DDT REFERENCE MANUAL

(for the Tymshare debugging system)

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1.0 General

DDT is the debugging system for the SDS 930 Time-Sharing System. It has facilities for symbolic reference to and typeout of memory locations and central registers. Furthermore, it permits the use of literals in the same manner as in the assembler. It can also insert breakpoints into programs, perform a trace, and search programs for specified words and specified effective addresses. There is a command to facilitate program patching. Finally, DDT can load both absolute and relocatable files in the format produced by the assembler.

The system has a language for communication between DDT and its users. The basic components of this language are symbols, constants, and commands.

1.1 Symbols

A <u>symbol</u> is any string of letters, digits, and dots (.) containing at least one letter. (However, a digit string followed by B or D is interpreted as an octal or decimal number respectively). In symbols of more than six characters, only the first six are significant: thus, ALPHABET is equivalent to ALPHAB. All opcodes recognized by the assembler are built-in symbols, except for some I/O instructions. Other symbols are ;1, ;2, ;A, ;B, ;F, ;L, :M, ;Q, ;X, and dot. Their meanings are explained below.

Every symbol may have a value. This value is a 24-bit integer; for most symbols it will be either an address in memory or the octal encoding of an operation code. Examples:

ABC AB124 12XYZ

The following are not symbols:

135B AB*CD

Symbols may be introduced to DDT in two basically different ways:(A) They may be written out by the assembler and read in from the binary program file by DDT.

(B) They may be typed in and assigned values during debugging.

It is possible for a symbol to be undefined. This may occur if a program is loaded which references an external symbol not defined in a previously - 10 C - 100 BB (1) loaded program. It may also occur if an undefined symbol is typed in an the state of the s expression. In general, undefined symbols are legal input to DDT except •n•13 1979 * 148 when their values would be required immediately for the execution of a 的复数 化石 化合物 网络白色教教 建氯化合物 command. Thus, for example, the ;G (GO TO) command could not have an un-AND THE MOTOR THE AND ADDRESS OF THE defined symbol as its argument.

And a state of the second state of the Undefined symbols may become defined in several ways. They may be C - HOTEL MALENAMED MALE HAVE defined as external in the assembler (i.e. with EXT, ENTRY, or \$) and read STORES ALCONTON by DDT as part of a binary program. Alternatively, they may be defined by one of the symbol definition commands available in DDT. When the definition enne catingte presidui to provide occurs, the value of the symbol will be substituted in all the expressions e e e construction de la in which the symbol has appeared.

If DDT type [U] after typing out the contents of a register, it means 1.000 5 5 that the register contains an undefined symbol. The register is closed at , we forgete the third of the states of the set once so that its contents cannot be erroneously changed. ma readence e deci and encadar

. . .

The only restriction on this facility is that, as for ARPAS, the AGLETERS SU COTEMAN TIDE undefined symbol must be the only thing in the address field of the word HELET SIME DELET in which it appears. Incorrect uses of undefined symbols will be detected and the state of the state by DDT and will result in the error comment (U).

DDT keeps track of references to undefined symbols by building a pointer chain through the address fields of the words referring to the symbol. Thus, suppose that the symbol A is undefined and appears as follows :electory:

> **S1** LDA Α **S**2 STA Α \$3 MRG Α

· 1995年1996年1998年19月1日,19月1日日,19月1日日19月1日日,19月1日日日 and nowhere else in the program. After loading, the entry for A in DDT's n na an Traistean

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symbol table will contain a flag indicating that it is undefined and a pointer to 3. The above locations will contain:

S1	LDA	0
S 2	; Sta	S 1
\$ 3	: MRG	S 2

When the symbol is defined, DDT goes through the pointer chain and fills in the value. It recognizes the end of the pointer chain by a O address.

From this description it should be obvious what will happen if the pointer chain is destroyed. A probable consequence is that a search down the pointer chain will not terminate. DDT does such searches whenever it prints an address. If the chain it is searching has more than 256 links, it will print the symbol followed by (U) and continue. Fixing up an undefined symbol pointer chain which has been clobbered is an exercise which we leave to the reader.

