codes.

Table 4-1. XLT86 Error Codes

Error Code	Meaning
Bad Flag	invalid parameter list [pl ... pn]
Balance	Unmatched right parenthesis or missing trailing string quote.
Boundary	Invalid program boundary, usually results from a branch to the middle of an instruction.
Convert	Cannot convert an operand to internal form.
End-Line	The end of a program line contains extraneous characters.
Exp Ovfl	Expression stack overflow; the expression is nested too deeply.
Gtr 7	An expression produced a value greater than 7, where a value from 0-7 is required.

Table 4-1. (continued)

Error Code	Meaning
Gtr 255	An expression produced a value greater than 255, where a value from 0-255 is required.
Mov M,M?	The source line contains the invalid instruction MOV M,M.
No Comma	Missing comma where comma is required.
No Value	A label or variable was encountered that does not have an assigned value.
Not Impl	The instruction or directive is not implemented in XLT86.
Phase	A label or variable has a different value on two passes through the source program.
Str Len	A string was encountered that exceeds the capacity of XLT86, check for missing right quote mark.
Value	The value produced by an expression is not compatible with the context in which it occurs.

4.2 Translate-86 Error Messages

The Translate-86 phase also produces a limited number of error messages. All errors produced by this phase are fatal, and cause immediate termination of XLT86. Table 4-2 lists these error messages.

Table 4-2. Translate-86 Error Messages

Error Message	Meaning
Bad Oper	Invalid 8080 operation code was encountered during translation; probably due to bad disk I/O operation. Check for hardware controller faults.
Not BDOS	A CALL or JMP occurred below the base origin of the program where the target is not 0000H (warm boot) or 0005H (BDOS entry).
Phase (B)	The Directed Graph does not correspond to the source program at the Basic Block level; usually due to a hardware malfunction.
Phase (S)	The Directed Graph does not correspond to the source program at the statement level; usually due to a hardware malfunction

An error produced by Translate-86 is accompanied by the console error message:

Fatal Error (See PRN file)

to indicate that such an error occurred.

4.3 Memory Overflow

The XLT86 program occupies approximately 30K bytes of main memory. The remainder of memory, up to the base of CP/M, stores the program graph that represents the 8086 program being translated. The error message:

ERROR (7) "Free Space Exhausted"

is issued if the program graph exceeds available memory. A 64K CP/M system allows translation of 8080 programs of up to approximately 6K.

The above error causes XLT86 to terminate. To continue, you must divide your source program into smaller modules and retry the translation.

APPENDIX A

SAMPLE PROGRAM TRANSLATIONS

The DUMP.ASM program presented here and normally included as a sample assembly language program with CP/M illustrates the translation process. The XLT86 command line:

```
XLT86 DUMP [8086]
```

produces the first example shown below. The "80" parameter selects the 8080 program listing option, while the "86" parameter selects the 8086 listing option. XLT86 places full lines of dashes ("----") between the Basic Blocks in the 8080 listing. This translation of the DUMP program, however, requires modification to run under CP/M-86. In particular, the DUMP.ASM program contains initialization code that saves the entry SP (statements 34 to 37) and resets the SP to a local stack (statement 39). The return statement following the FINIS label (statement 95) returns control to the CCP.

To perform an exactly equivalent sequence of operations, you must also save the stack segment register (SS) upon entry to the DUMP program, and restore this value before executing the return. Further, the simple RET operation must be replaced by a Far Return (RETF) to balance the original Far Call from the CCP. A simpler solution is to eliminate the initialization code (statements 33 through 39) and use the CCP's built-in 96 byte stack. Control returns to the CCP by executing a RETF at statement 95. If you want to use a local stack, set the SS register to the value of DS upon entry, and set SP to the Offset of STKTOP. Control returns to the CCP through execution of function call #0 in place of the RET in statement 95, as follows:

```
MOV CL,0  
MOV DL,0  
INT 224
```

The second listing shows the Basic Block information collected by the flow analyzer, and produced by the command line:

```
XLT86 DUMP [B]
```

where the "B" parameter selects the Basic Block trace. Under normal circumstances, either of the commands shown below are sufficient and reduce the amount of trace information:

```
XLT86 DUMP [N]
```

