UNIVERSITY OF ILLINOIS DIGITAL COMPUTER

LIBRARY ROUTINE X 12 - 235

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TITLE

Symbolic Address Decimal Order Input (SADOI) (SADOI Itself)

TYPE

Input Assembly Program

NUMBER OF WORDS

615 Words in the locked out section of the drum

WORDS USED

999-1023 in the high speed memory

11,058-12,799 on drum

TEMPORARY STORAGE

0, 1

DESCRIPTION

This routine is stored on the drum. It is automatically brought into the Williams memory and is not included on the program tape. It is an input routine designed to read instructions with addresses in decimal or symbolic form. It is compatible with the D.O.I. Any program written for the D.O.I. (library routine X 1) and which refers to the D.O.I. only in standard fashion may be used with SADOI. It may be necessary to recompute sum checks for programs to be input by SADOI.

During input the entire Williams memory is used by SADOI. The instructions in SADOI occupy locations 32-562, 999-1023. The remainder of the Williams memory is used for temporary storage. Instruction pairs of the program being read are initially stored on the drum. The correspondence used is drum address 11756+X Williams memory address X. When SADOI is instructed to jump to the program which has been read, words of the program are transferred from the drum to the Williams memory and the jump is executed. A more detailed description is given under interludes.

Initially the reading of a program tape must be started by a hold start or a clear start. In either case Williams memory 2-998 is transferred to drum locations 11,758-12,754. SADOI is then placed in the Williams memory. Input then begins.

Each instruction consists of a pair of <u>function digits</u> and an <u>address</u>. Addresses are written in numerical or symbolic form. <u>Numerical addresses</u> are written and punched in decimal form with non-significant zeros omitted. Thus we write 1023, 35, 7. Each numerical address is followed by a terminating symbol K, S, N, J, F, or L. This symbol indicates a method of treating the numerical address and the instruction. For example, L5 29F and 40 7L. A zero address is indicated by writing a terminating symbol immediately following the function digits. For example, OOK and 26F.

Symbolic addresses are written and punched as a pair of parentheses; any combination of symbols or characters may be placed between the parentheses. Examples of distinct symbolic addresses are (), (0), (1), (51372), (WQXR). Generally speaking, any two combinations of teletype characters which produce the same sequence of printed characters will represent the same symbolic address.

Numerical and symbolic addresses are most frequently written to refer to memory locations. When symbolic addresses are used this reference is made by using the same symbolic address twice. One use is as the address of an instruction. The other use is as a marker showing which instruction is being referred to. The instruction 40 (A) says store the contents of the accumulator at (A), using "(A)" as an address of the instruction. We call this an inside symbolic address or inside address. Some other instruction pair will use "(A)" as an outside symbolic address or outside address. This instruction pair might be (A) OOF OOF. When this instruction pair is read by SADOI during input the location at which it is to be stored is remembered by SADOI as the real address corresponding to "(A)". This real address is substituted whenever instructions are written with the inside symbolic address (A).

There is no limit to the number of references that may be made to a particular memory location using symbolic addresses. Further, the word being referred to may appear anywhere on the program tape. When symbolic addresses are used SADOI determines the real addresses completely in terms of the relative positions of instructions on the input tape. The programmer will find that he normally thinks in terms of symbolic addresses while writing a program. Numerical addresses can be written only after bookkeeping that is made unnecessary by writing symbolic addresses. Further, the programmer will find that use of symbolic addresses greatly facilitates making corrections in program tapes. For, instructions often may be added or deleted without having to change addresses.

TERMINATING SYMBOLS

with a pair of sexadecimal characters, a numerical address and terminated by K. It tells SADOI where to store words of program. In addition it provides an address which is called the relativizer. This address is the origin against which relative addresses are measured (see discussion of L, below). The first instruction on a tape must be a directive. The instruction following a directive will always be a left hand instruction. The directive is never stored as part of the program being read.

There are four types of directives which are written as follows:

00 n K with n a non-zero numerical address. The following instruction pairs on the tape will go into locations n, n + 1, n + 2, The relativizer is given the value n.

00 K If m is the location of the last instruction pair, then set the relativizer to m + l. The instruction pairs following will go into locations m + l, m + 2,

Ol nK If m is the value of the relativizer then the instruction pairs following will go into m + n, m + n + 1, The relativizer is not changed.

02 nK Set the relativizer to the address m of the preceding instruction. The following instruction pairs will go into m + n, m + n + 1, When using this type of directive the preceding instruction will most often be a dummy instruction written solely to provide a value for the 02 K directive. For example, 00 (A) 02 K will result in the real address corresponding to (A) being used as the relativizer. When the preceding instruction (00 (A)) appears on the tape as a <u>left hand instruction</u>, it will not be stored as part of the program being read.

The programmer should be careful to use symbolic addresses with the 02 K directive only in cases where the symbolic address is previously used as an outside address.

F means fixed address. Thus L5 29F 40F when executed instructs ILLIAC to put the number at location 29 into A then, store this number at location 0. The pairs of instructions 00F 00F, 80F 00F, and LL 4095F LL 4095F are particularly useful as they set a memory location to 0, -1, -2^{-39} respectively.

L means relative address. If n is the value of the relativizer, the instruction 50 3L will have address n + 3.

N is used with a jump instruction to cause SADOI to jump to a specified location. Let n be the address of the jump instruction. If $n \le 998$ then the words of program are read back from the drum and put into the Williams memory. Control then transfers to the word at location n. The jump instruction is stored at location 1. It is not otherwise stored as part of the program being read.

If the address $n \ge 999$ then a new address m is computed and substituted into the jump instruction. m = 93 + (n - 999). The jump (stored at location 1) is then executed. By this means control may be transferred to any part of the Williams memory during input. The following jump instructions are useful:

24 999N: Stop input. This instruction has no other effect.

26 999N: Use to mark an address (see example 4)

26 1000N: Special interlude marker (see interludes, below)

22 1000N: Special drum storage order (see same, below)

26 1469N: Jump to the first word of temporary storage.

Example 1

00 100 K

L5 L 82 40 F

OF F OO F

26 100 N.

In this example instructions are stored at locations 100, 101. Control transfers to location 100. The first pair of instructions prints itself in sexadecimal.

J terminations used after right-hand instructions tell SADOI to convert the pair of instructions into a fraction. It is thus a way of reading a few decimal fractions along with the instructions of a program. For this purpose instructions are written in a special form. For example, the fraction .3333 3333 3333 should be written as 00 F 00 3333 33333 and the fraction .1 should be written as 00 F 00 1000 0000 0000J. The function digits of both instructions are zero and the decimal digits of the fraction are written as the address of the right instruction. Twelve decimal digits must always follow the second pair of zeros except that zeros immediately following the decimal point may be omitted. For example, write .001 as 00 F 00 10 0000 0000J. Any fraction in the range 0 to .9999 9999 9999 may be written in this way.

We can extend the range by using the function digits of the left instruction. By writing 80 for the function digits we can subtract 1 from the fraction. For example, the number 8888 8889 would be written as 00F 00 8888 8889 while _0.8888 8889 = -1 +.1111 1111 1111 would be written as 80F 00 1111 1111 1111J. (see also page 17)

S This termination is used to refer to S-parameters. Discussion of S-parameters is deferred until later.

USE OF SYMBOLIC ADDRESSES

Example 2

The use of symbolic addresses can be illustrated by the following example. In this example library subroutine P 10 has been rewritten using symbolic addresses. This routine is used to print headings.

	myambie s	<u>-</u>	
(1)	O4F	L5 F	
	32 (2)	50 (3.)	5th hole character
(2)	26 (4)	50 (1)	non 5th hole character
(4)	00 6F	42 (3)	
(3)	02 1F	92 F	Substitute proper address in the 92 instruction
(5)	91 4F	40 F	Enter loop
	L7 F	32 (1)	5-hold delay? If not, repeat.
	92 7 07F	26 999F	Exit from loop. See discussion of interludes.
	26 (5)	26 lN	

For this example to be complete it should be preceded by a directive. Each pair of instructions referred to in the program is marked by an outside address. Thus, in the first instruction pair (1) 04F L5F, (1) is used as an outside address. Similarly (2), ..., (5) are used as outside addresses. Certain of the instructions use these symbolic addresses as addresses. For example, 32 (2) 50 (3) which instructs ILLIAC to jump to the right side of (2) if the accumulator is positive. Otherwise, put the contents of (3) into Q register.

It should be noted that forward and backward references can be made. The third word contains the instruction 50 (1). This refers back to an instruction pair already read by SADOI. The reference is called a backward reference. In the same word, the instruction 26 (4) makes a forward reference to an instruction pair not yet read by SADOI.

In example 2, location 0 is used for temporary storage. This means that no matter where the block of instructions is located in the Williams memory, location 0 will be used for storing results temporarily during the computation. Thus, fixed addresses are used. In the first word, for example, the instruction L5F refers to this temporary storage location.

Last, addresses of certain instructions do not refer to memory locations. In shift instructions, for example, the address further supplements the function digits by telling how far to shift. In such cases fixed addresses are almost always written.

WRITING SYMBOLIC ADDRESSES

An outside symbolic address is written before the function digits of either instruction of a pair. In example 2, the first instruction pair could also have been punched as

Further, an instruction pair may have several outside addresses. For example, (1) (2) (3) (4) 40 F (5) 81 4 F.

An inside symbolic address is written immediately following the function digits. Only one inside address may be used in an instruction. It must not be followed by a termination K, S, N, J, F, L.

Constants may be added to inside addresses. Thus, the instruction L5 21 (N12) will refer to the 21st word past the word marked by (N12) as an outside address. Generally the form "n ()" will be used to add n to the real address corresponding to ().

The full range of available teletype signs may be used in writing symbolic addresses: numbers, letters, punctuation marks, etc. During input the characters placed between parentheses are read and converted to numerical equivalents. From these numerical equivalents a symbolic address word is formed.

As stated earlier, two combinations of teletype characters which produce the same sequence of printed characters will represent the same symbolic address. Delays, spaces, and LF-CR characters do not affect the symbolic address. Figures-shift and letters-shift characters, on the other hand, may affect the symbolic

address. After reading the beginning parenthesis of a symbolic address, SADOI assumes each character following is in figures-shift until a letters-shift character is read. Thus, input of a symbolic address is not terminated by "A". The symbolic address (lA) would be punched using the following six characters: "(", "l", "LTR-SH", "A", "FIG-SH", ")". The symbolic address (l2) would be punched using four characters.