1.2 Block Structure

A limited facility called the <u>block structure</u> facility is provided to simplify the referencing of local symbols which are defined in more than one program. Note that DDT's block structure has only a tenuous connection with the block structure of ALGOL. The block structure of a program is organized in the following manner: every IDENT read by DDT as part of a binary program file begins a new block. Any local symbol known to DDT has a block number associated with it; global symbols do not have a block number. Undefined symbols are always treated as global.

The name of a block is the symbol in the label field of the IDENT. If two IDENTs with the same symbol are read, the message (ALREADY DEFINED) is printed, and the local symbol tables from the two blocks will be merged.

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Global symbols must be unique within an entire program and are recognized nade weels and the second the second at all times. If a multiple definition is encountered, the latest one takes interport file and precedence. Local symbols are recognized according to the following rules:

- (1) At any given time one block is called the primary block. All local symbols associated with the primary block will be recognized.
- (2) If a symbol is used which is neither global nor in the primary block, the entire symbol table is scanned for it. If it occurs in only one block, the symbol is recognized properly. If it occurs in more than one block, the error message (A) is printed.
- (3) A symbol may be explicitly qualified by writing:

SYMA&SYMB Second Same SYMA must be the name of a block. SYMB is then referenced as though the block whose name is SYMA were primary.

(4) When a register is opened (see section 2.1), the block to which the symbolic part of its location belongs becomes primary. Thus, NN&XYZ/ causes block NN to become primary; if ABC is a unique local symbol in block PQ, then ABC/ causes block PQ to become primary. 122 2411

1.3 Literals

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Literals have the same format and meaning in DDT as in the assembler, i.e. the two characters' =' signal the beginning of a literal, which is terminated by any of the characters which ordinarily terminate an expression. In contrast to the assembler, the expression in a DDT literal must be defined.

The literal is looked up in the literal table. If it is found, the address which has been assigned to it is the value of the symbol. If it does not appear in the literal table, it is stored at the address which is the current value of ;F, and this address is taken as the value of the literal. ;F is increased by 1. For example, if the literal -1 does not already exist in the literal table and ;F is 1000B, then LDA =-1 causes -1 to be stored at 1000B, and is equivalent to LDA 1000B; the new value of ;F

is 1001B. Exception: In patch mode, literals are saved and not stored until the patch is completed since otherwise they would interfere with the patch.

When DDT types out a symbol whose value is an address in the literal table, it will type out in the same format in which it would be input; that is, as = followed by the numeric value of the literal.

1.4 Constants

significance:

A <u>constant</u> is any string of digits, possibly followed by a B or D. The number represented by the string is evaluated, truncated to 24 bits and then used just like the value of a symbol. The radix for numbers is normally 8 (octal), but may be changed arbitrarily by the commands described in section 2.4 below. If a number is terminated by B or D, it is interpreted as octal or decimal respectively regardless of the current radix. Constants are always printed by DDT in the current radix.

It is possible to enter numeric op codes by typing the number followed by an & sign. Thus 100 = 14400000B if the current radix is decimal (100D=144B). 1.5 Commands

A <u>command</u> is an order typed to DDT which instructs it to do something. The commands are listed and their functions explained in the table below. 1.6 Expressions

An <u>expression</u> is a string of numbers or symbols connected with blanks, +, -, ;*, ;/, ;&, ;<, ;=, ;>, and ;%. These operators have the following

> + addition subtraction (integer) multiplication ;* ;/ (integer) division ;& (AND) ;< (LSS) (EQL) as in ARPAS ;= GTR) ;> (OR)

Expressions are evaluated strictly left to right: all operators have the same precedence. Parentheses are not allowed. The first symbol or number

may be preceded by a minus sign. Blank acts like plus, except that the following operand is truncated to 1^4 bits before being added to the accumulated value of the expression. The value of an expression is a 2^4 -bit integer. An expression may be a single symbol or constant.

