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CP/M 2.0 ALTERATION GUIDE

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1. INTRODUCTION

The standard CP/M system assumes operation on an Intel MDS-800 microcomputer development system, but is designed so that the user can alter a specific set of subroutines which define the hardware operating environment. In this way, the user can produce a diskette which operates with any IBM-3741 format compatible drive controller and other peripheral devices.

Although standard CP/M 2.0 is configured for single density floppy disks, field-alteration features allow adaptation to a wide variety of disk subsystems from single drive minidisks through high-capacity "hard disk" systems. In order to simplify the following adaptation process, we assume that CP/M 2.0 will first be configured for single density floppy disks where minimal editing and debugging tools are available. If an earlier version of CP/M is available, the customizing process is eased considerably. In this latter case, you may wish to briefly review the system generation process, and skip to later sections which discuss system alteration for non-standard disk systems.

In order to achieve device independence, CP/M is separated into three distinct modules:

BIOS - basic I/O system which is environment dependent BDOS - basic disk operating system which is not dependent upon the hardware configuration CCP - the console command processor which uses the BDOS

Of these modules, only the BIOS is dependent upon the particular hardware. That is, the user can "patch" the distribution version of CP/M to provide a new BIOS which provides a customized interface between the remaining CP/M modules and the user's own hardware system. The purpose of this document is to provide a step-by-step procedure for patching your new BIOS into CP/M.

If CP/M is being tailored to your computer system for the first time, the new BIOS requires some relatively simple software development and testing. The standard BIOS is listed in Appendix B, and can be used as a model for the customized package. A skeletal version of the BIOS is given in Appendix C which can serve as the basis for a modified BIOS. In addition to the BIOS, the user must write a simple memory loader, called GETSYS, which brings the operating system into memory. In order to patch the new BIOS into CP/M, the user must write the reverse of GETSYS, called PUTSYS, which places an altered version of CP/M back onto the diskette. PUTSYS can be derived from GETSYS by changing the disk read commands into disk Sample skeletal GETSYS and PUTSYS programs are write commands. described in Section 3, and listed in Appendix D. In order to make the CP/M system work automatically, the user must also supply a cold start loader, similar to the one provided with CP/M (listed in Appendices A and B). A skeletal form of a cold start loader is given in Appendix E which can serve as a model for your loader.

2. FIRST LEVEL SYSTEM REGENERATION

The procedure to follow to patch the CP/M system is given below in several steps. Address references in each step are shown with a following "H" which denotes the hexadecimal radix, and are given for a 20K CP/M system. For larger CP/M systems, add a "bias" to each address which is shown with a "+b" following it, where b is equal to the memory size - 20K. Values for b in various standard memory sizes are

24K:b = 24K - 20K = 4K = 1000H32K:b = 32K - 20K = 12K = 3000H40K:b = 40K - 20K = 20K = 5000H48K:b = 48K - 20K = 28K = 7000H56K:b = 56K - 20K = 36K = 9000H62K:b = 62K - 20K = 42K = A800H64K:b = 64K - 20K = 44K = B000H

Note: The standard distribution version of CP/M is set for operation within a 20K memory system. Therefore, you must first bring up the 20K CP/M system, and then configure it for your actual memory size (see Second Level System Generation).

(1) Review Section 4 and write a GETSYS program which reads the first two tracks of a diskette into memory. The data from the diskette must begin at location 3380H. Code GETSYS so that it starts at location 100H (base of the TPA), as shown in the first part of Appendix d.

(2) Test the GETSYS program by reading a blank diskette into memory, and check to see that the data has been read properly, and that the diskette has not been altered in any way by the GETSYS program.

(3) Run the GETSYS program using an initialized CP/M diskette to see if GETSYS loads CP/M starting at 3380H (the operating system actually starts 128 bytes later at 3400H).

(4) Review Section 4 and write the PUTSYS program which writes memory starting at 3380H back onto the first two tracks of the diskette. The PUTSYS program should be located at 200H, as shown in the second part of Appendix D.

(5) Test the PUTSYS program using a blank uninitialized diskette by writing a portion of memory to the first two tracks; clear memory and read it back using GETSYS. Test PUTSYS completely, since this program will be used to alter CP/M on disk.

(6) Study Sections 5, 6, and 7, along with the distribution version of the BIOS given in Appendix B, and write a simple version which performs a similar function for the customized environment. Use the program given in Appendix C as a model. Call this new BIOS by the name CBIOS (customized BIOS). Implement only the primitive disk operations on a single drive, and simple console input/output functions in this phase.

(7) Test CBIOS completely to ensure that it properly performs console character I/O and disk reads and writes. Be especially careful to ensure that no disk write operations occur accidently during read operations, and check that the proper track and sectors are addressed on all reads and writes. Failure to make these checks may cause destruction of the initialized CP/M system after it is patched.

(8) Referring to Figure 1 in Section 5, note that the BIOS is placed between locations 4A00H and 4FFFH. Read the CP/M system using GETSYS and replace the BIOS segment by the new CBIOS developed in step (6) and tested in step (7). This replacement is done in the memory of the machine, and will be placed on the diskette in the next step.

(9) Use PUTSYS to place the patched memory image of CP/M onto the first two tracks of a blank diskette for testing.

(10) Use GETSYS to bring the copied memory image from the test diskette back into memory at 3380H, and check to ensure that it has loaded back properly (clear memory, if possible, before the load). Upon successful load, branch to the cold start code at location 4A00H. The cold start routine will initialize page zero, then jump to the CCP at location 3400H which will call the BDOS, which will call the CBIOS. The CBIOS will be asked by the CCP to read sixteen sectors on track 2, and if successful, CP/M will type "A>", the system prompt.

When you make it this far, you are almost on the air. If you have trouble, use whatever debug facilities you have available to trace and breakpoint your CBIOS.

(11) Upon completion of step (10), CP/M has prompted the console for a command input. Test the disk write operation by typing

SAVE 1 X.COM

(recall that all commands must be followed by a carriage return).

CP/M should respond with another prompt (after several disk accesses):

A>

If it does not, debug your disk write functions and retry.

(12) Then test the directory command by typing

DIR

CP/M should respond with

A: X COM

(13) Test the erase command by typing

ERA X.COM

CP/M should respond with the A prompt. When you make it this far, you should have an operational system which will only require a bootstrap loader to function completely.

(14) Write a bootstrap loader which is similar to GETSYS, and place it on track \emptyset , sector 1 using PUTSYS (again using the test diskette, not the distribution diskette). See Sections 5 and 8 for more information on the bootstrap operation.

(15) Retest the new test diskette with the bootstrap loader installed by executing steps (11), (12), and (13). Upon completion of these tests, type a control-C (control and C keys simultaneously). The system should then execute a "warm start" which reboots the system, and types the A prompt.

(16) At this point, you probably have a good version of your customized CP/M system on your test diskette. Use GETSYS to load CP/M from your test diskette. Remove the test diskette, place the distribution diskette (or a legal copy) into the drive, and use PUTSYS to replace the distribution version by your customized version. Do not make this replacement if you are unsure of your patch since this step destroys the system which was sent to you from Digital Research.

(17) Load your modified CP/M system and test it by typing

DIR

CP/M should respond with a list of files which are provided on the initialized diskette. One such file should be the memory image for the debugger, called DDT.COM.

NOTE: from now on, it is important that you always reboot the CP/M system (ctl-C is sufficient) when the diskette is removed and replaced by another diskette, unless the new diskette is to be read only.

(18) Load and test the debugger by typing

DDT

(see the document "CP/M Dynamic Debugging Tool (DDT)" for operating procedures. You should take the time to become familiar with DDT, it will be your best friend in later steps.

(19) Before making further CBIOS modifications, practice using the editor (see the ED user's guide), and assembler (see the ASM user's guide). Then recode and test the GETSYS, PUTSYS, and CBIOS programs using ED, ASM, and DDT. Code and test a COPY program which does a sector-to-sector copy from one diskette to another to obtain back-up copies of the original diskette (NOTE: read your CP/M Licensing Agreement; it specifies your legal responsibilities when copying the CP/M system). Place the copyright notice

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on each copy which is made with your COPY program.

(20) Modify your CBIOS to include the extra functions for punches, readers, signon messages, and so-forth, and add the facilities for a additional disk drives, if desired. You can make these changes with the GETSYS and PUTSYS programs which you have developed, or you can refer to the following section, which outlines CP/M facilities which will aid you in the regeneration process.

You now have a good copy of the customized CP/M system. Note that although the CBIOS portion of CP/M which you have developed belongs to you, the modified version of CP/M which you have created can be copied for your use only (again, read your Licensing Agreement), and cannot be legally copied for anyone else's use.

It should be noted that your system remains file-compatible with all other CP/M systems, (assuming media compatiblity, of course) which allows transfer of non-proprietary software between users of CP/M.

3. SECOND LEVEL SYSTEM GENERATION

Now that you have the CP/M system running, you will want to configure CP/M for your memory size. In general, you will first get a memory image of CP/M with the "MOVCPM" program (system relocator) and place this memory image into a named disk file. The disk file can then be loaded, examined, patched, and replaced using the debugger, and system generation program. For further details on the operation of these programs, see the "Guide to CP/M Features and Facilities" manual.

Your CBIOS and BOOT can be modified using ED, and assembled using ASM, producing files called CBIOS.HEX and BOOT.HEX, which contain the machine code for CBIOS and BOOT in Intel hex format.

To get the memory image of CP/M into the TPA configured for the desired memory size, give the command:

MOVCPM xx *

where "xx" is the memory size in decimal K bytes (e.g., 32 for 32K). The response will be:

> CONSTRUCTING xxK CP/M VERS 2.0 READY FOR "SYSGEN" OR "SAVE 34 CPMxx.COM"

At this point, an image of a CP/M in the TPA configured for the requested memory size. The memory image is at location 0900H through 227FH. (i.e., The BOOT is at 0900H, the CCP is at 980H, the BDOS starts at 1180H, and the BIOS is at 1F80H.) Note that the memory image has the standard MDS-800 BIOS and BOOT on it. It is now necessary to save the memory image in a file so that you can patch your CBIOS and CBOOT into it:

SAVE 34 CPMxx.COM

The memory image created by the "MOVCPM" program is offset by a negative bias so that it loads into the free area of the TPA, and thus does not interfere with the operation of CP/M in higher memory. This memory image can be subsequently loaded under DDT and examined or changed in preparation for a new generation of the system. DDT is loaded with the memory image by typing:

DDT CPMxx.COM Load DDT, then read the CPM image

DDT should respond with

NEXT PC 2300 0100

(The DDT prompt)

You can then use the display and disassembly commands to examine

portions of the memory image between 900H and 227FH. Note, however, that to find any particular address within the memory image, you must apply the negative bias to the CP/M address to find the actual address. Track 00, sector 01 is loaded to location 900H (you should find the cold start loader at 900H to 97FH), track 00, sector 02 is loaded into 980H (this is the base of the CCP), and so-forth through the entire CP/M system load. In a 20K system, for example, the CCP resides at the CP/M address 3400H, but is placed into memory at 980H by the SYSGEN program. Thus, the negative bias, denoted by n, satisfies

3400H + n = 980H, or n = 980H - 3400H

Assuming two's complement arithmetic, n = D580H, which can be checked by

3400H + D580H = 10980H = 0980H (ignoring high-order overflow).

Note that for larger systems, n satisfies

(3400H+b) + n = 980H, or n = 980H - (3400H + b), or n = D580H - b.

The value of n for common CP/M systems is given below

memory size	bias b	negative offset n
20K	ØØØØH	D580H - 0000H = D580H
24K	1000H	D580H - 1000H = C580H
32K	ЗФФйн	D580H - 3000H = A580H
4 Ø K	5000H	D580H - 5000H = 8580H
4 8K	7000H	D580H - 7000H = 6580H
5 6K	9000H	D580H - 9000H = 4580H
62K	АЗЙЙН	D580H - A800H = 2D80H
64K	вøøйн	D580H - B000H = 2580H

Assume, for example, that you want to locate the address x within the memory image loaded under DDT in a 20K system. First type

Hx,n Hezadecimal sum and difference

and DDT will respond with the value of x+n (sum) and x-n (difference). The first number printed by DDT will be the actual memory address in the image where the data or code will be found. The input

H3400,D580

for example, will produce 980H as the sum, which is where the CCP is located in the memory image under DDT.

Use the L command to disassemble portions the BIOS located at (4A00H+b)-n which, when you use the H command, produces an actual address of 1F80H. The disassembly command would thus be

It is now necessary to patch in your CBOOT and CBIOS routines. The BOOT resides at location Ø900H in the memory image. If the actual load address is "n", then to calculate the bias (m) use the command:

H900,n Subtract load address from target address.

The second number typed in response to the command is the desired bias (m). For example, if your BOOT executes at 0080H, the command:

H900,80

will reply

0980 0880

88Ø Sum and difference in hex.

Therefore, the bias "m" would be 0880H. To read-in the BOOT, give the command:

ICBOOT.HEX Input file CBOOT.HEX

Then:

Rm

Read CBOOT with a bias of m (=900H-n)

You may now examine your CBOOT with:

L900

We are now ready to replace the CBIOS. Examine the area at 1F80H where the original version of the CBIOS resides. Then type

ICBIOS.HEX Ready the "hex" file for loading

assume that your CBIOS is being integrated into a 20K CP/M system, and thus is origined at location 4A00H. In order to properly locate the CBIOS in the memory image under DDT, we must apply the negative bias n for a 20K system when loading the hex file. This is accomplished by typing

RD580 Read the file with bias D580H

Upon completion of the read, re-examine the area where the CBIOS has been loaded (use an "LlF80" command), to ensure that is was loaded properly. When you are satisfied that the change has been made, return from DDT using a control-C or "G0" command.

Now use SYSGEN to replace the patched memory image back onto a diskette (use a test diskette until you are sure of your patch), as shown in the following interaction

	Start the SYSGEN program
SYSGEN VERSION 2.0 S	Sign-on message from SYSGEN
SOURCE DRIVE NAME (OR RE	ETURN TO SKIP)
	Respond with a carriage return
	to skip the CP/M read operation
S	since the system is already in
n	nemory.
DESTINATION DRIVE NAME	4 ⁻
	Respond with "B" to write the
r	new system to the diskette in
	drive [®] B.
DESTINATION ON B, THEN 1	
I	Place a scratch diskette in
ċ	lrive B, then type return.
FUNCTION COMPLETE	
DESTINATION DRIVE NAME	(OR RETURN TO REBOOT)

Place the scratch diskette in your drive A, and then perform a coldstart to bring up the new CP/M system you have configured.

Test the new CP/M system, and place the Digital Research copyright notice on the diskette, as specified in your Licensing Agreement:

Copyright (c), 1979 Digital Research The following program provides a framework for the GETSYS and PUTSYS programs referenced in Section 2. The READSEC and WRITESEC subroutines must be inserted by the user to read and write the specific sectors.

> GETSYS PROGRAM - READ TRACKS Ø AND 1 TO MEMORY AT 3380H ; REGISTER USE ; (SCRATCH REGISTER) Α ; В TRACK COUNT (0, 1) ; С SECTOR COUNT (1,2,...,26) ; (SCRATCH REGISTER PAIR) DE ; HL LOAD ADDRESS : SP SET TO STACK ADDRESS ; START: LXI SP,3380H ;SET STACK POINTER TO SCRATCH AREA н, ЗЗ80н ;SET BASE LOAD ADDRESS LXI MVI B, Ø START WITH TRACK Ø ; READ NEXT TRACK (INITIALLY Ø) RDTRK: C,1 MVI ; READ STARTING WITH SECTOR 1 ;READ NEXT SECTOR RDSEC: CALL READSEC ;USER-SUPPLIED SUBROUTINE ;MOVE LOAD ADDRESS TO NEXT 1/2 PAGE ;HL = HL + 128 D,128 LXI DAD D ;SECTOR = SECTOR + 1 INR C A,C MOV ;CHECK FOR END OF TRACK CPI 27 JC RDSEC ;CARRY GENERATED IF SECTOR < 27 ; ARRIVE HERE AT END OF TRACK, MOVE TO NEXT TRACK ; INR В MOV A,B ;TEST FOR LAST TRACK CPI 2 JC ;CARRY GENERATED IF TRACK < 2 RDTRK ; ARRIVE HERE AT END OF LOAD, HALT FOR NOW ; HLT ; USER-SUPPLIED SUBROUTINE TO READ THE DISK ; **READSEC:** ENTER WITH TRACK NUMBER IN REGISTER B, ; SECTOR NUMBER IN REGISTER C. AND ; ADDRESS TO FILL IN HL ; ; PUSH B ;SAVE B AND C REGISTERS PUSH H ;SAVE HL REGISTERS perform disk read at this point, branch to label START if an error occurs POP H ;RECOVER HL POP B ;RECOVER B AND C REGISTERS RET ;BACK TO MAIN PROGRAM END START

Note that this program is assembled and listed in Appendix C for reference purposes, with an assumed origin of 100H. The hexadecimal operation codes which are listed on the left may be useful if the program has to be entered through your machine's front panel switches.