Thus + and K, - and S are different. Similarly the letter "X" and times "x" are different. The number "0" and letter "0", because they look the same when printed, are given the same numerical equivalent when read. N, J, F, L print the same regardless of the print shift. Each is given the same numerical equivalent in both print shifts when read. The letter "0" or number "0" is different from the absence of a character. Thus, (), (0), and (1), (01) are four distinct symbolic addresses.

The symbolic address recognized by SADOI consists of the last five printed characters. Thus, "(1A2B3C)" is the same as "(A2B3C)" but different from "(2B3C)".

USE OF TEMPORARY STORAGE LOCATIONS O, 1 BY SADOI

Location 0 is used by SADOI to compute and store numerical addresses during the input of instructions. At the beginning of interludes it is used to compute a sum check.

Location 1 is used by SADOI to form instruction pairs. Once a pair of instructions has formed a word at location 1, the contents of location 1 are stored as part of the program being input. The only cases in which the contents of location 1 are not stored as part of the program being read is when K or N terminated instructions are read.

The right side of location 1 always contains the most recent instruction read. The left side of location 1 always contains the immediately preceding instruction. Consider the following examples:

Example 3

(a) (b)

00 F 26 100 F 26 100 F 26 1N

26 1N L5 F 40 F

L5 F 40 F

In example 3a, the instruction 26 100F appears as a right-hand instruction. The word formed from 00F 26 100F is stored as part of the program being read. When the jump 26 lN is executed, the contents of location 1 is 26064 26001 (sexadecimal). When input resumes, both instructions L5F and 40F must be read before the contents of location 1 are again stored as part of the program. In example 3b, the instruction 26 100F appears as a left-hand instruction. After the right hand instruction 26 lN is read the contents of location 1 appear as before: 26064 26001. This word is not stored as part of the program because of the N termination. As before, the next pair of instructions, L5F and 40F, are read before the contents of location 1 are stored as part of the program being read.

During input when an inside symbolic address is used and a backward reference is made, the real address is substituted immediately into location 1. If an instruction like L5 21 (N12) is written, the numerical part, 21, has been computed and is stored at location 0. When (N12) makes a backward reference, the real address is substituted. Then, N(0) + N(1) is computed and stored at location 1.

When forward references are made during input, sufficient information is retained to allow SADOI to form the correct addresses later. A pseudo address is temporarily substituted into the word at location 1.

S PARAMETERS

Words stored in locations 2-15 of the Williams memory are called S parameters. These numbers give a way of modifying instructions during input. Following the S termination must be a d=2, 3, 4, 5, 6, 7, 8, 9, K, S, N, J, F, L. This gives the address (in sexadecimal) of a location which is added to the S-terminated instruction before it is stored in the memory.

The <u>right-hand</u> instruction of an instruction pair stored in locations 2-15 may have an inside symbolic address and still be used as an S parameter. The left hand instruction <u>must not</u> have an inside address if it is to be used as an S parameter.

The directives OOK, OlK and O2K enable the programmer to minimize the number of memory locations needed to store his program. Use of the O2K directive for this purpose is illustrated in example 4. An S parameter is stored at location 5. After the interrupting instructions are read SADOI is set so that the following instructions will be stored immediately after the earlier instructions.

Example 4

L5. F 40 1F

(A) 26 999N

Introduce address (A)

00 5K

Store S parameter at 5

OOF OO 12F.

00 (A) 02K

Reset SADOI to begin storing at (A)

L5 1F 82 S5 ·

Instruction using S parameter

The instruction (A) 26 999N has no effect other than to introduce the symbolic address (A).

INTERLUDES

whenever SADOI is told to jump to instructions which have been read, an interlude begins. This is done by writing an N-terminated jump instruction with address less than 999. The jump instruction may appear on tape as if it were a left-hand or a right-hand instruction. Pairs of instructions may be written. For example, the instructions 26 L 26 lN or 26 (A) 26 lN may be used to start interludes. In these cases, the instruction pair is stored at location 1 only.

The system used with the Decimal Order Input (D.O.I.) in which OOF 26L

. . .

26 IN

is written to begin an interlude may also be used. During input the last two instructions read are stored at location 1. The next to last instruction is always the left-hand instruction of the pair stored at location 1.

There are three types of interludes possible with SADOI.

REGULAR INTERLUDES

Regular interludes happen when SADOI is not explicitly told differently (see below). To get ready for the interlude real addresses are substituted into instructions wherever necessary. After the substitutions are completed drum locations 11,758 to 12,754 inclusive are transferred to Williams memory locations 2-998 inclusive. Control then transfers to the right side of location 1. The interlude begins.

Since at the beginning of a regular interlude substitutions of real addresses are made, the programmer should be careful about putting on tape library routines having interludes. The programmer should put such routines at the beginning of his tape before inside addresses are first used or else convert the interlude to the special interlude discussed below.

Use of a symbolic address as an inside address but not as an outside address is a programming error. SADOI detects all errors of this type before stopping. Information is punched about each error (see below). After all possible substitutions are made and errors indicated, the instructions on the drum are transferred to the William's memory. An FFO31 stop then occurs.

The symbolic addresses used prior to an interlude are remembered. This information is stored on the drum during the interlude at addresses higher than 11,000.

A small part of the instructions in SADOI are stored at 999-1023. These instructions enable the programmer to jump to SADOI after an interlude for more input. Two types of jumps may be used. When a jump to the left side of location 999 is made the right hand address of the number in Q is used as a new relativizer. If this address has value m then the following instruction pairs will go into m, m + 1, m + 2,... until a directive is received. When a jump to the right side of location 1014 is made the old relativizer is retained. The right-hand address of A determines where the following instructions will go. If this address has value m the following instruction pairs will go into m, m + 1, m + 2, ...

When a regular interlude is ended by one of these jumps the contents of the Williams memory are put back on the drum. Locations 2-998 are moved to 11,758 - 12,754 respectively. SADOI is then put into the Williams memory, information about the previously used symbolic addresses is restored, and input begins.

SPECIAL INTERLUDES

Special interludes are provided so that a small amount of calculation may be accomplished without requiring substitution of real addresses for symbolic addresses and without transferring all of 1000 words from one memory to another. To write a special interlude precede the first instruction pair to be used by the instruction 26 1000N. The first instruction following an N. termination will always be a left-hand instruction. Follow the last instruction pair by a normal jump instruction terminated by N. Only the instructions lying between the two jump instructions will be brought from the drum for use in the interlude.

Example 5

Input one sexadecimal number

OOK

26 1000N

80 40F 40L

22 1L L5 1L

22 1014F 26L

26 IN

0000000000

The instruction 26 1000N marks the first word to be used in the special interlude. Let the Williams memory location for this word be n. When 26 1000N is read n is substituted into the right-hand address of the word in 1022. This address will have been set to zero previously, and by testing to see whether or not this address is zero SADOI determines whether a regular or a special interlude is to occur.

Let n_2 be the Williams memory location for the last word of the interlude which is read from the tape. Then the words brought from the drum will be placed in locations n_1 , $n_1 + 1$, ..., n_2 in the Williams memory. The programmer should be careful not to store instructions to be used outside this range. This may happen if directives are used between the 26 1000N instruction and the normal jump instruction.

Before a special interlude begins, information about symbolic addresses previously used is stored on the drum. After the interlude this information is brought back to the Williams memory again.

Special interludes, like regular interludes, are ended by jumps to 999 or 1014. The meanings of these jumps are the same as for regular interludes. After the end of the special interlude the S parameters 2-15 are copied onto the corresponding drum locations 11,758 - 11,771. Aside from the S parameters only the words in locations used for the special interlude are copied onto the corresponding drum locations.

FAST INTERLUDES

Fast interludes may be written when 44 instruction pairs or less are needed. During input, blocks of instruction pairs are accumulated in the Williams memory before being transferred to the drum. These transfers to the drum take place after

- (1) 45 instruction pairs are accumulated
 - (2) a directive is read
 - (3) an W terminated instruction with address less than 999 is read
 - (4) an S termination is used and the S parameter has a right inside address for which a real address must be substituted.

The addresses used for the temporary storage are 563 - 607. If the jump 26 1469 + k F is written then according to the rule cited earlier the fixed address 1469 + k is changed to

$$(1469 + k - 999) + 93 = 563 + k$$
.

Consequently the instruction 26 1469N will cause SADOI to jump to the left side of location 563.

Fast interludes must be ended by a jump to the left side of 93. Input resumes immediately. If the instructions used in the interlude go to addresses $m, m+1, \ldots, m+k$ then the instructions following will go into $m+k+1, m+k+2, \ldots$ until a directive is read.

Example 5 is rewritten for a fast interlude and presented in

Example 6

Input one sexadecimal number

00 K

80 40F 40 563F

26 93F 00 F

26 1469N

01 1K

In this example fixed addresses are used throughout. In general the instructions will not go to the same place as that occupied by them when the interlude is executed.

LIBRARY ROUTINES STORED ON DRUM

Certain of the more commonly used ILLIAC library routines have been stored in the locked out section of the drum memory. These routines are, at present, . Pl6, Yl, Nl2, Rl, S4, S5, T4, T5, Fl, Al.

By use of a modified directive, SADOI is instructed to take a specified routine from the drum and insert it in the program being read.

A modified directive like (Pl2) 00 100K consists of two parts:

- (1) An outside address which specifies the library routine.
- (2) A directive which specifies the relativizer and first memory location for the program.

Any of the other types of directive may be used. For example, (N12) 00K will cause N12 to be placed following the last instruction read. Similarly, the 01K and 02K type of directives will affect the relativizer and store addresses as described earlier.

When a library routine is obtained using a modified directive, the outside address is inserted into the list of symbolic addresses. The real address which corresponds is the location of the first word of the library routine. Thus, for example, when (N12) OOK is written to obtain N12, the entries to N12 may be written:

50 n

50 q

26 (N12).

The programmer may use instructions like L5 21(N12) to refer to words interior to a library routine.

When punching teletype tapes the programmer should be careful how he punches P16, Y1, etc. The letters must appear in <u>letters shift</u>. For example, (P16) is punched as

"(", "LTR-SH", "P", "FIG-SH", "1", "6", ")"

using seven teletype characters. Note, also that the use of a hyphen in writing P-16 is incorrect.