or

```
XLT86 DUMP
```

The first command is used only with a CRT-type device where the carriage-return character does not cause an automatic line-feed (see the description of the "N" parameter).

```

-----  

1      1 0100 ;      FILE DUMP PROGRAM, READS AN INPUT F  

2      2 0100 ;  

3      3 0100 ;      COPYRIGHT (C) 1975, 1976, 1977, 197  

4      4 0100 ;      DIGITAL RESEARCH  

5      5 0100 ;      BOX 579, PACIFIC GROVE  

6      6 0100 ;      CALIFORNIA, 93950  

7      7 0100 ;  

8      8 0100          ORG    100H  

9      9 0100 BDOS   EQU    0005H ;DOS ENTRY POINT  

10     10 0100 CONS   EQU    1      ;READ CONSOLE  

11     11 0100 TYPEF  EQU    2      ;TYPE FUNCTION  

12     12 0100 PRINTF EQU    9      ;BUFFER PRINT ENTRY  

13     13 0100 BRKF   EQU    11     ;BREAK KEY FUNCTION  

14     14 0100 OPENF  EQU    15     ;FILE OPEN  

15     15 0100 READF  EQU    20     ;READ FUNCTION  

16     16 0100 ;  

17     17 0100 FCB    EQU    5CH   ;FILE CONTROL BLOCK  

18     18 0100 BUFF   EQU    80H   ;INPUT DISK BUFFER  

19     19 0100 ;  

20     20 0100 ;      NON GRAPHIC CHARACTERS  

21     21 0100 CR     EQU    ODH   ;CARRIAGE RETURN  

22     22 0100 LF     EQU    OAH   ;LINE FEED  

23     23 0100 ;  

24     24 0100 ;      FILE CONTROL BLOCK DEFINITIONS  

25     25 0100 FCBDN  EQU    FCB+0 ;DISK NAME  

26     26 0100 FCBFN  EQU    FCB+1 ;FILE NAME  

27     27 0100 FCBFT  EQU    FCB+9 ;DISK FILE TYPE (3  

28     28 0100 FCBRL  EQU    FCB+12 ;FILE'S CURRENT REC  

29     29 0100 FCBRC  EQU    FCB+15 ;FILE'S RECORD COUN  

30     30 0100 FCBCR  EQU    FCB+32 ;CURRENT (NEXT) REC  

31     31 0100 FCBLN  EQU    FCB+33 ;FCB LENGTH  

32     32 0100 ;  

33     33 0100 ;      SET UP STACK  

34     34 0100 LXI    H,0  

35     35 0103 DAD    SP  

36     36 0104 ;      ENTRY STACK POINTER IN HL FROM THE  

37     37 0104 SHLD   OLDSP  

38     38 0107 ;      SET SP TO LOCAL STACK AREA (RESTORE  

39     39 0107 LXI    SP,STKTOP  

40     40 010A ;      READ AND PRINT SUCCESSIVE BUFFERS  

41     41 010A CALL   SETUP  ;SET UP INPUT FILE  

-----  

42     42 010D          CPI    255   ;255 IF FILE NOT PR  

43     43 010F          JNZ    OPENOK ;SKIP IF OPEN IS OK  

-----  

44     44 0112 ;  

45     45 0112 ;      FILE NOT THERE, GIVE ERROR MESSAGE  

46     46 0112 LXI    D,OPNMSG  

47     47 0115 CALL   ERR  

-----
```

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48	48 0118	JMP	FINIS	; TO RETURN
49	49 011B ;			
50	50 011B OPENOK: ;OPEN OPERATION OK, SET BUFFER INDE			
51	51 011B MVI A,80H			
52	52 011D STA IBP ;SET BUFFER POINTER			
53	53 0120 HL CONTAINS NEXT ADDRESS TO PRINT			
54	54 0120 LXI H,0 ;START WITH 0000			
55	55 0123 ;			
56	56 0123 GLOOP:			
57	57 0123 PUSH H ;SAVE LINE POSITION			
58	58 0124 CALL GNB			
59	59 0127 POP H ;RECALL LINE POSITI			
60	60 0128 JC FINIS ;CARRY SET BY GNB I			
61	61 012B MOV B,A			
62	62 012C ; PRINT HEX VALUES			
63	63 012C ; CHECK FOR LINE FOLD			
64	64 012C MOV A,L			
65	65 012D ANI OFH ;CHECK LOW 4 BITS			
66	66 012F JNZ NONUM			
67	67 0132 ; PRINT LINE NUMBER			
68	68 0132 CALL CRLF			
69	69 0135 ;			
70	70 0135 ; CHECK FOR BREAK KEY			
71	71 0135 CALL BREAK			
72	72 0138 ; ACCUM LSB = 1 IF CHARACTER READY			
73	73 0138 RRC ;INTO CARRY			
74	74 0139 JC FINIS ;DON'T PRINT ANY MO			
75	75 013C ;			
76	76 013C MOV A,H			
77	77 013D CALL PHEX			
78	78 0140 MOV A,L			
79	79 0141 CALL PHEX			
80	80 0144 NONUM:			
81	81 0144 INX H ;TO NEXT LINE NUMBE			
82	82 0145 MVI A,' '			
83	83 0147 CALL PCHAR			
84	84 014A MOV A,B			
85	85 014B CALL PHEX			
86	86 014E JMP GLOOP			
87	87 0151 ;			
88	88 0151 FINIS:			
89	89 0151 ; END OF DUMP, RETURN TO CCP			

```

90      90 0151 ;      (NOTE THAT A JMP TO 0000H REBOOTS)
91      91 0151     CALL    CRLF
-----
92      92 0154     LHLD    OLDSPL
93      93 0157     SPHL
94      94 0158 ;      STACK POINTER CONTAINS CCP'S STACK
95      95 0158     RET     ;TO THE CCP
-----
96      96 0159 ;      ;
97      97 0159 ;      ;
98      98 0159 ;      SUBROUTINES
99      99 0159 ;      ;
100     100 0159 BREAK: ;CHECK BREAK KEY (ACTUALLY ANY KEY
101     101 0159     PUSH H! PUSH D! PUSH B; ENVIRONMENT
102     104 015C     MVI     C,BRKF
103     105 015E     CALL    BDOS
-----
104     106 0161     POP B! POP D! POP H; ENVIRONMENT RE
105     109 0164     RET
-----
106     110 0165 ;      ;
107     111 0165 PCHAR: ;PRINT A CHARACTER
108     112 0165     PUSH H! PUSH D! PUSH B; SAVED
109     115 0168     MVI     C,TYPEF
110     116 016A     MOV     E,A
111     117 016B     CALL    BDOS
-----
112     118 016E     POP B! POP D! POP H; RESTORED
113     121 0171     RET
-----
114     122 0172 ;      ;
115     123 0172 CRLF:
116     124 0172     MVI     A,CR
117     125 0174     CALL    PCHAR
-----
118     126 0177     MVI     A,LF
119     127 0179     CALL    PCHAR
-----
120     128 017C     RET
-----
121     129 017D ;      ;
122     130 017D ;      ;
123     131 017D PNIB: ;PRINT NIBBLE IN REG A
124     132 017D     ANI     0FH     ;LOW 4 BITS
125     133 017F     CPI     10
126     134 0181     JNC     P10
-----
127     135 0184 ;      LESS THAN OR EQUAL TO 9
128     136 0184     ADI     '0'
129     137 0186     JMP     PRN
-----
130     138 0189 ;      ;
131     139 0189 ;      GREATER OR EQUAL TO 10
132     140 0189 P10:   ADI     'A' - 10
-----
```