The PUTSYS program can be constructed from GETSYS by changing only a few operations in the GETSYS program given above, as shown in Appendix D. The register pair HL become the dump address (next address to write), and operations upon these registers do not change within the program. The READSEC subroutine is replaced by a WRITESEC subroutine which performs the opposite function: data from address HL is written to the track given by register B and sector given by register C. It is often useful to combine GETSYS and PUTSYS into a single program during the test and development phase, as shown in the Appendix.

5. DISKETTE ORGANIZATION

The sector allocation for the standard distribution version of CP/M is given here for reference purposes. The first sector (see table on the following page) contains an optional software boot section. Disk controllers are often set up to bring track \emptyset , sector 1 into memory at a specific location (often location $\emptyset 0 \emptyset 0 H$). The program in this sector, called BOOT, has the responsibility of bringing the remaining sectors into memory starting at location 3400H+b. If your controller does not have a built-in sector load, you can ignore the program in track \emptyset , sector 1, and begin the load from track \emptyset sector 2 to location 3400H+b.

As an example, the Intel MDS-800 hardware cold start loader brings track 0, sector 1 into absolute address 3000H. Upon loading this sector, control transfers to location 3000H, where the bootstrap operation commences by loading the remainder of tracks 0, and all of track 1 into memory, starting at 3400H+b. The user should note that this bootstrap loader is of little use in a non-MDS environment, although it is useful to examine it since some of the boot actions will have to be duplicated in your cold start loader.

Track#	Sector#	Page#	Memory Address	CP/M Module name
ØØ	01		(boot address)	Cold Start Loader
ØØ	Ø2	ØØ	34ØØH+b	ССР
16	03	"	3480H+b	16
	Ø 4	Ø1 "	3500H+b	
	Ø 5 Ø 6	Ø 2	3580H+b	14
18	Ø 7	02	3600н+6 3680н+6	
	Ø8	Ø3	3700H+b	
11	Ø 9		378ØH+b	
+4	10	Ø 4	3800H+b	
84	11		388ØH+b	
67	12	Ø5	3900H+b	
14	13		398ØH+b	
41	14	ØG	3AØØH+b	
••	15		3A8ØH+b	
••	16	Ø7	3BØØH+b	
00	17		3B8ØH+b	ССР
ØØ	18	Ø 8	3СØØН+b	BDOS
"	19	ed	3C8ØH+b	"
	20	Ø 9	3DØØH+b	**
18	21	"	3D8ØH+b	
••	22	10	3EØØH+b	10
	23		3E8ØH+b	**
	24	11	3FØØH+b	••
11	25		3F8ØH+b	**
Ø1	26 Ø1	12	4000H+b	
ы Ю Т	Ø 2	13	4080н+ь 4100н+ь	••
14	Ø3		4180H+b	**
	Ø4	14	4200H+b	94
	ø5		4280H+b	11
**	ø6	15	4300H+b	
•1	Ø7		438ØH+b	
**	08	16	4400H+b	**
14	Ø 9	**	448ØH+b	*1
e4	10	17	4500H+b	14
	11) r	458ØH+b	"
	12	18	4600H+b	**
•1	13	44	468ØH+b	**
•1	14	19	4700H+b	49
18 18	15	**	478ØH+b	84
	16	20	4800H+b	44
	17		488ØH+b	•4
Ø1	18	21	4900H+b	
	19		4980H+b	BDOS
Ø1	20	22	4AØØH+b	BIOS
"	21		4A8ØH+b	•
	23	23	4BØØH+b	84 44
14	24 25		4B8ØH+b	42 54
Øl	25	24	4C00H+b	
			4C8ØH+b	BIOS
02-76	Ø1-26			(directory and data)

(All Information Contained Herein is Proprietary to Digital Research.)

6. THE BIOS ENTRY POINTS

The entry points into the BIOS from the cold start loader and BDOS are detailed below. Entry to the BIOS is through a "jump vector" located at 4A00H+b, as shown below (see Appendices B and C, as well). The jump vector is a sequence of 17 jump instructions which send program control to the individual BIOS subroutines. The BIOS subroutines may be empty for certain functions (i.e., they may contain a single RET operation) during regeneration of CP/M, but the entries must be present in the jump vector.

The jump vector at 4A00H+b takes the form shown below, where the individual jump addresses are given to the left:

4A00H+b	JMP	BOOT	;	ARRIVE HERE FROM COLD START LOAD
4AØ3H+b	JMP	WBOOT	;	ARRIVE HERE FOR WARM START
4A06H+b	JMP	CONST	;	CHECK FOR CONSOLE CHAR READY
4AØ9H+b	JMP	CONIN	;	READ CONSOLE CHARACTER IN
4AØCH+b	JMP	CONOUT	;	WRITE CONSOLE CHARACTER OUT
4A0FH+b	JMP	LIST	;	WRITE LISTING CHARACTER OUT
4A12H+b	JMP	PUNCH	;	WRITE CHARACTER TO PUNCH DEVICE
4A15H+b	JMP	READER	;	READ READER DEVICE
4A18H+b	JMP	HOME	;	MOVE TO TRACK ØØ ON SELECTED DISK
4AlBH+5	JMP	SELDSK	;	SELECT DISK DRIVE
4AlEH+o	JMP	SETTRK	;	SET TRACK NUMBER
4A21H+b	JMP	SETSEC	;	SET SECTOR NUMBER
4A24H+b	JMP	SETDMA	;	SET DMA ADDRESS
4A27H+b	JMP	READ	;	READ SELECTED SECTOR
4A2AH+b	JMP	WRITE	;	WRITE SELECTED SECTOR
4A2DH+b	JMP	LISTST	;	RETURN LIST STATUS
4A3ØH+b	JMP	SECTRAN	;	SECTOR TRANSLATE SUBROUTINE

Each jump address corresponds to a particular subroutine which performs the specific function, as outlined below. There are three major divisions in the jump table: the system (re)initialization which results from calls on BOOT and WBOOT, simple character I/O performed by calls on CONST, CONIN, CONOUT, LIST, PUNCH, READER, and LISTST, and diskette I/O performed by calls on HOME, SELDSK, SETTRK, SETSEC, SETDMA, READ, WRITE, and SECTRAN.

All simple character I/O operations are assumed to be performed in ASCII, upper and lower case, with high order (parity bit) set to zero. An end-of-file condition for an input device is given by an ASCII control-z (IAH). Peripheral devices are seen by CP/M as "logical" devices, and are assigned to physical devices within the BIOS.

In order to operate, the BDOS needs only the CONST, CONIN, and CONOUT subroutines (LIST, PUNCH, and READER may be used by PIP, but not the BDOS). Further, the LISTST entry is used currently only by DESPOOL, and thus, the initial version of CBIOS may have empty subroutines for the remaining ASCII devices.

The characteristics of each device are

- CONSOLE The principal interactive console which communicates with the operator, accessed through CONST, CONIN, and CONOUT. Typically, the CONSOLE is a device such as a CRT or Teletype.
- LIST The principal listing device, if it exists on your system, which is usually a hard-copy device, such as a printer or Teletype.
- PUNCH The principal tape punching device, if it exists, which is normally a high-speed paper tape punch or Teletype.
- READER The principal tape reading device, such as a simple optical reader or Teletype.

Note that a single peripheral can be assigned as the LIST, PUNCH, and READER device simultaneously. If no peripheral device is assigned as the LIST, PUNCH, or READER device, the CBIOS created by the user may give an appropriate error message so that the system does not "hang" if the device is accessed by PIP or some other user program. Alternately, the PUNCH and LIST routines can just simply return, and the READER routine can return with a IAH (ctl-Z) in reg A to indicate immediate end-of-file.

For added flexibility, the user can optionally implement the "IOBYTE" function which allows reassignment of physical and logical devices. The function creates a mapping of logical to IOBYTE physical devices which can be altered during CP/M processing (see the STAT command). The definition of the IOBYTE function corresponds to the Intel standard as follows: a single location in memory (currently location 0003H) is maintained, called IOBYTE, which defines the logical to physical device mapping which is effect at a particular time. The mapping is in performed by splitting the IOBYTE into four distinct fields of two bits each, called the CONSOLE, READER, PUNCH, and LIST fields, as shown below:

		most signi	ficant	least	significant
IOBYTE AT	ØØØ3H	LIST	PUNCH	READER	CONSOLE
		bits 6,7	bits 4,5	bits 2,3	bits Ø,1

The value in each field can be in the range $\emptyset-3$, defining the assigned source or destination of each logical device. The values which can be assigned to each field are given below

CONSOLE field (bits 0,1) Ø - console is assigned to the console printer device (TTY:) 1 - console is assigned to the CRT device (CRT:) 2 - batch mode: use the READER as the CONSOLE input, and the LIST device as the CONSOLE output (BAT:) 3 - user defined console device (UC1:) READER field (bits 2.3) Ø - READER is the Teletype device (TTY:) 1 - READER is the high-speed reader device (RDR:) - user defined reader # 1 (UR1:) 2 3 - user defined reader # 2 (UR2:) PUNCH field (bits 4,5) Ø - PUNCH is the Teletype device (TTY:) 1 - PUNCH is the high speed punch device (PUN:) 2 - user defined punch # 1 (UP1:) 3 - user defined punch # 2 (UP2:) LIST field (bits 6,7) \emptyset - LIST is the Teletype device (TTY:) 1 - LIST is the CRT device (CRT:) 2 - LIST is the line printer device (LPT:) 3 - user defined list device (UL1:)

> Note again that the implementation of the IOBYTE is optional, and affects only the organization of your CBIOS. NO CP/M systems use the IOBYTE (although they tolerate the existence of the IOBYTE at location $\emptyset \emptyset \emptyset \Im H$), except for PIP which allows access to the and STAT which allows physical devices. logical-physical assignments to be made and/or displayed (for more information, see the "CP/M Features and Facilities Guide"). In any case, the IOBYTE implementation should be omitted until your basic CBIOS is fully implemented and tested; then add the IOBYTE to increase your facilities.

> Disk I/O is always performed through a sequence of calls on the various disk access subroutines which set up the disk number to access, the track and sector on a particular disk, and the direct memory access (DMA) address involved in the I/O operation. After all these parameters have been set up, a call is made to the READ or WRITE function to perform the actual I/O operation. Note that there is often a single call to SELDSK to select a disk drive, followed by a number of read or write operations to the selected disk before selecting another drive for subsequent operations. Similarly, there may be a single call to set the DMA address, followed by several calls which read or write from the selected DMA address before the DMA address is changed. The track and sector subroutines are always called before the READ or WRITE operations are performed.

Note that the READ and WRITE routines should perform several retries (10 is standard) before reporting the error condition to the BDOS. If the error condition is returned to the BDOS, it will report the error to the user. The HOME subroutine may or may not actually perform the track 00 seek, depending upon your controller characteristics; the important point is that track 00 has been selected for the next operation, and is often treated in exactly the same manner as SETTRK with a parameter of 00.

The exact responsibilites of each entry point subroutine are given below:

BOOT

The BOOT entry point gets control from the cold start loader and is responsible for basic system initialization, including sending a signon message (which can be omitted in the first version). If the IOBYTE function is implemented, it must be set at this point. The various system parameters which are set by the WBOOT entry point must be initialized, and control is transferred to the CCP at 3400H+b for further processing. Note that reg C must be set to zero to select drive A.

WBOOT The WBOOT entry point gets control when a warm start occurs. A warm start is performed whenever a user program branches to location 0000H, or when the CPU is reset from the front panel. The CP/M system must be loaded from the first two tracks of drive A up to, but not including, the BIOS (or CBIOS, if you have completed your patch). System parameters must be initialized as shown below:

> location 0,1,2 set to JMP WBOOT for warm starts (0000H: JMP 4A03H+b) location 3 set initial value of IOBYTE, if implemented in your CBIOS location 5,6,7 set to JMP BDOS, which is the primary entry point to CP/M for transient programs. (0005H: JMP 3C06H+b)

(see Section 9 for complete details of page zero use) Upon completion of the initialization, the WBOOT program must branch to the CCP at 3400H+b to (re)start the system. Upon entry to the CCP, register C is set to the drive to select after system initialization.

CONST Sample the status of the currently assigned console device and return ØFFH in register A if a character is ready to read, and ØØH in register A if no console characters are ready.

CONIN Read the next console character into register A, and

set the parity bit (high order bit) to zero. If no console character is ready, wait until a character is typed before returning.

- CONOUT Send the character from register C to the console output device. The character is in ASCII, with high order parity bit set to zero. You may want to include a time-out on a line feed or carriage return, if your console device requires some time interval at the end of the line (such as a TI Silent 700 terminal). You can, if you wish, filter out control characters which cause your console device to react in a strange way (a control-z causes the Lear Seigler terminal to clear the screen, for example).
- LIST Send the character from register C to the currently assigned listing device. The character is in ASCII with zero parity.
- PUNCH Send the character from register C to the currently assigned punch device. The character is in ASCII with zero parity.
- READER Read the next character from the currently assigned reader device into register A with zero parity (high order bit must be zero), an end of file condition is reported by returning an ASCII control-z (1AH).
- HOME Return the disk head of the currently selected disk (initially disk A) to the track ØØ position. If your controller allows access to the track Ø flag from the drive, step the head until the track Ø flag is detected. If your controller does not support this feature, you can translate the HOME call into a call on SETTRK with a parameter of Ø.
- Select the disk drive given by register C for further SELDSK operations, where register C contains Ø for drive A, 1 for drive B, and so-forth up to 15 for drive P (the standard CP/M distribution version supports four drives). On each disk select, SELDSK must return in HL the base address of a 16-byte area, called the Disk Parameter Header, described in the Section 10. For standard floppy disk drives, the contents of the header and associated tables does not change, and thus the program segment included in the sample CBIOS performs this operation automatically. If there is an attempt to select a non-existent drive, SELDSK returns HL=0000H as an error indicator. Although SELDSK must return the header address on each call, it is advisable to postpone the actual physical disk select operation until an I/O function (seek, read or write) is actually performed, since disk selects often occur without utimately performing any disk I/O, and many controllers will unload the head of the current disk

before selecting the new drive. This would cause an excessive amount of noise and disk wear.

- SETTRK Register BC contains the track number for subsequent disk accesses on the currently selected drive. You can choose to seek the selected track at this time, or delay the seek until the next read or write actually occurs. Register BC can take on values in the range \emptyset -76 corresponding to valid track numbers for standard floppy disk drives, and \emptyset -65535 for non-standard disk subsystems.
- SETSEC Register BC contains the sector number (1 through 26) for subsequent disk accesses on the currently selected drive. You can choose to send this information to the controller at this point, or instead delay sector selection until a read or write operation occurs.
- Register BC contains the DMA (disk memory access) SETDMA address for subsequent read or write operations. For example, if $B = \emptyset \emptyset H$ and $C = 8 \emptyset H$ when SETDMA is called, then all subsequent read operations read their data 80H through 0FFH, and all subsequent write into operations get their data from 80H through ØFFH, until the next call to SETDMA occurs. The initial DMA address is assumed to be 80H. Note that the controller need not actually support direct memory access. If, for example, all data is received and sent through I/O ports, the CBIOS which you construct will use the 128 byte area starting at the selected DMA address for the memory buffer during the following read or write operations.
- READ Assuming the drive has been selected, the track has been set, the sector has been set, and the DMA address has been specified, the READ subroutine attempts to read one sector based upon these parameters, and returns the following error codes in register A:

Ø no errors occurred

1 non-recoverable error condition occurred

Currently, CP/M responds only to a zero or non-zero value as the return code. That is, if the value in register A is Ø then CP/M assumes that the disk operation completed properly. If an error occurs, however, the CBIOS should attempt at least 1Ø retries to see if the error is recoverable. When an error is reported the BDOS will print the message "BDOS ERR ON x: BAD SECTOR". The operator then has the option of typing <cr>> to ignore the error, or ctl-C to abort.

WRITE Write the data from the currently selected DMA address to the currently selected drive, track, and sector. The data should be marked as "non deleted data" to

maintain compatibility with other CP/M systems. The error codes given in the READ command are returned in register A, with error recovery attempts as described above.

- LISTST Return the ready status of the list device. Used by the DESPOOL program to improve console response during its operation. The value ØØ is returned in A if the list device is not ready to accept a character, and ØFFH if a character can be sent to the printer. Note that a ØØ value always suffices.
- SECTRAN Performs sector logical to physical sector translation in order to improve the overall response of CP/M. Standard CP/M systems are shipped with a "skew factor" of 6, where six physical sectors are skipped between each logical read operation. This skew factor allows enough time between sectors for most programs to load their buffers without missing the next sector. In particular computer systems which use fast processors. memory, and disk subsystems, the skew factor may be changed to improve overall response. Note, however, that you should maintain a single density IBM compatible version of CP/M for information transfer into and out of your computer system, using a skew factor of 6. In general, SECTRAN receives a logical sector number in BC, and a translate table address in The sector number is used as an index into the DE. translate table, with the resulting physical sector number in HL. For standard systems, the tables and indexing code is provided in the CBIOS and need not be changed.

(All Information Contained Herein is Proprietary to Digital Research.)