Miswriting the designating outside address, or asking for a library routine not stored on the drum will result in an immediate FF 030 stop.

Library routines, stored on the drum, which use interludes (Fl for example) execute their interludes without affecting in any way the symbolic addresses used by the programmer.

SELF CHECKING FEATURES

FF 824 This stop occurs because of a sum check failure in transfers of routines from the locked out section of the drum. Depending on the characters of the input tape last read before the stop, the routine may be

- (1) Part of SADOI which was incorrectly placed in the Williams memory
- (2) A library routine called for by the input tape which was incorrectly read back from the drum.

FF 030 This stop occurs when the programmer asks for a library routine not stored on the drum. Input stops immediately.

FF 031 This stop indicates that the programmer has misused his symbolic addresses. The nature of these errors is always indicated by output on the punch. The FF stop occurs at the beginning of the 1st regular interlude following an error. At this time all instructions read by SADOI are in the Williams memory.

Three types of coding error are covered by this stop.

S parameters: A word having an inside address is used as an S parameter and the real address corresponding to the inside address has not yet been given to SADOI (by use of an outside address). The punched output appears as S 004 355 H, consisting of "S" which identifies the type of error, two decimal addresses, and the symbolic address involved. The first decimal address specifies the location of the S parameter. The second decimal address specifies the location of the instruction using the S parameter.

If the order was, for example, L5 nS4, then this instruction appears in the Williams memory as L5 nF when the stop occurs.

Outside type error: The same symbolic address is used as an outside address for two instructions having different locations. In this case the punched output appears as:

0 367 528 Q.

"O" signifies "outside". The first decimal address is the location of the first instruction, the second decimal address is the location of the second instruction. The symbolic address involved is printed last.

During input, when the real address corresponding to a symbolic address is known, the real address is substituted into instructions immediately. In case of duplicate outside addresses, the most recent real address value is the one saved and used during subsequent input.

Inside type error: When a regular interlude is to begin, real addresses are substituted into instructions. When a zero real address or no real address is available, this type of error is indicated. The punched output appears as:

473 K3.

This first thing punched is the location of the instruction using the inside address. The symbolic address involved is punched second.

FF 032 This stop may occur when a directive appears as a right-hand instruction. The stop indicates that the left-hand instruction which preceded used an inside address which made a forward reference. This may mean that input is out of phase one-half word or that an 02K directive cannot be executed. Input stops immediately.

FF 033 This stop occurs when an N terminated order is read and

- (1) it appears as the right-hand order of an order pair, and
- (2) the left-hand order of the order pair has an inside address for which the corresponding real address is not known, and
- (3) the address of the right hand order is ≤ 998.

This error stop may occur at the start of special or regular interludes. The stop is made after the appropriate instructions have been transferred to the Williams memory.

There may also be output on the tape punch.

FF 034 This stop occurs when too many instructions of the type AB n (C) are written with $n \neq 0$. Each time such an instruction is read and the real address corresponding to "(C)" is not known (forward reference), a word is added to the <u>additions</u> list. If the real address is known, the additions list is not increased. This list is limited to 314 entries.

FF 035 This stop occurs when too many distinct symbolic addresses are used. The number of symbolic addresses is limited to 324.

SUM CHECKS

Library routine X-7, used for sum checks, is not well suited for use with SADOI. It is written as an interlude ending with a 36 999F jump. Ending an interlude causes SADOI to reload itself into the Williams memory, thus causing additional drum transfers which may be omitted by using a more efficient sum check.

At the beginning of interludes SADOI automatically computes a sum of the words (order pairs or numbers) transferred from the drum memory to the Williams memory. At the end of the transfer this sum is in location zero and may be used by the programmer for checking.

If the words transferred are x_1, \dots, x_n , the sum computed is defined by the relations $s_0 = 0$ $s_{i+1} = x_{i+1} + |s_i| \qquad i+1 = 1, 2, \dots, n.$

A sum check routine suitable for SADOI may be written as follows and placed at the end of a tape.

00996K

L3F 36 (A)

Let (A) represent the

FF F 26 (A)

location of the first

000000F000000F

instruction to be obeyed

26L 26 1N

in the program

The last word in this routine is the sum check and should be punched by hand in the form of an instruction after the program is tried the first time. This word is obtained then by reading the contents of the accumulator when the FF stop first occurs.

SPECIAL DRUM LOADING TECHNIQUE

A special instruction has been provided which will allow the programmer to load blocks of instructions directly into any desired drum locations. Instructions to be loaded this way must be written without symbolic addresses. A sum check is not computed.

The following sequence must be followed:

- (1) A directive, which sets the relativizer.
- (2) Follow the directive by OC nF 22 1000N. The order pairs following on the tape are stored on the drum at n, n + 1, n + 2,
- (3) To stop the direct loading, write a directive or begin an interlude (regular or special).

Library routines stored on the drum and containing <u>less than 45</u> words may be transferred by this technique. For example,

(R1) 00 100K

00 3000F 22 1000N

will put Rl on the drum starting at 3000. The relativizer used is 100.

Special care must be taken when this technique is used. All bookkeeping work of SADOI proceeds as if the instructions being read were being normally stored. Consequently, if instructions previously read used inside addresses and if these instructions would be overwritten were the new instructions still normally stored, then the bookkeeping will be incorrect.

DETAILS OF THE FORMATION OF INSTRUCTIONS DURING INPUT

During input, decimal addresses are computed and stored at location 0. The pair of instructions being formed is always stored at location 1. The following sequence is followed:

- (I) Instruction type AB n (C)
 - 1. Put N(1) in A.
 - 2. Read two function digits "A" and "B" shifting A left eight places.
 - 3. Shift A left 12 places and store at 1.
 - 4. Read the address n and store at 0.
 - 5. Read the symbolic address (C).
 - a. Forward reference (real address not known).
 - (1) Substitute a false address into the right-hand instruction at location 1.
 - (2) Test for n = 0. If zero, begin reading the next instruction.
 - (3) If n ≠ 0 put a new word into the additions list. Left most 30 bits are used for the number n. Right most 10 bits determine the symbolic address (c) used. Begin reading the next instruction.
 b. Backward reference (real address known).
 - (1) Substitute real address into right-hand instruction at 1.
 - (2) Form N(0) + N(1) and store at 1.
 - (3) Begin reading the next instruction.
- (II) Instruction type AB n D Da K, S, N, S, F, L termination.
 - 1. Repeat steps 1-4 of I above
 - 2. K, N, F; form N(0) + N(1) and store at 1. In case of F, begin reading the next instruction. In case of K, N, finish execution of the termination.
 - 3. L; form N(0) + N(1) + relativizer and store at 1. Begin reading the next instruction.
 - 4. Sk; form N(0) + N(1) + N(k) and store at 1. Begin reading the next instruction.
 - 5. J; form

$$N(1) + \frac{1}{2} \left[\frac{N(0) - 5 \times 10^{11} \times 2^{-39} + 5 \times 10^{11} \times 2^{-79}}{5 \times 10^{11} \times 2^{-39}} \right] + \frac{1}{2}$$

and store at 1. Begin reading the next instruction.

DATE 7/11/57 RT: 3/5/58:12/3/58
PROGRAMMED BY R. H. Farrell
APPROVED BY D. E. Muller

PROGRAM NOTES

INTRODUCTION

The following is by no means a complete analysis of SADOI. Some parts of the program are only named while others are analyzed in detail. The following tries to present the ideas which have made SADOI a usuable input routine. Input is something that happens in time. When one understands the general sequence of steps and a few of the minute details one should then be equipped to understand what actually takes place in ILLIAC.

INPUT STARTS

When instructions being read by SADOI are on tape the input taking place is called <u>tape input</u>. When the instructions are on the drum (reading a library routine) the input taking place is called <u>drum input</u>.

The main part of SADOI fills Williams memory locations 30-604. An <u>input start</u> occurs at any point that the main part of SADOI is written into the Williams memory. Immediately after each input start tape or drum input begins.

Interludes are classed as <u>tape controlled</u> or <u>drum controlled</u> according to the type of input at the start of the interlude. Fast interludes are not followed by input starts (and hence may be used to modify SADOI), while regular and special interludes are. SADOI distinguishes four classes of input start:

- (1) a hold or clear start
- (2) tape controlled regular interludes
- (3) tape controlled special interludes
- (4) drum controlled regular or special interludes.

During each input start approximately 200 instructions are used to set up SADOI to begin input. The overall sequence followed during input starts is shown in diagram 1.

TAPE INPUT SEQUENCE

To initiate a clear or hold start one character is read by ILLIAC. SADOI assumes this was the first of a pair of function digits and fakes in a zero for this digit. In this case input starts by reading the second function digit. Consequently, a tape may not be started by a symbolic address.

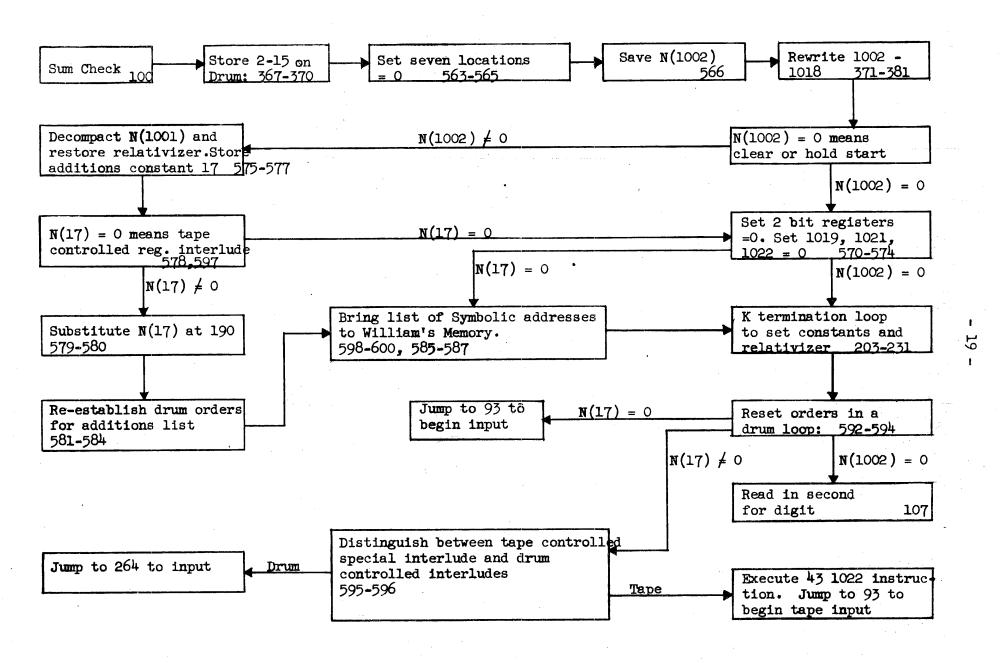


Diagram 1: INPUT START SEQUENCE

Tape input is started (following interludes) by a jump to location 93 after the set up operations. The jump instructions at 93 provide a variable link to the two main input sections of SADOI. A 26 999N instruction, when read by SADOI, also results in a jump to 93.

