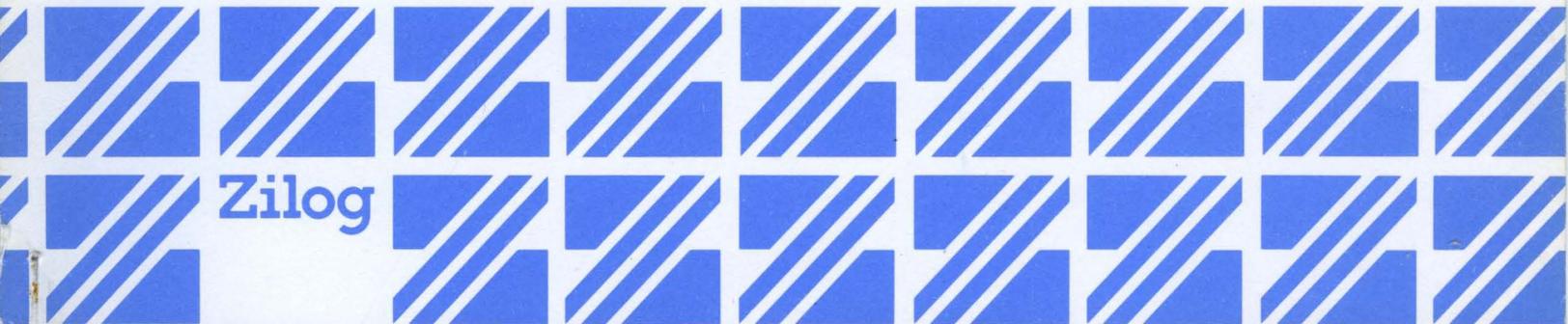




Z80TM-RIO

Relocating Assembler and Linker
User's Manual



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**Z80-RIO
RELOCATING ASSEMBLER AND LINKER
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PREFACE

This manual describes the relocating assembler (ASM) and linker (LINK) programs which support the translation of Z80 assembly language source code into executable machine code on the Zilog RIO operating system.

This manual contains two chapters -- Chapter 1 covers assembly language conventions, pseudo-ops, and how to invoke the Assembler with different command options. It does not describe the Z80 instruction set or mnemonics, since this information can be found in the publication

Z80-Assembly Language Programming Manual

Chapter 2 covers the functional characteristics of the Linker and how to invoke it with different command options.

RELATED PUBLICATIONS

RIO Operating System User's Manual
Z80-CPU Technical Manual

OVERVIEW

The relocating assembler ASM accepts a source file (a symbolic representation of a portion of a program in Z80 assembly language) and translates it into an object module. It may also produce a listing file which contains the source, assembled code, and optional symbol cross-reference. The features of the relocating assembler include macros and conditional assembly, as well as the capability to produce either relocatable or absolute object code.

All object code must be processed by the linker LINK before it is ready for execution. The linker accepts several object modules and combines them into a single program in the form of an executable procedure file. It may also produce a load map describing the assigned addresses and lengths of modules. An optional symbol file may be created for use by a symbolic debug package.

Relocation

Relocation refers to the ability to bind a program module and its data to a particular memory area at a time after the assembly process. The output of the assembler is an object module which contains enough information to allow the linker to assign that module a memory area. Since many modules may be loaded together to form a complete program, a need for inter-module communication arises. For example, one module may contain a call to a routine which was assembled as part of another module and may be located in some arbitrary part of memory. Therefore, the assembler must provide information in the object module which allows the linker to link inter-module references.

There are several major advantages to using the relocating assembler as compared to the absolute assembler:

- 1) Assignment of modules to memory areas can be handled by the linker rather than requiring the programmer to assign fixed absolute locations via the "ORG" pseudo-op, and thus modules can be relocated without requiring re-assembly.
- 2) If errors are found in one module, only that one module needs to be re-assembled and then re-linked with the other modules, thus increasing software productivity.
- 3) Programs can be structured into independent modules, coded separately and assembled, even though the other modules may not yet exist. This separation of

functionality also enhances program maintenance and verification.

- 4) Libraries of commonly used modules can be built and then linked with programs without requiring their re-assembly.
- 5) Communications between overlay segments can be achieved through methods similar to normal (non-overlay) inter-module references.

Relocatability is specified on a per module basis; in other words, a program may contain some relocatable modules and some absolute modules, but a single module is either entirely absolute or entirely relocatable. The notion of "absolute" is that there is a one-to-one direct correspondence between the assigned address of an instruction within a module during assembly, and the memory address of that instruction during program execution. Thus an absolute module whose origin is "absolute 0" will have its first instruction located at memory location 0. Absolute modules usually contain ORG pseudo-ops to set the origin of sections of code, since the default origin is 0 (see ORG pseudo-op).

Modules are by default assembled as relocatable modules. The notion of "relocatable" is that the assigned address of an instruction within a module during assembly, is a relative offset from an origin which is assigned an absolute address by the Linker. During program execution, the instruction will be located at the memory location specified by the assigned origin plus the relative offset. Thus, a relocatable module whose first instruction (which is located at "relocatable 0") will have that instruction located at the memory location which is the assigned origin of the module as determined by the Linker.

To achieve relocation in the Z80, instructions which reference memory addresses (such as LD, CALL, JP) and data values which serve as pointers to memory locations (such as PTR: DEFW DATA where DATA is a label in the program module) must be altered at linkage time to reflect the assigned absolute addresses. However, as far as the programmer is concerned, instructions and addressing remain the same as in the absolute assembler.

Inter-module Communication

To allow instructions to refer to "names" (either data values or entry points) in other separately assembled modules, the assembler provides two pseudo-ops, GLOBAL and EXTERNAL. The syntax is the pseudo-op, followed by a list of names.

The meaning of GLOBAL is that each of the listed names is defined in this module, and the name is made available to other modules which contain an EXTERNAL declaration for that name. The complementary form then is EXTERNAL, which declares each of the listed names to be defined in some other module. Notice that the GLOBAL name may be in either an absolute module or a relocatable module. A portion of the object module contains a list of both the GLOBALs which are defined in the module, as well as the EXTERNALs which the module references. One function of the Linker is to match all the EXTERNALs with the appropriate GLOBALs so that every instruction will reference the correct address during program execution.

Programming Hints

A general approach to designing a program, then, would be to divide it into manageable portions, each of which constitutes a module. Most modules can be relocatable, that is, it does not matter where the module is placed in memory as long as it is correctly referenced by other modules through the use of GLOBAL and EXTERNAL names. Parts of the program may be coded, assembled, linked and tested. Then other modules may be added in a step-wise fashion without re-assembly of already-tested modules, thereby increasing programming productivity.

For those applications which will ultimately be implemented in a system configuration where the program code is in ROM and the data area is in RAM, it is a good idea to segregate the data into one or two modules with GLOBAL declarations for each data address. During the program development stage, these data modules might be relocated anywhere, and then assigned a particular origin depending on the memory configuration only when the program is finished.

1.0 ASSEMBLER

1.1 THE ASSEMBLY LANGUAGE

The assembly language of the Z80 is designed to minimize the number of different opcodes corresponding to the set of basic machine operations and to provide for a consistent description of instruction operands. The nomenclature has been defined with special emphasis on mnemonic value and readability. For example, the movement of data is indicated primarily by a single opcode, LD (standing for Load), regardless of whether the movement is between different registers or between registers and memory.

The first operand of an LD instruction is the destination of the operation, and the second operand is the source of the operation. For example,

```
LD A,B
```

indicates that the contents of the second operand, register B, are to be transferred to the first operand, register A. Similarly,

```
LD C,3FH
```

indicates that the constant 3FH is to be loaded into the register C. Enclosing an operand wholly in parentheses indicates a memory location addressed by the contents of the parentheses. For example,

```
LD HL,(1200)
```

indicates the contents of memory locations 1200 and 1201 are to be loaded into the 16-bit register pair HL. Similarly,

```
LD (IX+6),C
```

indicates the contents of the register C are to be stored into the memory location addressed by the current value of the 16-bit index register IX plus 6.

The regular formation of assembly instructions minimizes the number of mnemonics and format rules that the user must learn and manipulate. Additionally, the resulting programs are easier to interpret, which in turn reduces programming errors and improves software maintenance.

1.2 ASSEMBLY LANGUAGE CONVENTIONS

An assembly language source program consists of a sequence of statements in the order which defines the user's program. Each statement corresponds to a line of text of up to 128 characters ended by a carriage return. Each statement may contain zero or more of the following four fields which are identified primarily by their order in the statement line and by specific delimiter characters as indicated below.

LABEL	OPCODE	OPERAND	COMMENT
-------	--------	---------	---------

For example, here is a statement which contains all four fields:

LOOP:	LD	HL,VALUE	;GET THE VALUE
-------	----	----------	----------------

a. DELIMITERS - Opcodes, operands, and labels must be separated from each other by one or more commas, spaces or tabs. Notice that anywhere a comma could be used, a space or tab will be equivalent, and vice-versa.

b. COMMENTS - Comments are not a functional part of an assembly language source program, but instead are used for program documentation to add clarity, and to facilitate software maintenance. A comment is defined as any string following a semicolon in a line, and is ignored by the assembler. Comments can begin in any column, so a statement might consist of only a comment.

c. LABELS - A label is a symbol representing up to 16 bits of information and is used to specify an address or data. If the programmer attempts to use a symbol that has been defined to have a value greater than 8 bits for an 8-bit data constant, the assembler will generate an error message.

A label is composed of a string of one or more characters, of which the first six characters must be unique. For example, the labels LONGNAME and LONGNAMEALSO will be considered to be the same label. The first character must be alphabetic and any following characters must be either alphanumeric (A...Z or 0...9), the question-mark character (?) or the underbar character (_). Any other characters within a label will cause an error message. A label can start in any column if immediately followed by a colon. It does not require a colon if started in column one.

Labels are entered into a symbol table along with the value assigned by the assembler. The symbol table normally resides in RAM, but automatically overflows to a disk file if necessary, so there is virtually no limit to the number of labels that can occur in a single program.

Labels are normally assigned the value of the reference counter for the statement in which they occur. The reference counter corresponds to the CPU program counter and is used to assign and calculate machine-language addresses within an object module. The symbol \$ is used to represent the value of the reference counter of the first byte of the current instruction, and can be used in general expressions (see Expressions).

d. OPCODES AND OPERANDS

Z80 machine instructions are represented by specific mnemonics, usually consisting of a single opcode (such as LD, CALL, etc.) and zero or more operands (such as a register designator like HL or a condition code like Z). See the Appendix for a complete listing of all opcode/operand combinations.

The Assembler also recognizes certain pseudo-ops (sometimes called directives) which occur in the opcode/operand positions (see Pseudo-ops). Certain opcodes allow a particular operand to be the value of an arbitrary expression using arithmetic and logical operators (such as addition or multiplication) and operands (such as labels, numbers or character values).

Numbers

The Assembler will accept numbers in the following different bases: including binary, octal, decimal and hexadecimal. Numbers must always start with a digit (leading zeros are sufficient), and may be followed immediately by a single letter which signifies the base of the numbers ('B' for binary, 'O' or 'Q' for octal, 'D' for decimal, and 'H' for hexadecimal). If no base is specified, decimal is assumed. For example, here is the same number represented in each of the four bases:

1011100B, 134Q, 1340, 92, 92D, 05CH

Character Values

The Assembler will accept as operands ASCII characters bounded by a single quote mark with the corresponding ASCII value (e.g., 'A' has value 41H). Because all operands are evaluated as 16-bit quantities, strings of characters longer than two will be truncated in value to the first two characters, so that 'AB' and 'ABCDE' will have the same value. There is a method to define data storage which is initialized to a string of up to 63 characters (see DEFM, DEFT pseudo-ops). The quote character itself is represented by two successive single quote marks within a string (e.g., LD A, ''' causes the A register to be loaded with the value 27H, which is the ASCII value for the single quote character). A special string called the null string (which has the value of 0) can be represented by just two successive single quotes as in ''.

Upper/Lower Case

The Assembler processes source text which contains both upper and lower case alphabetic characters in the following manner. All opcodes and keywords, such as register names or condition codes, must be either all capitals or all lower case. Label names may consist of any permutation of upper and lower case; however, two names which differ in case will be treated as two different names. Thus LABEL, label, and LaBeL will be considered as three different names. Notice that one could use a mixture of case to allow definition of labels or macros which look similar to opcodes, such as Push, or LdiR, without redefining the meaning of the opcode. All assembler commands, such as *List or *Include (see Assembler Commands), can be in any combination of upper or lower case, as can the arithmetic operators such as .NOT., .AND., or .EQ., and numbers can be any mixture of case, such as Offffh, 0AbCdH, or 011001b.

1.3 EXPRESSIONS

The Assembler will accept a wide range of expressions involving arithmetic and logical operations. All expressions are evaluated left to right, except that unary operators are performed first, exponentiation next, multiplication, division, modulo, and shifts next, followed by addition and subtraction, then logical operations and comparisons. The following table is arranged in a hierarchy of descending precedence.

Table of Operators

OPERATOR	FUNCTION	PRECEDENCE
+	UNARY PLUS	1
-	UNARY MINUS	1
.NOT. or \	LOGICAL NOT	1
.RES.	RESULT	1
**	EXPONENTIATION	2
*	MULTIPLICATON	3
/	DIVISION	3
.MOD.	MODULO	3
.SHR.	LOGICAL SHIFT RIGHT	3
.SHL.	LOGICAL SHIFT LEFT	3
+	ADDITION	4
-	SUBTRACTION	4
.AND. or &	LOGICAL AND	5
.OR. or ^	LOGICAL OR	6
.XOR.	LOGICAL XOR	6
.EQ. or =	EQUALS	7
.GT. or >	GREATER THAN	7
.LT. or <	LESS THAN	7
.UGT.	UNSIGNED GREATER THAN	7
.ULT.	UNSIGNED LESS THAN	7

Parentheses can be used to ensure correct expression evaluation. Note, however, that enclosing an expression wholly in parentheses indicates a reference to a memory address.

Delimiters such as spaces or commas are not allowed within an expression since they serve to separate the expression from other portions of the statement.

16-bit signed integer arithmetic is used throughout.

Notice that anywhere a number or character value can be used in an expression, a symbol (label) could be used as well. The value of the symbol is always a 16-bit quantity.

The negative of an expression can be formed by a preceding minus sign "-".

For example:

LD A,-2 is equivalent to LD A,0FEH,
since 0FEH is the 8-bit two's complement
representation of negative 2.

The five comparison operators (.EQ., .GT., .LT., .UGT., and .ULT.) evaluate to a logical True (all ones) if the comparison is true and a logical False (zero) otherwise. The operators .GT. and .LT. deal with signed numbers whereas .UGT. and .ULT. assume unsigned arguments.

For example:

1.EQ.2 is equivalent to 0,
(2+2).LT.5 is equivalent to OFFFH
and '0'=0 is equivalent to 0

The Result operator (.RES.) causes overflow to be suppressed during evaluation of its argument, thus overflow is not flagged with an error message.

For example:

LD BC,7FFFH+1 would cause an error message,
whereas LD BC,.RES.(7FFFH+1) would not.

The Modulo operator (.MOD.) is defined as:

$X.MOD.Y = X - Y * (X/Y)$
where the division (X/Y) is integer division.

For example:

8.MOD.3 is 2 and 0ABH.MOD.0BH is 6

The Shift operators (.SHR.,.SHL.) shift their first argument right or left by the number of positions given in their second argument. Zeros are shifted into the high-order or low-order bits, respectively.

Mode of Expressions

Arithmetic expressions have a mode associated with them: absolute, relocatable or external as defined below. A symbol in an expression is absolute if its value is a constant (either a number, character value, or a symbol which has been EQUated to a constant expression). A symbol in an expression is relocatable if its value is subject to relocation after assembly (either a label or the reference counter symbol '\$'). Simple relocation means that only a

single relocation factor needs to be added to the relocatable value at load time. GLOBAL names within a relocatable module have a relocatable mode. GLOBAL names within an absolute module have an absolute mode. A symbol is external if it is declared in an EXTERNAL pseudo-op.