133	141 018B	PRN:	CALL	PCHAR
134	142 018E		RET	
135	143 018F	;		
136	144 018F	PHEX:	;PRINT HEX CHAR IN REG A	
137	145 018F		PUSH PSW	
138	146 0190		RRC	
139	147 0191		RRC	
140	148 0192		RRC	
141	149 0193		RRC	
142	150 0194		CALL PNIB	;PRINT NIBBLE
143	151 0197		POP PSW	
144	152 0198		CALL PNIB	
145	153 019B		RET	
146	154 019C	;		
147	155 019C	ERR:	;PRINT ERROR MESSAGE	
148	156 019C	;	D,E ADDRESSES MESSAGE ENDING WITH "	
149	157 019C		MVI C,PRINTF	;PRINT BUFF
150	158 019E		CALL BDOS	
151	159 01A1		RET	
152	160 01A2	;		
153	161 01A2	;		
154	162 01A2	GNB:	;GET NEXT BYTE	
155	163 01A2		LDA IBP	
156	164 01A5		CPI 80H	
157	165 01A7		JNZ G0	
158	166 01AA	;	READ ANOTHER BUFFER	
159	167 01AA	;		
160	168 01AA	;		
161	169 01AA		CALL DISKR	
162	170 01AD		ORA A	;ZERO VALUE IF READ
163	171 01AE		JZ G0	;FOR ANOTHER BYTE
164	172 01B1	;	END OF DATA, RETURN WITH CARRY SET	
165	173 01B1		STC	
166	174 01B2		RET	
167	175 01B3	;		
168	176 01B3	GO:	;READ THE BYTE AT BUFF+REG A	
169	177 01B3		MOV E,A	;LS BYTE OF BUFFER
170	178 01B4		MVI D,0	;DOUBLE PRECISION I
171	179 01B6		INR A	;INDEX=INDEX+1
172	180 01B7		STA IBP	;BACK TO MEMORY
173	181 01BA	;	POINTER IS INCREMENTED	
174	182 01BA	;	SAVE THE CURRENT FILE ADDRESS	
175	183 01BA		LXI H,BUFF	
176	184 01BD		DAD D	

```

177      185 01BE    ;      ABSOLUTE CHARACTER ADDRESS IS IN HL
178      186 01BE    MOV     A,M
179      187 01BF    ;      BYTE IS IN THE ACCUMULATOR
180      188 01BF    ORA     A      ;RESET CARRY BIT
181      189 01C0    RET

-----
182      190 01C1    ;
183      191 01C1    SETUP: ;SET UP FILE
184      192 01C1    ;      OPEN THE FILE FOR INPUT
185      193 01C1    XRA     A      ;ZERO TO ACCUM
186      194 01C2    STA     FCBCR ;CLEAR CURRENT RECO
187      195 01C5    ;
188      196 01C5    LXI     D,FCB
189      197 01C8    MVI     C,OPENF
190      198 01CA    CALL    BDOS

-----
191      199 01CD    ;      255 IN ACCUM IF OPEN ERROR
192      200 01CD    RET

-----
193      201 01CE    ;
194      202 01CE    DISKR: ;READ DISK FILE RECORD
195      203 01CE    PUSH H! PUSH D! PUSH B
196      206 01D1    LXI     D,FCB
197      207 01D4    MVI     C,READF
198      208 01D6    CALL    BDOS

-----
199      209 01D9    POP B! POP D! POP H
200      212 01DC    RET

-----
201      213 01DD    ;
202      214 01DD    ;      FIXED MESSAGE AREA
203      215 01DD    SIGNON: DB      'FILE DUMP VERSION 1.4$'
204      216 01F3    OPNMSG: DB      CR,LF,'NO INPUT FILE PRESEN
205      217 0213    ;
206      218 0213    ;      VARIABLE AREA
207      219 0213    IBP:    DS      2      ;INPUT BUFFER POINT
208      220 0215    OLDSP:   DS      2      ;ENTRY SP VALUE FRO
209      221 0217    ;
210      222 0217    ;      STACK AREA
211      223 0217    DS      64      ;RESERVE 32 LEVEL S

-----
212      224 0257    STKTOP:
213      225 0257    ;
214      226 0257    END