The input sections are so called because they contain the instructions that read coded information and translate it into numbers. Various other sections are excluded because their functions are to execute the terminations or bookkeeping. Speed of operation while reading tapes was deemed critical. Consequently, some bookkeeping uses available time between reading of characters on the tape. In other cases similar bookkeeping operations will be found repeated several times in order to eliminate the need for testing which phase of the input was last completed.

- I. Tape input section: 43-69, 104-147
 - A. Read decimal address 43-69
 - B. Read symbolic address 115-147
 - C. Read second function digit 104-114
- II. Drum input section: 262-327.

The bookkeeping and terminations will be discussed in detail below. The following diagram shows the tape input sequence.

DRUM INPUT SEQUENCE

The maximum compacting by Y4 is to store N words of program in 4/5N locations. Consequently, 36 compacted words placed in the Williams memory will yield at most 45 words of program. For this reason 45 words temporary storage was allowed, locations 563-607. In operation the drum input section brings 36 words from the drum and stores them in locations 572-607. The words are then picked off one at a time and decoded. The resulting instruction pairs are stored in 563-607, starting at 563.

For many library routines stored on the drum 36 words are sufficient to contain the entire routine. When 36 words do not suffice and the current 36 words have been exhausted, the assembled part of the routine is transferred to the drum. Another block of 36 words is put into the Williams memory. The process thus continues until the entire routine has been obtained.

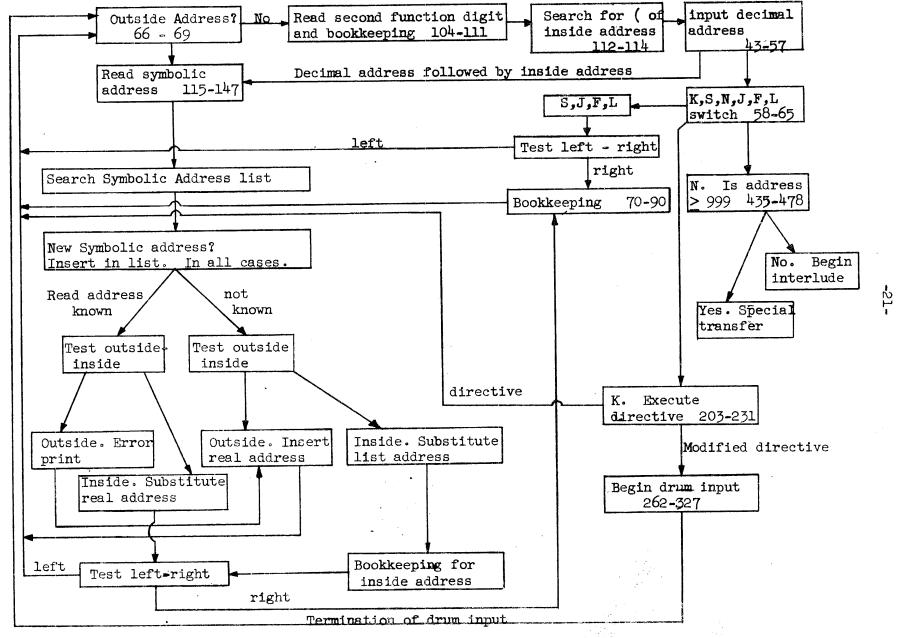


Diagram 2: TAPE INPUT SEQUENCE

If during drum input an interlude is initiated, certain information must be retained in order to restart the drum input after the interlude. Consequently, locations 1019-1021, 1023 are used for temporary storage by the drum input section. They are set equal zero at all times except during drum input. Their use is

1019: A counter for which of the 5 blocks of 8 bits was last taken.

1020: Word currently being decoded.

1021: Drum order set to bring next word from drum.

1023: Sum check on library routines.

In addition 1022 is used when a library routine calls for a special interlude.

Where possible the drum input section was interwoven with other parts of SADOI. This was done to shorten the program but also to keep the drum input compatible with tape input. Thus the K, S, and N termination sections of SADOI are used during both types of input. The bookkeeping section (69-92) is also used during drum input after each instruction pair of a routine is formed. Because of this sharing

- (1) S parameters may be automatically preset during drum input
- (2) The special transfers (N terminated jumps with address ≥ 999) may be used during drum input
- (3) The various types of directives may be used during drum input
- (4) S parameters having symbolic addresses will be properly handled. The following diagram shows the drum input sequence.

READING SYMBOLIC ADDRESSES

The parentheses "(" and ")" have no effect other than to initiate and terminate reading a symbolic address. Each character on the tape following "(" is examined by means of a 32 position switch (132-147) and has one of the following effects:

- (1) None. Spaces, Jelays and LF-CR characters are skipped.
- (2) Change from letters to figures shift or from figures to letters shift.

 Letters and figures shift characters.
- (3) Terminate reading of the symbolic address. Closing ")".
- (4) Add six bits to the symbolic address word being formed in ILLIAC

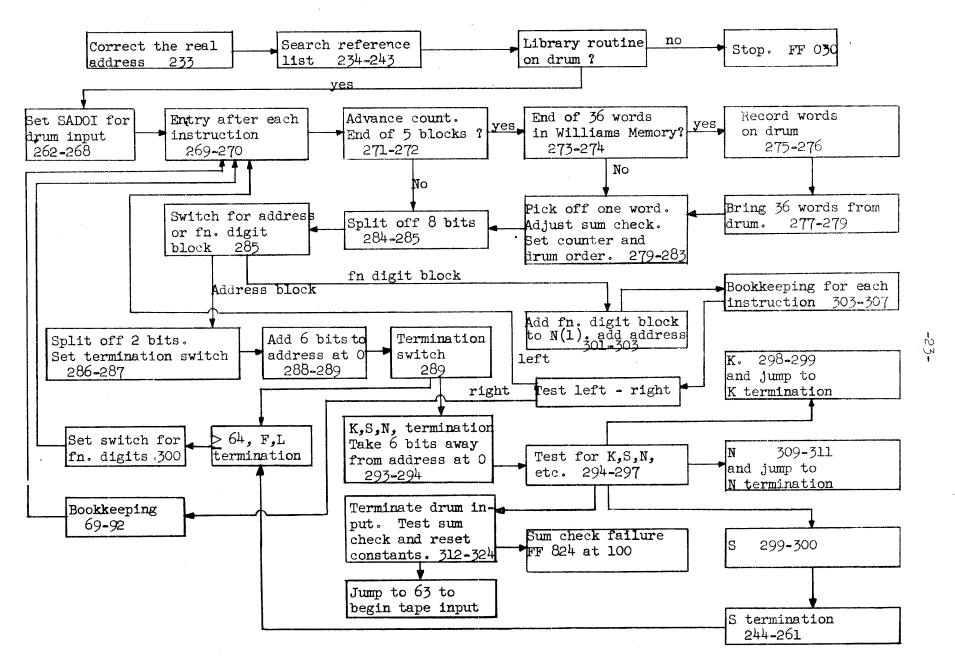


Diagram 3: DRUM INPUT SEQUENCE

In general any character (except parentheses) which will cause a printed mark causes the addition of six bits to the symbolic address word being formed. The composition of the block is as follows: let $b_0b_1...b_5$ be the digits of the block. Then

b₀ = 0 figures shift = 1 letters shift

b₁ = 0 non-5-hole character = 1 5-hole character.

 $b_2, ..., b_5$ correspond to the 8, 4, 2 and 1 hole positions of the tape respectively and $b_i = 1$ when the position is punched, $b_i = 0$ when the position is not punched.

Gertain exceptions mentioned in the description of SADOI are made to this rule.

The most recently added block of six bits is always added to the symbolic address word on the left end by means of a right shift into Q. Part of this right shift is done by the 90 instruction at 118. The shift of four positions right occurs before the new digits read from tape are gated into A. Since the 90 instruction does not clear A, care is taken to make sure N(A) = 0 after the right shift. In particular the constant $N(34) = -1 + 16 \times 2^{-39}$ which is added to N(A) for 5-hole characters eliminates the minus sign introduced by the 90 instruction.

SYMBOLIC ADDRESS COMPILER

The symbolic address word is formed by successive right shifts into Q. The first step following the reading is to execute the following instructions:

SJ F 00 1F 40 946F 43 946F 89 1F L4 946F 40945F.

These instructions complete the formation of the symbolic address word and stores it temporarily at 945, the head of the symbolic address list. The word with sign digit changed is stored at 946.

Each time a symbolic address is read the list is searched to determine whether a corresponding symbolic address word is in the list. In long programs the searching time becomes significant since the total time is an n² process. The arrangement used in SADOI results from the following arguments (see words 153-162).

- (1) Then coding for ILLIAC loop speeds are faster when addresses are increased than when decreased.
- (2) If a symbolic address was previously used it was most likely one of the more recently received addresses. Therefore the search through the list should begin at the end of the list containing the most recently received symbolic address words.
- (3) The first symbolic address word is therefore stored at 944. Each new word is stored at the next lower address.
- (4) By storing each new word temporarily at 945,946 and beginning the search at the low address end, the search will always be automatically terminated, either by finding a word in the list or by finding the new word at the head of the list.
- (5) A loop to accomplish the search requires 3-1/2 words with a subloop of 2-1/2 words. Read around ratio for this loop was well above 200. It is consequently duplicated to reduce the read around ratio value.

The right address position of the symbolic address word is used to store the <u>real address</u> corresponding to the symbolic address. When examining a symbolic address word SADOI determines whether the real address has been received by whether the right-hand address is non-zero or not.