In the following, A stands for an absolute symbol or expression, R stands for a relocatable symbol or expression, and X stands for an external symbol or expression. <operator> is one of the standard arithmetic symbols such as +, -, *, .OR., .AND., etc. <comparison> is either .EQ., .LT., .GT., .UGT., .ULT. -- the resulting value is either True (all ones) or False (all zeros).

A is defined as either:

- 1) absolute symbol
- 2) A <operator> A
- 3) +A or -A or .NOT.A or .RES.A (unary operators)
- 4) R - R
- 5) A <comparison> A
- 6) R <comparison> R

R is defined as either:

- 1) relocatable symbol
- 2) R + A
- 3) R - A
- 4) A + R
- 5) +R (unary plus)
- 6) .RES.R (ignores overflow in a relocatable expression)

X is defined as either:

- 1) external symbol
- 2) X + A
- 3) X - A
- 4) A + X
- 5) +X (unary plus)
- 6) .RES.X (ignores overflow in an external expression)

Invalid expressions include:

- | | |
|-------|---|
| R + R | ; this is not simple relocatable since the relocation factor would have to be added twice |
| R * A | |
| A - R | |
| X + R | ; cannot add external and relocatable |
| R = A | ; cannot compare absolute with relocatable |
| X = X | ; cannot compare externals |

Relative Addressing

In specifying relative addressing with either the JR (Jump Relative) or DJNZ (Decrement and Jump if Not Zero) instructions, the Assembler automatically subtracts the value of the next instruction's reference counter from the value given in the operand field to form the relative address for the jump instruction. For example:

```
JR      C,LOOP
```

will jump to the instruction labeled LOOP if the Carry flag is set. The limits on the range of a relative address is 128 bytes in either direction from the reference counter of the next instruction (notice that JR and DJNZ are two-byte instructions). Thus LOOP must be a label within the range $\langle -126, +129 \rangle$ from the JR instruction. An error message will be generated if this range is exceeded.

Notice that the mode of a label and the reference counter must match so that the result of the implied subtraction is absolute (both must be relocatable or both must be absolute). If not, an error message will be generated.

An expression which evaluates to a displacement in the range $\langle -126, +129 \rangle$ can be added to the reference counter to form a relative address. For example:

```
JR      C,$+5
```

will jump to the instruction which is 5 bytes beyond the current instruction.

1.4 LISTING FORMAT

The Assembler produces a paginated listing of the source program along with the generated object code, plus an optional cross-reference listing of all symbols. A description of the various fields in the listing format follows. Refer also to the sample listing.

Heading	Each page has a heading which contains the name of the source file, a page number, column headings as explained below, and the Assembler version number. In addition, the heading may contain a user-specified heading (see *Heading Assembler command) and a user-specified string which is usually the date of assembly (see Date option).
LOC	Location: This column contains the value of the reference counter for statements which generate object code. It is blank otherwise.
OBJ CODE	Object code: This column contains the value of generated object code, up to a maximum of four bytes of code. It is blank if a statement does not generate object code.
M	Mode: This single character column indicates the mode of an instruction. An 'R' indicates the instruction contains a reference to a relocatable address, an 'X' indicates the instruction contains a reference to an external address, and a blank indicates the instruction is absolute. Notice that an instruction which contains either a relocatable or external reference will be modified by the Linker; therefore, the value in the object code listing does not reflect the true value of the instruction during program execution.
STMT	Statement number: This column contains the sequence number of each source statement.

SOURCE STATEMENT

The remainder of the line contains the source statement text. Listing lines longer than 132 characters will be truncated.

Cross-reference

If the cross-reference option is specified, then an alphabetical list of all the symbols in the source program, along with the statement number of each line which contains a reference to that symbol, will be appended to the listing file. The format is a symbol, followed by its 16-bit value, a mode indicator, the statement line on which the symbol was defined, and a list of the statement lines which refer to the symbol.

The mode of a symbol follows its value, with blank=absolute, 'R'=relocatable, 'G'=GLOBAL, and 'X'=external. Names which are declared EXTERNAL are considered to be defined by the EXTERNAL pseudo-op statement. Names which are declared GLOBAL are considered to be defined where they appear as a label, and considered to be referenced by the GLOBAL statement.

LOC	OBJ CODE	M	STMT	SOURCE	STATEMENT	multiply	6/15/77	PAGE 1
			1		GLOBAL ENTRY,MLTPLY			ASM 4.0
			2		EXTERNAL OUTNUM			
			3					
			4		; TEST OF MULTIPLY ROUTINE			
0000	ED5B0D00	R	5	ENTRY:	LD DE,(DATA)			
0004	0E02		6		LD C,TWO			
0006	CD0F00	R	7		CALL MLTPLY			;DO MULTIPLICATION
0009	CD0000	X	8		CALL OUTNUM			;OUTPUT VALUE IN HL
000C	76		9		HALT			
			10					
000D	AA00		11	DATA:	DEFW OAAH			
			12					
			13	TWO	EQU 2			
			14					
			15		;8-BIT BINARY MULTIPLY ROUTINE			
			16		; MULTIPLIER IN C			
			17		; MULTIPLICAND IN DE			
			18		; PARTIAL SUM IN HL			
000F	210000		19	MLTPLY:	LD HL,0			
0012	0608		20		LD B,8			
			21					
0014	CB39		22	MLOOP:	SRL C			
0016	3001		23		JR NC,NOADD			
0018	19		24		ADD HL,DE			
			25					
0019	CB23		26	NOADD:	SLA E			
001B	CB12		27		RL D			
001D	10F5		28		DJNZ MLOOP			
			29					
001F	C9		30		RET			
			31		END			

CROSS REFERENCE

SYMBOL	VAL	M	DEFN	REFS
DATA	000D	R	11	5
ENTRY	0000	G	5	1
MLOOP	0014	R	22	28
MLTPLY	000F	G	19	1
NOADD	0019	R	26	23
OUTNUM	0000	X	2	8
TWO	0002		13	6

multiply

6/15/77

PAGE 2

1.5 PSEUDO-OPS

In addition to normal opcodes which generate executable machine instructions, the Assembler recognizes several pseudo-ops which are used to control the generation of object code. Pseudo-ops have the same basic format as instructions, with the pseudo-op located in the opcode field and any operands following. Labels and comments generally may be used subject to the restrictions listed with the individual pseudo-op descriptions.

Data Definition

DEFB -- Define Byte

The DEFB pseudo-op takes one operand and generates a single byte of object code located at the current value of the reference counter.

Label	Opcode	Operand
optional:	DEFB	expression

If a label is present, it will be assigned the current value of the reference counter, and thus can be used to refer to the data value. The expression must evaluate to a quantity which can be represented in a single byte (8 bits). An error message will be generated if the value is not in the range -128 to 255. Because the Assembler treats all relocatable and external expressions as 16-bit addresses, the mode of an expression used in a DEFB must be absolute, since a 16-bit value will generally not fit into a single byte.

Examples:

Object Code	Label	Opcode	Operand
BD	TRUE:	DEFB	OBDH
04		DEFB	2+2
30	CHAR:	DEFB	'0'

DEFW -- Define Word

The DEFW pseudo-op takes one operand and generates a single word (16 bits or 2 bytes) of object code located at the current value of the reference counter.

Label	Opcode	Operand
optional:	DEFW	expression

The least significant byte is located at the current value of the reference counter, while the most significant byte is located at the next higher location. The order for storing low and then high order bytes is the format for addresses stored in memory; thus the DEFW can be used to store values which "point" to a particular memory address. The mode of the expression can be either absolute, relocatable or external. If a label is present, it will be assigned the current value of the reference counter, and thus can be used to refer to the data value.

The expression may consist of a character value, but notice that only the first two characters in a string are used. The object code will contain the characters in the same order as they appear in the string, not in reverse order.

Examples:

Suppose that PROC1 is a label on an instruction which was located at the reference counter value of 40BDH.

Object Code	Label	Opcode	Operand
BD40	PTR:	DEFW	PROC1
0400		DEFW	2+2
3031	STRNG:	DEFW	'01'

DEFM -- Define Message

The DEFM pseudo-op takes one string operand and generates a sequence of bytes in the object code which represents the 7-bit ASCII code for each character in the string.

Label	Opcode	Operand
optional:	DEFM	string

A string is represented by a sequence of characters bounded by single quote marks. The number of characters must be in the range 0 to 63. The single quote character itself is represented by two successive single quote marks within the string. If a label is present, it will be assigned the current value of the reference counter, and thus can be used to refer to the first character in the string.

Notice that only the length byte and the first three characters of a long string will be listed in the object code field in the listing.

Examples:

Object Code	Label	Opcode	Operand
48454C4C	MSG:	DEFM	'HELLO THERE'
612762		DEFM	'a''b'

DEFT -- Define Text

The DEFT pseudo-op takes one string operand and generates a sequence of bytes in the object code which represents the 7-bit ASCII code for each character in the string. In addition, it inserts a byte which contains the length of the string as the first byte in the object code.

Label	Opcode	Operand
optional:	DEFT	string

A string is represented by a sequence of characters bounded by single quote marks. The number of characters must be in the range of 0 to 63. The single quote character itself is represented by two successive single quote marks within the string. If a label is present, it will be assigned the current value of the reference counter and thus can be used to refer to the length byte of the string.

Notice that only the length byte and the first three characters of a long string will be listed in the object code field in the listing.

Examples:

Object Code	Label	Opcode	Operand
0B48454C	MSG:	DEFT	'HELLO THERE'
03612762		DEFT	'a''b'

Storage Definition

DEFS -- Define Storage

The DEFS pseudo-op takes one operand which specifies a number of bytes to be reserved for data storage starting at the current value of the reference counter.

Label	Opcode	Operand
optional:	DEFS	expression

If a label is present, it will be assigned the current value of the reference counter, and thus can be used to refer to the first byte of data storage. The expression can evaluate to any 16-bit quantity; however, the mode must be absolute. Any symbol appearing in the expression must have been defined before the Assembler encounters the expression.

The DEFS pseudo-op reserves storage by incrementing the reference counter by the value of the expression. If the result exceeds a 16-bit quantity, an error is generated and the reference counter is not changed. Since no object code is generated into storage area, the contents of the storage during initial program execution are unpredictable.

Examples:

Notice that the symbol LINELN in the following example must have been previously defined.

Label	Opcode	Operand
BUFFER:	DEFS	128
	DEFS	LINELN+1

Source Termination

END -- End of Source

The END pseudo-op signifies the end of the source program, and thus any subsequent text will be ignored.

Label	Opcode	Operand
optional:	END	

If a label is present, it will be assigned the current value of the reference counter. Operands are ignored. If a source program does not contain an END pseudo-op, then the end-of-file mark in the last source file in the assembler command line will signify the end of the program.

Symbol Definition

Labels on instructions are automatically assigned the current value of the reference counter. The pseudo-ops EQU and DEFL can be used to assign arbitrary values to symbols. In addition, to facilitate inter-module communication, certain symbols can be declared to be either GLOBAL or EXTERNAL to a particular module.

EQU -- Equate

The EQU pseudo-op is used to assign the value of an expression to the symbol in the label field.

Label	Opcode	Operand
name:	EQU	expression

The name in the label field is required and must follow the rules for forming a label (remember the colon is optional if the name starts in column one). The name must not be redefined as a label anywhere else in the source program, or a multiple definition error will occur.

The value of the expression can be any 16-bit value, but any symbol appearing in the expression must have been defined before the Assembler encounters the expression. The symbol being EQUated inherits the mode of the expression, that is, relocatable or absolute. The mode of the expression must not be external. A name which is EQUated to a GLOBAL name in a module will be relocatable if the GLOBAL name is relocatable, otherwise it is absolute; however, the name will not be accessible to other separately assembled modules.

Examples:

Suppose the symbol TWO has the value 2.

Label	Opcode	Operand	
ASCICR	EQU	ODH	; ASCII carriage return
FOUR	EQU	TWO+2	
SEMICO	EQU	';'	; ASCII semicolon

DEFL -- Define Label

The DEFL pseudo-op is used to assign the value of an expression to the symbol in the label field. This symbol may be redefined by subsequent DEFL pseudo-ops.

Label	Opcode	Operand
name:	DEFL	expression

The name in the label field is required and must follow the rules for forming a label (remember the colon is optional if the name starts in column one). The DEFL pseudo-op is identical to the EQU pseudo-op except that "name" may appear in multiple DEFL pseudo-ops in the same program. Notice that both the value and the mode can be changed by each occurrence of the DEFL.

Example:

Object Code	Label	Opcode	Operand
	FLAG:	DEFL	0
3E00		LD	A,FLAG
	FLAG:	DEFL	OFFH
3EFF		LD	A,FLAG

Programming Hints

Using symbolic names instead of numbers for constant values enhances the readability of a program and tends to make the code self-documenting. For instance, the symbol BUFLLEN is a more descriptive name for a value than, say, 72. Furthermore, by using EQUates, if a value which is used throughout a program needs to be changed, only the EQU statement needs to be modified rather than finding all occurrences of the number 72.

It is generally preferable to use EQU for symbol definition, since the Assembler will generate error messages for multiply-defined symbols which may indicate spelling errors or some other oversight by the user. DEFL should be reserved for special cases where it is necessary to reuse the same symbol, for instance, in conjunction with conditional assembly.

A common use of EQUates is to symbolically represent offsets in a data structure. The following example demonstrates how EQUates might be used to reference certain fields in an I/O parameter vector as offsets from the IY index register:

```

IOLUN   EQU   0           ; logical unit designator
IOREQ   EQU   1           ; request code
IODTA   EQU   2           ; data transfer address
IODL    EQU   4           ; data length
OPEN    EQU   04H        ; open I/O device request code

LD (IY+IOREQ),OPEN      ; open the device
CALL SYSTEM

```

GLOBAL -- Global Symbol Declaration

The GLOBAL pseudo-op is used to declare that each of its operands are symbols defined in the module, and the name and value are made available to other modules which contain an EXTERNAL declaration for that name.

Label	Opcode	Operands
	GLOBAL	name1, name2,...

There must be no label. There may be zero or more names which are separated by delimiters and refer to labels which are defined elsewhere in the module. GLOBAL pseudo-ops may occur anywhere within the source text. Notice that the mode of a GLOBAL depends on how the label is defined, which, for labels on instructions or data values, depends on whether the module is relocatable or absolute.

A name may not be declared as both GLOBAL and EXTERNAL. If the name is never actually defined in the module (i.e., there is no corresponding label definition), an error message will be generated at the end of the listing.

Examples:

Label	Opcode	Operand
	GLOBAL	ENTRY
	GLOBAL	FEE,FI,FO,FUM

*Implementation Note

In order to distinguish between GLOBAL symbols which have been defined and those which have not, the Assembler maintains undefined GLOBAL symbols in the symbol table with the value OFFFFH. This implies that no GLOBAL symbol should be given this value, or it will be considered to be undefined.

EXTERNAL -- External Symbol Declaration

The EXTERNAL pseudo-op is used to declare that each of its operands are symbols defined in some other module but referenced in this module.

Label	Opcode	Operands
	EXTERNAL	name1,name2,...

There must be no label. There may be zero or more names which are separated by delimiters and refer to names which are defined in some other module. EXTERNAL pseudo-ops may occur anywhere within the source text. The EXTERNAL pseudo-op assigns each name an external mode, which allows the name to be used in arbitrary expressions elsewhere in the module, subject to the rules for external expressions.

A name may not be declared as both EXTERNAL and GLOBAL. If the name also appears as a label (i.e., the name is assigned a value within this module), an error message will be generated.

Examples:

Label	Opcode	Operands
	EXTERNAL	TTY
	EXTERNAL	SYSTEM, MEMTOP, MEMBOT

Sample Relocatable Program

This simple program is an incomplete specification of three separate modules -- a main routine and an input and output routine -- where each is maintained in a different source file and is assembled separately from the others. In addition, each source file contains an *INCLUDE of the same file of constants (EQUates) to maintain consistency throughout the program and avoid repetitious typing and editing.