```

```

1      0 M      EQU     Byte Ptr 0[BX]
1      1 ;      FILE DUMP PROGRAM, READS AN INPUT FILE AND
2      2 ;
3      3 ;      COPYRIGHT (C) 1975, 1976, 1977, 1978
4      4 ;      DIGITAL RESEARCH
5      5 ;      BOX 579, PACIFIC GROVE
6      6 ;      CALIFORNIA, 93950
7      7 ;
8      8 ORG     100H
9      9 BDOS   EQU     0005H ;DOS ENTRY
10    10 CONS   EQU     1 ;READ CONSO
11    11 TYPEF  EQU     2 ;TYPE FUNCT
12    12 PRINTF EQU     9 ;BUFFER PRI
13    13 BRKF   EQU     11 ;BREAK KEY
14    14 OPENF  EQU     15 ;FILE OPEN
15    15 READF  EQU     20 ;READ FUNCT
16    16 ;
17    17 FCB    EQU     5CH ;FILE CONTR
18    18 BUFF   EQU     80H ;INPUT DISK
19    19 ;
20    20 ;      NON GRAPHIC CHARACTERS
21    21 CR     EQU     0DH ;CARRIAGE R
22    22 LF     EQU     0AH ;LINE FEED
23    23 ;
24    24 ;      FILE CONTROL BLOCK DEFINITIONS
25    25 FCBDN  EQU     FCB+0 ;DISK NAME
26    26 FCBFN  EQU     FCB+1 ;FILE NAME
27    27 FCBFT  EQU     FCB+9 ;DISK FILE
28    28 FCBRL  EQU     FCB+12 ;FILE'S CUR
29    29 FCBRC  EQU     FCB+15 ;FILE'S REC
30    30 FCBCR  EQU     FCB+32 ;CURRENT (N
31    31 FCBLN  EQU     FCB+33 ;FCB LENGTH
32    32 ;
33    33 ;      SET UP STACK
34    34 MOV     BX,0
35    35 ADD     BX,SP
36    36 ;      ENTRY STACK POINTER IN HL FROM THE CCP
37    37 MOV     Word Ptr OLDSP,BX
38    38 ;      SET SP TO LOCAL STACK AREA (RESTORED AT FIN)
39    39 MOV     SP,(Offset STKTOP)
40    40 ;      READ AND PRINT SUCCESSIVE BUFFERS
41    41 CALL   SETUP ;SET UP INP
42    42 CMP     AL,255 ;255 IF FIL
43    43 JNZ     OPENOK ;SKIP IF OP
44    44 ;
45    45 ;      FILE NOT THERE, GIVE ERROR MESSAGE AND RETU
46    46 MOV     DX,(Offset OPNMSG)
47    47 CALL   ERR
48    48 JMPS   FINIS ;TO RETURN
49    49 ;
50    50 OPENOK:          ;OPEN OPERATION OK,
51    51 MOV     AL,80H
52    52 MOV     Byte Ptr IBP,AL ;SET BUFFER
53    53 ;      HL CONTAINS NEXT ADDRESS TO PRINT
54    54 MOV     BX,0 ;START WITH