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7. A SAMPLE BIOS

The program shown in Appendix C can serve as a basis for your first BIOS. The simplest functions are assumed in this BIOS, so that you can enter it through the front panel, if absolutely necessary. Note that the user must alter and insert code into the subroutines for CONST, CONIN, CONOUT, READ, WRITE, and WAITIO subroutines. Storage is reserved for user-supplied code in these regions. The scratch area reserved in page zero (see Section 9) for the BIOS is used in this program, so that it could be implemented in ROM, if desired.

Once operational, this skeletal version can be enhanced to print the initial sign-on message and perform better error recovery. The subroutines for LIST, PUNCH, and READER can be filled-out, and the IOBYTE function can be implemented.

8. A SAMPLE COLD START LOADER

The program shown in Appendix D can serve as a basis for your cold start loader. The disk read function must be supplied by the user. and the program must be loaded somehow starting at location 0000. Note that space is reserved for your patch so that the total amount of storage required for the cold start loader is 128 bytes. Eventually, you will probably want to get this loader onto the first disk sector (track Ø, sector 1), and cause your controller to load it into memory automatically upon system start-up. Alternatively, you may wish to place the cold start loader into ROM, and place it above the CP/M In this case, it will be necessary to originate the program system. at a higher address, and key-in a jump instruction at system start-up which branches to the loader. Subsequent warm starts will not require this key-in operation, since the entry point 'WBOOT' gets control, thus bringing the system in from disk automatically. Note also that the skeletal cold start loader has minimal error recovery, which may be enhanced on later versions.

9. RESERVED LOCATIONS IN PAGE ZERO

Main memory page zero, between locations ØØH and ØFFH, contains several segments of code and data which are used during CP/M processing. The code and data areas are given below for reference purposes.

Locations from to	Contents
0000H - 0002H	Contains a jump instruction to the warm start entry point at location 4A03H+b. This allows a simple programmed restart (JMP 0000H) or manual restart from the front panel.
0003H - 0003H	Contains the Intel standard IOBYTE, which is optionally included in the user's CBIOS, as described in Section 6.
0004H - 0004H	Current default drive number ($\emptyset=A,\ldots,15=P$).
0005H — 0007H	Contains a jump instruction to the BDOS, and serves two purposes: JMP 0005H provides the primary entry point to the BDOS, as described in the manual "CP/M Interface Guide," and LHLD 0006H brings the address field of the instruction to the HL register pair. This value is the lowest address in memory used by CP/M (assuming the CCP is being overlayed). Note that the DDT program will change the address field to reflect the reduced memory size in debug mode.
ØØØ8H - ØØ27H	(interrupt locations 1 through 5 not used)
ØØ3ØH - ØØ37H	(interrupt location 6, not currently used - reserved)
ØØ38H - ØØ3AH	Restart 7 - Contains a jump instruction into the DDT or SID program when running in debug mode for programmed breakpoints, but is not otherwise used by CP/M.
003BH - 003FH	(not currently used - reserved)
ØØ40H - ØØ4FH	<pre>16 byte area reserved for scratch by CBIOS, but is not used for any purpose in the distribution version of CP/M</pre>
0050H - 005BH	(not currently used - reserved)
005CH - 007CH	default file control block produced for a transient program by the Console Command Processor.
007DH - 007FH	Optional default random record position

ØØ8ØH - ØØFFH default 128 byte disk buffer (also filled with the command line when a transient is loaded under the CCP).

Note that this information is set-up for normal operation under the CP/M system, but can be overwritten by a transient program if the BDOS facilities are not required by the transient.

If, for example, a particular program performs only simple I/O and must begin execution at location 0, it can be first loaded into the TPA, using normal CP/M facilities, with a small memory move program which gets control when loaded (the memory move program must get control from location 0100H, which is the assumed beginning of all transient programs). The move program can then proceed to move the entire memory image down to location 0, and pass control to the starting address of the memory load. Note that if the BIOS is overwritten, or if location 0 (containing the warm start entry point) is overwritten, then the programmer must bring the CP/M system back into memory with a cold start sequence.

10. DISK PARAMETER TABLES.

Tables are included in the BIOS which describe the particular characteristics of the disk subsystem used with CP/M. These tables can be either hand-coded, as shown in the sample CBIOS in Appendix C, or automatically generated using the DISKDEF macro library, as shown in Appendix B. The purpose here is to describe the elements of these tables.

In general, each disk drive has an associated (16-byte) disk parameter header which both contains information about the disk drive and provides a scratchpad area for certain BDOS operations. The format of the disk parameter header for each drive is shown below

					2	ameter						
	XLT			 	 	DIRBUF			CSV		ALV	
	16b		16b	 16b	16b	16b	16b		16b		16b	

where each element is a word (16-bit) value. The meaning of each Disk Parameter Header (DPH) element is

- XLT Address of the logical to physical translation vector, if used for this particular drive, or the value 0000H if no sector translation takes place (i.e, the physical and logical sector numbers are the same). Disk drives with identical sector skew factors share the same translate tables.
- 0000 Scratchpad values for use within the BDOS (initial value is unimportant).
- DIRBUF Address of a 128 byte scratchpad area for directory operations within BDOS. All DPH's address the same scratchpad area.
- DPB Address of a disk parameter block for this drive. Drives with identical disk characteristics address the same disk parameter block.
- CSV Address of a scratchpad area used for software check for changed disks. This address is different for each DPH.
- ALV Address of a scratchpad area used by the BDOS to keep disk storage allocation information. This address is different for each DPH.

Given n disk drives, the DPH's are arranged in a table whose first row of 16 bytes corresponds to drive \emptyset , with the last row corresponding to drive n-1. The table thus appears as

DPBASE:

ØØ	XLT	ØØI	0000		0000		ØØØØ	DIRBUF DBP	ØØ CSV	ØØ ALV Ø	Ø
Ø1	XLT	Ø1	0000		0000		ØØØØ	DIRBUF DBP	Ø1 CSV	Ø1 ALV Ø	1
	(and so-forth through)										
n-]	XLTr	n-1	0000		0000		0000	DIRBUF DBP1	n-1 CSVr	n-1 ALVn-	1

where the label DPBASE defines the base address of the DPH table.

A responsibility of the SELDSK subroutine is to return the base address of the DPH for the selected drive. The following sequence of operations returns the table address, with a 0000H returned if the selected drive does not exist.

NDISKS	EQU	4 ;NUMBE	ER OF DISK DRIVES
SELDSK:			
	;SELEC'	r disk giv	EN BY BC
	LXI	н,0000н	;ERROR CODE
	MOV		DRIVE OK?
	CPI		CY IF SO
	RNC		RET IF ERROR
	;NO ER	ROR, CONTI	INUE
	MOV		;LOW(DISK)
	MOV		;HIGH(DISK)
	DAD	Н	;*2
	DAD	Н	;*4
	DAD	Н	;*8
	DAD	Н	;*16
	LXI	D,DPBASE	FIRST DPH
	DAD	D	;DPH(DISK)
	RET		

The translation vectors (XLT ØØ through XLTn-1) are located elsewhere in the BIOS, and simply correspond one-for-one with the logical sector numbers zero through the sector count-1. The Disk Parameter Block (DPB) for each drive is more complex. A particular DPB, which is addressed by one or more DPH's, takes the general form

1					DSM		-			OFF	Ι
	16b	 8b	8b	8b	16b	 16b		8b	16b	 16b	

where each is a byte or word value, as shown by the "8b" or "16b" indicator below the field.

SPT is the total number of sectors per track

BSH is the data allocation block shift factor, determined by the data block allocation size.

- EXM is the extent mask, determined by the data block allocation size and the number of disk blocks.
- DSM determines the total storage capacity of the disk drive
- DRM determines the total number of directory entries which can be stored on this drive ALØ,AL1 determine reserved directory blocks.
- CKS is the size of the directory check vector
- OFF is the number of reserved tracks at the beginning of the (logical) disk.

The values of BSH and BLM determine (implicitly) the data allocation size BLS, which is not an entry in the disk parameter block. Given that the designer has selected a value for BLS, the values of BSH and BLM are shown in the table below

BLS	BSH	BLM
1,024	3	7
2,048	4	15
4,096	5	31
8,192	6	63
16,384	7	127

where all values are in decimal. The value of EXM depends upon both the BLS and whether the DSM value is less than 256 or greater than 255, as shown in the following table

BLS	DSM < 256	DSM > 255
1,024	Ø	N/A
2,048	1	Ø
4,096	3	1
8,192	7	3
16,384	15	7

The value of DSM is the maximum data block number supported by this particular drive, measured in BLS units. The product BLS times (DSM+1) is the total number of bytes held by the drive and, of course, must be within the capacity of the physical disk, not counting the reserved operating system tracks.

The DRM entry is the one less than the total number of directory entries, which can take on a 16-bit value. The values of ALØ and ALl, however, are determined by DRM. The two values ALØ and ALl can together be considered a string of 16-bits, as shown below.

	1 1110					l	AL1					-				
				1		l										
ø	ø								Ø8						 15	-

where position 00 corresponds to the high order bit of the byte labelled AL0, and 15 corresponds to the low order bit of the byte labelled AL1. Each bit position reserves a data block for number of directory entries, thus allowing a total of 16 data blocks to be assigned for directory entries (bits are assigned starting at 00 and filled to the right until position 15). Each directory entry occupies 32 bytes, resulting in the following table

BLS	Dire	ectory	Entries			
1,024	32	times	#	bits		
2,048	64	times	#	bits		
4,096	128	times	#	bits		
8,192	256	times	#	bits		
16,384	512	times	#	bits		

Thus, if DRM = 127 (128 directory entries), and BLS = 1024, then there are 32 directory entries per block, requiring 4 reserved blocks. In this case, the 4 high order bits of AL0 are set, resulting in the values AL0 = 0F0H and AL1 = 00H.

The CKS value is determined as follows: if the disk drive media is removable, then CKS = (DRM+1)/4, where DRM is the last directory entry number. If the media is fixed, then set CKS = Ø (no directory records are checked in this case).

Finally, the OFF field determines the number of tracks which are skipped at the beginning of the physical disk. This value is automatically added whenever SETTRK is called, and can be used as a mechanism for skipping reserved operating system tracks, or for partitioning a large disk into smaller segmented sections.

To complete the discussion of the DPB, recall that several DPH's can address the same DPB if their drive characteristics are identical. Further, the DPB can be dynamically changed when a new drive is addressed by simply changing the pointer in the DPH since the BDOS copies the DPB values to a local area whenever the SELDSK function is invoked.

Returning back to the DPH for a particular drive, note that the two address values CSV and ALV remain. Both addresses reference an area of uninitialized memory following the BIOS. The areas must be unique for each drive, and the size of each area is determined by the values in the DPB.

The size of the area addressed by CSV is CKS bytes, which is sufficient to hold the directory check information for this particular drive. If CKS = (DRM+1)/4, then you must reserve (DRM+1)/4 bytes for directory check use. If CKS = \emptyset , then no storage is reserved.

The size of the area addressed by ALV is determined by the maximum number of data blocks allowed for this particular disk, and is computed as (DSM/8)+1.

The CBIOS shown in Appendix C demonstrates an instance of these tables for standard 8" single density drives. It may be useful to examine this program, and compare the tabular values with the definitions given above.

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11. THE DISKDEF MACRO LIBRARY.

A macro library is shown in Appendix F, called DISKDEF, which greatly simplifies the table construction process. You must have access to the MAC macro assembler, of course, to use the DISKDEF facility, while the macro library is included with all CP/M 2.0 distribution disks.

A BIOS disk definition consists of the following sequence of macro statements:

MACLIB DISKDEF DISKS n DISKDEF Ø,... DISKDEF 1,... DISKDEF n-1 ENDEF

where the MACLIB statement loads the DISKDEF.LIB file (on the same disk as your BIOS) into MAC's internal tables. The DISKS macro call follows, which specifies the number of drives to be configured with your system, where n is an integer in the range 1 to 16. A series of DISKDEF macro calls then follow which define the characteristics of each logical disk, Ø through n-1 (corresponding to logical drives A through P). Note that the DISKS and DISKDEF macros generate the in-line fixed data tables described in the previous section, and thus must be placed in a non-executable portion of your BIOS, typically directly following the BIOS jump vector.

The remaining portion of your BIOS is defined following the DISKDEF macros, with the ENDEF macro call immediately preceding the END statement. The ENDEF (End of Diskdef) macro generates the necessary uninitialized RAM areas which are located in memory above your BIOS.

The form of the DISKDEF macro call is

DISKDEF dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0]

where

dn is the logical disk number, Ø to n-1 fsc is the first physical sector number (Ø or 1) lsc is the last sector number skf is the optional sector skew factor is the data allocation block size bls dir is the number of directory entries is the number of "checked" directory entries cks is the track offset to logical track ØØ ofs [Ø] is an optional 1.4 compatibility flag

The value "dn" is the drive number being defined with this DISKDEF
macro invocation. The "fsc" parameter accounts for differing sector numbering systems, and is usually Ø or 1. The "lsc" is the last numbered sector on a track. When present, the "skf" parameter defines the sector skew factor which is used to create a sector translation table according to the skew. If the number of sectors is less than 256, a single-byte table is created, otherwise each translation table element occupies two bytes. No translation table is created if the skf parameter is omitted (or equal to \emptyset). The "bls" parameter specifies the number of bytes allocated to each data block, and takes on the values 1024, 2048, 4096, 8192, or 16384. Generally, performance increases with larger data block sizes since there are fewer directory references and logically connected data records are physically close on the disk. Further, each directory entry addresses more data and the BIOS-resident ram space is reduced. The "dks" specifies the total disk size in "bls" units. That is, if the bls = 2048 and dks = 1000, then the total disk capacity is 2.048.000 bytes. If dks is greater than 255, then the block size parameter bls must be The value of "dir" is the total number of greater than 1024. directory entries which may exceed 255, if desired. "cks" The parameter determines the number of directory items to check on each directory scan, and is used internally to detect changed disks during system operation, where an intervening cold or warm start has not occurred (when this situation is detected, CP/M automatically marks the disk read/only so that data is not subsequently destroyed). As stated in the previous section, the value of cks = dir when the media is easily changed, as is the case with a floppy disk subsystem. Τf the disk is permanently mounted, then the value of cks is typically \emptyset , since the probability of changing disks without a restart is quite low. The "ofs" value determines the number of tracks to skip when this particular drive is addressed, which can be used to reserve additional operating system space or to simulate several logical drives on a single large capacity physical drive. Finally, the $[\emptyset]$ parameter is included when file compatibility is required with versions of 1.4 which have been modified for higher density disks. This parameter ensures that only 16K is allocated for each directory record, as was the case for previous versions. Normally, this parameter is not included.

For convenience and economy of table space, the special form

DISKDEF i,j

gives disk i the same characteristics as a previously defined drive j. A standard four-drive single density system, which is compatible with version 1.4, is defined using the following macro invocations:

DISKS 4 DISKDEF Ø,1,26,6,1024,243,64,64,2 DISKDEF 1,0 DISKDEF 2,0 DISKDEF 3,0 ENDEF

with all disks having the same parameter values of 26 sectors per track (numbered 1 through 26), with 6 sectors skipped between each access, 1024 bytes per data block, 243 data blocks for a total of 243k byte disk capacity, 64 checked directory entries, and two operating system tracks.

The DISKS macro generates n Disk Parameter Headers (DPH's), starting at the DPH table address DPBASE generated by the macro. Each disk header block contains sixteen bytes, as described above, and correspond one-for-one to each of the defined drives. In the four drive standard system, for example, the DISKS macro generates a table of the form:

DPBASE	EQU	Ş
DPEØ:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV0,ALV0
DPEl:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV1,ALV1
DPE2:	DW	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV2,ALV2
DPE3:	D₩	XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV3,ALV3

where the DPH labels are included for reference purposes to show the beginning table addresses for each drive Ø through 3. The values contained within the disk parameter header are described in detail in the previous section. The check and allocation vector addresses are generated by the ENDEF macro in the ram area following the BIOS code and tables.

Note that if the "skf" (skew factor) parameter is omitted (or equal to \emptyset), the translation table is omitted, and a $\emptyset\emptyset\emptyset\emptyset$ H value is inserted in the XLT position of the disk parameter header for the disk. In a subsequent call to perform the logical to physical translation, SECTRAN receives a translation table address of DE = $\emptyset\emptyset\emptyset\emptyset$ H, and simply returns the original logical sector from BC in the HL register pair. A translate table is constructed when the skf parameter is present, and the (non-zero) table address is placed into the corresponding DPH's. The table shown below, for example, is constructed when the standard skew factor skf = 6 is specified in the DISKDEF macro call:

XLTØ: DB 1,7,13,19,25,5,11,17,23,3,9,15,21 DB 2,8,14,20,26,6,12,18,24,4,10,16,22

Following the ENDEF macro call, a number of uninitialized data areas are defined. These data areas need not be a part of the BIOS which is loaded upon cold start, but must be available between the BIOS and the end of memory. The size of the uninitialized RAM area is determined by EQU statements generated by the ENDEF macro. For a standard four-drive system, the ENDEF macro might produce

4C72 =	BEGDAT EQU Ş
	(data areas)
$4DB\emptyset =$	ENDDAT EQU \$
Ø13C =	DATSIZ EQU Ş-BEGDAT

which indicates that uninitialized RAM begins at location 4C72H, ends at 4DB0H-1, and occupies 013CH bytes. You must ensure that these addresses are free for use after the system is loaded.