File MAIN.S contains:

```
GLOBAL MAIN,BUFFER
EXTERNAL INPUT,OUTPUT
*INCLUDE EQUATES
FULL:    DEFB FALSE           ; LOCAL FLAG BYTE
BUFFER:  DEFS 80              ; GLOBAL DATA BUFFER
MAIN:
    LD A,(FULL)              ; TEST BUFFER FLAG
    CP FALSE
    CALL Z,INPUT             ; FILL BUFFER IF EMPTY

    LD A,(FULL)              ; TEST BUFFER FLAG
    CP TRUE
    CALL Z,OUTPUT           ; OUTPUT BUFFER IF FULL

    JP MAIN                  ; INFINITE LOOP
END
```

File INDRIVER.S contains:

```
GLOBAL INPUT
EXTERNAL BUFFER
*INCLUDE EQUATES
INPUT:   ; DRIVER ENTRY POINT
        LD HL,BUFFER        ; ACCESS GLOBAL DATA BUFFER

        RET
        END
```

File OUTDRIVER.S contains:

```
GLOBAL OUTPUT
EXTERNAL BUFFER
*INCLUDE EQUATES
OUTPUT:  ; DRIVER ENTRY POINT
        LD HL,BUFFER        ; ACCESS GLOBAL DATA BUFFER

        RET
        END
```

File EQUATES.S contains:

```
TRUE EQU    OFFH
FALSE EQU   0
```

Reference Counter Control

ORG -- Origin

The ORG pseudo-op is used to set the reference counter to the value of an expression. The reference counter serves the same function for the Assembler as the program counter does for the CPU in specifying where object code is located.

Label	Opcode	Operand
optional:	ORG	expression

The reference counter is set to the 16-bit value of the expression, so that the next machine instruction or data item will be located at the specified address. Any symbol appearing in the expression must have been defined before the Assembler encounters the expression. The reference counter is initially set to zero, so if no ORG statement precedes the first instruction or data byte in the module, that byte will be assembled at location zero. If a label is present, it will be assigned the same value as the expression. A module may contain any number of ORG statements.

The mode of the expression in an ORG pseudo-op must not be external and depends on the relocatability of the module. If a module is absolute, the ORG pseudo-op serves to assign an absolute address to both the reference counter and the label. The expression will be treated as relocatable in a relocatable module, since the effect is to change the relocatable reference counter. Thus the label on an ORG statement in a relocatable module will have a relocatable mode. For example, the effect of the statement

Label	Opcode	Operand
LAB:	ORG	100

within a relocatable module would be to set the reference counter to "relocatable 100", assign the label LAB the value 100, and give it a relocatable mode.

Relocatable modules do not generally contain ORG statements, since the pseudo-op is useful only to reserve space within the module (in a manner similar to the DEFS pseudo-op), and not to assign an absolute location to a section of code since this would defeat the purpose of relocation. Since modules are by default relocatable, a warning message will be printed if an ORG pseudo-op is encountered in a relocatable module to indicate to the user that he may have forgotten to specify the Absolute option

in the Assembler command line. This warning message may be suppressed by the NOW option.

Some applications require that a section of code must start on a particular address boundary, for instance, at an address which is some multiple of a memory page size (say, 256 bytes per page). A useful formula for incrementing the reference counter to the next page boundary is $N*((\$+N-1)/N)$ where N is the page size. For example, suppose the page size is 8 and it is required that an interrupt vector table INTVEC be located at the next page boundary following the current value of the reference counter. Then the following statement achieves this effect regardless of the current value of the reference counter.

Label	Opcode	Operand
INTVEC:	ORG	$8*((\$+7)/8)$

Notice that this formula is useful only in an absolute module, since multiplication and division are undefined for relocatable operands (the reference counter symbol \$ is relocatable in a relocatable module).

*Implementation Note

In order to assist the Linker in assigning origins to relocatable modules, the Assembler computes the size of a module so that other modules can be relocated immediately following it. This length is computed as the value of the reference counter at the end of assembly minus the origin of the module. The origin of a relocatable module is relocatable zero, while the origin of an absolute module is the lowest address in the module which contains object code. In either case, if a program contains an ORG statement which sets the reference counter to a value less than it was previously, unexpected results may occur. In particular, an absolute module whose reference counter at the end of assembly is less than the lowest address which contains object code will cause an erroneous length to be computed. The programmer is advised that, in general, the reference counter should be incremented in a monotonic fashion so that addresses are in ascending order.

Conditional Assembly

Conditional assembly allows the programmer to inhibit the assembly of portions of the source text provided certain conditions are satisfied. Conditional assembly is particularly useful when a program needs to contain similar code sequences for slightly different applications. Rather

than generating a multitude of programs to handle each application, the application-dependent sections of code can be enclosed by the conditional pseudo-ops within a single program. Then by changing the values of several symbols used to control the conditional assembly, the user can generate different object modules from subsequent assemblies of the same source.

Another common use of conditional assembly is in conjunction with macros to control generation of code dependent on the value of parameters (see Macros).

COND and ENDC -- Conditional Assembly

The pseudo-op COND is used to test the value of an operand and, depending on the result, inhibit assembly of subsequent statements until an ENDC pseudo-op is encountered.

Label	Opcode	Operand
optional:	COND	expression
optional:	ENDC	

The Assembler evaluates the expression in the operand field of the COND pseudo-op. If the 16-bit result is true (non-zero), the COND pseudo-op is ignored and subsequent statements are assembled normally. If the result is false (zero), the assembly of subsequent statements is inhibited until an ENDC pseudo-op is encountered. The mode of the expression can be either relocatable or absolute; however, it may not be external. There is no operand for the ENDC pseudo-op.

Notice that the definition of symbols within a conditional assembly block may be inhibited, and thus references to these symbols elsewhere in the module may cause undefined symbol errors. In particular, the label on an ENDC pseudo-op will not be defined if the assembly is inhibited when the pseudo-op is encountered.

Conditional assembly blocks cannot be nested, i.e., the occurrence of an ENDC causes resumption of assembly of subsequent statements. Therefore, each COND does not require a matching ENDC.

The user must be cautious of nesting conditional blocks within macro definitions (see Macros) and vice-versa. If a macro definition is only partially completed because a surrounding conditional assembly block has inhibited the definition, or conversely, if a macro contains an unmatched COND/ENDC, unexpected results may occur.

Examples:

Label	Opcode	Operand
	COND	FLAG ; these statements are ; assembled only if the ; value of FLAG is not ; equal to zero
	ENDC	

(see Macros for further examples)

1.6 MACROS

HORSESHIT
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IS REDEFINE MNEMONICS

INSTRUCTIONS

Macros provide a means for the user to define his own opcodes, or to redefine existing opcodes. A macro defines a body of text which will be automatically inserted in the source stream at each occurrence of a macro call. In addition, parameters provide a capability for making limited changes in the macro at each call. If a macro is used to redefine an existing opcode, a warning message is generated to indicate that future use of that opcode will always be processed as a macro call. If a program uses macros, then the assembly option M must be specified (described later).

Macro Definition

The body of text to be used as a macro is given in the macro definition. Each definition begins with a Macro pseudo-op and ends with an ENDM pseudo-op. The general forms are:

Label	Opcode	Operands
name:	MACRO	#P0, #P1, ..., #Pn
optional:	ENDM	

The name is required, and must obey all the usual rules for forming labels (remember the colon is optional if the name starts in column one).

There can be any number of parameters (including none at all), each starting with the symbol #. The rest of the parameter name can be any string not containing a delimiter (blank, comma, tab, semicolon) or the symbol #. However, parameters will be scanned left to right for a match, so the user is cautioned not to use parameter names which are prefix substrings of later parameter names. Parameter names are not entered in the symbol table.

The label on an ENDM is optional, but if one is given it must obey all the usual rules for forming labels.

Each statement between the MACRO and ENDM statements is entered into a temporary macro file. The only restriction on these statements is that they do not include another macro definition (nested definitions are not allowed). They may include macro calls. The depth of nested calls is limited only by the available buffer space in the macro parameter stack. The amount of buffer space required by a sequence of nested calls is dependent on the number and length of parameters, which in general allows calls nested to a depth of about 15.

The statements of the macro body are not assembled at definition time, so they will not define labels, generate code, or cause errors. Exceptions are the assembler commands such as *List, which are executed wherever they occur. Within the macro body text, the formal parameter names may occur anywhere that an expansion-time substitution is desired. This includes within comments and quoted strings. The symbol # may not occur except as the first symbol of a parameter name.

Macros must be defined before they are called.

Macro Calls and Macro Expansion

A macro is called by using its name as an opcode at any point after the definition. The general form is:

```
label: name 'S1', 'S2',... 'Sn'
```

The label is optional and will be assigned to the current value of the reference counter, while the name must be a previously-defined macro. There may be any number of argument strings, Si, separated by any number of blanks, tabs, or commas. Commas do not serve as parameter place holders, only as string delimiters. If there are too few parameters, the missing ones are assumed to be null. If there are too many, the extras are ignored. The position of each string in the list corresponds with the position of the macro parameter name it is to replace. Thus the third string in a macro call statement will be substituted for each occurrence of the third parameter name.

The strings may be of any length and may contain any characters. The outer level quotes around the string are generally optional, but are required if the string contains delimiters or the quote character itself. The quote character is represented by two successive quote marks at the inner level. The outer level quotes, if present, will not occur in the substitution. The null string, represented by two successive quote marks at the outer level, may be used in any parameter position.

After processing the macro call statement, the Assembler switches its input from the source file to the macro file. Each statement of the macro body is scanned for occurrences of parameter names, and for each occurrence found, the corresponding string from the macro call statement is substituted. After substitution, the statement is assembled normally.

Symbol Generator

Every macro definition has an implicit parameter named `#$YM`. This may be referenced by the user in the macro body, but should not explicitly appear in the `MACRO` statement. At expansion time, each occurrence of `#$YM` in the definition is replaced by a string representing a 4-digit hexadecimal constant.

This string is constant over a given level of macro expansion, but increases by one for each new macro call. The most common use of `#$YM` is to provide unique labels for different expansions of the same macro. Otherwise, a macro containing a label would cause multiple definition errors if it were called more than once. Notice that a generated label must start with an alphabetic character just like all other labels.

The following is a simple macro which fills a buffer specified by the parameter `#BUFFER` with the number of blanks specified by `#COUNT`. The use of `#$YM` is necessary to create a unique label for each call to the macro.

```
FILLBL    MACRO #BUFFER #COUNT
          LD  B,#COUNT          ; number of bytes
          LD  HL,#BUFFER         ; address of buffer
FL#$YM    LD  (HL), ' '         ; insert blank
          INC HL                 ; advance pointer
          DJNZ FL#$YM           ; loop until done
          ENDM
```

If this is the first macro call, it would generate the following expansion when called with parameters `INBUF` and `50`:

```
FILLBL INBUF 50
LD B,50
LD HL,INBUF
FLO000:  LD (HL),' '
          INC HL
          DJNZ FLO000
          ENDM
```

The next call would cause the label `FLO001` to be generated, and so on.

Recursion

Macros may include calls to other macros, including themselves. A macro which directly calls itself (or indirectly by calling a second macro which calls the first

macro) is said to be recursive. Each recursive call causes a new expansion of the macro, possibly with different parameters. In order to prevent the macro being called endlessly, conditional assembly is used to inhibit a recursive call when certain conditions are satisfied (see the following examples). A macro which calls itself more times than the macro parameter stack can accommodate will generate an error. This often indicates that the recursion halting conditions would never be satisfied and the macro would be endlessly calling itself.

Listing Format

By default, each expanded statement is listed with a blank STMT field to differentiate macro expansions from normal source statements. If the Maclist flag is turned off by the NOM option or *M OFF, then only the macro call is listed.

Examples:

The following example, MOVE, is a macro which is used to move (copy) a block of memory. The source address, destination address, and number of bytes are all specified by parameters to the macro call.

The macro also provides the option to save the registers before the move and restore them upon completion. This option is implemented through the use of conditional assembly and the fact that any missing parameters are assumed to be null. Thus, if a non-null parameter is substituted for #SAVE?, the true condition of the COND statement is met and the code to save and restore the registers is assembled. If no parameter is passed, the #SAVE? parameter will be equal to the null string and the code will be ignored.

```

1 ; MOVE MACRO--BLOCK MOVE OF #LENGTH BYTES
2 ;           FROM #SOURCE TO #DEST
3 ;           IF 4TH PARAMETER IS PRESENT,
4 ;           THEN REGISTERS ARE PRESERVED
5 ;
6 MOVE  MACRO #DEST #SOURCE #LENGTH #SAVE?
7       COND .NOT.('#SAVE?='' )
8       PUSH BC
9       PUSH DE
10      PUSH HL
11      ENDC
12      LD  BC,#LENGTH ;MOVE #LENGTH BYTES
13      LD  DE,#DEST
14      LD  HL,#SOURCE
15      LDIR
16      COND .NOT.('#SAVE?='' )
17      POP  HL
18      POP  DE
19      POP  BC
20      ENDC
21      ENDM

```

If the MOVE macro is called as follows:

```
MOVE OUTBUF INBUF 25 S
```

code would be generated to save the registers BC, DE and HL, to move 25 bytes from INBUF to OUTBUF and then to restore registers BC, DE and HL.

```

22 ; A SAMPLE USE OF 'MOVE' WHICH PRESERVES
23 ; REGISTERS
24 MOVE OUTBUF INBUF 25 S
      COND .NOT.('S'='')
0000 C5      PUSH BC
0001 D5      PUSH DE
0002 E5      PUSH HL
      ENDC
0003 011900  LD  BC,25 ;MOVE 25 BYTES
0006 111C00  LD  DE,OUTBUF
0009 213500  LD  HL,INBUF
000C EDB0    LDIR
      COND .NOT.('S'='')
000E E1      POP  HL
000F D1      POP  DE
0010 C1      POP  BC
      ENDC

```

If the last parameter were omitted and the call was made as follows:

```
MOVE OUTBUF INBUF 25
```

the code to preserve the registers would not be generated because the #SAVE? parameter is the null string.

```
25 ; THIS 'MOVE' DOES NOT SAVE REGISTER
26 ; CONTENTS
27 MOVE OUTBUF INBUF 25
      COND .NOT.(''='')
      PUSH BC
      PUSH DE
      PUSH HL
      ENDC
0011 011900 LD BC,25 ;MOVE 25 BYTES
0014 111C00 LD DE,OUTBUF
0017 213500 LD HL,INBUF
001A EDB0 LDIR
      COND .NOT.(''='')
      POP HL
      POP DE
      POP BC
      ENDC
28
001C 29 OUTBUF: DEFS 25
0035 30 INBUF: DEFS 25
```

The next example demonstrates the use of recursion, as well as the use of the DEFL pseudo-op, to redefine the value of a label to be used as a counter to control the recursion.

When DUP is called, it initializes the label to the value passed to #CNT and then calls DUP1 with the parameter that was passed to #OP.

The DUP1 macro decrements the value of the label and recursively calls itself until the label has a value of zero.

The second call to DUP demonstrates the use of the *M OFF and *M ON assembler commands to suppress the listing of the expansion of the macro (see Assembler commands description).

```

1 ; DUPLICATE MACRO -- DUPLICATES #OP FOR #CNT
2 ; LINES USES DLAB LABEL AS RECURSION FACTOR
3 DUP MACRO #OP #CNT
4 DLAB DEFL #CNT
5     DUP1 '#OP'
6     ENDM
7
8 ; DUP1 RECURSIVE MACRO -- OUTPUTS #OP EACH TIME
9 ; RECURS ON THE VALUE OF THE LABEL DLAB
10 DUP1 MACRO #OP
11     COND DLAB>0
12 DLAB DEFL DLAB-1
13     DUP1 '#OP'
14     #OP
15     ENDC
16     ENDM
17
18 ; ROTATE RIGHT ACCUMULATOR 3 TIMES
19     DUP RRA 3
    DLAB DEFL 3
        DUP1 'RRA'
        COND DLAB>0
    DLAB DEFL DLAB-1
        DUP1 'RRA'
        RRA
        ENDC
0000     1F        RRA
        ENDC
0001     1F        RRA
        ENDC
0002     1F        RRA
        ENDC
20 *M OFF
0003     21        DUP 'INC HL' 4 ; INCREMENT HL BY 4
        22 *M ON

```

The last example, DB, is a macro which accepts up to 10 parameters, and issues a DEFB pseudo-op for each value.