```

```

55      55  ;
56      56  GLOOP:
57      57      PUSH    BX          ;SAVE LINE
58      58      CALL    GNB
59      59      POP     BX          ;RECALL LIN
60      60      JB     FINIS      ;CARRY SET
61      61      MOV     CH,AL
62      62  ;   PRINT HEX VALUES
63      63  ;   CHECK FOR LINE FOLD
64      64      MOV     AL,BL
65      65      AND    AL,OFH      ;CHECK LOW
66      66      JNZ    NONUM
67      67  ;   PRINT LINE NUMBER
68      68      CALL   CRLF
69      69  ;
70      70  ;   CHECK FOR BREAK KEY
71      71      CALL   BREAK
72      72  ;   ACCUM LSB = 1 IF CHARACTER READY
73      73      ROR    AL,1      ;INTO CARRY
74      74      JB     FINIS      ;DON'T PRIN
75      75  ;
76      76      MOV    AL,BH
77      77      CALL   PHEX
78      78      MOV    AL,BL
79      79      CALL   PHEX
80      80  NONUM:
81      81      LAHF
81      81      INC    BX          ;TO NEXT LI
81      81      SAHF
82      82      MOV    AL,' '
83      83      CALL   PCHAR
84      84      MOV    AL,CH
85      85      CALL   PHEX
86      86      JMPS  GLOOP
87      87  ;
88      88  FINIS:
89      89  ;   END OF DUMP, RETURN TO CCP
90      90  ;   (NOTE THAT A JMP TO 0000H REBOOTS)
91      91      CALL   CRLF
92      92      MOV    BX,Word Ptr OLDSP
93      93      MOV    SP,BX
94      94  ;   STACK POINTER CONTAINS CCP'S STACK LOCATION
95      95      RET              ;TO THE CCP
96      96  ;
97      97  ;
98      98  ;   SUBROUTINES
99      99  ;
100     100  BREAK:           ;CHECK BREA
101     101      PUSH   BX
101     102      PUSH   DX
101     103      PUSH   CX          ; ENVIRONME
102     104      MOV    CL,BRKF
103     105      INT    224
104     106      POP    CX
104     107      POP    DX

```

```

104    108      POP     BX          ; ENVIRONME
105    109      RET
106    110      ;          ; PRINT A CH
107    111      PCHAR:
108    112      PUSH    BX
108    113      PUSH    DX
108    114      PUSH    CX          ; SAVED
109    115      MOV     CL,TYPEF
110    116      MOV     DL,AL
111    117      INT    224
112    118      POP     CX
112    119      POP     DX
112    120      POP     BX          ; RESTORED
113    121      RET
114    122      ;
115    123      CRLF:
116    124      MOV     AL,CR
117    125      CALL   PCHAR
118    126      MOV     AL,LF
119    127      CALL   PCHAR
120    128      RET
121    129      ;
122    130      ;
123    131      PNIB:           ; PRINT NIBB
124    132      AND    AL,0FH        ; LOW 4 BITS
125    133      CMP    AL,10
126    134      JNB    P10
127    135      ;      LESS THAN OR EQUAL TO 9
128    136      ADD    AL,'0'
129    137      JMPS   PRN
130    138      ;
131    139      ;      GREATER OR EQUAL TO 10
132    140      P10:           ADD    AL,'A' - 10
133    141      PRN:           CALL   PCHAR
134    142      RET
135    143      ;
136    144      PHEX:           ; PRINT HEX
137    145      LAHF
137    145      XCHG   AL,AH
137    145      PUSH    AX
137    145      XCHG   AL,AH
138    146      ROR    AL,1
139    147      ROR    AL,1
140    148      ROR    AL,1
141    149      ROR    AL,1
142    150      CALL   PNIB        ; PRINT NIBB
143    151      POP    AX
143    151      XCHG   AL,AH
143    151      SAHF
144    152      CALL   PNIB
145    153      RET
146    154      ;
147    155      ERR:           ; PRINT ERRO
148    156      ;      D,E ADDRESSES MESSAGE ENDING WITH "$"
149    157      MOV    CL,PRINTF    ; PRINT BUFF

```

```

150    158      INT     224
151    159      RET
152    160      ;
153    161      ;
154    162      GNB:          ;GET NEXT B
155    163      MOV     AL,Byte Ptr IBP
156    164      CMP     AL,80H
157    165      JNZ     GO
158    166      ;      READ ANOTHER BUFFER
159    167      ;
160    168      ;
161    169      CALL    DISKR
162    170      OR      AL,AL          ;ZERO VALUE
163    171      JZ      GO          ;FOR ANOTHE
164    172      ;      END OF DATA, RETURN WITH CARRY SET FOR EOF
165    173      STC
166    174      RET
167    175      ;
168    176      GO:          ;READ THE B
169    177      MOV     DL,AL          ;LS BYTE OF
170    178      MOV     DH,0           ;DOUBLE PRE
171    179      INC     AL           ;INDEX=INDE
172    180      MOV     Byte Ptr IBP,AL ;BACK TO ME
173    181      ;      POINTER IS INCREMENTED
174    182      ;      SAVE THE CURRENT FILE ADDRESS
175    183      MOV     BX,BUFF
176    184      ADD     BX,DX
177    185      ;      ABSOLUTE CHARACTER ADDRESS IS IN HL
178    186      MOV     AL,M
179    187      ;      BYTE IS IN THE ACCUMULATOR
180    188      OR      AL,AL          ;RESET CARR
181    189      RET
182    190      ;
183    191      SETUP:        ;SET UP FIL
184    192      ;      OPEN THE FILE FOR INPUT
185    193      XOR     AL,AL          ;ZERO TO AC
186    194      MOV     Byte Ptr .FCBCR,AL ;CLEAR CURR
187    195      ;
188    196      MOV     DX,FCB
189    197      MOV     CL,OPENF
190    198      INT     224
191    199      ;      255 IN ACCUM IF OPEN ERROR
192    200      RET
193    201      ;
194    202      DISKR:        ;READ DISK
195    203      PUSH    BX
195    204      PUSH    DX
195    205      PUSH    CX
196    206      MOV     DX,FCB
197    207      MOV     CL,READF
198    208      INT     224
199    209      POP     CX
199    210      POP     DX
199    211      POP     BX
200    212      RET