After modification, you can use the STAT program to check your drive characteristics, since STAT uses the disk parameter block to decode the drive information. The STAT command form

STAT d:DSK:

decodes the disk parameter block for drive d (d=A,...,P) and displays the values shown below:

r: 128 Byte Record Capacity
k: Kilobyte Drive Capacity
d: 32 Byte Directory Entries
c: Checked Directory Entries
e: Records/ Extent
b: Records/ Block
s: Sectors/ Track
t: Reserved Tracks

Three examples of DISKDEF macro invocations are shown below with corresponding STAT parameter values (the last produces a full 8-megabyte system).

DISKDEF 0,1,58,,2048,256,128,128,2 r=4096, k=512, d=128, c=128, e=256, b=16, s=58, t=2

DISKDEF 0,1,58,,2048,1024,300,0,2 r=16384, k=2048, d=300, c=0, e=128, b=16, s=58, t=2

DISKDEF Ø,1,58,,16384,512,128,128,2 r=65536, k=8192, d=128, c=128, e=1024, b=128, s=58, t=2

12. SECTOR BLOCKING AND DEBLOCKING.

Upon each call to the BIOS WRITE entry point, the CP/M BDOS includes information which allows effective sector blocking and deblocking where the host disk subsystem has a sector size which is a multiple of the basic 128-byte unit. The purpose here is to present a general-purpose algorithm which can be included within your BIOS which uses the BDOS information to perform the operations automatically.

Upon each call to WRITE, the BDOS provides the following information in register C:

Ø = normal sector write
1 = write to directory sector
2 = write to the first sector
of a new data block

Condition Ø occurs whenever the next write operation is into a previously written area, such as a random mode record update, when the write is to other than the first sector of an unallocated block, or when the write is not into the directory area. Condition 1 occurs when a write into the directory area is performed. Condition 2 occurs when the first record (only) of a newly allocated data block is written. In most cases, application programs read or write multiple 128 byte sectors in sequence, and thus there is little overhead involved in either operation when blocking and deblocking records since pre-read operations can be avoided when writing records.

Appendix G lists the blocking and deblocking algorithms in skeletal form (this file is included on your CP/M disk). Generally, the algorithms map all CP/M sector read operations onto the host disk through an intermediate buffer which is the size of the host disk sector. Throughout the program, values and variables which relate to the CP/M sector involved in a seek operation are prefixed by "sek," while those related to the host disk system are prefixed by "hst." The equate statements beginning on line 29 of Appendix G define the mapping between CP/M and the host system, and must be changed if other than the sample host system is involved.

The entry points BOOT and WBOOT must contain the initialization code starting on line 57, while the SELDSK entry point must be augmented by the code starting on line 65. Note that although the SELDSK entry point computes and returns the Disk Parameter Header address, it does not physically selected the host disk at this point (it is selected later at READHST or WRITEHST). Further, SETTRK, SETTRK, and SETDMA simply store the values, but do not take any other action at this point. SECTRAN performs a trivial trivial function of returning the physical sector number.

The principal entry points are READ and WRITE, starting on lines 110 and 125, respectively. These subroutines take the place of your previous READ and WRITE operations.

The actual physical read or write takes place at either WRITEHST or READHST, where all values have been prepared: hstdsk is the host

disk number, hsttrk is the host track number, and hstsec is the host sector number (which may require translation to a physical sector number). You must insert code at this point which performs the full host sector read or write into, or out of, the buffer at hstbuf of length hstsiz. All other mapping functions are performed by the algorithms.

This particular algorithm was tested using an 80 megabyte hard disk unit which was originally configured for 128 byte sectors, producing approximately 35 megabytes of formatted storage. When configured for 512 byte host sectors, usable storage increased to 57 megabytes, with a corresponding 400% improvement in overall response. In this situation, there is no apparent overhead involved in deblocking sectors, with the advantage that user programs still maintain the (less memory consuming) 128-byte sectors. This is primarily due, of course, to the information provided by the BDOS which eliminates the necessity for pre-read operations to take place.

APPENDIX A: THE MDS COLD START LOADER

		;	MDS-800	Cold Sta	art Loade	er for CP/M 2.0
		; ;	Version	2.Ø Augu	ıst, 1979)
0000	_	; false	0.011	Ø		
ffff		true	equ equ	not fals	20	
0000		testing		false		
0000		;	cqu	14100		
		,	if	testing		
		bias	equ	Ø3400h		
			endif			
			if	not test	ting	
0000	=	bias	equ	0000h		
			endif			
0000		cpmb	equ	bias		;base of dos load
0806		bdos	equ	806h+bia		;entry to dos for calls
1880		bdose	equ	1880h+b:		;end of dos load ;cold start entry point
1600		boot	equ	1600h+b: boot+3	las	;warm start entry point
1603	-	rboot	egu	000043		; walm statt entry point
3000		;	org	3000h	:loaded	here by hardware
5000		;	019	00000	,	
1880	=	bdosl	equ	bdose-c	omb	
0002		ntrks	egu	2	-	;tracks to read
ØØ31	=	bdoss	egu	bdosl/1	28	;# sectors in bdos
ØØ19		bdosØ	egu	25		;# on track Ø
0018	=	bdosl	equ	bdoss-bo	dosØ	;# on track l
60.00	_	;		acoaab	.intol .	nonitor hage
f800 ff0f		mon8Ø rmon8Ø	equ	Øf8ØØh ØffØfh	•	nonitor base t location for mon80
0078		base	equ equ	078h		used by controller
0079		rtype	equ	base+1	;result	
ØØ7b		rbyte	equ	base+3		
007f		reset	equ	base+7		controller
		;	1			
0078	=	dstat	equ	base		tatus port
0079		ilow	equ	base+l		pb address
ØØ7a		ihigh	equ	base+2		opb address
ØØff		bsw	equ	Øffh	;boot si	
0003		recal	equ	3h		brate selected drive
0004		readf	egu	4h Laab		ead function d of boot for stack
0100	=	stack	equ	100h	;use end	d of boot for stack
		; rstart:				
3000	310001	rstart.	lxi	sp.stac	k:in case	e of call to mon8Ø
		;		isk stat		
3003	db79		in	rtype		
	db7b		in	rbyte		
		;		f boot s	witch is	off
		coldsta				
3007			in	bsw		
3888	820730		ani jnz	coldsta	rtswitch	onr
			-			

clear the controller ; 300e d37f ;logic cleared out reset ; ; 3010 0602 mvi b, ntrks ; number of tracks to read 3012 214230 lxi h,iopbØ ; start: ; read first/next track into cpmb ; a,1 3Ø15 7d mov 3016 d379 out ilow 3018 7c a.h mov 3019 d37a out ihiqh 301b db78 wait0: in dstat ani jz waitØ ; check disk status ; 3022 db79 in rtype 3024 e603 ani 11b 2 3026 fe02 cpi ; if testing ;go to monitor if 11 or 10 rmon8Ø cnc endif if. not testing 3028 d20030 jnc rstart ;retry the load endif ; ;i/o complete, check status 302b db7b rbyte in if not ready, then go to mon80 ; 302d 17 ral 302e dc0fff rmon8Ø ;not ready bit set сc 3Ø31 lf ;restore rar 3Ø32 e61e ;overrun/addr err/seek/crc ani 1111Øb ; i f testing ;go to monitor cnz rmon8Ø endif if not testing 3Ø34 c2ØØ3Ø jnz rstart ;retry the load endif ; ; 3037 110700 1xi d, iopbl ; length of iopb 3Ø3a 19 dad d ;addressing next iopb 3Ø3b Ø5 dcr b ; count down tracks 303c c21530 jnz start ; ; jmp boot, print message, set-up jmps ; 3Ø3f c3ØØ16 jmp boot ; parameter blocks ;

3Ø42	80	iopb0:	đb	8Øh	;iocw, no update
3043	Ø 4		db	readf	;read function
3Ø44	19		db	bdosØ	;# sectors to read trk Ø
3045	ØØ		db	Ø	;track Ø
3046	Ø 2		db	2	;start with sector 2, trk Ø
3Ø47	0000		đw	cpmb	;start at base of bdos
0007	=	iopbl	equ	\$-iopbØ	
		;			
3Ø49	8Ø	iopbl:	đb	8Øh	
304a	Ø 4		đb	readf	
304b	18		db	bdosl	;sectors to read on track 1
304c	Ø1		đb	1	;track l
304d	Ø1		db	1	;sector 1
304e	800c		dw	cpmb+bdc	sø*128 ;base of second rd
3050			end	-	

APPENDIX B: THE MDS BASIC I/O SYSTEM (BIOS)

	;		i/o drivers for cp/m 2.0 rive single density version)
	; ;	version	2.0 august, 1979
ØØ14 =	vers	equ	20 ;version 2.0
	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	digital box 579	ht (c) 1979 research , pacific grove nia, 93950
4a00 3400 = 3c06 = 1600 = 002c = 0002 = 0002 = 0004 = 0080 = 000a =	cpmb bdos cpml nsects offset cdisk buff retry ;	equ equ equ equ	<pre>4a00h ;base of bios in 20k system 3400h ;base of cpm ccp 3c06h ;base of bdos in 20k system \$-cpmb ;length (in bytes) of cpm system cpml/128;number of sectors to load 2 ;number of disk tracks used by cp 0004h ;address of last logged disk 0080h ;default buffer address 10 ;max retries on disk i/o before e</pre>
		boot wboot (boot an const conin conout list	nd wboot are the same for mds)
	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	mds, wh seldsk settrk setsec	<pre>llowing calls set-up the io parameter bloc ich is used to perform subsequent reads an select disk given by reg-c (0,1,2) set track address (0,76) for sub r/w set sector address (1,,26) set subsequent dma address (initially 80h</pre>
	;;;;	read/wr read write	
4a00 c3b34a 4a03 c3c34a 4a06 c3614b 4a09 c3644b 4a0c c36a4b	wboote:	jmp	ctor for indiviual routines boot wboot const conin conout

4a0f c36d4b 4a12 c3724b 4a15 c3754b 4a18 c3784b 4a1b c37d4b 4a1e c3a74b 4a21 c3ac4b 4a21 c3ac4b 4a24 c3bb4b 4a27 c3c14b 4a2a c3ca4b 4a2d c3704b 4a30 c3b14b		jmp jmp jmp jmp jmp jmp jmp jmp jmp jmp	list punch reader home seldsk settrk setsec setdma read write listst sectran	;list	status
4a33+= 4a33+824a00 4a37+000000 4a3b+6e4c73 4a3f+0d4dee 4a43+824a00 4a47+000000 4a4b+6e4c73 4a4f+3c4d1d 4a53+824a00 4a55+6e4c73 4a5f+6b4d4c 4a63+824a00 4a65+6e4c73 4a6f+9a4d7b	dpel: dpe2:	maclib disks equ dw dw dw dw dw dw dw dw dw dw dw dw dw	4 \$ xlt0,000 0000h,00 dirbuf,d csv0,alv xlt1,000 0000h,00 dirbuf,d csv1,alv xlt2,000 0000h,00 dirbuf,d csv2,alv xlt3,000 0000h,00 dirbuf,d csv3,alv	; four ; base Øh ØØh pbØ Ø Ø b 0 b 0 b 0 b 2 Ø h ØØh pb2 2 Ø h ØØh pb3 3	<pre>the disk definition library disks of disk parameter blocks ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors ;translate table ;scratch area ;dir buff,parm block ;check, alloc vectors</pre>
4a73+= 4a73+1a00 4a75+03 4a76+07 4a77+00 4a77+00 4a78+f200 4a7a+3f00 4a7c+c0 4a7c+c0 4a7c+1000 4a80+0200 4a80+0200 4a82+= 4a82+01 4a83+07 4a84+0d 4a85+13 4a86+19 4a87+05 4a88+0b 4a89+11 4a8a+17 4a8b+03	dpb0 xlt0	diskdef equ dw db db db db dw dw dw db db db db db db db db db db db db db	0,1,26,6 \$ 26 3 7 0 242 63 192 0 16 2 \$ 1 7 13 19 25 5 11 17 23 3	,1024	,243,64,64,offset ;disk parm block ;sec per track ;block shift ;block mask ;extnt mask ;disk size-1 ;directory max ;alloc0 ; 1loc1 ;check size ;offset ;translate table