The macro issues a DEFB for its first parameter and then calls itself recursively with the same parameters it received minus the first. Thus with each call the parameters are shifted to the left by one, and any missing parameters become the null string. After all of the values have been processed, the call to DB will pass the null string as parameter #1 and the condition of the COND statement will not be met. Notice how the outer level of single quotes is stripped from the parameter strings.

```

25 ; DB MACRO -- ACCEPTS UP TO 10 VALUES ON SAME L
26 ; OUTPUTS A DEFB FOR EACH ONE -- SELF-RECURSIV
27 DB MACRO #0 #1 #2 #3 #4 #5 #6 #7 #8 #9
28 DEFB #0
29 COND '#1'
30 DB '#1' '#2' '#3' '#4' '#5' '#6' '#7' '#8' '#9'
31 ENDC
32 ENDM
33
34 DB 1,2,3,4,5,6,7,8,9,10
0007 01 DEFB 1
COND '2'
0008 02 DB '2' '3' '4' '5' '6' '7' '8' '9' '10'
DEFB 2
COND '3'
0009 03 DB '3' '4' '5' '6' '7' '8' '9' '10' ''
DEFB 3
COND '4'
000A 04 DB '4' '5' '6' '7' '8' '9' '10' '' ''
DEFB 4
COND '5'
000B 05 DB '5' '6' '7' '8' '9' '10' '' '' ''
DEFB 5
COND '6'
000C 06 DB '6' '7' '8' '9' '10' '' '' '' ''
DEFB 6
COND '7'
000D 07 DB '7' '8' '9' '10' '' '' '' '' ''
DEFB 7
COND '8'
000E 08 DB '8' '9' '10' '' '' '' '' '' ''
DEFB 8
COND '9'
000F 09 DB '9' '10' '' '' '' '' '' '' ''
DEFB 9
COND '10'
0010 0A DB '10' '' '' '' '' '' '' '' '' ''
DEFB 10
COND ''
DB '' '' '' '' '' '' '' '' '' ''
ENDC

```

```

0011 02 36          DB LAB,LAB*3+1,TRUE,FALSE,OFFH
          DEFB LAB
          COND 'LAB*3+1'
          DB 'LAB*3+1' 'TRUE' 'FALSE' 'OFFH' '' '' '' '' ''
0012 07          DEFB LAB*3+1
          COND 'TRUE'
          DB 'TRUE' 'FALSE' 'OFFH' '' '' '' '' '' '' ''
0013 FF          DEFB TRUE
          COND 'FALSE'
          DB 'FALSE' 'OFFH' '' '' '' '' '' '' '' ''
0014 00          DEFB FALSE
          COND 'OFFH'
          DB 'OFFH' '' '' '' '' '' '' '' '' ''
0015 FF          DEFB OFFH
          COND ''
          DB '' '' '' '' '' '' '' '' '' ''
          ENDC
          ENDC
          ENDC
          ENDC
          ENDC
0016 37 *M OFF
          38          DB 0 0 0 0 0 0 0 0 0 0
          39 *M ON
          40 LAB: EQU 2
          41 TRUE: EQU OFFH
          42 FALSE: EQU 0

```


Assembler command line options description. The included file may also contain an *Include command, up to a nested level of four.

If the *Include command appears within a macro definition, the Assembler will always try to shoe-horn the file inside the macro definition, and although the *Include statement will appear in a macro expansion, the file will not be included again at the point of expansion. *Include works in the expected manner in conjunction with conditional assembly. For example,

```
                COND      exp
*Include        FILE1
                ENDC
;FILE1 is included only if the value of exp is non-zero.
```

1.8 ASSEMBLER COMMAND LINE OPTIONS

The Assembler interacts with the RIO operating system to permit the user complete control over the routing of all input and output (I/O) during assembly. The Assembler makes use of both the console device and disk files for I/O. The reader may wish to refer to the RIO Operating System User's Manual mentioned in the preface for a more complete description of files and I/O handling.

Using the Assembler

The Assembler is invoked from the RIO Executive by a command in the general form:

```
ASM file_name* [(options)]
```

where * means that one or more file names may be specified in the standard RIO format, and the square brackets mean that the options are optional. However, if present, they must be enclosed by parentheses. The Assembler processes the given source files in the order specified to produce a single object module and a listing file. It is not necessary to specify options to control assembly since defaults exist, but when included they must be enclosed by parentheses. The command is terminated by either a return or a semicolon.

When invoked, the Assembler first identifies itself by printing on the console its name and version number, where n.n is the current version specification:

```
ASM n.n
```

The Assembler processes the source program in two passes. The first pass identifies instructions, tests for correctness of operand combinations, defines the value of symbols and writes an intermediate file to the disk. At the conclusion of this process, a message is printed on the console:

```
PASS 1 COMPLETE
```

During the second assembly pass, listing and object files are created as output. The listing file contains the source language statements in their original format as well as the corresponding machine language and memory addresses. The listing file may optionally include a symbol cross-reference. The listing file is an ASCII type file with 128 bytes per record.

The object module consists of a binary type file with 128 bytes per record which is meaningful only to the Linker. It consists of four basic parts: a Header, the External Symbol Dictionary (ESD), Code blocks, and an optional Internal Symbol Dictionary (ISD). The Header contains information describing the module such as the number of ESD entries and the number of bytes of code. The ESD consists of all the GLOBAL and EXTERNAL names and possibly some indication of where the names are defined in the following Code blocks. The Code blocks contain both the actual machine code, as indicated in the listing, as well as a bit-map which indicates which instructions are to be relocated or linked to EXTERNAL names. The ISD is created only if the Symbol option in the command line is specified. The ISD contains all the symbols which are internal to the module (i.e., not GLOBAL or EXTERNAL) and is used by the Linker to create a symbol table file for use by a symbolic debug package. See the Appendix for a complete description of an object module.

Both files are created by default on the same disk drive with the first source file. If files already exist with the same name, they are erased and replaced with the newly generated files.

Errors in the source program detected during assembly are indicated by appropriate messages following the specific line both on the console and in the listing file (see the Appendix). In addition, the number of errors is indicated on the console at the end of the second pass by the following message where n is a decimal number:

n ASSEMBLY ERRORS

Normally, the assembly is complete after two passes. However, if either the cross-reference or the symbol option are specified, the Assembler starts a third pass by printing the following message:

PASS 2 COMPLETE

At the conclusion of the assembly, a message is printed on the console:

ASSEMBLY COMPLETE

Both the second and third pass are implemented as overlays to a portion of the first pass. The overlays are in procedure type files named ASM2 and ASM3, respectively. If these files cannot be found when required, the Assembler will abort.

At any point during the assembly, entering a '?' character will cause the Assembler to pause until another '?' character is entered. This allows the user to scan a portion of output while it is generated. In addition, if the ESCape character (ASCII 1BH) is entered, the Assembler will close all files immediately and abort the assembly. Notice that if the assembly is aborted, output files may contain incomplete information.

Several temporary scratch files are created by the Assembler, but are erased upon completion of an assembly. They include, depending on the options specified, the intermediate file, macro file, cross-reference file, and a symbol table overflow file. The user can control the assignment of these files if necessary (see below).

The Assembler uses a buffered I/O technique for handling the assembly language source file, listing file, object file and temporary files. The Assembler automatically determines the available work space and allocates the buffer sizes accordingly. Hence there are no constraints on the size of the assembly language source file that can be assembled, provided there is adequate space for the files.

Module Identification

The Assembler follows the general principle that anywhere a file name may be specified, that file name may be completely or partially qualified by either a device driver name or a drive specification.

For example:

- 1) ASM PROG1.S
- 2) ASM \$MYDOS:2/multiply

The first example demonstrates the use of an unqualified file name, PROG1.S, which can be found on the master device. If that device is ZDOS, then the standard disk drive search sequence of drive 1,2,...0 will be used. The second example demonstrates a fully qualified file name, multiply, which can be found on drive 2 of the device MYDOS.

Each source file must be ASCII type with 128 bytes per record.

When the period character is used to separate a file name into several parts, those parts are referred to as file name "extensions". For instance, the file name PROG1.S in the above example has the extension .S. The notion of file name extensions is a useful convention for the user who

wishes to categorize certain files by their names. To assist in this categorization, and to provide a uniform method for naming generated files such as the associated object and listing files, the Assembler insists that all source file names must end in the extension .S (for source). Notice there may be several extensions, as long as the last one is .S, and either upper or lower case is acceptable. Furthermore, the extension does not have to be typed in the command line, so that if the last two characters are not .S, the Assembler will automatically append the extension before attempting to open the file.

By default, the object module and list file will be given the name of the first source file which comprises the module, as indicated in the command line, but with the extensions .OBJ and .L, respectively. The case of the extension is the same as the first character of the file name.

For example:

- 1) ASM YOURPROG
looks for source YOURPROG.S, creates
YOURPROG.OBJ and YOURPROG.L
- 2) ASM myprog1.s myprog2 TTY.INT
looks for sources myprog1.s, myprog2.s, and
TTY.INT.S, creates myprog1.obj and myprog1.l

It is possible to override the name of either the object or listing file by appropriate options in the command line as described below.

Options

Several options can be used to modify the ASM command; however, the list of options must be enclosed in parentheses following the source file names. The options may be specified in any order, with capital letters in the following list signifying the minimum allowable abbreviation. Notice that either upper or lower case can be used.

In order to maximize the effectiveness of allocation of file buffers, and thus reduce I/O activity, certain processing must be explicitly requested:

Macro	Enables macro processing. If not specified, macro definitions and calls will be flagged as errors.
-------	--

Symbol	Causes a third pass to be executed, producing a binary representation of the internal symbols (i.e., not GLOBAL or EXTERNAL) to be appended to the object module. This information may be used by a symbolic debug package.
Xref	Causes a third pass to be executed, producing a sorted cross-reference table at the end of the list file.

The following options will override default settings:

Absolute	Causes the entire module to be assembled with absolute addresses. Default is relocatable.
NOList	Suppresses the creation of the listing file.
NOMaclist	Suppresses the listing and printing of macro expansions.
NOObject	Suppresses the creation of the object file.
NOWarning	Suppresses the generation of opcode redefinition warning messages, and relocatable ORG warning messages.

The following option aids documentation:

D=string	Places up to an 18-character string, not including blanks, tabs, commas, or semicolons, in the heading of each page of the listing (normally used to specify a date).
----------	---

Note the following interactions:

NOL overrides *List and *Maclist commands in the text.

*Maclist commands in the text override NOM.

X will override NOL for the third pass only. That is, (NOL X) will produce a list file containing only the cross-reference listing.

NOO overrides S.

I/O Routing Options

The Assembler uses several logical I/O units, some of which depend on the requirements of the module being assembled. The user can control the definition of these units either through interaction with the RIO Executive before giving the ASM command, or by specifying certain options in the command line. Pass completion notices and all fatal error messages which cause the Assembler to abort (see Fatal Error Messages) are always routed to the console (logical unit CONOUT). Normally, errors in the source statements will also be routed to the console, unless overridden by the Print option.

Print	Causes the listing output as it is generated to be routed to the logical unit SYSLST, in addition to being directed to a listing file. This is particularly useful if SYSLST is defined to be a line printer driver, for instance. The device driver must recognize the WRITE ASCII request code to operate properly. The Print option inhibits source statement errors from being sent to the console to avoid duplicate messages when SYSLST and CONOUT are both assigned to the same device.
-------	---

The following options will override the default definition of output and temporary files used by the Assembler.

file_name can be a completely or partially qualified file name.

I=file_name	defines intermediate file
L=file_name	defines listing file
M=file_name	defines macro file
O=file_name	defines object file
T=file_name	defines symbol table overflow file
X=file_name	defines cross-reference file

The symbol table overflow file is created automatically but only if necessary. M= and X= imply the Macro and Xref options, respectively.

Notice that giving a complete file name to a file which would normally be a temporary scratch file will cause that file to remain after assembly. For instance, I=3 causes the intermediate file to be assigned to a scratch file (zero-length name) on drive 3 of the master device, whereas I=3/I.TMP causes the intermediate file to remain after assembly as the file I.TMP on drive 3 of the master device.

Examples:

- 1) ASM PROG1
produces PROG1.OBJ and PROG1.L on the master device and same drive as the source program, PROG1.S, will use the same drive for both the intermediate file and the symbol table overflow file if necessary, and will not print the listing on SYSLST.
- 2) ASM PROG1 (O=\$MYDOS:0/A.OUT S X=2 P NOL)
produces an object file with attached symbol table named A.OUT on drive 0 of device \$MYDOS, prints the listing with a cross-reference on SYSLST, uses drive 2 of the master device for the temporary cross-reference file, and produces no listing file. The intermediate file and symbol table overflow file will be created as scratch files on the same device and drive as the source, PROG1.S.

Logical Unit Definitions

The Assembler uses the following logical I/O units:

2	CONOUT	Error messages
3	SYSLST	Print option output
4		List file
5		Object file
6		Intermediate file
7		Macro file
8		Xref file
9		Symbol table overflow file
10-13		Source files (multiple units used to support nested *Include commands)

Fatal Error Messages

There are several conditions that may occur which will cause the Assembler to terminate further processing, close all files, and abort by returning to the RIO Executive. The following messages may appear on the console:

INVALID OPTION: s

Indicates that the string s was not recognized as a valid command line option. The cause may be, for instance, a misspelled option specifier or invalid use of delimiters within options.

MEMORY TOO SMALL

Indicates that the amount of available buffer space is not adequate to permit assembly. Buffers are allocated during each pass depending on the various options requested in the command line.

LINE TOO LONG: filename

Indicates that a source text line in "filename" was encountered that was longer than 128 characters. Due to buffer allocation, longer lines are not acceptable. This error may indicate that the contents of the source file have been damaged.

FILE NOT FOUND: filename

Indicates that the source "filename" could not be found in a device directory. Remember that the source file name must end in the extension .S.

INVALID ATTRIBUTES: filename

Indicates that the source "filename" had attributes other than type ASCII (subtype is ignored) and a record length of 128 bytes.

INVALID FILE: filename

Indicates that the string "filename" is not a proper file name specifier. This may include illegal characters within the name or an invalid device or drive designator (see RIO Operating System User's Manual for further details). The string "filename" is printed as it appeared in the command line.

I/O ERROR e ON UNIT u

All other fatal errors involve I/O errors (hex value e) on a particular logical unit (decimal value u). Refer to the Logical Unit Definitions for the Assembler and the RIO Operating System User's Manual for an explanation of specific errors.

2.0 LINKER

2.1 FUNCTION

The function of the Linker is to process one or more object modules created by the Assembler (version 4.0 or greater), or any other language translator which generates relocatable object modules, and output a single program in the form of an executable procedure file. The Linker provides relocation of modules and resolves inter-module references to allow linking of separately assembled modules. In addition, a load map file and a binary symbol table file may be created.

The Linker interacts with the RIO operating system to permit the user complete control over the routing of all input and output (I/O) during linkage. The Linker makes use of both the console device and disk files for I/O. The reader may wish to refer to the RIO Operating System User's Manual mentioned in the preface for a more complete description of files and I/O handling.

Features

- 1) Creates a segmented procedure type file of either 128, 256, 512, or 1024 bytes per record. Code may be assigned to any location in memory, and up to 16 different segments of code can be generated which allows for discontinuous sections of memory which are not overlaid when the procedure file is loaded.
- 2) Provides ability to link both absolute and relocatable modules together.
- 3) Allows optional assignment of some or all relocatable modules to specific absolute origins.
- 4) Permits specification of the execution starting address as either a hexadecimal number or a GLOBAL symbol.
- 5) Checks for multiply-defined GLOBALS and unresolved EXTERNALS.
- 6) Provides the capability to specify "link-only" modules, that is, modules which are used to control assignment of relocatable module origins and resolve EXTERNAL references, but are not included as part of the generated procedure file.

- 7) Creates a load map with the assigned origin and length of each module, followed by an alphabetically sorted list of all GLOBALs, their assigned addresses and the module which contains them. The map file also contains a list of any error messages generated during linkage.
- 8) Allows optional creation of a binary symbol table file, structured by module, for use by a symbolic debug package.
- 9) Provides user control over the routing of all input and output during linkage.