```

```

200    212  L@1      EQU      $
200    212          DSEG
200    212          ORG      Offset L@1
201    213  ;
202    214  ;      FIXED MESSAGE AREA
203    215  SIGNON  DB      'FILE DUMP VERSION 1.4$'
204    216  OPNMSG   DB      CR,LF,'NO INPUT FILE PRESENT ON DIS
206    218  ;      VARIABLE AREA
207    219  IBP     RS      2           ;INPUT BUFF
208    220  OLDSP   RS      2           ;ENTRY SP V
209    221  ;
210    222  ;      STACK AREA
211    223          RS      64          ;RESERVE 32
212    224  STKTOP   RS      0
213    225  ;
214    226          END

```

L I S T O F B A S I C B L O C K S

Block At 0005 (subr), A86 = 0005

Entry Active: ----- Exit Active: -----

|stmt#| opcode uses | op | v1 | v2 | opcode kills | live regs |

Block At 0100 (code), A86 = 0100

Entry Active: BCDE---A---- Exit Active: BCDEHL-A----

|stmt#| opcode uses | op | v1 | v2 | opcode kills | live regs |

34	---	LXI	H 0000	-----	HL-----	BCDEHL-A-----
35	---HL	DAD	SP	-----	HL--O-----	BCDEHL-A-----
37	---HL	SHLD	0215	-----	-----	BCDEHL-A-----
39	-----	LXI	M 0257	-----	-----	BCDEHL-A-----
41	-----	CALL	01C1	-----	-----	BCDEHL-A-----

Block At 010D (code), A86 = 0115

Entry Active: BCDEHL-A---- Exit Active: BCDEHL-AOZSPI

|stmt#| opcode uses | op | v1 | v2 | opcode kills | live regs |

42	---A---	CPI	FF	-----	OZSPI	BCDEHL-AOZSPI
43	---Z---	JNZ	011B	-----	-----	BCDEHL-AOZSPI

Block At 0112 (code), A86 = 011C

Entry Active: BC---HL-AOZSPI Exit Active: BCDEHL-AOZSPI

|stmt#| opcode uses | op | v1 | v2 | opcode kills | live regs |

46	-----	LXI	D 01F3	--DE-----	-----	BCDEHL-AOZSPI
47	-----	CALL	019C	-----	-----	BCDEHL-AOZSPI

Block At 0118 (code), A86 = 0122
 Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
48	JMP	0151					BCDEHL-AOZSPI

Block At 011B (code), A86 = 0125
 Entry Active: BCDE---AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
51	MVI	A	A	80		A	BCDE---AOZSPI
52	STA	A	0213				BCDE---AOZSPI
54	LXI	H	0000			HL	BCDEHL-AOZSPI

Block At 0123 (code), A86 = 012D
 Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
57	PUSH	HL	H				BCDEHL-AOZSPI
58	CALL	01A2					BCDEHL-AOZSPI

Block At 0127 (code), A86 = 0136
 Entry Active: BCDE---AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
59	POP	HL	H			HL	BCDEHL-AOZSPI
60	JC	O	0151				BCDEHL-AOZSPI

Block At 012B (code), A86 = 0141
 Entry Active: -CDEHL-A---- Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
61	MOV	A	B	A	B		BCDEHL-----
64	MOV	L	A	L	A		BCDEHL-A----
65	ANI	A	OF			AOZSPI	BCDEHL-AOZSPI
66	JNZ	Z	0144				BCDEHL-AOZSPI

Block At 0132 (code), A86 = 014C
 Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
68	CALL	0172					BCDEHL-AOZSPI