Examples:	LDA	has the value	7600000
	LDA 10	has the value	7600010 if the
			radix is octal
	LDA 10D	has the value	7600012
If SYM is a sy	mbol with the	value 1212, then	
	SYM	has the value	1212
	SYM 10	has the value	1555
	LDA SYM	has the value	07601212
		· · · · · · · · · · · · · · · · · · ·	

If this last expression were put into a memory register and later executed by the program the effect would be to load the contents of SYM, register 1212, into the A register.

When DDT types out expressions, two mode switches control the format of the output. Commands for setting these modes are described in section 2.4 below. The C-S mode determines whether quantities will be printed as constants or as symbolic expressions. In the latter case, the opcode (if any) and the address will be put into symbolic form. If the first nine bits of the value are 0 or 1, no opcode will be printed; in the latter case a negative integer will be printed. If the opcode is not recognizable as a symbol, it will be typed as a number followed by an (e) sign.

The R-V mode controls the format in which addresses are typed. DDT types addresses when asked to open the previous or the next register, when it reports the results of word and address searches, and on breakpoints. In relative mode, addresses are typed in symbolic form, i.e., as the largest defined symbol smaller than the address plus a constant if necessary. If the constant is bigger than 200 octal, or if the value of the symbol is less than the first location of the program, the entire address is typed as a constant. In absolute mode, addresses are always typed as constant.

1.7 The Open Register

One other major ingredient of the DDT language is the <u>open register</u>. Certain commands cause a register to be "opened". This means that its contents are typed out (except in enter mode, for which see the $\$ command), followed by a tab. Any expression the user types will then be inserted into the open register in place of its current contents. After this insertion the register is closed at once. Note that the string LDA ABC= is a command, and does not cause LDA ABC to be entered into the current open register. The <u>current location</u> is given by the symbol "." (dot) which always has as its value the address of the last register opened, whether or not it is still open.

Note:

- (1) Comma and star (for indirect addressing) may be used in expressions as they are used in the assembler; e.g. LDA* 0,2 has the value 27640000.
- (2) DDT will respond to any illegal input with the character ? followed by a tab (if a register is open) or carriage return (otherwise), after which it will behave as if nothing had been typed since the last tab or carriage return. The command ? also erases everything typed since the last tab or carriage return.

1.8 Memory Allocation and DDT

DDT may cause the time-sharing system to assign memory for use either by DDT itself or by the user's program. DDT's memory is used to hold the symbol table, which starts in block O and grows upward in memory. The symbol table contracts at the end of each load of a binary file and when symbols are killed; this contraction may cause memory to be released.

DDT grabs program memory when it is required for loading a binary file, or when a ;U (execute) command is given and the value of ;F is such that

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a new block is needed to hold the instruction to be executed. For executing an instruction, DDT requires location ;**F**, ;**F**+1 and ;**F**+2. Memory is never grabbed for examination of a register; however, entering information with $\$ can cause memory to be assigned. Attempts to open locations not assigned will cause DDT to type ?. This means that upon initial entry to DDT no registers are available for examination. The easiest way to obtain memory is to simply start typing in a program using the $\$ command.

If an attempt to acquire or reference memory leads to a trap, DDT types (M) and abandons whatever it is doing. This can happen if the machine size is exceeded, or if an attempt is made to change read-only memory.

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2.0 DDT Commands

In the following descriptions of DDT commands, <S> will be used to denote an arbitrary symbol. <E> or <W> will be used to denote an arbitrary expression which may be typed by the user: <E> will be used when the value of this expression is truncated to 14 bits before it is used by DDT, while <W> will denote a full 24-bit expression. <A> will be used to denote an <u>optional</u> 14-bit expression. If none is typed, the last expression printed out will usually be used; deviations from this rule will be described under the individual commands. <T> will denote a file name followed by a dot: DDT will type a tab whenever it expects a file name.