<pre>4a73+= dpb1 equ dpb0 ;equivalent parameters 0016+= als1 equ als0 ;same allocation vector size 4a82+= xlt1 equ xlt0 ;same translate table diskdef 2,0 4a73+= dpb2 equ dpb0 ;equivalent parameters 0016+= als2 equ als0 ;same allocation vector size 0016+= css2 equ css0 ;same checksum vector size 4a82+= xlt2 equ xlt0 ;same translate table diskdef 3,0 4a73+= dpb3 equ dpb0 ;equivalent parameters 0016+= css2 equ css0 ;same checksum vector size 4a82+= xlt2 equ xlt0 ;same translate table diskdef 3,0 4a73+= dpb3 equ als0 ;same allocation vector size 0016+= css3 equ css0 ;same checksum vector size 0016+= css3 equ css0 ;same checksum vector size 4a82+= xlt3 equ xlt0 ;same translate table ; endef occurs at end of assembly ; t end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; t he following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ;interrupt revert port 00fc = intc equ 0fdh ;interrupt mask port 00fc = inte equ 0111\$1110b;enable rst 0(warm boot),rst 7 ; mds monitor equates f000 = mon80 equ 0ff06fh ;restart mon80 (boot error) ff06 = ri equ 0ff06fh ;console character to reg-a f006 = ri equ 0ff06fh ;list from c to list device f012 = csts equ 0ff0fh ;list from c to list device f012 = csts equ 0ff0fh ;list from c to list device</pre>	4a8c+Ø9 4a8d+Øf 4a8e+15 4a8f+Ø2 4a9Ø+Ø8 4a91+Øe 4a92+14 4a93+1a 4a94+Ø6 4a95+Øc 4a96+12 4a97+18 4a98+Ø4 4a99+Øa 4a99+Øa 4a9b+16		db db db db db db db db db db db db db d	9 15 21 2 8 14 20 26 6 12 18 24 4 10 16 22	
<pre>001f+= alsl equ alsØ ;same allocation vector size 0010+= cssl equ cssØ ;same checksum vector size 4a82+= xlt1 equ xltØ ;same translate table diskdef 2,0 4a73+= dpb2 equ dpbØ ;equivalent parameters 001f+= als2 equ alsØ ;same allocation vector size 4a82+= xlt2 equ xltØ ;same translate table diskdef 3,0 4a73+= dpb3 equ dpbØ ;equivalent parameters 001f+= als3 equ alsØ ;same checksum vector size 4a82+= xlt2 equ xltØ ;same checksum vector size 0010+= css3 equ cssØ ;same checksum vector size 4a82+= xlt3 equ xltØ ;same translate table ; endef occurs at end of assembly ; i end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; i the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; i we also assume the mds system has four disk drive 00fc = intc equ 0fch ;interrupt mask port 00fc = inte equ 0ffAh ;interrupt mask port 00fc = inte equ 0ffAh ;interrupt mask port i the gov 0ffAh ;interrupt mask port 00fc = inte equ 0ffAh ;interrupt mask port i mds monitor equates f800 = mon80 equ 0ff80h ;restart mon80 (boot error) ff0f = revnt 00ffAh ;restart mon80 (boot error) ff0f = ri equ 0ff80h ;restart mon80 (boot error) ff803 = ci equ 0ff80h ;restart mon80 (boot error) ff804 = ri equ 0ff80h ;restart mon80 (boot error) ff805 = inte equ 0ff80h ;restart mon80 (boot error) ff805 = co equ 0ff80h ;restart mon80 (boot error) ff806 = ri equ 0ff80h ;restart mon80 (boot error) ff809 = co equ 0ff80h ;restart mon80 (boot error) ff806 = ri equ 0ff80h ;resta</pre>	1-72+	dabl		-	acquivelent neremators
Ø01Ø+=csslequcssØ;same checksum vector size4a82+=xlt1equxltØ;same translate tablediskdef 2,04a73+=dpb2equdpbØ;equivalent parametersØ010+=css2equalsØ;same allocation vector sizeØ10+=css2equxltØ;same translate tableda82+=xlt2equxltØ;same translateØ010+=css2equcssØ;same checksum vector sizeØ010+=css3equalsØ;same allocation vector sizeØ010+=css3equcssØ;same translate tableda82+=xlt3equcssØ;same translate tableiendef occurs at end of assembly;endef occurs at end of assemblyiend of controller - independent code, the remaini;are tailored to the particular operating environmibe altered for any system which differs from the;we also assume the mds system has four disk driveØ0fd =revrtequØfA;interrupt revert portiinteequØfA;interrupt control portØ0fe =inteequØfA;interrupt control portimds monitor equatesinteinteimds monitor equatesinteinteimds monitor equatesif ØfA;restart mon80 (boot error)iinteequØfAØA;restart mon80 (boot error)iinteequØfBØA;res			-	-	
<pre>4a82+= xlt1 equ xlt0 ;same translate table diskdef 2,0 4a73+= dpb2 equ dpb0 ;equivalent parameters 001f+= als2 equ als0 ;same allocation vector size 0010+= css2 equ css0 ;same checksum vector size 4a82+= xlt2 equ xlt0 ;same translate table diskdef 3,0 4a73+= dpb3 equ dpb0 ;equivalent parameters 001f+= als3 equ als0 ;same allocation vector size 0010+= css3 equ css0 ;same checksum vector size 4a82+= xlt3 equ xlt0 ;same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ;interrupt revert port 00fc = intc equ 0ffch ;interrupt mask port 00fc = inte equ 0ffch ;interrupt mask port 00fc = inte equ 0ffsh ;interrupt control port for = inte equ 0ffsh ;interrupt control port 00fa = ci equ 0ffsh ;interrupt control port 00fe = revrt equ 0ffsh ;interrupt control port foff = rmon80 equ 0ffs0h ;restart mon80 (boot error) fa06 = ri equ 0ffs0h ;restart mon80 (boot error) fa06 = ri equ 0ffs0h ;restart mors0 (boot error) fa06 = ri equ 0ffs0h ;restart mors0 (boot error) fa07 = inte equ 0ffs0h ;restart mors0 (boot error) fa06 = ri equ 0ffs0h ;restart mors0 (boot error) fa07 = inte equ 0ffs0h ;restart mors0 (boot error) fa07 = inte equ 0ffs0h ;restart mors0 (boot error) fa08 = ri equ 0ffs0h ;restart mors0 (boot error) fa09 = co equ 0ffs0h ;restart mors0 (boot error) fa09 = co equ 0ffs0h ;restart mors0 (boot error) fa09 = co equ 0ffs0h ;restart mors0 co console o fa06 = ri equ 0ffs0h ;restart from c to console o fa07 = lo equ 0ffs0h ;list from c to list device</pre>					
<pre>diskdef 2,0 4a73+= dp2 equ dpb0 ;equivalent parameters 001f+= als2 equ als0 ;same allocation vector size 0010+= css2 equ css0 ;same checksum vector size 4a82+=</pre>			-		•
<pre>4a73+= dpb2 equ dpbØ ;equivalent parameters 0011f+= als2 equ alsØ ;same allocation vector size 0010+= cs2 equ csØ ;same checksum vector size 4a82+= xlt2 equ xltØ ;same translate table diskdef 3,0 4a73+= dpb3 equ dpbØ ;equivalent parameters 001f+= als3 equ alsØ ;same allocation vector size 0010+= cs3 equ csØ ;same checksum vector size 4a82+= xlt3 equ xltØ ;same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fc = intc equ 0fch ;interrupt revert port 00fc = intc equ 0ffh ;interrupt mask port it equ 0ffh ;interrupt control port 00fe = mon80 equ 0ff0fh ;restart mon80 (boot error) ff0ff = rmon80 equ 0ff0fh ;restart mon80 (boot error) ff0ff = rime equ 0ff80fh ;reader in to reg-a ff806 = ri equ 0ff80fh ;reader in to reg-a ff806 = ri equ 0ff80fh ;reader in to reg-a ff807 = lo equ 0ff80fh ;list from c to list device</pre>	40027-	XILI			; Same clansiale cable
<pre>001f+= als2 equ als0 ; same allocation vector size 0010+= css2 equ css0 ; same checksum vector size 4a82+= xlt2 equ xlt0 ; same translate table diskdef 3,0 4a73+= dpb3 equ dpb0 ; equivalent parameters 001f+= als3 equ als0 ; same allocation vector size 0010+= css3 equ css0 ; same checksum vector size 4a82+= xlt3 equ xlt0 ; same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ; interrupt revert port 00ff = intc equ 0ff3h ; interrupt control port 00fe = inte equ 0ff3h ; interrupt control port 00fe = mon80 equ 0ff806h ; mds monitor ff0f = rmon80 equ 0ff806h ; restart mon80 (boot error) ff806 = ri equ 0ff806h ; reader in to reg-a f806 = ri equ 0ff806h ; reader in to reg-a f806 = ri equ 0ff806h ; reader in to reg-a f806 = ri equ 0ff806h ; reader in to reg-a f806 = ri equ 0ff806h ; reader in to reg-a f806 = ri equ 0ff809h ; console char from c to punch devic f807 = lo equ 0ff806h ; punch char from c to punch devic f809 = lo equ 0ff806h ; list from c to list device</pre>	1-73+-	dob2		•	aquivalant parameters
ØØ1Ø+=css2equcss0;same checksum vector size4a82+=xlt2equxltØ;same translate tablediskdef 3,04a73+=dpb3equdpbØ;equivalent parametersØØ16+=als3equalsØ;same allocation vector size4a82+=xlt3equcss0;same checksum vector size4a82+=xlt3equcss0;same translate table;endef occurs at end of assembly;end of controller - independent code, the remaini;are tailored to the particular operating environm;be altered for any system which differs from the;the following code assumes the mds monitor exists;and uses the i/o subroutines within the monitor;we also assume the mds system has four disk driveØØfd =revrtequØff3 =iconiconequØff3 =icon;mds monitor equatesf800 =mon80inteequØff0f =rmon80(bi =integuifequifequifequifids monitorifequifids monitor equatesfa00 =ciequØff0fh ;restart mon80 (boot error)ifids mon80ids mon80equifequifequifequifequifids monitor <t< td=""><td></td><td></td><td></td><td></td><td></td></t<>					
<pre>4a82+= xlt2 equ xlt0 ;same translate table diskdef 3,0 4a73+= dpb3 equ dpb0 ;equivalent parameters 0016+= als3 equ als0 ;same allocation vector size 0010+= css3 equ css0 ;same checksum vector size 4a82+= xlt3 equ xlt0 ;same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ;interrupt revert port 00fc = intc equ 0ffh ;interrupt control port 00f7 = inte equ 0f11\$1100;enable rst 0(warm boot),rst 7 ; mds monitor equates f800 = mon80 equ 0ff00fh ;restart mon80 (boot error) f603 = ci equ 0ff0fh ;restart mon80 (boot error) f604 = ri equ 0f809h ;console character to reg-a f806 = ri equ 0f809h ;console character to reg-a f806 = ri equ 0f809h ;punch char from c to console o f807 = book equ 0f809h ;punch char from c to punch devic f809 = co equ 0f809h ;punch char from c to punch devic f809 = lo equ 0f800h ;list from c to list device</pre>			-		
<pre>diskdef 3,0 4a73+= dpb3 equ dpb0 ;equivalent parameters 001f+= als3 equ als0 ;same allocation vector size 4a82+= xlt3 equ xlt0 ;same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ;interrupt revert port 00fc = intc equ 0fch ;interrupt mask port 00f3 = icon equ 0ff3h ;interrupt control port 00f6 = mon80 equ 0ff800h ;mds monitor ; fo0f = revns00 equ 0ff800h ;restart mon80 (boot error) ff0f = rime equ 0ff800h ;reader in to reg-a f800 = co equ 0ff800h ;console character to reg-a f800 = co equ 0ff80h ;list from c to list device </pre>			-		•
<pre>4a73+= dpb3 equ dpb0 ;equivalent parameters 001f+= als3 equ als0 ;same allocation vector size 0010+= css3 equ css0 ;same checksum vector size 4a82+= xlt3 equ xlt0 ;same translate table ; endef occurs at end of assembly ; ; end of controller - independent code, the remaini ; are tailored to the particular operating environm be altered for any system which differs from the ; ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; 00fd = revrt equ 0fdh ;interrupt revert port 00ff = intc equ 0ffch ;interrupt mask port 00ff = intc equ 0ffch ;interrupt control port 00ff = inte equ 0111\$1110b;enable rst 0(warm boot),rst 7 ; mds monitor equates f800 = mon80 equ 0f800h ;mds monitor ff0f = rimen80 equ 0f800h ;restart mon80 (boot error) f803 = ci equ 0f806h ;restart mon80 (boot error) f806 = ri equ 0f806h ;restart mon80 (boot error) f806 = ri equ 0f806h ;restart mon80 (boot error) f806 = ri equ 0f806h ;restart mon80 (boot error) f807 = inte equ 0f806h ;restart mon80 (boot error) f808 = ri equ 0f806h ;restart mon80 (boot error) f806 = ri equ 0f806h ;restart mon80 (boot error) f807 = bo equ 0f806h ;restart from c to console o f806 = ri equ 0f806h ;restart from c to punch devic f806 = lo equ 0f806h ;list from c to list device</pre>	4027-	XICZ			; same translate table
<pre>001f+= als3 equ als0 ;same allocation vector size 0010+= css3 equ css0 ;same checksum vector size 4a82+= xlt3 equ xlt0 ;same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ;interrupt revert port 00f3 = icon equ 0f3h ;interrupt mask port 00f7 = inte equ 0f11111100;enable rst 0(warm boot),rst 7 ; mds monitor equates f800 = mon80 equ 0f800h ;mds monitor ff0f = rime equ 0f60 ;console character to reg-a f806 = ri equ 0f806 ;reader in to reg-a f809 = co equ 0f806h ;reader in to reg-a f809 = co equ 0f806h ;punch char from c to console o f805 = po equ 0f806h ;pusch char from c to punch devic f806 = po equ 0f806h ;pusch char from c to punch devic f806 = lo equ 0f806h ;list from c to list device</pre>	4-72	Jack 2			
<pre>Ø010+= 4a82+= xlt3 equ xltØ ;same checksum vector size 4a82+= xlt3 equ xltØ ;same translate table ; endef occurs at end of assembly ; end of controller - independent code, the remaini ; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists ; and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive Ø0fd = revrt equ Øfch ;interrupt revert port Ø0f3 = icon equ Øfch ;interrupt mask port inte equ Øff3 ; interrupt control port Ø07e = inte equ Øff3h ;interrupt control port ; mds monitor equates f800 = mon80 equ Øff0fh ;restart mon80 (boot error) f60f = ri equ Øff0fh ;restart mon80 (boot error) f806 = ri equ Øff0fh ;restart mon80 (boot error) f806 = ri equ Øff0fh ;reader in to reg-a f806 = ri equ Øff0fh ;reader in to reg-a f809 = co equ Øff0fh ;punch char from c to console o f80c = po equ Øff0fh ;punch char from c to punch devic f80f = lo equ Øf80fh ;list from c to list device</pre>			+		
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<pre> if end of controller - independent code, the remaini are tailored to the particular operating environm be altered for any system which differs from the if the following code assumes the mds monitor exists and uses the i/o subroutines within the monitor if we also assume the mds system has four disk drive if equ Øfdh ; interrupt revert port if we also assume the mds system has four disk drive if equ Øfch ; interrupt mask port if equ Øf3h ; interrupt mask port if equ Øf3h ; interrupt control port if equ Øf1sh ; interrupt control port if equ Øf800h ; mds monitor if mds monitor equates f800 = mon80 equ Øf800h ; mds monitor if equ Øf803h ; console character to reg-a f806 = ri equ Øf806h ; reader in to reg-a f806 = ri equ Øf806h ; console char from c to console o f80c = po equ Øf806h ; punch char from c to punch devic f80f = lo equ Øf806h ; list from c to list device </pre>	4a82+=				-
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<pre>; are tailored to the particular operating environm ; be altered for any system which differs from the ; the following code assumes the mds monitor exists and uses the i/o subroutines within the monitor ; we also assume the mds system has four disk drive 00fd = revrt equ 0fdh ; interrupt revert port 00fc = intc equ 0fch ; interrupt mask port 00f3 = icon equ 0f3h ; interrupt control port 007e = inte equ 0111\$1110b; enable rst 0(warm boot), rst 7 ; mds monitor equates f800 = mon80 equ 0f800h ; mds monitor ff0f = rmon80 equ 0f801h ; restart mon80 (boot error) f803 = ci equ 0f803h ; console character to reg-a f806 = ri equ 0f806h ; reader in to reg-a f809 = co equ 0f809h ; console char from c to console o f80c = po equ 0f80ch ; punch char from c to punch devic f80f = lo equ 0f80fh ; list from c to list device</pre>					
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<pre>ØØfd = revrt equ Øfdh ;interrupt revert port ØØfc = intc equ Øfch ;interrupt mask port ØØf3 = icon equ Øf3h ;interrupt control port ØØ7e = inte equ Ø111\$111Øb;enable rst Ø(warm boot),rst 7 ; ; mds monitor equates f8ØØ = mon8Ø equ Øf8ØØh ;mds monitor ffØf = rmon8Ø equ ØffØfh ;restart mon8Ø (boot error) f8Ø3 = ci equ ØffØfh ;restart mon8Ø (boot error) f8Ø3 = ci equ Øf8Ø3h ;console character to reg-a f8Ø6 = ri equ Øf8Ø6h ;reader in to reg-a f8Ø9 = co equ Øf8Ø9h ;console char from c to console o f8Øc = po equ Øf8Øch ;punch char from c to punch devic f8Øf = lo equ Øf8Øfh ;list from c to list device</pre>		;	-		
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<pre>ØØf3 = icon equ Øf3h ;interrupt control port ØØ7e = inte equ Øf3h ;interrupt control port ; ; mds monitor equates f800 = mon80 equ Øf800h ;mds monitor ff0f = rmon80 equ Øf80fh ;restart mon80 (boot error) f803 = ci equ Øf803h ;console character to reg-a f806 = ri equ Øf806h ;reader in to reg-a f809 = co equ Øf809h ;console char from c to console o f80c = po equ Øf80ch ;punch char from c to punch devic f80f = lo equ Øf80fh ;list from c to list device</pre>					
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; mds monitor equates f800 = mon80 equ 0f800h ;mds monitor ff0f = rmon80 equ 0ff0fh ;restart mon80 (boot error) f803 = ci equ 0f803h ;console character to reg-a f806 = ri equ 0f806h ;reader in to reg-a f809 = co equ 0f809h ;console char from c to console o f80c = po equ 0f80ch ;punch char from c to punch devic f80f = lo equ 0f80fh ;list from c to list device	ØØ7e =	inte	equ	0111\$11	løb;enable rst Ø(warm boot),rst 7
f800 =mon80equØf800h;mds monitorff0f =rmon80equØff0fh;restart mon80(boot error)f803 =ciequØf803h;console character to reg-af806 =riequØf806h;reader in to reg-af809 =coequØf809h;console char from c to console of80c =poequØf80ch;punch char from c to punch devicf80f =loequØf80fh;list from c to list device		;			
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f80c =poequ0f80ch;punch char from c to punch devicf80f =loequ0f80fh;list from c to list device		ri	equ	Øf806h	;reader in to reg-a
f80c =poequ0f80ch;punch char from c to punch devicf80f =loequ0f80fh;list from c to list device		co	equ	Ø£8Ø9h	; console char from c to console o
f80f = lo equ 0f80fh ;list from c to list device		ро	equ	Øf8Øch	; punch char from c to punch devic
		10	equ	Øf8Øfh	
	f812 =	csts	-		
					· –