When a symbolic address word is already in the list, then, when the search terminates, the right address of N(A) is the right address of the symbolic address word. This address is saved by executing a 42 17F instruction.

An inside symbolic address makes a <u>forward reference</u> when the real address is not known and a <u>backward reference</u> when the real address is known. These cases are treated separately. At some point real addresses must be substituted wherever inside addresses were used. This is done immediately when the real address is known. The compiler thus treats the following cases:

- 1. Real address known $N(17) \neq 0$
 - A. Symbolic address read was an outside address. Then an error has been made. (165-166). Indicate error and proceed as in II B.
 - B. Inside symbolic address was read. Substitute real address into right address of N(1). Then form N(0) + N(1) and store at 1. (172-174)

II. Real address not known

- A. New symbolic address. Determine whether word found in the list was N(945). If so put N(945) at the bottom of the list and proceed as in II B or II C.
- B. Outside symbolic address was read. Substitute the real address into the symbolic address word.
- C. Inside symbolic address was read
 - N(0) = 0 so no constant is to be added. Substitute into N(1) the location of the symbolic address word in the Williams memory (a number ≥ 512).
 - 2. N(0) ≠ 0. Form a new word composed as follows: The left 30 bits equal the right 30 bits of N(0); the right address is the location of the symbolic address word. Store this word on the drum in the additions list until needed. Substitute into N(1) the location the additions word will have in the Williams memory when it is used (a number < 512).</p>

BACK SUBSTITUTION PROCESS (479-562)

of where real addresses must be substituted later. A digit 1 means a substitution must be made while a digit 0 means no substitution is required. One register thus serves 40 instructions or 20 instruction pairs that have been read. Let an instruction pair have location n (the address it will ultimately have in the Williams memory) and let

$$n = 20q + r$$
 with $0 \le r < 20$.

Then 947 + q is the location of the 2 bit register. After each directive this constant is computed by the K-termination section (see 217-226) and stored at 26. If $a_0a_1\cdots a_39$ is the 2 bit register then a_{2r} is the digit position for left instructions, $a_{2r}+1$ for right instructions.

Throughout input a <u>counter word</u> is stored at location 25. This word has one non zero digit a_i where i = 2r or 2r + 1 according as the instruction being read is a left or right instruction. The counter word is used to add or delete non zero digits from the 2 bit registers as each new instruction is read.

During the back substitution for real addresses the 2 bit registers are successively tested to determine whether they are equal zero. When a non-zero 2 bit register is found the corresponding 20 words are brought from the drum. The required substitutions are made and the 20 words replaced on the drum.

The parts of SADOI required during the back substitution are the temporary storage locations 16-29, the numerical constants 30-42, and instructions 357-562. The list of additions words stored on the drum are transferred to the Williams memory overwriting locations 43-356 (the upper limit is variable). The list of symbolic addresses is then transferred to the drum overwriting the additions list (which is not to be saved). The back substitution then begins.

When a substitution is required the address n of the instruction involved is examined:

- (1) n < 512 means an addition is required.
- (2) n > 512 means no addition is required. The following diagram shows in detail the operation of instructions 479-562.

USE OF THE WILLIAMS AND DRUM MEMORIES AND SUM CHECKS

SADOI is permanently stored at locations 1027-1641 on the drum. It is set so that the first word of the reference list is stored at 1850. To change this address it is necessary to change the drum instruction at 1252 and adjust the sum check at 1612.

When a clear or hold start is used the drum bootstraps load locations 1614-1640 into the Williams memory locations 999-1025. Control transfers to 1020 of the Williams memory. The words at 1019-1025 are used during hold or clear starts to restore Williams memory locations 2-32 after the drum bootstraps have finished their work. They are not reloaded into the Williams memory during other input starts.

The sum check at 1641 is used by the drum bootstrap routines but not otherwise. The sum check at 1612 is used each time the main part of SADOI is put into the Williams memory. Let $N(n) = X_n$ be the number stored at location n while the main part of SADOI is being brought to the William's memory. The drum loop computes the following sum

$$Y_{31} = X_{31}$$

 $Y_{n+1} = X_{n+1} + |Y_n|$

The values Y_n are stored at 31. Let $X_{603} = - |Y_{602}| + X_{604}$. X_{603} is the sum check constant stored on the drum at 1612. X_{604} is the last word brought from the drum. Therefore $Y_{604} = X_{604} + |Y_{603}| = X_{604} + |X_{603}| + |Y_{602}| = X_{604} + |X_{604}|$. Since X_{604} is a negative number, $Y_{604} = 0$.

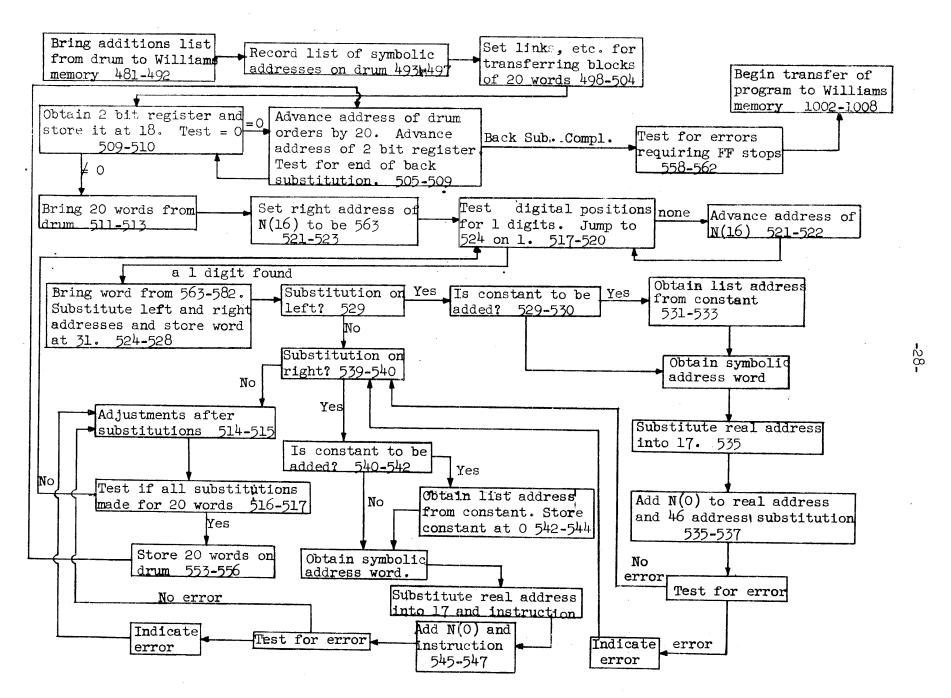


Diagram 4: BACK SUBSTITUTION

During input instruction pairs are stored on the drum. Locations 11,756-12,799 are allocated for this purpose. The correspondence used is location $X \longleftrightarrow drum \text{ address } 11,756 + X$. The bookkeeping of SADOI is so set that any location $0 \le X \le 1039$ may be used. However, only words having locations $2 \le X \le 998$ will be brought from the drum during interludes. The extra words which are available on the drum can be obtained only by special coding.

Storage locations with addresses l1,000 < u \leq l1,755 are used to store the various lists computed by SADOI.

During input the only other material stored on the drum are the additions constants. Small blocks (13 words) are held temporarily in the Williams memory at 608-620. When a full block is accumulated the words are transferred to the drum reversing their order during the transfer. The first block of 13 are stored at 11,742-11,754.

During interludes the manner in which the various lists are moved depends on the type of interlude.

- I. Regular tape controlled interludes

 The additions list and 2 bit registers are not saved. The list of symbolic addresses is stored on the drum. The order of the list is reversed during the transfer so that the low address end of the list is stored at 11,754.
- II. During all other types of interludes (excluding fast interludes)

 The three lists are stored on the drum. The additions list, already there, is not moved. The Williams memory from the low end of the symbolic address list to 998 is stored on the drum "below" the additions list, the order being reversed during the transfer. The combined length of 2 bit registers, symbolic address list and additions list cannot exceed 700 words.

The use of the Williams memory is:

- 0. Sum check and decimal addresses computed during input
- 1. Instruction pairs being formed
- 2-15. S parameters
- 16-31. Temporary storage for SADOI
 - 16. General purpose temporary storage
 - 17. During input of symbolic addresses, $N(17) = 32 \times 2^{-39}$ for letters shift, = 0 for figures shift. Used only for right address substitution otherwise.
 - 13. Used during reading of symbolic addresses and back substitution.

- 19. General purpose storage
- 20. = 0 if no outside addresses for current instruction pair. Otherwise contains the location of the most recently received outside address for the instruction pair.
- 21. \(\neq 0 \) when reading outside addresses = 0 when reading inside addresses
- 22. A word with at most 2 non zero digits. These digits are the positions in the 2 bit register corresponding to current instruction pair.
- 23. Bits to be added to the two bit register
- 24. Temporary storage for the current 2 bit register
- 25. Counter word
- 26. Address for current 2 bit register in list
- 27-28. Set aside for use by the error print sections
 - 29. Counter word delayed one right shift
 - 30. The constant zero
 - 31. N(31) = -1/2 when reading right instructions = 1/2 when reading left instructions
- 32-562. Instructions of SADOI
- 563-604. Instructions of SADOI used only during input starts
- 563-607. Temporary storage for blocks of 45 instruction pairs
- 43-357. William's memory storage for additions list during back substitution
- 608-620. Temporary storage while 13 additions words are accumulated
- 621-944. List of symbolic addresses
- 945-946. Temporary storage for symbolic address word during search of the list
- 947-998. The 2 bit registers
- 999-1025. The part of SADOI put into the Williams memory by the drum bootstrap routines
 - 1001. Relativizer
- 1019-1025. Used only during input starts after clear or hold starts to restore 2-32 of the Williams memory. Otherwise
 - 1019. = 0 except during drum input. Used for counter
 - 1020. = 0 except during drum input. Store word being decoded
 - 1021. = 0 except during drum input. Store drum order
 - 1022. Used to indicate errors have been made. Also used to indicate special interludes

- 1023. Use to store accumulated sum during computation of sum check when reading back a library routine.
- 1002-1018. Is spoken of as the control block. These locations are rewritten at <u>each</u> input start.
 - 84. Right address = location for the current instruction pair
 - 86. Right address = location in 563-607 where next instruction pair will be stored temporarily
 - 180. Right address = one less than the location in 608-620 where next additions word will be stored
 - 190. Right address = one less than the location the next additions word will have in the Williams memory.
 - 999. Right address = location of last word in symbolic address list.