2.2 USING THE LINKER

The linker is invoked from the RIO Executive by issuing the command LINK followed by a list of module names and options as described below. The general form is:

```
LINK file_name* [(options)]
```

where * means that one or more file names may be specified in the standard RIO format, and the square brackets mean that the options are optional; however, if present they must be enclosed by parentheses. The command is terminated by either a return or a semicolon.

When invoked, the Linker first identifies itself by printing on the console its name and version number, where n.n is the current version specification:

```
LINK n.n
```

There are two distinct passes to the Linker. During the first pass the Linker must build what is called the Link Directory, which allows it to assign absolute addresses to modules and GLOBAL names. EXTERNAL names must be partially resolved -- i.e., assigned the absolute address of the corresponding GLOBAL name. Modules are normally assigned ascending addresses in the order they are specified.

During the second pass, the object code of each module is processed, local references are relocated, EXTERNAL references are completely resolved, and both a procedure file and a load map file are generated. An optional third pass will occur if a symbol file is to be generated. At the end of linkage, if there were any non-fatal errors, an error count is printed on the console:

```
n ERRORS
```

where n is the total number of multiply-defined GLOBALs and unresolved EXTERNALs. Other messages are considered only warning messages.

The final message on the console is:

LINK COMPLETE

At any time during linkage, the '?' character can be entered, and the Linker will pause until another '?' character is entered. If the ESCape character is entered, the Linker will close all files immediately and abort by returning to the RIO Executive. If the Linker is aborted, output files may contain incomplete information.

Non-fatal Errors

There are two non-fatal error conditions which can occur during the first phase of linking. Multiply-defined GLOBAL names will be listed and should be considered a significant error. Unresolved EXTERNALs will be listed but may not indicate an error, since the module may not have actually used the EXTERNAL name in a reference, for instance, during the initial checkout of a program. These errors will cause the Linker to suppress creation of a procedure file unless the LET option is specified (see Options).

If either of these errors exist, an appropriate message is output to both the console and map file, followed by a list of the symbol names and the appropriate object module which caused the error:

MULTIPLY-DEFINED GLOBAL IN MODULE:
symbol_name module_name

This message indicates that symbol_name was defined as a GLOBAL in two different modules, where module_name contains the second definition. If a load map is generated, the user can determine the other module from the alphabetic list of GLOBALs. References to a multiply-defined GLOBAL will always refer to the first definition.

UNRESOLVED EXTERNAL IN MODULE:
symbol_name module_name

This message indicates that symbol_name which was declared to be EXTERNAL to module_name could not be matched to a corresponding GLOBAL name. If a procedure file is created with unresolved EXTERNAL references, those references will contain unpredictable addresses.

Assignment of Module Origins

The Linker maintains a reference counter which is similar to the CPU program counter and is used to assign sections of machine code to memory locations. By default, relocatable modules are located in contiguous memory areas automatically by assigning the current value of the reference counter to be the module's origin and then incrementing the reference counter by the module's length. The reference counter is always initialized to zero.

In order to control the assignment of relocatable modules to certain absolute addresses, it is possible to intersperse assignments to the reference counter between module names. The form for these assignments is either:

1) $\$=X$

X is a hexadecimal number (no trailing 'H')

2) $\$=\$+X$

where the first allows direct assignment to the reference counter which causes the next relocatable module to have that address as its origin. The second form allows the reference counter to be incremented by an appropriate amount, say to allocate some buffer space between modules. For example, to link three modules together with an origin of 4000 and a data buffer module at 9000:

```
LINK    $=4000    MODA    MODB    MODC    $=9000    DATA
```

MODA, MODB and MODC will follow immediately after each other in memory starting at location 4000.

If the reference counter assignment is poorly formed (embedded delimiters or an invalid hexadecimal number) an invalid option message will be issued and the Linker will abort. Notice that absolute modules may be linked to relocatable modules or to each other. Absolute modules contain their own address information, so it makes little sense to use \$ when linking absolute modules. The reference counter is reset to the lowest address of the absolute module which contains object code and then incremented by the length of the module. Both the lowest address and the length of the module are calculated by the language translator which created the object module.

Segment Generation

Often the need arises in a program to separate it into several non-contiguous segments of machine code due either

to certain hardware requirements--such as the application memory configuration or peripheral device interface restrictions--or certain software requirements--such as the need to implement an overlay scheme. The Linker provides an automatic facility for generating up to 16 non-contiguous memory segments, each of which consists of an arbitrary number of bytes located at an arbitrary starting point. Each segment contains a number of bytes which is always an integral multiple of the procedure file's record length (default of 128 bytes--see Record Length option).

When the resulting procedure file is loaded, only those areas of memory occupied by the segments will be modified. The Linker uses two strategies in determining when to start generating a new segment:

- 1) A new segment is started for machine code which is to be located at an address which is less than the previous value of the reference counter. This may occur as the result of an ORG statement in assembly language or from the assignment of the Linker's reference counter using the \$=X construct. The Linker will output the following message:

```
POSSIBLE CODE OVERLAY AT n IN module_name
```

which indicates that machine code may be overlaid at memory location n when the procedure file is subsequently loaded. The module_name indicates the object module which causes the new segment to be generated. If the maximum allowable number of segments (16) is exhausted, the Linker will abort.

Example:

```
LINK $=5000 MODA $=4000 MODB
```

module MODB will start a new segment

- 2) A new segment is started for machine code which is to be located at an address which is greater than the procedure file's record length beyond the current reference counter value, that is, the address of the last generated byte of machine code. This avoids entire records of meaningless data in the procedure file which might be caused by either a large DEFS statement in assembly language or an assignment to the Linker's reference counter using the \$=X construct. If the maximum allowable number of segments is exhausted, the Linker is forced to generate all remaining machine code in the last segment.

Example:

```
LINK $=1000 MODA    $=2000 MODB
```

module MODB will start a new segment unless MODA is larger than 400 bytes, assuming a record length of 400.

Link-only Modules

A standard technique for implementing a program which is larger than the available memory space is to "overlay" one portion on top of another which is no longer needed. The various portions can be maintained as separate procedure files which are loaded by program interaction with the RIO Executive (see Overlays for further details). In order to facilitate the construction of overlay programs, the Linker provides a method to specify that certain modules are "link-only"; that is, the object modules will be used to help assign module starting addresses and to resolve EXTERNAL references, but will not cause generation of machine code in the resulting procedure files.

Any module name that is immediately preceded by the minus character '-' (ASCII 2DH) will be considered to be a link-only module. A link-only module will appear in the load map along with its GLOBALs, but it will not be loaded as a part of the procedure file.

Example:

```
LINK  $=2000  RESIDENT  -OVERLAY.2
```

would cause the creation of a procedure file containing only the machine code of the module RESIDENT and not that of OVERLAY.2. However, any EXTERNAL references within RESIDENT which are satisfied by GLOBALs in OVERLAY.2 will be correctly resolved.

Module Identification

The Linker follows the general principle that anywhere a file name may be specified, it may be completely or partially qualified by either a device driver name or a drive specification.

For example:

- 1) LINK MODA MODB.OBJ
- 2) LINK \$MYDOS:2/multiply

The first example demonstrates the use of unqualified file names, MODA and MODB.OBJ, which can be found on the master device. If that device is ZDOS, then the standard disk drive search sequence of drive 1,2,...,0 will be used. The second example demonstrates a fully qualified file name, multiply, which can be found on drive 2 of the device MYDOS.

Each object module must be a binary type file with 128 bytes per record. In addition, an object module contains information within it which identifies it, and the Linker will abort if a module does not contain the correct identification.

Each object module file name must end in the extension .OBJ, and either upper or lower case is acceptable. Furthermore, the extension does not have to be typed in the command line, so that if the last four characters are not .OBJ, the Linker will automatically append the extension before attempting to open the file.

By default, the resulting procedure file will have the same name as the first object file which comprises the program as indicated in the command line, but without the .OBJ extension. The map file and symbol file will have the same name as the procedure file, but with the extensions .MAP and .SYM, respectively. The case of the extension is the same as the first character of the file name.

For example:

- 1) LINK YOURPROG
looks for object YOURPROG.OBJ; creates
YOURPROG, YOURPROG.MAP, and YOURPROG.SYM
- 2) LINK myprog1.obj myprog2 TTY.INT
looks for object modules myprog1.obj, myprog2.obj
and TTY.INT.OBJ; creates myprog1, myprog1.map
and myprog1.sym

It is possible to override the name of either the procedure file, map file or symbol file by appropriate options in the command line as described below.

2.3 LINKER COMMAND LINE OPTIONS

As the name implies, options are optional and default values will be assumed as indicated. If options are desired, they must be enclosed in parentheses following the object module file names. The format for options is a keyword--usually just the first letter as indicated by the capitals in the following list--sometimes followed by an equal sign (=), and an appropriate value or name as specified for each option. Options can be specified in any order and the last occurrence of a particular option will override previous ones.

Notice that several of the command options are used to control the definition of the logical I/O units used by the Linker. Thus, through interaction with the RIO Executive before giving the LINK command and by using the options which specify a file name (and the Print option), the user can direct the routing of all I/O during linkage. Notice that "filename" in the following descriptions means a standard RIO file name which may be completely or partially qualified.

All fatal error messages which cause the Linker to abort (see Fatal Error Messages) are always routed to the console (logical unit CONOUT). Normally, non-fatal errors and warning messages are also routed to the console, unless overridden by the Print option, in which case they are routed to SYSLSST.

Entry=n

Used to specify the starting execution address for the procedure file. n may be either a hex value (must start with a numeric digit, a leading zero is sufficient) or a GLOBAL name. Default--the starting address is taken to be the assigned origin of the first object module. If a GLOBAL name is specified but not found, a warning message will be output and the default origin will be used.

LET

Used to force the second phase of linking even if there were errors such as unresolved EXTERNALS or multiply-defined GLOBALS. Default--errors inhibit the generation of a procedure file.

Map=filename

Used to cause generation of a load map which indicates module starting addresses and lengths. In addition, all GLOBALs are listed (alphabetically sorted) with their assigned addresses and the module which contains them. The filename specifies a file where the load map and any non-fatal error messages are to be output. Default--a load map (ASCII type with 128 bytes per record) is generated and output to a file with extension .MAP with the same name as the procedure file.

Name=filename

Used to specify the name of the procedure file to be created. The name defaults to that of the first module name minus the .OBJ extension. The name will also be given to the map and symbol files (unless overridden by specific options) with the appropriate extensions .MAP and .SYM.

NOMap

Suppresses creation of the load map file. Default--a load map is output to a file.

NOWarning

Used to inhibit warning messages about multiply-defined GLOBALs or unresolved EXTERNALs. Suppresses warnings that possible overlay of code may occur when the reference counter is decremented. Also inhibits the checking for multiply-defined GLOBALs and thus speeds up the first phase of linking. Default--warning messages are output and multiply-defined GLOBALs are checked for.

Print

Causes the output of the load map and error messages to be routed to SYSLST. Default--only error messages are output to CONOUT and nothing is output to SYSLST.

RLength=n

Used to specify the record length of the procedure file. n is a hexadecimal number equal to either 80H, 100H, 200H, or 400H (no trailing 'H'). Default--record length is 80H (128 bytes decimal).

Stacksize=n

Used to specify the size of the stack space required by the procedure file when executed. n is a hexadecimal number of bytes. Default--a stack size of 80H bytes is specified.

Symbol

Symbol=filename

Used to create a binary symbol file (128 bytes per record) combining the absolutized GLOBALs and the ISD (Internal Symbol Dictionary) of each object module for use by a symbolic debug package. Notice that some modules may have no ISD. The file name, if not specified, defaults to the same name as the procedure file, with extension .SYM. Default--no symbol file is created.

Note the following interactions:

NOM overrides Map=filename.

NOW is an implied LET since errors will be suppressed.

If a procedure file is not successfully generated, there will be no symbol file and the map file will only contain error messages.

If the user specifies a scratch file (i.e., a file name with a zero length name) for an output file, that file will disappear after linkage.

Examples:

- 1) LINK \$=1000 PROG1 PROG2
produces a procedure file PROG1 (with entry point 1000) and a map file PROG1.MAP on the master device and same drive as the object module, PROG1.OBJ, and will not print the load map on SYSLST.
- 2) LINK PROG1 (N=\$MYDOS:2/PROG.RUN S NOM P E=MAIN)
produces a procedure file PROG.RUN (with entry point equal to the assigned address of GLOBAL MAIN) and a symbol file PROG.RUN.SYM on drive 2 of the device MYDOS, will not create a map file but will print the load map and any error messages on SYSLST.

Logical Unit Definitions

The Linker uses the following logical I/O units:

2	CONOUT	Error messages
3	SYSLST	Print option output
4		Object files
5		Procedure file
6		Map file
7		Symbol file

Fatal Error Messages

There are several conditions that may occur which cause the Linker to terminate further processing, close all files, and abort by returning to the RIO Executive. The following messages may appear on the console:

INVALID OPTION: s

Indicates that the string s was not recognized as a valid command line option. The cause may be, for instance, a misspelled option specifier or invalid use of delimiters within options. Notice that all numbers must be specified in hexadecimal notation (but no trailing 'H').

LINK DIRECTORY OVERFLOW

Indicates that there is insufficient memory space to permit linkage. The Link Directory is basically a composite list of GLOBALS and EXTERNALS for each object module which the Linker uses to relocate and resolve object code references. File buffers are allocated in proportion to available space; however, there must be room for both the Link Directory and the minimum size buffer for the various files. The size of the Link Directory may be kept small by reducing the number and length of GLOBAL symbols in a program and by not including EXTERNAL declarations for symbols never referenced in a module. This message will also occur if the user attempts to link more than 127 modules together.

PROGRAM TOO BIG

Indicates that the reference counter was incremented past FFFFH, which is beyond the address space of the CPU. The program size or origin should be reduced.

TOO MANY SEGMENTS

Indicates that more than the maximum number (16) of procedure file segments was created. This may occur when the reference counter is decremented too many times to allow a new segment to be generated. The reference counter is decremented by either an assignment to \$ in the Linker command line or by an ORG statement in assembly language within an object module. This problem can be avoided by rearranging modules or sections of code so that the reference counter is incremented in ascending order.

FILE NOT FOUND: filename

Indicates that the object "filename" could not be found in a device directory. Remember that the object file name must end in the extension .OBJ.

INVALID FILE: filename

Indicates that the string "filename" is not a proper file name specifier. This may include illegal characters within the name or an invalid device or drive designator (see RIO Operating System User's Manual for further details). The string "filename" is printed as it appeared in the command line.

INVALID FORMAT: filename

Indicates that the object filename had attributes other than type binary (subtype is ignored) and a record length of 128 bytes. This message will also appear if the file does not contain a header of data which is consistent with an object module. This may indicate that the file was damaged and should be regenerated by reassembling the module.

INVALID DATA: filename

Indicates that the object filename contains invalid data, such as a checksum error, or a symbol name whose length is either zero or greater than the maximum allowed (currently 127). This usually occurs as a result of a damaged file which should be regenerated by reassembling the module.

I/O ERROR e ON UNIT u

All other fatal errors involve I/O errors (hex value e) on a particular logical unit (decimal value u). Refer to the Logical Unit Definitions for the Linker and the RIO Operating System User's Manual for an explanation of specific errors.

2.4 MAP FORMAT

The Linker produces a load map which indicates the assigned origin and length of each module, and an alphabetic list of all GLOBALs with their assigned addresses and the modules that contain them. In addition, the map contains the program's name, its length in bytes, and its entry point. Refer to the following example which is, hopefully, self-explanatory.

```
LINK $=1000 MODA MODB $=2000 COS.MATH SIN.MATH (N=COMPUTE
E=MAIN)
```

```
LINK 1.0
```

```
LOAD MAP
```

MODULE	ORIGIN	LENGTH
MODA	1000	0B32
MODB	1B32	01AA
COS.MATH	2000	00B0
SIN.MATH	20B0	00B6

GLOBAL	ADDRESS	MODULE
--------	---------	--------

COS	2000	COS.MATH
MAIN	10CC	MODA
SIN	20B0	SIN.MATH
SQRT	1B32	MODB
VECTOR	1669	MODA

```
PROGRAM COMPUTE -- 1166 BYTES
ENTRY: 10CC
```

The total number of bytes for a program is computed as the difference between the lowest and highest addresses which actually contain machine code in the procedure file. Thus the number of bytes may not accurately reflect the actual amount of memory a program occupies.