2.1 Register Opening Commands

<A> / This opens the register addressed by the value of <A>. DDT will give a tab, type an expression whose value is equal to the contents of the register, give another tab and await further commands. The precise form of the expression typed is dependent on the setting of the S-C and R-V modes. If the user types in an expression, DDT will insert its value into the register. Typing another command closes the register, unless it is a type value or symbol definition command. Note that in a command that requires a preceding expression, the expression is regarded as part of the command and would not, for instance, be inserted into the open register. If another / is given as the next command with no preceding expression the contents of the register addressed by the expression typed by DDT are typed out. A further / repeats this process. Note, however, that the original register opened remains the open register; any changes made will go into that register.

carriage This command does not necessarily have any effect. If the specified return conditions are present, however, any of the following actions may occur:

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- (1) If there is an open register, the register is closed.
- (?) If DDT is in enter mode, it leaves it.
- (3) If DDT is in patch mode, the patch is terminated (for a fuller description of this effect, see the patch command).
- <A>] This command has the same effect as /, except that the contents of the register opened are always typed in symbolic form.
- <A> [This command has the same effect as /, except that the contents of the register opened are typed in constant form.
- <A> \$ This command has the same effect as /, except that the contents of the register opened are typed as a signed integer.
- <E> " This command acts like /, except that the register constants are typed in ASCII. Unprintable characters, as in QED, are preceded by &, e.g. 141 (control-A) prints out as &A.
- line This command opens the register whose address is the current location
 feed
 plus one, iJe. the register after the one just opened. The output of DDT on
 this command is carriage return, register address (format controlled by
 the R-V mode), /, tab, value of contents, tab.
- ; $\omega(\omega=\text{space})$ This is equivalent to line feed except that nothing is printed. Its main use is in entering programs or data, e.g.

1000 1;.2;.3 (carriage return)

is equivalent to

1

1000\1(carriage return)1001\2(carriage return)1002\3(carriage return)

This command opens the register whose address is the current location minus one, i.e. the previous register. The output is the same as for the line feed command.

Example:				
ABC/	LDA	ALPHA		(line feed)
ABC+1/	STA	BETA	STA GAMMA	(line feed)
ABC+2/	LDB	DELTA	↑	
ABC+1/	STA	GAMMA		

This command opens the register whose address is the last 14 bits of the value of the last expression typed. The output is the same as for line feed.

This command is the same as /, except that the contents of the register are not typed. DDT goes into enter mode, in which the contents of registers opened by line feed, \dagger , or (are not typed. Any other command caused DDT to go out of enter mode. In particular, carriage return has this effect. When a register has been opened with \backslash , DDT thinks that it has typed out the contents. The type value commands will, therefore, work on the contents of the register.

The type register in special mode characters [,], \$ (type as a negative integer), " (type in ASCII) are also preserved by line feed, up arrow and (.

;\ This command suppresses typeout of register addresses during line feed, up arrow and (chains. Carriage return cancels the command.

2.2 Type Value Commands

This command types the value of the last expression typed (;Q) in constant form. It may appear in the form <W> =, in which case the value of the <W> is typed. Otherwise, the expression referred to is the one most recently typed, either by DDT or by the user.

This command types the value of ;Q as a signed integer.

This command types the value of ;Q in symbolic form.

This command types the value of ;Q typed as a word of text (see " command on previous page).

(à

#

1

This command types the address part of ;Q in symbolic form. If, for instance, the program has executed BRM X, then $X\setminus @$ will cause DDT to print the address of the BRM.

(

/

Example:

LDA=	76 00000		
LDA 10=	7600010		
LDA ←	LDA		
76 00000 ←	LDA		
-1=	77777777		
-1#	-1		
	-1		
10221043'	ABC		

2.3 Symbol Definition Commands

<\$>:

...

This command defines the value of the symbol \leq to be the current location. If \leq has been used but is undefined, it becomes global; otherwise it becomes local and associated with the block which is primary when the : command is given.

<S @ This command defines the value of <S> to be the address of the last expression typed by DDT or the user. The symbol is local and associated with the block which is primary when the @ command is given.

 $\langle < S \rangle$ This defines < S to have the value of < S, and to be global.

2.4 Mode Changing Commands

This command is followed by a string of arbitrary characters terminated by D^{C} (control D). If a register is open, the string will be inserted into successive locations packed 3 characters per word; otherwise characters beyond the third will be thrown away. For example, if no register is open, "ABCDED^C = yields 10221043.