; disk ports and commands ; ;base of disk command io ports 0078 =equ 78h base base ;disk status (input) 0078 =dstat equ base+1 ;result type (input) $\emptyset \emptyset 79 =$ rtype equ ;result byte (input) 007b =base+3 rbyte equ ; ; iopb low address (output) 0079 =ilow eau base+1 ØØ7a = ihigh equ base+2 ; iopb high address (output) ; 0004 =4 h ;read function readf equ 0006 =writf 6h :write function equ 3h :recalibrate drive 0003 =recal equ 0004 =iordy equ 4h ;i/o finished mask Ødh $\emptyset \emptyset \emptyset d =$;carriage return cr equ 000a = lf Øah ;line feed equ ; signon: ;signon message: xxk cp/m vers y.y 4a9c ØdØaØa db cr,lf,lf 4a9f 323Ø db '2Ø' ;sample memory size 4aal 6b2043f 'k cp/m vers db vers/10+'0','.',vers mod 10+'0' 4aad 322e30 db 4abØ ØdØaØØ āb cr.lf.Ø boot: print signon message and go to ccp (note: mds boot initialized iobyte at 0003h) ; 4ab3 310001 lxi sp,buff+80h 4ab6 219c4a lxi h,signon prmsq 4ab9 cdd34b call ;print message 4abc af xra ; clear accumulator а 4abd 320400 sta cdisk ;set initially to disk a 4acØ c3Øf4b jmp gocpm ; go to cp/m ; ; wboot:; loader on track Ø, sector 1, which will be skippe read cp/m from disk - assuming there is a 128 byt ; start. ; ; 4ac3 318000 lxi sp, buff ; using dma - thus 80 thru ff ok f ; 4ac6 ØeØa mvi c, retry ; max retries 4ac8 c5 push b wboot0: ;enter here on error retries 4ac9 Ø1ØØ34 lxi b,cpmb ;set dma address to start of disk 4acc cdbb4b call setdma 4acf ØeØØ mvi c,Ø ;boot from drive Ø 4adl cd7d4b call seldsk 4ad4 ØeØØ mvi c.Ø 4ad6 cda74b call settrk ;start with track Ø 4ad9 ØeØ2 mvi c,2 ;start reading sector 2 4adb cdac4b call setsec ; ; read sectors, count nsects to zero 4ade cl ;10-error count pop b 4adf Ø62c mvi b,nsects

		rdsec:	;read ne	ext secto	
4ael			push	b	;save sector count
4ae2	cdcl4b		call	read	
4ae5	c2494b		jnz		;retry if errors occur
4ae8	2a6c4c		lhld	iod	
4aeb	118000		lxi	d,128	;sector size
4aee	19		dad	đ	; incremented dma address in hl
4aef	44		mov	b,h	
4afØ	4d		mov		;ready for call to set dma
4afl	cdbb4b		call	setdma	
4af4	3a6b4c		lda		;sector number just read
4af7	fela		cpi	26	;read last sector?
4af9	daØ54b		jc	rdl	
		;	must be	sector 2	26, zero and go to next track
4afc	3a6a4c		lda	iot	;get track to register a
4aff	3c		inr	a	
4bØØ	4f		mov	c,a	;ready for call
4bØl	cda74b		call	settrk	
4bØ4	af		xra	a	;clear sector number
4bØ5	3c	rdl:	inr	а	;to next sector
4bØ6	4f		mov	c,a	;ready for call
4bØ7	cdac4b		call	setsec	
4b0a	cl		gog		;recall sector count
4bøb			dcr	b	;done?
4bØc	c2el4a		jnz	rdsec	
		;			
		;			bad, reset default buffer address
					om cold start boot)
		;		stØ and	rst/
4bØf			di.	1.01	
	3el2		mvi	a,12h	;initialize command
	d3fd		out	revrt	
4b14			xra	a · ·	
	d3fc		out		;cleared
	3e7e				;rst0 and rst7 bits on
	d3fc		out	intc	
4blb			xra	a	
401C	d3f3		out	icon	;interrupt control
		;			for oddrogg to 00b
4-1-	a10000	;			fer address to 80h
	018000		lxi	b, buff	
4021	cdbb4b		call	setdma	
		;	roact m	onitor of	atry points
		;		Surren	ntry points
11-21	2 2				
	3ec3		mvi	a,jmp	
4b26	320000		sta	Ø	
4b26 4b29	320000 21034a		sta lxi	Ø h,wboote	
4b26 4b29 4b2c	320000 21034a 220100		sta lxi shld	Ø h,wboote 1	e ;jmp wboot at location ØØ
4b26 4b29 4b2c 4b2f	320000 21034a 220100 320500		sta lxi shld sta	Ø h,wboote 1 5	
4b26 4b29 4b2c 4b2f 4b32	320000 21034a 220100 320500 21063c		sta lxi shld sta lxi	Ø h,wboote 1 5 h,bdos	;jmp wboot at location ØØ
4b26 4b29 4b2c 4b2f 4b32 4b35	320000 21034a 220100 320500 21063c 220600		sta lxi shld sta lxi shld	Ø h,wboote 1 5 h,bdos 6	;jmp wboot at location 00; ;jmp bdos at location 5
4b26 4b29 4b2c 4b2f 4b32 4b35 4b38	320000 21034a 220100 320500 21063c 220600 323800		sta lxi shld sta lxi shld sta	Ø h,wboote 1 5 h,bdos 6 7*8	;jmp wboot at location ØØ
4b26 4b29 4b2c 4b2f 4b32 4b35 4b38 4b38	320000 21034a 220100 320500 21063c 220600 323800 2100f8		sta lxi shld sta lxi shld sta lxi	Ø h,wboote 1 5 h,bdos 6 7*8 h,mon8Ø	;jmp wboot at location 00; ;jmp bdos at location 5
4b26 4b29 4b2c 4b2f 4b32 4b35 4b38 4b38	320000 21034a 220100 320500 21063c 220600 323800	•	sta lxi shld sta lxi shld sta lxi shld	Ø h,wboote 1 5 h,bdos 6 7*8 h,mon8Ø 7*8+1	;jmp wboot at location 00 ;jmp bdos at location 5 ;jmp to mon80 (may have been chan
4b26 4b29 4b2c 4b2f 4b32 4b35 4b38 4b38	320000 21034a 220100 320500 21063c 220600 323800 2100f8	;	sta lxi shld sta lxi shld sta lxi shld	Ø h,wboote 1 5 h,bdos 6 7*8 h,mon8Ø	;jmp wboot at location 00 ;jmp bdos at location 5 ;jmp to mon80 (may have been chan

previously selected disk was b, send parameter to 1da cdisk ;last logged disk number 4b41 3a0400 4b44 4f ;send to ccp to log it in mov c,a 4b45 fb ei 4b46 c30034 jmp cpmb ; error condition occurred, print message and retry ; booterr: 4b49 cl b ;recall counts pop 4b4a Ød dcr С 4b4b ca524b booterØ jz try again ; 4b4e c5 push b 4b4f c3c94a jmp wbootØ booter0: otherwise too many retries ; 4b52 215b4b lxi h,bootmsq 4b55 cdd34b call prmsq rmon80 ;mds hardware monitor 4b58 c30fff jmp bootmsg: 4b5b 3f626f4 '?boot',Ø db ; ; ; console status to reg-a const: (exactly the same as mds call) ; 4b61 c312f8 jmp csts ; ; console character to reg-a conin: 4b64 cdØ3f8 call ci 4b67 e67f ani 7fh ;remove parity bit 4b69 c9 ret ; conout: ; console character from c to console out 4b6a c3Ø9f8 jmp CO list: ;list device out (exactly the same as mds call) ; 4b6d c30ff8 jmp 10 listst: ;return list status 4b70 af xra а 4b71 c9 ret ;always not ready ; punch: ; punch device out (exactly the same as mds call) ; 4b72 c30cf8 jmp po ; reader: ;reader character in to reg-a (exactly the same as mds call) ; 4b75 c306f8 jmp ri home: ; move to home position

treat as track ØØ seek ; 4b78 ØeØØ c.Ø mvi settrk 4b7a c3a74b jmp seldsk: ;select disk given by register c h.0000h ;return 0000 if error 4b7d 210000 lxi 4b8Ø 79 mov a,c ndisks ;too large? 4b81 feØ4 cpi ; leave hl = 00004b83 dØ rnc ; ;00 00 for drive 0,1 and 10 10 fo 4b84 e602 ani 10b 4b86 32664c dbank :to select drive bank sta 4b89 79 ;00, 01, 10, 11 mov a,c ;mds has Ø,1 at 78, 2,3 at 88 4b8a e601 ani 1b;result ØØ? ora 4b8c b7 а 4b8d ca924b jΖ setdrive ;selects drive 1 in bank 4b90 3e30 mvi a,00110000b setdrive: ;save the function 4b92 47 mov b.a 4b93 21684c ; io function lxi h,iof mov a,m 4b96 7e ;mask out disk number 11001111b 4b97 e6cf ani ;mask in new disk number 4b99 bØ ora b 4b9a 77 ;save it in iopb mov m,a 4898 5800 h:Ø moy :hl=disk number 4b9e 29 dad h ;*2 ;*4 dad h 4b9f 29 4baØ 29 dad ;*8 h ;*16 4bal 29 dad h lxi d,dpbase 4ba2 11334a ;hl=disk header table address d 4ba5 19 dad ret 4ba6 c9 ; ; settrk: ;set track address given by c 4ba7 216a4c h.iot lxi 4baa 71 mov m,c 4bab c9 ret ; setsec: ;set sector number given by c h,ios 4bac 216b4c lxi 4baf 71 mov m.c 4bbØ c9 ret sectran: ;translate sector bc using table at de ;double precision sector number i 4bbl 0600 mvi b,Ø ;translate table address to hl 4bb3 eb xchq 4bb4 Ø9 ;translate(sector) address dad b 4bb5 7e mov ;translated sector number to a a,m 4bb6 326b4c sta ios ;return sector number in 1 4669 6f moy l,a ; setdma: ;set dma address given by regs b,c

4bbb 69 mov 1,c h.b 4bbc 60 mov iod 4bbd 226c4c shld 4bcØ c9 ret ; ;read next disk record (assuming disk/trk/sec/dma read: c, readf ; set to read function mvi 4bcl ØeØ4 setfunc 4bc3 cde04b call ;perform read function call waitio 4bc6 cdfØ4b ;may have error set in reg-a ret 4bc9 c9 ; ; ;disk write function write: c.writf 4bca ØeØ6 mvi setfunc ;set to write function 4bcc cdeØ4b call waitio 4bcf cdfØ4b call ;may have error set 4bd2 c9 ret ; ; utility subroutines ; ;print message at h,l to Ø prmsq: 4bd3 7e mov a,m 4bd4 b7 ora а ;zero? 4bd5 c8 rz more to print ; 4bd6 e5 push h c,a 4bd7 4f mov conout 4bd8 cd6a4b call 4bdb el pop h 4bdc 23 inx h 4bdd c3d34b jmp prmsg ; setfunc: set function for next i/o (command in reg-c) ; ; io function address 4beØ 21684c lxi h,iof ;get it to accumulator for maskin 4be3 7e mov a,m 11111000b ; remove previous command 4be4 e6f8 ani 4be6 bl ;set to new command ora С 4be7 77 ;replaced in iopb mov m,a the mds-800 controller reg's disk bank bit in sec ; mask the bit from the current i/o function ; 4be8 e620 ØØ100000b ;mask the disk select bit ani 4bea 216b4c lxi h,ios ;address the sector selec 4bed b6 ;select proper disk bank ora m 4bee 77 ;set disk select bit on/o mov m,a 4bef c9 ret ; waitio: 4bfØ ØeØa mvi c, retry ; max retries before perm error rewait: start the i/o function and wait for completion ; 4bf2 cd3f4c call ; in rtype intype 4bf5 cd4c4c call inbyte ;clears the controller ; 4bf8 3a664c lda dbank ;set bank flags

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4bfb b7 ;zero if drive 0,1 and nz ora а 4bfc 3e67 a, iopb and Øffh ; low address for iopb mvi 4bfe Ø64c b,iopb shr 8 ; high address for iopb mvi iodrl ;drive bank 1? 4c00 c20b4c jnz 4cØ3 d379 ilow out ;low address to controlle 4cØ5 78 a,b mov 4cØ6 d37a ihigh out ;high address 4cØ8 c31Ø4c jmp waitØ ; to wait for complete iodrl: ;drive bank 1 4cØb d389 out ilow+lØh ;88 for drive bank 10 4cØd 78 mov a,b 4c0e d38a out ihigh+lØh ; 4cl0 cd594c wait0: call instat ;wait for completion 4cl3 e604 ani iordy ;ready? 4cl5 cal04c waitØ jz ; check io completion ok ; 4c18 cd3f4c call intype ; must be io complete (00) 00 unlinked i/o complete, 01 linked i/o comple ; 10 disk status changed ll (not used) ; 4clb fe02 cpi lØb ;ready status change? 4cld ca324c jΖ wready ; must be $\emptyset\emptyset$ in the accumulator ; 4c20 b7 ora а 4c21 c2384c jnz werror ; some other condition, re ; check i/o error bits ; 4c24 cd4c4c call inbyte 4c27 17 ral 4c28 da324c jс wready ;unit not ready 4c2b lf rar 4c2c e6fe ani 11111110b ; any other errors? 4c2e c2384c jnz werror ; read or write is ok, accumulator contains zero ; 4c31 c9 ret ; wready: ;not ready, treat as error for now 4c32 cd4c4c call inbyte ;clear result byte 4c35 c3384c jmp trycount ; werror: ;return hardware malfunction (crc, track, seek, e the mds controller has returned a bit in each pos ; of the accumulator, corresponding to the conditio ; Ø - deleted data (accepted as ok above) ; - crc error 1 ; 2 - seek error ; 3 - address error (hardware malfunction) ; 4 - data over/under flow (hardware malfunct ; 5 - write protect (treated as not ready) ; 6 - write error (hardware malfunction) ; 7 not ready ;

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(accumulator bits are numbered 7 6 5 4 3 2 1 \emptyset) ; ; it may be useful to filter out the various condit : but we will get a permanent error message if it i ; recoverable, in any case, the not ready conditio ; treated as a separate condition for later improve ; trycount: register c contains retry count, decrement 'til z ; 4c38 Ød dcr С rewait ; for another try 4c39 c2f24b jnz ; cannot recover from error ; 4c3c 3eØl mvi a.l ;error code 4c3e c9 ret : intype, inbyte, instat read drive bank 00 or 10 4c3f 3a664c intype: lda dbank 4c42 b7 ora а 4c43 c2494c jnz intypl ;skip to bank 10 4c46 db79 in rtype 4c48 c9 ret 4c49 db89 intypl: in rtype+10h ;78 for Ø,1 88 for 2,3 4c4b c9 ret 4c4c 3a664c inbyte: lda dbank 4c4f b7 ora а 4c50 c2564c inz inbvtl 4c53 db7b in rbyte 4c55 c9 ret 4c56 db8b inbytl: in rbyte+10h 4c58 c9 ret 4c59 3a664c instat: 1da dbank 4c5c b7 ora а 4c5d c2634c inz instal 4c60 db78 in dstat 4c62 c9 ret 4c63 db88 instal: in dstat+10h 4c65 c9 ret ; ; ; data areas (must be in ram) : 4c66 00 dbank: db Ø ;disk bank 00 if drive 0.1 10 if drive 2.3 ; iopb: ; io parameter block 4c67 8Ø db 80h ;normal i/o operation 4c68 Ø4 iof: db readf ; io function, initial read 4c69 Ø1 ion: db 1 ;number of sectors to read 4c6a Ø2 iot: db offset ;track number 4c6b Ø1 ios: db 1 ;sector number 4c6c 8000 iod: dw buff ;io address ; ; define ram areas for bdos operation ;

4c6e+= begdat 4c6e+ dirbuf: 4ce+ alv0: 4d0d+ csv0: 4d1d+ alv1: 4d3c+ csv1: 4d4c+ alv2: 4d6b+ csv2: 4d7b+ alv3: 4d9a+ csv3: 4daa+= enddat Øl3c+= datsiz 4daa datsiz	endef equ ds ds ds ds ds ds ds ds ds equ equ end	\$ 128 31 16 31 16 31 16 31 16 \$ \$-begdat	;directory	access	buffer
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APPENDIX C: A SKELETAL CBIOS

	;	skeletal cbios for first level of cp/m 2.0 altera
ØØ14 =	; msize	equ 20 ;cp/m version memory size in kilo
	; ; ;	"bias" is address offset from 3400h for memory sy than 16k (referred to as "b" throughout the text)
	bias ccp bdos bios cdisk iobyte	equ (msize-20)*1024 equ 3400h+bias ;base of ccp equ ccp+806h ;base of bdos equ ccp+1600h ;base of bios equ 0004h ;current disk number 0=a,,15=p equ 0003h ;intel i/o byte
4a00 002c =	nsects	org bios ;origin of this program equ (\$-ccp)/128 ;warm start sector count
4a@@ c39c4a 4a@3 c3a64a 4a@6 c3114b 4a@9 c3244b 4a@c c3374b 4a@f c3494b 4a12 c34d4b 4a15 c34f4b 4a15 c34f4b 4a16 c35a4b 4a16 c37d4b 4a21 c3924b 4a21 c3924b 4a24 c3ad4b 4a27 c3c34b 4a2a c3d64b 4a3@ c3a74b		<pre>jump vector for individual subroutines jmp boot ;cold start jmp wboot ;warm start jmp const ;console status jmp conin ;console character in jmp conout ;console character out jmp list ;list character out jmp punch ;punch character out jmp reader ;reader character out jmp home ;move head to home positi jmp seldsk ;select disk jmp settrk ;set track number jmp setsec ;set sector number jmp setdma ;set dma address jmp read ;read disk jmp write ;write disk jmp listst ;return list status jmp sectran ;sector translate</pre>
4a33 734a00 4a37 000000 4a3b f04c8d 4a3f ec4d70 4a43 734a00 4a45 f04c8d 4a4f fc4d8f 4a53 734a00 4a57 000000 4a5b f04c8d 4a5f 0c4eae	; ; dpbase: ;	<pre>fixed data tables for four-drive standard ibm-compatible 8" disks disk parameter header for disk 00 dw trans,0000h dw 0000h,0000h dw dirbf,dpblk dw chk00,all00 disk parameter header for disk 01 dw trans,0000h dw 0000h,0000h dw dirbf,dpblk dw chk01,all01 disk parameter header for disk 02 dw trans,0000h dw 0000h,0000h dw 0000h,0000h dw 0000h,0000h</pre>