Block At 0135 (code), A86 = 014F

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

71			CALL	0159			BCDEHL-AOZSPI
----	--	--	------	------	--	--	---------------

Block At 0138 (code), A86 = 0152

Entry Active: BCDEHL-A-ZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

73	A		RRC			AO	BCDEHL-AOZSPI
74	O		JC	0151			BCDEHL-AOZSPI

Block At 013C (code), A86 = 015A

Entry Active: BCDEHL--OZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

76	H		MOV	A	H	A	BCDEHL-AOZSPI
77			CALL	018F			BCDEHL-AOZSPI

Block At 0140 (code), A86 = 015F

Entry Active: BCDEHL--OZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

78	L		MOV	A	L	A	BCDEHL-AOZSPI
79			CALL	018F			BCDEHL-AOZSPI

Block At 0144 (code), A86 = 0164

Entry Active: BCDEHL--OZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

81	HL		INX	H		HL	BCDEHL--OZSPI
82			MVI	A	20	A	BCDEHL-AOZSPI
83			CALL	0165			BCDEHL-AOZSPI

Block At 014A (code), A86 = 016C

Entry Active: BCDEHL--OZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

84	B		MOV	A	B	A	BCDEHL-AOZSPI
85			CALL	018F			BCDEHL-AOZSPI

Block At 014E (code), A86 = 0171

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
86	JMP		0123				BCDEHL-AOZSPI

Block At 0151 (code), A86 = 0174

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
91	CALL		0172				BCDEHL-AOZSPI

Block At 0154 (subr), A86 = 0177

Entry Active: BCDE---AOZSPI Exit Active: -----

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
92	LHLD	0215				HL	BCDEHL-AOZSPI
93	SPHL						BCDEHL-AOZSPI
95	RET						

Block At 0159 (subr), A86 = 017E

Entry Active: BCDEHL-AOZSPI Exit Active: -----AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
101	PUSH	H					BCDE---AOZSPI
102	PUSH	D					BC---AOZSPI
103	PUSH	B					AOZSPI
104	MVI	C	OB	C			AOZSPI
105	CALL	0005					AOZSPI

Block At 0161 (subr), A86 = 0195

Entry Active: -----AOZSPI Exit Active: -----

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
106	POP	B				BC	BC---AOZSPI
107	POP	D				DE	BCDE---AOZSPI
108	POP	H				HL	BCDEHL-AOZSPI
109	RET						

Block At 0165 (subr), A86 = 01A8

Entry Active: BCDEHL-AOZSPI Exit Active: -----AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
-------	--------	------	----	----	----	--------------	-----------

112	-----HL-----	PUSH	H		-----	BCDE---AOZSPI	
113	--DE-----	PUSH	D			BC-----AOZSPI	
114	BC-----	PUSH	B			-----AOZSPI	
115	-----	MVI	C	02	-C-----	-----AOZSPI	
116	-----A-----	MOV	E	A	-----E-----	-----AOZSPI	
117	-----	CALL	0005			-----AOZSPI	

Block At 016E (subr), A86 = 01C1

Entry Active: -----AOZSPI Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs	
118	-----	POP	B		BC-----	BC-----AOZSPI	
119	-----	POP	D		-----DE-----	BCDE---AOZSPI	
120	-----	POP	H		-----HL-----	BCDEHL-AOZSPI	
121	BCDEHL-AOZSPI	RET					

Block At 0172 (subr), A86 = 01D4

Entry Active: BCDEHL--OZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode uses	op	v1	v2	opcode kills	live regs	
124	-----	MVI	A	0D	-----A-----	BCDEHL-AOZSPI	
125	-----	CALL	0165			BCDEHL-AOZSPI	

Block At 0177 (code), A86 = 01D9

Entry Active: BCDEHL--OZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode uses	op	v1	v2	opcode kills	live regs	
126	-----	MVI	A	0A	-----A-----	BCDEHL-AOZSPI	
127	-----	CALL	0165			BCDEHL-AOZSPI	

Block At 017C (subr), A86 = 01DE

Entry Active: BCDEHL-AOZSPI Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs	
128	BCDEHL-AOZSPI	RET					

Block At 017D (code), A86 = 01DF

Entry Active: BCDEHL-A----- Exit Active: BCDEHL-A-----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs	
132	-----A-----	ANI	0F		-----AOZSPI	BCDEHL-A-----	
133	-----A-----	CPI	0A		-----OZSPI	BCDEHL-AO-----	
134	-----O-----	JNC	0189			BCDEHL-A-----	

Block At 0184 (code), A86 = 01E8

Entry Active: BCDEHL-A----- Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
136	-----A-----	ADI	30			-----AOZSPI	BCDEHL-AOZSPI
137	-----A-----	JMP	018B			-----	BCDEHL-AOZSPI

Block At 0189 (code), A86 = 01ED

Entry Active: BCDEHL-A----- Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
140	-----A-----	ADI	37			-----AOZSPI	BCDEHL-AOZSPI

Block At 018B (code), A86 = 01EF

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
141	-----	CALL 0165				-----	BCDEHL-AOZSPI

Block At 018E (subr), A86 = 01F2

Entry Active: BCDEHL-AOZSPI Exit Active: -----

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
142	BCDEHL-AOZSPI	RET				-----	-----