;D (DECIMAL) This command changes the current radix (see section 1.4).
;O (OCTAL) This changes the current radix to octal.

(RADIX) sets the current radix to the value of the expression, which must be >2.

;[(CONSTANT) This command changes the S-C mode to constant, i.e. makes / equivalent to [.

(SYMBOLIC) This command changes the S-C mode to symbolic, i.e. makes / equivalent to].

;" (ASCII) This makes / equivalent to ".

;\$ (SIGNED INTEGER) This makes / equivalent to \$.

;R (RELATIVE) This command changes the R-V mode to relative. This mode determines the format for the output of addresses, both in symbolic expression and when generated by line feed and t.

;V

(ABSOLUTE) This command changes the R-V mode to absolute.

2.5 Breakpoint Commands

- <A>,<E>! (BREAKPOINT) <E>! sets breakpoint 0 at the address given by the value of the expression; <N>, <E>! sets breakpoint N (N must be between 0 and 3 inclusive). The effect is that if the program executes the instruction at this address control returns to DDT, which will print the address and the contents of the A, B and X registers and await further commands (see below). The break occurs before execution of the instruction in the breakpoint location. ;L is set to the location at which the break occurred.
- (CLEAR ALL BREAKPOINTS). ! alone causes all breakpoints to be cleared.
 <A>;! (LIST OR CLEAR BREAKPOINTS)
- <N>;! causes breakpoint N to be removed, where N lies between O and 3
 inclusive. ;! alone causes all breakpoints to be listed: if breakpoint 1
 is set at ABC+3, and no other breakpoints are set, then ;! produces the
 printout * ABC+3 * * .
- <A>;P (PROCEED) This command restarts the program after a break. The program executes the instruction at the break and goes on from there. No breakpoint is removed unless this is specifically done by ! or ;! so that, if the program arrives at this location again, another break will occur. If <E>;P is given, another break will not occur until some breakpoint has been reached that many times.

<A>;N (NEXT) This command executes the instruction at ;L and breaks. This program provides a trace facility in that repeated executions of ;N will provide a running print out of the contents of the significant internal registers, instruction by instruction. The function is essentially the same as that of the step switch on the console. <E>;N will cause <E> instructions to be executed before the next break occurs.

The ;N command follows the flow of control in the user's program. In particular, it will normally trace the execution of users' POPs (see ;O below). The execution of SYSPOPs, however, is not traced. In other words, a SYSPOP such as FAD (floating add) is regarded as one instruction by ;N. Cells ;F, ;F+1, and ;F+2 are used by ;N and ;P.

- <E>;S (STEP). This is equivalent to <E> repetitions of ;N. Note that this is not the same as <E>;N.
- <E>;V (ADVANCE). This is equivalent to <E> repetitions of <P, and is not the same as <E>;P.
- (N>;0 (POP TRACE MODE). If (N>>0, programmed operators (POPs) together with their associated subroutines will be treated like machine instructions for the ;N and ;S commands, i.e. the break will not occur until control. returns to the location following the POP. Since DDT determines when it should break by counting POPs, BRMs, SBRMs, BRRs and SBRRs, it can be fooled by POPs which do sufficiently peculiar things. If (N><0, POP subroutines will be traced, i.e. the first break after the POP will be at the first instruction of the subroutine.
- N>;U (SUBROUTINE TRACE MODE). If <N>=1, BRMs or SBRMs together with the subroutine called will be treated as single instructions by ;N. The same algorithm is used as in ;O to determine when to break. If <N>=O, subroutines will be traced explicitly.

Attempts to proceed through certain instructions having to do with forks will produce erroneous results, and breakpoints encountered when the program is running in a fork will not do the right thing. Attempts to proceed through unreasonable instructions will cause the error comment

| \$ > >.