BOOKKEEPING

Each new instruction (except K and N terminated) is treated by SADOI as overwriting instructions already stored. The memory location may be zero as after clear starts. This case is also included. The 2 bit registers must therefore be modified continually during input in order to keep track of where substitutions of real addresses must be made.

Since a word is added to the instructions being accumulated on the drum only after complete instruction pairs are received modification of two positions in a 2 bit register is made simultaneously after the complete instruction pair is received.

After the instruction pair is received the previous two digits of the register are irrelevant. These positions are therefore set to zero. New 1 digits are then added for each instruction of the pair that will require a substitution later.

Accumulation of information for the bookkeeping is accomplished by use of five storage locations 22-26.

- 26. The location of the current 2 bit register
- 24. Temporary storage of the 2 bit register (to avoid many address substitutions).
- 25. The counter word. This word is as described earlier. During input of the function digits it is shifted right one shift. A copy of the unshifted word is stored at 29.

- 22. Before the right shift of N(25), the sum N(22) + N(25) is formed and stored at 22. After two instructions are read $N(24) - [N(24) \cdot N(22)]$ sets two positions of the 2 bit register = 0, where "." means logical and.
- 23. For each instruction which requires a substitution later, N(23) + N(29)is formed and stored. The total correction of the 2 bit register after an instruction pair has the form $N(24) - [N(24) \cdot N(22)] + N(23)$. See 69-73 When the right instruction of a pair is terminated by K or N

N(1) is not stored as part of the program being read. Modification of the 2 bit register does not take place. However, N(22) and N(23) are computed. Since $N(23) \neq 0$ means the use of a forward reference in the instruction pair the test $N(23) \neq 0$ results in the FF032 and FF033 stops.

During input starts, after K, and after instruction pairs are read, N(22) and N(23) are set to zero. N terminations with addresses \geq 999 get special treatment. It was decided that before executing a special jump SADOI should be reset so that the next instruction would be a left instruction and in general that the N terminated instruction was not read. Words 439-444 therefore

- (1) back up the counter word to its previous position,
- (2) destroy information about inside and outside addresses having been used: $N(20) = 0 \quad N(22) = 0 \quad N(23) = 0$
- (3) set for left-hand instruction: N(31) = 1/2.

The 2 bit register held at 24 is replaced in the list when K, S, N terminations are read or when N(25) = 0.

S TERMINATIONS

Use of S parameters is made almost unnecessary by the details of SADOI. They were included, however, to make SADOI compatible with the D.O.I. In certain cases it is convenient not only to use S parameters but to have a facility allowing use of inside addresses. Some provision is made for this in SADOI.

First, any number which may be used as an S parameter is stored simultareously in the Williams memory and on the drum. These are the words having locations 2, 3, ..., L. If the number is not used as an S parameter then it may have inside audresses on both sides which make forward references. Backward references never cause difficulty because the real address is substituted immediately. If it is used as a S parameter only the right instruction may use an inside address making a forward reference.

when an S termination is written SADOI checks the 2 bit register at 947 to determine whether a substitution is needed. The current 2 bit register at 25 is replaced in the list just to make sure the register at 947 is actually there. In case a substitution is needed, then the 2 copies of the S parameter must be changed. The instruction pairs (563-607) being held temporarily in the Williams memory are put on the drum to make sure the second copy of the S parameter is actually on the drum. SADOI then determines whether the required real address has been received. If so corrections are made. Otherwise, an error is indicated by the tape punch.

After special interludes which end with 999 or 1014 jumps to SADOI locations 2-15 are transferred to the corresponding drum locations regardless of the locations specified for the special interlude.

LOCATION	ORDER
	40 1012
1027	26 1012
1028	F5 607
1028	40 1015
1029	40 1016
1027	26 1012
1030	40 604
10,0	26 607
1031	00 0
10)1	F5 607
1032	40 605
,	26 607
1033	40 607
	L5 608
1034	40 606
	26 607
1035	L4 604
	46 608
1036	40 609
	26 607
1037	40 31
	22 604
1038	40 608
	26 607
1039	40 31
	L6 31
1040	00 0
31	00 0
1041	00 0
32	00 2

Words of bootstrap stored on drum at 1027-1039. These set up a loop at Williams memory locations 604-609 to read main body of SADOI into the Williams memory. Remainder of bootstrap is in Williams memory 1013-1016. Start at 1013.

NOTES

X 12

PAGE 1

First words of main part of SADOI.
Stored at 31, 32.

WILLIAMS MEMORY LOCATIONS ONLY ARE INDICATED. DRUM LOCATIONS ARE 1009 GREATER THAN WILLIAMS MEMORY LOCATIONS

LOCATION	ORI	DER		NOTES	PAGE 2 X 12
3 0	00 0	00 0			
1	00 0	00 0	Nur	merical Constants	
2	00 0	00 2			
3	00 0	00 10			
4	80 0	00 16			
5	00 0	00 20			
. 6	00 0	00 32			
7	00 0	00 40		•	•
8	00 0	00 999	•		
9	9 741701 282048				
	•				
40	7L4095	LL3072			
1	7L4095	LL2048			
2	00 0	00 132			
3	40 16	41 0		Read in Decimal	Address
4	49 21	L5 16	Se	t up	
5	LO 33	36 56			
6	L4 33	40 0			
7	91 4	32 48	Re	ad Loop	
8	26 53	LO 33			
9	36 56	10 3			
50	F4 0	00 2			
1	F 4 0	00 1			
2	40 0	26 47			
3	FO 32	40 16	Cl	neck 5-hole character:	s for
4	L7 16	36 47	st	cart of inside address	9
5	22 115	FF 53			
6	F4 56	42 57	Se	et switch	
7	L 5 0	26 0			
8	L5 70	22 74	K		
9	L5 251	22 74	S		

LOCATION	ION ORDER		NOTES PAGE 3 X 12
60	L5 158	22 74	NOTES PAGE 3 X 12
. 1	50 39	22 97	J
2	L4 1	26 64	F
3	L41001	L4 1	L and final operations of
4	40 1	Ll 31	J and S.
5	40 31	32 69	Left-right switch. Jump after right inst.
6	91 4	36 104	Search for (of
7 -	FO 32	40 16	outside address.
8	L7 16	36 66	
9	26 116	50 24	
			Bookkeeping
70	J 0 22	-1 203	Modify current 2-bit reg. after
1	L4 24	L4 23	inst. pair
2	40 24	L3 25	Test counter word for zero
3	36 74	22 81	
4 .	F5 74	42 77	Replace 2-bit register in list
5	L5 26	42 76	when counter = 0 or for K, S, N
6	L5 24	40 0	
7 =	50 30	26 0	
8	89 1	40 25	When counter word = 0, get next
9	F5 26	42 80	2-bit register and reset counter word.
d٥	12 2/	T# : 0	
80	42 26	L5 0	
2	40 24 41 23	41 22	
3	LO 91	L5 84 36 85	Store S parameters in Williams Memory.
4	L5 1	30 85 40 0	
5	F5 84	40 84	
6	L5 1	40 563	Stand of the standard of the s
7	41 20	F5 86	Store inst. pair temporarily at
8	40 86	LO 92	563 – 607
9	32 90	50 89	
•	J - / V)	

LOCATION	ORD	ER	NOTES PAGE 4	X 12
90	26 357	26 66		
1	L5 1	40 16		
2	75 1	40 608		
3	26 66	22 63	26 999N	
4	22 101	26 95	26 1000N	
5	L5 1	10 20	Set drum input	
6	L4 103	40 362		
7	26 93	LO 39	J termination	
8	66 39	10 1		
9	- J 0	22 63		
100	L3 31 ·	32 361	Sum check test	
1	FF2084	L5 84	Set special interlude address	
2	421022	26 93		-
3	86 11			
4	10 4	L5 25	Read second function digit;	
5	40 29	L4 22	Save old counter word and	
6	40 22	L5 1	advance counter word one	
7	00 4	80 4	shift.	
8	50 30	00 12		
9 .	40 1	L5 25		
110	10 1	32 111		
1	L4 25	40 25		
2	91 4	36 43	Search for (of	
3	FO 32	40 16	inside address.	
4	L7 16	36 112	Input Symbolic Address	alle files
5	41 0	41 21		
6	L1 42	40 18	Set for fig. shift	
7	41 17	51 17		
8	10 2	90 4	Core of input loop	
9	L4 42	42 120		
j.				
1				

LOCATION	ORD	ŒR	NOTES PAGE 5	X 12
120	L4 18	26 0	Jump to switch	
120	LO 17	26 118	Entry for N, J, F, L	
2	L4 34	26 118	Entry for "normal" 5-hole char.	
3	L5 36	40 17	Change of letters-figures shift;	
•	LO 42	40 18	entry when character is skipped.	
4 5	01 4	22 118		
,6	40 16	L3 17) or A	
7	36 148	L5 16		
8	26 122	40 16	zero or P	
9	L3 17	36 131		
,				
130	L5 16	26 118		
1	F5 37	26 118		
2	32 128	26 125	O delay P delay	
3	36 118	26 122	1 S Q D	
4	36 118	26 125	2 CR-LF W CR-LF	
5	36 118	26 122	3 (E B	
6	36 118	26 123	4 Ltr.Sh. R Ltr.Sh.	
7	36 118	26 122	5 , T V	•
8	36 118	26 126	6) Y A	
9	36 118	26 122	7 / U X	
140	36 118	26 125	8 dela y I dela y	
1	36 118	26 122	9 = 0 G	
2	36 118	26 122	+ . K M	
3	36 118	23 123	- Nr.sh. S Nr.Sh.	
4	36 121	26 122	N 'N H	
5	36 121	26 122	J : J C	
6	36 121	26 122	F x F Z	
7	36 121	26 125	L space L space	
8	- J 0	30 1	Symbolic Address List Comp	
9	40 946	43 946		

LOCAT SON	ORDI	ER		
150	89 1	36 151		
- P	L4 946	40 945		
2	L5 999	22 153		
3	F5 160	42 154		:
4	L1 946	L4 0		
5	36 159	L4 40		
6	36 159	42 17		
7	L5 154	42 160		
8	26 163	00 435		
9	F5 154	42 160		
	•			
160	Ll 946	L4 0		
1	36 153	L4 40		
2	36 153	42 17		
3	L3 17	36 167	•	
4	L3 21	32 172		
5	L5 160	42 20		
6	42 171	26 328		
7	L5 160	LO 198		
8	36 192	L3 21		
9	36 175	L5 160		
170	42 171	42 20		
1	L5 84	42 0		
2	26 66	L5 17		
3	L4 0	L4 1		
4	40 1	26 138		
5	L3 0	32 185		
6	F5 180	42 180		
7	50 30	L5 160		
8	42 17	L5 0		
9	00 10	L4 17		·

Search list of symbolic addresses.