4a63 734a00 4a67 000000 4a6b f04c8d 4a6f lc4ecd	;	disk parameter header for disk Ø3 dw trans,0000h dw 0000h,0000h dw dirbf,dpblk dw chk03,all03	
4a73 01070d 4a77 19050b 4a7b 170309 4a7f 150208 4a83 141a06 4a87 121804 4a8b 1016	; trans:	Sector translate vector1125,5,11,173b23,3,9,153b21,2,8,143b20,26,6,123b18,24,4,103b16,22	
4a8d la00 4a8f 03 4a90 07 4a91 00 4a92 f200 4a94 3f00 4a96 c0 4a97 00 4a98 l000 4a9a 0200	; dpblk:	disk parameter block, common to all disksdw26db3db3db7db7db0flw242dw63db192db0db16gw24db16grade2dw16grade2dw2dw16grade3dw16grade3dw16grade3dw16grade3dw16grade3dw16grade3dw16grade3dw16grade3dw16grade3dw16grade3grade </td <td></td>	
1	; ; ; ; boot:	end of fixed tables individual subroutines to perform each function ;simplest case is to just perform parameter initi	
4a9c af 4a9d 320300 4aa0 320400 4aa3 c3ef4a	;	xraa; zero in the accumstaiobyte; clear the iobytestacdisk; select disk zerojmpgocpm; initialize and go to cp/	
4aa6 318000 4aa9 0e00 4aab cd5a4b 4aae cd544b	wboot:	<pre>;simplest case is to read the disk until all sect lxi sp,80h ;use space below buffer f mvi c,0 ;select disk 0 call seldsk call home ;go to track 00</pre>	
4abl Ø62c 4ab3 ØeØØ 4ab5 16Ø2	; ; ;	mvib,nsects;b counts # of sectors tomvic,0;c has the current trackmvid,2;d has the next sector tonote that we begin by reading track 0, sector 2 scontains the cold start loader, which is skipped	
4ab7 210034 4aba c5 4abb d5 4abc e5 4abd 4a 4abe cd924b 4acl cl	loadl:	<pre>lxi h,ccp ;base of cp/m (initial lo ;load one more sector push b ;save sector count, current track push d ;save next sector to read push h ;save dma address mov c,d ;get sector address to register c call setsec ;set sector address from register pop b ;recall dma address to b,c</pre>	

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;replace on stack for later recal 4ac2 c5 push b ;set dma address from b,c setdma 4ac3 cdad4b call ; drive set to Ø, track set, sector set, dma addres ; call read 4ac6 cdc34b cpi ØØh ; any errors? 4ac9 fe00 ;retry the entire boot if an erro wboot 4acb c2a64a jnz ; no error, move to next sector ; 4ace el ;recall dma address h pop ;dma=dma+128 4acf 118000 d.128 lxi ;new dma address is in h,l 4ad2 19 dad d 4ad3 d1 pop d ;recall sector address 4ad4 cl b ;recall number of sectors remaini pop ;sectors=sectors-1 4ad5 Ø5 dcr b ;transfer to cp/m if all have bee 4ad6 caef4a jz gocpm ; more sectors remain to load, check for track chan ; 4ad9 14 inr d 4ada 7a ;sector=27?, if so, change tracks mov a,d 4adb felb cpi 27 4add daba4a ;carry generated if sector<27 loadl jc ; end of current track, go to next track ; 4aeØ 16Ø1 mvi d.1 ; begin with first sector of next 4ae2 Øc С ;track=track+1 inr ; save register state, and change tracks ; 4ae3 c5 push b 4ae4 d5 d push 4ae5 e5 push h 4ae6 cd7d4b ;track address set from register call settrk 4ae9 el h pop 4aea dl d pop 4aeb cl b pop : for another sector 4aec c3ba4a loadl jmp ; end of load operation, set parameters and go to c ; gocpm: a,Øc3h 4aef 3ec3 ;c3 is a jmp instruction mvi ; for jmp to wboot 4afl 320000 sta Ø 4af4 21034a lxi h,wboote ;wboot entry point 4af7 220100 shld 1 ;set address field for jmp at Ø 4afa 320500 sta 5 ; for jmp to bdos 4afd 21063c lxi h,bdos ;bdos entry point 4600 220600 shld 6 ;address field of jump at 5 to bd ; 4b03 018000 lxi b,80h ;default dma address is 80h 4b06 cdad4b call setdma ; 4b09 fb ei ;enable the interrupt system 4b0a 3a0400 lda cdisk ;get current disk number 4bØd 4f mov ;send to the ccp c,a 4b0e c30034 jmp CCD ;go to cp/m for further processin

; ; simple i/o handlers (must be filled in by user) ; in each case, the entry point is provided, with s ; to insert your own code ; ; ; console status, return Øffh if character ready, const: 4b11 ;space for status subroutine ds lØh 4b21 3eØØ mvi a.00h 4b23 c9 ret ; conin: ; console character into register a 4b24 ds lØh ;space for input routine 4b34 e67f ani 7fh ;strip parity bit 4b36 c9 ret ; conout: ; console character output from register c 4b37 79 ;get to accumulator mov a.c 4b38 ds 10h ;space for output routine 4b48 c9 ret list: ;list character from register c 4b49 79 mov a,c ; character to register a 4b4a c9 ret ;null subroutine listst: ;return list status (Ø if not ready, l if ready) ;Ø is always ok to return 4b4b af xra a 4b4c c9 ret ; punch: ; punch character from register c 4b4d 79 ; character to register a mov a,c 4b4e c9 ret ;null subroutine ; ; reader: ; read character into register a from reader devic 4b4f 3ela mvi a,lah ;enter end of file for now (repla 4b51 e67f ani 7fh ; remember to strip parity bit 4b53 c9 ret ; ; i/o drivers for the disk follow ; for now, we will simply store the parameters away ; in the read and write subroutines ; ; ;move to the track 00 position of current drive home: translate this call into a settrk call with param ; 4b54 ØeØØ с.Ø mvi ;select track Ø 4b56 cd7d4b call settrk 4b59 c9 ret ;we will move to 00 on first read ; seldsk: ;select disk given by register c 4b5a 210000 h,0000h ;error return code lxi 4b5d 79 mov a.c 4b5e 32ef4c sta diskno 4b61 feØ4 cpi 4 ;must be between Ø and 3

;no carry if 4,5,... 4b63 dØ rnc disk number is in the proper range ; ;space for disk select 4b64 ds 10 compute proper disk parameter header address ; diskno 4b6e 3aef4c lda ;l=disk number Ø,1,2,3 4b71 6f mov 1,a 4b72 2600 ; high order zero mvi h,Ø 4b74 29 ;*2 dađ h :*4 h 4b75 29 dad 4b76 29 dad h :*8 4b77 29 ;*16 (size of each header) dad h 4b78 11334a lxi d,dpbase ;hl=.dpbase(diskno*16) 4b7b 19 dad d 4b7c c9 ret ; settrk: ;set track given by register c 4b7d 79 mov a.c 4b7e 32e94c sta track 4b81 lØh ;space for track select ds 4b91 c9 ret ; setsec: ;set sector given by register c 4b92 79 mov a,c 4b93 32eb4c sector sta 4b96 lØh ;space for sector select ds 4ba6 c9 ret ; sectran: ;translate the sector given by bc using the ;translate table given by de 4ba7 eb xchq ;hl=.trans 4ba8 Ø9 dad b ; hl=.trans(sector) 4ba9 6e **1**,m ;1 = trans(sector) mov 4baa 2600 ;hl= trans(sector) h,Ø mvi 4bac c9 ret ;with value in hl ; setdma: ;set dma address given by registers b and c 4bad 69 1,c ;low order address mov 4bae 60 mov h.b ; high order address 4baf 22ed4c ;save the address shld dmaad 4bb2 đs 10h ;space for setting the dma addres 4bc2 c9 ret ; read: ;perform read operation (usually this is similar so we will allow space to set up read command, th ; common code in write) ; 4bc3 ds lØh ;set up read command 4bd3 c3e64b jmp waitio ; to perform the actual i/o ; write: ;perform a write operation 4bd6 ;set up write commanu 10h ds : waitio: ;enter here from read and write to perform the ac operation. return a ØØh in register a if the ope ; properly, and Ølh if an error occurs during the r ;

4be6 4ce6 3eØ1 4ce8 c9	;;;;;	ds mvi ret	256 a,1	e have saved the disk number in 'd the track number in 'track' (Ø-76 the sector number in 'sector' (1- the dma address in 'dmaad' (Ø-655 ;space reserved for i/o drivers ;error condition ;replaced when filled-in
	;;;;	data ar system n	ea, and o memory in	f the cbios is reserved uninitiali does not need to be a part of the mage (the space must be available, n "begdat" and "enddat").
4ce9	, track:	ds	2	;two bytes for expansion
4ceb	sector:		2	; two bytes for expansion
4ced	dmaad:		2	;direct memory address
4cef	diskno:		ī	;disk number 0-15
	;			
	;		ram area	a for bdos use
$4cf\emptyset =$	begdat		\$;beginning of data area
4cfØ	dirbf:	ds	128	scratch directory area;
4d7Ø	al100:		31	;allocation vector Ø
4d8f	all01:		31	;allocation vector 1
4dae	al102:		31	;allocation vector 2
4dcd	al103:		31	;allocation vector 3
4dec	chk00:	ds	16	;check vector Ø
4dfc	chkØ1:	ds	16	; check vector 1
4eØc	chk02:	ds	16	; check vector 2
4elc	chkØ3:	ds	16	;check vector 3
4e2c =	, enddat	equ	Ş	;end of data area
Ø13c =	datsiz	equ		t;size of data area
4e2c		end	-	

	APPENDIX	D: A SKELETAL	GETSYS/PUTSYS PROG	GRAM
	; ;		s and putsys progra rams at the base of	
0100		org ØlØØh		
0014 =	msize	egu 20	; size of	cp/m in Kbytes
	; "bias ;		to add to addresse s "b" throughout th	
0000 = 3400 = 3c00 = 4a00 =	bias ccp bdos bios	equ (msize- equ 3400h+ equ ccp+086 equ ccp+166	ðøh	
	;	getsys programs 3880h + bias	s tracks Ø and l to	memory at
	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	register a b c d,e h,l sp	usage (scratch register track count (Ø sector count (l (scratch register load address set to stack addr	,76) .26) : pair)
Ø10Ø 318033 Ø103 218033 Ø106 Ø6ØØ		lxi sp,ccp- lxi h,ccp-6 mvi b,0	-0080h 3080h	start of getsys convenient plac set initial loa start with trac read next track
Ø108 ØeØ1	rd\$sec:	mvi c,l	;	each track star
010a cd0003 010d 118000 0110 19 0111 0c 0112 79 0113 felb 0115 da0a01	1	call read\$se lxi d,128 dad d inr c mov a,c cpi 27 jc rdsec	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	get the next se offset by one s (h1=h1+128) next sector fetch sector nu and see if la <, do one more
	; arriv	e here at end of	f track, move to ne	ext track
Ø118 Ø4 Ø119 78 Ø11a feØ2 Ø11c daØ8Ø1		inr b mov a,b cpi 2 jc rd\$trk	;	track = track+l check for last track = 2 ? <, do another
	; arriv	e here at end of	f load, halt for la	ack of anything b
Øllf fb Øl20 76		ei hlt		

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putsys program, places memory image starting at 3880h + bias back to tracks 0 and 1 start this program at the next page boundary

0200 org (\$+0100h) and 0ff00h

0200 318033	put\$sys: lxi	sp,ccp-0080h	; convenient plac
0203 218033	lxi	h,ccp-0080h	; start of dump
0206 0600	mvi	b,Ø	; start with trac
	wr\$trk:		
Ø2Ø8 ØeØl	mvi	c,1	; start with sect
	wr\$sec:		
020a cd0004	call	write\$sec	; write one secto
Ø2Ød 118ØØØ	lxi	d,128	; length of each
Ø21Ø 19	dad	d	; <hl>=<hl> + 128</hl></hl>
Ø211 Øc	inr	С	; $\langle c \rangle = \langle c \rangle + 1$
Ø212 79	mov	a,c	; see if
Ø213 felb	cpi	27	; past end of t
Ø215 daØaØ2	jc	wr\$sec	; no, do another
	_		

; arrive here at end of track, move to next track

Ø218 Ø4 Ø219 78 Ø21a feØ2 Ø21c daØ8Ø2	inr b mov a,b cpi 2 jc wr\$trk	<pre>; track = track+1 ; see if ; last track ; no, do another</pre>
;	done with putsys, halt for lack	of anything bette
Ø2lf fb	ei	

Ø22Ø 76 hlt

;

;

;

;

; user supplied subroutines for sector read and write

move to next page boundary

0300 org (\$+0100h) and 0ff00h

read\$sec:
 ; read the next sector
 ; track in ,
 ; sector in <c>
 ; dmaaddr in <h1>

Ø300 c5 push b Ø301 e5 push h ; user defined read operation goes here Ø302 ds 64

0342 el pop h 0343 cl pop b

Ø344 c9	ret
0400	org (\$+0100h) and 0ff00h ; another page bo
	write\$sec:
	; same parameters as read\$sec
Ø400 c5	push b
Ø4Ø1 e5	push h
0402	; user defined write operation goes here ds 64
Ø442 el	pop h
Ø443 cl	pop b
Ø444 c9	ret
	; end of getsys/putsys program
Ø445	end

; this is a sample cold start loader which, when modified ; resides on track ØØ, sector Øl (the first sector on the ; diskette). we assume that the controller has loaded ; this sector into memory upon system start-up (this pro-; gram can be keyed-in, or can exist in read/only memory ; beyond the address space of the cp/m version you are ; running). the cold start loader brings the cp/m system ; into memory at "loadp" (3400h + "bias"). in a 20k ; memory system, the value of "bias" is 0000h, with large ; values for increased memory sizes (see section 2). afte ; loading the cp/m system, the clod start loader branches ; to the "boot" entry point of the bios, which begins at ; "bios" + "bias." the cold start loader is not used un-; til the system is powered up again, as long as the bios ; is not overwritten. the origin is assumed at 0000h. an ; must be changed if the controller brings the cold start ; loader into another area, or if a read/only memory area ; is used.

0000		org	Ø	;	base of ram in cp/m
ØØ14 =	msize	equ	20	;	min mem size in kbytes
	bias ccp bios biosl boot size sects	egu egu egu egu egu egu	(msize-20)*1024 3400h+bias ccp+1600h 0300h bios bios+biosl-ccp size/128	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;	offset from 20k system base of the ccp base of the bios length of the bios size of cp/m system # of sectors to load
	;	begin tl	ne load operation	n	
0000 010000	cold:	1	h 0	_	b-a

0000 01020	00 lxi	b,2	;	$b=\emptyset$, c=sector 2
ØØØ3 1632	mvi	d,sects	;	d=# sectors to load
0005 21003	34 lxi	h,ccp	;	base transfer address

lsect: ; load the next sector

; insert inline code at this point to ; read one 128 byte sector from the ; track given in register b, sector ; given in register c, ; into the address given by <hl> ; ; branch to location "cold" if a read error occurs

; * ; * user supplied read operation goes here ... ; * ; ****** ; ØØØ8 c36bØØ past\$patch ; remove this when patche jmp 60h ØØØb ds past\$patch: ; go to next sector if load is incomplete ØØ6b 15 dcr d ; sects=sects-1 006c ca004a ΊZ boot ; head for the bios more sectors to load ; ; ; we aren't using a stack, so use <sp> as scratch registe to hold the load address increment ; ; 128 bytes per sector ØØ6f 318ØØØ lxi sp,128 0072 39 dad sp ; $\langle hl \rangle = \langle hl \rangle + 128$ ØØ73 Øc inr С ; sector = sector + 1 0074 79 mov a,c 0075 felb 27 cpi ; last sector of track? 0077 da0800 jc lsect ; no, go read another ; end of track, increment to next track ØØ7a ØeØl c,1 mvi ; sector = 1Ø07c Ø4 b inr ; track = track + 1007d c30800 ; for another group jmp lsect ØØ8Ø ; of boot loader end

APPENDIX F: CP/M DISK DEFINITION LIBRARY 1: ; CP/M 2.0 disk re-definition library 2: ; 3: ; Copyright (c) 1979 4:; Digital Research 5:; Box 579 6: ; Pacific Grove, CA 7: ; 9395Ø 8:; CP/M logical disk drives are defined using the 9: ; 10: ; macros given below, where the sequence of calls 11: ; is: 12: ; 13: ; disks n diskdef parameter-list-Ø 14: ; 15: ; diskdef parameter-list-l 16: ; diskdef parameter-list-n 17: : 18: ; endef 19: ; 20:; where n is the number of logical disk drives attached to the CP/M system, and parameter-list-i defines the 21: ; 22: : characteristics of the ith drive $(i=\emptyset, 1, \ldots, n-1)$ 23: ; 24: ; each parameter-list-i takes the form dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0] 25: ; 26: ; where is the disk number Ø,1,...,n-1 27: ; dn 28: ; fsc is the first sector number (usually \emptyset or 1) is the last sector number on a track 29: : lsc is optional "skew factor" for sector translate 30: ; skf 31: ; is the data block size (1024,2048,...,16384) bls 32: ; dks is the disk size in bls increments (word) 33: ; dir is the number of directory elements (word) 34: ; is the number of dir elements to checksum cks is the number of tracks to skip (word) 35:; ofs 36: ; is an optional Ø which forces 16K/directory en [Ø] 37: ; 38: ; for convenience, the form 39: ; dn,dm 40:; defines disk dn as having the same characteristics as 41: ; a previously defined disk dm. 42: : 43: ; a standard four drive CP/M system is defined by 44: ; disks 45: ; diskdef 0,1,26,6,1024,243,64,64,2 46: : dsk set Ø 47: ; rept 3 48: ; dsk+l dsk set 49: ; diskdef %dsk,Ø 50:; endm 51: ; endei 52: ; the value of "begdat" at the end of assembly defines t 53: ;