2.6 Input/Output Commands

- <A>;Y<P>DDT expects to find a binary program on the file <P>. If the program is absolute it is read in. If it is relocatable it is read in and relocated at the location specified by <A>. If the expression is omitted, relocatable loading commences at location 240B and continues by beginning each program in the first available location after the preceding one. After reading is complete, the first location not used by the program is typed out. Any local symbols on the binary file are ignored.
- <A>;T<F> This command is identical to ;Y except that is also reads local
 symbols from the file and adds them to DDT's symbol table. Any symbols on
 the file will be recognized by DDT thereafter.

The following two points should be noted in connection with ;Y and ;T commands.

- The use of an expression before ;T or ;Y when the file is absolute (i.e. SAVE file or self-loading paper tape) is in error.
- 2) The block read in becomes the primary block.
- ;W<F> Causes all global symbols to be written on the specified file, in a format which can be read back in with ;T.

:C<F> Causes all symbols to be written on the specified file.

2.7 Search Commands

<W>;W (WORD SEARCH) <W>;W searches memory between the limits ;1 and ;2 for cells whose contents match <W> when both are masked by the value of ;M. The locations and contents of all such cells are typed out.

(NOT-WORD SEARCH). This is the same as ;W, except that all registers which do not match <W> will be printed. This is useful, for example, in finding and printing all non-zero registers in a given part of memory.
(EFFECTIVE ADDRESS SEARCH). <E>;E searches memory between the limits ;1 and ;? for effective addresses equal to <E>. Indexing, if specified, is done with the value of ;X. Indirect address chains are followed to a depth of 64. The addresses and contents of all words found are typed out. When ;W or ;E is complete, . is left pointing to the last register where the expression was found.

2.8 The Patch Command

- (A>) <A>) causes a patch to be inserted. If a register is open and an expression is given, the expression is entered into the register. If a register is open, or if no expression is typed, the patch is made at . Otherwise, the patch is made at <A>. DDT inserts in this location a branch to the current value of ;F. When the patch is done, ;F is updated. It then gives a carriage return and a) and waits for the user to type in the patch. Legal input consists of a series of expressions whose values are inserted in successive locations in memory. Each of these expressions should be terminated by line feed or ; ω , exactly as though the program were being typed in with the \ command instead of as a patch. The t command may be given in place of the line feed and has its usual meaning, except that the contents of the previous location are not typed. Two other commands are legal in patch mode. They are:
 - Colon, which may be used to define a local symbol with value equal to the current location.

- (2) Carriage return, which terminates the patch. When the patch is terminated, DDT inserts in the next available location the original contents of the location at which the patch was inserted. It then inserts in the following two locations branch instructions to the first and second locations following the patch. This means that if the patch command is a skip instruction, the program will continue to operate correctly. Any other command given in patch mode may cause unpredictable errors.
- <A>;I Is identical to the) command except that it puts the instruction
 being patched before the new code inserted by the programmer instead of after.

2.9 Miscellaneous Commands

- ;? and ? This commands erase everything typed since the last tab or carriage return. It is always legal.
- <E>;G (GO TO) <E>;G restores the A, B and X registers which were saved when DDT was entered (unless they have been modified) and transfers to the location specified by the value of the expression.
- ;K (KILL) This command resets DDT's symbol table to its initial state. DDT will type back --OK and wait for a confirming dot. Any other character will abort the command.

<S>:K (KILL). Removes only the symbol <S> from the table.

- <E>,<E>;L Sets ;1 and ;2 (the lower and upper brounds for searches) to the values of the first and second expressions respectively.
- ;U (UNDEFINED). This command causes all undefined symbols to be listed.
 <E>;U (EXECUTE). This causes the value of the expression to be executed as an instruction. If it is a branch, control goes to the location branched to. In all other cases control remains with DDT. A single carriage return is typed before execution of the instruction. If the instruction does not

branch and does not skip, or returns to the following location, a \$ and another carriage return are typed after its execution. If the instruction does skip, two dollar signs (\$\$) are typed followed by a carriage return.

;Z (ZERO) <E>,<E>;Z sets to zero all locations between the value of the first expression and that of the second. ;Z alone releases all memory accessible to the user's program. DDT will type back --OK and wait for a confirming dot. Any other characters will abort the command. If this memory is returned, due to later access by DDT or a program, it will be cleared to zero.