Note double loop because of high read around ratio.

Test real add. = 0

Test inside-outside

Outside. Word in list has

non-zero address.

Test for new sym. Add.

Test inside-outside.

lst time outside. Substitute

real address.

Inside when real address is known.

Inside when real address is not known. Jump to 185 when no constant to be added.
Otherwise add new word to additions list.

LOCATION		ORDER		NOTES	Page 7 X 12
180	50 30	40 607	***		
1	L5 180	LO 199			
2	32 447	26 184			
3	L5 261	42 180			
4	F5 190	42 190	•		e e
5	26 186	L5 160		For inside, subst.	list address
6	42 1	L5 . 29		and add bit to thos	e to be added
7	L4 . 23	40 23		to 2 bit register.	
8	49 21	L1 31		Bookkeeping after i	nside add.
9	40 31	32 69			
190	26 66	00 42			
1	+6 66	00 357			
2	L5 9 99	F O 30		New symbolic addres	s. Decrease
3	42 999	42 160		list address by one	and add new
4	42 195	89 1		word to the list.	
5	L4 946	40 0			
6	Ll 999	L4 200		Test if list capaci	ty is exceeded.
7	32 204	22 168			
8	L1 946	L4 945			
9	50 30	40 620			
	•				
200	491002	- 5 620			
1	L5 151	42 999			
. 2	26 66	00 93		V Mormin	nation
3	L3 23	32 204		Test for error	UR CIOU
ŢŤ.	26 480	L5 1.			
5	10 16	40 1		Determine function	digit and set switch
6	01 4	L4 221			
7	42 208	50 207			
8	26 357	26 0			·
9	L5 0	22 214		00 K	

LOCATION	OR	DER	NOTES PAGE 8 X 12
210	L51001	22 212	Ol K
1	L5 1	10 4	02 K
2	421001	ľ ₇ O	
3	42 84	42 0	
4	22 217	401001	
5	L 3 0	32 216	
6	23 212	L5 84	
7	26 212	5 0 0	Compute address and position of bit
8	09 1	40 25	in counter word.
9	67 35	40 18	
220	F 0 30	42 224	
1	11 1	- 5 209	
2	I4 254	42 26	
3	42 225	L3 18	
74	32 225	19 0	
5	40 25	L5 0	
6	40 54	L5 367	Set drum order
7	L4 O	40 362	
8	49 31	49 21	
9	41 22	41 23	
		76 07	Test whether outside address calling
230	L3 20	36 93 42 232	for library routine. Set correct real
1	L5 20	42 0	address. Search reference word list on
2	151001 15 242	40 235	drum. End search on -1.
3 4	411023	41 20	
	85 11	41 20	
5 6	36 238	LO 32	
. 7	36 479	14 32	
8	LO 945	32 239	
9	221 240	<u>r.</u> 4 41	

LOCATION	OF	RDER	·	notes	PAGE 9	X 12
240	3 6 262	F5 235				
1	40 235	26 235				
2	85 11	001850				
3	81 5	42 247		S Termi	nation	- en qu tu
4	10 1	42 248		Read extra characte	r and set addre	esses.
5	42 260	L4 367				
. 6	40 259	49 16			•	
7	50 16	11 0				
8	JO 947	L5 0		Determine is inside	address used.	
9	42 251	40 16				
25 0	22 250	- 3 243		Is correction to be	made?	
1	36 260	L5 0	·			
2	42 17	42 16				
3	L3 17	36 341		Can correction be m	ade?	
4	L5 947	- 0 947		Change 2 bit regist	er	
5	40 947	L5 26				
6	42 257	50 256		Record words on dru	m.	
**************************************	26 357	L5 0		Replace 2 bit regis	ter and correct	•
8	40 24	L5 16		copies of S paramet	er.	
9	86 11					
					•	
260	L5 . 16	40 0				
1	L4 0	22 93		Modifie	d Y2	•=
2	LO 41	L4 242				
3	401021	491019		Store drum order an	d modify counter	er
. 4	L51022	40 19		Save info. on spec.	interludes and	errors.
5	411022	L5 296		Modify important ju	mp addresses.	
6	42 90	L5 325				
7	40 93	L5 92		Set end constant.		
8	42 379	26 269				
9	41 0	50 0	•	·		
						ı
					•	

LOCATION	OR	DER	NOTES PAGE 10 I 12
270	F5 270	42 285	Set addfn. dig. switch.
1	L51019 ·	00 1	Advance count. End of 5 blocks?
2	401019	32 283	
3	F5 279	40 279	End of 36 words from drum?
4	LO 327	32 279	
5	L5 294	42 279	Record words on drum.
6	12 375	50 276	
7	26 357	L51021	Bring next group of 36
8	40 374	50 278	words from the drum.
9	26 373	L5 608	Pick off one word. Adjust
280	401020	L61023	sum check, drum order, counter.
1	401023	F51021	
2	401021	19 4	
3	401019	L51020	
4	50 32	10 8	Split off 8 bits.
5	401020	22 0	Addfn. dig. switch.
6	LN O	01 2	Set termination switch.
7	F4 287	42 289	
8	L5 0	00 6	Add 6 bits to address at 0.
9	40 0	26 0	
290	26 271	00 607	≥ 64
1	L41001	40 0	L L
2	40 16	22 300	F
3	L5 0	10 6	K, S, N.
4	40 0	-3 572	Test for interlude
5	32 309	L4 286	Test for termination of drum input.
6	32 312	- 5 269	Test for K, S.
7	L4 286	32 299	
8	L5 16	40 0	K termination.
9	26 58	01 7	S termination.

LOCATION	OR	DER	NOTES PAGE 11 X 12
300	22 244	F5 300	Set switch for fn. digits.
1	22 270	L5 1	Read fn. digit block.
2	00 20	L4 O	Add to word and bookkeeping.
3	40 1	L5 25	
4	L4 22	40 22	
5	L5 25	10 1	
6	36 307	L4 25	
7	40 25	Ll 31	
8	40 31	32 69	
9	26 269	L5 1	N termination.
100 miles			
310	LO 16	40 1	
1	L5 16	40 0	
2	26 60	49 21	Test sumcheck. Reset constants,
3	L5 327	42 279	addresses. Terminate drum input.
4	L5 19	401022	
5	L51021	40 317	
6	F5 317	401021	
7	85 11		
8	F41023	401023	
9	L31023	32 320	
320	26 101	L5 324	
1	42 90	L5 326	
2	40 93	411019	
3	411020	411021	
4	49 31	26 66	
5	26 269	22 291	Constants
6	26 66	22 63	
7	+6 373	L5 608	Duink October Brown
8	41 27	L5 84	Print Outside Type Error
9	LO 17	42 27	·

LOCATION	OI	RDER	NOTES PAGE 12	v 10
330	L3 27	36 66	INGB 12	A 12
1	92 131	92 2		
2	92 963	L5 20		
3	42 382	L5 17		
4	42 27	50 334		
5	26 407	L5 84		
6	42 27	50 336		
7	26 407	F5 337		
8	26 382	89 1		
9	L41022	36 171		
340	401022	26 171		
1	92 131	92 259	Print S type error	
2	92 706	92 707		
3	92 963	41 27		
4	L5 253	42 382		
5	50 32	L5 248		
6	42 27	5 0 34 6	·	
7	26 407	L5 84		
8	42 27	50 348	•	
9	26 407	F5 349		
250				
350	26 382	L31019		
1	36 354	89 1		
2	L4 19	36 356		
3	40 19	26 356		
4	89 1	L41022		
5	3 6 3 56	401022		
6 ,	L5 0	22 93		
7	+5 0	42 366		
8	L5 368	42 361	Store words on the drum.	
9	L5 86	42 369	Closed subroutine. Set	

	•			
LOCATION	<u> </u>	ORDER		NOTES PAGE 13 X 12
360	L5 361	LO 369		initially to store 2 - 15
1	32 365	L5 2		on the drum.
2	86 11	023566		
3	F5 362	40 362		
4	F5 361	40 361	•	
5	22 360	L5 363		
6	42 86	22 563		Link.
7	86 11	023564		
8	00 0	00 563		
9	32 365	L5 16		
370	86 11	023 56 3		
1	36 374	22 0		Link. Bring words
2	00 0	42 375		from the drum. Closed
3 % i	+5 507	42 371		subroutine. Set initially
4	85 11	001617		to rewrite 1002 - 1018.
5	32 375	401002		
6	F5 374	40 374		
7	F5 375	40 375		
8	LO 379	26 371		
9	- 2 375	401019		
380	85 11	023564		
1	85 11	023563	· · · · · · · · · · · · · · · · · · ·	
2	42 388	L 5 0		Entry. Print symbolic address.
3	10 10	40 27		Closed subroutine. Address of word
4	41 28	19 3		to be printed is substituted in
5	40 21	26 390		382.
6	L5 21	36 389		
7	L5 28	32 388		
8	92 707	22 0		Link.
9	04 1	40 21		

LOCATION	OR	DER	NOTES PAGE 14 X
390	50 28	L5 27	
1	10 6	40 27	
2	- 5 0	00 2	
3	40 946	10 2	
4	43 946	- 3 0	
5	36 386	- J 0	
6	36 400	89 1	
7	L4 28	40 28	
8	32 399	92 259	
9	26 400	92 707	
400	51 946	01 4	· ·
1	-4 0	36 403	
2	50 405	22 403	•
3	5 0 389	00 3	
4	00 3	42 405	
5	02 1	92 0	
6	26 386	00 549	
7	+5 0	42 416	Entry. Convert binary
8	50 32	L5 27	address to decimal and print.
9	00 20	50 417	Closed subroutine. Address
410	- 0 0	36 410	to be printed is stored at 27.
1	-4 0	50 418	
2	- 0 0	36 412	
3	-4 0	46 27	·
4	00 12	L4 27	
5	00 8	82 12	
6	92 963	22 0	Link.
7	00 99	ro o	
8	00 9	LL O	
9	41 27	L5 532	Print inside type error.