2.5 OVERLAYS

This section provides some guidelines for the user whose program must be implemented as an overlay structure due to memory space constraints. By utilizing certain facilities of the Linker, such as assignment of relocatable module origins and link-only modules, the user may build an effective overlay scheme. Notice, however, that overlay management must be built into the user's program in the sense that the program must explicitly initiate the loading of an overlay segment before any references are made to that segment.

An overlay structure normally has a resident segment which contains data and routines which are common to the various overlay segments. In particular, the routine which initiates the loading of the other segments would normally be in the resident segment. The other segments will normally be maintained as separate procedure files which have been created through a series of linkages as described below.

The method for structuring an overlay program is best illustrated by a simple example. Consider a program (EXAMPLE) which gathers a large amount of data from a particular source (READ), reorganizes and calculates new data (PROCESS), and then outputs this data to a particular destination (WRITE). Assume that this sequence only needs to be done once. Presumably, there is only room for one of these segments to be resident at a time, along with the data and a few common routines, but by simply over-laying the previous segment, the program can accomplish its desired goal.

By properly declaring data and routines as GLOBAL and EXTERNAL in each of the four modules which comprise the program, the user can then combine them in separate linkages as follows. Since READ can fit in memory with the resident segment, EXAMPLE, the two modules are linked together into a single procedure file called EXAMPLE. Notice that because the resident module contains references to both PROCESS and WRITE, these modules are included as link-only modules to allow the proper references to be resolved.

```
1) LINK EXAMPLE $=5000 READ $=5000 -PROCESS
    $=5000 -WRITE
```

Then the two overlay segments are created by linking the resident segment with the appropriate module, but specifying the resident as a link-only module.

```
2) LINK -EXAMPLE $=5000 PROCESS (N=PROCESS)
3) LINK -EXAMPLE $=5000 WRITE (N=WRITE)
```

The program can be executed by issuing the command EXAMPLE, and provided the resident portion initiates the loading of the two overlay segments at the proper time, the entire program will run to completion.

Overlay procedure files may be loaded in several ways. A program can contain code to open the correct file and read the contents into memory at the right location. Alternatively, a command string which loads a file without executing it can be passed to the RIO operating system by making an I/O request to logical unit 0 with a pointer to a section of memory which contains a string such as 'PROCESS,' followed by a carriage return. See the RIO Operating System User's Manual for further details.

Much more complicated overlay structures can be designed in the same general fashion. For instance, an overlay segment may cause a portion of itself to be overlaid. Or an overlay segment may be loaded several times, instead of just once, and if it contains data that must be saved, it must be written to a temporary file before it is overlaid. By careful planning and programming, the user may utilize the facilities of the Linker to create an overlay schema suitable to the application.

APPENDIX A
OPCODE LISTING

LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
0000	00		1	nop	; 00
0001	018605	R	2	ld bc,nn	; 01
0004	02		3	ld (bc),a	; 02
0005	03		4	inc bc	; 03
0006	04		5	inc b	; 04
0007	05		6	dec b	; 05
0008	0620		7	ld b,n	; 06
000A	07		8	rlca	; 07
000B	08		9	ex af,af'	; 08
000C	09		10	add hl,bc	; 09
000D	0A		11	ld a,(bc)	; 0a
000E	0B		12	dec bc	; 0b
000F	0C		13	inc c	; 0c
0010	0D		14	dec c	; 0d
0011	0E20		15	ld c,n	; 0e
0013	0F		16	rrca	; 0f
0014	102E		17	djnz \$+dis	; 10
0016	118605	R	18	ld de,nn	; 11
0019	12		19	ld (de),a	; 12
001A	13		20	inc de	; 13
001B	14		21	inc d	; 14
001C	15		22	dec d	; 15
001D	1620		23	ld d,n	; 16
001F	17		24	rla	; 17
0020	182E		25	jr \$+dis	; 18
0022	19		26	add hl,de	; 19
0023	1A		27	ld a,(de)	; 1a
0024	1B		28	dec de	; 1b
0025	1C		29	inc e	; 1c
0026	1D		30	dec e	; 1d
0027	1E20		31	ld e,n	; 1e
0029	1F		32	rra	; 1f
002A	202E		33	jr nz,\$+dis	; 20
002C	218605	R	34	ld hl,nn	; 21
002F	228605	R	35	ld (nn),hl	; 22
0032	23		36	inc hl	; 23
0033	24		37	inc h	; 24
0034	25		38	dec h	; 25
0035	2620		39	ld h,n	; 26
0037	27		40	daa	; 27
0038	282E		41	jr z,\$+dis	; 28
003A	29		42	add hl,hl	; 29
003B	2A8605	R	43	ld hl,(nn)	; 2a
003E	2B		44	dec hl	; 2b
003F	2C		45	inc l	; 2c
0040	2D		46	dec l	; 2d
0041	2E20		47	ld l,n	; 2e
0043	2F		48	cpl	; 2f
0044	302E		49	jr nc,\$+dis	; 30
0046	318605	R	50	ld sp,nn	; 31
0049	328605	R	51	ld (nn),a	; 32
004C	33		52	inc sp	; 33
004D	34		53	inc (hl)	; 34
004E	35		54	dec (hl)	; 35
004F	3620		55	ld (hl),n	; 36
0051	37		56	scf	; 37
0052	382E		57	jr c,\$+dis	; 38
0054	39		58	add hl,sp	; 39

LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT
0055	3A8605	R	59	ld a,(nn) ; 3a
0058	3B		60	dec sp ; 3b
0059	3C		61	inc a ; 3c
005A	3D		62	dec a ; 3d
005B	3E20		63	ld a,n ; 3e
005D	3F		64	ccf ; 3f
005E	40		65	ld b,b ; 40
005F	41		66	ld b,c ; 41
0060	42		67	ld b,d ; 42
0061	43		68	ld b,e ; 43
0062	44		69	ld b,h ; 44
0063	45		70	ld b,l ; 45
0064	46		71	ld b,(hl) ; 46
0065	47		72	ld b,a ; 47
0066	48		73	ld c,b ; 48
0067	49		74	ld c,c ; 49
0068	4A		75	ld c,d ; 4a
0069	4B		76	ld c,e ; 4b
006A	4C		77	ld c,h ; 4c
006B	4D		78	ld c,l ; 4d
006C	4E		79	ld c,(hl) ; 4e
006D	4F		80	ld c,a ; 4f
006E	50		81	ld d,b ; 50
006F	51		82	ld d,c ; 51
0070	52		83	ld d,d ; 52
0071	53		84	ld d,e ; 53
0072	54		85	ld d,h ; 54
0073	55		86	ld d,l ; 55
0074	56		87	ld d,(hl) ; 56
0075	57		88	ld d,a ; 57
0076	58		89	ld e,b ; 58
0077	59		90	ld e,c ; 59
0078	5A		91	ld e,d ; 5a
0079	5B		92	ld e,e ; 5b
007A	5C		93	ld e,h ; 5c
007B	5D		94	ld e,l ; 5d
007C	5E		95	ld e,(hl) ; 5e
007D	5F		96	ld e,a ; 5f
007E	60		97	ld h,b ; 60
007F	61		98	ld h,c ; 61
0080	62		99	ld h,d ; 62
0081	63		100	ld h,e ; 63
0082	64		101	ld h,h ; 64
0083	65		102	ld h,l ; 65
0084	66		103	ld h,(hl) ; 66
0085	67		104	ld h,a ; 67
0086	68		105	ld l,b ; 68
0087	69		106	ld l,c ; 69
0088	6A		107	ld l,d ; 6a
0089	6B		108	ld l,e ; 6b
008A	6C		109	ld l,h ; 6c
008B	6D		110	ld l,l ; 6d
008C	6E		111	ld l,(hl) ; 6e
008D	6F		112	ld l,a ; 6f
008E	70		113	ld (hl),b ; 70
008F	71		114	ld (hl),c ; 71
0090	72		115	ld (hl),d ; 72
0091	73		116	ld (hl),e ; 73

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LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
0092	74	117	ld (hl),h	; 74	
0093	75	118	ld (hl),l	; 75	
0094	76	119	halt	; 76	
0095	77	120	ld (hl),a	; 77	
0096	78	121	ld a,b	; 78	
0097	79	122	ld a,c	; 79	
0098	7A	123	ld a,d	; 7a	
0099	7B	124	ld a,e	; 7b	
009A	7C	125	ld a,h	; 7c	
009B	7D	126	ld a,l	; 7d	
009C	7E	127	ld a,(hl)	; 7e	
009D	7F	128	ld a,a	; 7f	
009E	80	129	add a,b	; 80	
009F	81	130	add a,c	; 81	
00A0	82	131	add a,d	; 82	
00A1	83	132	add a,e	; 83	
00A2	84	133	add a,h	; 84	
00A3	85	134	add a,l	; 85	
00A4	86	135	add a,(hl)	; 86	
00A5	87	136	add a,a	; 87	
00A6	88	137	adc a,b	; 88	
00A7	89	138	adc a,c	; 89	
00A8	8A	139	adc a,d	; 8a	
00A9	8B	140	adc a,e	; 8b	
00AA	8C	141	adc a,h	; 8c	
00AB	8D	142	adc a,l	; 8d	
00AC	8E	143	adc a,(hl)	; 8e	
00AD	8F	144	adc a,a	; 8f	
00AE	90	145	sub b	; 90	
00AF	91	146	sub c	; 91	
00B0	92	147	sub d	; 92	
00B1	93	148	sub e	; 93	
00B2	94	149	sub h	; 94	
00B3	95	150	sub l	; 95	
00B4	96	151	sub (hl)	; 96	
00B5	97	152	sub a	; 97	
00B6	98	153	sbc a,b	; 98	
00B7	99	154	sbc a,c	; 99	
00B8	9A	155	sbc a,d	; 9a	
00B9	9B	156	sbc a,e	; 9b	
00BA	9C	157	sbc a,h	; 9c	
00BB	9D	158	sbc a,l	; 9d	
00BC	9E	159	sbc a,(hl)	; 9e	
00BD	9F	160	sbc a,a	; 9f	
00BE	A0	161	and b	; a0	
00BF	A1	162	and c	; a1	
00C0	A2	163	and d	; a2	
00C1	A3	164	and e	; a3	
00C2	A4	165	and h	; a4	
00C3	A5	166	and l	; a5	
00C4	A6	167	and (hl)	; a6	
00C5	A7	168	and a	; a7	
00C6	A8	169	xor b	; a8	
00C7	A9	170	xor c	; a9	
00C8	AA	171	xor d	; aa	
00C9	AB	172	xor e	; ab	
00CA	AC	173	xor h	; ac	
00CB	AD	174	xor l	; ad	

LOC	OBJ CODE	M	STMT	SOURCE	STATEMENT
00CC	AE		175	xor (hl)	; ae
00CD	AF		176	xor a	; af
00CE	B0		177	or b	; b0
00CF	B1		178	or c	; b1
00D0	B2		179	or d	; b2
00D1	B3		180	or e	; b3
00D2	B4		181	or h	; b4
00D3	B5		182	or l	; b5
00D4	B6		183	or (hl)	; b6
00D5	B7		184	or a	; b7
00D6	B8		185	cp b	; b8
00D7	B9		186	cp c	; b9
00D8	BA		187	cp d	; ba
00D9	BB		188	cp e	; bb
00DA	BC		189	cp h	; bc
00DB	BD		190	cp l	; bd
00DC	BE		191	cp (hl)	; be
00DD	BF		192	cp a	; bf
00DE	C0		193	ret nz	; c0
00DF	C1		194	pop bc	; c1
00E0	C28605	R	195	jp nz,nn	; c2
00E3	C38605	R	196	jp nn	; c3
00E6	C48605	R	197	call nz,nn	; c4
00E9	C5		198	push bc	; c5
00EA	C620		199	add a,n	; c6
00EC	C7		200	rst 0	; c7
00ED	C8		201	ret z	; c8
00EE	C9		202	ret	; c9
00EF	CA8605	R	203	jp z,nn	; ca
00F2	CC8605	R	204	call z,nn	; cc
00F5	CD8605	R	205	call nn	; cd
00F8	CE20		206	adc a,n	; c
00FA	CF		207	rst 8	; cf
00FB	D0		208	ret nc	; d0
00FC	D1		209	pop de	; d1
00FD	D28605	R	210	jp nc,nn	; d2
0100	D320		211	out (n),a	; d3
0102	D48605	R	212	call nc,nn	; d4
0105	D5		213	push de	; d5
0106	D620		214	sub n	; d6
0108	D7		215	rst 10h	; d7
0109	D8		216	ret c	; d8
010A	D9		217	exx	; d9
010B	DA8605	R	218	jp c,nn	; da
010E	DB20		219	in a,(n)	; db
0110	DC8605	R	220	call c,nn	; dc
0113	DE20		221	sbc a,n	; de
0115	DF		222	rst 18h	; df
0116	E0		223	ret po	; e0
0117	E1		224	pop hl	; e1
0118	E28605	R	225	jp po,nn	; e2
011B	E3		226	ex (sp),hl	; e3
011C	E48605	R	227	call po,nn	; e4
011F	E5		228	push hl	; e5
0120	E620		229	and n	; e6
0122	E7		230	rst 20h	; e7
0123	E8		231	ret pe	; e8
0124	E9		232	jp (hl)	; e9

LOC	OBJ CODE	M	STMT	SOURCE	opcode STATEMENT	
0125	EA8605	R	233	jp pe,nn		; ea
0128	EB		234	ex de,hl		; eb
0129	EC8605	R	235	call pe,nn		; ec
012C	EE20		236	xor n		; ee
012E	EF		237	rst 28h		; ef
012F	F0		238	ret p		; f0
0130	F1		239	pop af		; f1
0131	F28605	R	240	jp p,nn		; f2
0134	F3		241	di		; f3
0135	F48605	R	242	call p,nn		; f4
0138	F5		243	push af		; f5
0139	F620		244	or n		; f6
013B	F7		245	rst 30h		; f7
013C	F8		246	ret m		; f8
013D	F9		247	ld sp,hl		; f9
013E	FA8605	R	248	jp m,nn		; fa
0141	FB		249	ei		; fb
0142	FC8605	R	250	call m,nn		; fc
0145	FE20		251	cp n		; fe
0147	FF		252	rst 38h		; ff
0148	CB00		253	rlc b		; 00
014A	CB01		254	rlc c		; 01
014C	CB02		255	rlc d		; 02
014E	CB03		256	rlc e		; 03
0150	CB04		257	rlc h		; 04
0152	CB05		258	rlc l		; 05
0154	CB06		259	rlc (hl)		; 06
0156	CB07		260	rlc a		; 07
0158	CB08		261	rrc b		; 08
015A	CB09		262	rrc c		; 09
015C	CB0A		263	rrc d		; 0a
015E	CB0B		264	rrc e		; 0b
0160	CB0C		265	rrc h		; 0c
0162	CB0D		266	rrc l		; 0d
0164	CB0E		267	rrc (hl)		; 0e
0166	CB0F		268	rrc a		; 0f
0168	CB10		269	rl b		; 10
016A	CB11		270	rl c		; 11
016C	CB12		271	rl d		; 12
016E	CB13		272	rl e		; 13
0170	CB14		273	rl h		; 14
0172	CB15		274	rl l		; 15
0174	CB16		275	rl (hl)		; 16
0176	CB17		276	rl a		; 17
0178	CB18		277	rr b		; 18
017A	CB19		278	rr c		; 19
017C	CB1A		279	rr d		; 1a
017E	CB1B		280	rr e		; 1b
0180	CB1C		281	rr h		; 1c
0182	CB1D		282	rr l		; 1d
0184	CB1E		283	rr (hl)		; 1e
0186	CB1F		284	rr a		; 1f
0188	CB20		285	sla b		; 20
018A	CB21		286	sla c		; 21
018C	CB22		287	sla d		; 22
018E	CB23		288	sla e		; 23
0190	CB24		289	sla h		; 24
0192	CB25		290	sla l		; 25

LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
0194	CB26	291	sla (hl)		; 26
0196	CB27	292	sla a		; 27
0198	CB28	293	sra b		; 28
019A	CB29	294	sra c		; 29
019C	CB2A	295	sra d		; 2a
019E	CB2B	296	sra e		; 2b
01A0	CB2C	297	sra h		; 2c
01A2	CB2D	298	sra l		; 2d
01A4	CB2E	299	sra (hl)		; 2e
01A6	CB2F	300	sra a		; 2f
01A8	CB38	301	srl b		; 38
01AA	CB39	302	srl c		; 39
01AC	CB3A	303	srl d		; 3a
01AE	CB3B	304	srl e		; 3b
01B0	CB3C	305	srl h		; 3c
01B2	CB3D	306	srl l		; 3d
01B4	CB3E	307	srl (hl)		; 3e
01B6	CB3F	308	srl a		; 3f
01B8	CB40	309	bit 0,b		; 40
01BA	CB41	310	bit 0,c		; 41
01BC	CB42	311	bit 0,d		; 42
01BE	CB43	312	bit 0,e		; 43
01C0	CB44	313	bit 0,h		; 44
01C2	CB45	314	bit 0,l		; 45
01C4	CB46	315	bit 0,(hl)		; 46
01C6	CB47	316	bit 0,a		; 47
01C8	CB48	317	bit 1,b		; 48
01CA	CB49	318	bit 1,c		; 49
01CC	CB4A	319	bit 1,d		; 4a
01CE	CB4B	320	bit 1,e		; 4b
01D0	CB4C	321	bit 1,h		; 4c
01D2	CB4D	322	bit 1,l		; 4d
01D4	CB4E	323	bit 1,(hl)		; 4e
01D6	CB4F	324	bit 1,a		; 4f
01D8	CB50	325	bit 2,b		; 50
01DA	CB51	326	bit 2,c		; 51
01DC	CB52	327	bit 2,d		; 52
01DE	CB53	328	bit 2,e		; 53
01E0	CB54	329	bit 2,h		; 54
01E2	CB55	330	bit 2,l		; 55
01E4	CB57	331	bit 2,a		; 57
01E6	CB56	332	bit 2,(hl)		; 56
01E8	CB58	333	bit 3,b		; 58
01EA	CB59	334	bit 3,c		; 59
01EC	CB5B	335	bit 3,e		; 5a
01EE	CB5B	336	bit 3,e		; 5b
01F0	CB5C	337	bit 3,h		; 5c
01F2	CB5D	338	bit 3,l		; 5d
01F4	CB5E	339	bit 3,(hl)		; 5e
01F6	CB5F	340	bit 3,a		; 5f
01F8	CB60	341	bit 4,b		; 60
01FA	CB61	342	bit 4,c		; 61
01FC	CB62	343	bit 4,d		; 62
01FE	CB63	344	bit 4,e		; 63
0200	CB64	345	bit 4,h		; 64
0202	CB65	346	bit 4,l		; 65
0204	CB66	347	bit 4,(hl)		; 66
0206	CB67	348	bit 4,a		; 67

LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
0208	CB68	349	bit 5,b		; 68
020A	CB69	350	bit 5,c		; 69
020C	CB6A	351	bit 5,d		; 6a
020E	CB6B	352	bit 5,e		; 6b
0210	CB6C	353	bit 5,h		; 6c
0212	CB6D	354	bit 5,l		; 6d
0214	CB6E	355	bit 5,(hl)		; 6e
0216	CB6F	356	bit 5,a		; 6f
0218	CB70	357	bit 6,b		; 70
021A	CB71	358	bit 6,c		; 71
021C	CB72	359	bit 6,d		; 72
021E	CB73	360	bit 6,e		; 73
0220	CB74	361	bit 6,h		; 74
0222	CB75	362	bit 6,l		; 75
0224	CB76	363	bit 6,(hl)		; 76
0226	CB77	364	bit 6,a		; 77
0228	CB78	365	bit 7,b		; 78
022A	CB79	366	bit 7,c		; 79
022C	CB7A	367	bit 7,d		; 7a
022E	CB7B	368	bit 7,e		; 7b
0230	CB7C	369	bit 7,h		; 7c
0232	CB7D	370	bit 7,l		; 7d
0234	CB7E	371	bit 7,(hl)		; 7e
0236	CB7F	372	bit 7,a		; 7f
0238	CB80	373	res 0,b		; 80
023A	CB81	374	res 0,c		; 81
023C	CB82	375	res 0,d		; 82
023E	CB83	376	res 0,e		; 83
0240	CB84	377	res 0,h		; 84
0242	CB85	378	res 0,l		; 85
0244	CB86	379	res 0,(hl)		; 86
0246	CB87	380	res 0,a		; 87
0248	CB88	381	res 1,b		; 88
024A	CB89	382	res 1,c		; 89
024C	CB8A	383	res 1,d		; 8a
024E	CB8B	384	res 1,e		; 8b
0250	CB8C	385	res 1,h		; 8c
0252	CB8D	386	res 1,l		; 8d
0254	CB8E	387	res 1,(hl)		; 8e
0256	CB8F	388	res 1,a		; 8f
0258	CB90	389	res 2,b		; 90
025A	CB91	390	res 2,c		; 91
025C	CB92	391	res 2,d		; 92
025E	CB93	392	res 2,e		; 93
0260	CB94	393	res 2,h		; 94
0262	CB95	394	res 2,l		; 95
0264	CB96	395	res 2,(hl)		; 96
0266	CB97	396	res 2,a		; 97
0268	CB98	397	res 3,b		; 98
026A	CB99	398	res 3,c		; 99
026C	CB9A	399	res 3,d		; 9a
026E	CB9B	400	res 3,e		; 9b
0270	CB9C	401	res 3,h		; 9c
0272	CB9D	402	res 3,l		; 9d
0274	CB9E	403	res 3,(hl)		; 9e
0276	CB9F	404	res 3,a		; 9f
0278	CBA0	405	res 4,b		; a0
027A	CBA1	406	res 4,c		; a1

LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
027C	CBA2	407	res	4,d	; a2
027E	CBA3	408	res	4,e	; a3
0280	CBA4	409	res	4,h	; a4
0282	CBA5	410	res	4,l	; a5
0284	CBA6	411	res	4,(hl)	; a6
0286	CBA7	412	res	4,a	; a7
0288	CBA8	413	res	5,b	; a8
028A	CBA9	414	res	5,c	; a9
028C	CBAA	415	res	5,d	; aa
028E	CBAB	416	res	5,e	; ab
0290	CBAC	417	res	5,h	; ac
0292	CBAD	418	res	5,l	; ad
0294	CBAE	419	res	5,(hl)	; ae
0296	CBAF	420	res	5,a	; af
0298	CBB0	421	res	6,b	; b0
029A	CBB1	422	res	6,c	; b1
029C	CBB2	423	res	6,d	; b2
029E	CBB3	424	res	6,e	; be
02A0	CBB4	425	res	6,h	; b4
02A2	CBB5	426	res	6,l	; b5
02A4	CBB6	427	res	6,(hl)	; b6
02A6	CBB7	428	res	6,a	; b7
02A8	CBB8	429	res	7,b	; b8
02AA	CBB9	430	res	7,c	; b9
02AC	CBBA	431	res	7,d	; ba
02AE	CBBB	432	res	7,e	; bb
02B0	CBBC	433	res	7,h	; bc
02B2	CBBD	434	res	7,l	; bd
02B4	CBBF	435	res	7,a	; bf
02B6	CBC0	436	set	0,b	; c0
02B8	CBC1	437	set	0,c	; c1
02BA	CBC2	438	set	0,d	; c2
02BC	CBC3	439	set	0,e	; c3
02BE	CBC4	440	set	0,h	; c4
02C0	CBC5	441	set	0,l	; c5
02C2	CBC6	442	set	0,(hl)	; c6
02C4	CBC7	443	set	0,a	; c7
02C6	CBC8	444	set	1,b	; c8
02C8	CBC9	445	set	1,c	; c9
02CA	CBCA	446	set	1,d	; ca
02CC	CBCB	447	set	1,e	; cb
02CE	CBCC	448	set	1,h	; cc
02D0	CBCD	449	set	1,l	; cd
02D2	CBCE	450	set	1,(hl)	; ce
02D4	CBCF	451	set	1,a	; cf
02D6	CBD0	452	set	2,b	; d0
02D8	CBD1	453	set	2,c	; d1
02DA	CBD2	454	set	2,d	; d2
02DC	CBD3	455	set	2,e	; d3
02DE	CBD4	456	set	2,h	; d4
02E0	CBD5	457	set	2,l	; d5
02E2	CBD6	458	set	2,(hl)	; d6
02E4	CBD7	459	set	2,a	; d7
02E6	CBD8	460	set	3,b	; d8
02E8	CBD9	461	set	3,c	; d9
02EA	CBDA	462	set	3,d	; da
02EC	CBDB	463	set	3,e	; db
02EE	CBDC	464	set	3,h	; dc

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LOC      OBJ CODE M STMT SOURCE      opcode
STATEMENT

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02F0    CBDD          465      set 3,l          ; dd
02F2    CBDE          466      set 3,(hl)       ; de
02F4    CBDF          467      set 3,a          ; df
02F6    CBE0          468      set 4,b          ; e0
02F8    CBE1          469      set 4,c          ; e1
02FA    CBE2          470      set 4,d          ; e2
02FC    CBE3          471      set 4,e          ; e3
02FE    CBE4          472      set 4,h          ; e4
0300    CBE5          473      set 4,l          ; e5
0302    CBE6          474      set 4,(hl)      ; e6
0304    CBE7          475      set 4,a          ; e7
0306    CBE8          476      set 5,b          ; e8
0308    CBE9          477      set 5,c          ; e9
030A    CBEA          478      set 5,d          ; ea
030C    CBEB          479      set 5,e          ; eb
030E    CBEC          480      set 5,h          ; ec
0310    CBED          481      set 5,l          ; ed
0312    CBEE          482      set 5,(hl)     ; ee
0314    CBEF          483      set 5,a          ; ef
0316    CBF0          484      set 6,b          ; f0
0318    CBF1          485      set 6,c          ; f1
031A    CBF2          486      set 6,d          ; f2
031C    CBF3          487      set 6,e          ; f3
031E    CBF4          488      set 6,h          ; f4
0320    CBF5          489      set 6,l          ; f5
0322    CBF6          490      set 6,(hl)     ; f6
0324    CBF7          491      set 6,a          ; f7
0326    CBF8          492      set 7,b          ; f8
0328    CBF9          493      set 7,c          ; f9
032A    CBFA          494      set 7,d          ; fa
032C    CBFB          495      set 7,e          ; fb
032E    CBFC          496      set 7,h          ; fc
0330    Cbfd          497      set 7,l          ; fd
0332    CBFE          498      set 7,(hl)     ; fe
0334    CBFF          499      set 7,a          ; ff
0336    DD09          500      add ix,bc       ; 09
0338    DD19          501      add ix,de       ; 19
033A    DD218605 R    502      ld ix,nn        ; 21
033E    DD228605 R    503      ld (nn),ix     ; 22
0342    DD23          504      inc ix          ; 23
0344    DD29          505      add ix,ix       ; 29
0346    DD2A8605 R    506      ld ix,(nn)     ; 2a
034A    DD2B          507      dec ix          ; 2b
034C    DD3405        508      inc (ix+ind)   ; 34
034F    DD3505        509      dec (ix+ind)   ; 35
0352    DD360520     510      ld (ix+ind),n  ; 36
0356    DD39          511      add ix,sp       ; 39
0358    DD4605        512      ld b,(ix+ind)  ; 46
035B    DD4E05        513      ld c,(ix+ind)  ; 4e
035E    DD5605        514      ld d,(ix+ind)  ; 56
0361    DD5E05        515      ld e,(ix+ind)  ; 5e
0364    DD6605        516      ld h,(ix+ind)  ; 66
0367    DD6E05        517      ld l,(ix+ind)  ; 6e
036A    DD7005        518      ld (ix+ind),b  ; 70
036D    DD7105        519      ld (ix+ind),c  ; 71
0370    DD7205        520      ld (ix+ind),d  ; 72
0373    DD7305        521      ld (ix+ind),e  ; 73
0376    DD7405        522      ld (ix+ind),h  ; 74

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LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
0379	DD7505		523	ld (ix+ind),l	; 75
037C	DD7705		524	ld (ix+ind),a	; 77
037F	DD7E05		525	ld a,(ix+ind)	; 7e
0382	DD8605		526	add a,(ix+ind)	; 86
0385	DD8E05		527	adc a,(ix+ind)	; 8e
0388	DD9605		528	sub (ix+ind)	; 96
038B	DD9E05		529	sbc a,(ix+ind)	; 9e
038E	DDA605		530	and (ix+ind)	; a6
0391	DDAE05		531	xor (ix+ind)	; ae
0394	DDB605		532	or (ix+ind)	; b6
0397	DDBE05		533	cp (ix+ind)	; be
039A	DDE1		534	pop ix	; e1
039C	DDE3		535	ex (sp),ix	; e3
039E	DDE5		536	push ix	; e5
03A0	DDE9		537	jp (ix)	; e9
03A2	DDF9		538	ld sp,ix	; f9
03A4	DDCB0506		539	rlc (ix+ind)	; 06
03A8	DDCB050E		540	rrc (ix+ind)	; 0e
03AC	DDCB0516		541	rl (ix+ind)	; 16
03B0	DDCB051E		542	rr (ix+ind)	; 1e
03B4	DDCB0526		543	sla (ix+ind)	; 26
03B8	DDCB052E		544	sra (ix+ind)	; 2e
03BC	DDCB053E		545	srl (ix+ind)	; 3e
03C0	DDCB0546		546	bit 0,(ix+ind)	; 46
03C4	DDCB054E		547	bit 1,(ix+ind)	; 4e
03C8	DDCB0556		548	bit 2,(ix+ind)	; 56
03CC	DDCB055E		549	bit 3,(ix+ind)	; 5e
03D0	DDCB0566		550	bit 4,(ix+ind)	; 66
03D4	DDCB056E		551	bit 5,(ix+ind)	; 6e
03D8	DDCB0576		552	bit 6,(ix+ind)	; 76
03DC	DDCB057E		553	bit 7,(ix+ind)	; 7e
03E0	DDCB0586		554	res 0,(ix+ind)	; 86
03E4	DDCB058E		555	res 1,(ix+ind)	; 8e
03E8	DDCB0596		556	res 2,(ix+ind)	; 96
03EC	DDCB059E		557	res 3,(ix+ind)	; 9e
03F0	DDCB05A6		558	res 4,(ix+ind)	; a6
03F4	DDCB05AE		559	res 5,(ix+ind)	; ae
03F8	DDCB05B6		560	res 6,(ix+ind)	; b6
03FC	DDCB05BE		561	res 7,(ix+ind)	; be
0400	DDCB05C6		562	set 0,(ix+ind)	; c6
0404	DDCB05CE		563	set 1,(ix+ind)	; ce
0408	DDCB05D6		564	set 2,(ix+ind)	; d6
040C	DDCB05DE		565	set 3,(ix+ind)	; de
0410	DDCB05E6		566	set 4,(ix+ind)	; e6
0414	DDCB05EE		567	set 5,(ix+ind)	; ee
0418	DDCB05F6		568	set 6,(ix+ind)	; f6
041C	DDCB05FE		569	set 7,(ix+ind)	; fe
0420	ED40		570	in b,(c)	; 40
0422	ED41		571	out (c),b	; 41
0424	ED42		572	sbc hl,bc	; 42
0426	ED438605 R		573	ld (nn),bc	; 43
042A	ED44		574	neg	; 44
042C	ED45		575	retn	; 45
042E	ED46		576	im 0	; 46
0430	ED47		577	ld i,a	; 47
0432	ED48		578	in c,(c)	; 48
0434	ED49		579	out (c),c	; 49
0436	ED4A		580	adc hl,bc	; 4a