& (LIST BLOCKS). The names of all blocks are printed.

2.10 Special Symbols

The value of "." is the current location, i.e. the address of the last register opened.

The following symbols refer to various special registers of the machine. Their value is the contents of these registers as saved by DDT: ;X= will print the saved contents of the X register. To change the contents of a register, a command of the form $\langle E \rangle$;A is used. This command sets the A register to the value of the expression. Whenever DDT executes any command involving execution of instructions in the user's program, it restores the values of all machine registers. If any of these values have been changed by the user, it is the changed value which will be restored.

;A The value of this symbol is the contents of the A register.
;B The value of this symbol is the contents of the B register.
;X The value of this symbol is the contents of the X register.
;L The value of this symbol is the contents of the program counter.
The only reason for changing ;L is to set the location from which ;N will begin execution.

The values of the following special symbols are used by DDT in certain commands or are available to the programmer for his general enlightenment. These values may be changed in the same way that the values of the symbols for the central registers of the machine may be changed.

;M The value of this symbol is the mask for word searches.

;1 The value of this symbol is the lower bound for word and effective address searches. It may also be set by the ;L command.

;2 The value of this symbol is the upper bound for word and effective address searches. It may also be set by using ;L.

;Q This symbol has a value equal to the value of the last expression typed by DDT or the user. It is useful, for instance, if the programmer wishes to add one to the contents of the open register; he need only type ;Q + 1.

;F The value of this symbol is the address of the lowest location in core not used by the program. New literals and patches are inserted starting at this address. Note: like all other special symbols, ;F may be changed by the command $\langle E \rangle$;F. It is also updated as necessary by patches and literal definitions.

2.11 Panics

DDT recognizes four kinds of panic conditions:

- (1) Illegal instruction panics from the user's program.
- (2) Memory allocation exceeded panics from the user's program.
- (3) Panics generated by pushing the rubout button.
- (4) Panics generated by the execution of BRS 10 in the user's program.

For the first two of these conditions DDT prints out a message, the location of the instruction at which the panic occurred, and the contents of this location. The messages are as follows:

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- (1) Illegal instruction panic I > >
- (2) Memory allocation exceeded M > >
- (3) The other two types of panics cause DDT to type bell and carriage return. ;L and . will both be equal to the location at which the panic occurred.

If a memory allocation exceeded panic is caused by a transfer to an illegal location, the contents of the location causing the panic is not available and DDT, therefore, types a ?.

Two other panic conditions are possible in DDT.

- If the rubout button is pushed twice with no intervening typing by the user, control returns to the executive.
- (?) If the rubout button is pushed while DDT is executing a command, execution and typeout are terminated and DDT types carriage return and bell and then awaits further commands.

2.12 Multiple Program Debugging

It is occasionally desirable to hold several programs with different maps and symbol tables in DDT simultaneously. This situation could be approximated using the DUMP and RECOVER commands in the time-sharing executive, but several commands are provided in DDT itself to facilitate the process.

- $\langle W_1 \rangle, \langle W_2 \rangle$; R (SET MAP). The pseudo-relabeling for the program is set according to the value of $\langle W_1 \rangle$ and $\langle W_2 \rangle$. This command is essentially equivalent to executing BRS 44 with $\langle W_1 \rangle$ in A and $\langle W_2 \rangle$ in B.
- %E (ERASE). DDT types --OK and waits for a confirming dot. Any other character will abort the command. DDT then resets itself to its initial state, i.e. the symbol table, program map, breakpoints and modes are all reset. The program memory, however, is not released.

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- (DUMP). This command also requires a confirming dot. The entire state of DDT is saved away and a number typed out which will allow this state to be retrieved by the %R command (see below). DDT then resets itself as described under %E above.
- STR (RECOVER). This command requires a confirming dot. If the present state of DDT has ever been dumped (i.e. was produced by %R), it is dumped again. Then the state is restored exactly as it was when the %D was given, whose number was the value of <E>. Using an illegal number for %R can lead to chaos.

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