54: ; beginning of the uninitialize ram area above the bios, 55:; while the value of "enddat" defines the next location 56: ; following the end of the data area. the size of this 57: : area is given by the value of "datsiz" at the end of t 58: ; assembly. note that the allocation vector will be qui 59: ; large if a large disk size is defined with a small blo 60: : size. 61: ; 62: dskhdr macro dn 63: ;; define a single disk header list 64: dpe&dn: dw xlt&dn.0000h ;translate table dw 65: 0000h.0000h ;scratch area 66: dw dirbuf,dpb&dn ;dir buff,parm block 67: dw csv&dn,alv&dn ;check, alloc vectors 68: endm 69: ; 70: disks macro nd 71: ;; define nd disks 72: ndisks ;; for later reference set nd 73: dpbase Ŝ equ ; base of disk parameter blocks 74: ;; generate the nd elements 75: dsknxt Ø set 76: rept nd 77: dskhār %dsknxt 78: dsknxt dsknxc+1 set 79: endm 80: endm 81: ; 82: dpbhdr macro dn 83: dpb&dn equ \$;disk parm block 84: endm 85: ; 86: ddb macro data, comment 87: ;; define a db statement 88: db data comment 89: endm 90: ; 91: ddw macro data, comment 92: ;; define a dw statement 93: dw data comment 94: endm 95: ; 96: gcd macro m,n 97: ;; greatest common divisor of m,n 98: ;; produces value gcdn as result 99: ;; (used in sector translate table generation) 100: gcdm set m ;;variable for m 101: gcdn set n ;;variable for n 102: gcdr set Ø ;;variable for r 103: 65535 rept 104: gcdxset gcdm/gcdn 105: gcdr set gcdm - gcdx*gcdn 106: if $gcdr = \emptyset$ 107: exitm 108: endif

109: gcdm set qcdn llØ: gcdn set gcdr 111: endm 112: endm 113: ; 114: diskdef macro dn,fsc,lsc,skf,bls,dks,dir,cks,bfs,kl6 115: ;; generate the set statements for later tables nul lsc 116: if 117: ;; current disk dn same as previous fsc 118: dpb&dn equ dpb&fsc ;equivalent parameters 119: als&dn equ als&fsc ;same allocation vector size 120: css&dn equ css&fsc ;same checksum vector size 121: xlt&dn equ xlt&fsc ;same translate table 122: else ;;sectors Ø...secmax 123: secmax set lsc-(fsc) secmax+1;;number of sectors 124: sectors set (dks)/8 ;;size of allocation vector 125: als&dn set ((dks) mod ε) ne Ø 126: if 127: als&dn set als&dn+1 128: endif 129: css&dn set (cks)/4 ;;number of checksum elements 130: ;; gene 131: blkval set generate the block shift value bls/128 ;;number of sectors/block 132: blkshf set ;;counts right 0's in blkval Ø ;; rills with 1's from right 133: blkmsk set Ø 134: ;;once for each bit position rept 16 135: if blkval=1 136: exitm 137: endif 138: ;; otherwise, high order 1 not found yet 139: blkshf set blkshf+l 140: blkmsk set (blkmsk shl 1) or 1 141: blkval set blkval/2 142: endm 143: ;; generate the extent mask byte 144: blkval set bls/1024 ;;number of kilobytes/block 145: extmsk set Ø ;;fill from right with 1's 146: rept 16 147: if blkval=1 148: exitm 149: endif 150: ;; otherwise more to shift 151: extmsk set (extmsk shl 1) or 1 152: blkval set blkval/2 153: endm 154: ;; may be double byte allocation (dks) > 256155: if 156: extmsk set (extmsk shr 1) 157: endif 158: ;; may be optional $[\emptyset]$ in last position 159: if not nul k16 160: extmsk set k16 161: endif 162: ;; now generate directory reservation bit vector 163: dirrem set dir ;;# remaining to process

bls/32 ;;number of entries per block 164: dirbks set ;;fill with l's on each loop Ø 165: dirblk set 166: 16 rept 167: if dirrem=Ø 168: exitm endif 169: not complete, iterate once again 170: ;; shift right and add 1 high order bit 171: ;; 172: dirblk (dirblk shr i) or 8000h set dirrem > dirbks 173: if 174: dirrem set dirrem-dirbks 175: else 176: dirrem Ø set endif 177: 178: endm 179: dpbhdr dn ;;generate equ \$ 180: ddw %sectors,<;sec per track> %blkshf,<;blcck shift> ddb 181: ddb %blkmsk,<;block mask> 182: ddb 183: %extmsk,<;extnt mask> 184: ddw %(dks)-l,<;uisk size-l> 185: adw %(dir)-l,<;oirectory max> 186: dãb %dirblk shr 8,<;alloc0> %dirblk and Øffh,<;allocl> ddb 187: %(cks)/4,<;check size> 188: ddw %ofs,<;offset> 189: ddw generate the translate table, if requested 190: ;; 191: if nul skf 192: xlt&dn ;no xlate table equ Ø 193: else 194: if $skf = \emptyset$ 195: xlt&dn equ Ø ;no xlate table 196: else 197: ;; generate the translate table 198: nxtsec Ø ;;next sector to fill set 199: nxtbas set Ø ;; moves by one on overflow 200: gcd %sectors,skf 201: ;; gcdn = gcd(sectors,skew) 202: neltst sectors/gcdn set 203: ;; neltst is number of elements to generate 204: ;; before we overlap previous elements 205: nelts neltst ;;counter set ;translate table 206: xlt&dn \$ equ 207: sectors ;; once for each sector rept 208: if sectors < 256209: ddb %nxtsec+(fsc) 210: else 211: ddw %nxtsec+(fsc) 212: endif 213: nxtsec set nxtsec+(skf)214: if nxtsec >= sectors 215: nxtsec set nxtsec-sectors 216: endif 217: nelts set nelts-1 218: if nelts = \emptyset

220:	nxtbas nxtsec nelts	set set endif endm	nxtbas+i nxtbas neltst
224: 225: 226: 227:	;	endif endif endm	;;end of nul fac test ;;end of nul bls test
	defds lab:	macro ds endm	lab,space space
232: 233: 234:	lds	macro äefds endm	lb,dn,val lb&dn,%val&dn
235: 236:	; endef	macro	
237: 238:	;; begdat		e the necessary ram data areas \$
239:	dirbuf: dsknxt		<pre>128 ;directory access buffer Ø</pre>
241: 242: 243:	GDAMAC	rept lås lås	ndisks ;;once for each disk alv,%dsknxt,als csv,%dsknxt,css
244: 245:	dsknxt	set endm	dsknxt+1
	enddat datsiz ;;	egu egu db Ø at endm	\$ \$-begdat this point forces hex record

1: :***** 2: :* Sector Deblocking Algorithms for CP/M 2.0 * 3: ;* 4: ;* 6: : 7: ; utility macro to compute sector mask 8: smask hblk macro 9: ;; compute log2(hblk), return @x as result 10: ;; (2 ** @x = hblk on return)11: @y hblk set 12: @x set Ø 13: ;; count right shifts of @y until = 1 14: rept 8 15: if @y = 116: exitm 17: endif 18: ;; @y is not 1, shift right one position 19: @y set @y shr l 20: @x 0x + 1set 21: endm 22: enâm 23: ; 25: ;* * 26: :* CP/M to host disk constants * 27: ;* * 29: blksiz equ 2048 ;CP/M allocation size 30: hstsiz equ 512 ;host disk sector size 31: hstspt equ 20 ;host disk sectors/trk 32: hstblk equ hstsiz/128 ;CP/M sects/host buff hstblk * hstspt ;CP/M sectors/track 33: cpmspt equ 34: secmsk equ hstblk-1 ;sector mask 35: smask hstblk :compute sector mask 36: secshf equ **@x** ;log2(hstblk) 37: ; 39: ;* * 40: ;* * BDOS constants on entry to write 41: ;* * 43: wrall Ø equ ;write to allocated 44: wrdir 1 equ ;write to directory 45: wrual egu 2 ;write to unallocated 46: ; 48: ;* * 49: ;* The BDOS entry points given below show the * 50: ;* code which is relevant to deblocking only. * 51: ;* 53: ;

54:; DISKDEF macro, or hand coded tables go here ;disk param block base 55: dpbase equ \$ 56: ; 57: boot: 58: wboot: ;enter here on system boot to initialize 59: 6Ø: xra а ;Ø to accumulator hstact 61: sta ;host buffer inactive 62: ;clear unalloc count sta unacnt 63: ret 64: ; 65: seldsk: 66: ;select disk 67: ;selected disk number mov a,c 68: sta sekdsk ;seek disk number 69: ;disk number to HL mov 1,a 7Ø: h,Ø mvi 71: rept 4 ; multiply by 16 72: h dad 73: endm 74: ;base of parm block lxi d,dpbase 75: dad d ;hl=.dpb(curdsk) 76: ret 77:; 78: settrk: 79: ;set track given by registers BC 80: mov h,b 81: mov 1,c 82: shld sektrk ;track to seek 83: ret 84: ; 85: setsec: 86: ;set sector given by register c 87: mov a,c 88: sta seksec ;sector to seek 89: ret 90: ; 91: setdma: 92: ;set dma address given by BC 93: h,b mov 94: 1,c mov 95: shld dmaadr 96: ret 97:; 98: sectran: 99: ;translate sector number BC 100: mov h,b 101: mov 1,c 102: ret 103: ;

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* 105: :* * 106: :* The READ entry point takes the place of * 107: ;* the previous BIOS definition for READ. 108: ;* * 110: read: ;read the selected CP/M sector 111: 112: mvi a.l 113: sta readop ;read operation 114: sta rsflag ;must read data 115: mvi a,wrual 116: sta wrtype ;treat as unalloc 117: jmp rwoper ; to perform the read 118: ; 120: ;* * 121: ;* The WRITE entry point takes the place of 122: ;* the previous BIOS definition for WRITE. * + 123: ;* 125: write: 126: ;write the selected CP/M sector 127: xra ;Ø to accumulator а 128: sta readop ;not a read operation 129: mov ;write type in c a,c 130: sta wrtype 131: cpi wrual ;write unallocated? 132: chkuna ;check for unalloc jnz 133: ; 134: ; write to unallocated, set parameters 135: mvi a,blksiz/128 ;next unalloc recs 136: sta unacnt 137: lda sekdsk disk to seek 138: sta unadsk ;unadsk = sekdsk 139: lhld sektrk 140: shld unatrk ;unatrk = sectrk 141: lda seksec 142: sta unasec ;unasec = seksec 143: ; 144: chkuna: ; check for write to unallocated sector 145: 146: lda unacnt ;any unalloc remain? 147: ora а 148: jz alloc ;skip if not 149: ; 150:; more unallocated records remain 151: dcr а ;unacnt = unacnt-1 152: sta unacht 153: lda sekdsk ;same disk? 154: lxi h,unadsk 155: cmp m ;sekdsk = unadsk? 156: jnz alloc ;skip if not 157: : 158: ; disks are the same

159: lxi h,unatrk 160: call sektrkcmp ;sektrk = unatrk? 161: alloc ;skip if not jnz 162: ; 163: ; tracks are the same 164: lda seksec ;same sector? 165: lxi h,unasec 166: ;seksec = unasec? cmp m 167: alloc ;skip if not jnz 168: ; match, move to next sector for future ref 169: ; 170: ;unasec = unasec+1 inr m 171: ;end of track? mov a,m 172: cpi cpmspt ;count CP/M sectors 173: jc noovf ;skip if no overflow 174: ; 175: ; overflow to next track 176: mvi m,Ø ;unasec = Ø 177: lhld unatrk 178: inx h 179: shld ;unatrk = unatrk+1 unatrk 18Ø: ; 181: noovf: ;match found, mark as unnecessary read 182: 183: ;Ø to accumulator xra а 184: rsflag ;rsflag = \emptyset sta 185: ; to perform the write jmp rwoper 186: ; 187: alloc: ;not an unallocated record, requires pre-read 188: ;0 to accum 189: xra а 190: sta ; unacht = \emptyset unacnt 191: inr ;1 to accum а 192: sta rsflag ;rsflag = 1 193: ; * 195: ;* * 196: ;* Common code for READ and WRITE follows 197: ;* * 199: rwoper: 200: ;enter here to perform the read/write 201: xra ;zero to accum а 202: sta erflag ;no errors (yet) 203: lda seksec ; compute host sector 204: rept secshf 205: ora ; carry = \emptyset а 206: ;shift right rar 207: endm 208: sekhst sta ;host sector to seek 209: ; 210:; active host sector? 211: h,hstact ;host active flag lxi 212: mov a,m 213: mvi ;always becomes 1 m,1

а 214: ;was it already? ora filhst 215: jz ;fill host if not 216: ; 217: ; host buffer active, same as seek buffer? 218: lda sekdsk lxi cmp 219: h,hstdsk :same disk? 220: ;sekdsk = hstdsk? m jnz 221: nomatch 222: ; same disk, same track? 223: ; 224: lxi h,hsttrk 225: call sektrkcmp ;sektrk = hsttrk? 226: jnz nomatch 227:; 228: ; same disk, same track, same buffer? 229: lda sekhst 230: lxi h,hstsec ;sekhst = hstsec? 231: cmp m 232: jz match ;skip if match 233: ; 234: nomatch: 235: ;proper disk, but not correct sector 236: lda hstwrt ;host written? 237: ora а 238: cnz writehst ;clear host buff 239: ; 240: filhst: 241: ;may have to fill the host buffer lda sekdsk sta hstdsk 242: 243: sta hstdsk lhld sektrk shld hsttrk lda sekhst sta hstsec lda rsflag 244: 245: 246: 247: 248: ;need to read? ora a cnz readhst xra a sta hstwrt 249: 250: ;yes, if l 251: ;0 to accum 252: ;no pending write 253: ; 254: match: 255: ; copy data to or from buffer 256: lda seksec ;mask buffer number 257: ;least signif bits ani secmsk 258: ;ready to shift mov l,a h,Ø 259: mvi ;double count 260: rept 7 ;shift left 7 261: dad h 262: endm 263: ; hl has relative host buffer address lxi 264: d,hstbuf 265: dad d ;hl = host address 266: xchq ;now in DE 267: lhld dmaadr ;get/put CP/M data 268: mvi c,128 ;length of move

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269: lda readop ;which way? 270: ora а ;skip if read 271: jnz rwmove 272: ; write operation, mark and switch direction 273: ; 274: mvi a,1 275: hstwrt ; hstwrt = 1sta ;source/dest swap 276: xchq 277: ; 278: rwmove: ;C initially 128, DE is source, HL is dest 279: 28Ø: ldax d ;source character 281: inx d 282: mov m,a ;to dest 283: inx h 284: ;loop 128 times dcr С rwmove 285: jnz 286: ; 287: ; data has been moved to/from host buffer 288: lda wrtype ;write type cpi wrdir ;to directory? 289: ; in case of errors 290: lda erflag 291: ;no further processing rnz 292: ; 293: ; clear host buffer for directory write 294: ;errors? ora а 295: ;skip if so rnz ;Ø to accum 296: xra а hstwrt ; buffer written 297: sta 298: call writehst 299: lda erflag 300: ret 301: ; * 303: ;* * Utility subroutine for 16-bit compare 304: ;* * 305: ;* 307: sektrkcmp: 308: ;HL = .unatrk or .hsttrk, compare with sektrk 309: xchq 310: lxi h,sektrk 311: ldax d ;low byte compare 312: cmp m ;same? 313: rnz ;return if not low bytes equal, test high ls 314: ; 315: inx d 316: h inx 317: ldax d 318: cmp m ;sets flags 319: ret 320: ;

322: ;* * × 323: :* WRITEHST performs the physical write to * 324: ;* the host disk, READHST reads the physical * 325: ;* disk. * 326: ;* 328: writehst: ;hstdsk = host disk #, hsttrk = host track #, 329: 330: ;hstsec = host sect #. write "hstsiz" bytes 331: ; from hstbuf and return error flag in erflag. 332: ;return erflag non-zero if error 333: ret 334: ; 335: readhst: ;hstdsk = host disk #, hsttrk = host track #, 336: 337: ;hstsec = host sect #. read "hstsiz" bytes 338: ; into hstbuf and return error flag in erflag. 339: ret 340: ; 342: ;* * * 343: ;* Unitialized RAM data areas 344: :* * 346: ; 347: sekdsk: ds 1 ;seek disk number 348: sektrk: ds 2 ;seek track number 349: seksec: ds 1 ;seek sector number 350:; 351: hstdsk: ds 1 ;host disk number 352: hsttrk: ds 2 ;host track number 353: hstsec: ds 1 ;host sector number 354: ; ;seek shr secshf 355: sekhst: ds 1 356: hstact: ds 1 ;host active flag 357: hstwrt: ds 1 ;host written flag 358: ; 359: unacnt: ds 1 ;unalloc rec cnt 360: unadsk: ds 1 ;last unalloc disk 361: unatrk: ds 2 ;last unalloc track 362: unasec: ds 1 ;last unalloc sector 363: ; 364: erflag: ds 1 ;error reporting 1 365: rsflag: ds ;read sector flag 366: readop: ds 1 ;1 if read operation 1 367: wrtype: ds ;write operation type 368: dmaadr: ds 2 ;last dma address 369: hstbuf: ds hstsiz ;host buffer 370:;

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