Block At 018F (subr), A86 = 01F3

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-A-----

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
145	-----AOZSPI	PUSH	PSW			-----	BCDEHL-A-----
146	-----A-----	RRC				-----AO-----	BCDEHL-A-----
147	-----A-----	RRC				-----AO-----	BCDEHL-A-----
148	-----A-----	RRC				-----AO-----	BCDEHL-A-----
149	-----A-----	RRC				-----AO-----	BCDEHL-A-----
150	-----A-----	CALL 017D				-----	BCDEHL-A-----

Block At 0197 (code), A86 = 0208

Entry Active: BCDEHL----- Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
151	-----	POP	PSW			-----AOZSPI	BCDEHL-AOZSPI
152	-----	CALL	017D			-----	BCDEHL-AOZSPI

Block At 019B (subr), A86 = 0211

Entry Active: BCDEHL-AOZSPI Exit Active: -----

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
153	BCDEHL-AOZSPI		RET				

Block At 019C (subr), A86 = 0212

Entry Active: B-DEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
157	MVI	C	09	-C		BCDEHL-AOZSPI	
158	CALL	0005				BCDEHL-AOZSPI	

Block At 01A1 (subr), A86 = 0217

Entry Active: BCDEHL-AOZSPI Exit Active: -----

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
159	BCDEHL-AOZSPI		RET				

Block At 01A2 (subr), A86 = 0218

Entry Active: BCDEHL----- Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
163	LDA	0213				A	BCDEHL-A-----
164	CPI	80				OZSPI	BCDEHL-AOZSPI
165	JNZ	01B3				Z	BCDEHL-AOZSPI

Block At 01AA (code), A86 = 0222

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
169	CALL	01CE					BCDEHL-AOZSPI

Block At 01AD (code), A86 = 0225

Entry Active: BCDEHL-A----- Exit Active: BCDEHL-A-ZSPI

stmt#	opcode	uses	op	v1	v2	opcode kills	live regs
170	ORA	A				AZOZSPI	BCDEHL-A-ZSPI
171	JZ	01B3				Z	BCDEHL-A-ZSPI

Block At 01B1 (subr), A86 = 022C

Entry Active: BCDEHL-A-ZSPI Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
173		STC			0	BCDEHL-AOZSPI
174	BCDEHL-AOZSPI	RET				

Block At 01B3 (subr), A86 = 022E

Entry Active: BC-----A----- Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
177	A	MOV	E	A	E	BC-E---A---
178		MVI	D	00	D	BCDE---A---
179	A	INR	A		A-Z	BCDE---A---
180	A	STA	0213			BCDE-----
183		LXI	H	0080	HL	BCDEHL-----
184	DEHL	DAD	D		HL-O	BCDEHL-----
186		MOV	A	M	A	BCDEHL-A---
188	A	ORA	A		AOZSPI	BCDEHL-AOZSPI
189	BCDEHL-AOZSPI	RET				

Block At 01C1 (subr), A86 = 0247

Entry Active: B---HL-A----- Exit Active: BCDEHL-AOZSPI

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
193	A	XRA	A		AOZSPI	B---HL-AOZSPI
194	A	STA	007C			B---HL-AOZSPI
196		LXI	D	005C	DE	B-DEHL-AOZSPI
197		MVI	C	OF	C	BCDEHL-AOZSPI
198		CALL	0005			BCDEHL-AOZSPI

Block At 01CD (subr), A86 = 0254

Entry Active: BCDEHL-AOZSPI Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
200	BCDEHL-AOZSPI	RET				

Block At 01CE (subr), A86 = 0255

Entry Active: BCDEHL-AOZSPI Exit Active: -----AOZSPI

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
203	HL	PUSH	H			BCDE---AOZSPI
204	DE	PUSH	D			BC---AOZSPI
205	BC	PUSH	B			-----AOZSPI
206		LXI	D	005C	DE	-----AOZSPI
207		MVI	C	14	C	-----AOZSPI
208		CALL	0005			-----AOZSPI

Block At 01D9 (subr), A86 = 026F

Entry Active: -----AOZSPI Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
209	-----	POP	B		BC-----	BC-----AOZSPI
210	-----	POP	D		--DE-----	BCDE---AOZSPI
211	-----	POP	H		--HL-----	BCDEHL-AOZSPI
212	BCDEHL-AOZSPI	RET				

Block At 01DD (data), A86 = 0282

Entry Active: ----- Exit Active: -----

stmt#	opcode uses	op	v1	v2	opcode kills	live regs
215	-----	DB	0016		-----	-----
216	-----	DB	0020		-----	-----
219	-----	DS	0002		-----	-----
220	-----	DS	0002		-----	-----
223	-----	DS	0040		-----	-----

Block At 0257 (code), A86 = 02FC

Entry Active: BCDEHL-AOZSPI Exit Active: BCDEHL-AOZSPI

stmt#	opcode uses	op	v1	v2	opcode kills	live regs

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