LOC	OBJ CODE M	STMT	SOURCE	opcode STATEMENT	
0438	ED4B8605 R	581	ld bc,(nn)		; 4b
043C	ED4D	582	reti		; 4d
043E	ED4F	583	ld r,a		; 4f
0440	ED50	584	in d,(c)		; 50
0442	ED51	585	out (c),d		; 51
0444	ED52	586	sbc hl,de		; 52
0446	ED538605 R	587	ld (nn),de		; 53
044A	ED56	588	im 1		; 56
044C	ED57	589	ld a,i		; 57
044E	ED58	590	in e,(c)		; 58
0450	ED59	591	out (c),e		; 59
0452	ED5A	592	adc hl,de		; 5a
0454	ED5B8605 R	593	ld de,(nn)		; 5b
0458	ED5E	594	im 2		; 5e
045A	ED5F	595	ld a,r		; 5f
045C	ED60	596	in h,(c)		; 60
045E	ED61	597	out (c),h		; 61
0460	ED62	598	sbc hl,hl		; 62
0462	ED67	599	rrd		; 67
0464	ED68	600	in l,(c)		; 68
0466	ED69	601	out (c),l		; 69
0468	ED6A	602	adc hl,hl		; 6a
046A	ED6F	603	rld		; 6f
046C	ED72	604	sbc hl,sp		; 72
046E	ED738605 R	605	ld (nn),sp		; 73
0472	ED78	606	in a,(c)		; 78
0474	ED79	607	out (c),a		; 79
0476	ED7A	608	adc hl,sp		; 7a
0478	ED7B8605 R	609	ld sp,(nn)		; 7b
047C	EDA0	610	ldi		; a0
047E	EDA1	611	cpi		; a1
0480	EDA2	612	ini		; a2
0482	EDA3	613	outi		; a3
0484	EDA8	614	ldd		; a8
0486	EDA9	615	cpd		; a9
0488	EDAA	616	ind		; aa
048A	EDAB	617	outd		; ab
048C	EDB0	618	ldir		; b0
048E	EDB1	619	cpir		; b1
0490	EDB2	620	inir		; b2
0492	EDB3	621	otir		; b3
0494	EDB8	622	laddr		; b8
0496	EDB9	623	cpdr		; b9
0498	EDBA	624	indr		; ba
049A	EDBB	625	otdr		; bb
049C	FD09	626	add iy,bc		; 09
049E	FD19	627	add iy,de		; 19
04A0	FD218605 R	628	ld iy,nn		; 21
04A4	FD228605 R	629	ld (nn),iy		; 22
04A8	FD23	630	inc iy		; 23
04AA	FD29	631	add iy,iy		; 29
04AC	FD2A8605 R	632	ld iy,(nn)		; 2a
04B0	FD2B	633	dec iy		; 2b
04B2	FD3405	634	inc (iy+ind)		; 34
04B5	FD3505	635	dec (iy+ind)		; 35
04B8	FD360520	636	ld (iy+ind),n		; 36
04BC	FD39	637	add iy,sp		; 39
04BE	FD4605	638	ld b,(iy+ind)		; 46

LOC	OBJ CODE M STMT	SOURCE	opcode STATEMENT
04C1	FD4E05	639	ld c,(iy+ind) ; 4e
04C4	FD5605	640	ld d,(iy+ind) ; 56
04C7	FD5E05	641	ld e,(iy+ind) ; 5e
04CA	FD6605	642	ld h,(iy+ind) ; 66
04CD	FD6E05	643	ld l,(iy+ind) ; 6e
04D0	FD7005	644	ld (iy+ind),b ; 70
04D3	FD7105	645	ld (iy+ind),c ; 71
04D6	FD7205	646	ld (iy+ind),d ; 72
04D9	FD7305	647	ld (iy+ind),e ; 73
04DC	FD7405	648	ld (iy+ind),h ; 74
04DF	FD7505	649	ld (iy+ind),l ; 75
04E2	FD7705	650	ld (iy+ind),a ; 77
04E5	FD7E05	651	ld a,(iy+ind) ; 7e
04E8	FD8605	652	add a,(iy+ind) ; 86
04EB	FD8E05	653	adc a,(iy+ind) ; 8e
04EE	FD9605	654	sub (iy+ind) ; 96
04F1	FD9E05	655	sbc a,(iy+ind) ; 9e
04F4	FDA605	656	and (iy+ind) ; a6
04F7	FDAE05	657	xor (iy+ind) ; ae
04FA	FDB605	658	or (iy+ind) ; b6
04FD	FDBE05	659	cp (iy+ind) ; be
0500	FDE1	660	pop iy ; e1
0502	FDE3	661	ex (sp),iy ; e3
0504	FDE5	662	push iy ; e5
0506	FDE9	663	jp (iy) ; e9
0508	FDF9	664	ld sp,iy ; f9
050A	FDCB0506	665	rlc (iy+ind) ; 06
050E	FDCB050E	666	rrc (iy+ind) ; 0e
0512	FDCB0516	667	rl (iy+ind) ; 16
0516	FDCB051E	668	rr (iy+ind) ; 1e
051A	FDCB0526	669	sla (iy+ind) ; 26
051E	FDCB052E	670	sra (iy+ind) ; 2e
0522	FDCB053E	671	srl (iy+ind) ; 3e
0526	FDCB0546	672	bit 0,(iy+ind) ; 46
052A	FDCB054E	673	bit 1,(iy+ind) ; 4e
052E	FDCB0556	674	bit 2,(iy+ind) ; 56
0532	FDCB055E	675	bit 3,(iy+ind) ; 5e
0536	FDCB0566	676	bit 4,(iy+ind) ; 66
053A	FDCB056E	677	bit 5,(iy+ind) ; 6e
053E	FDCB0576	678	bit 6,(iy+ind) ; 76
0542	FDCB057E	679	bit 7,(iy+ind) ; 7e
0546	FDCB0586	680	res 0,(iy+ind) ; 86
054A	FDCB058E	681	res 1,(iy+ind) ; 8e
054E	FDCB0596	682	res 2,(iy+ind) ; 96
0552	FDCB059E	683	res 3,(iy+ind) ; 9e
0556	FDCB05A6	684	res 4,(iy+ind) ; a6
055A	FDCB05AE	685	res 5,(iy+ind) ; ae
055E	FDCB05B6	686	res 6,(iy+ind) ; b6
0562	FDCB05BE	687	res 7,(iy+ind) ; be
0566	FDCB05C6	688	set 0,(iy+ind) ; c6
056A	FDCB05CE	689	set 1,(iy+ind) ; ce
056E	FDCB05D6	690	set 2,(iy+ind) ; d6
0572	FDCB05DE	691	set 3,(iy+ind) ; de
0576	FDCB05E6	692	set 4,(iy+ind) ; e6
057A	FDCB05EE	693	set 5,(iy+ind) ; ee
057E	FDCB05F6	694	set 6,(iy+ind) ; f6
0582	FDCB05FE	695	set 7,(iy+ind) ; fe
		696	

LOC OBJ CODE M STMT SOURCE STATEMENT

opcode

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697 nn defs 2
698 ind equ 5
699 m equ 10h
700 n equ 20h
701 dis equ 30h
702 end

APPENDIX B
OBJECT MODULE FORMAT

Header - - - - -	ESD - - - - -	Code Blocks - - - - -	ISD
16 Bytes	Variable	128 Bytes/Block	Variable

I. Header

Byte #0	:	80 hex (absolute object file mark) 81 81 hex (relocatable object file mark)
1	:	Reserved for future use
2...3	:	# of bytes of code in module
4...5	:	# of ESD entries
6...7	:	Record # of first code block
8...9	:	Lowest absolute address used if absolute module
10...11	:	Reserved for future use
12...13	:	# of ISD entries
14...15	:	Record # of first ISD block

II. ESD (EXTERNAL SYMBOL DICTIONARY)

Each entry is of variable length, with zero-filled entries, if necessary, to fill the last physical record of the ESD. The format of an entry of size n:

Byte #0	:	Attribute byte - high order bit indicates a GLOBAL (bit=1) or EXTERNAL (bit=0) Low order bit indicates whether GLOBAL is relocatable (bit=1) or absolute (bit=0). All other bits are reserved for future use (currently=0).
1	:	Symbol name length in bytes (i.e., n-4)
2...n-3	:	Symbol name (ASCII characters)
n-2...n-3	:	16 bit value In the case of a GLOBAL name, the value is the address which defines the symbol. The value of an EXTERNAL name is its index (relative byte position) in the ESD which will be modified by the Linker.

III. CODE BLOCK

Byte #0 : Code Mark
(83 hex if this is the last code block in the module, 82 hex otherwise)

1...2 : Reference counter value of first byte in block

3 : # of code bytes in block

4...17 : Bitmap (only 5 high-order bits of #17 are used). The bitmap is organized with 1 bit per byte of code, where 0=absolute; 1,0=relocatable; 1,1=EXTERNAL reference.

18...126 : Object code
The two bytes for a relocatable reference is a 16 bit signed quantity to be added to the assigned starting location of the module at link time. The 2 bytes of code for an EXTERNAL reference is a 16 bit signed quantity to be added to the final assigned absolute address of the EXTERNAL. In addition, there are 2 extra bytes which contain a relative pointer to the attribute byte of the EXTERNAL name in the ESD (these two bytes are the relative byte position of the EXTERNAL in the ESD). These two bytes are not actually converted to machine code during the linking process, and thus have zero bits in the corresponding bitmap entries.
*** Relocatable and EXTERNAL references will not be split across a record boundary; instead, when necessary, they will start a new record.

127 : Checksum (sum of bytes up to Checksum + Checksum =0). In fact, the checksum byte actually follows the last good code byte in the block, any following bytes are ignored. Thus, in a block with zero bytes of code, the checksum byte will actually occur in byte #18.

IV. ISD (INTERNAL SYMBOL DICTIONARY - appended by Symbol Option in Assembler)

Same variable format as ESD - only difference is there are no GLOBALs and thus the GLOBAL bit in the attribute byte is always reset (=0). The low order bit of attribute byte indicates whether internal symbol is relocatable (bit=1) or absolute (bit=0).

APPENDIX C

ASSEMBLY ERROR MESSAGES AND EXPLANATIONS

- 1) **WARNING - OPCODE REDEFINED**
Indicates that an opcode has been redefined by a macro so that future uses of the opcode will result in the appropriate macro call. This message may be suppressed by the NOW option.
- 2) **NAME CONTAINS INVALID CHARACTERS**
Indicates that a name (either a label or an operand) contains illegal characters. Names must start with an alphabetic character and any following characters must be either alphanumeric (A...Z or 0...9), a question mark (?) or an underbar (_).
- 3) **INVALID OPCODE**
Indicates that the opcode was not recognized. Occurs when the opcode contains an illegal character (including non-printing control characters), when the opcode is not either all upper case or all lower case, or when macros are used and the M option is not specified.
- 4) **INVALID NUMBER**
Indicates an invalid character in a number. Occurs when a number contains an illegal character (including non-printing control characters) or a number contains a digit not allowed in the specified base (e.g., 8 or 9 in an octal number or a letter in a hexadecimal number where the trailing H was omitted).
- 5) **INVALID OPERATOR**
Indicates use of an invalid operator in an expression. Occurs when an operator such as .AND. or .XOR. is misspelled or contains illegal characters.
- 6) **SYNTAX ERROR**
Indicates the syntax error of the statement is invalid. Occurs when an expression is incorrectly formed, unmatched parentheses are found in an operand field, or a DEFM or DEFT string is either too long (greater than 63 characters) or contains unbalanced quotes.
- 7) **ASSEMBLER ERROR**
Indicates that the assembler has failed to process this instruction. Usually occurs when an expression is incorrectly formed.

- 8) UNDEFINED SYMBOL
Indicates that a symbol in an operand field is misspelled or not declared as a label for an instruction or pseudo-op, or was not declared EXTERNAL.
- 9) INVALID OPERAND COMBINATION
Indicates that the operand combination for this opcode is invalid. Occurs when a register name or condition code is misspelled or incorrectly used with the particular opcode.
- 10) EXPRESSION OUT OF RANGE
Indicates that the value of an expression is either too large or too small for the appropriate quantity. Occurs on 16-bit arithmetic overflow or division by zero in an expression, incrementing the reference counter beyond a 16-bit value, or trying to use a value which will not fit into a particular bit-field--typically a byte.
- 11) MULTIPLE DECLARATION
Indicates that an attempt was made to redefine a label. Occurs when a label is misspelled, or mistakenly used several times. The pseudo-op DEFL can be used to assign a value to a label which can then be redefined by another DEFL.
- 12) MACRO DEFINITION ERROR
Indicates that a macro is incorrectly defined. Occurs when the M option is not specified but macros are used, when a macro is defined within another macro definition, when the parameters are not correctly specified, or an unrecognized parameter is found in the macro body.
- 13) UNBALANCED QUOTES
Indicates that a string is not properly bounded by single quote marks, or quote marks inside a string are not properly matched in pairs.
- 14) ASSEMBLER COMMAND ERROR
Indicates that an assembler command is not recognized or correctly formed. The command must begin with an asterisk (*) in column one, the first letter identifies the command, and any parameters such as 'ON', 'OFF' or a filename must be properly delimited. The command will be ignored.

- 15) MACRO EXPANSION ERROR
Indicates that the expansion of a single line in a macro has overflowed the expansion buffer. Occurs when substitution of a parameter causes the line to increase in length beyond the capacity of the buffer (currently 128 bytes). The line will be truncated.
- 16) MACRO STACK OVERFLOW
Indicates that the depth of nesting of macro calls has exceeded the macro parameter stack buffer capacity. Occurs when the sum of the parameter string lengths (plus some additional information for each macro call) is longer than the buffer (currently 256 bytes), which often happens if infinitely recursive macro calls are used. The macro call which caused the error will be ignored.
- 17) INCLUDE NESTED TOO DEEP
Indicates that an *Include command was found which would have caused a nesting of included source files to a depth greater than four, where the original source file is considered to be level one. The command will be ignored.
- 18) GLOBAL DEFINITION ERROR
Indicates that either a label was present on a GLOBAL pseudo-op statement, or there was an attempt to give an absolute value to a GLOBAL symbol in a relocatable module. The latter case is not allowed since all GLOBALs in a relocatable module will be relocated by the Linker. May occur either after a GLOBAL pseudo-op or after an EQU or DEFL statement which is attempting to absolutize a relocatable GLOBAL symbol.
- 19) EXTERNAL DEFINITION ERROR
Indicates that either a label was present on an EXTERNAL pseudo-op statement, or there was an attempt to declare a symbol to be EXTERNAL which had previously been defined within the module to have an absolute value. May occur due to a misspelling or other oversight.
- 20) NAME DECLARED GLOBAL AND EXTERNAL
Indicates that the name was found in both a GLOBAL pseudo-op and an EXTERNAL pseudo-op, which is contradictory. May occur due to a misspelling or other oversight.

- 21) LABEL DECLARED AS EXTERNAL
Indicates that a name has been declared in both an EXTERNAL pseudo-op and as a label in this module. May occur due to a misspelling or other oversight.
- 22) INVALID EXTERNAL EXPRESSION
Indicates that a symbol name which has been declared in an EXTERNAL pseudo-op is improperly used in an expression. May occur when invalid arithmetic operators are applied to an external expression or when the mode of an operand must be either absolute or relocatable.
- 23) INVALID RELOCATABLE EXPRESSION
Indicates that an expression which contains a relocatable value (either a label or the reference counter symbol \$ in a relocatable module) is improperly formed or used. May occur when invalid arithmetic operators are applied to a relocatable expression or when the mode of an operand must be absolute. Remember that all relocatable values (addresses) must be represented in 16 bits.
- 24) EXPRESSION MUST BE ABSOLUTE
Indicates that the mode of an expression is not absolute when it should be. May occur when a relocatable or external expression is used to specify a quantity that must be either constant or representable in less than 16 bits.
- 25) UNDEFINED GLOBAL(S)
Indicates that one or more symbols which were declared in a GLOBAL pseudo-op were never actually defined as a label in this module. A cross-reference listing may be helpful in determining which GLOBALS are undefined (value=FFFF).
- 26) WARNING - ORG IS RELOCATABLE
Indicates that an ORG statement was encountered in a relocatable module. This warning is issued to remind the user that the reference counter is set to a relocatable value, not an absolute one. May occur when the Absolute option is not specified for an absolute module. This warning may be suppressed by the NOW option.

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