LOCATION	OF	RDER	NOTES PAGE 15 X 12
420	42 434	L5 531	TAME IN A 12
1	LO 551	36 423	·
2	L5 533	26 424	
3	L5 535	10 20	
4	42 382	22 427	
5	41 27	L5 406	
6	42 434	L5 544	
7	42 382	92 131	
8	L5 381	LO 367	
9	LO 368	L4 16	
430	42 27	50 430	
1	26 407	F5 431	
2	26 382	89 1	
3	L41022	32 434	
4	401022	26 0	N Termination
5	L5 1	L4 O	remination
6	40 1	L5 0	Construct jump instruction.
7	LO 38	32 438	
8	26 445	L4 202	Jumps with address \geq 999.
9	42 1	41 22	
1 40	41 23	L1 31	
1	36 442	23 442	
2	L5 29	L4 29	
3	40 25	49 31	
4	41 20	22 1	
5	41 0	50 445	Record words on drum.
6	26 357	F5 456	Set link and record remainder
7	42 457	L5 180	of additions list on drum.
8		L5 370	
9	L4 458	r∪ 452	

LOCATION	ori	DER	V	NOTES	PAGE 16	X 12
450	40 453	40 370		Note that the order is	reversed	
1	L5 452	5 0 30		during the transfer.		
2	22 456	L5 0				
3	86 11				•	
4	F5 453	40 453				
5	£5 452	FO 30	•		·	
6	42 452	LO 458				
7	32 452	26 183				
8	22 456	L5 608				
9	L51022	42 17		Test for special interlu	ide.	
460	L3 17	32 476				
1	L5 17°	421008		Set addresses in control	block	
2	421003	L4 367		(999 - 1018) for special	interlude.	
3	401009	LO 478				
4	401002	L5 84				
5	421017	421018				
6	L31019	36 468		Test for tape-drum input	•	
7	L5 19	401022		•		
8	L3 23	36 470		Test for error in jump a	it	
9	L5 480	4 0 1		location 1.		
470	L5 190	50 32		Compact address at 190 w	rith	
1	10 1 0	- 5 0		relativizer and store at	1001.	
2	L41001	401001				
3	L5 523	42 452		Store 2 bit registers an	d symbolic	
4	L5 999	42 458	·	address list on drum.		
5	L51005	42 457				
6	22 448	L31019		Test for drum or regular	interlude.	
7	36 481	26 467				,
8	01 0			Back Substituti	on and	_
9	FF 48	FF 49		Regular interludes for t	ape input	_

LOCATION	OF	DER	NOTES PAGE 17 X 12
480	FF 50	FF 51	
1	L5 190	42 488	Transfer additions list from drum to
2	LO 191	32 483	Williams memory. This loop is set
3 ઙૢ૾ૢ૽,	FF 52	L5 381	initially to bring symbolic address
4	FO 488	L4 497	list, etc., from the drum during
5	40 487	L5 488	input starts.
6	5 0 30	22 491	
7	85 11	023563	
8	32 488	40 0	
. 9	F5 487	4 0 487	

490	L5 488	F O 30	
1	42 488	L 9 497	
2	36 487	26 588	See word 593. Changed during input starts.
3	42 458	L5 513	Store list of symbolic addresses
4	42 452	F5 491	on drum.
5	42 457	L5 478	
6	L4 381	26 449	
7	32 488	40 43	
8	L5 367	40 381	
9	L5 380	40 374	
500	L5 368	L4 35	Set links, etc., for transferring
1	42 369	42 379	blocks of 20 words to and
2	L5 373	46 361	from the drum.
, 3	46 369	L5 556	
4	42 371	22 509	
5	L5 381	L4 35	Advance addresses.
6	40 381	LO 478	
7	40 374	F5 509	Entry after transfer to drum.
8	40 509	LO 557	
9	36 558	Fl 947	Test 2 bit register for 0.

LOCATION	ORL	DER	NOTES PAGE 18	X 12
510	40 18	F 3 18		
1	36 505	L5 368	In case 2 bit reg. $\neq 0$,	
2	42 361	42 375	bring 20 words from the	
3	26 374	00 3/1/1	drum.	
4	40 16	L5 18	Adjustment after subst.	
5	80 2	F4 32	real addresses.	
6	40 18	F 3 18	Test for completion of subst.	
7	36 553	L5 18	Test for inside addresses requiring	,
8	36 524	80 1	substitution of real addresses. Both	
9	36 524	80 1	instructions of a pair tested each	
			time through loop.	
520	F4 32	40 18		
1	F5 16	40 16		
2	22 517	L5 368	Entry after bringing 20 words.	
3	22 521	00 998		
4	L5 16	42 525	Set up for address substitution.	
5	42 549	L5 0		
6	46 531	42 542		į .
7	46 535	42 544		
8	40 31	Fl 18	Substitution for left instruction.	
9	36 539	L5 531	Test is constant to be added?	
530	LO 551	32 534		•
1	L5 0	42 533	Set up if constant to be added.	
2	40 0	50 539		
3	43 0	L5 0		
4	22 535	41 0		
5	L5 0	42 17	Substitution of address in left	
6	00 10	L4 O	instruction.	
7	00 10	46 31		
8	L3 17	36 419	Test for error.	
9	Fl 18	80 1	Substitution for right instruction.	
			•	

LOCATION			ORDER		NOTES PAGE 19
540	36	549	41	0	Test is constant to be added to
2	L5	542	LO	552	right instruction?
Ž	32	544	L5	0	Set up if constant is to be added.
3	42	544	10	10	
4	40	0	L5	0	
5	42	17	43	31	Substitute address.
6	L5	17	L4	31	
7	L4	0	40	31	
8	L3	17	36	425	Test for error.
9	L5	31	40	0	
550	F5	16	26	514	
1	L5	512	42	533	Constants
2	32	544	L5	512	
3	L5	381	40	362	Record modified words on drum.
. 4.	·L4	35	40	381	
5	L5	16	42	369	
0	22	361	00	522	
7	36	558	Fl	999	Test for error in jump instruction
8	L51	022	36	560	at location 1. Begin interlude.
9	L5	479	40	1	
560	L3	23	361	.002	
1	L5	480	40	1	
2	261	002			Turnet Oliver
3	23	107	41	0	Input Starts
4	41	30	41	17	Entry at right of 563 from 366.
5	41	20	41	1	
6	L510	002	42	0	Save information on type of start.
7	40	16	50	567	Rewrite 1002 - 1018.
8	26	373	L3	16	Hold or clear start when N(16)=0
9	32 5	570	26	575	

X 12

LOCATION	OR	DER	NOTES PAGE 20 X 12
570	42 574	41 947	Set 2 bit registers = 0
1	F5 570	40 570	
2	LO 601	32 570	
3	411019	411021	Set locations for special interludes,
4	411022	22 588	errors and drum input.
5	501001	01 10	In case of interludes decompact 1001.
6	40 17	11 10	Re-establish relativizer.
7	- 5 0	401001	
8	L3 17	32 597	N(17)= 0 means regular interlude.
9	F5 579	42 594	For special and drum input interludes
580	L5 17	42 190	replace constant at 190. Re-establish
1	42 488	L5 487	drum orders for additions list and
2	L4 497	FO 488	bring back symbolic address list and 2 bit
3	40 487	L4 478	registers.
4	40 370	L5 523	Transfer lists to Williams memory.
5	42 488	L5 999	
6	42 497	L5 487	
7	26 484		
8	41 1	L5 600	Enter K termination loop to set up
9	42 231	L5 16	constants for SADOI
590	36 210	26 214	
1	L5 202	42 231	
2	L5 112	42 497	Reset some orders and start input.
3	L5 602	40 492	
4	50 30	25 563	
5	151021	32 596	Distinguish between special and
Ė	25 2m2	4,31.027	drum input interludes.
7	2f 93	F5 597	Segular interlades
8	26 570	L5 202	
9	42 594	L5 513	

LOCATION		OR	ŒR		NOTES PAGE 21	X
600	26	585	00	591		
1	N2	574	41	999	Constants	
2	36	487	L5	999		
3	00	0	00	0	sum check constant.	
4	80	0	26	100	Overwrite bootstrap to start SADOI	
1614	49	1002			999 type interludes	
999	S5	945				
1615	42	1002				*
1000	22	1008				
1616	00	0			Relativizer	
1001	00	0				
1617	85	11				
1002	02	3566			Bring assembled	
1618	32	1003				
1003	40	2			program to	
1619	L 6	0				
1004	40	0			Williams memory.	
1620	F 5	1002				
1005	40	1002				
1621	F5	1003				
1006	40	1003				
1622	LO	1018				
1007	36	1002				
1623	22	1				
.1008	L5	16			Put contents of	
1624	86	11				
1009	02	3580°			Williams memory	
1625	F5	1009		•		
1010	40	1009			on drum.	
1626	F5	1008				
1011	40	1008				

	ORDER	NOTES	PAGE 22	2
LOCATION				
1627	LO 1017			
1012		Bootstrap start		
1628	F5 1015			
1013	40 607	Waste		
1629	26 1015	1014 interludes		
1014	26 1000			
1630	85 11	for main part of		
1015	00 1027			
1631	40 608	SADOI		
1016	26 607			
1632	K2 1	constants.		
1017	L5 999	33.133.133	.*	
1633	S2 1003			
1018	40 999	2 - 31		
1634	05 11	Restore		
1019	03 512	1653027).	
1635	85 11	after SADOI		
1020	03 482	arter bibor	*	
1636	32 1021	is first	;	
1021	40 2	12 11130	•	
1637	F5 1021	called for. Used		
1022	40 1021	Called 101.	c	
1638	F5 1020	as storage locations	4	
1023	40 1020	as stolage rooming		
1639	LO 1019	by SADOI.		
1024	36 1020	by Sador.		
1640	41 1002		ţ	
1025	22 1008	Ch -ale	,	
1641	8S 1759	Sum check		
	13 2340			
		•		